

Algorithm Questions

BIG O

SEARCHING

Risk and Pricing Solutions

Big O

Properties

What is the asymptotic running time of the following?

```
public static int SearchRecursive(int[] arr, int searchKey)
{
    if (arr == null) throw new ArgumentNullException();

    return SearchRecursive(arr, 0, arr.Length - 1, searchKey);
}

private static int SearchRecursive(int[] arr, int lo, int hi, int searchKey)
{
    // The search key is not in the array. Return the complement of the index
    // at which it should be inserted.
    if (lo > hi) return ~lo;

    int median = lo + (hi - lo) / 2;
    int comparisonResult = arr[median].CompareTo(searchKey); ;

    // a direct hit
    if (comparisonResult == 0) return median;

    return comparisonResult < 0
        ? SearchRecursive(arr, median + 1, hi, searchKey)
        : SearchRecursive(arr, lo, median - 1, searchKey);
}
```

The running time is $O(\log n)$

Why do we not care about the base of the logarithm?

$$\text{Because } \log_a x = \frac{\log_b x}{\log_b a} = \log_b x \frac{1}{\log_b a} =$$

So $\log_a x$ and $\log_b x$ differ by a constant factor and we don't worry about constant factors in asymptotic notation

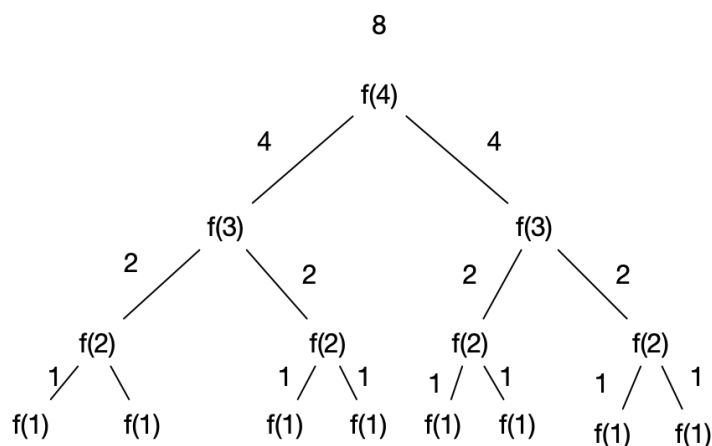
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What is the running time of the following function and what does it do?

```
public static int Function(int n)
{
    if (n == 1)
        return 1;

    return Function(n - 1) + Function(n - 1);
}
```

Look at the call graph of the specific case of Function(4). We get the total run time by summing the number of calls at each level. $1+2+4+8$. As we move down from one level to the other each level has double the number of calls as the level before. In our case we have $1 + 2 + 4 + 8 = 2^0 + 2^1 + 2^2 + 2^3 = 2^4 - 1$



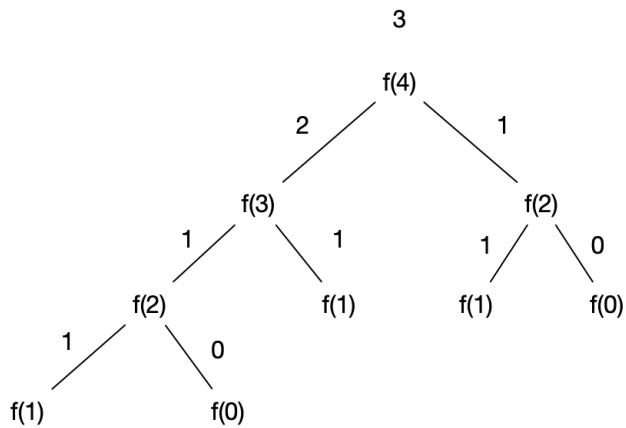
The running time is $O(2^{n-1}) = \frac{O(2^n)}{2} = O(2^n)$ Why? Because $2^{n-1} = \frac{1}{2^{n-1}2}$ and we can ignore constant factors.

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What is the asymptotic running time of the following?

```
public static int FibonacciRecursive(int n)
{
    if (n == 0)
        return 0;
    if (n == 1)
        return 1;

    return FibonacciRecursive(n - 1) + FibonacciRecursive(n - 2);
}
```



The call graph for fibonacci is very similar to the previous question so $O(2^n)$ is a correct upper bound. It is possible to prove a tighter upper bound as $O(\phi^n)$ where $\phi = (1 + \sqrt{5})/2$

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What is the asymptotic running time of the following?

```
public static int FibonacciRecursive(int n)
{
    if (n == 0)
        return 0;
    if (n == 1)
        return 1;

    return FibonacciRecursive(n - 1) + FibonacciRecursive(n - 2);
}
```

What is the asymptotic running time of the following?

```
public static int Function2(int n)
{
    int res = 0;
    for (int i = 0; i < n; i++)
        res += i;

    for (int i = 0; i < n; i++)
        res += i;

    return res;
}
```

The running time is $O(n)$ We ignore the fact it is $2n$ as we drop the constant factors

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What is the asymptotic running time of the following?

```
public static int PairCount(int[] a)
{
    int count = 0;
    for (int i = 0; i < a.Length; i++)
    {
        for (int j = i + 1; j < a.Length; j++)
        {
            if (a[i] == a[j]) count++;
        }
    }

    return count;
}
```

The running time is $O(n^2)$. Remember that the sum of the first n integers is given by

$$s = 1 + 2 + \dots + (n - 2) + (n - 1) + n$$

We can write $2s$ as

$$1 + 2 + \dots + (n - 2) + (n - 1) + n$$

$$n + (n - 1) + (n - 2) + \dots + 2 + 1$$

$$2s = n(n + 1) \therefore s = \frac{n(n + 1)}{2}$$

So in our we are replacing n with $n-1$

$$s = \frac{(n - 1)((n - 1) + 1)}{2} = \frac{(n - 1)n}{2}$$

In our asymptotic notation we call this $O(n^2)$ by dropping the lower order terms.

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What is the asymptotic running time of the following?

```
public static void PrintPairs(int[] a, int[] b)
{
    for (int i = 0; i < a.Length; i++)
    {
        for (int j = 0; j < b.Length; j++)
        {
            Console.WriteLine($"{a[i]}, {b[j]}");
        }
    }
}
```

$O(xy)$ where x is the number of elements in a and y is the number of elements in b

What is the asymptotic running time of the following?

```
public static void PrintPairsManyTimes(int[] a, int[] b)
{
    for (int i = 0; i < a.Length; i++)
    {
        for (int j = 0; j < b.Length; j++)
        {
            for (int k = 0; k < 10; k++)
            {
                Console.WriteLine($"{a[i]}, {b[j]}");
            }
        }
    }
}
```

$O(10xy)$ where x is the number of elements in a and y is the number of elements in b which is of course just $O(xy)$

What is the asymptotic running time of the following?

```
public void ReverseArray(int[] a)
{
    for (int i = 0; i < a.Length/2; i++)
    {
        int temp = a[i];
        a[i] = a[a.Length-1-i];
        a[a.Length-1-i] = temp;
    }
}
```

$O(n)$ We ignore the constant factor of $\frac{n}{2}$

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What is the asymptotic running time of sorting each string in an array and then sorting the array itself?

If we let the number of strings in the array be n and the length of each string be l then sorting each string takes $O(l \log l)$. We have to do this n time so we get $O(n \times l \log l)$. The sorting of the array itself is $O(n \log n)$ but each string comparison requires l character compares in the work case so it is actually $O(l \times n \times \log n)$. Adding the two things together we obtain

$$O(n \times l \log l + l \times n \times \log n) = O(nl(\log l + \log n)) =$$

What is the asymptotic running time of the following code?

```
public static bool IsPrimeNaive(int x)
{
    if (x <= 1) return false;

    for (int i = 2; i < x; i++)
    {
        if (x % i == 0)
            return false;
    }
    return true;
}
```

The runtime is then $O(x)$

What is the asymptotic running time of the following code?

```
public bool IsPrimeUsingSquareRoot(int n)
{
    if (n < 2)
        return false;

    if (n == 2)
        return true;

    // The definition of a prime is an integer x
    // which is not exactly divisible by any
    // number other than itself and one. If a
    // number x is not prime it can be written as
    // the product of two factors a x b. If both
    // a and b were greater than the square root of
    // x then a x b would also be greater than x and hence
    // a x b is not x. SO testing all factors up to floor(sqrt(x))
    // is sufficient as if one factor is floor(sqrt(x)) the other factor must
    // be less than that

    // hence test the n-2 integers from
    // 2,..., Floor(sqrt(N))
    return Enumerable.Range(2, (int)Math.Floor(Math.Sqrt(n)))
        .All(i => n % i > 0);
}
```

The runtime is then $O(\sqrt{n})$

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What is the asymptotic running time of the following code?

```
public static int Factorial(int x)
{
    if (x == 0) return 1;
    return x * Factorial(x-1);
}
```

The running time is simple $O(x)$

Risk and Pricing Solutions

Searching

ITERATIVE BINARY SEARCH

Implement Iterative Binary Search and state its asymptotic run time?

```
public int SearchIterative<T>(IList<T> arr, T val)
    where T : IComparable<T>
{
    if (arr == null)
        throw new ArgumentNullException();

    int loIdx = 0;
    int hiIdx = arr.Count - 1;
    while (loIdx <= hiIdx)
    {
        int miIdx = loIdx + (hiIdx - loIdx) / 2;
        int comp = val.CompareTo(arr[miIdx]);

        if (comp == 0)
            return miIdx;

        if (comp > 0)
            loIdx = miIdx + 1;
        else
            hiIdx = miIdx - 1;
    }

    return ~loIdx;
}
```

The runtime is $O(\log N)$