CS 314, Assignment 1 - Report

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1 Abstract

The objective of this task is to simulate breadth-first search, depth-first search, and DFID in the state space. The state-space consists of an $(m \times n)$ grid. The start state is (0,0). The goal state is the position of [*] in the grid. The Pacman is allowed to move UP, DOWN, LEFT and RIGHT (except for boundary). A comparison of the path length and the number of states explored between the different search methods and, also between the orders in which neighbours are added, are performed.

2 Pseudo Codes

In the following subsections, pseudo codes for important functions in the code are explained.

2.1 movGen(state)

The function takes a state as input and returns a set of states that are reachable from the input state in one step. [Algorithm 1]

Algorithm 1 moveGen(state)

- 1: procedure MOVEGEN(state)
- 2: $nextStates \leftarrow ()$

- ▷ initialize nextStates to empty set
- 3: **for** neighbour n of state in order(DOWN,UP,RIGHT,LEFT) **do**
- 4: **if** n is not boundary **then**
- 5: nextStates.append(n)
- 6: **return** nextStates

▷ nextStates are required moves generated

2.2 GoalTest(state)

This function returns True if the input state is goal and False otherwise. [Algorithm 2]

Algorithm 2 goalTest(state)

```
1: procedure GOALTEST(state)
2: if state.value == '*' then
3: return true
4: return false ▷ state is not goal
```

2.3 BFS

This function is the implementation of BFS (Breadth First Search) Algorithm. [Algorithm 3]

Algorithm 3 final_BFS()

```
1: procedure FINAL_BFS
2:
      if Source \&\& goalCell == VALID then
3:
          pushSource(queue)
      while !queue do
4:
          statesExplored++;
5:
          extractCell(queue);
6:
          if currentCell == qoalCell then
                                                                     ▶ Using Coordinates
7:
             bool b = goalTest(goalCell);
8:
9:
             if b == 1 then
                 tracePath(source, currentCell);
10:
                 return distance
11:
             pop(extractCell)
12:
          movGen(state)
13:
```

2.4 DFS

This function is the implementation of DFS (Depth First Search) Algorithm. [Algorithm 4]

2.5 DFID

This function is the implementation of DFID (Iterative Deepening Depth-First Search) Algorithm. [Algorithm 5]

With the code developed and some random customized Mazes were generated Delorie, it could be confirmed that our code was working fine. With Cell Width = 3 and Cell Height = 2 configuration, the values of path length and the number of states explored for each of the algorithm for each maze was noted at three different orders of visiting adjacent cells viz Down>Up>Right>Left, Down>Up>Left>Right, and Right>Down>Up>Left. An analysis of the results obtained is furnished below.

Algorithm 4 final_DFS()

```
1: procedure FINAL_DFS
      if Source \&\& goalCell == VALID then
         pushSource(queue)
3:
      while !stack do
4:
         statesExplored++;
5:
         extractCell(stack);
6:
         if currentCell == goalCell then
                                                                 7:
            bool b = goalTest(goalCell);
8:
            if b == 1 then
9:
                tracePath(source, currentCell);
10:
11:
                return distance
12:
            pop(extractCell)
         movGen(state)
13:
```

3 Data - Variation in Mazes

With variation in sizes of mazes, the corresponding data has been tabulated to further analyze different algorithms across various sizes of mazes.

The following **Table 1** shows the statistics related to mazes of different sizes created and corresponding number of states explored and path length across different algorithms using the Order: Down > Up > Right > Left.

The following **Table 2** shows the statistics related to mazes of different sizes created and corresponding number of states explored and path length across different algorithms using the Order: Right > Left > Down > Up.

4 Graphical Interpretation & Inference

4.1 Graphs

With the help of above data given in the tables, the following graphs are compared to analyse various components and efficiency of algorithms with varying sizes of mazes.

In *Graph 1*, No. of States Explored vs. Size of Maze is plotted for the Order: Down > Up > Right > Left.

In *Graph 2*, No. of States Explored vs. Size of Maze is plotted for the Order: Right > Left > Down > Up.

In *Graph 3*, Path Length vs. Size of Maze is plotted for the Order: Down > Up > Right > Left.

Algorithm 5 final_DFID()

```
1: procedure FINAL_DFID
      if Source \&\& goalCell == VALID then
2:
          if currentDepth == 0 then
3:
             return -1:
4:
          distancefromSource(currentCell)
      if currentCell == goalCell then
                                                                     ▶ Using Coordinates
5:
          bool b = goalTest(goalCell);
6:
7:
          if b == 1 then
             tracePath(source, currentCell);
8:
             return distance
9:
          statesExplored++;
10:
      movGen(state)
11:
      currentDepth = 1
12:
      while TRUE do
13:
          final_DFIS();
14:
          if currentCell == goalCell then
15:
             break
16:
          currentDepth++;
17:
```

In *Graph* 4, Path Length vs. Size of Maze is plotted for the Order: Right > Left > Down > Up.

4.2 Inference

- 1. BFS always yields paths of shorter (or equal to) lengths than those yielded by DFS.
- 2. BFS and DFID yield paths of same length.
- 3. Number of states explored by BFS and DFS are generally different but do not vary appreciably among themselves.
- 4. Number of states explored by DFID is quite large when compared to those explored by BFS and DFS due to the fact that finding optimal path in DFID with cycles in the graph may require visiting a node multiple times.
- 5. Order of visiting adjacent cells does influence the result obtained, as is evident from the plots of cumulative analysis seen above, though the difference in the path lengths obtained and number of states explored in each case is not appreciably large.

	Order: Down > Up > Right > Left				
Algorithm	, ,				
	Horizontal Cells	Vertical Cells	No. of States Explored	Path Length	
BFS	2	2	13	10	
DFS	2	2	14	10	
DFID	2	2	62	10	
BFS	4	4	43	28	
DFS	4	4	35	32	
DFID	4	4	790	28	
BFS	6	6	117	38	
DFS	6	6	71	46	
DFID	6	6	5183	38	
BFS	8	8	137	88	
DFS	8	8	194	102	
DFID	8	8	7612	88	
BFS	10	10	347	126	
DFS	10	10	257	148	
DFID	10	10	50437	126	

Table 1: Order: Down > Up > Right > Left

5 Conclusion

The results of the dependence of the path length and number of states explored, as seen in the previous section, are summarized in the table below. For small inputs in DFID, we observe that the increase in the number of explored states is due to the small branching factor and high constant attached with the time complexity.

	${ m Order:\ Right> Left> Down> Up} \ { m Cell\ Statistics}\ ({ m Cell\ Width}=3,\ { m Cell\ Height}=2)$				
Algorithm					
	Horizontal cells	Vertical cells	No. of States Explored	Path Length	
BFS	2	2	13	10	
DFS	2	2	13	10	
DFID	2	2	63	10	
BFS	4	4	41	28	
DFS	4	4	48	28	
DFID	4	4	1150	28	
BFS	6	6	117	38	
DFS	6	6	93	38	
DFID	6	6	5648	38	
BFS	8	8	137	88	
DFS	8	8	184	88	
DFID	8	8	20689	88	
BFS	10	10	345	126	
DFS	10	10	247	126	
DFID	10	10	174250	126	

Table 2: Order: Right > Left > Down > Up

Algorithm	Dependence on order of Neighbors Added			
Aigoritiiii	No. of States Explored	Path Length		
BFS	True	False		
DFS	True	True		
DFID	True	False		

Table 3: Dependence of order of Neighbours added

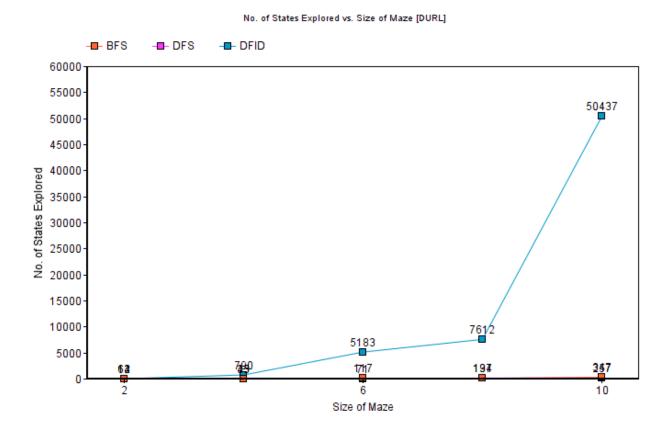


Figure 1: No. of States Explored vs. Size of Maze is plotted for the Order: Down > Up > Right > Left

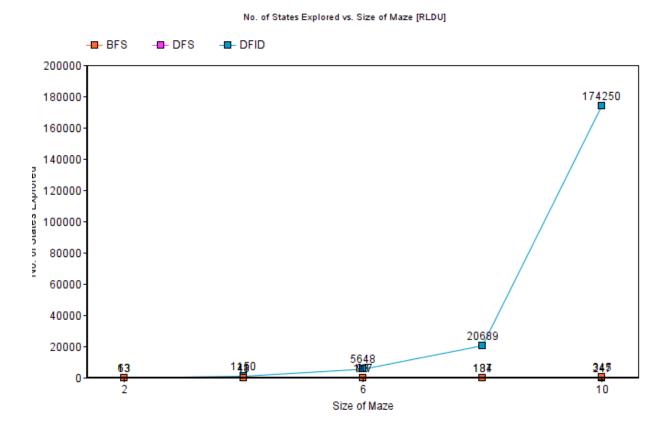


Figure 2: No. of States Explored vs. Size of Maze is plotted for the Order: Right > Left > Down > Up

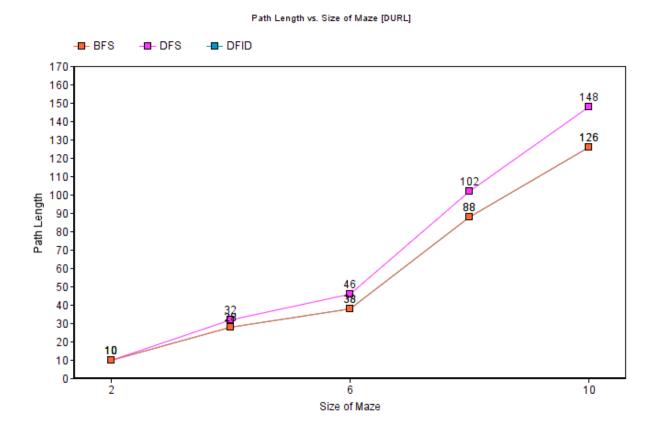


Figure 3: Path Length vs. Size of Maze is plotted for the Order: Down > Up > Right > Left

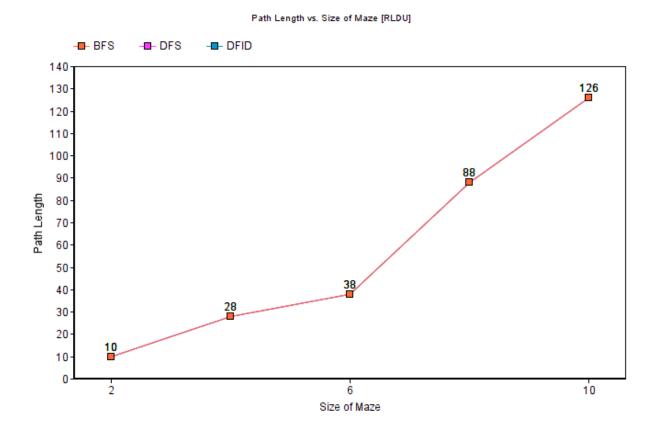


Figure 4: Path Length vs. Size of Maze is plotted for the Order: Right > Left > Down > Up