GFSR - Lewis, Payne 1973

Uses the same LFSR across all bits of the words : if the degree of the LFSR is n then the period is 2^n-1 .

The initialization is problematic and slow especially in the original rng.

If we see the state as made by n vertical words then fill each row of bits one after the other with the LFSR skipping some outcomes after any

$$[\mathbf{w_0},\mathbf{w_1},\ldots] = egin{bmatrix} w_{0,0} & \ldots & w_{0,n-1} \ w_{1,0} & \ldots & w_{1,n-1} \ w_{2,0} & \ldots & w_{2,n-1} \end{bmatrix}$$

```
In [2]: p:2
Out[2]: 2
In [3]: n:3
Out[3]: 3
In [4]: gf_primitive_poly(p,3)
Out[4]: x^3 + x + 1
```

Therefore a $LFSR(3,1+x+x^3)$ will have maximal period $p^n - 1 = 2^3 - 1 = 8 - 1 = 7$

```
In [5]: seed:[1,0,1]
Out[5]: [1,0,1]
In [6]: lfsr:matrix([1,1,0],
                      [1,0,0])
In [7]: for i:1 thru p^n do ( seed:mod(seed . lfsr,p), print(i,seed))
         2 [ 0 0 1 ]
3 [ 1 0 0 ]
         7 [ 1 0 1 ]
8 [ 0 1 0 ]
Out[7]: done
```

Let's now build a GFSR with the same polynomial using numbers $0 \leq x \leq m:2^3$

```
In [17]: m:2<sup>3</sup>
Out[17]: 8
```

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In [20]: mseed:matrix([7,3,2])
Out[20]: (7 3 2)
```

In [21]: for i:1 thru p^m-1 do (mseed:mod(mseed . lfsr,m), print(i,mseed))

Out[21]:	done
In []:	