

## Assignment 2

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### Part A) Problems and Proofs (6 points)

Rewrite the following claim as a universal proposition in symbolic form:

"The algorithm bruteForceSearch() is a linear-time algorithm".

and use the strategy of generalizing from generic particular to prove that the claim is correct. For a detailed explanation of the claim and related concepts in algorithm analysis, see the figure on the next page.

Generalizing:

Let  $A(x)$ :  $x$  is an algorithm  
 $B(x)$ :  $x$  is brute ForceSearch()  
 $L(x)$ :  $x$  is a linear algorithm.

Then the claim can be rewritten as

$$(\forall x \in \mathbb{Z}), (A(x) \wedge B(x)) \Rightarrow L(x)$$

Generic particular : brute Force Search()  $(\star)$

From Figure 1, It shows that  $(\star)$  Iterates over array values [] until it finds the value. So, time complexity is  $O(n)$  where  $n$  is size of array, given worst case, it will still be  $O(n)$  where value is at the last element of values []. Therefore, it is a Linear - time algorithm.

$$\therefore (A(x) \wedge B(x)) \Rightarrow L(x) \quad \square$$

### Part B) Proof by Cases (6 points)

Use the method of proof by cases to prove the following statement

$$\forall n \geq 3, n^2 - 3n + 3 \text{ is odd.}$$

```
public int bruteForceSearch(int[] values, int value){  
    int numOfSteps = 0;  
    boolean found = false;  
  
    while( !found && numOfSteps < values.size()){  
        numOfSteps++;  
        if(values[numOfSteps] == value) found = true;  
    }  
  
    return numOfSteps;  
}
```

Figure 1

Proof by cases.

Case 1:  $n \in \mathbb{Z}_{\text{odd}} : n \geq 3$ ,

Then,  $n = 2k + 1$ .

We have,  $(2k+1)^2 - 3(2k+1) + 3 \quad \text{for } k \geq 3$

$$\Leftrightarrow 4k^2 + 4k + 1 - 6k - 3 + 3$$

$$\Leftrightarrow 4k^2 - 2k + 1$$

$$\Leftrightarrow 2(\underbrace{2k^2 - k}_{\in \mathbb{Z}}) + 1 \quad \square$$

Case 2: Let  $n \in \mathbb{Z}_{\text{even}} : n \geq 3$ ,

Then, by def,  $n = 2k$ .

$$\Rightarrow (2k)^2 - 3(2k) + 3 \quad \text{for } k \geq 3$$

$$\Leftrightarrow 4k^2 - 6k + 3$$

$$\Leftrightarrow 2(2k^2 - 3k) + 3$$
  
$$\in \mathbb{Z}.$$

$\therefore$  The statement in Part B is true by case 1 & 2,  
for  $\forall n \geq 3$ .  $\square$

**4.52** Let  $n$  and  $m$  be integers. If  $nm$  is not evenly divisible by 3, then neither  $n$  nor  $m$  is evenly divisible by 3. (In fact, the converse is true too, but you don't have to prove it.)

Proof by contra positive

" $n$  and  $m$  is evenly divisible by 3  $\Rightarrow nm$  is divisible by 3."

Let  $(n, m) \in \mathbb{R}$ ,

then by def,  $n = 3k$  and  $m = 3k$ . ( $k \in \mathbb{Z}$ )

$$nm = 3k \cdot 3k = (3k)^2 = 9k^2 = 3(\underbrace{3k^2}_{\in \mathbb{Z}})$$

$\therefore (n \in (3k) \wedge m \in (3k)) \Rightarrow (nm) \in (3k)$  where  $k \in \mathbb{Z}$ .  $\square$

4.56 Let  $x, y$  be positive real numbers. If  $x^2 - y^2 = 1$ , then  $x$  or  $y$  (or both) is not an integer.

Proof by contradiction

Let's assume  $(x, y) \in \mathbb{Z}^+ : x^2 - y^2 = 1$  (for the sake of contradiction)

$$\Rightarrow (x+y)(x-y) = 1$$

$$\Leftrightarrow (x+y) = 1 \wedge (x-y) = 1$$

$$\Leftrightarrow (x = 1-y) \wedge (x = 1+y)$$

$$\Leftrightarrow 1-y = 1+y \Leftrightarrow 2y = 0 \Leftrightarrow y = 0 !!!$$

which is absurd because we assumed  $(x, y) \in \mathbb{Z}^+$

$\therefore$  The original claim is true b/c assuming  $(x, y) \in \mathbb{Z}^+$

creates a contradiction  $\square$

$$5.1 \sum_{i=0}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

Proof by induction

Base Case: Let  $S(n) = (5.1)$  where  $n=0$ .

$$S(n) = \sum_{i=0}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

$$S(0) = \sum_{i=0}^0 i^3 = \frac{0(0+1)(2(0)+1)}{6}$$

$$0^3 = \frac{0}{6}$$

$$0 = 0 \Rightarrow LHS = RHS$$

Inductive Hypothesis: If  $S(n)$  holds, then  $S(n+1)$  will hold  
for some  $n \geq 0$ .

Inductive Step:  $n = n+1$

$$S(n+1) \Rightarrow \text{RHS} = \frac{(n+1)((n+1)+1)(2(n+1)+1)}{6} = \frac{(n+1)(n+2)(2n+3)}{6}$$
$$\text{RHS} = \frac{(n^2 + 3n + 2)(2n + 3)}{6} = \frac{2n^3 + 9n^2 + 13n + 6}{6}$$

By Ind Hypo, LHS =

$$\sum_{i=0}^n i^2 + (n+1)^2 = \frac{n(n+1)(2n+1)}{6} + (n+1)^2$$

Now, we simplify RHS

$$= \frac{(n^2+n)(2n+1)}{6} + \frac{6(n^2+2n+1)}{6}$$

$$= \frac{2n^3 + n^2 + 2n^2 + n + 6n^2 + 12n + 6}{6}$$

$$= \frac{2n^3 + 9n^2 + 13n + 6}{6}$$

RHS = LHS.

∴ By Principle of Math Induction, we have proven that  
scr is true for some  $n \geq 0$ .  $\square$