

Problem §1 Choose any three problems from 1.2(a-i) as a warmup. Then complete Exercise 1.3.

- 1.2: Use truth tables to prove:
 - 1.2.a: $P \iff \neg(\neg P)$.
 - 1.2.c: $(P \Rightarrow Q) \iff (\neg Q \Rightarrow \neg P)$.
 - 1.2.d: $(P \Rightarrow Q) \iff (\neg P) \vee Q$.
- 1.3: Let P and Q be statements.

(a) Prove that

$$P \vee \neg P$$

is true, and explain why this justifies the Law of the Excluded Middle (which states that exactly one of P and $\neg P$ is true).

(b) Prove that

$$(\neg Q \Rightarrow \neg P) \Rightarrow (P \Rightarrow Q)$$

is true, and explain why this justifies the method of Proof by Contradiction (which states that in order to prove that P is true, it suffices to show that $\neg P$ is false).

Solution:

1.2.a

P	$\neg P$	$\neg(\neg P)$
T	F	T
F	T	F

•

1.2.c

P	Q	$P \Rightarrow Q$	$\neg Q \Rightarrow \neg P$
T	T	T	T
T	F	F	F
F	T	T	T
F	F	T	T

•

1.2.d

P	Q	$\neg P$	$P \Rightarrow Q$	$\neg P \vee Q$
T	T	F	T	T
T	F	F	F	F
F	T	T	T	T
F	F	T	T	T

•

1.3.a

P	$\neg P$	$P \underline{\vee} \neg P$
T	F	T
F	T	T

- Since the statement is true regardless of P , $P \underline{\vee} \neg P$ is true. This also justifies the Law of the Excluded Middle, as an XOR statement is true only when one, but not both, of the statements are true; hence only one of P and $\neg P$ may be true at once in order for $P \underline{\vee} \neg P$ to be true.

1.3.b

P	Q	$\neg P$	$\neg Q$	$P \Rightarrow Q$	$\neg Q \Rightarrow \neg P$	$(\neg Q \Rightarrow \neg P) \Rightarrow (P \Rightarrow Q)$
T	T	F	F	T	T	T
T	F	F	T	F	F	T
F	T	T	F	T	T	T
F	F	T	T	T	T	T

- Since the statement is true regardless of P or Q , $(\neg Q \Rightarrow \neg P) \Rightarrow (P \Rightarrow Q)$ is true. This also justifies Proof by Contradiction: if $\neg P$ is false, then P is necessarily true in order for the above statement to be true as well (alternatively, from the Law of the Excluded Middle, $\neg P$ being false necessarily implies that P is true).

Problem §2 Complete Exercise 1.7: *Prove each of the following formulas:*

(a) $S \cap (T \cup U) = (S \cap T) \cup (S \cap U)$

(b) $S \cup (T \cap U) = (S \cup T) \cap (S \cup U)$

(c) Suppose $S, T \subset U$. Then

$$(S \cup T)^c = S^c \cap T^c \text{ and } (S \cap T)^c = S^c \cup T^c.$$

(d) $S \Delta T = (S \cup T) \setminus (S \cap T) = (S \setminus T) \cup (T \setminus S)$

Solution:

(a) *Proof.* Let $e \in S \cap (T \cup U)$. Then

$$\begin{aligned} & (e \in S) \wedge (e \in T \vee e \in U) \\ &= (e \in S \wedge e \in T) \vee (e \in S \wedge e \in U). \end{aligned}$$

Thus $e \in (S \cap T) \cup (S \cap U)$, and so $S \cap (T \cup U) \subset (S \cap T) \cup (S \cap U)$.

Conversely, let $e \in (S \cap T) \cup (S \cap U)$. Then

$$\begin{aligned} & (e \in S \wedge e \in T) \vee (e \in S \wedge e \in U) \\ &= e \in S \wedge (e \in T \vee e \in U). \end{aligned}$$

Thus, $e \in S \cap (T \cup U)$, and so $(S \cap T) \cup (S \cap U) \subset S \cap (T \cup U)$.

Since both are subsets of each other, $S \cap (T \cup U) = (S \cap T) \cup (S \cap U)$. \square

(b) *Proof.* Let $e \in S \cup (T \cap U)$. Then

$$\begin{aligned} & e \in S \vee (e \in T \wedge e \in U) \\ &= (e \in S \vee e \in T) \wedge (e \in S \vee e \in U). \end{aligned}$$

Thus $e \in (S \cup T) \cap (S \cup U)$, and so $S \cup (T \cap U) \subset (S \cup T) \cap (S \cup U)$.
Conversely, let $e \in (S \cup T) \cap (S \cup U)$. Then

$$\begin{aligned} & (e \in S \vee e \in T) \wedge (e \in S \vee e \in U) \\ & e \in S \vee (e \in T \wedge e \in U). \end{aligned}$$

Thus, $e \in S \cup (T \cap U)$, and so $(S \cup T) \cap (S \cup U) \subset S \cup (T \cap U)$.

Since both are subsets of each other, $s \cup (t \cap u) = (s \cup t) \cap (s \cup u)$. \square

(c) *Proof.* Let $e \in (S \cup T)^c$. Then

$$\begin{aligned} & e \in U \wedge \neg(e \in S \vee e \in T) \\ & = e \in U \wedge (e \notin S \wedge e \notin T) \\ & = (e \in U \wedge e \notin S) \wedge (e \in U \wedge e \notin T). \end{aligned}$$

Thus, $e \in S^c \cap T^c$, and so $(S \cup T)^c \subset S^c \cap T^c$.

Conversely, let $e \in S^c \cap T^c$. Then

$$\begin{aligned} & (e \in U \wedge e \notin S) \wedge (e \in U \wedge e \notin T) \\ & = e \in U \wedge (e \notin S \wedge e \notin T) \\ & = e \in U \wedge \neg(e \in S \vee e \in T). \end{aligned}$$

Thus, $e \in (S \cup T)^c$, and so $S^c \cap T^c \subset (S \cup T)^c$.

Since both subsets are equal, $(S \cup T)^c = S^c \cap T^c$.

Now, let $e \in (S \cap T)^c$. Then

$$\begin{aligned} & e \in U \wedge \neg(e \in S \wedge e \in T) \\ & e \in U \wedge (e \notin S \vee e \notin T) \\ & (e \in U \wedge e \notin S) \vee (e \in U \wedge e \notin T). \end{aligned}$$

Thus, $e \in S^c \cup T^c$, and so $(S \cap T)^c \subset S^c \cup T^c$.

Conversely, let $e \in S^c \cup T^c$. Then

$$\begin{aligned} & (e \in U \wedge e \notin S) \vee (e \in U \wedge e \notin T) \\ & e \in U \wedge (e \notin S \vee e \notin T) \\ & e \in U \wedge \neg(e \in S \wedge e \in T). \end{aligned}$$

Thus, $e \in (S \cap T)^c$, and so $S^c \cup T^c \subset (S \cap T)^c$.

Since both are subsets of each other, $(S \cap T)^c = S^c \cup T^c$. \square

(d) *Proof.* Let $e \in (S \cup T) \setminus (S \cap T)$. Then

$$\begin{aligned} & (e \in S \vee e \in T) \wedge \neg(e \in S \wedge e \in T) \\ & (e \in S \vee e \in T) \wedge (e \notin S \vee e \notin T) \\ & ((e \in S \vee e \in T) \wedge e \notin S) \vee ((e \in S \vee e \in T) \wedge e \notin T) \\ & ((e \in S \wedge e \notin S) \vee (e \in T \wedge e \notin S)) \vee ((e \in S \wedge e \notin T) \vee (e \in T \wedge e \notin T)) \\ & (e \in T \wedge e \notin S) \vee (e \in S \wedge e \notin T). \end{aligned}$$

Thus, $e \in (S \setminus T) \cup (T \setminus S)$, and so $(S \cup T) \setminus (S \cap T) \subset (S \setminus T) \cup (T \setminus S)$.

Conversely, let $e \in (S \setminus T) \cup (T \setminus S)$. Then

$$\begin{aligned} & (e \in T \wedge e \notin S) \vee (e \in S \wedge e \notin T) \\ & (e \in S \wedge e \notin T) \vee (e \in T \wedge e \notin S) \\ & ((e \in S \wedge e \notin S) \vee (e \in T \wedge e \notin S)) \vee ((e \in S \wedge e \notin T) \vee (e \in T \wedge e \notin T)) \\ & ((e \in S \vee e \in T) \wedge e \notin S) \vee ((e \in S \vee e \in T) \wedge e \notin T) \\ & (e \in S \vee e \in T) \wedge (e \notin S \vee e \notin T) \\ & (e \in S \vee e \in T) \wedge \neg(e \in S \wedge e \in T). \end{aligned}$$

Thus, we get the statement $e \in (S \cup T) \setminus (S \cap T)$, and so $(S \setminus T) \cup (T \setminus S) \subset (S \cup T) \setminus (S \cap T)$. Since both are subsets of each other, $(S \cup T) \setminus (S \cap T) = (S \setminus T) \cup (T \setminus S)$. \square

From these problems, we observe that sets and logical statements are quite similar. A set is analogous to a logical statement, and the operators union and intersection resemble the logical “or” and “and” respectively (specifically, given sets S, T , $e \in S \cup T$ is equivalent to $e \in S \vee e \in T$, and $e \in S \cap T$ is equivalent to $e \in S \wedge e \in T$). Given a well defined complement of S , the complement S^c is analogous to the logical “not” (just as only one of P and $\neg P$ may be true, only one of $e \in S^c$ and $e \in S$ may be true). The symmetric difference is analogous to the “xor” operator in the sense that an element e being in $S \Delta T$ meaning e is in S or T , but not both, is similar in structure to notion that $P \vee Q$ means that in order to be true, either P or Q could be true, but not both.

Problem §3 Complete Exercise 1.16:

- Let S, T be finite sets with $|S| = |T|$, and let $f : S \rightarrow T$ be a function from S to T . Prove the following are equivalent:
 - f is injective.
 - f is surjective.
 - f is bijective.

and Exercise 1.17:

- Give an example of a function $f : \mathbb{N} \rightarrow \mathbb{N}$ that is injective, but not surjective.
- Give an example of a function $f : \mathbb{N} \rightarrow \mathbb{N}$ that is surjective, but not injective.

Solution:

- (1.16)

Proof. Let $n = |S| = |T|$. We start by showing f injective implies f surjective. Let f be an injective function, and suppose that f is not surjective. Then $\exists t \in T$ such that $\forall s \in S, f(s) \neq t$; and so $|\text{im } S| < n$. By the definition of a function, every $s \in S$ is mapped to some element $f(s) \in T$; and since $|S| = n$ and $|\text{im } S| < n$, at least one $e \in \text{im } S$ is mapped to by at least two distinct elements $s, s' \in S$ (analogously, imagine each $e \in \text{im } S$ represents a “hole”, and each $s \in S$ a pigeon; since there are at most $n - 1$ holes, and n pigeons, by the PHP, at least one hole must have at least two distinct pigeons). But this implies that $e = f(s) = f(s'), s \neq s'$, a contradiction to injectivity. Thus, if f is injective, then f must be surjective as well.

Now, we show that f surjective implies f injective. Let f be a surjective function, and suppose that f is not injective. Then $\exists s, s' \in S$ such that $f(s) = f(s'), s \neq s'$. By definition of a function, each $s \in S$ is mapped to one and only one $f(s) \in \text{im } S$. But since f is not injective, at least one $f(s) \in \text{im } S$ is mapped to by at least two distinct $s, s' \in S$ (i.e. $\exists s, s' \in S, \exists f(s), f(s') \in \text{im } S, f(s) = f(s'), s \neq s'$), which implies that $|\text{im } S| < n$ (equivalently, at least one $t \in T$ is not mapped to by any $s \in S$), a contradiction to surjectivity. Thus, if f is surjective, then f must be injective as well.

Since f injective implies f surjective, and f surjective implies f injective, if f is either injective or surjective, it is bijective as well; and trivially, f bijective implies both injective and surjective. Thus the three statements are equivalent. \square

- (1.17)

– Let

$$f : \mathbb{N} \longrightarrow \mathbb{N}$$

$$n \longmapsto f(n) = n + 1.$$

f is injective, as no two $n_1, n_2 \in \mathbb{N}$ share a $\text{succ}(n)$ unless $n_1 = n_2$ (equivalently, $n_1 + 1 = n_2 + 1$ implies $n_1 = n_2$). f is also not surjective, as $1 \notin \text{im } f$.

– Let

$$\begin{aligned} f : \mathbb{N} &\longrightarrow \mathbb{N} \\ n &\longmapsto f(n) = \left\lceil \frac{n}{2} \right\rceil. \end{aligned}$$

f is surjective, as for any $k \in \mathbb{N}$, take $n = 2k \in \mathbb{N}$; then we get $f(n) = \left\lceil \frac{2k}{2} \right\rceil = k$. On the other hand, f is not injective. Let $n_1, n_2 \in \mathbb{N}$, $n_1 = 1, n_2 = 2$. Then $f(n_1) = f(n_2) = 1$, but $n_1 \neq n_2$.

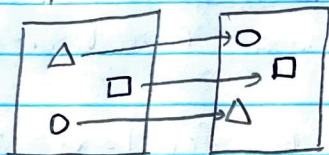
2.1 Define permutations σ_1, σ_2 on $\{\Delta, \square, \circ\}$

by:

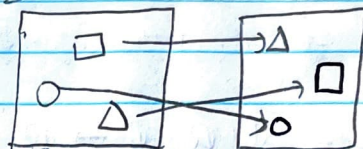
$$\sigma_1(\Delta) = \circ \quad \sigma_1(\square) = \square \quad \sigma_1(\circ) = \Delta$$

$$\sigma_2(\Delta) = \square \quad \sigma_2(\square) = \Delta \quad \sigma_2(\circ) = \circ$$

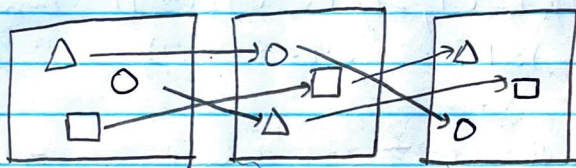
a) σ_1 :



σ_2 :



b)

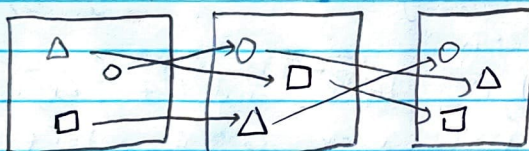


$$\sigma_2 \circ \sigma_1$$

First do σ_1 . Then σ_2 .

$$\sigma_2 \circ \sigma_1(\Delta) = \circ \quad \sigma_2 \circ \sigma_1(\square) = \Delta \quad \sigma_2 \circ \sigma_1(\circ) = \square$$

c)

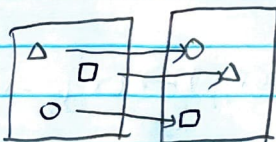


$$\sigma_1 \circ \sigma_2$$

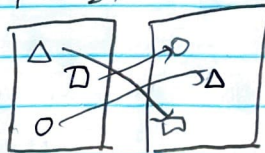
First do σ_2 . Then σ_1 .

$$\sigma_1 \circ \sigma_2(\Delta) = \square \quad \sigma_1 \circ \sigma_2(\square) = \circ \quad \sigma_1 \circ \sigma_2(\circ) = \Delta$$

d) $\sigma_2 \circ \sigma_1$:



$\sigma_1 \circ \sigma_2$:



No, $\sigma_2 \circ \sigma_1 \neq \sigma_1 \circ \sigma_2$.

2.4) Fill in the blanks,

