

Algorithmics	Student information	Date	Number of session
	UO: 284185	1/3/2023	3
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## Activity 1. Basic recursive models.

### DIVISION

- $O(n^k)$  if  $a < b^k$
- $O(n^k * \log n)$  if  $a = b^k$
- $O(n^{\log_b a})$  if  $a > b^k$

### SUBTRACTION

- $O(n^k)$  if  $a < 1$
- $O(n^{k+1})$  if  $a = 1$
- $O(a^{n/b})$  if  $a > 1$

- A brief explanation for each of the given classes indicating how you calculated the complexity of that class

### 1. SUBTRACTION1

$a = 1, b = 1, k = 0$

$a = 1$  because we only make one recursive call.  $b = 1$  because we subtract one in the call, and  $k = 0$  because excluding recursive calls, the complexity of the method is  $O(1)$ . Then, as we the recursive call is a subtraction, we are going to use the scheme by subtraction and thus, as  $a = 1$ , we have a complexity of  $O(n^{k+1}) = O(n)$ .

### 2. SUBTRACTION2

$a = 1, b = 1, k = 1$

$a = 1$  because we only make one recursive call.  $b = 1$  because we subtract one in the call, and  $k = 1$  because excluding recursive calls, the complexity of the method is  $O(n)$ . Then, as we the recursive call is a subtraction, we are going to use the scheme by subtraction and thus, as  $a = 1$ , we have a complexity of  $O(n^{k+1}) = O(n^2)$ .

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### 3. SUBTRACTION3

$a = 2, b = 1, k = 0$

$a = 1$  because we make two recursive call.  $b = 1$  because we subtract one in the call, and  $k = 0$  because excluding recursive calls, the complexity of the method is  $O(1)$ . Then, as we the recursive call is a subtraction, we are going to use the scheme by subtraction and thus, as  $a = 1$ , we have a complexity of  $O(a^{(n/b)}) = O(2^n)$ .

### 4. DIVISION1

$a = 1, b = 3, k = 1$

$a = 1$  because we make two recursive call.  $b = 3$  because we divide by three in the call, and  $k = 1$  because excluding recursive calls, the complexity of the method is  $O(n)$ . Then, as we the recursive call is a division, we are going to use the scheme by division and thus, as  $1 < 3^1$  ( $a < b^k$ ), we have a complexity of  $O(n^k) = O(n)$ .

### 5. DIVISION2

$a = 2, b = 2, k = 1$

$a = 2$  because we make two recursive calls.  $b = 2$  because we divide by two in the call, and  $k = 1$  because excluding recursive calls, the complexity of the method is  $O(n)$ . Then, as we the recursive call is a division, we are going to use the scheme by division and thus, as  $2 = 2^1$  ( $a = b^k$ ), we have a complexity of  $O(n * \log n)$ .

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## 6. DIVISION3

$a = 2, b = 2, k = 0$

$a = 2$  because we make two recursive calls.  $b = 2$  because we divide by two in the call, and  $k = 0$  because excluding recursive calls, the complexity of the method is  $O(1)$ . Then, as the recursive call is a division, we are going to use the scheme by division and thus, as  $2 > 2^0$  ( $a > b^k$ ), we have a complexity of  $O(n)$ .

- A brief explanation for each of the 2 new classes indicating how you calculate the complexity to get the requested one.

### 1. SUBSTRATION4

$a = 3, b = 2, k = 0$

$a = 3$  because we make two recursive calls.  $b = 2$  because we subtract two in the call, and  $k = 0$  because excluding recursive calls, the complexity of the method is  $O(1)$ . Then, as the recursive call is a subtraction, we are going to use the scheme by subtraction and thus, as  $3 > 1$  ( $a > 1$ ), we have a complexity of  $O(3^{(n/2)})$ .

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## 2. DIVISION4

$a = 4, b = 2, k = 1$

$a = 4$  because we make two recursive calls.  $b = 2$  because we divide by two in the call, and  $k = 1$  because excluding recursive calls, the complexity of the method is  $O(n)$ . Then, as we the recursive call is a division, we are going to use the scheme by division and thus, as  $4 > 2^1$  ( $a > b^k$ ), we have a complexity of  $O(n^2)$ .