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## Regular Expressions

- Regular expressions are used for representing certain sets of strings in an algebraic fashion.

{Rules}

$\rightarrow \boxed{a, b, c, \dots, \Lambda, \phi}$

1. Any terminal symbol i.e symbols  $a, b, c, \dots$

including  $\Lambda$  and  $\phi$  are regular expressions.  
 $\{a, b, c, \Lambda, \phi\}$

2. The union of two regular expressions is also a regular expression  
 $R_1, R_2 \Rightarrow (R_1 + R_2)$

3. Concatenation is also RE.  $R_1, R_2 \Rightarrow (R_1 \cdot R_2)$

4. Closure of RE is RE  $R \rightarrow R^*$

$$a^* = \{a, aa, aaa, \dots\}$$

$$a^+ = \{a, aa, aaa, \dots\}$$

$$a^+ \cdot a = a^*$$

# 10.1.1 THEOREM ( $\text{FA} \rightarrow \text{RE}$ )

(see Ex. 10.1.1)

## Examples:

describe the following sets as Regular Expressions:

$$1) \{0, 1, 2\} : 0 \text{ or } 1 \text{ or } 2$$

$$R = 0 + 1 + 2$$

$$2) \{1, ab\}$$

$$R = 1 \mid ab$$

↑  
this is power set

\*)

$$3) \{abb, a, b, bba\} : abb \text{ or } a \text{ or } b \text{ or } bba$$

$$R = abb + a + b + bba$$

$$4) \{1, 0, 00, 000, \dots\} : R = 0^* \text{ all strings that can be formed with no including empty symbol.}$$

$$5) \{1, 11, 111, 1111, \dots\} : R = 1^+$$

$\cdot \text{P} \text{P}^2$

Part Identities of Regular Expressions

and

1)  $\phi + R = R$

for

$\phi = \text{empty set} \therefore \phi \cup R = R$

h

2)  $\phi \cdot R + R \phi = \phi$

Proof

3)  $\phi R \doteq R \phi = R$

$$\{\Lambda = \epsilon\}$$

4)  $\epsilon^* = \epsilon$  and  $\phi^* = \epsilon$

5)  $R + R = R$

6)  $R^* \cdot R^* = R^*$

7)  $RR^* = R^*R$

8)  $(R^*)^* = R^*$

9)  $\epsilon + RR^* = \epsilon + R^*R = R^*$

10)  $(PQ^*)P = P(QP)^*$

11)  $(P+Q)^* = (P^*Q^*)^* = (P^*+Q^*)^*$

12)  $R(P+Q) = RP + RQ$  and  $(P+Q)R = PR + QR$

design RE for following language over  $\Sigma$ .

1) language accepting string of length exactly 2.

$$L_1 = \{aa, ab, ba, bb\}$$

$$R_1 = aa + ab + ba + bb$$

$$a(a+b) + b(a+b) \Rightarrow \boxed{(a+b)(a+b)}$$

2) language accepting of length atleast 2.

$$L_2 = \{aa, ab, ba, bb, aaa, bbb, abab, \dots\}$$

$$R_2 = \underbrace{(a+b)(a+b)}_{\text{atleast 2}} (a+b)^*$$

↑  
anything more than length 2.

abc

3) language accepting string of length atleast 2.

$$L_3 = \{e, a, b, aa, bb, ab, ba\}$$

$$R_3 = e + a + b + aa + ab + bb + ba$$

$$= (e + a + b)(e + a + b)$$

①

ab      ba

1 2      ↑ ↑  
          a, b, c → 0, 1, 2

$$\underbrace{(a+bc)^* a (a+bc)^* b (a+bc)^* + (a+bc)^* b c}_{\text{---}}$$

$$(0+1)^* 1 \textcircled{0} 1 \text{---} \text{---} \text{---}$$

$$(0+1)^* 1 (0+1) (0+1)^0 \text{---}$$

$$(0+1)^9$$

$$101$$

$$\begin{matrix} \Sigma^1 \\ \Sigma^2 \\ \Sigma^3 \end{matrix}$$

$$(0^* 1^* 0^*)^* 0^* 1^*$$

$$\textcircled{101}$$

$$\underline{\epsilon, 0, 00, \epsilon, 0, \dots}$$

$$1 \textcircled{1101} \underline{0100}$$

$$0^* = \{ \epsilon, 0^n b^{2m+1} : n \geq 0 \wedge m \geq 0 \}$$

$$(0^* 1^* 0^*)^*$$

2x0+1

$$\begin{matrix} 1 \\ 0 \\ 1 \\ 0 \\ 1 \end{matrix} \begin{matrix} b^0 \\ b^0 \\ b^3 \\ b^5 \\ b^7 \end{matrix}$$

$$\epsilon, 0100, 01000100, 0$$

$$(aa)^* (bb)^* b$$

$$\begin{matrix} 0, aa, aaaa \\ \underline{bb} \\ \underline{bbbb} \\ \underline{bbbbb} \end{matrix} b$$

abba

$(a+ab)^*$

a, a, b, aa, bbb, ab

$R^*R$

$P^+$