# Prime Fields : $\mathbb{Z}_p$

```
In [2]: p:7
Out[2]: 7
In [3]: n:1
Out[3]: 1
In [4]: gf_set_data(p,n)
Out[4]: Structure [GF-DATA]
In [5]: gf_info()
         characteristic = 7
         reduction polynomial = x
         primitive element = 3
        nr of elements = 7
nr of units = 6
         nr of primitive elements = 2
Out[5]: false
In [6]: gf_add_table()
Out[6]:
          /0 1 2 3 4 5 6
           2\quad 3\quad 4\quad 5\quad 6\quad 0\quad 1
          3 4 5 6 0 1 2
4 5 6 0 1 2 3
In [7]: gf_mult_table()
Out[7]: /1 2 3 4 5 6
```

#### If an element $\alpha$ generates with its powers all the elements of the group $\mathbb{Z}_n^*$

then lpha is a generator of the cyclic multiplicative group and is called **primitive**.

In  $\mathbb{Z}_7^*$  , 3 and 5 are the **primitive** elements.

```
In [27]: for i:1 thru p-1 do print(i,gf_order(i))

1 1
2 3
3 6
4 3
5 6
6 2
Out[27]: done
```

1 of 3 4/10/18, 1:25 PM

#### Fermat's little theorem

```
p \mid (a^p - a)
```

or, said in another way

$$a^p - a \equiv 0 \mod p$$

```
In [29]: for a:1 thru p-1 do print(a,mod(a^p-a,p))
         1 0
         2 0
         5 0
         6 0
Out[29]: done
```

#### if a eq 0 then we can multiply by $a^{-1}$

```
a^{p-1}-1\equiv 0\mod p \quad \text{or} \quad a^{p-1}\equiv 1\mod p
```

```
In [30]: for a:1 thru p-1 do print(a,mod(a^(p-1),p))
         1 1
         3 1
         4 1
         5 1
         6 1
```

#### Out[30]: **done**

#### Inverse computed with Fermat's little theorem

```
let's multiply again by a^{-1}
a^{p-2} \equiv a^{-1} \mod p
  In [11]: for a:2 thru p-1 do print(a, "^-1=", mod(a^(p-2),p))
           2 ^-1= 4
           3 ^-1= 5
           4 ^-1= 2
          5 ^-1= 3
           6 ^-1= 6
  Out[11]: done
  In [12]: gf_primitive()
  Out[12]: 3
  In [16]: gf_make_logs()
  In [24]: for j:0 thru p-2 do print(gf_primitive(),"^",j,"=",gf_powers[j])
          3 ^ 0 = 1
3 ^ 1 = 3
3 ^ 2 = 2
           3 ^ 3 = 6
  Out[24]: done
```

2 of 3 4/10/18, 1:25 PM

In	]:	
	 • •	

3 of 3 4/10/18, 1:25 PM

# Prime Power Fields : $\mathbb{F}_q=\mathbb{Z}_p[x]/m(x)=GF(p^n)$

```
In [ ]: p:3
In [ ]: n:2
In [3]: gf_set_data(p,n)
Out[3]: Structure [GF-DATA]
In [4]: gf_info()
         characteristic = 3
         reduction polynomial = x^2+1
         primitive element = x+1
         \frac{1}{1} nr of elements = 9
        nr of units = 8
        nr of primitive elements = 4
Out[4]: false
In [18]: atable:gf_add_table()
          0 1 2 3 4 5
                              6 7
                                    8
           1 2 0 4 5 3 7 8 6
In [19]: for i:1 thru p^n do for j:1 thru p^n do atable[i,j]:gf_n2p(atable[i,j])
Out[19]: done
```

In [20]:	print(atal	ble)\$							
		[ 0	]	[ 1		[ 2 ]			
		[ 1	]	2		[ 0 ]			
		[ [ 2	] ]	0	 				
		[ x	ו ] ]	x + 1	 	[ x + 2 ]			
	Col 1 =	[ x + 1	] Col 2 =	x + 2	Col 3 =				
	İ	x + 2	]	x		x + 1 ]			
	i	2 x	] ]	2 x + 1	i i	2 x + 2 ]			
	İ	[ 2 x + 1	]	2 x + 2		2 x ]			
		2 x + 2 x	] ]	2 x [ x + 1 ]	i i	[ 2 x + 1 ] [ x + 2 ]			
		[ [ x + 1	]	[ x + 2	] 	[ x ]			
		[ x + 2	] ]	X		[			
		l [ 2 x	] ]	2 x + 1	 	[ 2 x + 2 ]			
	Col 4 =	[ 2 x + 1	] ] Col 5 =   1	2 x + 2	   Col 6 =	[ 2 x ]			
		[ 2 x + 2	, ] ]	2 x	   	[ 2 x + 1 ]			
		0	]	1					
	į	1	] ]	2	i i	0 ]			
	İ	2 [ 2 x	]	0 [ [ 2 x + 1 ]	j   	1 ] [2x+2]			
		[ [ 2 x + 1	]	2 x + 2		[ 2 x ]			
		l [ 2 x + 2	]	2 x		[ 2 x + 1 ]			
		[ 0	]	1		[ 2 ]			
	Col 7 =	[ 1	] ] Col 8 =   1	2	Col 9 =	[ 0 ]			
		[ 2 [	, ] ]	0					
		[ X	]	x + 1		[ x + 2 ]			
	į	[ x + 1 ]	]	x + 2		x ]			
		[ x + 2	]	[ x ]	l 1	[ x + 1 ]			
Out[20]:	$\begin{pmatrix} 0 \\ 1 \end{pmatrix}$	1	2	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	+1  x +	-2 $2x$	$2x+1\\2x+2$	1	
	2	0	1 :	x+1 $x+2$	$egin{array}{cccccccccccccccccccccccccccccccccccc$	2x+1 $2x+2$	2x + 2 $2x$	$\begin{bmatrix} 2 & x \\ 2 & x + 1 \end{bmatrix}$	
	x	x + 1	x + 2	2x $2x$	x + 1  2x -	+2 0	1	2	
	x+1	x+2	x 2	x+1 2 $x$	z+2 2:	$egin{array}{cccccccccccccccccccccccccccccccccccc$	2	0	
	$\begin{vmatrix} x+2\\2x \end{vmatrix}$	x $2x + 1$	x+1 2 $2x+2$	x+2 2	$egin{array}{cccc} 2x & 2x & -1 & 2 \end{array}$	$+ 1 \qquad 2 \qquad \qquad x$	$0 \ x+1$	$egin{array}{c} 1 \ x+2 \end{array}$	
	2x+1	2x+2	$\frac{2x+2}{2x}$	1	2 0	x + 1	$x+1 \\ x+2$	$\begin{bmatrix} x + 2 \\ x \end{bmatrix}$	
	$\sqrt{2x+2}$	2  x	2x+1		0 1	x+2	x	x+1	

```
In [21]: | mtable:gf_mult_table()$
Out[21]:
                        7
                    8
                           3
                               5
                                  4
                                  7
                        1
                               2
                                  3
                                  2
                    1
                               8
                                  5
                        6
                           8
                               3
               5
                     ^{2}
                                  1
In [23]: for i:1 thru p^n-1 do for j:1 thru p^n-1 do mtable[i,j]:gf_n2p(mtable[i,j])
Out[23]: done
In [24]: print(mtable)
                                       x + 1
                                              x + 2
                                                          2 x
                                                                 2 \times + 1 \quad 2 \times + 2 
                       1
                               2 x
                                      2 \times + 2 \quad 2 \times + 1
                                2
                                       x + 2
                                               2 x + 2
                                                           1
                                        2 x
                                                  1
                                                                    2
                    2 x + 2
                              x + 2
                                                        2 x + 1
                                                                             Х
                    2 \times + 1 \quad 2 \times + 2
                                         1
                                                  Х
                                                         x + 1
                                                                   2 x
                                                                             2
                                      2 x + 1
                                                x + 1
                                                           2
                     x + 2
                                         2
                                                 2 x
                                                        2 x + 2
                                                                             1
         [2x + 2]
                     x + 1
                             2 x + 1
                                                  2
                                                         x + 2
                                                                    1
Out[24]:
                       2
                                                         2x
                                      x + 1
                                               x + 2
                                                               2x + 1
                                                                        2x + 2
                               \boldsymbol{x}
                       1
                               2x
                                     2x+2 2x+1
                                                                         x + 1
                                                         \boldsymbol{x}
                                      x + 2
                                              2x + 2
                                                                x + 1
                                                                        2x + 1
                                                                  2
                             x + 2
                                       2x
                                                 1
                                                      2x + 1
                                                                           \boldsymbol{x}
                                                                           2
                                        1
                                                       x + 1
                                                                 2x
                                                 \boldsymbol{x}
                                                         2
                               1
                                     2x + 1
                                              x+1
                                                               2x + 2
                                                                         x + 2
                       \boldsymbol{x}
                                                2x
                                        2
                                                                          1
           2x + 1
                    x + 2
                             x + 1
                                                       2x + 2
                                                                  \boldsymbol{x}
           2x + 2
                    x + 1
                            2x + 1
                                                 2
                                                       x + 2
                                                                  1
                                                                          2x
                                        \boldsymbol{x}
In [25]: gf_make_logs()
In [31]: | for i:1 thru p^n-1 do print("(",gf primitive(),")^",i,"=",gf n2p(gf powers[i]))
         (x + 1)^1 = x + 1
         (x + 1)^2 = 2x
         (x + 1)^3 = 2x + 1
         (x + 1)^4 = 2
         (x + 1)^5 = 2x + 2
         (x + 1)^6 = x
         (x + 1)^7 = x + 2
         (x + 1)^{8} = 1
Out[31]: done
```

```
In [32]: for i:1 thru p^n-1 do print(gf_n2p(i)," has order ",gf_order(gf_n2p(i)))

1 has order 1
2 has order 2
x has order 4
x + 1 has order 8
2 x has order 8
2 x has order 4
2 x + 1 has order 8
2 x + 2 has order 8
0 x + 2 has order 8
0 x + 2 has order 8
1 x + 2 has order 8
2 x + 1 has order 8
2 x + 2 has order 8
0 x + 2 has order 8
0 x + 2 has order 8
```

## Universal Polynomial : $x^{p^n} - x$

factors in all irreducible monic polynomials of degree  $k\mid n$  In this case n=2 and therefore it factors in all irreducibles of degree 1 or 2

```
In [64]: gf_{-}factor(x^p^n-x)
Out[64]: x(x+1)(x+2)(x^2+1)(x^2+x+2)(x^2+2x+2)

In [64]: gf_{-}factor(x^2+1)
Out[64]: x^2+1

In [64]: for i:0 thru p^n-1 do (print(gf_{-}add(x,-gf_{-}n2p(i))))
\begin{array}{c} x \\ x+2 \\ x+1 \\ 0 \\ 2 \\ x+1 \\ 0 \end{array}
\begin{array}{c} 2 \\ 2 \\ 2 \\ x+1 \\ 0 \end{array}
Out[64]: done

In []:
```

# **Primitive polynomials**

 $primitive\ polynomials \subset irreducible\ polynomials \subset polynomials$ 

Is  $f(x)=x^3+x+1$  a primitive polynomial over  $\mathbb{Z}_2$  ?

```
In [1]: f:x^3+x+1;

Out[1]: x^3+x+1

In [2]: n:hipow(f,x);

Out[2]: 3

In [3]: p:modulus:2;

Out[3]: 2
```

To be primitive f(x) should divide  $x^{p^n-1}-1$  and no other  $x^e-1$  for  $e< p^n-1$ 

```
In [4]: \operatorname{divide}(x^{(p^n-1)-1,f)};

Out[4]: [x^4 + x^2 + x + 1, 0]

In [5]: \operatorname{for e:1 thru p^n-1 do print("e=",e,",",divide(x^e-1,f))};

e=1, [0, x+1]
2
e=2, [0, x+1]
e=3, [1, x]
2
e=4, [0, x+1]
2
e=5, [0, x+1]
2
2
2
3
4
4
5
4
5
4
7
8
8
9

Out[5]: \operatorname{done}

In [6]: \operatorname{gf\_primitive\_poly\_p(f,p)};

Out[6]: \operatorname{true}

In [7]: \operatorname{factor}(x^{(p^n-1)-1});

Out[7]: (x+1)(x^3+x+1)(x^3+x^2+1)
```

1 of 3 4/15/18, 10:31 AM

## Universal polynomial : $U(x) = x^{p^n} - x$

The universal polynomial is the product of all irreducible polynomials of degree d for  $\forall d:d\mid n$ 

```
In [8]: factor(x^p^n - x); Out[8]: x(x+1)(x^3+x+1)(x^3+x^2+1)
```

# If f(x) is a primitive polynomial of degree n over $\mathbb{Z}_p$ then x is a generator of

$$\mathbb{F}_{p^n}=\mathbb{Z}_p[x]/f(x)$$

# The above should give you a hint of why the LFSR works. Multiplying by $\boldsymbol{x}$ is shifting left.

If x is a generator then the order of x should be  $p^n-1$ 

2 of 3 4/15/18, 10:31 AM

# If we know the factorization of $p^n-1$ then we can check that

 $orall primes \ q|(p^n-1),f 
mid x^{rac{p^n-1}{q}}-1$ 

```
In [12]: modulus:3;
Out[12]: 3
In [13]: p:3;
Out[13]: 3
In [14]: n:4;
Out[14]: 4
In [15]: f:gf_primitive_poly(p,n);
Out[15]: x<sup>4</sup> + x + 2
In [16]: ifactors(p^n-1);
Out[16]: [[2,4],[5,1]]
```

#### Prime factors of 80 are 2,5

3 of 3 4/15/18, 10:31 AM

#### LFSR Linear Feedback Shift Register

An LFSR has maximal period if its associated/connection polynomial is primitive.

In that case if n is the length of the LFSR, the period will be  $2^n-1$ 

```
In [1]: p:2
Out[1]: 2
In [2]: n:4
Out[2]: 4
```

## A poly of degree m over $\mathbb{Z}_p$ is primitive if its order is $p^m-1$

```
Here p=2, n=4, therefore p^m-1=2^4-1=15
```

```
In [3]: f:gf_primitive_poly(p,n)

Out[3]: x^4 + x + 1

In [4]: modulus:2

Out[4]: 2

In [5]: gf_set_data(p,n)

Out[5]: Structure [GF-DATA]

In [6]: gf_order(f(x))

Out[6]: 15
```

# f(x) is a primitive polynomial so we can expect a period of $2^4-1=15$ from its LFSR

```
In [7]: modulus:2
Out[7]: 2
In [8]: seed:[0,1,0,1]
Out[8]: [0,1,0,1]
```

# This matrix does exactly what a LFSR does: shifts right and replaces first bit with the xor of the taps

1 of 3 4/10/18, 1:09 PM

```
In [10]: expand(charpoly(mlfsr,lambda))
Out[10]: \lambda^4 - \lambda^3 - 1
In [11]: gf_primitive_poly_p(%,p)
Out[11]: true
In [14]: for i:1 thru p^n-1 do ( seed:seed . mlfsr, seed:mod(seed,2),print(i,seed) )
         2 [ 0 3 [ 1
                 0
                       1]
                    1
                 0 0 1]
          4 [ 0
                1 0 0 ]
         5 [ 0 0 1 0 ]
6 [ 0 0 0 1 ]
7 [ 1 0 0 0 ]
8 [ 1 1 0 0 ]
          9[1110]
          10 [ 1 1 1 1 ]
          11 [ 0 1
         12 [ 1 0 1 1 ]
13 [ 0 1 0 1 ]
          14 [ 1 0 1 0 ]
          15 [ 1 1 0 1 ]
Out[14]: done
In [64]: p:3
Out[64]: 3
In [64]: n:3
Out[64]: 3
In [64]: gf_primitive_poly(p,n)
Out[64]: x^3 + 2x + 1
In [64]: seed:[0,2,1]
\texttt{Out[64]:}\ [0,2,1]
Out[64]: /2 \ 1 \ 0
            0 \ 0 \ 1
In [64]: expand(charpoly(mlfsr,lambda))
Out[64]: -\lambda^3+2\,\lambda^2+1
In [64]: gf_primitive_poly_p(%,p)
Out[64]: false
```

2 of 3 4/10/18, 1:09 PM

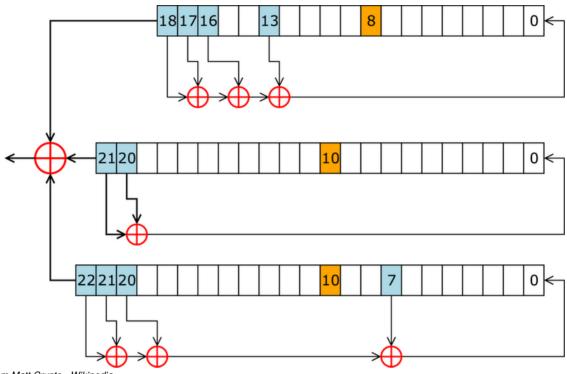
3 of 3 4/10/18, 1:09 PM

### GSM A5/1 cipher

#### 3 different LFSR are initialized with 64 bits

and their output is xored to produce 114 bits at a time

these bits are xored with 114 bits of the data stream and sent every 4.615 msec



pic from Matt Crypto - Wikipedia

```
In [1]: lfsr1:x^19+x^18+x^17+x^14+1
```

Out[1]: 
$$x^{19} + x^{18} + x^{17} + x^{14} + 1$$

Out[2]: **true** 

Out[3]: 
$$x^{22} + x^{21} + 1$$

Out[4]:  $\mathbf{true}$ 

Out[5]: 
$$x^{23} + x^{22} + x^{21} + x^8 + 1$$

Out[6]: **true** 

1 of 1 4/10/18, 2:27 PM

#### GFSR - Lewis, Payne 1973

Uses the same LFSR across all bits of the words : if the degree of the LFSR is n then the period is  $2^n-1$ .

The initialization is problematic and slow especially in the original rng.

If we see the state as made by n vertical words then fill each row of bits one after the other with the LFSR skipping some outcomes after any

$$[\mathbf{w_0},\mathbf{w_1},\ldots] = egin{bmatrix} w_{0,0} & \ldots & w_{0,n-1} \ w_{1,0} & \ldots & w_{1,n-1} \ w_{2,0} & \ldots & w_{2,n-1} \end{bmatrix}$$

```
In [2]: p:2
Out[2]: 2
In [3]: n:3
Out[3]: 3
In [4]: gf_primitive_poly(p,3)
Out[4]: x^3 + x + 1
```

# Therefore a $LFSR(3,1+x+x^3)$ will have maximal period $p^n - 1 = 2^3 - 1 = 8 - 1 = 7$

```
In [5]: seed:[1,0,1]
Out[5]: [1,0,1]
In [6]: lfsr:matrix([1,1,0],
                      [1,0,0])
In [7]: for i:1 thru p^n do ( seed:mod(seed . lfsr,p), print(i,seed))
         2 [ 0 0 1 ]
3 [ 1 0 0 ]
         7 [ 1 0 1 ]
8 [ 0 1 0 ]
Out[7]: done
```

# Let's now build a GFSR with the same polynomial using numbers $0 \leq x \leq m:2^3$

```
In [17]: m:2<sup>3</sup>
Out[17]: 8
```

```
In [20]: mseed:matrix([7,3,2])
Out[20]: (7 3 2)
```

In [21]: for i:1 thru p^m-1 do ( mseed:mod(mseed . lfsr,m), print(i,mseed))

Out[21]:	: done	
In [ ]:	:	

#### **T400 Twisted GFSR**

based on  $GFSR(25, x^{25} + x^{11} + 1)$ 

n=25 words x w=16 bits = 400 bits, twisting vector  $\mathbf{a} = \mathbf{0xA875}$  (a 16 bits vector)

twisting matrix 
$$A=\left[egin{array}{cc} 0_{15x1} & I_{15x15} \\ \mathbf{a_{1x16}} \end{array}
ight]$$
 (a 16x16 bits array)

if we let  $x = [x_0 \dots x_{w-2} x_{w-1}]$  then the block multiplication is

$$egin{align*} \left[x_0 \ldots x_{w-2} \mid x_{w-1}
ight] & \cdot \left[egin{array}{cccc} 0_{15x1} & I_{15x15} \ a_0 & a_1 \ldots a_{15} \end{array}
ight] = \left[egin{array}{cccc} x_{w-1} . \, a_0 & x_0 + x_{w-1} * a_1 \ldots x_{w-2} + x_{w-1} * a_{16} \end{array}
ight] \ & = \left[x_{w-1} . \, \mathbf{a}_{1x16} \oplus shiftright(\mathbf{x})
ight] \end{aligned}$$

The form of  $\boldsymbol{A}$  is dictated by the necessity to make it simple to multiply by it :

$$\mathbf{x}.\mathbf{A} = if(x_{w-1} = 0) then shiftright(\mathbf{x}) else shiftright(\mathbf{x}) \oplus \mathbf{a}$$

(Matsumoto, Kurita, 1992) Theorem : if  $arphi_A(x)$  is the characteristic polynomial of the w imes w bits matrix A and  $arphi_A(t^n+t^m)$  is of degree nw and is primitive then the period of :

$$x_{l+n} = x_{l+m} \oplus x_l$$
.  $A$ 

is  $2^{nw}-1$ .

This generator returns the random floats  $\frac{x}{2^{16}}$ 

In [26]: a\_15\_ident:diagmatrix(15,1)\$

In [27]: a\_15\_zero\_row:[0,0,0,0,0,0,0,0,0,0,0,0,0,0,0]

Out[27]: [0,0,0,0,0,0,0,0,0,0,0,0,0,0]

1 of 4 4/10/18, 2:07 PM

```
In [28]: a_15_zero_col:transpose(a_15_zero_row)
Out[28]:
          0
          0
          0
          0
          0
          0
          0
          0
          0
In [29]: a_16_zero_row:addcol(matrix(a_15_zero_row),matrix([0]))
In [ ]:
In [30]: a_16_vector_a:[1,0,1,0,1,0,0,0,0,1,1,1,0,1,0,1]
Out[30]: [1,0,1,0,1,0,0,0,1,1,1,0,1,0,1]
In [31]: a_15x16:addcol(a_15_zero_col,a_15_ident)$
Out[31]:
                0 \quad 0
                                             0
                                                 0
                                                   0 0
                                                      0
                                                       0
                                     0 \quad 0
                                     0 \ 0 \ 0 \ 1
                                                   0 \quad 0
                                          0 \ 0 \ 1 \ 0 \ 0
                                    0 \quad 0
             0
                                     0 \quad 0
                                          0 \ 0 \ 0 \ 1
                                                      0
```

2 of 4 4/10/18, 2:07 PM

```
In [32]: | shiftright_16x16:addrow(a_15x16,a_16_zero_row)
In [33]: a 16x16:addrow(a 15x16,a 16 vector a)
Out[33]:
In [34]: phi:charpoly(a_16x16,t^25+t^11)
                                       Out[34]:
In [35]: phi:expand(phi)
\texttt{Out[35]:} \quad t^{400} + 16\,t^{386} - t^{375} + 120\,t^{372} - 15\,t^{361} + 560\,t^{358} - 105\,t^{347} + 1820\,t^{344} - 455\,t^{333} + 4368\,t^{330} - t^{325}
                                                -\,13\,t^{311} - 3003\,t^{305} + 11440\,t^{302} - 78\,t^{297} - 5005\,t^{291} + 12870\,t^{288} - 286\,t^{283} - 6435\,t^{277} - t^{275} +
                                          -\,6435\,t^{263}-11\,t^{261}+8008\,t^{260}-1287\,t^{255}-t^{250}-5005\,t^{249}-55\,t^{247}+4368\,t^{246}-1716\,t^{241}-10
                                          t^{233} + 1820\,t^{232} - 1716\,t^{227} - t^{225} - 45\,t^{222} - 1365\,t^{221} - 330\,t^{219} + 560\,t^{218} - 1287\,t^{213} - 9\,t^{211} - 12
                                      t^{205} + 120\,t^{204} - 715\,t^{199} - 36\,t^{197} - 210\,t^{194} - 105\,t^{193} - 462\,t^{191} + 16\,t^{190} - 286\,t^{185} - 84\,t^{183} - 252\,t^{191} + 16\,t^{190} - 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{190} + 100\,t^{1
```

3 of 4 4/10/18, 2:07 PM

 $+\,t^{176} - 78\,t^{171} - 126\,t^{169} - 210\,t^{166} - t^{165} - 165\,t^{163} - 13\,t^{157} - 126\,t^{155} - 120\,t^{152} - 55\,t^{149} - t^{143} - t^{143} + t^{144} + t^{$ 

 $t^{135} - 36t^{127} - 10t^{124} - t^{121} - 9t^{113} - t^{110} - t^{100} - t^{99} - 4t^{86} - 6t^{72} - 4t^{58} - t^{50} - t^{44} - t^{120} - t^$ 

4 of 4 4/10/18, 2:07 PM

0 0

#### T800 TGFSR A twisted GFSR

#### 25 words x 32 bits = 800 bits

0 0

based on  $LGFSR(25, x^{25} + x^{18} + 1)$  twisting vector  $\mathbf{a} = \mathbf{0x8ebfd028}$ 

1 of 4 4/10/18, 2:14 PM

 $0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0$ 

 $0 \quad 0 \quad 0 \quad 0$ 

```
Out[16]:
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
       0
In [20]: arow32:[1,0,0,0,1,1,1,0,1,0,1,1,1,1,1,1,1,1,0,0,0,0,0,0,0,1,0,1,0,0,0]$
```

2 of 4 4/10/18, 2:14 PM

 $0 \\ \texttt{ut[20]:} \ \overline{[1,0,0,0,1,1,1,0,1,0,1,1,1,1,1,1,1,1,1,0,1,0,0,0,0,0,0,1,0,1,0,0,0] }$ 

In [18]: | a:addcol(acol1,a)\$

Out[18]:

3 of 4 4/10/18, 2:14 PM

In [ ]: d:charpoly(a,t^25+t^18)
In [ ]:

```
In [1]: ident 4x4:ident(4);
Out[1]: /1 \ 0 \ 0 \ 0
In [2]: zero_4x1:transpose([0,0,0,0]);
Out[2]: /0
In [3]: vec_a:[1,1,0,1,1];
Out[3]: [1,1,0,1,1]
In [4]: m:ident_4x4;
Out[4]: /1 0 0 0
In [5]: m:addcol(zero_4x1,m);
 Out[5]: (0 \ 1 \ 0 \ 0)
In [6]: m:addrow(m,vec_a);
Out[6]: (0 \ 1 \ 0 \ 0)
In [9]: gf_primitive_poly(2,5);
Out[9]: x^5 + x^2 + 1
In [10]: m;
Out[10]: (0 \ 1 \ 0 \ 0)
```

1 of 2 4/12/18, 6:49 PM

2 of 2 4/12/18, 6:49 PM

#### **Tempering**

The new very long period Twisted GFSR can have not so good statistical properties. Matsumoto and Kurita developed a shuffling of the bits that can help with this.

For T800 they resubmitted a TT800 version (Tempered Twisted GFSR). The Marsenne Twister MT19337 is instead already equipped with a similar pattern.

```
T800 twisting:
   #include <stdio.h>
   #include <unistd.h>
   int main(int argc, char* argv[])
   unsigned int x,y,z;
   int s = 7, t=15;
   unsigned int a = 0x8ebfd028 ,b = 0x2b5b2500 ,c = 0xdb8b0000;
     while (1) {
       if (read(0,&x,4) != 4) _exit(1);
       y = x ^ ((x << s) & b);
       z = y ^ ((y << t) & c);
       write(1,&z,4);
     }
   }
MT19337 twisting:
   #include <stdio.h>
   #include <unistd.h>
   int main(int argc, char* argv[])
   unsigned int x,y,z;
   int u = 11, s = 7, t=15, l = 18;
   unsigned int a = 0x9908b0df ,b = 0x9d2c5680 ,c = 0xefc60000;
     while (1) {
       if (read(0,&x,4) != 4) _exit(1);
       y = x ^ (x>>u);
       y = y ^ ((y << s) & b);
       y = y ^ ((y << t) & c);
       z = y ^ (y>>l) ;
       write(1,&z,4);
     }
   }
    In [ ]:
```

1 of 1 4/10/18, 3:53 PM

#### **Design an MT**

```
In [1]: load("bitwise");
 {\tt Out[1]:} \quad /usr/local/share/maxima/5.41.0/share/contrib/bitwise/bitwise.lisp \\
 In [2]: seed:matrix([27,17,21,5,30,14,16]);
 {\tt Out[2]:}\ (27\ 17\ 21\ 5\ 30\ 14\ 16)
 In [4]: n:matrix_size(seed)[2];
 Out[4]: 7
 In [5]: ident_4x4:ident(4);
 Out[5]: /1 0 0 0
           0 \ 1 \ 0 \ 0
           0 \ 0 \ 1 \ 0
 In [6]: zero_4x1:transpose([0,0,0,0]);
 Out[6]:
            0
 In [7]: matrix A:ident 4x4;
 Out[7]: /1 \ 0 \ 0 \ 0
           0 \ 1 \ 0 \ 0
           0 \ 0 \ 1 \ 0
 In [8]: matrix_A:addcol(zero_4x1,matrix_A)$
 Out[8]: (0 \ 1 \ 0 \ 0)
            0 \ 0 \ 1 \ 0 \ 0
           0 \ 0 \ 0 \ 1 \ 0
           (0 \ 0 \ 0 \ 0 \ 1)
 In [9]: vec_a:[1,0,1,1,0]$
 Out[9]: [1,0,1,1,0]
In [10]: matrix_A:addrow(matrix_A,vec_a)$
Out[10]: (0 \ 1 \ 0 \ 0 \ 0)
            0 \ 0 \ 1 \ 0 \ 0
            0 \ 0 \ 0 \ 0 \ 1
In [11]: matrix_A[5]:vec_a;
Out[11]: [1,0,1,1,0]
In [12]: m:2;
Out[12]: 2
```

1 of 3 4/16/18, 10:09 AM

```
In [17]: bit_and(seed[1][m],16);
Out[17]: 16
In [18]: bit_rsh(%,4);
Out[18]: 1
In [19]: bottom_tap:[ bit_rsh(bit_and(seed[1][m],16),4),
                        bit_rsh(bit_and(seed[1][m],8),3),
                        bit_rsh(bit_and(seed[1][m],4),2),
bit_rsh(bit_and(seed[1][m],2),1),
                                 bit_and(seed[1][m],1)];
Out[19]: [1,0,0,0,1]
In [20]: top_tap:[
                       bit_rsh(bit_and(seed[1][n],16),4),
                       bit_rsh(bit_and(seed[1][n-1],8),3),
bit_rsh(bit_and(seed[1][n-1],4),2),
                       bit_rsh(bit_and(seed[1][n-1],2),1),
                               bit_and(seed[1][n-1],1)];
Out[20]: [1,1,1,1,0]
In [21]: bottom_tap:matrix(bottom_tap);
Out[21]: (1 0 0 0 1)
In [22]: new_bits:mod(bottom_tap + top_tap . matrix_A,2);
Out[22]: (1 1 1 1 0)
In [23]: new_numb:mod(new_bits[1][5]*16+
                        new_bits[1][4]*8+
                       new_bits[1][3]*4+
                       new_bits[1][2]*2+
                       new_bits[1][1],32);
Out[23]: 15
In [24]: shr7:matrix([0,1,0,0,0,0,0],
                        [0,0,1,0,0,0,0],
                        [0,0,0,1,0,0,0],
                        [0,0,0,0,1,0,0],
                        [0,0,0,0,0,1,0],
                        [0,0,0,0,0,0,1],
                       [0,0,0,0,0,0,0]);
Out[24]:
           /0 1 0 0 0 0 0
              0 \ 0 \ 1
                         0 \ 0 \ 0
               0 \ 0 \ 0 \ 1 \ 0 \ 0
            0
               0 \ 0 \ 0 \ 0 \ 1 \ 0
               0 0 0
                        0 \ 0 \ 1
              0 0
                     0
                            0
In [25]: seed:seed . shr7;
{\tt Out[25]:}\ (0\ 27\ 17\ 21\ 5\ 30\ 14)
In [26]: seed[1][1]:new_numb;
Out[26]: 15
```

2 of 3 4/16/18, 10:09 AM

3 of 3 4/16/18, 10:09 AM