

Searching in a List

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Programming, Data Structures and Algorithms using Python
Week 2

Search problem

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- Naive solution scans the list

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            return(True)  
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- Is value v present in list l ?
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- Input size n , the length of the list
- Worst case is when v is not present in l
- Worst case complexity is $O(n)$

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Searching a sorted list

- What if 1 is sorted in ascending order?

Searching a sorted list

- What if l is sorted in ascending order?
- Compare v with the midpoint of l

Searching a sorted list

- What if `l` is sorted in ascending order?
- Compare `v` with the midpoint of `l`
 - If midpoint is `v`, the value is found
 - If `v` less than midpoint, search the first half
 - If `v` greater than midpoint, search the second half
 - Stop when the interval to search becomes empty

```
def binarysearch(v,l):  
    if l == []:  
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    m = len(l)//2  
  
    if v == l[m]:  
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    if v < l[m]:  
        return(binarysearch(v,l[:m]))  
    else:  
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■ Binary search

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Binary search

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 - Each call halves the interval to search
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- $\log n$ — number of times to divide n by 2 to reach 1
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- $O(\log n)$ steps

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Alternative calculation

- $T(n)$: the time to search a list of length n
 - If $n = 0$, we exit, so $T(n) = 1$
 - If $n > 0$, $T(n) = T(n // 2) + 1$

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def bsearch(v,l):  
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 - Worst case is when the value is not present in the list
- For a sorted list, binary search takes time $O(\log n)$
 - Halve the interval to search each time
- In a sorted list, we can determine that v is absent by examining just $\log n$ values!