

Recursion

- Recursion is a mathematical technique that evaluates a function by calling the same function repeatedly on smaller inputs.
- Most programming languages support such a style of programming.
 - Often very elegant to study.
- Helps in problem solving too.

Recursion

- **Q:** How many twists does it take to screw in a light bulb?
A: Is it already screwed in? Then zero. If not, then twist it once, ask me again, and add 1 to my answer.

Recursion

- Relates to mathematical induction
- Divide and Conquer algorithms



Lets start with an examples

- A mathematical view of computer science
- $\text{factorial}(n) = n * \text{factorial}(n-1)$ *equation*

$$\text{factorial}(1) = 1$$



09:17



Factorial – deeper look

$\text{if } (n == 1) \text{ return } 1$

$\text{else return } n \times \text{fact}(n-1)$

$$T(n) = T(n-1) + 3$$

$$T(n) = T(n-2) + 6$$

$$T(n) = T(n-k) + k \cdot 3$$

$$T(n) = T(0) + n \cdot 3$$

$\text{fact}(5)$

	return	state
$\rightarrow F(5)$	$5 \times F(4)$	R
$\rightarrow F(4)$	$4 \times F(3)$	R
$\rightarrow F(3)$	$3 \times F(2)$	R
$\rightarrow F(2)$	$2 \times F(1)$	R
$F(1)$	1	

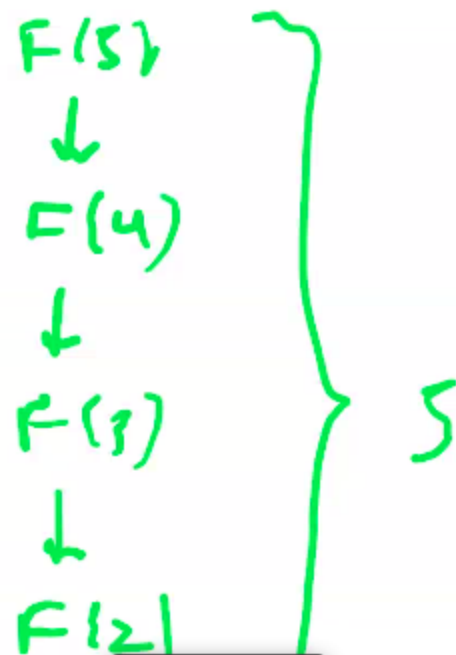
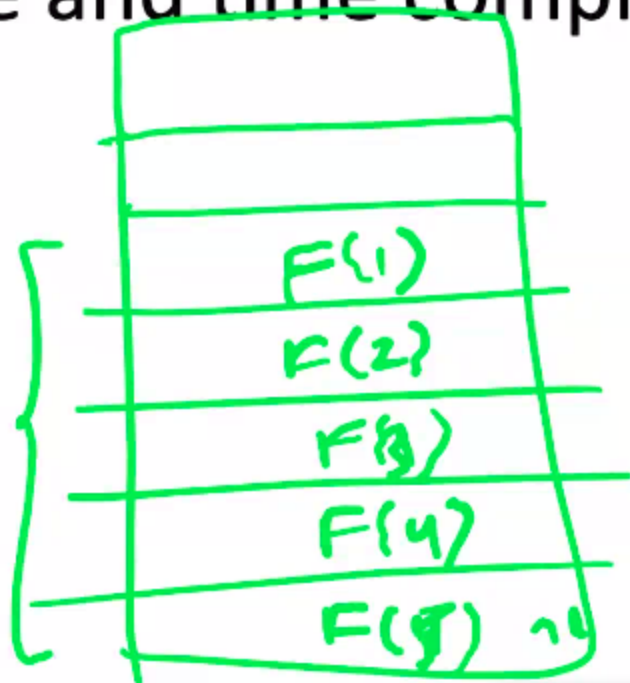


19:19



Factorial – deeper look

- Space and time complexity



Show participants



21:03



Inductive Reasoning

- A recursive algorithm must have a **base case**.
- If I call `factorial(n)` with $n=1$, I am done
- If I call `factorial(n)` with $n>1$, it makes a recursive call with a smaller value of n ; must eventually reach $n=1$.

Recursion and Induction

$$\rightarrow 1 + 2 + 3 + 4 + \dots + n = n(n+1)/2$$

L.H.S R.H.S

$$\rightarrow n=0 \quad 0 = 0 \quad \checkmark$$

$$\rightarrow \underbrace{1 + 2 + 3 + \dots + k + k+1}_{k+1} = \frac{(k+1)(k+2)}{2}$$

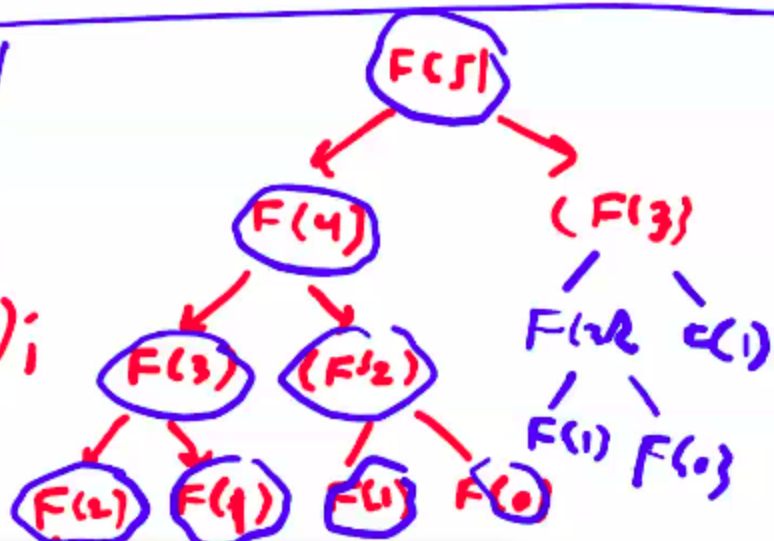
$$\frac{k(k+1)}{2} + (k+1)$$

$$= \frac{k(k+1) + 2(k+1)}{2} = \frac{(k+1)(k+2)}{2}$$

Recursion with multiple base cases

$$\text{fib}(n) = \begin{cases} \text{fib}(n-1) + \text{fib}(n-2) & \text{if } n > 1 \\ n & \text{if } n = 0, 1 \end{cases}$$

```
fib(n) {  
  if (n <= 1) return n;  
  else return fib(n-1) + fib(n-2);  
}
```



35:32



Fibonacci deeper look

- Time and space complexity

$$T(n) = T(n-1) + T(n-2) + c$$

$$T(0) = T(1) = 1$$

$$T(n-1) \approx T(n-2)$$

$$T(n) = 2T(n-2) + c$$

$$T(n) = 2[2 \cdot T(n-4) + c] + c$$

$$T(n) = 2^k T(n-2k) + (2^k - 1)c$$

$$2^{n/2}$$

Recursion with memoization

```
Fib(n)
{
  if (n <= 1) return n; ✓
  if  $F_n$  is in memory return  $F_n$ 
  // return ← Fib(n-1) + Fib( $n-2$ )
  Save  $F_n$  in memory
  return  $F_n$ 
}
```



55:44



Exponentiation

$$x^n \cdot / \cdot M$$

$$x^n = \begin{cases} n=0 & x_n = 1 \\ x \cdot x^{n-1} & \text{otherwise} \end{cases}$$

$$x^n = \begin{cases} \frac{x^{n/2} \cdot x^{n/2}}{n \text{ is even}} \\ x \cdot x^{n-1} & n \text{ is odd} \\ 1 & n \leq 0 \end{cases}$$

log n