

# COSC221 Assignment 3

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Part A) **Remark:**

- $\mathbb{Z}^+ = \{\forall n \in \mathbb{Z} : n \geq 1\} = \{1, 2, 3, \dots\}$
- $\mathbb{Z}_{even}^+ = \{n \in \mathbb{Z}^+ : n \bmod 2 = 0\} = \{2, 4, 6, \dots\}$
- $\mathbb{Z}_{odd}^+ = \mathbb{Z}^+ \setminus \mathbb{Z}_{even}^+ = \{1, 3, 5, \dots\}$

*Proof.* By induction on  $n \in \mathbb{Z}^+$ . Here, we choose  $\mathbb{Z}^+$  because it is more intuitive to have a positive integer of  $n \times n$  chessboard.

Let  $S(n) = n \times n$  chessboard, where there exists a rook's tour from top-left to bottom-left.

**Base Case:**  $n = 1$ . It is known that a  $1 \times 1$  chessboard is trivial, but ambiguous, so we can still proceed:

$S(1) = 1 \times 1$  chessboard, so there exists one cell only on that chessboard, which makes it the top-left and the bottom-left cell at the same time, where the rook starts and ends at the same cell. Hence, there exists the rook's tour.

**Inductive Hypothesis:** If our constructed  $S(n)$  holds, then there exists some  $n \in \mathbb{Z}^+$  such that it satisfies  $S(n + 1)$  as well.

**Inductive Step:**  $n \rightarrow n + 1$ .

$S(n + 1) = (n + 1) \times (n + 1)$  chessboard. By the given information on page 234 for rook's tour and rook's move, we let  $(x, y, k) \in \mathbb{Z}^+$ , then it is true that rook can start at  $\langle x, y \rangle$ . Then from  $n \times n$  chessboard to  $(n + 1) \times (n + 1)$  chessboard, there are two possible movements:

- **Vertical movement:**

$$\langle x, y \pm k \rangle \rightarrow \langle x, n + 1 \rangle \Rightarrow y \pm k = n + 1 \Leftrightarrow y = n + 1 \pm k$$

- **Horizontal movement:**

$$\langle x \pm k, y \rangle \rightarrow \langle n + 1, y \rangle \Rightarrow x \pm k = n + 1 \Leftrightarrow x = n + 1 \pm k$$

Now, we know that, if we let

$$k = \min(\mathbb{Z}^+) = 1$$

and,

$$\min(x) = \min(y) = n + 1 - k$$

therefore,

$$\min(x) = \min(y) = n + 1 - 1 = n$$

Then, our movement for both vertical and horizontal are both within the bounds of  $\mathbb{Z}^+$ , as we said above,  $k \in \mathbb{Z}^+$ . Now, let these coordinates be true that  $\langle 1, 1 \rangle \Rightarrow$  top-left,  $\langle n + 1, 1 \rangle \Rightarrow$  top-right,  $\langle 1, n + 1 \rangle \Rightarrow$  bottom-left and,  $\langle n + 1, n + 1 \rangle \Rightarrow$  bottom-right.

Now to traverse through the board, there are two possible cases:

**Case  $\alpha$ :**  $(n + 1) \times (n + 1)$  chessboard where  $(n + 1) \in \mathbb{Z}_{\text{even}}^+$

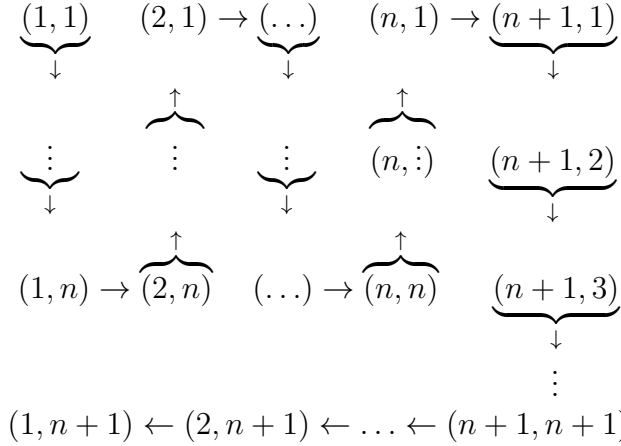
Then we let the rook start at  $(1, 1)$ , we follow a horizontal “zig-zag” pattern below:

$$\begin{array}{c} (1, 1) \rightarrow (2, 1) \rightarrow \dots \rightarrow (n, 1) \rightarrow \underbrace{(n + 1, 1)}_{\downarrow} \\ \underbrace{(1, 2)}_{\downarrow} \leftarrow (2, 2) \leftarrow \dots \leftarrow (n, 2) \leftarrow (n + 1, 2) \\ (1, 3) \rightarrow (2, 3) \rightarrow \dots \rightarrow (n, 3) \rightarrow \underbrace{(n + 1, 3)}_{\downarrow} \\ \vdots \\ (1, n + 1) \leftarrow (2, n + 1) \leftarrow \dots \leftarrow (n + 1, n + 1) \end{array}$$

With this visual, we can clearly see that the rook arrives at the left side every time  $y \in \mathbb{Z}_{\text{even}}^+$ , so it is entirely possible that the rook can end up on the bottom-left side, because  $\max(y) = (n + 1) \in \mathbb{Z}_{\text{even}}^+$ .

**Case  $\beta$ :**  $(n + 1) \times (n + 1)$  chessboard where  $(n + 1) \in \mathbb{Z}_{\text{odd}}^+$

Then we let the rook start at  $(1, 1)$ , we follow a vertical “zig-zag” pattern below:



With this visual, we can clearly see that whenever  $x \in \mathbb{Z}_{\text{odd}}^+$ , it will be making a downwards movement, so since  $\max(x) = (n+1) \in \mathbb{Z}_{\text{odd}}^+$ , then we know that  $x = n \in \mathbb{Z}_{\text{even}}^+$ , that means that it will be making an upwards movement so that it can turn right towards  $(n+1, 1)$  where it can move down towards  $(n+1, n)$ . Note that we have never really visited  $y = n+1$  for these traversal, so let it be known that, after reaching  $(n+1, n+1)$ , it can go directly to  $(1, n+1)$ , where the rook will be at the bottom-left of the chessboard.

$\therefore$  By mathematical induction, and by **Case  $\alpha$**  and  **$\beta$** , it is true that the rook's tour exists for the  $n \times n$  chessboard, for some  $n \in \mathbb{Z}^+$ .  $\square$

Part B) **Remark:** We know that in a simple graph,

$$directed \Rightarrow density = \frac{|E|}{\frac{n(n-1)}{2}} = \frac{2|E|}{n(n-1)}$$

$$undirected \Rightarrow density = \frac{|E|}{n(n-1)}$$

Where  $n$  is the vertex or node, and  $|E|$  is the edges.

**11.36** As a function of  $n$  the density of an  $n$ -node path means,  $|E| \Rightarrow (n-1)$ , where  $n$  is vertices, so we have,

$$f(n) = \frac{2(n-1)}{n(n-1)} = \frac{2}{n}$$

**11.37** As a function of  $n$  the density of an  $n$ -node cycle means,  $|E| \Rightarrow n$ , where  $n$  is vertices, so we have,

$$f(n) = \frac{2n}{n(n-1)} = \frac{2}{n-1}$$

**11.38** As a function of  $n$  the density of a graph that consists of  $\frac{n}{3}$  disconnected triangles,  $|E| \Rightarrow n$ , where  $n$  is vertices, so we have,

$$f(n) = \frac{2n}{n(n-1)} = \frac{2}{n-1}$$

which is essentially the same as a cycle.

Part C) **Remark:** Let it be known that clique 1, clique 2 and clique 3 represents the “clusters” of nodes from left to right respectively in **Figure 11.33d**

**11.54** We can see that each “clusters” or cliques represents its own movie, then there must be at least, 3 movies, but it is also shown that there is 2 curves on the top of the graph, where the actor from the clique 3, knows actor from clique 1 and clique 2. which means that there must exists at least another 2 movies. Therefore, there must exists, at the minimum, five movies that could be generated by this collaboration network.

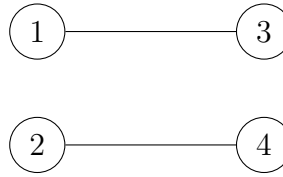
**11.55** No, it is uncertain. It is entirely possible that there are more than 5 movies generated in this graph because each cliques can have its own subset, where the actors are in another movie different than the minimum 5 movies.

Part D) **11.95**

*Proof.* By counterexample, assume  $G$  is a bipartite graph where,

$$G = \langle L \cup R, E \rangle$$

and where,  $|L| = |R|$ , so we can take,  $L = \{1, 2\}$  and  $R = \{3, 4\}$  and since all nodes from  $L$  and  $R$  has at least one neighbours, we can construct  $G$  such that  $E = \{\{1, 3\}, \{2, 4\}\}$ , visually:



We can see that the graph we made,  $G$ , consists of two disjoint sets which is absurd, because the original statement say that it is connected.

$\therefore$  It is disproved that if  $G = \langle L \cup R, E \rangle$  where  $G$  is a bipartite graph with  $|L| = |R|$ , then it is impossible for  $G$  to be a connected graph, if there is at minimum, one neighbour for every node in  $L$  and  $R$ .  $\square$

End of Assignment 3.