

COSC 221 - Introduction to Discrete Structures

Lecture - Logic-04

Readings

- ▶ Propositional Logic: Sections 3.1, 3.2, 3.3
- ▶ Predicate Logic: Sections 3.4
- ▶ Computer Science Connections
 1. Computational Complexity (Section 3.3)
 2. Modern Compilers (Section 3.3)
 3. Game Trees (Section 3.4)

- ▷ Argument Form
- ▷ Inference Rules

Not covered in detail in the textbook,
but important!

- ▷ Argument Form
- ▷ Inference Rules

Not covered in detail in the textbook,
but important!

If $3 > 1$, then $\sqrt{3} > 1$

$3 > 1$

$\therefore \sqrt{3} > 1$

If 2 is even, then 3 is even

2 is even

\therefore 3 is even

Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
- ▷ Inference Rules

Same pattern, different contexts

	$p \Rightarrow q$
	p
\therefore	q

<p>If $3 > 1$, then $\sqrt{3} > 1$</p> <p>$3 > 1$</p> <p>$\therefore \sqrt{3} > 1$</p>
<p>If 2 is even, then 3 is even</p> <p>2 is even</p> <p>\therefore 3 is even</p>

Propositional Logic

- ▷ Argument Form
- ▷ Inference Rules

Arguments and Logic Inference

Same pattern, different contexts

$p \Rightarrow q$
p
$\therefore q$

conclusion

If $3 > 1$, then $\sqrt{3} > 1$

$3 > 1$

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If 2 is even, then 3 is even

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Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
- ▷ Inference Rules

premises (or, hypotheses)

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Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
- ▷ Inference Rules

premises (or, hypotheses)

Same pattern, different contexts

$p \Rightarrow q$
p
$\therefore q$

conclusion

Argument Form: $\varphi_1, \varphi_2, \dots, \varphi_{n-1}, \varphi_n$

A sequence of propositions, containing one/more proposition variables.

Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
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premises (or, hypotheses)

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Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
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premises (or, hypotheses)

Same pattern, different contexts

$p \Rightarrow q$
p
$\therefore q$

conclusion

Valid Argument Form:

Whenever the premises are all true,
the conclusion is true.

Argument Form: $\varphi_1, \varphi_2, \dots, \varphi_{n-1}, \varphi_n$

A sequence of propositions, containing one/more proposition variables.

Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
- ▷ Inference Rules

Same pattern, different contexts

$$\begin{array}{l} p \Rightarrow q \\ p \\ \therefore q \end{array}$$

Q.1) Is this argument form valid?

- A) Yes
- B) No
- C) You haven't taught us yet!
- D) Whatever!

Valid Argument Form:

Whenever the premises are all true, the conclusion is true.

Argument Form: $\varphi_1, \varphi_2, \dots, \varphi_{n-1}, \varphi_n$

A sequence of propositions, containing one/more proposition variables.

Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
- ▷ Inference Rules

p	q	$p \Rightarrow q$	p	q
T	T	T	T	T
T	F	F	T	F
F	T	T	F	T
F	F	T	F	F

Same pattern, different contexts

$$\begin{array}{l}
 p \Rightarrow q \\
 p \\
 \therefore q
 \end{array}$$

Valid Argument Form:

Whenever the premises are all true, the conclusion is true.

Testing Validity of Argument Forms

- ▷ Truth-Table
- ▷ Tautology

Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
- ▷ Inference Rules

premises (or, hypotheses)

p	q	$p \Rightarrow q$	p	q
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Same pattern, different contexts

$p \Rightarrow q$
p
$\therefore q$

Testing Validity of Argument Forms

- ▷ Truth-Table
- ▷ Tautology

Propositional Logic

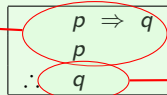
Arguments and Logic Inference

- ▷ Argument Form
- ▷ Inference Rules

premises (or, hypotheses)

p	q	$p \Rightarrow q$	p	q
T	T	T	T	T
T	F	F	T	F
F	T	T	F	T
F	F	T	F	F

Same pattern, different contexts



Conclusion

Testing Validity of Argument Forms

- ▷ Truth-Table
- ▷ Tautology

Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
- ▷ Inference Rules

premises (or, hypotheses)

p	q	$p \Rightarrow q$	p	q
T	T	T	T	T
T	F	F	T	F
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Same pattern, different contexts

$p \Rightarrow q$
 p
 $\therefore q$

Conclusion

Critical Row: premises are all true

Testing Validity of Argument Forms

- ▷ Truth-Table
- ▷ Tautology

Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
- ▷ Inference Rules

p	q	$p \Rightarrow q$	p	q
T	T	T	T	T
T	F	F	T	F
F	T	T	F	T
F	F	T	F	F

Same pattern, different contexts

$$\begin{array}{l}
 p \Rightarrow q \\
 p \\
 \therefore q
 \end{array}$$

Critical Row: premises are all true

Testing Validity of Argument Forms

- ▷ Truth-Table
- ▷ Tautology

An argument form is valid iff the conclusion is true in every **critical row**

Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
- ▷ Inference Rules

Same pattern, different contexts

$$\begin{array}{l} p \Rightarrow q \\ p \\ \therefore q \end{array}$$

Testing Validity of Argument Forms

- ▷ Truth-Table
- ▷ Tautology

An argument form is valid iff it is a tautology when treated as a conditional proposition.

Propositional Logic

Arguments and Logic Inference

- ▷ Argument Form
- ▷ Inference Rules

For example,

$$(p \Rightarrow q) \wedge p \Rightarrow q$$

is indeed a tautology.

Same pattern, different contexts

$$\begin{array}{l} p \Rightarrow q \\ p \\ \therefore q \end{array}$$

Testing Validity of Argument Forms

- ▷ Truth-Table
- ▷ Tautology

An argument form is valid iff it is a tautology when treated as a conditional proposition.

Propositional Logic

Inference Rules

- ▷ Modus Ponens
- ▷ Modus Tollens
- ▷ Elimination
- ▷ Transitivity

Inference Rule: Valid Argument Form

Propositional Logic

Inference Rules

- ▷ Modus Ponens
- ▷ Modus Tollens
- ▷ Elimination
- ▷ Transitivity

Inference Rule: Valid Argument Form

Modus Ponens

(Latin meaning: method of affirming)

$p \Rightarrow q$	If $3 > 1$, then $\sqrt{3} > 1$
p	$3 > 1$
$\therefore q$	$\therefore \sqrt{3} > 1$

Propositional Logic

Inference Rules

- ▷ Modus Ponens
- ▷ Modus Tollens
- ▷ Elimination
- ▷ Transitivity

Inference Rule: Valid Argument Form

Sound Argument

- ▷ Valid argument form
- ▷ Premises are all true

Modus Ponens

From Argument Form to Argument:

Substitute propositions for proposition variables

$p \Rightarrow q$	If $3 > 1$, then $\sqrt{3} > 1$
p	$3 > 1$
$\therefore q$	$\therefore \sqrt{3} > 1$

Propositional Logic

Inference Rules

- ▷ Modus Ponens
- ▷ Modus Tollens
- ▷ Elimination
- ▷ Transitivity

Inference Rule: Valid Argument Form

Sound Argument

- ▷ Valid argument form
- ▷ Premises are all true

Modus Tollens

(Latin meaning: method of denying)

$p \Rightarrow q$	If $\sqrt{.5} > 1$, then $.5 > 1$
$\neg q$	$.5 \leq 1$
$\therefore \neg p$	$\therefore \sqrt{.5} \leq 1$

Propositional Logic

Inference Rules

- ▷ Modus Ponens
- ▷ Modus Tollens
- ▷ Elimination
- ▷ Transitivity

Inference Rule: Valid Argument Form

Sound Argument

- ▷ Valid argument form
- ▷ Premises are all true

Elimination

$p \vee q$	$(x > 1) \vee (y > 1)$
$\neg q$	$x \leq 1$
$\therefore p$	$\therefore y > 1$

Propositional Logic

Inference Rules

- ▷ Modus Ponens
- ▷ Modus Tollens
- ▷ Elimination
- ▷ Transitivity

Inference Rule: Valid Argument Form

Sound Argument

- ▷ Valid argument form
- ▷ Premises are all true

Transitivity

$p \Rightarrow q$	If $x > 1$, then $ x > 1$
$q \Rightarrow r$	If $ x > 1$, then $x \neq 0$
$\therefore p \Rightarrow r$	\therefore If $x > 1$, then $x \neq 0$

Additional Notes Other Useful Inference Rules

Generalization

p	\parallel	If $x > 1$, then $ x > 1$
$\therefore p \vee r$	\parallel	$\therefore (x > 1) \vee (x < 1)$

Specialization

$p \wedge q$	\parallel	$(x > 1) \wedge x \text{ even}$
$\therefore p$	\parallel	$\therefore x > 1$

Additional Notes Other Useful Inference Rules

Proof by Division into Cases

$p \vee q$	$(x \geq 1) \vee (x \leq -1)$
$p \Rightarrow r$	If $x \geq 1$, then $x^2 \geq 1$
$q \Rightarrow r$	If $x \leq -1$, then $x^2 \geq 1$
$\therefore r$	$\therefore x^2 \geq 1$

Exercise: Use the truth-table method to show the above is valid.

Contradiction Rule

$\neg p \Rightarrow F$	
$\therefore p$	

The Story: Finding my glasses?

RK = "I was reading in the kitchen"

RL = "I was reading in the living room"

GK = "My glasses are on the kitchen table"

GC = "My glasses are on the coffee table"

SB = "I saw them at breakfast"

Propositional Logic

Use Inference Rules to Construct Complex Arguments

The Story: Finding my glasses?

RK = "I was reading in the kitchen"

RL = "I was reading in the living room"

GK = "My glasses are on the kitchen table"

GC = "My glasses are on the coffee table"

SB = "I saw them at breakfast"

Known Facts

▶ $RK \Rightarrow GK$

▶ $GK \Rightarrow SB$

▶ $\neg SB$

▶ $RK \vee RL$

▶ $RL \Rightarrow GC$

Propositional Logic

Use Inference Rules to Construct Complex Arguments

The Story: Finding my glasses?

RK = "I was reading in the kitchen"

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Known Facts

- ▶ $RK \Rightarrow GK$
- ▶ $GK \Rightarrow SB$
- ▶ $\neg SB$
- ▶ $RK \vee RL$
- ▶ $RL \Rightarrow GC$

Q.2) Inference rule used is

- A) transitivity
- B) modus tollens
- C) elimination
- D) modus ponens

Step 1

$$\begin{array}{l}
 RK \Rightarrow GK \\
 GK \Rightarrow SB \\
 \therefore RK \Rightarrow SB
 \end{array}$$

Propositional Logic

Use Inference Rules to Construct Complex Arguments

The Story: Finding my glasses?

RK = "I was reading in the kitchen"

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Known Facts

- ▶ $RK \Rightarrow GK$
- ▶ $GK \Rightarrow SB$
- ▶ $\neg SB$
- ▶ $RK \vee RL$
- ▶ $RL \Rightarrow GC$

Q.3) Inference rule used is

- A) transitivity
- B) modus tollens
- C) elimination
- D) modus ponens

Step 2

$$\begin{array}{l}
 RK \Rightarrow SB \\
 \neg SB \\
 \therefore \neg RK
 \end{array}$$

Propositional Logic

Use Inference Rules to Construct Complex Arguments

The Story: Finding my glasses?

RK = "I was reading in the kitchen"

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Known Facts

- ▶ $RK \Rightarrow GK$
- ▶ $GK \Rightarrow SB$
- ▶ $\neg SB$
- ▶ $RK \vee RL$
- ▶ $RL \Rightarrow GC$

Q.4) Inference rule used is

- A) transitivity
- B) modus tollens
- C) elimination
- D) modus ponens

Step 3

$$\begin{array}{l} \neg RK \\ RL \vee RK \\ \hline \therefore RL \end{array}$$

Propositional Logic

Use Inference Rules to Construct Complex Arguments

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GK = "My glasses are on the kitchen table"

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Known Facts

- ▶ $RK \Rightarrow GK$
- ▶ $GK \Rightarrow SB$
- ▶ $\neg SB$
- ▶ $RK \vee RL$
- ▶ $RL \Rightarrow GC$

Q.5) Inference rule used is

- A) transitivity
- B) modus tollens
- C) elimination
- D) modus ponens

Step 4

$$\begin{array}{l} RL \Rightarrow GC \\ RL \\ \therefore GC \end{array}$$

Propositional Logic

Use Inference Rules to Construct Complex Arguments

The Story: Finding my glasses?

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Known Facts

- ▶ $RK \Rightarrow GK$
- ▶ $GK \Rightarrow SB$
- ▶ $\neg SB$
- ▶ $RK \vee RL$
- ▶ $RL \Rightarrow GC$

Put Together: the argument

(1)	$RK \Rightarrow GK$	Given
(2)	$GK \Rightarrow SB$	Given
(3)	$RK \Rightarrow SB$	transitivity
(4)	$\neg SB$	Given
(5)	$\neg RK$	modus tollens: (3), (4)
(6)	$RL \vee RK$	Given
(7)	RL	elimination: (5), (6)
(8)	$RL \Rightarrow GC$	Given
\therefore	GC	modus ponens: (7), (8)

Propositional Logic

Use Inference Rules to Construct Arguments

Let x, y, z be three particular integers, satisfying

▷ $x < y < z$

▷ $z \leq 1$

▷ $|x| > 1$

To show that $x < -1$.

True Propositions

(a) $(x > 1) \Rightarrow (y > 1)$

(b) $(y > 1) \Rightarrow (z > 1)$

(c) $(z \leq 1)$

(d) $(x < -1) \vee (x > 1)$

Propositional Logic

Use Inference Rules to Construct Arguments

Let x, y, z be three particular integers, satisfying

- ▷ $x < y < z$
- ▷ $z \leq 1$
- ▷ $|x| > 1$

To show that $x < -1$.

True Propositions

- (a) $(x > 1) \Rightarrow (y > 1)$
- (b) $(y > 1) \Rightarrow (z > 1)$
- (c) $(z \leq 1)$
- (d) $(x < -1) \vee (x > 1)$

The Argument

(1)	$(x > 1) \Rightarrow (y > 1)$	Given
(2)	$(y > 1) \Rightarrow (z > 1)$	Given
(3)	$(x > 1) \Rightarrow (z > 1)$	Transitivity
(4)	$\neg(z > 1)$	Given
(5)	$\neg(x > 1)$	Modus Tollens: (3),(4)
(6)	$(x < -1) \vee (x > 1)$	Given
∴	$(x < -1)$	Elimination: (5), (6)

Propositional Logic



Use Inference Rules to Construct Complex Arguments

Knights always tell truth and Knaves always tell lie. *A* says “*B* is a knight”

B says “*A* and I are of opposite type”.

What are *A* and *B*?

Propositional Logic



Use Inference Rules to Construct Complex Arguments

Knights always tell truth and Knaves always tell lie. A says “ B is a knight”
 B says “ A and I are of opposite type”.

What are A and B ?

- ▶ Suppose A is a knight. B is a knight and “ A and I are of opposite types” is true. A contradiction!

Propositional Logic



Use Inference Rules to Construct Complex Arguments

Knights always tell truth and Knaves always tell lie. A says “ B is a knight”
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What are A and B ?

- ▶ Suppose A is a knight. B is a knight and “ A and I are of opposite types” is true. A contradiction!
- ▶ $\therefore A$ is NOT a knight

Knights always tell truth and Knaves always tell lie. A says “ B is a knight”

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What are A and B ?

- ▶ Suppose A is a knight. B is a knight and “ A and I are of opposite types” is true. A contradiction!
- ▶ $\therefore A$ is NOT a knight

Q.6) The inference rule used is

- A) Modus Ponens
- B) Elimination
- C) Contradiction
- D) Generalization

Propositional Logic



Use Inference Rules to Construct Complex Arguments

Knights always tell truth and Knaves always tell lie. A says “ B is a knight”
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What are A and B ?

- ▶ Suppose A is a knight. B is a knight and “ A and I are of opposite types” is true. A contradiction!
- ▶ $\therefore A$ is NOT a knight
- ▶ A is a knight or A is a knave

Knights always tell truth and Knaves always tell lie. A says “ B is a knight”
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What are A and B ?

- ▶ Suppose A is a knight. B is a knight and “ A and I are of opposite types” is true. A contradiction!
- ▶ $\therefore A$ is NOT a knight
- ▶ A is a knight or A is a knave
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Knights always tell truth and Knaves always tell lie. A says “ B is a knight”
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What are A and B ?

- ▶ Suppose A is a knight. B is a knight and “ A and I are of opposite types” is true. A contradiction!
- ▶ $\therefore A$ is NOT a knight
- ▶ A is a knight or A is a knave
- ▶ $\therefore A$ is a knave

Q.7) The inference rule used is

- A) Modus Ponens
- B) Elimination
- C) Contradiction
- D) Generalization

Additional Notes Common Errors: Logical Fallacies

Converse Error	Inverse Error
$p \Rightarrow q$ q $\therefore p$	$p \Rightarrow q$ $\neg p$ $\therefore \neg q$
(affirming the consequence)	(denying the antecedent)

Concepts

- ▷ Propositions and Connectives
- ▷ Truth Values and Truth Tables
- ▷ Logic Equivalence and Laws

Arguments and Logic Inference

- ▷ Argument and Argument Form
- ▷ Inference Rules

Modelling and Inferences

- ▷ Wumpus-world, graph coloring
- ▷ iClikier exercises
- ▷ Reasoning and programming

- ▷ Statements \Rightarrow Propositions
- ▷ Negation/Contrapositive of Compound (Conditional) Propositions
- ▷ Use of truth tables and/or logic laws

Concepts

- ▷ Propositions and Connectives
- ▷ Truth Values and Truth Tables
- ▷ Logic Equivalence and Laws

Arguments and Logic Inference

- ▷ Argument and Argument Form
- ▷ Inference Rules

- ▷ Truth tables and Valid Argument Forms
- ▷ Construct logic arguments using inference rules

Modelling and Inferences

- ▷ Wumpus-world, graph coloring
- ▷ iClikier exercises
- ▷ Reasoning and programming

Additional Notes

Announcement 1: Quiz 1 (Monday, Feb 5th)

▷ Topics:

- Propositional Logic
- Basics of Predicate Logic

Logics and Automated Reasoning in AI

Knowledge Representation: KB (Knowledge Base)

Query and Question

Algorithmic Problem

Inference Engine

Logics and Automated Reasoning in AI

Knowledge Representation: KB (Knowledge Base)

Query and Question

Conjunction of facts and inference rules

Algorithmic Problem

Inference Engine

Logics and Automated Reasoning in AI

Knowledge Representation: KB (Knowledge Base)

Query and Question

Given a statement Q , is $KB \Rightarrow Q$ true?

Algorithmic Problem

Inference Engine

Logics and Automated Reasoning in AI

Knowledge Representation: KB (Knowledge Base)

Query and Question

Given a statement Q , is $KB \Rightarrow Q$ true?

Algorithmic Problem

Inference Engine

Is $KB \cap \neg Q$ False?

Why does this work?

Logics and Automated Reasoning in AI

Knowledge Representation: KB (Knowledge Base)

Query and Question

Given a statement Q , is $KB \Rightarrow Q$ true?

Algorithmic Problem

Inference Engine

A solver for the satisfiability problem.

E.g., $F = (x_1 \vee x_2) \wedge (x_8 \vee \neg x_{15} \vee x_{50}) \cdots (\neg x_{99} \vee x_{120})$



- ▶ **Imperative Programming**
- ▶ **Declarative Programming**

Propositional Logic and Beyond

Logic and Programming

- ▶ **Imperative Programming**
- ▶ **Declarative Programming**

Java, C, C#, etc.

Instructions telling computers “how to do things”

Propositional Logic and Beyond

Logic and Programming

- ▶ **Imperative Programming**
- ▶ **Declarative Programming**

Java, C, C#, etc.

Instructions telling computers “how to do things”

Prolog, LISP.

Statements telling computers “what to do”

Propositional Logic and Beyond

Logic and Programming

- ▶ **Imperative Programming**
- ▶ **Declarative Programming**

Java, C, C#, etc.

Instructions telling computers “how to do things”

Prolog, LISP.

Statements telling computers “what to do”

Prolog - Programming in Logic

- ▶ Developed in 1970s
- ▶ Used in some areas of AI

Propositional Logic and Beyond

Logic and Programming

- ▶ **Imperative Programming**
- ▶ **Declarative Programming**

```
% A Mini Knowledgebase about Java Classes

%Rules
is-a(X, Y):- extends(X, Y).
is-a(X, Y):- extends(X, Z), is-a(Z, Y).
is-a(X, Y):- interface(Y), implements(X, Y).

%Facts
extends(mylist, arraylist).
extends(arraylist, abstractlist).
extends(error, throwable).
extends(exception, throwable).

implements(abstractlist, list).
implements(throwable, serializable).
interface(list).
```

Prolog - Programming in Logic

- ▶ Developed in 1970s
- ▶ Used in some areas of AI

Propositional Logic and Beyond

Logic and Programming

- ▶ **Imperative Programming**
- ▶ **Declarative Programming**

Propositions
E.g. "mylist extends arraylist"

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```


Propositional Logic and Beyond

Logic and Programming

- **Imperative Programming**
- **Declarative Programming**

Propositions?

E.g. "X is a Y if X extends Y"

Remember "He is a student"

Propositions

E.g. "mylist extends arraylist"

% A Mini Knowledgebase about Java Classes

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Propositional Logic and Beyond

Logic and Programming

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```

Prolog Program

A collection of facts and rules

- ▶ **universal conditional statements**
- ▶ Next topic

Demo time! Go to SWISH

Prolog: Program Examples

% A Mini Knowledgebase about Java Classes

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interface(list).

Prolog: Program Examples - Sudoku

```
%-----
%Program Example: sudoku with CLP. It solves every
%sudoku puzzle and to each puzzle, it can find all possible
%solutions. (Basic idea from the SWI-Prolog Manual)
%-----

:- use_module(library(clpfd)). %import the CLP module

board([[_,_,3,_,_,_,_,_,_],
        [_,_,_,_,_,3,_,8,5],
        [_,_,1,_,2,_,_,_,_],
        [_,_,_,5,_,7,_,_,_],
        [_,_,4,_,_,_,1,_,_],
        [_,9,_,_,_,_,_,_,_],
        [5,_,_,_,_,_,_,7,3],
        [_,_,2,_,1,_,_,_,_],
        [_,_,_,_,4,_,_,_,6]]).
```

```
%-----  
% the predict to solve soduko  
% -----  
solve:-  
    board(Rows),  
    length(Rows, 9), maplist(length_(9), Rows),  
    append(Rows, Vs), Vs ins 1..9,  
  
    sudoku(Rows),  
  
    %solve the puzzle by labelling the variables  
    label(Vs),  
  
    %output the solution  
    maplist(writeln, Rows).
```

