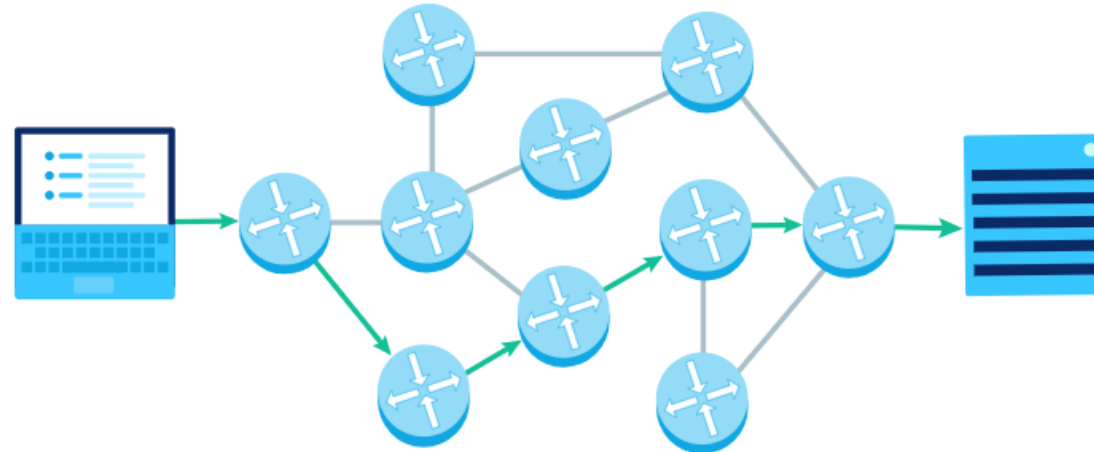


# Computer Networks

( RCS-601 )

## Routing - III



# Link State Routing

## Key Features:

- ❑ **The routers share the knowledge only about their neighbors** compared to all the routers in the autonomous system.
- ❑ **Sharing of information takes place only with all the routers** in the network, by sending small updates using flooding compared to sending larger updates to their neighbors
- ❑ Sharing of information takes place only when there is a change, which leads to lesser internet traffic compared to distance vector routing.

# Link State Routing (contd.)



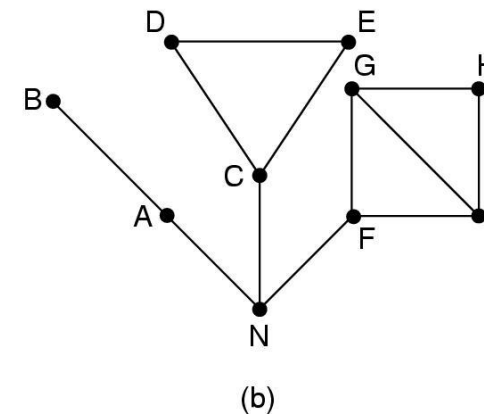
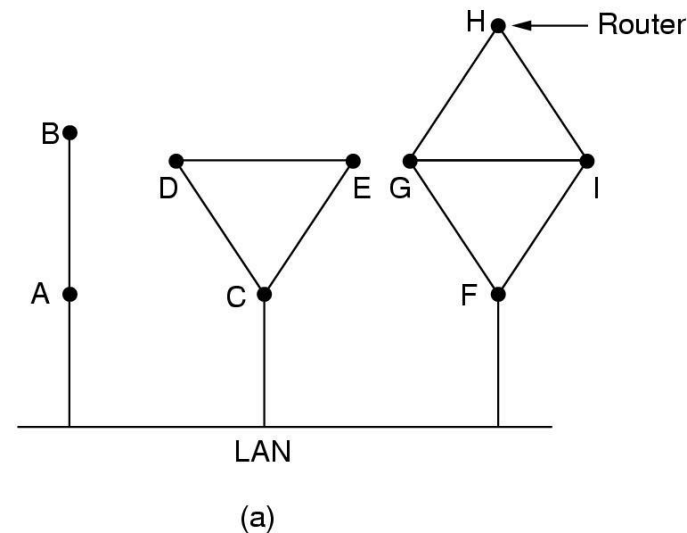
## Overview of algorithm:

Each router must:

- 1) Discover its **neighbours** and **learn** their network addresses
- 2) Measure the delay or **cost** to each of its neighbours
- 3) Construct a **link state packet** with these distances
- 4) Send this packet to **all** other routers
- 5) Compute the **shortest path** to every other router

## 1) Learning about neighbours:

- Upon boot of router
  - Send HELLO packet on each point-to-point line
  - Routers are supposed to send reply with a globally unique name
- LAN model





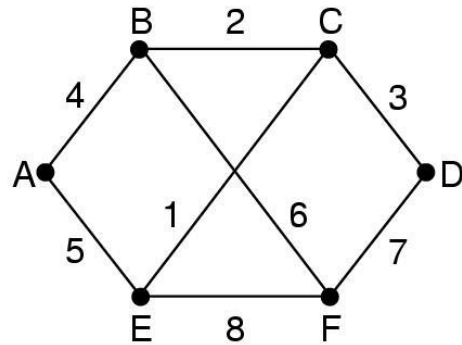
## 2) Measuring Line Cost

- The link state routing algorithm requires each router to know, or at least have a reasonable estimate of, the delay to each of its neighbors. The most direct way to determine this delay is to send over the line a special ECHO packet that the other side is required to send back immediately. By measuring the round-trip time and dividing it by two, the sending router can get a reasonable estimate of the delay. For even better results, the test can be conducted several times, and the average used.

### 3) Building link state packets

- Packet containing:
  - Identity of sender
  - Sequence number + age
  - For each neighbour: name + distance

- When to build?
  - periodically
  - when significant events occur



(a)

		Link	State			Packets	
A		B	C	D		E	F
Seq.		Seq.	Seq.	Seq.		Seq.	Seq.
Age		Age	Age	Age		Age	Age
B	4	A	B	C	3	A	B
E	5	C	D	F	7	C	D
		F	E			F	E

(b)

# Link State Routing (contd.)



## 4) Distributing link state packets

### ○ Trickiest part of algorithm

- Arrival time for packets different
- How to keep consistent routing tables

### ○ Basic algorithm

- Flooding + Sequence number (in each packet) to limit duplicates

### ○ Manageable problems

- Wrap around of sequence numbers:
- Wrong sequence number used:
  - lost in case of crash
  - Corruption

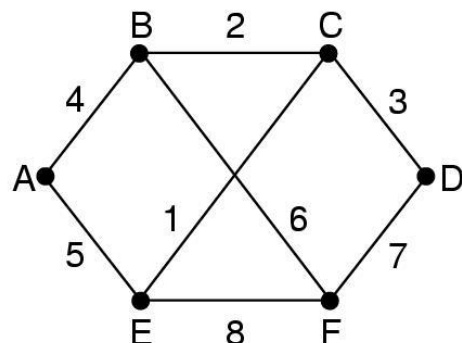
### ○ Refinements

- Link state packets are not forwarded immediately
- During holding time:
  - duplicates are discarded
  - Old packets are thrown out

## Link State Routing (contd.)

Packet buffer for router B

- ACK flag: ACK to send
- Send flag: packet to forward



Link		State		Packets	
A		B		C	
Seq.		Seq.		Seq.	
Age		Age		Age	
B	4	A	4	B	2
E	5	C	2	D	3
		F	6	E	1

Source	Seq.	Age	Send flags			ACK flags			Data
			A	C	F	A	C	F	
A	21	60	0	1	1	1	0	0	
F	21	60	1	1	0	0	0	1	
E	21	59	0	1	0	1	0	1	
C	20	60	1	0	1	0	1	0	
D	21	59	1	0	0	0	1	1	

***The packet buffer for router B***





## 5) Computing new routes:

- With a full set of link state packets, a router can:
  - Construct the entire subnet graph
  - Run Dijkstra's algorithm to compute the shortest path to each destination
- Problems for large subnets
  - Memory to store data
  - Compute time

*Thank  
you*

