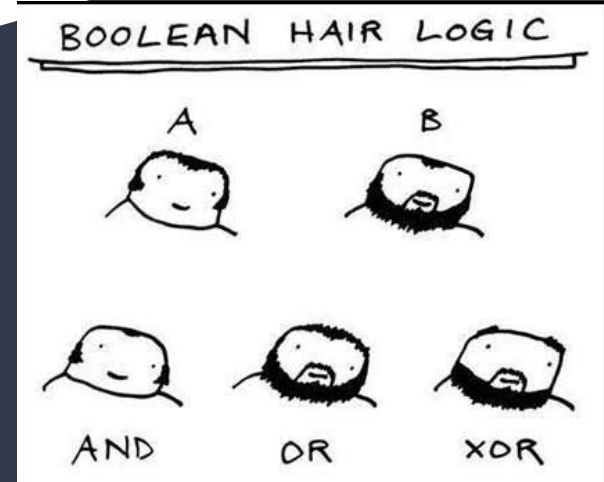


# CS 2340 – Computer Architecture

9 Procedures (Leaf/Non-Leaf), Stack  
Dr. Alice Wang



# Review

## Last Lecture

- MIPS has 3 instruction formats
  - R-, I- and J- type
- Practiced being an Assembler and converted assembly code to machine code and back

## Today

- Procedures (Leaf/Non-leaf)
  - Stack

By end of today's lecture you will have all of the skills to do the  
Term Project!

# Procedures



pro·ce·dure

/prə'sējər/

*noun*

noun: **procedure**; plural noun: **procedures**

an established or official way of doing something.

"the police are now reviewing procedures"

**Similar:**

course of action

line of action

plan of action

policy

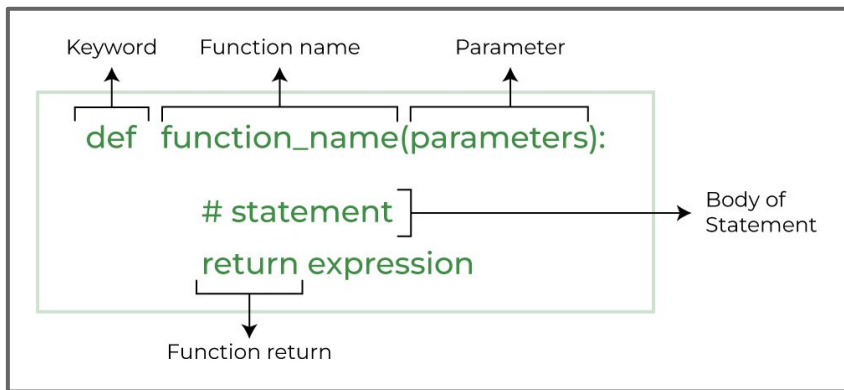
series of steps



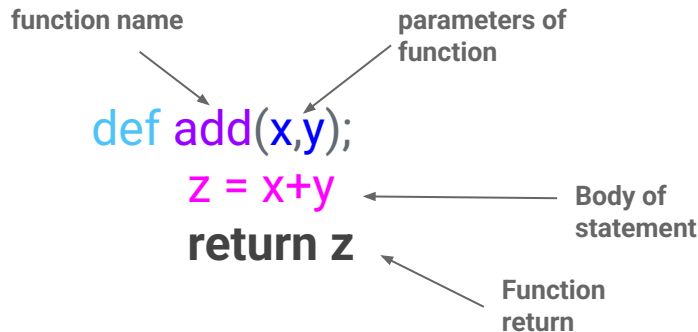
- a series of actions conducted in a certain order or manner.  
"the standard procedure for informing new employees about conditions of work"
- ~~a surgical operation.~~  
~~"the procedure is carried out under general anesthesia"~~
- **COMPUTING**  
another term for subroutine.

# Procedures enable modularity

- Most important advance in computer science
- Procedures enable structured programming
- Enable very complex programs
- They enable programmers to develop and test parts of a program in isolation
- Procedures help define interfaces between system components



## Python Example

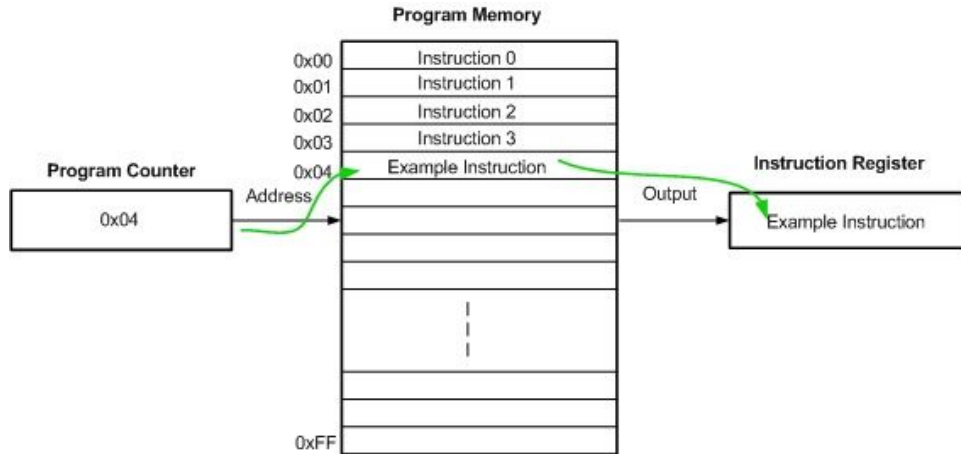
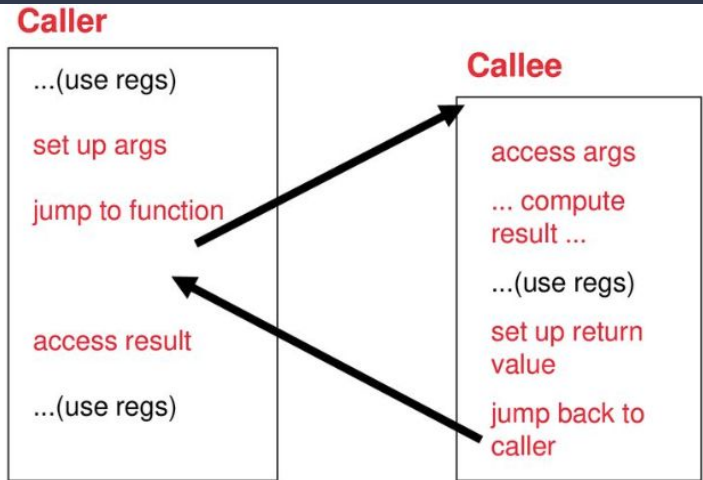


# Procedure terminology

**Caller** - The program that calls or instigates a procedure

**Callee** - A procedure that executes a series of instructions based on parameters provided by the caller

**Program counter (PC)** - The register that contains the address of the current instruction being executed



# Procedure Call and Return Instructions

- Procedure call: jump and link
  - **jal ProcedureLabel**
  - Address of following instruction put in \$ra
  - Jumps to target address ProcedureLabel
- Procedure return: jump register
  - **jr \$ra**
  - Copies \$ra to program counter
  - Can also be used for computed jumps to any register
    - e.g., for case/switch statements

# Procedure Calling

## Steps required

1. Place parameters in passed registers - **\$a0~\$a3**
2. Transfer control to procedure - **jal**
3. Acquire storage for procedure - **addi \$sp, \$sp, -4**
4. Perform procedure's operations
5. Place result in register for caller - **\$v0,\$v1**
6. Restore variables and return to place of call - **jr**

First: passing arguments and return values through registers

# Registers

- Can be overwritten by callee
  - \$a0 – \$a3: arguments
  - \$v0, \$v1: result values
  - \$t0 – \$t9: temporaries
  - ...
- Must be saved/restored by callee
  - \$s0 – \$s7: saved
  - \$gp: global pointer for static data
  - \$sp: stack pointer
  - \$fp: frame pointer

REGISTER NAME, NUMBER, USE, CALL CONVENTION

NAME	NUMBER	USE	PRESERVED ACROSS A CALL?
\$zero	0	The Constant Value 0	N.A.
\$at	1	Assembler Temporary	No
\$v0-\$v1	2-3	Values for Function Results and Expression Evaluation	No
\$a0-\$a3	4-7	Arguments	No
\$t0-\$t7	8-15	Temporaries	No
\$s0-\$s7	16-23	Saved Temporaries	Yes
\$t8-\$t9	24-25	Temporaries	No
\$k0-\$k1	26-27	Reserved for OS Kernel	No
\$gp	28	Global Pointer	Yes
\$sp	29	Stack Pointer	Yes
\$fp	30	Frame Pointer	Yes
\$ra	31	Return Address	No

**When you write a procedure you must follow the last column and preserve certain registers across a call, it's not done automatically**



# Preserved across procedure calls

“Not preserved across procedure calls”

- Register may change value if a procedure is called

“Preserved across procedure calls”

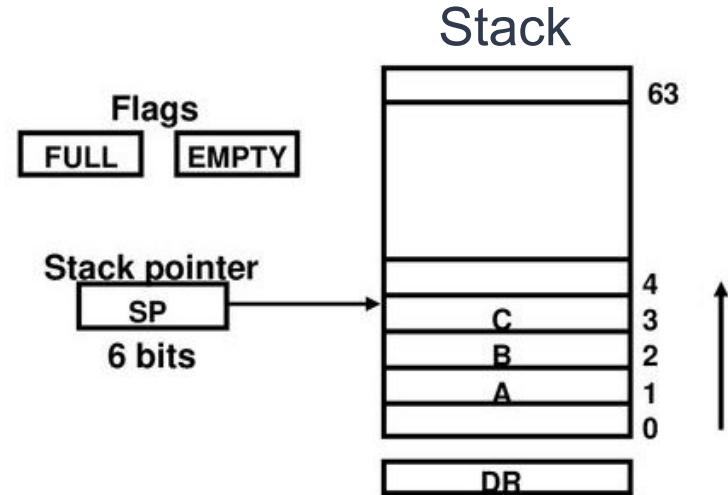
- Assume the register value will not be changed by a procedure
- Store the register value on the stack before you use it
- Restore the register from the stack before jumping back

# Procedure terminology – Stacks

**Stack** - A data structure of memory organized as last-in-first-out queue.

**Stack pointer** - A register that contains the memory address of the last data element added to the stack

**Push/Pop** - push adds an item to the top of the stack, pop removes the item from the top.

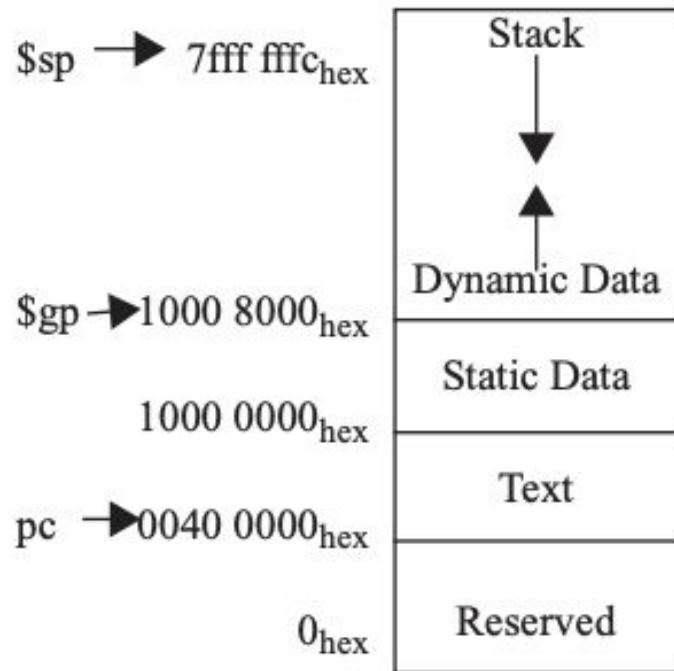


# MIPS Memory Layout

From bottom to top

- Text: program code (.text)
- Static data: global variables
  - static variables in C, constant arrays and strings (.data)
- Dynamic data: heap
  - malloc in C, new in Java, SysAlloc syscall in MIPS
- Stack: last-in first-out
  - allocated by procedures

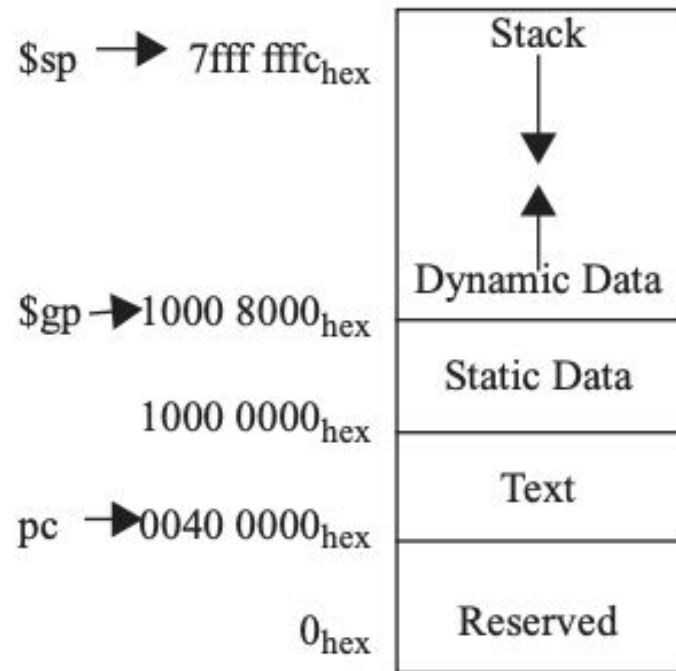
## MEMORY ALLOCATION



# Memory Layout

- Forgetting to free space using `free()` can case a “memory leak”
- Freeing space too early can turn to “dangling pointers”
- Java has automatic memory allocation to avoid these bugs
- Stack overflow - stack pointer exceeds the stack bound

## MEMORY ALLOCATION



# How can you view the Stack? (Mars Demo)

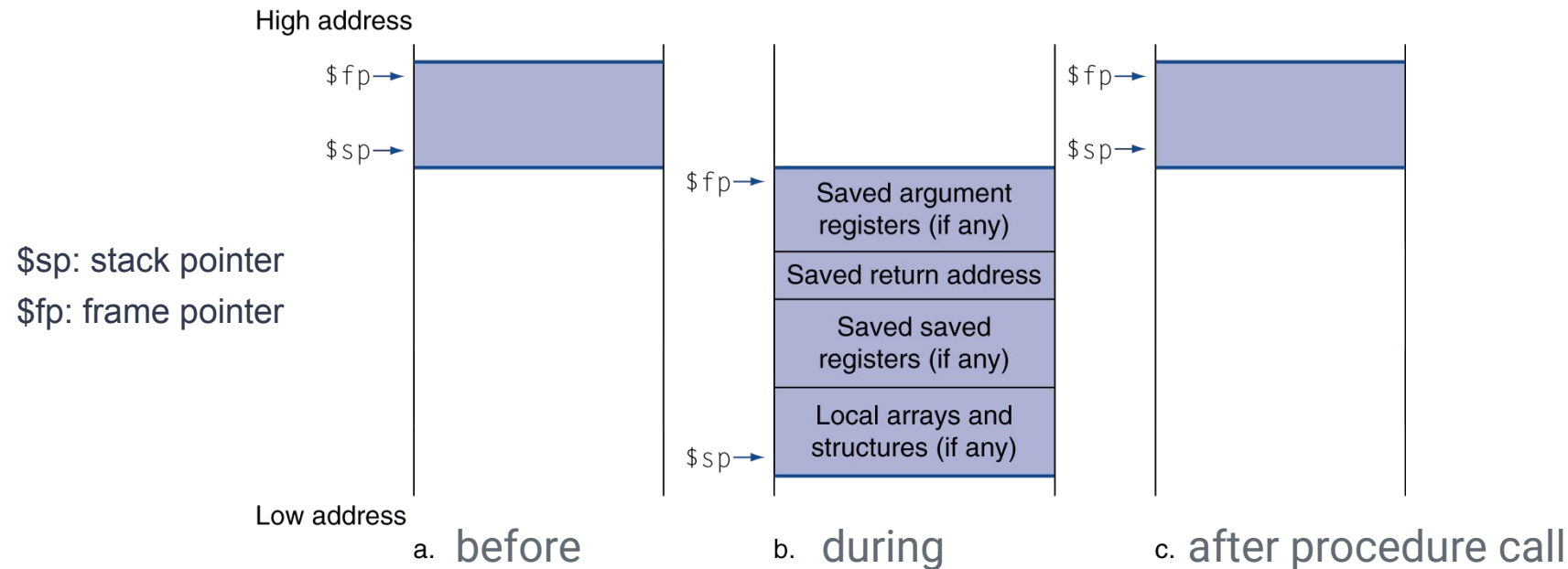
The screenshot shows the Mars Demo IDE interface. At the top is a menu bar (File, Edit, Run, Settings, Tools, Help) and a toolbar. Below the toolbar is a status bar indicating "Run speed at max (no interaction)". The main window is divided into several panes. The top pane, titled "Text Segment", displays assembly code with columns for Bkpt, Address, Code, Basic, and Source. The bottom pane, titled "Data", shows memory addresses and their corresponding values. A red box highlights the "current \$sp" option in the "Data" pane's dropdown menu. A red arrow points from the text in the yellow box to this option. The right pane shows a table of memory addresses and values for "Coprocc 1" and "Coprocc 0".

This pulldown is how you can view different parts of the memory including the instruction (.text) or the dynamic data (heap) or the stack (\$sp)

Address	Value (+0)	Value (+4)	Value (+8)	Value (+12)
0x7ffffef0	0x004002d0	0x00000000	0x00000001	0x00000000
0x7fffff00	0x00000000	0x00000000	0x00000000	0x00000000
0x7ffff020	0x00000000	0x10000000 (.extern)	0x10010000 (.data)	0x00000000
0x7ffff040	0x00000000	0x10040000 (heap)	0x00000000	0x00000000
0x7ffff060	0x00000000	current \$gp	0x00000000	0x00000000
		✓ current \$sp		
		0x00400000 (.text)		
		0x90000000 (.kdata)		
		0xffff0000 (MMIO)		

Coprocc 1	Coprocc 0
0	0x00000000
1	0x10010000
2	0x0000000a
3	0x0000005f
4	0x00000003
5	0x1001144c
6	0x00000005
7	0x00000000
8	0x00000000
9	0x00000001
10	0x00000003
11	0x00000000
12	0x00000000
13	0x0000000a
14	0x0000005f
15	0x00000052
16	0x00000052
17	0x0000000b
18	0x10010000
19	0x10010000
20	0x0000005f
21	0x00000001
22	0x00000003
23	0x00000000
24	0x0000000b
25	0x00000000
26	0x00000000
27	0x00000000
28	0x10008000
29	0x7ffffeffc
30	0x00000000
31	0x00400000
pc	0x004000f0
hi	0x00000000
lo	0x00000000

# Procedures are enabled by a stack



- Stack is a local data storage in memory for the callee
  - Data that doesn't fit in registers can be stored on the Stack

# Leaf Procedure example – not using the Stack

## Python Example

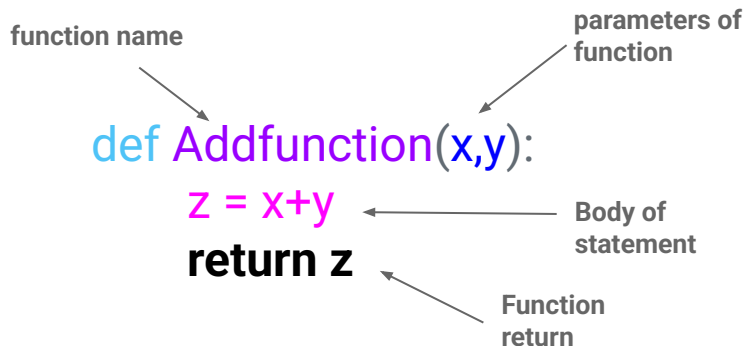
function name

parameters of function

```
def Addfunction(x,y):  
    z = x+y  
    return z
```

Body of statement

Function return



## MIPS Example

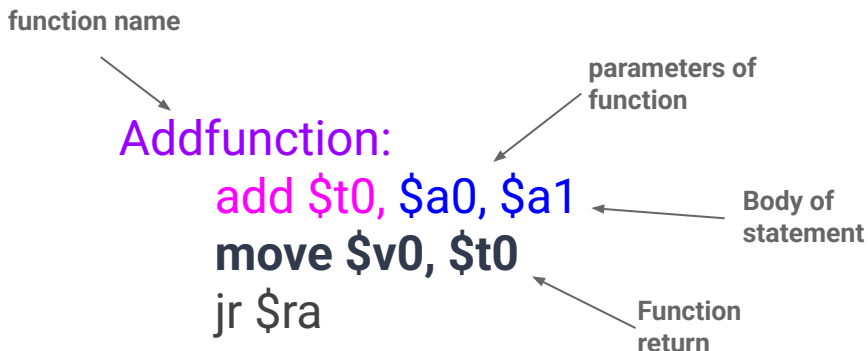
function name

parameters of function

```
Addfunction:  
    add $t0, $a0, $a1  
    move $v0, $t0  
    jr $ra
```

Body of statement

Function return



Leaf Procedure - does not call other procedures

Doesn't use any registers that need to be preserved across a call, doesn't use stack

# Leaf Procedure example – using the Stack

## Python Example

```
def Addfunction(x,y):  
    z = x+y  
    return z
```

- Use **\$s1** in the body of the statement, then I need to store **\$s1** on the stack beforehand
- **\$s1** needs to be preserved across a call

## MIPS Example

Addfunction:

```
addi $sp, $sp, -4  
sw $s1, 0($sp)
```

Allocate storage on  
stack: -4\*# of elements

Store data on the stack  
before it is used

```
add $s1, $a0, $a1  
move $v0, $s1
```

Body of  
statement

```
lw $s1, 0($sp)  
addi $sp, $sp, 4  
jr $ra
```

Restore data from the  
stack

Free storage on stack:  
4\*# of elements



# Leaf Procedure Example

Python code:

```
def Leaf_example (g, h, i, j):  
    f = (g + h) - (i + j)  
    return f
```

Step 1: Arguments g, h, i, j placed in \$a0, \$a1, \$a2, \$a3  
result f passed back through \$v0

# Leaf Procedure – My Turn

Python code:

```
def Leaf_example (g, h, i, j):  
    f = (g+h) - (i+j)  
    return f
```

***Used \$s2 to calculate f,  
thus need to save \$s2  
on the Stack***

***Restore \$s2 before  
jumping back***

MIPS code:

Leaf\_example:

```
addi $sp, $sp, -4  
sw    $s2, 0($sp)
```

```
add    $t0, $a0, $a1  
add    $t1, $a2, $a3  
sub    $s2, $t0, $t1
```

```
add    $v0, $s0, $0
```

```
lw    $s2, 0($sp)  
addi  $sp, $sp, 4  
jr    $ra
```

```
# Step 2.Transfer Control  
# Step 3. Acquire storage  
# save $s0 on stack
```

```
# Step 4. Perform  
# procedure's operation
```

```
# Step 5. Place result  
# in register for caller
```

```
# Step 6. Restore $s0 and  
# return to place of  
# call
```

# Leaf Procedure – My Turn

## Registers

\$v0	0
\$a0	3
\$a1	4
\$a2	5
\$a3	6
\$t0	0
\$t1	0
\$s0	50
\$ra	rtn_addr

MIPS code:

Leaf\_example:

```
addi $sp, $sp, -4
```

```
sw    $s0, 0($sp)
```

```
# Step 2. Transfer Control
```

```
# Step 3. Acquire storage
```

```
# save $s0 on stack
```

```
add   $t0, $a0, $a1
```

```
add   $t1, $a2, $a3
```

```
sub   $s0, $t0, $t1
```

```
# Step 4. Perform
```

```
# procedure's operation
```

```
add   $v0, $s0, $0
```

```
# Step 5. Place result
```

```
# in register for caller
```

```
lw    $s0, 0($sp)
```

```
addi  $sp, $sp, 4
```

```
jr    $ra
```

```
# Step 6. Restore $s0 and
```

```
# return to place of call
```

## Stack

	20
	1C
	18
	14
	10
	0C

\$sp

# Leaf Procedure – My Turn

## Registers

\$v0	0
\$a0	3
\$a1	4
\$a2	5
\$a3	6
\$t0	3+4=7
\$t1	5+6=11
\$s0	7-11= -4
\$ra	rtn_addr

MIPS code:

```
Leaf_example:      # Step 2. Transfer Control
    addi $sp, $sp, -4  # Step 3. Acquire storage
    sw    $s0, 0($sp)  #   save $s0 on stack

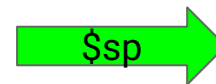
    add    $t0, $a0, $a1  # Step 4. Perform
    add    $t1, $a2, $a3  #   procedure's operation
    sub    $s0, $t0, $t1

    add    $v0, $s0, $0   # Step 5. Place result
                        #   in register for caller

    lw     $s0, 0($sp)    # Step 6. Restore $s0 and
    addi   $sp, $sp, 4    #   return to place of call
    jr     $ra
```

## Stack

	20
50	1C
	18
	14
	10
	0C



# Leaf Procedure – My Turn

## Registers

\$v0	-4
\$a0	3
\$a1	4
\$a2	5
\$a3	6
\$t0	7
\$t1	11
\$s0	-4
\$ra	rtn_addr

MIPS code:

```
Leaf_example:           # Step 2. Transfer Control
    addi $sp, $sp, -4    # Step 3. Acquire storage
    sw    $s0, 0($sp)    #   save $s0 on stack

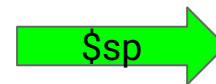
    add    $t0, $a0, $a1  # Step 4. Perform
    add    $t1, $a2, $a3  #   procedure's operation
    sub    $s0, $t0, $t1

    add    $v0, $s0, $0   # Step 5. Place result
                           #   in register for caller

    lw     $s0, 0($sp)    # Step 6. Restore $s0 and
    addi   $sp, $sp, 4    #   return to place of call
    jr     $ra
```

## Stack

	20
50	1C
	18
	14
	10
	0C



# Leaf Procedure – My Turn

MIPS code:

```
Leaf_example:          # Step 2. Transfer Control
    addi $sp, $sp, -4   # Step 3. Acquire storage
    sw    $s0, 0($sp)   #   save $s0 on stack

    add    $t0, $a0, $a1 # Step 4. Perform
    add    $t1, $a2, $a3 #   procedure's operation
    sub    $s0, $t0, $t1

    add    $v0, $s0, $0  # Step 5. Place result
                        #   in register for caller

    lw     $s0, 0($sp)   # Step 6. Restore $s0 and
    addi   $sp, $sp, 4   #   return to place of call
    jr     $ra
```

## Registers

\$v0	-4
\$a0	3
\$a1	4
\$a2	5
\$a3	6
\$t0	7
\$t1	11
\$s0	50
\$ra	rtn_addr

## Stack

	20
50	1C
	18
	14
	10
	0C



# Leaf Procedure – Your Turn


MIPS code:

Leaf\_example:

```
<Store data on stack>      # Step 2. Transfer Control
                             # Step 3. Acquire storage

add  $s2, $a0, $a1
add  $s1, $a2, $a3
sub  $s0, $s1, $s2

add  $v0, $s0, $0           # Step 5. Place result
                             #   in register for caller
<Restore data on stack>    # Step 6. Restore data and
addi $sp, $sp, 4           #   return to place of call
jr   $ra
```



Based on the program how many bytes do we need to allocate on the stack?

# Leaf Procedure – Your Turn

MIPS code:

```
Leaf_example:                # Step 2. Transfer Control
    addi $sp, $sp, -12        # Step 3. Acquire storage
    sw    $s0, 0($sp)         # save $s0 on stack
    sw    $s1, 4($sp)         # save $s1 on stack
    sw    $s2, 8($sp)         # save $s2 on stack

    add    $s2, $a0, $a1
    add    $s1, $a2, $a3
    sub    $s0, $s1, $s2

    add    $v0, $s0, $0        # Step 5. Place result
                                # in register for caller
    lw     $s0, 0($sp)        # Step 6. Restore $s0 and
    addi   $sp, $sp, 4        # return to place of call
    jr     $ra
```

What's wrong with this statement?

←



# Leaf Procedure – Your Turn

MIPS code:

```
Leaf_example:                # Step 2. Transfer Control
    addi $sp, $sp, -12        # Step 3. Acquire storage
    sw    $s0, 0($sp)         #   save $s0 on stack
    sw    $s1, 4($sp)         #   save $s1 on stack
    sw    $s2, 8($sp)         #   save $s2 on stack

    add    $s2, $a0, $a1
    add    $s1, $a2, $a3
    sub    $s0, $s1, $s2

    add    $v0, $s0, $0        # Step 5. Place result
                                # in register for caller
    lw     $s2, 8($sp)         # Restore $s2
    lw     $s1, 4($sp)         # Restore $s1
    lw     $s0, 0($sp)         # Restore $s0
    addi   $sp, $sp, 12
    jr     $ra                 # Step 6. Return to
                                # place of call
```

MIPS rules : We must restore all saved registers to original values before jumping back. We must have the stack pointer at the same place before jumping back


# Non-Leaf Procedure

- A Non-Leaf procedure calls other procedures
- Non-leaf procedures always need to use the stack, at minimum to store the return address of the calling procedure (\$ra).
- What happens if we don't store the return address? ...

# Non-Leaf Procedure – Example (bad)

*main:*

```
la $s0, A
lw $a0, 0($s0)
jal myproc1
sw $v0, 4($s0)
j exit
```



*myproc1: #non-leaf*

```
addi $v0, $a0, 5
move $a0, $v0
jal myproc2
jr $ra
```

*myproc2: #leaf*

```
addi $v0, $a0, 10
jr $ra
```

- *myproc1* is a non-leaf procedure because it calls *myproc2*
- *\$ra* is set to the address of *sw* instruction

# Non-Leaf Procedure – Example (bad)

*main:*


```
la $s0, A
lw $a0, 0($s0)
jal myproc1
sw $v0, 4($s0)
j exit
```

*myproc1:* *#non-leaf*

```
addi $v0, $a0, 5
move $a0, $v0
jal myproc2
jr $ra
```

*myproc2:* *#leaf*

```
addi $v0, $a0, 10
jr $ra
```



- *myproc1* is a non-leaf procedure because it calls *myproc2*
- *\$ra* is set to the address of *sw* instruction
- When *myproc2* is called *\$ra* is set to the address of *jr* instruction

# Non-Leaf Procedure – Example (bad)

*main:*


```
la $s0, A
lw $a0, 0($s0)
jal myproc1
sw $v0, 4($s0)
j exit
```

*myproc1:* *#non-leaf*

```
addi $v0, $a0, 5
move $a0, $v0
jal myproc2
jr $ra
```

*myproc2:* *#leaf*

```
addi $v0, $a0, 10
jr $ra
```



- *myproc1* is a non-leaf procedure because it calls *myproc2*
- *\$ra* is set to the address of *sw* instruction
- When *myproc2* is called *\$ra* is set to the address of *jr* instruction
- Jump return from *myproc2* is OK


# Non-Leaf Procedure – Example (bad)

*main:*

```
la $s0, A
lw $a0, 0($s0)
jal myproc1
sw $v0, 4($s0)
j exit
```

*myproc1:* *#non-leaf*

```
addi $v0, $a0, 5
move $a0, $v0
jal myproc2
jr $ra
```



*myproc2:* *#leaf*

```
addi $v0, $a0, 10
jr $ra
```

- *myproc1* is a non-leaf procedure because it calls *myproc2*
- *\$ra* is set to the address of *sw* instruction
- When *myproc2* is called *\$ra* is set to the address of *jr* instruction
- Jump return from *myproc2* is OK
- Jump return from *myproc1* is an infinite loop!

# Non-Leaf Procedure – Example (good)

*main:*

```
la $s0, A
lw $a0, 0($s0)
jal myproc1
sw $v0, 4($s0)
j exit
```

Solution is to store the \$ra to the stack before calling myproc2

*myproc1:* # non-leaf

```
addi $sp, $sp, -4
sw $ra, 0($sp)
```

} Store \$ra to stack before calling myproc2

```
addi $v0, $a0, 5
move $a0, $v0
jal myproc2
```

```
lw $ra, 0($sp)
addi $sp, $sp, 4
```

} Restore \$ra to stack after calling myproc2

```
jr $ra
```

*myproc2:* #leaf

```
addi $v0, $a0, 10
jr $ra
```

# Recursive Procedure



re·cur·sive

/rəˈkərsɪv/

*adjective*

characterized by recurrence or repetition.

- **MATHEMATICS • LINGUISTICS**

relating to or involving the repeated application of a rule, definition, or procedure to successive results.

"this restriction ensures that the grammar is recursive"

- **COMPUTING**

relating to or involving a program or routine of which a part requires the application of the whole, so that its explicit interpretation requires in general many successive executions.

"a recursive subroutine"

- Recursive Procedure or nested call is an example of a non-leaf procedure
  - One of the steps of the procedure involves invoking the procedure itself

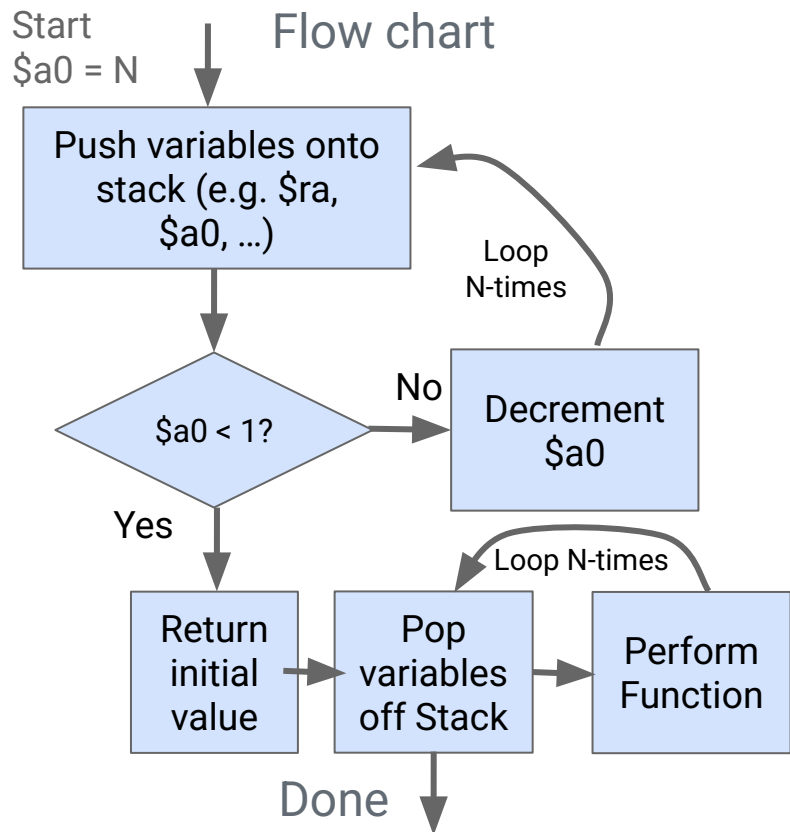


# Procedure Calling – Recursive

## Steps required

1. Place parameters in passed registers - **\$a0~\$a3**
2. Transfer control to procedure - **jal**
3. Acquire storage for procedure - **addi \$sp, \$sp, -4**
4. Perform procedure's operations
  - Save return address (**\$ra**)
  - Any arguments and temporaries needed after the call
  - Call the procedure itself
5. Place result in register for caller - **\$v0, \$v1**
6. Restore variables and return to place of call - **jr**
  - Restore from stack after the call
  - May return to place of call multiple times

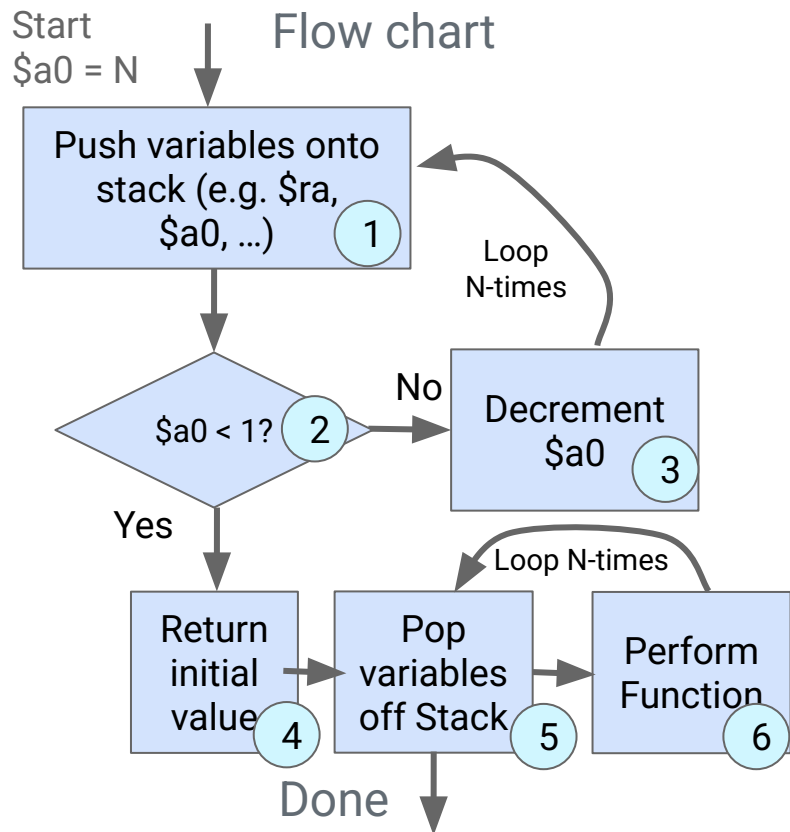
# Recursive MIPS program



Loop 1: Push variables onto the Stack

Loop 2: Pop values and perform functions

# Recursive MIPS program



NonLeafProc:

```
addi $sp, $sp, -8
sw    $ra, 4($sp)
sw    $a0, 0($sp)
```

```
slti $t0, $a0, 1
beq  $t0, $zero, Loop1
<Return initial value>
```

```
addi $sp, $sp, 8
jr    $ra
```

Loop1:

```
addi $a0, $a0, -1
jal  NonLeafProc
```

```
lw    $a0, 0($sp)
lw    $ra, 4($sp)
addi $sp, $sp, 8
<Perform Function>
jr    $ra
```

# Factorial C code – Recursive procedure

C code:

```
int fact (int n){  
    if (n < 1) return 1;  
    else return n * fact(n - 1);  
}
```

Python code:

```
def fact (n):  
    if n < 1:  
        return 1  
    else: return n * fact(n-1)
```

n = 3, execute fact (3)

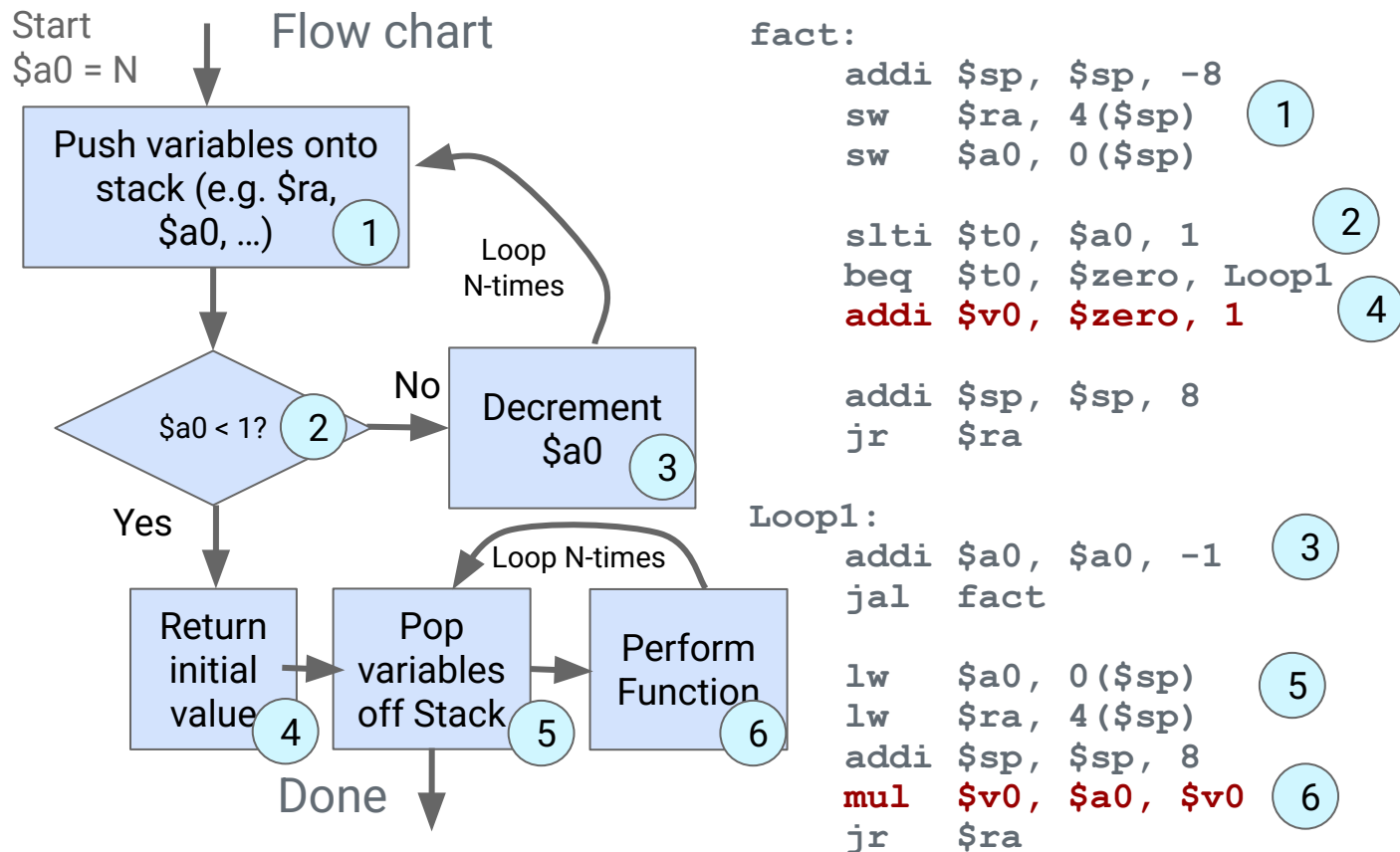
3 \* fact (2)

3 \* 2 \* fact (1)

3 \* 2 \* 1 \* fact (0)

3 \* 2 \* 1 \* 1 = 6

# Recursive MIPS program – Factorial Example



# Recursive Procedure Example

C code:

```
int fact (int n){  
    if (n < 1) return 1;  
    else return n *  
        fact(n - 1);  
}
```

Implemented in  
MIPS with 2  
procedures

MIPS code:

```
fact:                                     # Transfer control  
    addi $sp, $sp, -8                    # Acquire storage  
    sw   $ra, 4($sp)                     # save return address  
    sw   $a0, 0($sp)                     # save argument  
  
    slti $t0, $a0, 1                     # test for n < 1  
    beq  $t0, $zero, L1                  # branch if !(n < 1)  
    addi $v0, $zero, 1                   # if n < 1 result is 1  
  
    addi $sp, $sp, 8                     # pop 2 items from stack  
    jr   $ra                             # and return  
  
L1:  addi $a0, $a0, -1                    # Decrement n  
     jal  fact                           # Recursive call  
     lw   $a0, 0($sp)                     # restore original n  
     lw   $ra, 4($sp)                     # and return address  
     addi $sp, $sp, 8                     # pop 2 items from stack  
     mul  $v0, $a0, $v0                   # multiply to get result  
     jr   $ra                             # and return
```

# Recursive Procedure Example

C code:

```
int fact (int n){  
    if (n < 1) return 1;  
    else return n *  
        fact(n - 1);  
}
```

Step 1: Put argument  
n in \$a0

Expect result in \$v0

MIPS code:

```
fact:  
    addi $sp, $sp, -8      # Transfer control  
    sw   $ra, 4($sp)      # Acquire storage  
    sw   $a0, 0($sp)      # save return address  
                           # save argument  
  
    slti $t0, $a0, 1      # test for n < 1  
    beq  $t0, $zero, L1   # branch if !(n < 1)  
    addi $v0, $zero, 1    # if n < 1 result is 1  
  
    addi $sp, $sp, 8      # pop 2 items from stack  
    jr   $ra              # and return  
  
L1:  addi $a0, $a0, -1    # Decrement n  
     jal  fact            # Recursive call  
     lw   $a0, 0($sp)     # restore original n  
     lw   $ra, 4($sp)     # and return address  
     addi $sp, $sp, 8     # pop 2 items from stack  
     mul  $v0, $a0, $v0   # multiply to get result  
     jr   $ra            # and return
```

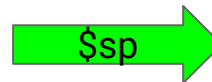
# Unrolling the Recursive Call

## Registers

\$v0	
\$a0	3
...	
\$t0	
...	
\$ra	Done

## Stack

	20
	1C
	18
	14
	10
	0C
	08
	04
	00



```
li $a0, 3           # n = 3
jal fact            # Calling fact
Done:...

fact:               # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw    $ra, 4($sp)  # save return address
    sw    $a0, 0($sp)  # save argument

    slti  $t0, $a0, 1  # test for n < 1
    beq   $t0, $zero, L1 # branch if !(n < 1)
    addi  $v0, $zero, 1 # if n < 1 result is 1

    addi  $sp, $sp, 8  # pop 2 items from stack
    jr    $ra          # and return

L1:  addi  $a0, $a0, -1 # Decrement n
     jal   fact        # Recursive call
     lw    $a0, 0($sp)  # restore original n
     lw    $ra, 4($sp)  # and return address
     addi  $sp, $sp, 8  # pop 2 items from stack
     mul   $v0, $a0, $v0 # multiply to get result
     jr    $ra          # and return
```



# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw   $ra, 4($sp)   # save return address
    sw   $a0, 0($sp)   # save argument

    slti $t0, $a0, 1   # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1 # if n < 1 result is 1

    addi $sp, $sp, 8   # pop 2 items from stack
    jr   $ra           # and return

L1: addi $a0, $a0, -1   # Decrement n
    jal  fact          # Recursive call
    lw   $a0, 0($sp)   # restore original n
    lw   $ra, 4($sp)   # and return address
    addi $sp, $sp, 8   # pop 2 items from stack
    mul  $v0, $a0, $v0 # multiply to get result
    jr   $ra           # and return
```

Registers

\$v0	
\$a0	3
...	
\$t0	
...	
\$ra	0

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
	14
	10
	0C
	08
	04
	00



\$sp

# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw   $ra, 4($sp)   # save return address
    sw   $a0, 0($sp)   # save argument

    slti $t0, $a0, 1   # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1 # if n < 1 result is 1

    addi $sp, $sp, 8   # pop 2 items from stack
    jr   $ra           # and return

L1: addi $a0, $a0, -1   # Decrement n
    jal  fact          # Recursive call
    lw   $a0, 0($sp)   # restore original n
    lw   $ra, 4($sp)   # and return address
    addi $sp, $sp, 8   # pop 2 items from stack
    mul  $v0, $a0, $v0 # multiply to get result
    jr   $ra           # and return
```

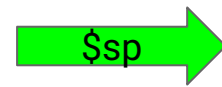
Registers

\$v0	
\$a0	3
...	
\$t0	0
...	
\$ra	0

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
	14
	10
	0C
	08
	04
	00

\$sp



# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw   $ra, 4($sp)   # save return address
    sw   $a0, 0($sp)   # save argument

    slti $t0, $a0, 1   # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1 # if n < 1 result is 1

    addi $sp, $sp, 8   # pop 2 items from stack
    jr   $ra           # and return

L1:  addi $a0, $a0, -1 # Decrement n
     jal  fact         # Recursive call
     lw   $a0, 0($sp)  # restore original n
     lw   $ra, 4($sp)  # and return address
     addi $sp, $sp, 8  # pop 2 items from stack
     mul  $v0, $a0, $v0 # multiply to get result
     jr   $ra          # and return
```

Registers

\$v0	
\$a0	2
...	
\$t0	0
...	
\$ra	L1+8

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
	14
	10
	0C
	08
	04
	00

\$sp

# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
addi $sp, $sp, -8  # Acquire storage
sw    $ra, 4($sp)  # save return address
sw    $a0, 0($sp)  # save argument

    slti $t0, $a0, 1 # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1 # if n < 1 result is 1

    addi $sp, $sp, 8 # pop 2 items from stack
    jr   $ra        # and return

L1: addi $a0, $a0, -1 # Decrement n
    jal  fact        # Recursive call
    lw   $a0, 0($sp) # restore original n
    lw   $ra, 4($sp) # and return address
    addi $sp, $sp, 8 # pop 2 items from stack
    mul  $v0, $a0, $v0 # multiply to get result
    jr   $ra        # and return
```

Registers

\$v0	
\$a0	2
...	
\$t0	0
...	
\$ra	L1+8

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
<b>\$ra = L1+8</b>	<b>14</b>
<b>\$a0 = 2</b>	<b>10</b>
	0C
	08
	04
	00



\$sp

# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw    $ra, 4($sp)  # save return address
    sw    $a0, 0($sp)  # save argument

    slti $t0, $a0, 1   # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1 # if n < 1 result is 1

    addi $sp, $sp, 8   # pop 2 items from stack
    jr    $ra          # and return

L1: addi $a0, $a0, -1   # Decrement n
    jal  fact          # Recursive call
    lw   $a0, 0($sp)   # restore original n
    lw   $ra, 4($sp)   # and return address
    addi $sp, $sp, 8   # pop 2 items from stack
    mul  $v0, $a0, $v0 # multiply to get result
    jr   $ra          # and return
```

Registers

\$v0	
\$a0	2
...	
\$t0	0
...	
\$ra	L1+8

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
\$ra = L1+8	14
\$a0 = 2	10
	0C
	08
	04
	00

\$sp



# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw    $ra, 4($sp)  # save return address
    sw    $a0, 0($sp)  # save argument

    slti  $t0, $a0, 1  # test for n < 1
    beq   $t0, $zero, L1 # branch if !(n < 1)
    addi  $v0, $zero, 1 # if n < 1 result is 1

    addi  $sp, $sp, 8  # pop 2 items from stack
    jr    $ra          # and return

L1:  addi  $a0, $a0, -1 # Decrement n
     jal   fact         # Recursive call
     lw    $a0, 0($sp)  # restore original n
     lw    $ra, 4($sp)  # and return address
     addi  $sp, $sp, 8  # pop 2 items from stack
     mul   $v0, $a0, $v0 # multiply to get result
     jr    $ra          # and return
```

Registers

\$v0	
\$a0	1
...	
\$t0	0
...	
\$ra	L1+8

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
\$ra = L1+8	14
\$a0 = 2	10
	0C
	08
	04
	00

\$sp

# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
addi $sp, $sp, -8  # Acquire storage
sw    $ra, 4($sp)  # save return address
sw    $a0, 0($sp)  # save argument

    slti $t0, $a0, 1  # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1 # if n < 1 result is 1

    addi $sp, $sp, 8  # pop 2 items from stack
    jr   $ra          # and return

L1: addi $a0, $a0, -1  # Decrement n
    jal  fact         # Recursive call
    lw   $a0, 0($sp)  # restore original n
    lw   $ra, 4($sp)  # and return address
    addi $sp, $sp, 8  # pop 2 items from stack
    mul  $v0, $a0, $v0 # multiply to get result
    jr   $ra          # and return
```

Registers

\$v0	
\$a0	1
...	
\$t0	0
...	
\$ra	L1+8

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
\$ra = L1+8	14
\$a0 = 2	10
\$ra = L1+8	0C
\$a0 = 1	08
	04
	00



\$sp

# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw    $ra, 4($sp)  # save return address
    sw    $a0, 0($sp)  # save argument

    slti $t0, $a0, 1   # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1 # if n < 1 result is 1

    addi $sp, $sp, 8   # pop 2 items from stack
    jr   $ra           # and return

L1: addi $a0, $a0, -1  # Decrement n
    jal  fact         # Recursive call
    lw   $a0, 0($sp)  # restore original n
    lw   $ra, 4($sp)  # and return address
    addi $sp, $sp, 8  # pop 2 items from stack
    mul  $v0, $a0, $v0 # multiply to get result
    jr   $ra         # and return
```

Registers

\$v0	
\$a0	1
...	
\$t0	0
...	
\$ra	L1+8

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
\$ra = L1+8	14
\$a0 = 2	10
\$ra = L1+8	0C
\$a0 = 1	08
	04
	00

\$sp





# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw    $ra, 4($sp)  # save return address
    sw    $a0, 0($sp)  # save argument

    slti  $t0, $a0, 1  # test for n < 1
    beq   $t0, $zero, L1 # branch if !(n < 1)
    addi  $v0, $zero, 1 # if n < 1 result is 1

    addi  $sp, $sp, 8   # pop 2 items from stack
    jr    $ra          # and return

L1:  addi $a0, $a0, -1  # Decrement n
     jal  fact         # Recursive call
     lw   $a0, 0($sp)  # restore original n
     lw   $ra, 4($sp)  # and return address
     addi $sp, $sp, 8   # pop 2 items from stack
     mul  $v0, $a0, $v0 # multiply to get result
     jr   $ra          # and return
```

Registers

\$v0	
\$a0	0
...	
\$t0	0
...	
\$ra	L1+8

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
\$ra = L1+8	14
\$a0 = 2	10
\$ra = L1+8	0C
\$a0 = 1	08
	04
	00

\$sp

# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
addi $sp, $sp, -8  # Acquire storage
sw    $ra, 4($sp)  # save return address
sw    $a0, 0($sp)  # save argument

    slti $t0, $a0, 1  # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1 # if n < 1 result is 1

    addi $sp, $sp, 8  # pop 2 items from stack
    jr   $ra          # and return

L1: addi $a0, $a0, -1  # Decrement n
    jal  fact          # Recursive call
    lw   $a0, 0($sp)   # restore original n
    lw   $ra, 4($sp)   # and return address
    addi $sp, $sp, 8   # pop 2 items from stack
    mul  $v0, $a0, $v0 # multiply to get result
    jr   $ra          # and return
```

## Registers

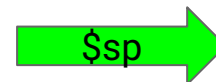
\$v0	
\$a0	0
...	
\$t0	0
...	
\$ra	L1+8

## Stack

	20
\$ra = 0	1C
\$a0 = 3	18
\$ra = L1+8	14
\$a0 = 2	10
\$ra = L1+8	0C
\$a0 = 1	08
\$ra = L1+8	04
\$a0 = 0	00

Pushed onto  
the stack all of  
the data needed  
for the recursive  
call. Time to  
pop...

\$sp



# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw   $ra, 4($sp)   # save return address
    sw   $a0, 0($sp)   # save argument

    slti $t0, $a0, 1   # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1 # if n < 1 result is 1

    addi $sp, $sp, 8   # pop 2 items from stack
    jr   $ra           # and return

L1: addi $a0, $a0, -1   # Decrement n
    jal  fact          # Recursive call
    lw   $a0, 0($sp)   # restore original n
    lw   $ra, 4($sp)   # and return address
    addi $sp, $sp, 8   # pop 2 items from stack
    mul  $v0, $a0, $v0 # multiply to get result
    jr   $ra           # and return
```

Registers

\$v0	1
\$a0	0
...	
\$t0	1
...	
\$ra	L1+8

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
\$ra = L1+8	14
\$a0 = 2	10
\$ra = L1+8	0C
\$a0 = 1	08
<b>\$ra = L1+8</b>	<b>04</b>
<b>\$a0 = 0</b>	<b>00</b>

\$sp

# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw   $ra, 4($sp)   # save return address
    sw   $a0, 0($sp)   # save argument

    slti $t0, $a0, 1   # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1 # if n < 1 result is 1
```

```
addi $sp, $sp, 8    # pop 2 items from stack
jr   $ra             # and return
```

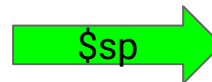
```
L1: addi $a0, $a0, -1 # Decrement n
    jal  fact         # Recursive call
    lw   $a0, 0($sp)  # restore original n
    lw   $ra, 4($sp)  # and return address
    addi $sp, $sp, 8  # pop 2 items from stack
    mul  $v0, $a0, $v0 # multiply to get result
    jr   $ra          # and return
```

Registers

\$v0	1
\$a0	0
...	
\$t0	1
...	
\$ra	L1+8

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
\$ra = L1+8	14
\$a0 = 2	10
\$ra = L1+8	0C
\$a0 = 1	08
\$ra = L1+8	04
\$a0 = 0	00



# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw   $ra, 4($sp)   # save return address
    sw   $a0, 0($sp)   # save argument

    slti $t0, $a0, 1   # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1  # if n < 1 result is 1

    addi $sp, $sp, 8    # pop 2 items from stack
    jr   $ra           # and return

L1: addi $a0, $a0, -1   # Decrement n
    jal fact           # Recursive call
    lw   $a0, 0($sp)   # restore original n
    lw   $ra, 4($sp)   # and return address
    addi $sp, $sp, 8    # pop 2 items from stack
    mul  $v0, $a0, $v0  # multiply to get result
    jr   $ra           # and return
```

Registers

\$v0	1=1*1
\$a0	1
...	
\$t0	1
...	
\$ra	L1+8

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
\$ra = L1+8	14
\$a0 = 2	10
\$ra = L1+8	0C
\$a0 = 1	08
\$ra = L1+8	04
\$a0 = 0	00

\$sp



# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw   $ra, 4($sp)   # save return address
    sw   $a0, 0($sp)   # save argument

    slti $t0, $a0, 1   # test for n < 1
    beq  $t0, $zero, L1 # branch if !(n < 1)
    addi $v0, $zero, 1 # if n < 1 result is 1

    addi $sp, $sp, 8   # pop 2 items from stack
    jr   $ra           # and return

L1: addi $a0, $a0, -1   # Decrement n
    jal fact           # Recursive call
    lw   $a0, 0($sp)   # restore original n
    lw   $ra, 4($sp)   # and return address
    addi $sp, $sp, 8   # pop 2 items from stack
    mul  $v0, $a0, $v0 # multiply to get result
    jr   $ra           # and return
```

Registers

\$v0	2=2*1
\$a0	2
...	
\$t0	1
...	
\$ra	L1+8

Stack

	20
\$ra = 0	1C
\$a0 = 3	18
\$ra = L1+8	14
\$a0 = 2	10
\$ra = L1+8	0C
\$a0 = 1	08
\$ra = L1+8	04
\$a0 = 0	00

\$sp



# Unrolling the Recursive Call

```
li $a0, 3          # n = 3
jal fact           # Calling fact
Done:...

fact:              # Transfer control
    addi $sp, $sp, -8  # Acquire storage
    sw    $ra, 4($sp)  # save return address
    sw    $a0, 0($sp)  # save argument

    slti  $t0, $a0, 1  # test for n < 1
    beq   $t0, $zero, L1 # branch if !(n < 1)
    addi  $v0, $zero, 1 # if n < 1 result is 1

    addi  $sp, $sp, 8  # pop 2 items from stack
    jr    $ra          # and return

L1:  addi  $a0, $a0, -1 # Decrement n
     jal   fact         # Recursive call
     lw    $a0, 0($sp)  # restore original n
     lw    $ra, 4($sp)  # and return address
     addi  $sp, $sp, 8  # pop 2 items from stack
     mul   $v0, $a0, $v0 # multiply to get result
     jr    $ra          # and return
```

$\$v0 = 6$   
 $\$ra, \$sp$  are  
restored

## Registers

$\$v0$	$6=3*2$
$\$a0$	3
...	
$\$t0$	1
...	
$\$ra$	Done

## Stack

	20
$\$ra = 0$	1C
$\$a0 = 3$	18
$\$ra = L1+8$	14
$\$a0 = 2$	10
$\$ra = L1+8$	0C
$\$a0 = 1$	08
$\$ra = L1+8$	04
$\$a0 = 0$	00

$\$sp$

# Leaf Procedure example – My turn

Execute the following procedure:

*MyProc:*

```
addi $sp, $sp, -12
sw $s0, 8($sp)
sw $s1, 4($sp)
sw $s2, 0($sp)
```

```
and $s0, $a0, $a1
srl $s1, $a2, 3
slt $s2, $s1, $a2
```

```
add $v0, $s0, $0
add $v1, $s1, $0
```

```
lw $s2, 0($sp)
lw $s1, 4($sp)
lw $s0, 8($sp)
addi $sp, $sp, 12
jr $ra
```

What is the final  
value of \$v1?

What is the final  
value of \$s0?

Registers (Initial)

Name	Number	Value
\$v0	2	1
\$v1	3	15
\$a0	4	5
\$a1	5	9
\$a2	6	15
\$a3	7	7
\$t0	8	63
\$t1	9	8
\$t2	10	10
\$t3	11	5
\$t4	12	15
\$t5	13	15
\$t6	14	0
\$t7	15	0
\$s0	16	8
\$s1	17	60
\$s2	18	0



# Leaf Procedure example – My turn

Execute the following procedure:

*MyProc:*

```
addi $sp, $sp, -12    Allocate storage on stack
sw $s0, 8($sp)        Store variables on stack
sw $s1, 4($sp)
sw $s2, 0($sp)
```

```
and $s0, $a0, $a1
srl $s1, $a2, 3
slt $s2, $s1, $a2
```

```
add $v0, $s0, $0
add $v1, $s1, $0
```

```
lw $s2, 0($sp)
lw $s1, 4($sp)
lw $s0, 8($sp)
addi $sp, $sp, 12
```

Stored onto stack (stack not shown)

Registers

Name	Number	Value
\$v0	2	1
\$v1	3	15
\$a0	4	5
\$a1	5	9
\$a2	6	15
\$a3	7	7
\$t0	8	63
\$t1	9	8
\$t2	10	10
\$t3	11	5
\$t4	12	15
\$t5	13	15
\$t6	14	0
\$t7	15	0
\$s0	16	8
\$s1	17	60
\$s2	18	0

# Leaf Procedure example – My turn

Execute the following procedure:

*MyProc:*

`addi $sp, $sp, -12`      Allocate storage on stack

`sw $s0, 8($sp)`      Store variables on stack

`sw $s1, 4($sp)`

`sw $s2, 0($sp)`

`and $s0, $a0, $a1`

`srl $s1, $a2, 3`

`slt $s2, $s1, $a2`

Execute Program:

$\$s0 = \$a0 \& \$a1 = 1$

$\$s1 = \$a2 \gg 3 = 1$

$\$s2 = 1 \text{ if } \$s1 < \$a2 = 0$

`add $v0, $s0, $0`

`add $v1, $s1, $0`

`lw $s2, 0($sp)`

`lw $s1, 4($sp)`

`lw $s0, 8($sp)`

`addi $sp, $sp, 12`

`jr $ra`

Registers

Name	Number	Value
\$v0	2	1
\$v1	3	15
\$a0	4	5
\$a1	5	9
\$a2	6	15
\$a3	7	7
\$t0	8	63
\$t1	9	8
\$t2	10	10
\$t3	11	5
\$t4	12	15
\$t5	13	15
\$t6	14	0
\$t7	15	0
\$s0	16	1
\$s1	17	1
\$s2	18	1

# Leaf Procedure example – My turn

Execute the following procedure:

*MyProc:*

```
addi $sp, $sp, -12    Allocate storage on stack
sw $s0, 8($sp)        Store variables on stack
sw $s1, 4($sp)
sw $s2, 0($sp)

and $s0, $a0, $a1      Execute Program:
                        $s0 = $a0 & $a1 = 1
srl $s1, $a2, 3         $s1 = $a2 >> 3 = 1
slt $s2, $s1, $a2       $s2 = 1 if $s1 < $a2 = 0

add $v0, $s0, $0        Return results in $v0 and $v1
add $v1, $s1, $0

lw $s2, 0($sp)
lw $s1, 4($sp)
lw $s0, 8($sp)
addi $sp, $sp, 12
jr $ra
```

Registers

Name	Number	Value
\$v0	2	1
\$v1	3	1
\$a0	4	5
\$a1	5	9
\$a2	6	15
\$a3	7	7
\$t0	8	63
\$t1	9	8
\$t2	10	10
\$t3	11	5
\$t4	12	15
\$t5	13	15
\$t6	14	0
\$t7	15	0
\$s0	16	1
\$s1	17	1
\$s2	18	1

# Leaf Procedure example – My turn

Execute the following procedure:

*MyProc:*

<code>addi \$sp, \$sp, -12</code>	Allocate storage on stack
<code>sw \$s0, 8(\$sp)</code>	Store variables on stack
<code>sw \$s1, 4(\$sp)</code>	
<code>sw \$s2, 0(\$sp)</code>	
<code>and \$s0, \$a0, \$a1</code>	Execute Program:
<code>srl \$s1, \$a2, 3</code>	<code>\$s0 = \$a0 &amp; \$a1 = 1</code>
<code>slt \$s2, \$s1, \$a2</code>	<code>\$s1 = \$a2 &gt;&gt; 3 = 1</code>
	<code>\$s2 = 1 if \$s1 &lt; \$a2 = 0</code>
<code>add \$v0, \$s0, \$0</code>	Return results in \$v0 and \$v1
<code>add \$v1, \$s1, \$0</code>	
<code>lw \$s2, 0(\$sp)</code>	Restore variables
<code>lw \$s1, 4(\$sp)</code>	Free Storage
<code>lw \$s0, 8(\$sp)</code>	
<code>addi \$sp, \$sp, 12</code>	Return to Caller
<code>jr \$ra</code>	

## Registers

Name	Number	Value
\$v0	2	1
\$v1	3	1
\$a0	4	5
\$a1	5	9
\$a2	6	15
\$a3	7	7
\$t0	8	63
\$t1	9	8
\$t2	10	10
\$t3	11	5
\$t4	12	15
\$t5	13	15
\$t6	14	0
\$t7	15	0
\$t8	16	8
\$t9	17	60
\$s2	18	0

Restored from stack (stack not shown)

# Leaf Procedure example – My turn

Execute the following procedure:

*MyProc:*

```
addi $sp, $sp, -12
sw $s0, 8($sp)
sw $s1, 4($sp)
sw $s2, 0($sp)
```

```
and $s0, $a0, $a1
srl $s1, $a2, 3
slt $s2, $s1, $a2
```

```
add $v0, $s0, $0
add $v1, $s1, $0
```

```
lw $s2, 0($sp)
lw $s1, 4($sp)
lw $s0, 8($sp)
addi $sp, $sp, 12
jr $ra
```

What is the final value  
of \$v1? 1

What is the final value  
of \$s0? 8

(all saved registers  
must be restored to  
original value before  
returning MyProc)

Registers (final)

Name	Number	Value
\$v0	2	1
\$v1	3	1
\$a0	4	5
\$a1	5	9
\$a2	6	15
\$a3	7	7
\$t0	8	63
\$t1	9	8
\$t2	10	10
\$t3	11	5
\$t4	12	15
\$t5	13	15
\$t6	14	0
\$t7	15	0
\$s0	16	8
\$s1	17	60
\$s2	18	0

# Procedure-Your Turn

Data Segment								
Address	Value (+0)	Value (+4)	Value (+8)	Value (+c)	Value (+10)	Value (+14)	Value (+18)	Value (+1c)
0x10010000	5	10	8	2	3	10	5	8
0x10010020	2	15	34	56	23	68	75	0
0x10010040	0	0	0	0	0	0	0	0

← → 0x10010000 (.data) ☒ Hexadecimal Addresses ☐ Hexadecimal Values ☐ ASCII

Execute the following procedure called SProc:

SProc:

```
addi $sp, $sp, -8
sw $s0, 4($sp)
sw $s1, 0($sp)
```

```
la $t0, NumArray
sll $t1, $a0, 2
sll $t2, $a1, 2
add $t1, $t1, $t0
add $t2, $t2, $t0
```

```
lw $s0, 0($t1)
lw $s1, 0($t2)
sw $s1, 0($t1)
sw $s0, 0($t2)
```

```
lw $s1, 0($sp)
lw $s0, 4($sp)
addi $sp, $sp, 8
jr $ra
```

NumArray base address is 0x10010008  
argument i and j are stored in \$a0, \$a1

What is the final value of NumArray[9]?  
What is the final value of \$s1?  
What is the final value of \$t1?

Registers		
Name	Number	Value
\$zero	0	0
\$at	1	0
\$v0	2	11
\$v1	3	15
\$a0	4	1
\$a1	5	9
\$a2	6	15
\$a3	7	7
\$t0	8	-7
\$t1	9	0
\$t2	10	12
\$t3	11	0
\$t4	12	0
\$t5	13	0
\$t6	14	0
\$t7	15	0
\$s0	16	5
\$s1	17	75
\$s2	18	0
\$s3	19	0
\$s4	20	0
\$s5	21	0
\$s6	22	0
\$s7	23	0
\$t8	24	0
\$t9	25	0
\$k0	26	0
\$k1	27	0
\$gp	28	268468224
\$sp	29	0x7ffefffc
\$fp	30	0
\$ra	31	0

# Summary

- Procedures:
  - Caller, Callee, Program Counter
  - Stack, Stack pointer, Push/Pop
  - Leaf Procedure Example
  - Non-leaf Procedure, Recursive Procedure Example

Now you have all of the skills to do the Term Project!

Next

# CPU Performance

