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# I. Project Discription

# DBSCAN

DBSCAN requires two parameters: ε (eps) and the minimum number of points required to form a dense region(minPts). It starts with an arbitrary starting point that has not been visited. This point's ε-neighborhood is retrieved, and if it contains sufficiently many points, a cluster is started. Otherwise, the point is labeled as noise. Note that this point might later be found in a sufficiently sized ε-environment of a different point and hence be made part of a cluster.

If a point is found to be a dense part of a cluster, its ε-neighborhood is also part of that cluster. Hence, all points that are found within the ε-neighborhood are added, as is their own ε-neighborhood when they are also dense. This process continues until the density-connected cluster is completely found. Then, a new unvisited point is retrieved and processed, leading to the discovery of a further cluster or noise.

--[Wikipedia](<http://en.wikipedia.org/wiki/DBSCAN>)

# II.Analysis of Sample given

## dbscan.h:

Given a data structure Point and some assignments of functions.

struct Point{

double x;

double y;

int lable; // -1 unvisited, 0 noise, >0 cluster index

};

## 2. 3spiral.txt

Data set given as Point.

## 3.main.cpp

Data input to vector<Point> dataset;

Call function dbscan();

Evaluate CPU time used;

## 4.dbscan.cpp

Function **int dbscan(vector<Point> &dataset, double eps, int min\_pts)** do dbscan, given radius and minimum number of neigborhoods.

For each point **p** in dataset, we call **region\_query()** to search its neighborhood point which is legal and then push it into its queue called **neighborhood.** Then we decide whether this point **p** is a noise point. Is not we mark it a new cluster and expand this cluster which means push other legal point into this cluster.

Function **expand\_cluster()** is an algorithm which just like Breadth-First-Search.

## 5.Time Evaluation

We have run this sample and it costs **150 second.**

## 6. Which part can be parallelized?

Let’s get started from a small part!

We found that we can do better in calculating Euclidean distance between points and query if this point belong to other points’ neighborhood by parallelization.

Implement and other parts of parallelizations’ implements will be discuss below.

# III. Implements of Parallelizations(CUDA)

### 1．We found that we can do better in calculating Euclidean distance between points and query if this point belong to other points’ neighborhood by parallelization likely.

Firstly in fuction **region\_query()**

We initialized data structure **nodeX[ ], nodeY[ ]** which mark every point’s location x, y.

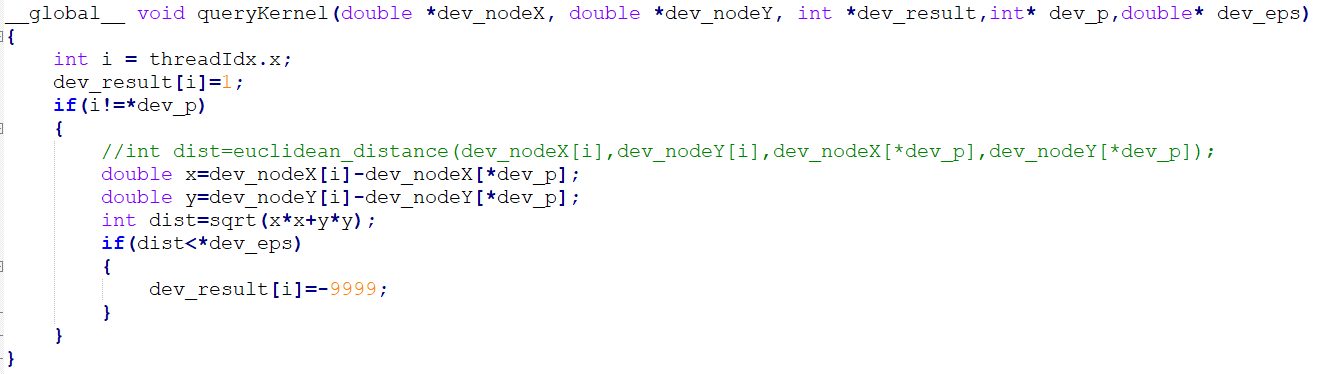
Memory initialization: cudaMalloc, cudaMemcpy(cudaMemcpyHostToDevice).

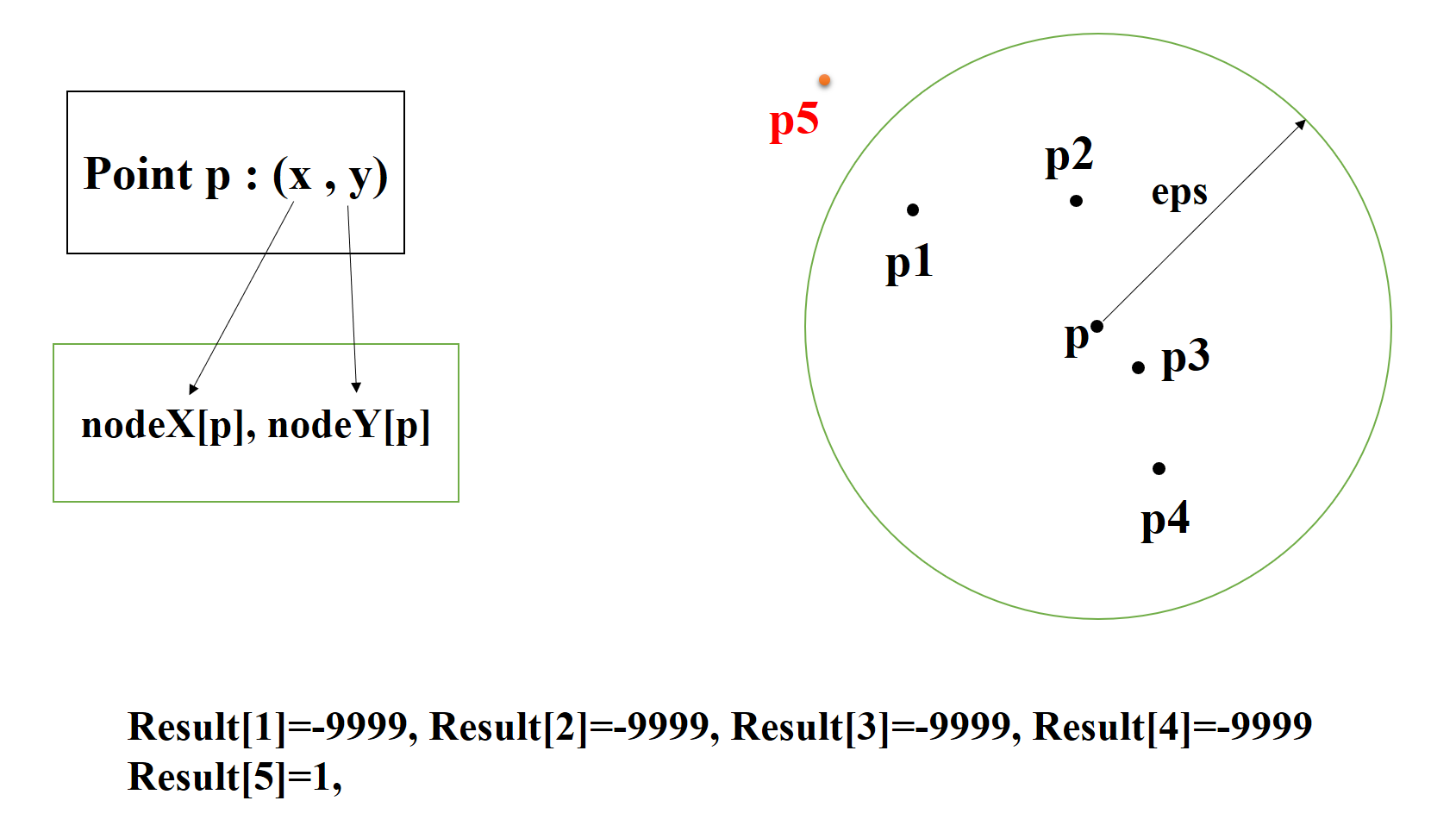


Call a query function on GPU:



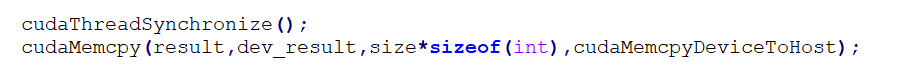
Which is defined as below:

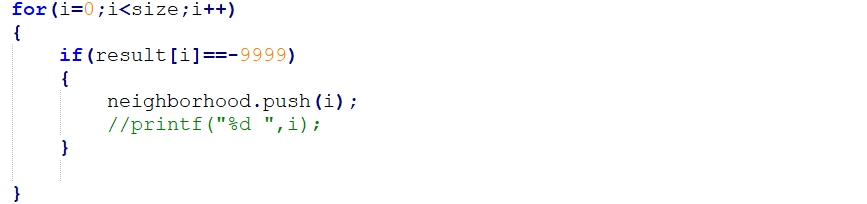




If dev\_result[i] is marked as **-9999** which means they are neighborhood between this point **p** and point **i.** If not , it will be marked as **1.**

**Then we copy this result set back to host and update points’ neighborhood:**

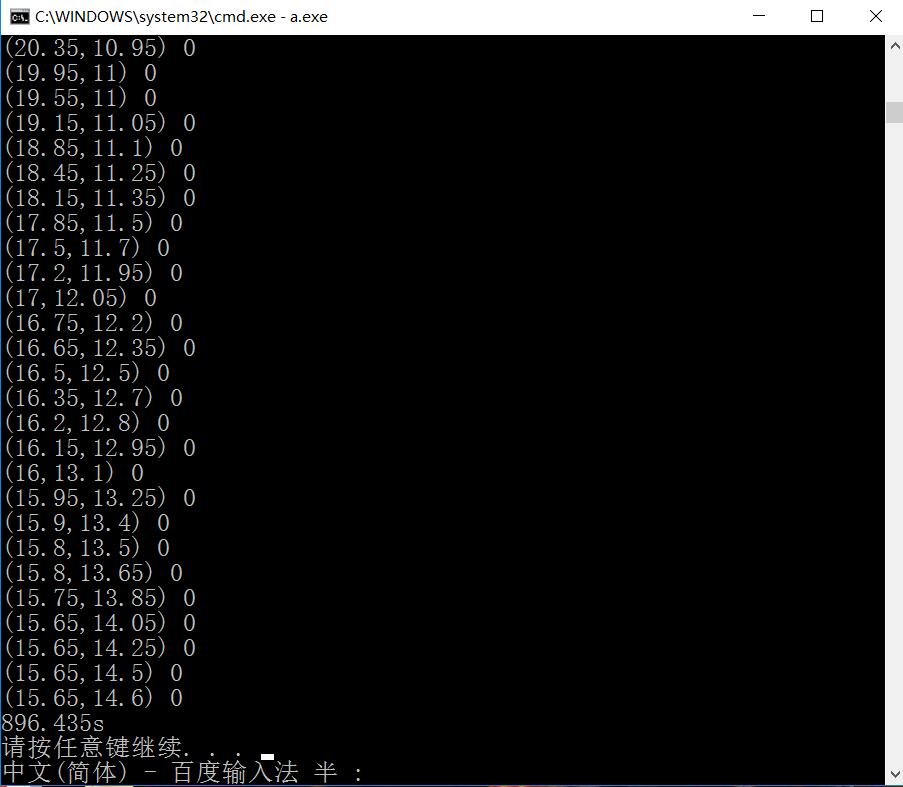




**Time evaluation:**

We found this parallel algorithm is much slower than serial algorithm about **6 times slower** than sample’s.

//screenshot



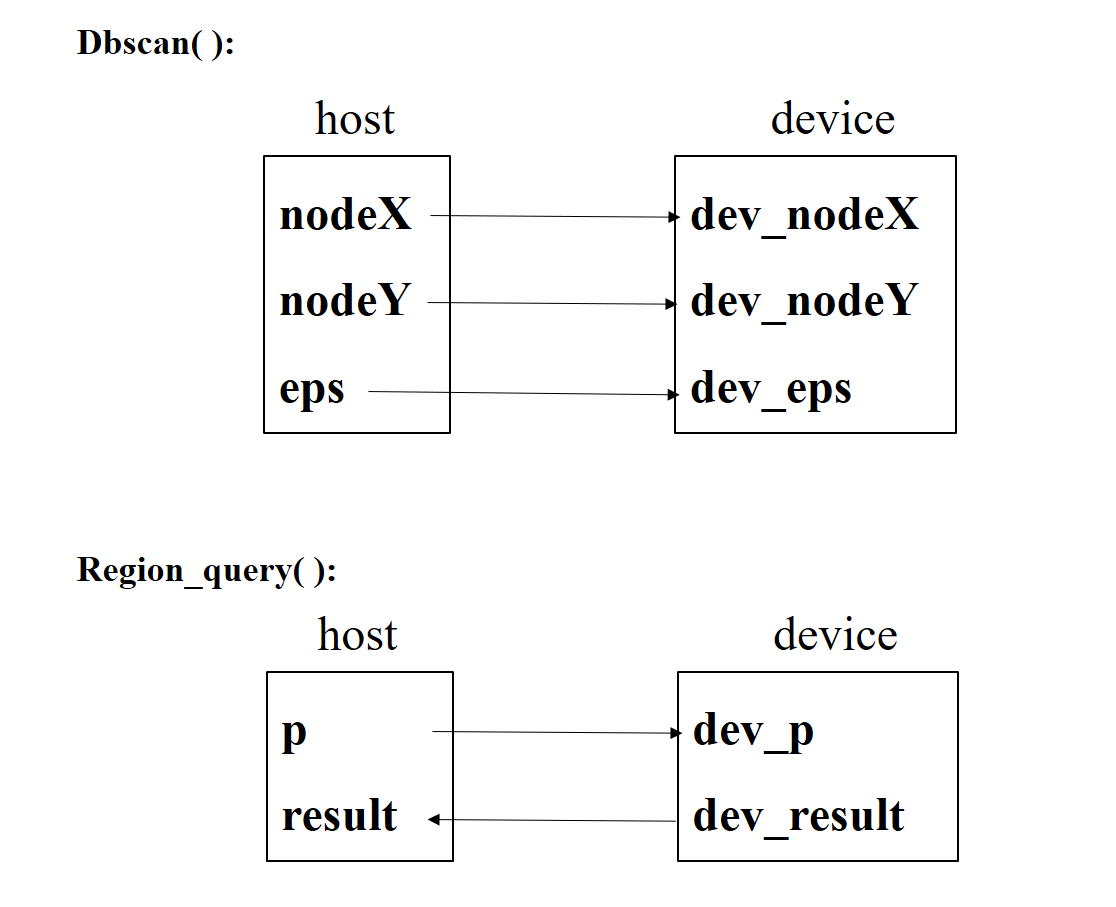
**Speedup: S = 150/896.435 = 0.167**

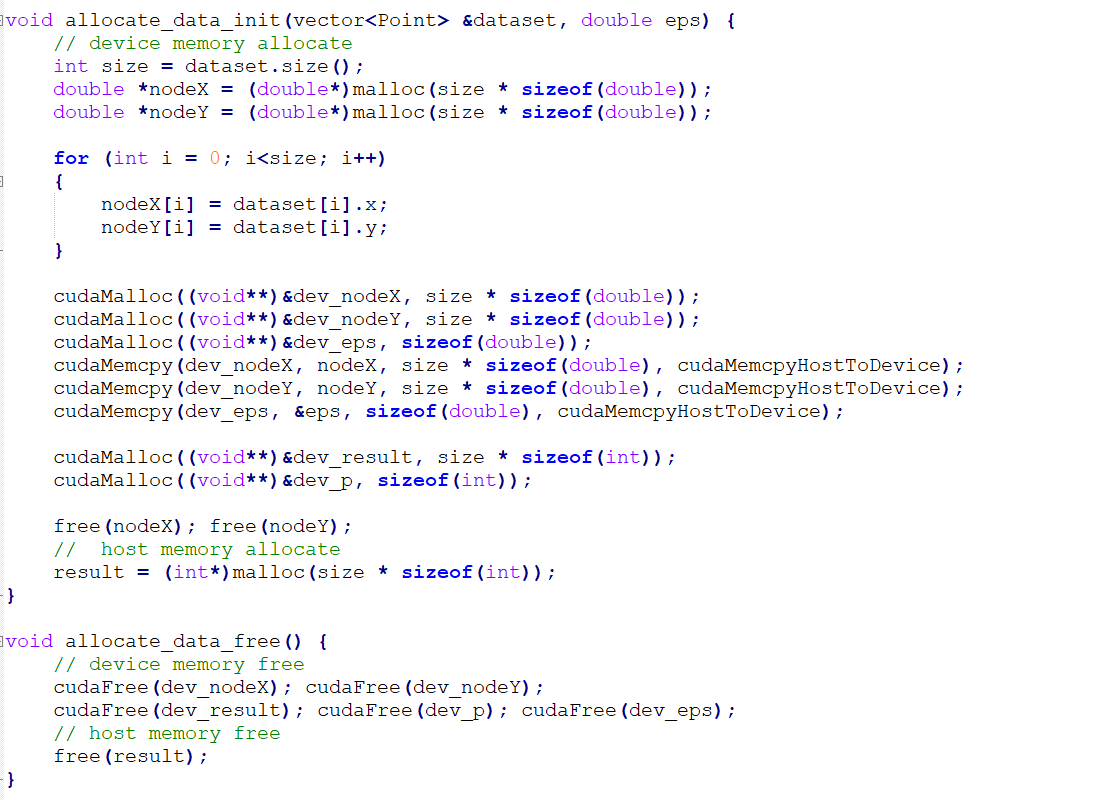
CPU：Intel core i7-6700HQ，GPU: NVIDIA GTX950M

We thought that it costs much of CPU and GPU time on memcpy between host and device. Because every time we call function **region\_query()** we do memcopy.

So we work out another solution described below.

### one time malloc, memcpy and free between host and device. We made malloc, memcpy and free between host and device as a series of functions.



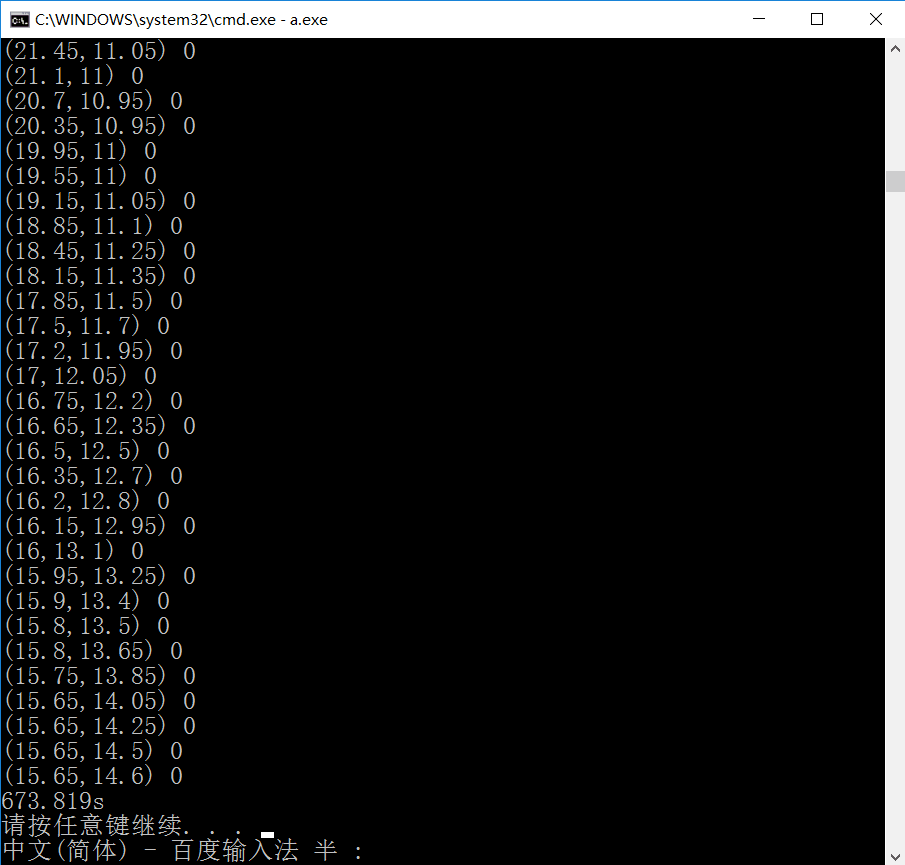


This functions will be only called for one time in function dbscan() and copy all data including nodeX/Y and eps into device one time for all.

**Time evaluation:**

We found this parallel algorithm(version2) is still slower than serial algorithm about  **4 to 5 times slower** than sample’s. we are not satisfied about that.

//screenshot



**Speedup: S = 150/673.819 = 0.223**

CPU：Intel core i7-6700HQ，GPU: NVIDIA GTX950M

We thought that every time we call function expand\_cluster() we will call another function region\_query() at least 2 times. Actually, it is not necessary to call function region\_query() so many time in function expand\_cluster().

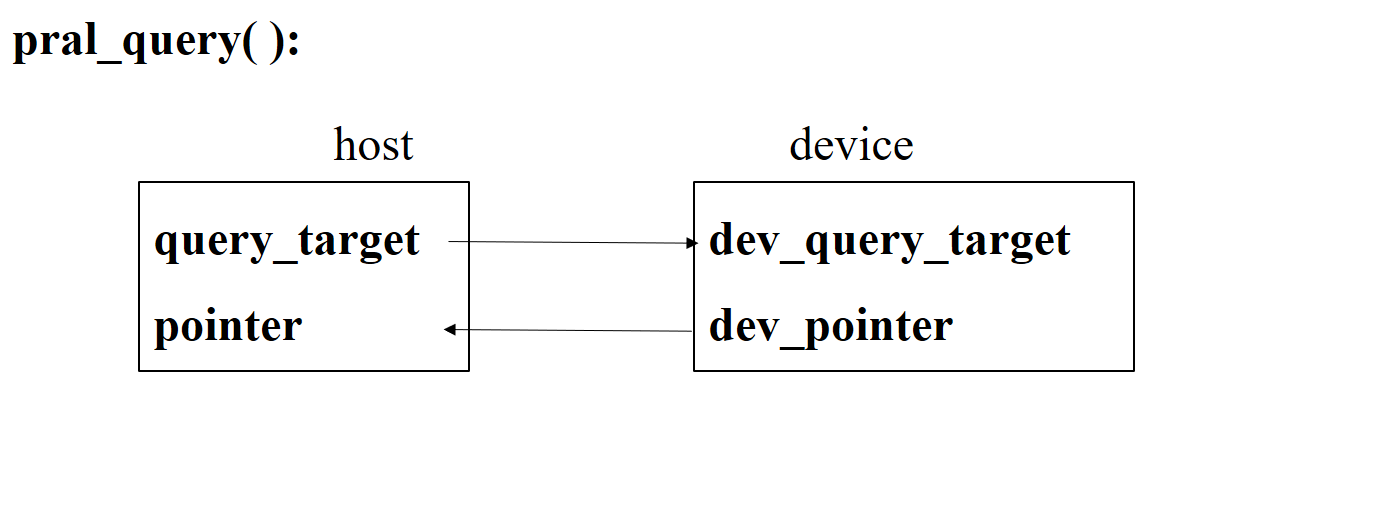
### 3.optimization for function region\_query() and get the result once for all

Just as mentioned above, it is not necessary to call function **region\_query()** so many time in function **expand\_cluster().** So we optimize function **region\_query()** and get the query result once for all**.**

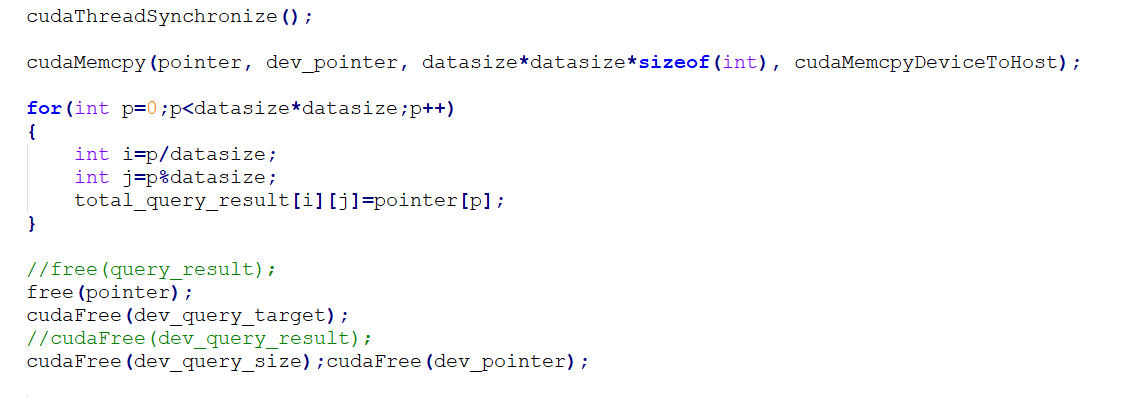
We already parallelized calculating Euclidean distance between points.

And now we will parallelize **region\_query()**!!!

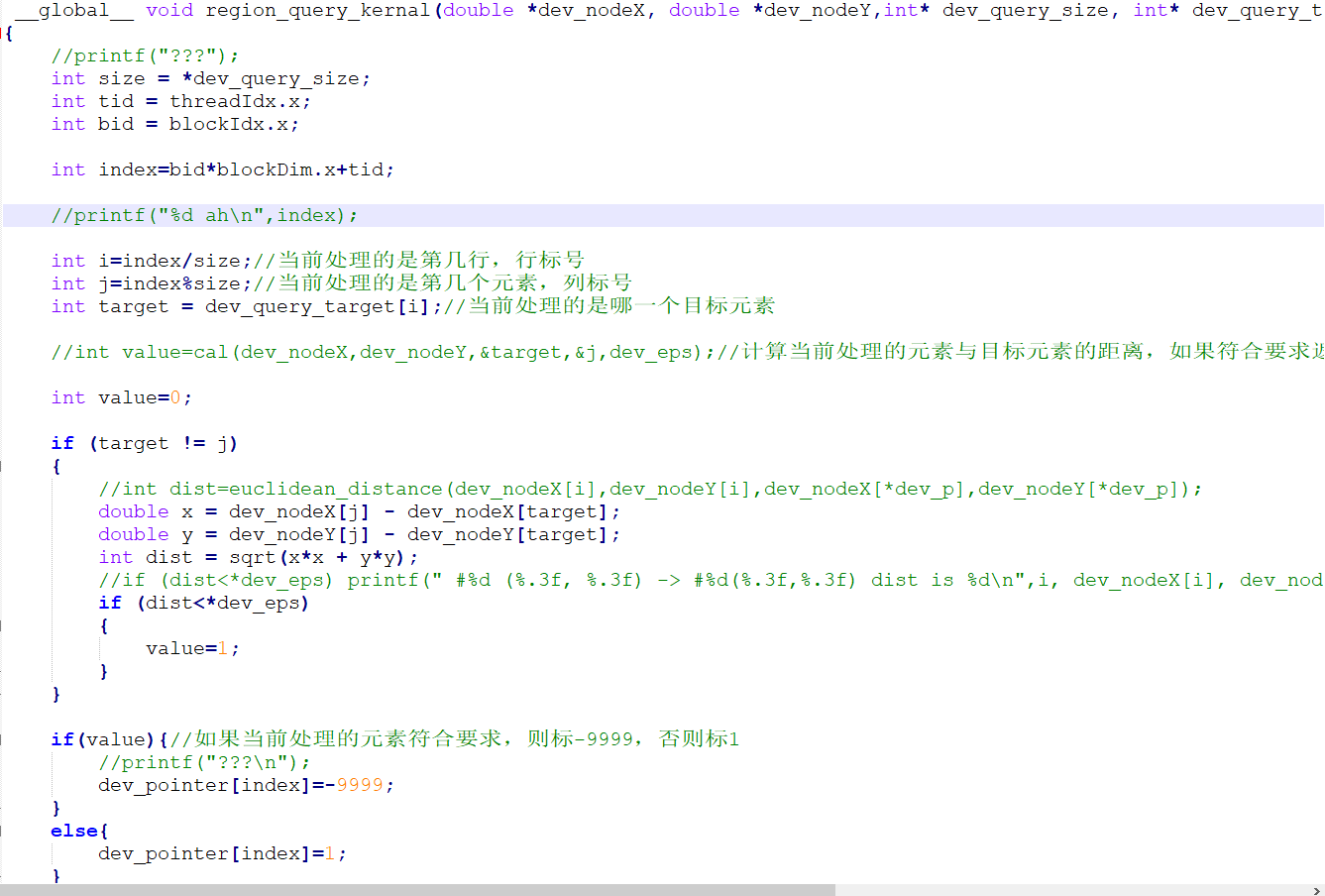
We do memcpy from host to device for region query in function **pral\_query(int\*query\_target,int\*\*total\_query\_result, double eps)**  and call **region\_query\_kernal()** on device(GPU).

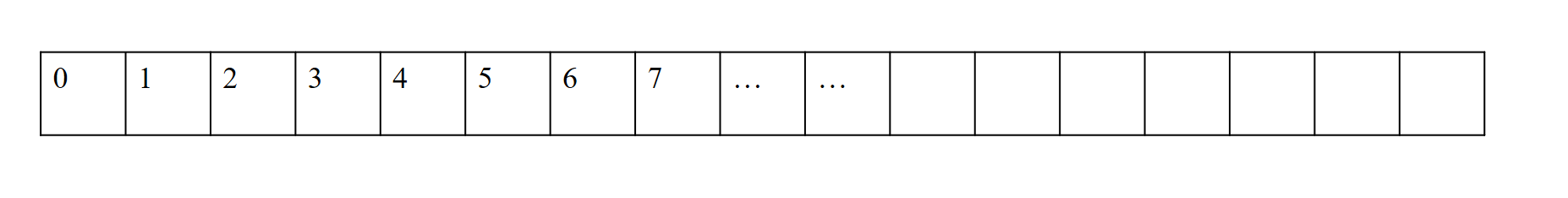




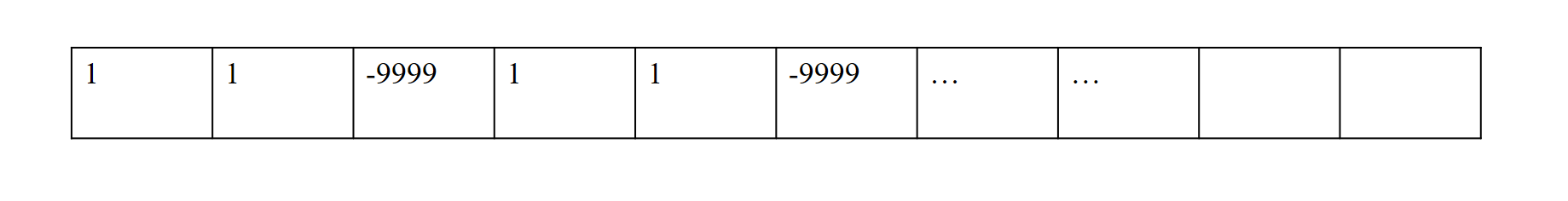


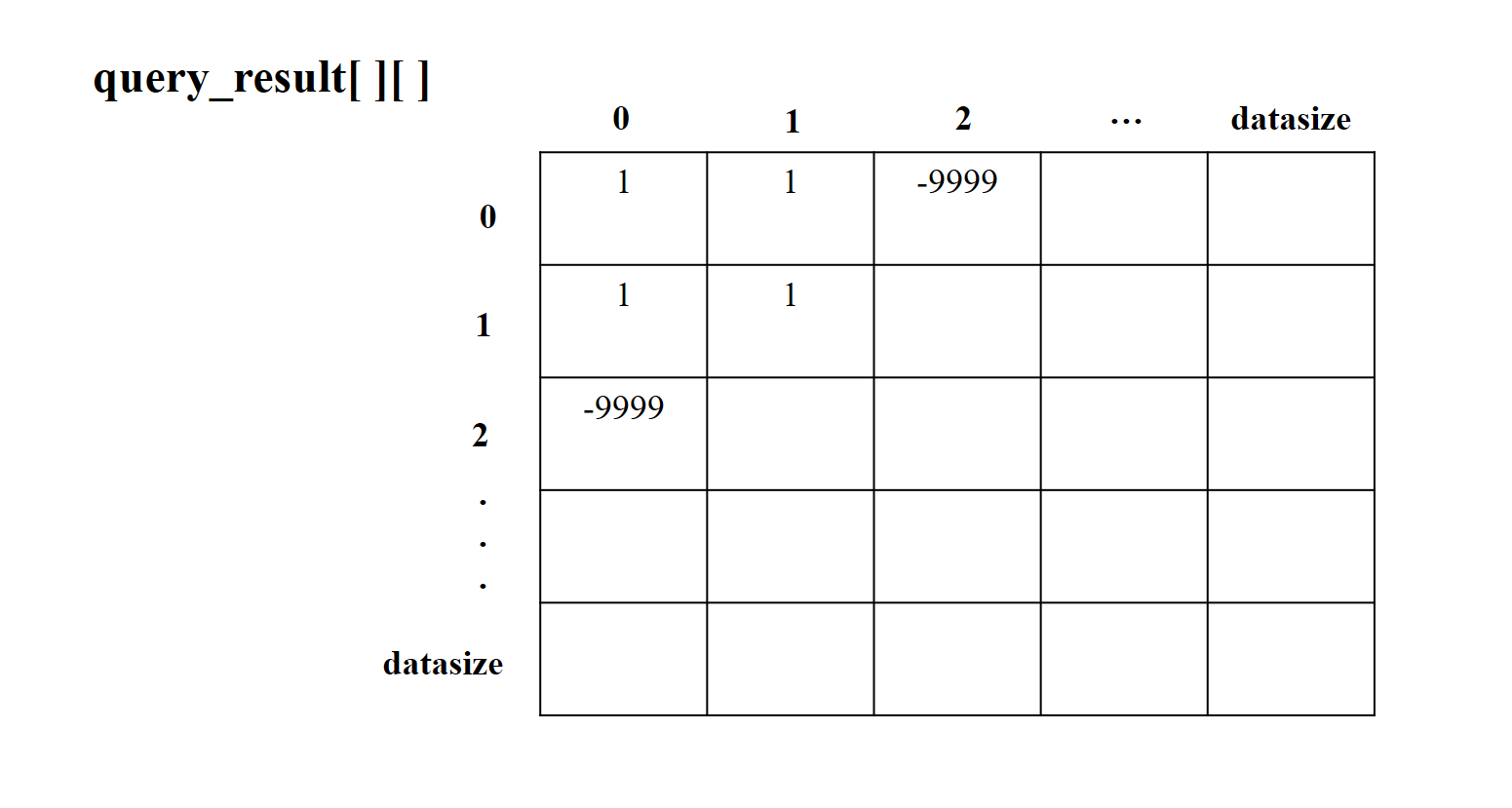
Function **region\_query\_kernal()** is define as below:



**index=blockIdx.x\*blockDim.x+threadIdx.x**

**pointer[ ]:**





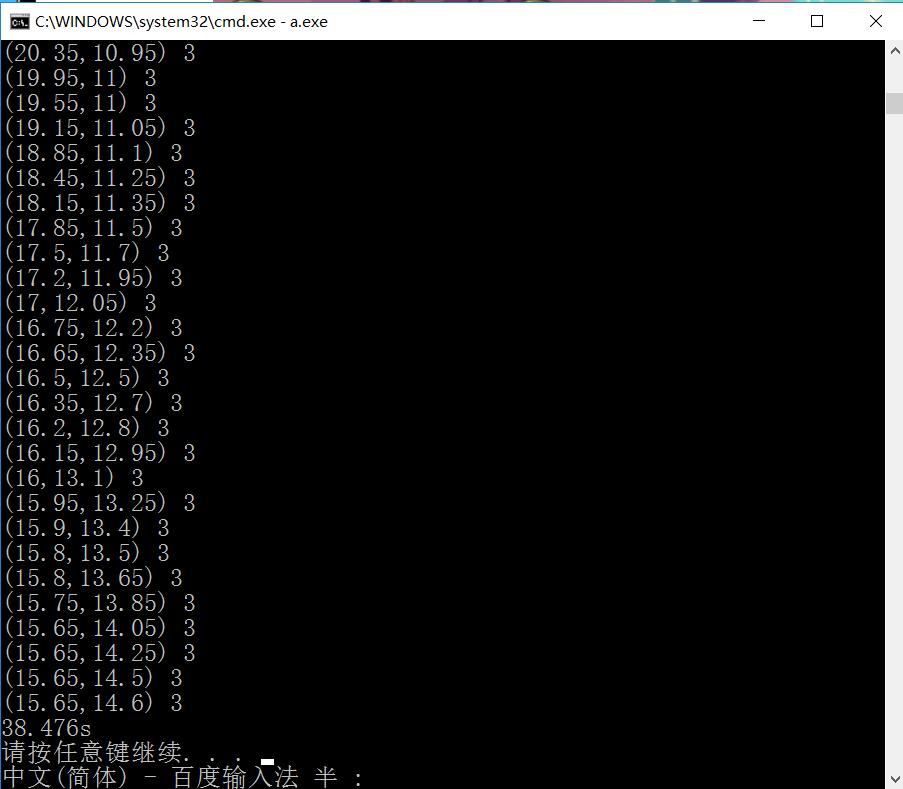
We copy those point to memory of device and compute the distance between every two points. If point **i** and point **j** are neighborhood we will mark in matrix **query\_result[i][j]** as **-9999**. If not, it will be marked as **1.**

Then in function dbscan() , we will check **query\_result[i][j]** to cnnfirm whether point I and point j are neighborhood.

**Time evaluation:**

We found this parallel algorithm(version3) is faster than serial algorithm about **4 times faster** than sample’s.

//screenshot



**Speedup: S = 150/38.476 = 3.899**

CPU：Intel core i7-6700HQ，GPU: NVIDIA GTX950M

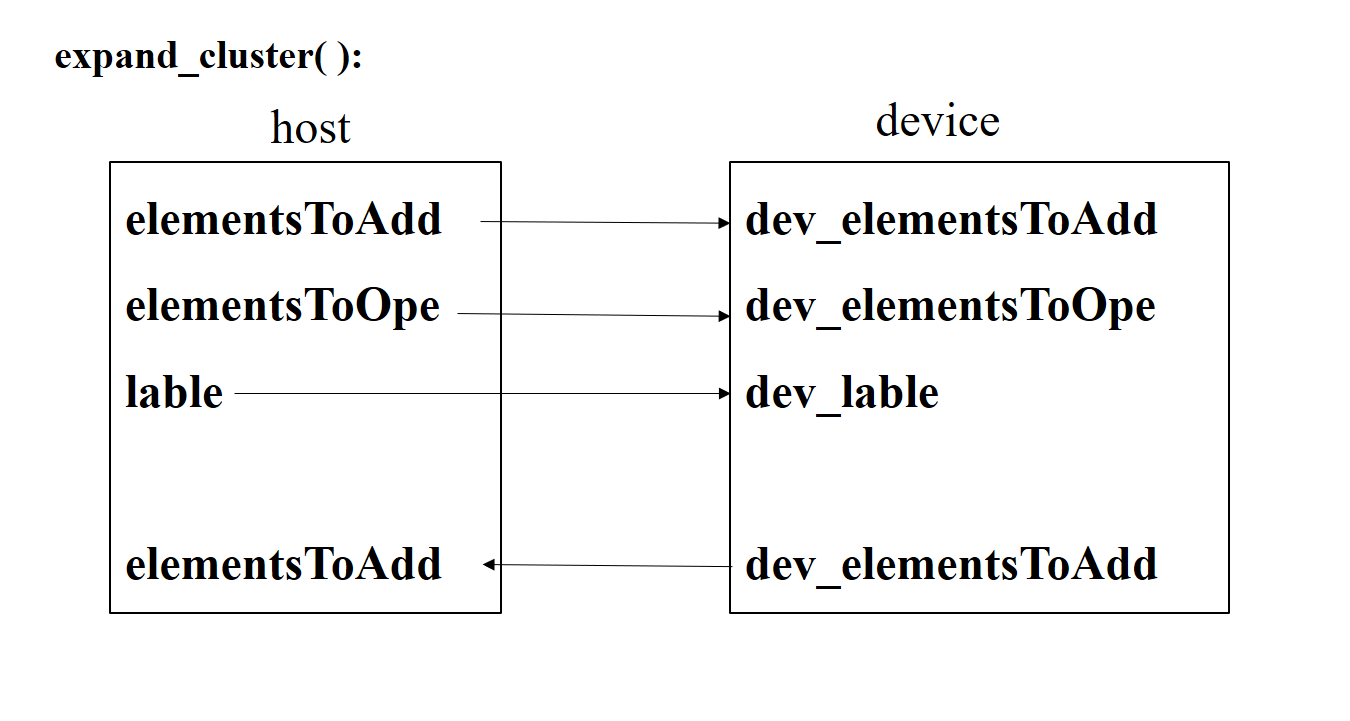
### 4.parallelize part of expand\_cluster()

In function **expand\_cluster()** we will find which points or element to push in queue by parallelization.

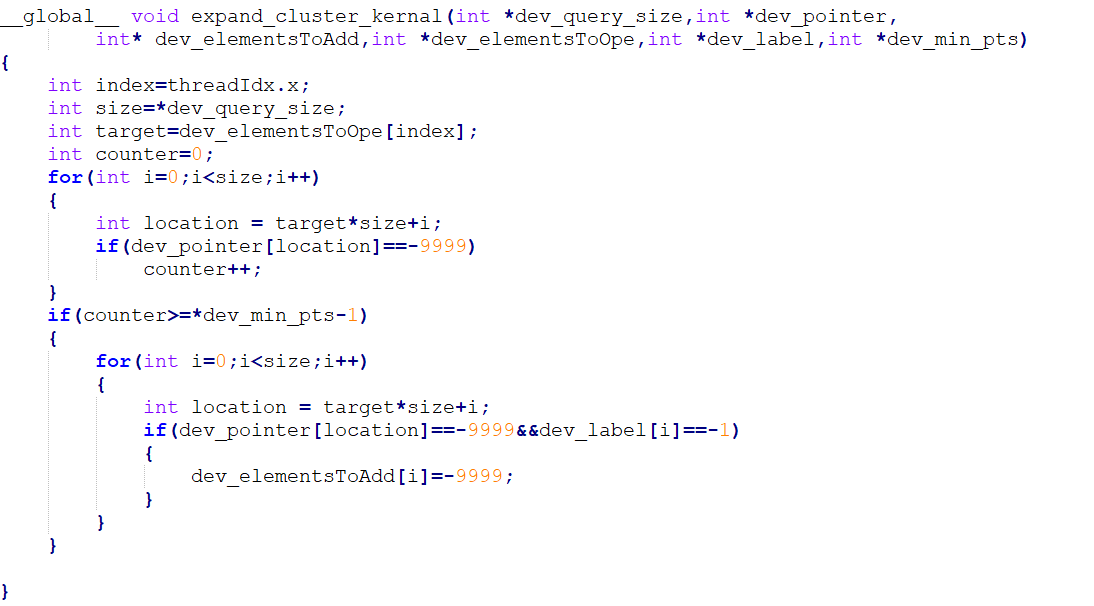
**elementsToOpe[ ]** means elements which need to be operated, and we find those elements’ neighborhood and save them to **elementsToAdd[ ]** by function **expand\_cluster\_kernal().**

We push it into a neighbor queue when we get the elementToAdd[].





Function **expand\_cluster\_kernal()** is defined as below:



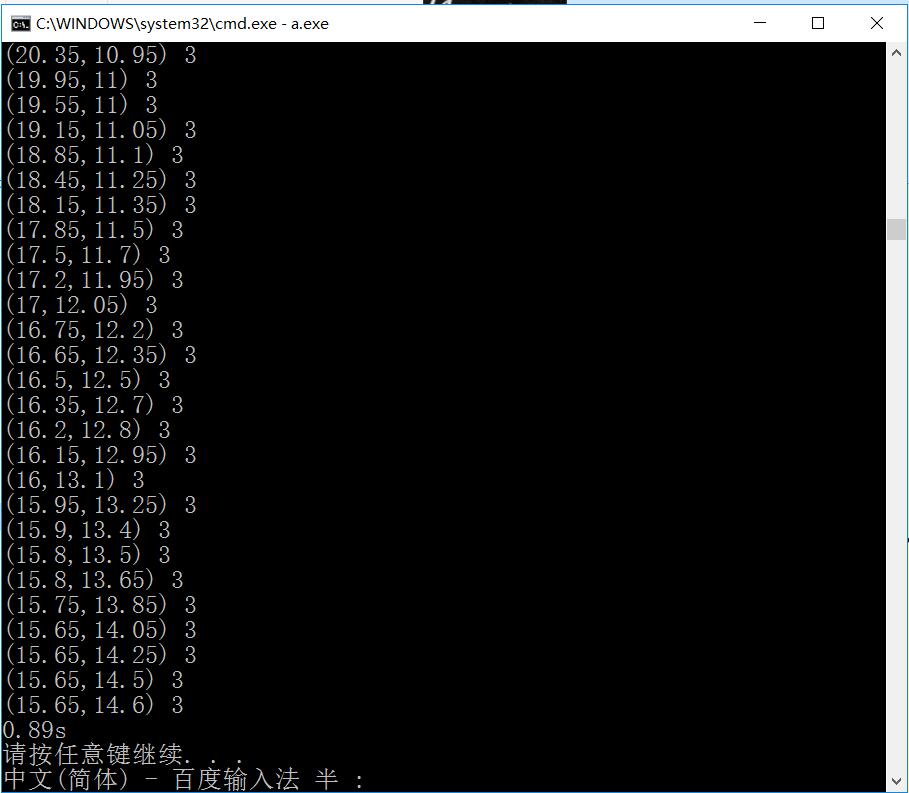
If the number of its neighbors is more than **min\_pst** and the neighbor dose not belongs to any cluster, we mark the neighbor **elementToAdd** and set the value **-9999.**

**Time evaluation:**

We found this parallel algorithm(version4) is faster than serial algorithm about **169 times faster** .

**We are satisfied about that.**

//screenshot

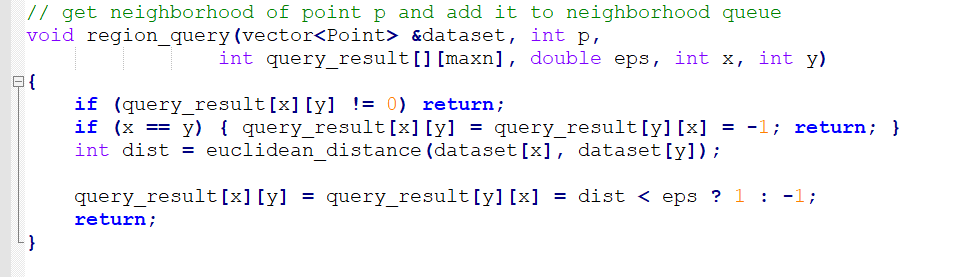


**Speedup: S = 150/0.89 = 168.539**

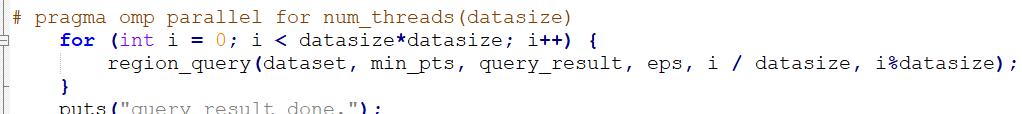
CPU：Intel core i7-6700HQ，GPU: NVIDIA GTX950M

# IV. Implements of Parallelizations(OpenMP)

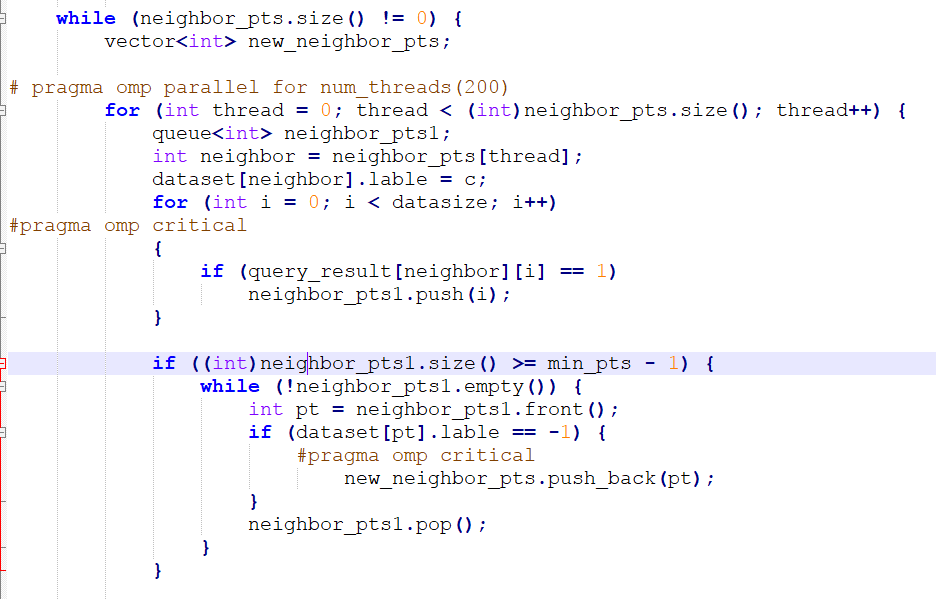
We set a matrix **query\_result[ ][ ]** which holds the “distance” between points for the same reason above. If the distance between point **i** and point **j** have not been computed **query\_result[i][j]** and **query\_result[j][i]** will be marked as **0**. If point **i** and point **j** are neighbors **query\_result[i][j]** and **query\_result[j][i]** will be marked as **1**. If not it will be marked as **-1**.



The section of querying whether they are beighbors between points can be parallelized as below:



Parallelization of **for** loop in BFS:

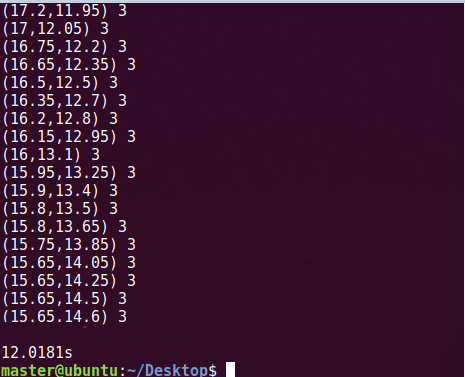


Critical section:

The operation to queue **neighbor\_pts1** and **new\_neighbor\_pts.**

**Time evaluation:**

We found this parallel algorithm(OpenMP) is faster than serial algorithm about **12.5 times faster** .



**Speedup: S = 150/12 = 12.5**

VM Ubuntu, core: 4

# V. Conclusion

We found that it’s more complex to use CUDA to implement a parallel program because of memory copy between host and device and other reasons. But the speedup of CUDA is awesome.

OpenMP is also a nice way to implement parallelization. The most important thing is OpenMP is much easier than CUDA and you can check the code line to find that. But it’s a bit slower than CUDA.