

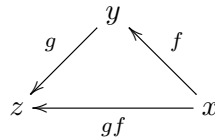
# TOPOLOGICAL QUANTUM FIELD THEORIES

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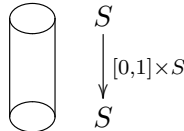
## CATEGORY THEORY

**Definition.** A category  $\mathcal{C}$  consists of

- a class of **objects**, denoted  $\text{Ob}(\mathcal{C})$ .
- given two objects  $x, y \in \mathcal{C}$ , a set  $\text{Hom}(x, y)$  of **morphisms**. Generalizing from the categories where  $\text{Hom}(x, y)$  is a set of functions, we denote  $f \in \text{Hom}(x, y)$  by  $f: x \rightarrow y$ . Morphisms satisfy the following properties:
  - given morphisms  $f: x \rightarrow y$  and  $g: y \rightarrow z$ , we can compose them and obtain  $g \circ f: x \rightarrow z$ . When there is no possibility of confusion  $g \circ f$  is abbreviated  $gf$ .



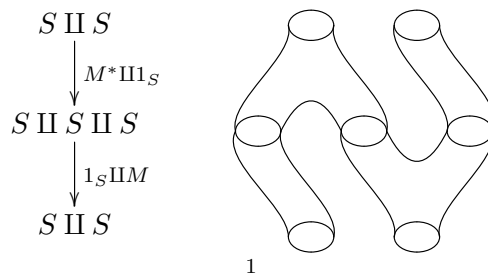
- for any  $x \in \mathcal{C}$ , there is an **identity** morphism  $1_x: x \rightarrow x$  such that, for any  $f: x \rightarrow y$ , we have  $f1_x = f = 1_y f$ . For example,



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Examples of categories are:

- Set, where objects are sets and morphisms are functions.
- $n\text{Cob}$ , where objects are  $(n - 1)$ -dimensional oriented compact manifolds, and morphisms are  $n$ -dimensional cobordisms.



- Vect, where objects are (finite-dimensional, complex) vector spaces, and morphisms are linear operators.
- Hilb, where objects are (finite-dimensional, complex) Hilbert spaces, and morphisms are linear operators.

Quantum mechanics uses Hilb rather than Vect because (among other things)

- given state vectors (i.e. unit vectors) in a Hilbert space, say  $\phi$  and  $\psi$ , then  $\langle \phi | \psi \rangle$  is the **amplitude** and  $|\langle \phi | \psi \rangle|^2$  is the **probability** that a system prepared in state  $\psi$  will be found in state  $\phi$ . There is no such structure in Vect.
- given an operator  $T: \mathcal{H} \rightarrow \mathcal{H}'$ , the condition  $\langle T^* \phi | \psi \rangle = \langle \phi | T \psi \rangle$  defines an **adjoint** operator  $T^*: \mathcal{H}' \rightarrow \mathcal{H}$ . In Vect, the best we can get is the dual  $T^*: \mathcal{H}'^* \rightarrow \mathcal{H}^*$ .
- **observables** in quantum mechanics are represented by self-adjoint operators  $A: \mathcal{H} \rightarrow \mathcal{H}$ , where  $\mathcal{H}$  is the space of states of the system and  $A = A^*$ . Such an operator<sup>1</sup> has associated an orthonormal basis  $\{\psi_i\}$  of  $\mathcal{H}$  such that  $A\psi_i = a_i\psi_i$  with  $a_i \in \mathbb{R}$ . The interpretation is that  $\psi_i$  is a state in which  $A$  will always be measured to be  $a_i$ .

The fact that in Hilb we have a canonical antiisomorphism  $\mathcal{H} \rightarrow \mathcal{H}^*$  induced by  $\langle \cdot | \cdot \rangle$  is very different from Vect or Set, but a lot like  $n\text{Cob}$ , where the “dual” of a space is the same space with the opposite orientation, and the “adjoint” of an  $n$ -cobordism is its time-reversal. Time reversal is of utmost importance in physics.

**Definition.** Given categories  $\mathcal{C}$  and  $\mathcal{D}$ , a **functor**  $F: \mathcal{C} \rightarrow \mathcal{D}$  consists of:

- A map sending any object  $x \in \mathcal{C}$  to an object  $F(x) \in \mathcal{D}$ .
- For any pair of objects  $x$  and  $y$ , a map sending morphisms  $f: x \rightarrow y$  to morphisms  $F(f): F(x) \rightarrow F(y)$ , such that these laws hold:
  - for any object  $x \in \mathcal{C}$ , we have  $F(1_x) = 1_{F(x)}$ .
  - for any pair of morphisms  $f: x \rightarrow y$  and  $g: y \rightarrow z$ , we have  $F(gf) = F(g)F(f)$ .

In short:  $F$  sends objects to objects, morphisms to morphisms, and preserves sources, targets, identities, and composition. ★

**Definition.** We say that a category  $\mathcal{C}$  has **adjoints** or **duals for morphisms** or is a **\*-category** if there is a contravariant functor  $*$ :  $\mathcal{C} \rightarrow \mathcal{C}$  which takes objects to themselves and such that  $*^2 = 1$  (the identity functor). For any object  $x$  or morphism  $f$ , the dual is denoted  $*(x) = x^*$  or  $*(f) = f^*$ . ★

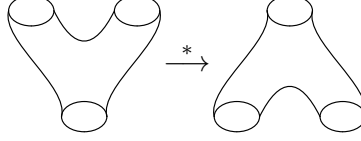
Spelling out the definition,  $*$  has to satisfy the following properties:

- $x^* = x$  for any  $x \in \mathcal{C}$ ,
- for any  $f: x \rightarrow y$  there is a morphism  $f^*: y \rightarrow x$  (this is what “contravariant” means),
- for any  $x \in \mathcal{C}$ ,  $(1_x)^* = 1_{x^*} = 1_x$ ,
- for any morphisms  $f: x \rightarrow y$  and  $g: y \rightarrow z$ , we have  $(gf)^* = f^*g^*$ , and
- $(f^*)^* = f$  for any morphism  $f$ .

<sup>1</sup>More generally, any **normal** operator, i.e. any operator such that  $NN^* = N^*N$ , has an orthonormal basis of eigenvectors with complex eigenvalues.

Examples of  $*$ -categories are:

- $n\text{Cob}$ , where  $M^*$  is obtained by exchanging the roles of input and output. If the cobordism is imbedded, this can be represented as reflection along the “time” direction.

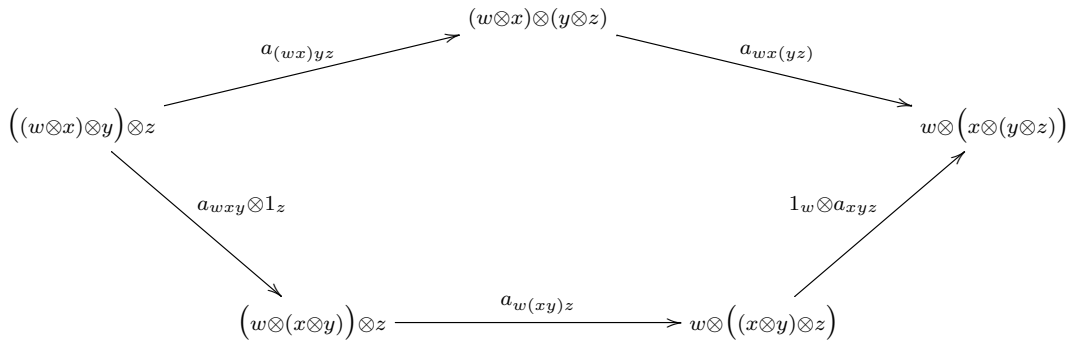


- $\text{Hilb}$ , where the adjoint  $T^*$  of a linear operator  $T: \mathcal{H} \rightarrow \mathcal{H}'$  is defined by  $\langle T^*\phi | \psi \rangle_{\mathcal{H}} = \langle \phi | T\psi \rangle_{\mathcal{H}'}$ .
- any **groupoid** (a category where every morphism is invertible), as then the inverse has the properties required of  $*$ .

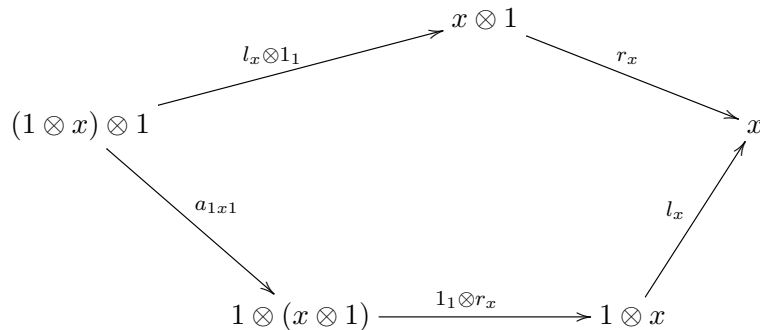
**Definition.** We say that  $F: \mathcal{C} \rightarrow \mathcal{D}$  is a  $*$ -functor if given  $f: x \rightarrow y$ , we have  $F(f^*) = F(f)^*: F(x) \rightarrow F(y)$ . ★

**Definition.** A category  $\mathcal{C}$  is **monoidal** if it is equipped with an operation  $\otimes$  with the following properties:

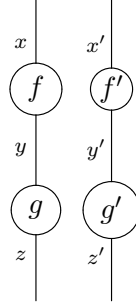
- for any  $x, y \in \mathcal{C}$ , there is an object  $x \otimes y \in \mathcal{C}$ ;
- for any  $f: x \rightarrow x'$  and  $g: y \rightarrow y'$ , there is a morphism  $f \otimes g: x \otimes y \rightarrow x' \otimes y'$ .
- for any objects  $x, y, z \in \mathcal{C}$  there is an isomorphism  $a_{xyz}: (x \otimes y) \otimes z \rightarrow x \otimes (y \otimes z)$  called the **associator** and satisfying the **pentagon identity**:



- there is an object  $1$  such that, for any object  $x \in \mathcal{C}$ , there are isomorphisms  $l_x: 1 \otimes x \rightarrow x$  and  $r_x: x \otimes 1 \rightarrow x$  called **units** satisfying the **other identity**:



- finally, given  $f: x \rightarrow y$ ,  $g: y \rightarrow z$ ,  $f': x' \rightarrow y'$  and  $g': y' \rightarrow z'$ , we require that  $(g \otimes g')(f \otimes f') = (gf) \otimes (g'f')$ , which just says that the following diagram is unambiguous:



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MacLane's theorem guarantees that if the above two diagrams commute, then any diagram that can be constructed from the associator and the units commutes.

Examples of monoidal categories are

- Grp: objects are groups, morphisms are group homomorphisms and  $\otimes$  is the direct product of groups.
- nCob: the  $\otimes$ , both for objects and for morphisms, is the disjoint union of manifolds.
- Vect or Hilb: the  $\otimes$  is the tensor product. This is how, in quantum mechanics, two things are put together.
- Elect: it has just one object, morphisms are electrical circuit elements, composition is serial combination of components, and  $\otimes$  is parallel or shunted combination of components.

*Note:* An algebra equipped with a nondegenerate trace is called a **Frobenius algebra**. We've just seen that given any TQFT, the Hilbert space  $Z(S^1)$  is a commutative Frobenius algebra with multiplication given by

$$Z(m): Z(S^1) \otimes Z(S^1) \rightarrow Z(S^1),$$

unit given by

$$Z(i): \mathbb{C} \rightarrow Z(S^1),$$

and trace given by

$$Z(i^*): Z(S^1) \rightarrow \mathbb{C}.$$

But the cool part is the converse: for any Hilbert space with the structure of a commutative Frobenius algebra, there exists a unique 2d TQFT. Actually, uniqueness isn't hard. The only real work is to figure out a formula for  $Z(m^*)$  in terms of the 3 maps just listed.

## FAITHFULNESS, FULLNESS, EQUIVALENCES

**Definition.** A functor  $F: C \rightarrow D$  is called **faithful** if for each pair of objects  $X, Y \in C$ , the map  $F_{X,Y}: \text{Hom}_C(X, Y) \rightarrow \text{Hom}_D(F(X), F(Y))$  is injective.  $F$  is called **full** if the maps  $F_{X,Y}: \text{Hom}_C(X, Y) \rightarrow \text{Hom}_D(F(X), F(Y))$  are all surjective. If  $C$  is a subcategory of  $D$ , then the inclusion functor is always faithful. If it is full,  $C$  is called a **full subcategory**. ★

**Definition.** A functor is called **essentially surjective** if every object in  $D$  is isomorphic to an image under  $F$  of an object of  $C$ . A functor is called an **equivalence** if it is faithful, full, and essentially surjective. ★