

# 1 Ordinals and transfinite induction

**Definition 1** (The ordinals,  $\text{Ord}$ ) *These are defined inductively:*

- $0 := \emptyset \in \text{Ord}$
- If  $n \in \text{Ord}$  then

$$\text{succ}(n) := n \cup \{n\} \in \text{Ord}$$

- If  $\mathcal{X} \subseteq \text{Ord}$  is a set, then

$$\sup \mathcal{X} := \bigcup \mathcal{X} \in \text{Ord}$$

*In other words, an ordinal is the set of all smaller ordinals. In particular, the last definition precludes the ordinals from being a set.  $\text{Ord}$  is bigger than any set. They form a class.*

*If  $\alpha = \text{succ}(n)$  then  $\alpha$  is a successor ordinal. Otherwise,  $\alpha$  is a limit ordinal.*

**Theorem 2** ( $\text{Ord}$  is well-ordered) *Every nonempty class of ordinals contains a least element.*

**Proof.** Let  $\mathcal{X}$  be a nonempty class of ordinals. Then it suffices to show

$$I := \inf \mathcal{X} = \bigcap \mathcal{X} \in \mathcal{X}$$

Suppose  $I \notin \mathcal{X}$ . Then for any  $x \in \mathcal{X}$ ,  $I < x$ . Consequently, for any  $x \in \mathcal{X}$ ,  $\text{succ}(I) \leq x$ . But then  $\text{succ}(I) = I$ , a contradiction.  $\square$

**Theorem 3** (Transfinite induction) *For a logical statement  $\phi$ , if*

1.  $\phi(0)$  holds,
2.  $\phi(n) \implies \phi(\text{succ}(n))$  (strictly speaking, this follows from (3) but is usually easier to prove first)
3. If, for all  $n < \alpha$ ,  $\phi(n)$  is true then  $\phi(\alpha)$  is true:

$$(\forall n < \alpha) \phi(n) \implies \phi(\alpha)$$

*Then  $\phi$  holds for all ordinals.*

**Proof.** Let  $\mathcal{X}$  be the class where  $\phi$  does not hold. Suppose  $\mathcal{X} \neq \emptyset$ . Then let for all ordinals  $n < \inf \mathcal{X}$ ,  $\phi(n)$  by hypothesis. Hence, by (3),  $\phi(\inf \mathcal{X})$  holds. But then  $\inf \mathcal{X} \notin \mathcal{X}$ , contradicting **theorem 2** ( $\text{Ord}$  is well-ordered). Therefore,  $\mathcal{X}$  is empty.  $\square$

**Theorem 4 (Ord fixed point theorem)** *Let  $(X, \leq)$  be a complete ordered space. Suppose  $F : X \rightarrow X$  is nondecreasing. Define*

$$F^0(x) := x$$

$$F^{\text{succ}(n)}(x) := f(F^n(x))$$

*If  $\alpha$  is a limit ordinal,*

$$F^\alpha(x) := \sup_{\beta < \alpha} F^\beta(x)$$

*Then for any  $x \in X$ ,  $F^\alpha(x)$  is eventually the smallest fixed point of  $F$  greater than or equal to  $x$ .*

**Proof.** Fix  $x$ . Then  $F^\square(x) : \mathbf{Ord} \rightarrow X$  is nondecreasing by construction. Because  $\mathbf{Ord}$  is bigger than any set,  $F^\square(x)$  cannot be injective. Hence there is some smallest  $\alpha \in \mathbf{Ord}$  such that  $F^\alpha(x) = F^\gamma(x)$  with  $\gamma > \alpha$ . But  $\alpha < \text{succ}(\alpha) \leq \beta$ , so by monotonicity,

$$F^\alpha(x) \leq F^{\text{succ}(\alpha)}(x) \leq F^\beta(x) = F^\alpha(x)$$

hence

$$F(F^\alpha(x)) = F^{\text{succ}(\alpha)}(x) = F^\alpha(x)$$

so  $F$  fixes  $F^\alpha(x)$ .

Suppose  $y$  is a fixed point and  $x \leq y \leq F^\alpha(x)$ . Then there is some smallest  $\beta$  such that  $y \leq F^\beta(x)$ :

- If  $\beta = \text{succ}(\gamma)$ , then

$$F^\gamma(x) \leq y \leq f(F^\gamma(x)) = F^\beta(x)$$

But by monotonicity,  $F(F^\gamma(x)) \leq f(y) = y$  by assumption  $y$  is fixed. Thus  $y = F^\gamma(x)$  for some  $\gamma \leq \alpha$ .

- If  $\beta$  is a limit ordinal, then for all  $\gamma < \beta$ ,

$$F^\gamma(x) \leq y \leq F^\beta(x) = \sup_{\hat{\gamma} < \beta} F^{\hat{\gamma}}(x) = y$$

so  $F^\beta(x) = y$ . □

By characterizing  $\sigma$ -algebras and  $\lambda$ -systems as fixed points of set functions, this result becomes a powerful tool for proving results about generated  $\sigma$ -algebras and  $\lambda$ -systems.

## 2 $\pi$ - $\lambda$ systems and theorem

Special notation:

- I'll use  $\amalg$  to represent *disjoint* unions,
- $\bigcup\uparrow$  to represent *increasing* unions,

- and  $\setminus_{\supseteq}$  to represent *proper* relative complements:  $A \setminus B$  when  $B \subseteq A$ .

Definition 5 ( $\lambda$ -system aka Dynkin system)  $\mathcal{D}$  is a Dynkin system on  $\Omega$  iff

- $\Omega \in \mathcal{D}$
- Closure under proper relative complements:

$$A \subseteq B \implies B \setminus_{\supseteq} A \in \mathcal{D}$$

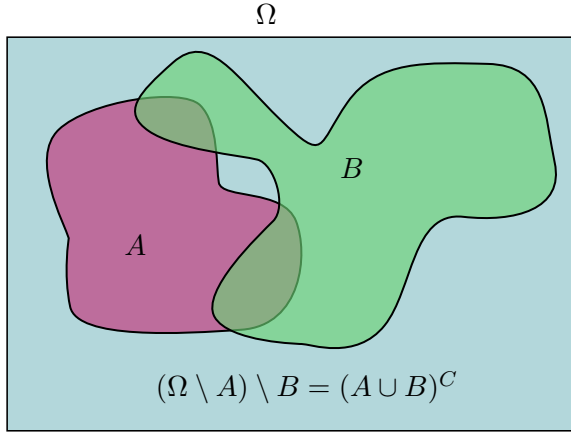
- Closure under increasing  $\sigma$ -unions:

$$A_1 \subseteq A_2 \subseteq \dots \in \mathcal{D} \implies \bigcup_{\uparrow} A_n \in \mathcal{D}$$

Allowing arbitrary complements would give a  $\sigma$ -algebra as

$$(\Omega \setminus A) \setminus B = (A \cup B)^C = A^C \cap B^C$$

allowing arbitrary  $\sigma$ -unions and intersections.



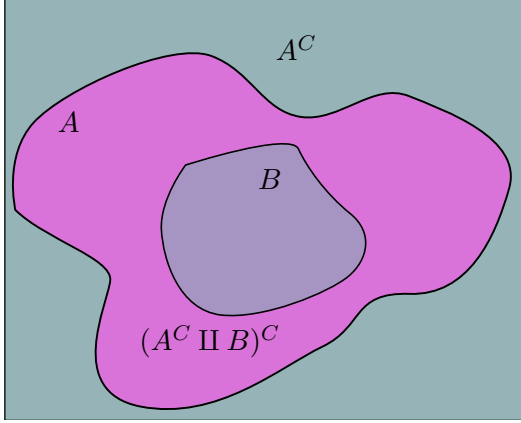
Theorem 6 (Alternative definition of  $\lambda$ -system) *Equivalently,*

- $\Omega \in \mathcal{D}$
- Closure under global complements:

$$A \in \mathcal{D} \implies A^C \in \mathcal{D}$$

- Closure under disjoint  $\sigma$ -unions

Proof. Suppose  $\mathcal{D}$  meets the criteria of [theorem 6 \(Alternative definition of  \$\lambda\$ -system\)](#). Then  $\mathcal{D}$  contains proper relative complements as  $A \setminus_{\supseteq} B = (A^C \cup B)^C$ :



But by taking proper relative complements, we can convert an ascending  $\sigma$ -union into a disjoint  $\sigma$ -union: if  $A_n \uparrow A$ , let

$$B_1 = A_1 \in \mathcal{D}$$

$$B_n = A_n \setminus \bigcup_{i < n} A_i \in \mathcal{D}$$

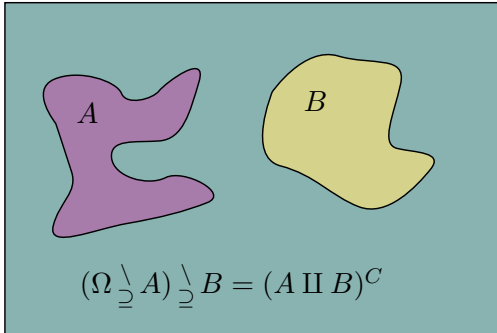
Then  $\{B_n\}$  are pairwise-disjoint and

$$A = \bigcup_n A_n = \coprod_n B_n \in \mathcal{D}$$

proving sufficiency.

For necessity, suppose  $\mathcal{D}$  satisfies the criteria of [definition 5](#) ( $\lambda$ -system aka Dynkin system). As  $A^C = \Omega \setminus A$ , a proper complement, it suffices to show closure under disjoint unions. First, note that it is closed under finite disjoint unions:

$$A \amalg B = ((\Omega \setminus A) \setminus B)^C \in \mathcal{D}$$



because the complements are proper by disjointness. Suppose  $B_i$  are pairwise disjoint. Then  $B = \coprod B_i \in \mathcal{D}$  by hypothesis.

$$A_1 = B_1 \in \mathcal{D}$$

$$A_n = B_n \amalg A_{n-1} \in \mathcal{D}$$

therefore

$$B = \coprod B_n = \bigcup \uparrow A_n \in \mathcal{D} \quad \square$$

**Definition 7 ( $\pi$ -system)** A  $\pi$ -system is closed under finite intersections.

**Theorem 8** If  $\mathcal{D}$  is a  $\pi$ -system and  $\lambda$ -system then it is a  $\sigma$ -algebra.

**Proof.** First, prove  $\mathcal{D}$  is an algebra. Note that  $\mathcal{D}$  is closed under binary union as  $A \cup B = (A^C \cap B^C)^C$  and  $A^C, B^C \in \mathcal{D}$  because  $\mathcal{D}$  is a  $\lambda$ -system, so  $A^C \cap B^C \in \mathcal{D}$  because it is a  $\pi$ -system, and its complement is in  $\mathcal{D}$  by closure under complements. In particular,  $\mathcal{D}$  must be closed under *arbitrary* relative complements as

$$A \setminus B = A \setminus_{\supseteq} (A \cap B)$$

Suppose  $A_i$  is a countable sequence of sets in  $\mathcal{D}$ . Then let

$$\begin{aligned} B_1 &= A_1 \\ B_n &= A_n \setminus \bigcup_{i < n} A_i \end{aligned}$$

so  $B_i$  are pairwise-disjoint and

$$\bigcup A_n = \coprod B_n \in \mathcal{D}$$

proving closure under  $\sigma$ -unions.  $\square$

**Theorem 9 (Dynkin's  $\pi$ - $\lambda$  theorem)** If  $P \subseteq \mathcal{D}$  where  $\pi$  is a  $\pi$ -system and  $\mathcal{D}$  is a  $\lambda$ -system, then

$$\sigma(P) \subseteq \mathcal{D}$$

**Proof.** As  $\lambda(P)$  is the minimal  $\lambda$ -system containing  $P$ ,

$$\lambda(P) \subseteq \mathcal{D}$$

It suffices to show  $\lambda(P)$  is a  $\pi$ -system because then it must be a  $\sigma$ -algebra.

Let  $U(X)$  indicate the set of disjoint  $\sigma$ -unions in  $X$ :

$$U(X) := \left\{ \prod_{i \in \mathbb{N}} A_i : A_i \in X, A_i \cap A_j = \emptyset \right\}$$

Let  $C(X)$  indicate the set of complements in  $X$ :

$$C(X) := \{A^C : A \in X\}$$

And define

$$\Lambda(X) := X \cup U(X) \cup C(X) = U(X) \cup C(X) : (2^{2^\Omega}, \subseteq) \rightarrow (2^{2^\Omega}, \subseteq)$$

By construction,  $\Lambda(X) = X$  iff  $X$  is a  $\lambda$ -system. By monotonicity ( $X \subseteq \Lambda(X)$ ), the minimum fixed point of  $\Lambda$  containing  $X$  is the  $\lambda$ -system generated by  $X$ . By [theorem 4 \(Ord fixed point theorem\)](#),

$$\Lambda^\omega(P) = \lambda(P)$$

for some  $\omega \in \mathbf{Ord}$ .

We will show by induction that for any  $\alpha \in \mathbf{Ord}$ , the intersection of any two elements in  $\Lambda^\alpha(P)$  is in  $\Lambda^{3\alpha}(P)$ .

1. This is trivially true for  $\alpha = 0$  by the hypothesis that  $P$  is a  $\pi$ -system.
2. Any element in  $\Lambda^{\text{succ}(\alpha)}(P)$  can be written as  $A$ ,  $A^C$ , or  $\coprod A_i$  with  $A, A_i \in \Lambda^\alpha$ . Consider writing an intersection of elements in  $\Lambda^{\text{succ}(\alpha)}$  in terms of proper complements and disjoint unions of intersections of elements in  $\Lambda^\alpha(P)$ :

$$\begin{aligned} A \cap B^C &= A \setminus \overbrace{(B \cap A)}^{\in \Lambda^{3\alpha}(P)} && \in \Lambda^{3\alpha+1}(P) \\ A^C \cap B^C &= (A \cup B)^C = ((A \cap B^C) \amalg B)^C && \in \Lambda^{3\alpha+3}(P) \\ \left( \coprod_i A_i \right) \cap \left( \coprod_j B_j \right) &= \coprod_{i,j} \overbrace{A_i \cap B_j}^{\in \Lambda^{3\alpha}(P)} && \in \Lambda^{3\alpha+1}(P) \\ A^C \cap \coprod B_j &= \coprod \underbrace{A^C \cap B_j}_{\in \Lambda^{3\alpha+1}(P)} && \in \Lambda^{3\alpha+2}(P) \end{aligned}$$

which are all in  $\Lambda^{3\alpha+3}(P) = \Lambda^{3\text{succ}(\alpha)}(P)$ .

3. If  $\alpha$  is a limit ordinal, then set

$$3\alpha \geq 3 \sup_{\beta < \alpha} \beta = \sup_{\beta < \alpha} 3\beta$$

For any two  $A, B \in \Lambda^\alpha$  there must be some  $n_A, n_B < \alpha$  such that  $A \in \Lambda^{n_A}(P)$  and  $B \in \Lambda^{n_B}(P)$ . Hence, if  $n = \max(n_A, n_B)$ , then  $A \cap B \in \Lambda^{3n}(P) \subseteq \Lambda^{3\alpha}(P)$ .

But this process terminates at  $\omega$ , hence  $\Lambda^{3\omega}(P) = \Lambda^\omega(P) = \lambda(P)$  and  $\lambda(P)$  is closed under intersections.  $\square$