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Q	Process	Allocation				Max				Available			
		A	B	C	D	A	B	C	D	A	B	C	D
	P ₀	0	0	1	2	0	0	1	2	1	5	2	0
	P ₁	1	0	0	0	1	7	5	0				
	P ₂	1	3	5	4	2	3	5	6				
	P ₃	0	6	3	2	0	6	5	2				
	P ₄	0	0	1	4	0	6	5	6				

a) Need matrix
(max - Allocation)

Process	A	B	C	D
P ₀	0	0	0	0
P ₁	0	7	5	0
P ₂	1	0	0	2
P ₃	0	0	2	0
P ₄	0	6	4	2

b) Safety Algorithm

⇒ if need ≤ available

then

execute resource process

new available = available + allocation

else

don't execute & move forward to next process

① P₀ → need ≤ available

(0, 0, 0, 0) ≤ (1, 5, 2, 0) So, execute P₀

new available = (1, 5, 2, 0) + (0, 0, 1, 2)
= (1, 5, 3, 2)

① $P_1 \rightarrow \text{need} > \text{available}$
 $(0, 7, 5, 0) > (1, 5, 3, 2)$ So, don't execute P_1

② $P_2 \rightarrow \text{need} \leq \text{available}$
 $(1, 0, 0, 2) \leq (1, 5, 3, 2)$ So, execute P_2

$$\begin{aligned} \text{new available} &= \text{available} + \text{allocation} \\ &= (1, 5, 3, 2) + (1, 3, 5, 4) \\ &= (2, 8, 8, 6) \end{aligned}$$

③ $P_3 \rightarrow \text{need} \leq \text{available}$
 $(0, 0, 2, 0) \leq (2, 8, 8, 6)$ So, execute P_3

$$\begin{aligned} \text{new available} &= \text{available} + \text{allocation} \\ &= (2, 8, 8, 6) + (0, 6, 3, 2) \\ &= (2, 14, 11, 8) \end{aligned}$$

④ $P_4 \rightarrow \text{need} \leq \text{available}$
 $(0, 6, 4, 2) \leq (2, 14, 11, 8)$ So, execute P_4

$$\begin{aligned} \text{new available} &= (2, 14, 11, 8) + (0, 0, 1, 4) \\ &= (2, 14, 12, 12) \end{aligned}$$

Again,

⑤ $P_1 \rightarrow \text{need} \leq \text{available}$
 $(0, 7, 5, 0) \leq (2, 14, 12, 12)$ So, execute P_1

$$\begin{aligned} \text{new available} &= (2, 14, 12, 12) + (1, 0, 0, 0) \\ &= (3, 14, 12, 12) \end{aligned}$$

Safe Sequence ;

$\langle P_0, P_2, P_3, P_4, P_1 \rangle$

c) If process P_1 arrives for $(0, 4, 2, 0)$

Resource Request Alg.

① If $\text{request} \leq \text{need}$, goto step 2
else error

② If $\text{request} \leq \text{available}$, goto step 3
else wait

③ $\text{Available} = \text{Available} - \text{Request}$
 $\text{Allocation} = \text{Allocation} + \text{Request}$
 $\text{Need} = \text{Need} - \text{Request}$

④ Check now state is safe or not?

\Rightarrow Need for P_1 $(0, 7, 5, 0)$

① $\text{Request} \leq \text{need}$; $(0, 4, 2, 0) \leq (0, 7, 5, 0)$

② $\text{Request} \leq \text{Available}$; $(0, 4, 2, 0) \leq (1, 5, 2, 0)$

③ $\text{Available} = \text{Available} - \text{Request}$
 $= (1, 5, 2, 0) - (0, 4, 2, 0)$
 $= (1, 1, 0, 0)$

$\text{Allocation} = \text{Allocation} + \text{Request}$
 $= (1, 0, 0, 0) + (0, 4, 2, 0)$
 $= (1, 4, 2, 0)$

$\text{Need} = \text{Need} - \text{Request}$
 $= (0, 7, 5, 0) - (0, 4, 2, 0)$
 $= (0, 3, 3, 0)$

④ Now, new snapshot will be

Process	Allocation				max				Available			
	A	B	C	D	A	B	C	D	A	B	C	D
P ₀	0	0	1	2	0	0	1	2	1	1	0	0
P ₁	1	4	2	0	1	7	5	0				
P ₂	1	3	5	4	2	3	5	6				
P ₃	0	6	3	2	0	6	5	2				
P ₄	0	0	1	4	0	6	5	6				

Need Matrix

Process	A	B	C	D
P ₀	0	0	0	0
P ₁	0	3	3	0
P ₂	1	0	0	2
P ₃	0	0	2	0
P ₄	0	6	4	2

Again, doing safety algorithm

P₀ ⇒ need ≤ available ; (0,0,0,0) ≤ (1,1,0,0) so, execute P₀
new available = (1,1,0,0) + (0,0,1,2) = (1,1,1,2)

P₁ ⇒ (0,3,3,0) > (1,1,1,2) don't execute P₁

P₂ ⇒ (1,0,0,2) ≤ (1,1,1,2) execute P₂
new available = (1,1,1,2) + (1,3,5,4) = (2,4,6,6)

P₃ ⇒ (0,0,2,0) ≤ (2,4,6,6) execute P₃
new available = (2,4,6,6) + (0,6,3,2) = (2,10,9,8)

P₄ ⇒ (0,6,4,2) ≤ (2,10,9,8) execute P₄
new available = (2,10,9,8) + (0,0,1,4) = (2,10,10,12)

P₁ ⇒ (0,3,3,0) ≤ (2,10,10,12) execute P₁
new available = (2,10,10,12) + (1,4,2,0) = (3,14,12,12)

< P₀, P₂, P₃, P₄, P₁ >