

**Definition 1 (Sublinear functional)**  $\rho : X \rightarrow \mathbb{R}$  is a sublinear functional if

$$\begin{aligned}\rho(x + y) &\leq \rho(x) + \rho(y) & x, y \in X \\ \rho(\lambda x) &= \lambda \rho(x) & \lambda > 0\end{aligned}$$

I'm splitting the proof into two parts because the infinite-dimensional part of the extension is independent from ZF.

**Lemma 2 (one-dimensional extension)** *Suppose  $E$  is a subspace of  $X$ . Let  $E' = E \oplus \text{span}\{e\}$  be a one-dimension-larger subspace of  $X$ . If  $\rho : X \rightarrow \mathbb{R}$  is a sublinear functional and  $f : E \rightarrow \mathbb{R}$  is a linear functional bounded by  $\rho$ , then there is an extension  $f'$  of  $f$  to  $E'$  such that  $f'$  is also bounded by  $\rho$ .*

**Proof.** Proceed by finding a bound on potential extensions to  $E'$  where  $f' \leq \rho$  still. Suppose we already have a suitable  $f'$ . Let  $f'(e) = \alpha$ . To extract a bound on  $\alpha$ , consider the following inequalities:

$$f'(x - e) = f(x) - \alpha \leq \rho(x - e) \quad (1)$$

$$f'(y + e) = f(y) + \alpha \leq \rho(y + e) \quad (2)$$

Combining these gives

$$f(x) - \rho(x - e) \leq \alpha \leq \rho(y + e) - f(y)$$

Because  $x$  and  $y$  are arbitrary, this requires

$$\sup \{f(x) - \rho(x - e)\} \leq \alpha \leq \inf \{\rho(y + e) - f(y)\}. \quad (3)$$

It is necessary to show the left hand side of [eq. \(3\)](#) is, in fact, less than or equal to the right hand side. But

$$f(x) + f(y) = f(x + y) \leq \rho(x + y) \leq \rho(x - e) + \rho(y + e).$$

Rearranging gives  $f(x) - \rho(x - e) \leq \rho(y + e) - f(y)$ . Recalling, again, that  $x$  and  $y$  are arbitrary gives

$$\sup \{f(x) - \rho(x - e)\} \leq \inf \{\rho(y + e) - f(y)\}.$$

Therefore there is at least one such  $\alpha$ .

Now, we need to check that satisfying [eq. \(3\)](#) is sufficient to give a bounded extension of  $f$ . Suppose  $\alpha$  satisfies [eq. \(3\)](#) and, if  $x \in E$ , define the extension  $f'(x + \lambda e) = f(x) + t\alpha$ . If  $\lambda > 0$ ,

$$f'(x + \lambda e) = \lambda(f(x/\lambda) + \alpha) \leq \underbrace{\lambda\rho(x/\lambda + e)}_{\text{by eq. (2)}} = \rho(x + \lambda e)$$

If  $\lambda = -\mu < 0$ ,

$$f'(x - \mu e) = \mu(f(x/\mu) - \alpha) \leq \underbrace{\mu\rho(x/\mu - e)}_{\text{by eq. (1)}} = \rho(x - \mu e)$$

proving sufficiency. □

By induction, this proves we can extend to a finite number of dimensions. The typical approach uses Zorn's lemma to extend to the infinite dimensional case:

**Theorem 3 (Hahn-Banach theorem)** *If  $f : E \rightarrow \mathbb{R}$  is bounded by a sublinear functional  $\rho : X \rightarrow \mathbb{R}$ , then  $f$  extends to a functional  $f' : X \rightarrow \mathbb{R}$  which is bounded by  $\rho$  on the whole space.*

**Proof.** The set of linear extensions of  $f$  from subspaces of  $X$  to  $\mathbb{R}$  is a poset ordered by inclusion. But if  $f_\alpha$  are chain of extensions of  $f$ , they must agree on every subset of  $X$  on which they are defined. Hence  $\sup f_\alpha := \bigcup f_\alpha$  is well-defined and gives an upper bound on the chain. By Zorn's lemma, there is a maximal  $f^*$ .

Suppose  $f^*$  is not defined on  $E$ , not the entirety of  $X$ . Then pick some  $e$  on which  $f^*$  is not defined. By [lemma 2 \(one-dimensional extension\)](#), there is some  $f^{*'}$  extending  $f^*$  to  $E \oplus \text{span}\{e\}$ . This contradicts maximality; hence,  $f^* : X \rightarrow \mathbb{R}$ .  $\square$

Here is a sketch of an alternate proof using the compactness of first order logic, which is strictly weaker than the axiom of choice.

**Proof.** We need to axiomatize the statement of the theorem. Consider the domain of discourse of points of the form  $X \times \mathbb{R}$ . Let each needed  $X \times \mathbb{R}$  be a constant in this language. Write the axioms of a vector space and a linear map in terms of the domain of discourse. Take the predicate  $\square \in E$  for every subspace  $E$ . Then consider the axiom schema  $\forall(x,y) x \in E \implies y \leq \rho(x)$  encoding “bounded by  $\rho$ ”. Inductively applying [lemma 2 \(one-dimensional extension\)](#) proves that each finite subtheory has a model. By compactness, there must be a model of the whole theory, which is a function defined on all of  $X$ , and this model is the desired extension.  $\square$

Although Hahn-Banach is weaker than logical compactness, it is still strong enough to construct the Banach-Tarski paradox.