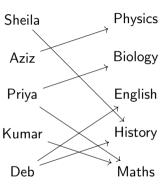
Madhavan Mukund

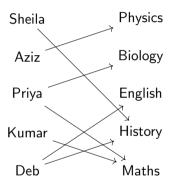
https://www.cmi.ac.in/~madhavan

Programming, Data Structures and Algorithms using Python Week 11

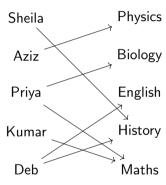
 Each instructor is willing to teach a set of courses



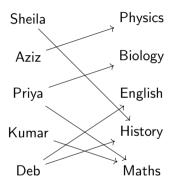
- Each instructor is willing to teach a set of courses
- Find an allocation so that
 - Each course is taught by a single instructor
 - Each instructor teaches only one course, which he/she is willing to teach



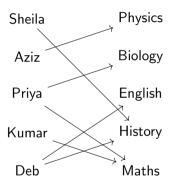
• V partitioned into V_0 , V_1



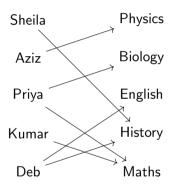
- \blacksquare V partitioned into V_0 , V_1
- All edges from V_0 to V_1



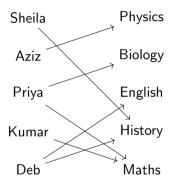
- V partitioned into V_0 , V_1
- All edges from V_0 to V_1
- Matching: subset of edges so that no two of them share an endpoint



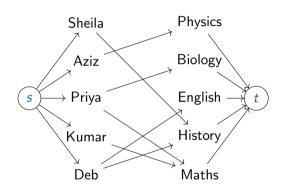
- V partitioned into V_0 , V_1
- All edges from V_0 to V_1
- Matching: subset of edges so that no two of them share an endpoint
- Find largest matching



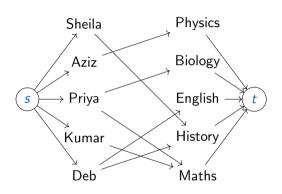
- V partitioned into V_0 , V_1
- All edges from V_0 to V_1
- Matching: subset of edges so that no two of them share an endpoint
- Find largest matching
- If possible, a perfect matching, all nodes covered



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- All edges from V_0 to V_1
- Matching: subset of edges so that no two of them share an endpoint
- Find largest matching
- If possible, a perfect matching, all nodes covered
- Add a source and a sink
 - All edge capacities 1



- V partitioned into V_0 , V_1
- All edges from V_0 to V_1
- Matching: subset of edges so that no two of them share an endpoint
- Find largest matching
- If possible, a perfect matching, all nodes covered
- Add a source and a sink
 - All edge capacities 1
- Find a maximum flow from s to t!

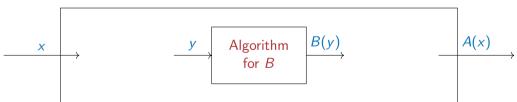


PDSA using Python Week 11

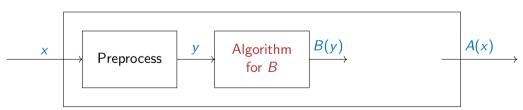
■ We want to solve problem *A*



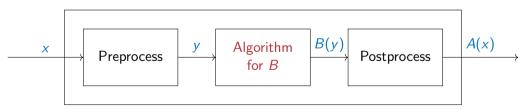
- We want to solve problem A
- We know how to solve problem *B*



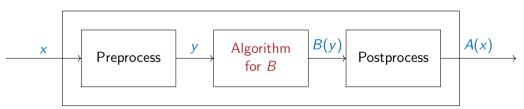
- We want to solve problem A
- We know how to solve problem *B*
- Convert input for *A* into input for *B*



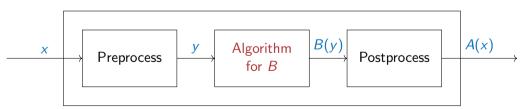
- We want to solve problem A
- We know how to solve problem *B*
- Convert input for *A* into input for *B*
- Interpret output of *B* as output of *A*



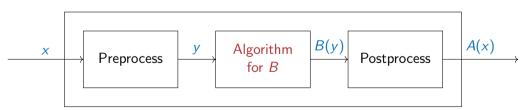
 \blacksquare A reduces to B



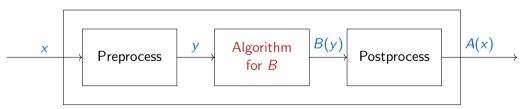
- \blacksquare A reduces to B
- Can transfer efficient solution from B to A



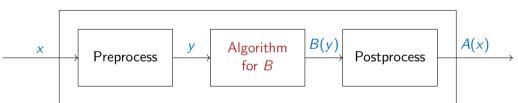
- \blacksquare A reduces to B
- Can transfer efficient solution from B to A
- But preprocessing and postprocessing must also be efficient!



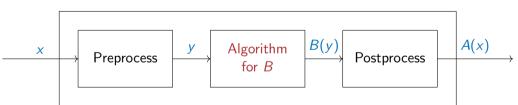
- A reduces to B
- Can transfer efficient solution from B to A
- But preprocessing and postprocessing must also be efficient!
- Typically, both should be polynomial time



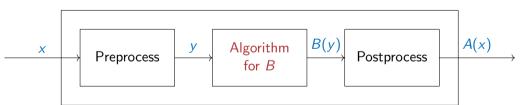
■ Bipartite matching reduces to max flow



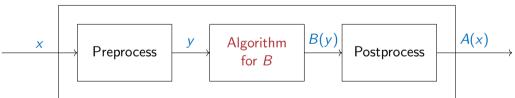
- Bipartite matching reduces to max flow
- Max flow reduces to LP



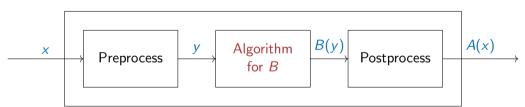
- Bipartite matching reduces to max flow
- Max flow reduces to LP
- Number of variables, constraints is linear in the size of the graph



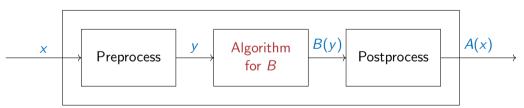
■ Reverse interpretation is also useful



- Reverse interpretation is also useful
- If A is known to be intractable and A reduces to B, then B must also be intractable



- Reverse interpretation is also useful
- If A is known to be intractable and A reduces to B, then B must also be intractable
- Otherwise, efficient solution for B will yield efficient solution for A



■ LP and network flows are powerful tools

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- Many algorithmic problems can be reduced to them
- Efficient, off-the-shelf implementations are available
- Useful to understand what can (and cannot) be modelled in terms of LP and flows