

COSC221 Assignment 3

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Part A) *Proof.* By induction on $n \in \mathbb{Z}^+$. Here, we choose \mathbb{Z}^+ because it is intuitive to have a positive integer of $n \times n$ chessboard.

Remark: \mathbb{Z}^+ is the set of all positive integers; $\mathbb{Z}^+ = \{1, 2, 3, \dots, n\}_{n \rightarrow \infty}$.

Base Case: $n = 1$

$S(1) = 1 \times 1$ chessboard, so there exists one cell only on that chessboard, which makes it the top-left and the bottom-left cell, where the rook starts and ends at the same cell. Hence, there exists the rook's tour, even if $n = 1$ is a little bit trivial.

Inductive Hypothesis: If our constructed $S(n)$ holds, then there exists some $n \in \mathbb{Z}^+$ such that it satisfies $S(n + 1)$ as well.

Inductive Step: $n \rightarrow n + 1$

$S(n + 1) = (n + 1) \times (n + 1)$ chessboard. By the given information on page 234 for rook's tour and rook's move, we let $(r, c, x) \in \mathbb{Z}^+$, then it is true that rook can start at $\langle r, c \rangle$. Then for $n \times n$ chessboard to $(n + 1) \times (n + 1)$ chessboard, there are two possible movement,

Vertical movement $\Rightarrow \langle r, c \pm x \rangle \rightarrow \langle r, n + 1 \rangle$, so this means that $c \pm x = n + 1 \Leftrightarrow c = n + 1 \pm x$

Horizontal movement $\Rightarrow \langle r \pm x, c \rangle \rightarrow \langle n + 1, c \rangle$, so this also means that $r \pm x = n + 1 \Leftrightarrow r = n + 1 \pm x$

Now, we know that, if we let

$$x = \min(\mathbb{Z}^+) = 1$$

and,

$$\min(r) = \min(c) = n + 1 - x$$

therefore,

$$\min(r) = \min(c) = n + 1 - 1 = n$$

Then, our movement for both vertical and horizontal are both within the bounds of \mathbb{Z}^+ , as we said above, $n \in \mathbb{Z}^+$. Now, let these coordinates be true that $\langle 1, 1 \rangle \Rightarrow$ top-left, $\langle n + 1, 1 \rangle \Rightarrow$ top-right, $\langle 1, n + 1 \rangle \Rightarrow$ bottom-left and, $\langle n + 1, n + 1 \rangle \Rightarrow$ bottom-right.

Now to traverse through the board, there are two possible cases:

Case α : $(n + 1) \times (n + 1)$ chessboard where $(n + 1) \in \mathbb{Z}_{\text{even}}$

Then if we let rook start at $\langle 1, 1 \rangle$, then we follow a horizontal “zig-zag” pattern below:

$$\begin{array}{c}
 (1, 1) \rightarrow (2, 1) \rightarrow \dots \rightarrow (n, 1) \rightarrow \underbrace{(n + 1, 1)}_{\downarrow} \\
 \underbrace{(1, 2)}_{\downarrow} \leftarrow (2, 2) \leftarrow \dots \leftarrow (n, 2) \leftarrow (n + 1, 2) \\
 (1, 3) \rightarrow (2, 3) \rightarrow \dots \rightarrow (n, 3) \rightarrow \underbrace{(n + 1, 3)}_{\downarrow} \\
 \vdots \\
 (1, n + 1) \leftarrow (2, n + 1) \leftarrow \dots \leftarrow (n + 1, n + 1)
 \end{array}$$

By this visual, we can clearly see that the rook arrives at the left side every time $c \in \mathbb{Z}_{\text{even}}$, so it is entirely possible that the rook can end up on the bottom-left side, because $c = (n + 1) \in \mathbb{Z}_{\text{even}}^+$.

Case β : $(n + 1) \times (n + 1)$ chessboard where $(n + 1) \in \mathbb{Z}_{\text{odd}}^+$

Then if we let rook start at $\langle 1, 1 \rangle$, then we follow a vertical “zig-zag” pattern below:

$$\begin{array}{ccccc}
 \underbrace{(1, 1)}_{\downarrow} & (2, 1) \rightarrow \underbrace{(\dots)}_{\downarrow} & (n, 1) \rightarrow \underbrace{(n + 1, 1)}_{\downarrow} & & \\
 & \uparrow & & \uparrow & \\
 \vdots & \vdots & \vdots & (n, \vdots) & \underbrace{(n + 1, 2)}_{\downarrow} \\
 & \downarrow & & &
 \end{array}$$

$$\begin{array}{c}
(1, n) \rightarrow \overbrace{(2, n)}^{\uparrow} \quad (\dots) \rightarrow \overbrace{(n, n)}^{\uparrow} \quad \underbrace{(n+1, n)}_{\downarrow} \\
\vdots \\
(1, n+1) \leftarrow (2, n+1) \leftarrow \dots \leftarrow (n+1, n+1)
\end{array}$$

By this visual, we can clearly see that whenever $r \in \mathbb{Z}_{\text{odd}}$, it will be making a downwards movement, so since $(n+1) \in \mathbb{Z}_{\text{odd}}^+$, we know that $n \in \mathbb{Z}_{\text{even}}^+$, that means that it will be making an upwards movement so that it can turn right towards $(n+1, 1)$ where it can start moving down again. Note that we have never really visited $c = n+1$ for these traversal, so let it be known that, after reaching $(n+1, n+1)$, it can go directly to $(1, n+1)$, where the rook will be at the bottom-left of the chessboard.

\therefore By mathematical induction, and by **Case α** and **β** , it is true that the rook's tour exists for the $n \times n$ chessboard, for some $n \in \mathbb{Z}^+$.

□

Part B) **11.36** We know that in a simple graph,

$$directed \Rightarrow density = \frac{|E|}{\frac{n(n-1)}{2}} = \frac{2|E|}{n(n-1)}$$

$$undirected \Rightarrow density = \frac{|E|}{n(n-1)}$$

Where n is the vertex or node, and $|E|$ is the edges.

As a function of n the density of an n -node path means, $|E| \Rightarrow (n-1)$, where n is vertices, so we have,

$$f(n) = \frac{2(n-1)}{n(n-1)} = \frac{2}{n}$$

11.37 As a function of n the density of an n -node cycle means, $|E| \Rightarrow n$, where n is vertices, so we have,

$$f(n) = \frac{2n}{n(n-1)} = \frac{2}{n-1}$$

11.38 As a function of n the density of a graph that consists of $\frac{n}{3}$ disconnected triangles, $|E| \Rightarrow n$, where n is vertices, so we have,

$$f(n) = \frac{2n}{n(n-1)} = \frac{2}{n-1}$$

which is essentially the same as a cycle.

Part C) **Remark:** Let it be known that clique 1, clique 2 and clique 3 represents the “clusters” of nodes from left to right respectively in **Figure 11.33d**

11.54 We can see that each “clusters” or cliques represents its own movie, then there must be at least, 3 movies, but it is also shown that there is two curves on the top of the graph, where the actor from the clique 3, knows actor from clique 1 and clique 2. which means that there must exists at least another 2 movies. Therefore, there must exists, at the minimum, five movies that could be generated by this collaboration network.

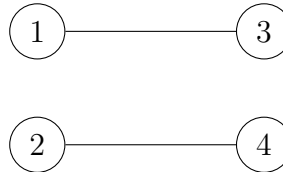
11.55 No, it is uncertain. It is entirely possible that there are more than 5 movies generated in this graph because each cliques can have its own subset, where the actors are in another movie different than the minimum 5 movies.

Part D) **11.95**

Proof. By counterexample, assume G is a bipartite graph where,

$$G = \langle L \cup R, E \rangle$$

and where, $|L| = |R|$, so we can take, $L = \{1, 2\}$ and $R = \{3, 4\}$ and since all nodes from L and R has at least one neighbours, we can construct G such that $E = \{\{1, 3\}, \{2, 4\}\}$, visually:



and we can see that the graph we made, G , consists of two disjoint sets which is absurd, because the original statement say that it is connected.

\therefore It is disproved that if $G = \langle L \cup R, E \rangle$ where G is a bipartite graph with $|L| = |R|$, then it is impossible for G to be a connected graph, if there is at minimum, one neighbour for every node in L and R .x' \square

End of Assignment 3.