

Insertion Sort

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Programming, Data Structures and Algorithms using Python
Week 2

Sorting a list

- You are the TA for a course
 - Instructor has a pile of evaluated exam papers
 - Papers in random order of marks
 - Your task is to arrange the papers in descending order of marks

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 - **Insert** into correct position with respect to first two
- Do this for the remaining papers
 - **Insert** each one into correct position in the second pile

Sorting a list

74 32 89 55 21 64

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~~74~~ 32 89 55 21 64

74

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32 74

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def InsertionSort(L):  
    n = len(L)  
    if n < 1:  
        return(L)  
    for i in range(n):  
        # Assume L[:i] is sorted  
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        j = i  
        while(j > 0 and L[j] < L[j-1]):  
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Analysis of iterative insertion sort

- Correctness follows from the invariant

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- $T(n)$ is $O(n^2)$

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Analysis of recursive insertion sort

- For input of size n , let
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Analysis of recursive insertion sort

- For input of size n , let
 - $TI(n)$ be the time taken by `Insert`
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- First calculate $TI(n)$ for `Insert`
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- Unwind to get $1 + 2 + \dots + n - 1$

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- Worst case complexity is $O(n^2)$
 - Unlike selection sort, not all cases take time n^2
 - If list is already sorted, **Insert** stops in 1 step
 - Overall time can be close to $O(n)$