

(17)

Regular Expression

- Priority of operators of RE :

$$()^* , a^+ , a , \cdot , +$$

- Definition :

1.) RE is said to be valid iff it can be derived from RE by a finite no. of applications of the rule a^* , a^+ , $a_1 \cdot a_2$, $a_1 + a_2$.

2.) If Σ is a given alphabet then, $\phi, \epsilon, 1$, $a \in \Sigma$ are REs.

obj. 1

$\Phi = \text{rule}$
 $\Sigma = \text{string of length } n$

generating
 families

① $a + \phi = a$

② $a + \epsilon = \{a, \epsilon\}$

③ $\mu = a + b, d(\mu) = \{a, b\}$

④ $\mu = a \cdot b, d(\mu) = \{a \cdot b\}$

⑤ $\mu = a + b + c, d(\mu) = \{a, b, c\}$

⑥ $\mu = (ab + a) \cdot b, d(\mu) = \{a, b, ab\}$

⑦ $\mu = a^+, d(\mu) = \{a, aa, aaa, aaaa, \dots\}$

⑧ $\mu = a^*, d(\mu) = \{a, a, aa, aaa, \dots\}$

⑨ $\mu = (a + ba)(b + a), d(\mu) = \{a, b, aa, bab, baq\}$

⑩ $\mu = (a + \epsilon)(\underbrace{b + \phi}_b), d(\mu) = \{a, b, b\}$

⑪ $\mu = (a + b)(a + b), d(\mu) = \{aa, ab, ba, bb\}$

⑫ $\mu = (a + b)^*, d(\mu) = \{a, b\}$

③

$$\mu(\mu) \Rightarrow (a+b)^0 = \epsilon$$

$$(a+b)^1 = \{a, b\}$$

$$(a+b)^2 = \{aa, ab, ba, bb\}$$

$$(a+b)^3 = \{~~aaa~~, b, \dots\}$$

$$\therefore \{ \epsilon, ab, aabb, \dots \}$$

⑬ $\mu = (a+b)^* \cdot (a+b)$

$$\{ \epsilon, \dots \} \cdot (a+b) = (a+b)^+$$

⑭ $\mu = a^* \cdot a^* = a^*$

$$\{ \epsilon, aa, \dots \} \cdot \{ \epsilon, aa, \dots \}$$

⑮ $\mu = (ab)^* = \{ \epsilon, ab, abab, ababab, \dots \}$

$$(ab)^0$$

$$(ab)^1$$

$$(ab)^2 = (ab)(ab)$$

$$(ab)^3 = (ababab)$$

eg 1: $u_1 = a^*$
 $u_2 = a^* + (aa)^*$

- (a) $d(u_1) \subseteq d(u_2)$
 (b) $d(u_1) \supseteq d(u_2)$
 ✓ (c) $d(u_1) = d(u_2)$
 (d) $d(u_1) \neq d(u_2)$

Ans 1
 same language can be generated by
 one or more Regular Expressions.

Ques 2: ϕ = empty set / null set $\therefore = \emptyset$
 ϵ = empty string.

(a) $u = \epsilon^*, d(u) = \{\epsilon^0, \epsilon^1, \epsilon^2, \dots\}$
 $= \{\epsilon, \epsilon, \epsilon, \dots\}$

to remove
 duplicacy
 in set

(b) $u = \epsilon^+, d(u) = \{\epsilon^1, \epsilon^2, \dots\} = \{\epsilon\}$

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$$(c) \mu = \phi^*, d(\mu) = \{\phi^0, \phi^1, \phi^2, \dots\} = \{\epsilon\}$$

$$A = \{a, b\}$$

eg:- $A^0 = \{\epsilon\}$ string of length 0

$$A^1 = \{a, b\}$$

$$A^2 = \{aa, ab, ba, bb\}$$

$$(d) \mu = \phi^+, d(\mu) = \{\phi^1, \phi^2, \dots\} = \{\phi\}$$

Questions:

$$① \mu^+ \cup \mu^* = \mu^*$$

$$② \mu^+ \cap \mu^* = \mu^+$$

$$③ \mu^* \cdot \mu^+ = \mu^+$$

$$④ (\mu^*)^* = \{\epsilon^0, a^1, a^2, \dots\}^* = \{\epsilon, a, aa, \dots\}^*$$

$$\Rightarrow \{\epsilon^*, a^*, aa^*, \dots\} = \{\epsilon, a, aa, \dots\}$$

$$\Rightarrow \mu^*$$

$$⑤ (\mu^*)^+ = \mu^*$$

versions

$$(6) (u^+)^* = u^*$$

$$(7) (u^*)^+ = u^*$$

$$(8) (a+b)^* = (a^*+b^*)^*$$

$$(9) (a+b)^* \neq (a^*+b^*)^*$$

$$(10) (a+b)^* \neq (a \cdot b)^*$$

$$(11) (a+b)^* \neq (a^* \cdot b)^*$$

* if you want ba from $(a^* b^*)^*$

$$(a^* b^*)^2 \Rightarrow (a^* b^*) (a^* b^*)$$

for rep
algebra

$\langle \phi \rangle$

symbol

regular

expression

R_1

ε

$a^* - \{$

$a^+ = \{$

$a^+ a =$

Regul# How to design RE from language (7)

follows

① Start with ab:

$\{ab, aba, abb, abba, \dots\}$

$\Rightarrow ab(a+b)^*$

② Start with bba:

$\Rightarrow bba(a+b)^*$

③ Ends with abb:

$\Rightarrow (a+b)^*abb$

④ Contains substring abb.

$\Rightarrow (a+b)^*abb(a+b)^*$

⑤ Starts and ends with a

- if a is starting and ending ~~both~~
 - if in start + ending
 - if in start & ending
- $\Rightarrow a + a(a+b)^*a$

⑥ start and ends with same symbol

$$\Rightarrow a + a(a+b)^*a + b + b(a+b)^*b.$$

⑦ starts and end with different symbol.

$$\Rightarrow a(a+b)^*b + b(a+b)^*a.$$

⑧ $|w| = 3$

$$(a+b)(a+b)(a+b)$$

⑨ $|w| \geq 3$

$$(a+b)^3(a+b)^*$$

⑩ $|w| \leq 3$

$$\Rightarrow \left\{ \underbrace{\varepsilon}_0, \underbrace{a}_1, \underbrace{b}_1, \underbrace{aa}_{2}, \underbrace{bb}_{2}, \underbrace{ba}_{2}, \dots \right\}$$

$$\Rightarrow \underbrace{\varepsilon}_0 + \underbrace{(a+b)}_1 + \underbrace{(a+b)^2}_2 + \underbrace{(a+b)^3}_3$$

$$\Rightarrow (a+b+\varepsilon)$$

$$① |w|_a = 2$$

$$b^* a b^* a b^*$$

$$② |w|_a \geq 2$$

$$(a+b)^* a (a+b)^* a (a+b)^*$$

$$③ |w|_a \leq 2$$

$$e + b^* a b^* + b^* a b^* a b^*$$

④ Symbol from left end is 1

$$(a+b)^2 b(a+b)^*$$

$$⑤ |w| = 0 \pmod{3} \rightarrow \text{for } 0$$

0 3 6 9

$$[(a+b)^3]^*$$

in context of regular

$$⑥ |w| = 2 \pmod{3} \rightarrow (a+b)^2 [(a+b)^3]^*$$

Regular

string

→

=

=

$q_1 a$

$q_1 a$

at

=

$\Rightarrow (q_1 c$

$\Rightarrow q_1 a$

④

ing

view.

, z,

various.

is also
 $\Rightarrow (R_1 + R_2)$

$\Rightarrow (R_1 \cdot R_2)$

*

, a, aa

, aaa

sign

$$(7) \quad |w|_b = 0 \pmod{2}$$

$$b = 0 \quad 2 \quad 4 \quad 6 \quad 8 \dots$$

$$a^* + (a^* b \ a^* b \ a^*)^*$$

ne

$$(8) \quad |w|_a = 1 \pmod{3}$$

$$\checkmark \quad b^* a \ b^* \quad \bullet \quad (b^* a \ b^* a \ b^* a \ b^*)^*$$

3

$$(9) \quad |w|_b = 2 \pmod{3}$$

bb

$$(7) \quad a^* \left(\frac{bb}{\downarrow} \right)^* a^* \quad \downarrow \quad \begin{matrix} b b & b b & b b \end{matrix}$$

~~concatenative~~

$$(8) \quad \left(\frac{b^* a \cdot (a a a)^* b^*}{\downarrow} \right)^*$$

1 4 7 10 13

$$a^* b b a^* \cdot (a^* b \ a^* b \ a^* b \ a^*)^*$$

ut

=

=> (9

9