

Logical and Shift Operations



Up Until Now

- Up until now, we've done
 - Arithmetic: `add`, `sub`, `addi`
 - Memory access: `lw` and `sw`
 - branches and jumps: `j`, `jr`, `jal`, `beq`, `bne`
- These instructions view contents of register as a **single quantity** (such as a signed or unsigned integer)

Bitwise Operations

- View contents of register as 32 **independent bits**
 - Since registers are composed of 32 bits, we may want to access individual bits (or groups of bits) rather than the whole.
- Two new classes of MIPS instructions for bitwise operations:
 - Logical Operators
 - Shift Operators

Bitwise Operations

- Truth Table: lists all combinations of inputs and outputs

A	B	AND	OR	NOR
0	0	0	0	1
0	1	0	1	0
1	0	0	1	0
1	1	1	1	0

AND: outputs 1 if both inputs are 1

OR : outputs 1 if at least one input is 1

NOR: outputs 1 if both inputs are 0

Bitwise Operations

- Bitwise result **applies function to each bit independently**
- The i^{th} bit of inputs produce the i^{th} bit of outputs

A	B	AND	OR	NOR
0	0	0	0	1
0	1	0	1	0
1	0	0	1	0
1	1	1	1	0

Bit-wise AND

AND(01101001,
11001100)
= 01001000

Bit-wise OR

OR (01101001,
11001100)
= 11101101

Boolean function applied
at each bit position


MIPS Logical Operators

MIPS Logical Operations

- MIPS Logical Operators are *bitwise operations*
- Basic MIPS logical operators

Syntax

```
and TargetReg, SourceReg1, SourceReg2  
or  TargetReg, SourceReg1, SourceReg2  
nor TargetReg, [SourceReg1, SourceReg2]
```



- Like many MIPS instructions, logical operations accept exactly 2 inputs and produce 1 output

Use Logical Operator in Conditional Statement

- Conditional statements

```
// C
if ( A < 0 && B < 0 ) {
    ...
}
```

```
$t0 = 0000 0000 0000 0000 0000 0000 0000 0001
$t1 = 0000 0000 0000 0000 0000 0000 0000 0000
$t0 and $t1 = 0000 0000 0000 0000 0000 0000 0000 0000
```

```
                # MIPS code
slt $t0, $s0, $zero # $t0 = $s0 < 0?
slt $t1, $s1, $zero # $t1 = $s1 < 0?
and $t2, $t0, $t1   # $t2 = ( $s0 < 0 && $s1 < 0 ) ?
```


Use Logical Operator in Conditional Statement

- Conditional statements

```
// C
if ( A < 0 || B < 0 ) {
    ...
}
```

```
$t0 = 0000 0000 0000 0000 0000 0000 0000 0001
$t1 = 0000 0000 0000 0000 0000 0000 0000 0000
$t0 or $t1 = 0000 0000 0000 0000 0000 0000 0000 0001
```

```
                # MIPS code
slt $t0, $s0, $zero # $t0 = $s0 < 0?
slt $t1, $s1, $zero # $t1 = $s1 < 0?
or  $t2, $t0, $t1   # $t2 = ( $s0 < 0 || $s1 < 0 ) ?
```

Logical Operators with Immediate

- Similar to `and`, `or`, but the third argument is an immediate

Syntax

```
andi TargetReg, SourceReg, Immediate
```

```
ori  TargetReg, SourceReg, Immediate
```

NOR for NOT

- Boolean expressions are made with **AND** and **OR** and **NOT**
- Why is **NOT** not a MIPS instruction?
 - **NOT** takes one operand and produces one result, which is not in keeping with the three-operand format of other instructions
 - How do we do **NOT** with **NOR**?

```
nor $t1, $t0, $zero
```

\$zero	Input	NOR
0	0	1
0	1	0

Mask

An Application of Logical Operators



No, not this one

What is a Mask?

- **A mask help us:**
 - Query a bit to find out if it is a 1 or a 0
 - Change the value of a bit to 1 or a 0
- Based on observation of **AND** & **OR** logical operation properties.

Use AND for Mask

- Any bit *and* 0 produces an output 0

0	Input	AND
0	0	0
0	1	0

This can be used to
create a *mask*.

- Any bit *and* 1 produces the original bit

1	Input	AND
1	0	0
1	1	1

Use AND for Mask

Example:

A = 1011 0110 1010 0100 0011 1101 1001 1010

B = 0000 0000 0000 0000 0000 0000 1111 1111

A and B = 0000 0000 0000 0000 0000 0000 1001 1010

- In this example, **B** is called a **mask**.
- **B** is used to isolate the rightmost 8 bits of **A** by *masking* out the rest of the string (e.g., setting it to all 0s)
- Thus, the **and** operator can be used to set certain portions of a bitstring to 0s, while leaving the rest alone.

Use AND for Mask

Example: If A = 0xB6A43D9A is saved in \$t0, then what is \$t1 and \$t2 after the following instructions?

```
andi $t1, $t0, 0xFF
```

```
$t0 = 1011 0110 1010 0100 0011 1101 1001 1010  
0xFF = 0000 0000 0000 0000 0000 0000 1111 1111  
$t1 = 0000 0000 0000 0000 0000 0000 1001 1010 = 0x9A
```

```
andi $t2, $t0, 0x000000FF
```

```
$t0 = 1011 0110 1010 0100 0011 1101 1001 1010  
0x000000FF = 0000 0000 0000 0000 0000 0000 1111 1111  
$t2 = 0000 0000 0000 0000 0000 0000 1001 1010 = 0x9A
```


Use AND for Mask

Example: If A = 0xB6A43D9A is saved in \$t0, then what is \$t1 and \$t2 after the following instructions?

```
andi $t1, $t0, 0xFF
```

```
$t0 = 1011 0110 1010 0100 0011 1101 1001 1010  
0xFF = 0000 0000 0000 0000 0000 0000 1111 1111  
$t1 = 0000 0000 0000 0000 0000 0000 1001 1010 = 0x9A
```

```
andi $t2, $t0, 0x000000FF
```

```
$t0 = 1011 0110 1010 0100 0011 1101 1001 1010  
0x000000FF = 0000 0000 0000 0000 0000 0000 1111 1111  
$t2 = 0000 0000 0000 0000 0000 0000 1001 1010 = 0x9A
```

Uses OR for Mask

- Any bit **or** 0 produces the original bit

0	Input	OR
0	0	0
0	1	1

- Any bit **or** 1 produces 1

0	Input	OR
1	0	1
1	1	1

**This can also be used
to create a *mask*.**

Uses OR for Mask

- Can be used to force certain bits of a string to 1s.

Example: if \$t0 contains 0x12345678, then after

```
ori $t1, $t0, 0xFFFF
```

\$t1 contains 0x1234FFFF (high-order 16 bits are untouched, low-order 16 bits are forced to 1s).

\$t0	=	0001	0010	0011	0100	0101	0110	0111	1000
0xFFFF	=	0000	0000	0000	0000	1111	1111	1111	1111
\$t1	=	0001	0010	0011	0100	1111	1111	1111	1111

Why it's important

Controller:

- A card that interfaces with a peripheral
- It uses three special purpose registers
 - Status – bit flags, 1 – true, 0 – false
 - Command – bit flags or integers to send commands



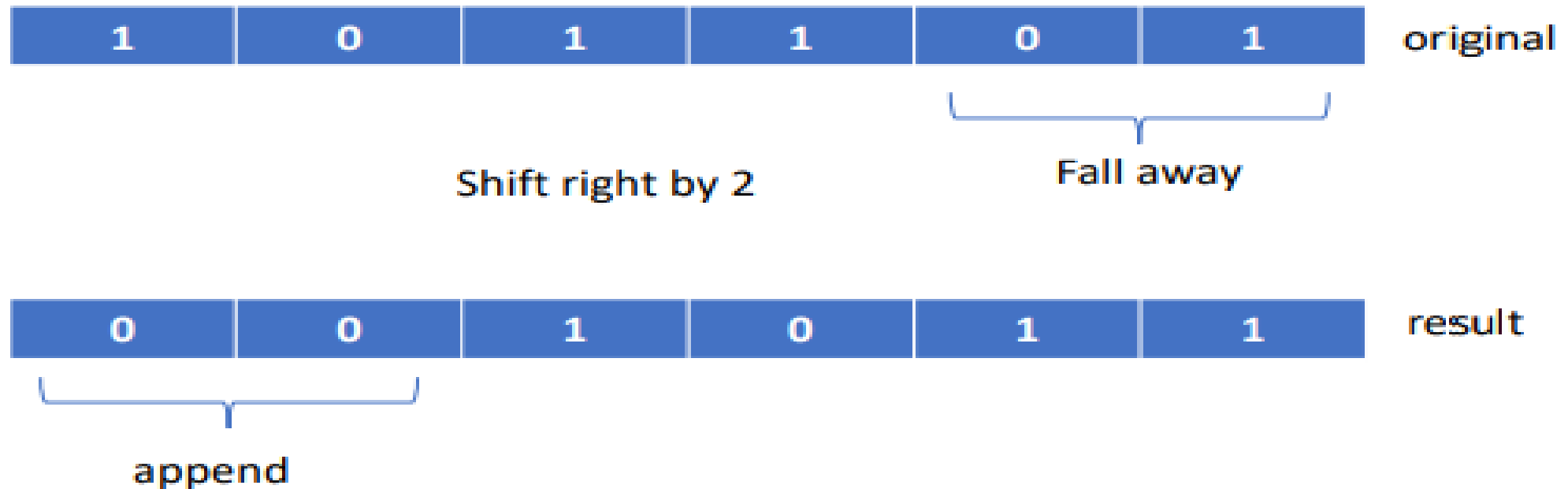
- Check the individual bits in ALU status register

Shift Operators

<<<<<<<<<

What is shifting?

- Moving the bits in a register a certain number of positions to the left or right.



Shift Logical Instructions

- Shift Instruction Syntax:

Operation **TargetReg**, **SourceReg**, **ShiftAmount**

- 1) **Operation**: operation name
- 2) **TargetReg**: register that will receive value
- 3) **SourceReg**: register that contains the original value
- 4) **ShiftAmount**: shift amount (non-negative constant < 32)

Shift Logical Instructions

- Shift left logical **sll**
 - shifts left and fills emptied bits with 0s

```
sll TargetReg, SourceReg, ShiftAmount
```

- Shift right logical **srl**
 - shifts right and fills emptied bits with 0s

```
srl TargetReg, SourceReg, ShiftAmount
```


Shift Logical Instructions

Example:

Assume \$t0 contains 0001 0010 0011 0100 0101 0110 0111 1000

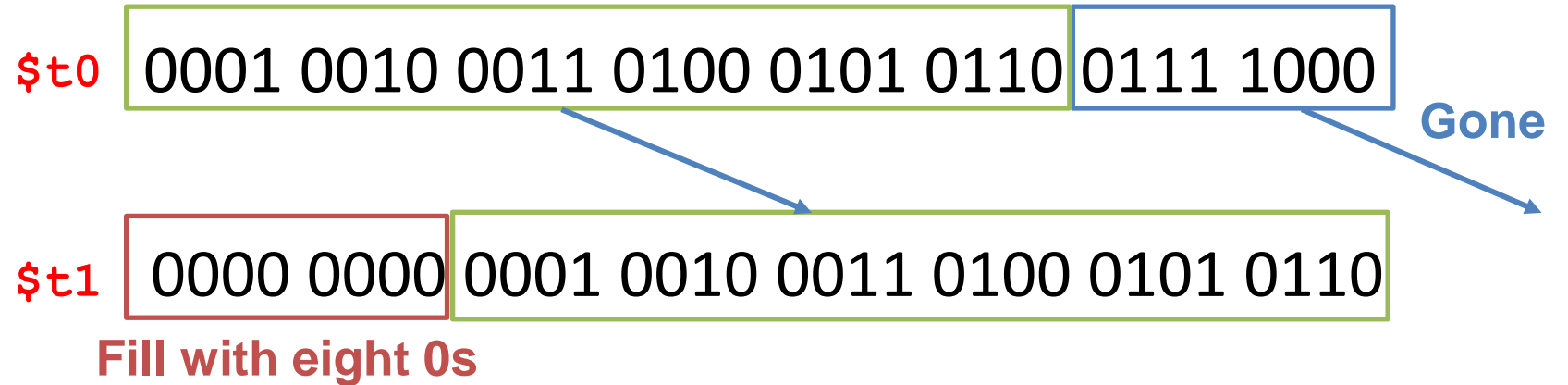
What are \$t1 and \$t2

```
# shift right  
srl $t1, $t0, 8
```

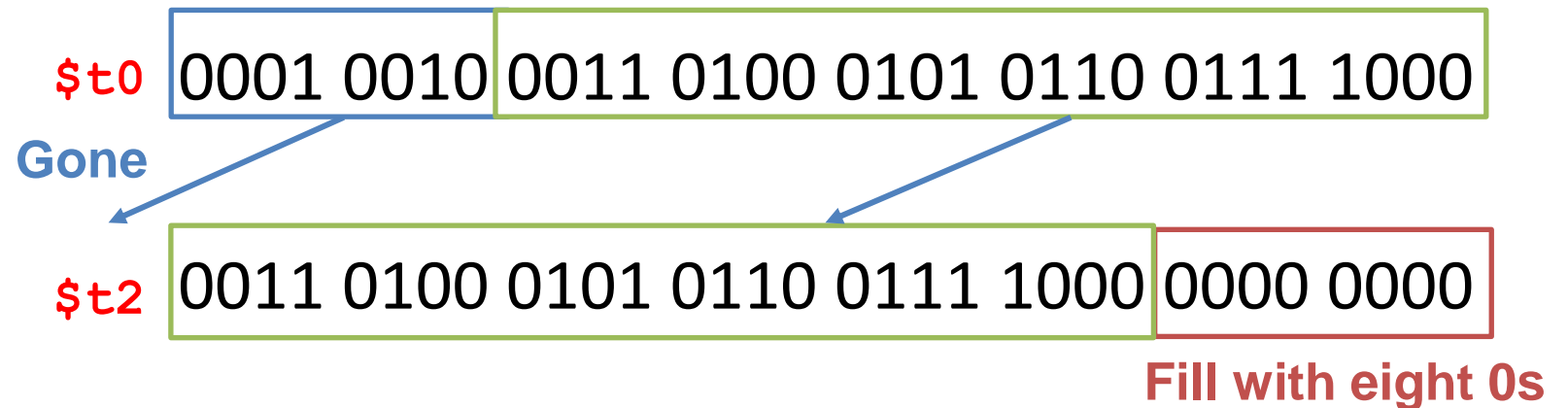
```
# shift left  
sll $t2, $t0, 8
```

Shift Logical Instructions

```
# shift right  
srl $t1, $t0, 8
```



```
# shift left  
sll $t2, $t0, 8
```



Shift Arithmetic Instructions

- Shift right arithmetic **sra**
 - Shifts right and fills emptied bits by sign extending (sensitive to sign bit)

```
sra TargetReg, SourceReg, ShiftAmount
```

- Why? A negative number should stay negative after shifting
 - If MSB = 0, shift and fill the new bits with 0s
 - If MSB = 1, shift and fill the new bits with 1s

Shift Arithmetic Instructions

Example: SRA (shift right arithmetic) by 8 bits

\$t3 = 0001 0010 0011 0100 0101 0110 0111 1000

\$t4 = 1001 0010 0011 0100 0101 0110 0111 1000

What happen after shift right arithmetic by 8 bits?

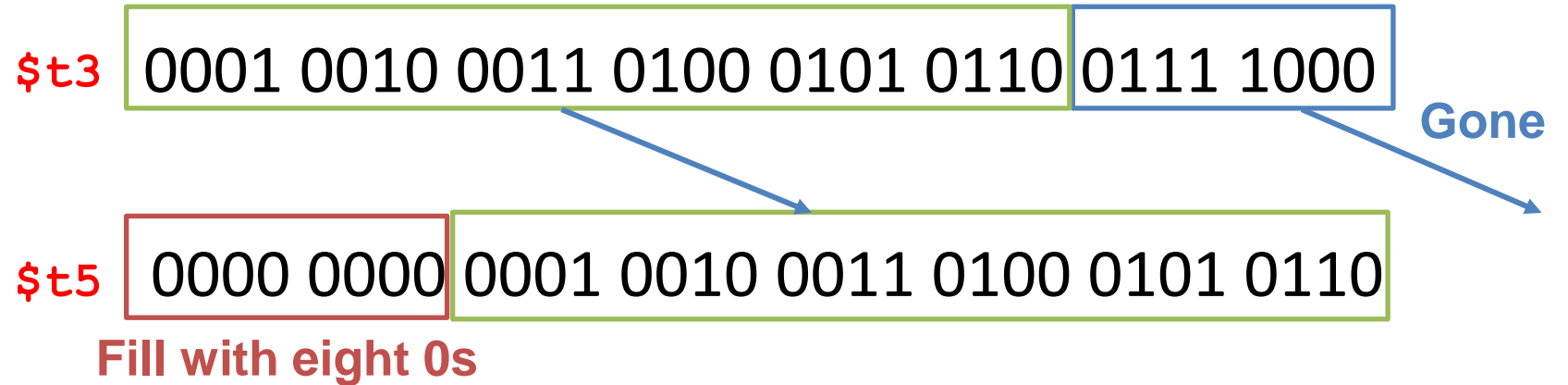
```
sra $t5, $t3, 8
```

```
sra $t5, $t4, 8
```

Shift Arithmetic Instructions

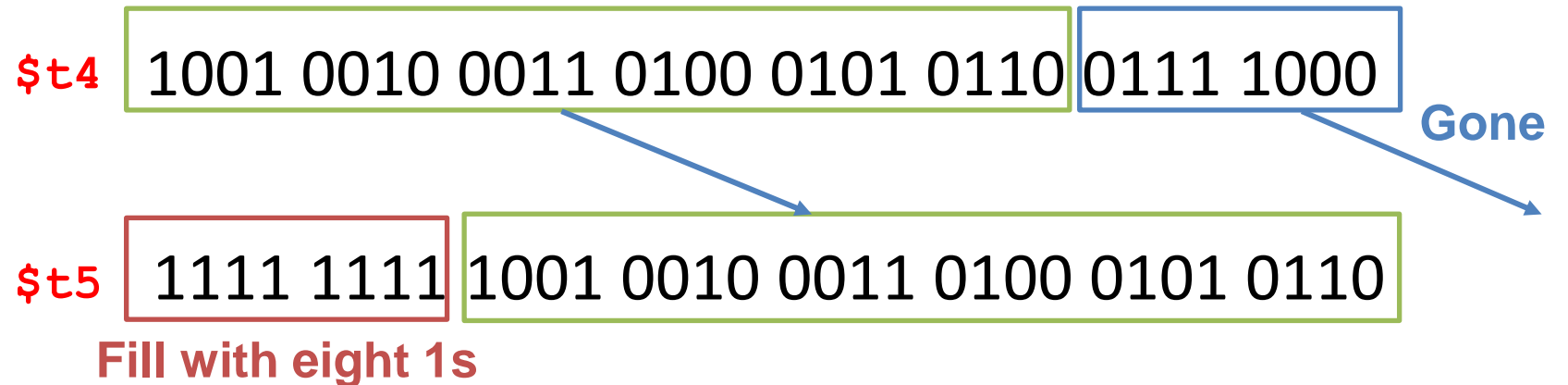
```
# shift right  
sra $t5, $t3, 8
```

If MSB = 0, the new bit
after shifting = 0



```
# shift right  
sra $t5, $t4, 8
```

If MSB = 1, the new bit
after shifting = 1



Use Shift Instructions in Multiplication

- In decimal:
 - Multiplying by 10 = shifting left by 1: $714_{10} \times 10_{10} = 7140_{10}$
 - Multiplying by 100 = shifting left by 2: $714_{10} \times 100_{10} = 71400_{10}$
 - Multiplying by 10^n = shifting left by n
- In binary:
 - Multiplying by 2 = shifting left by 1: $11_2 \times 10_2 = 110_2$
 - Multiplying by 4 = shifting left by 2: $11_2 \times 100_2 = 1100_2$
 - Multiplying by 2^n = shifting left by n

Use Shift Instructions in Multiplication

- Shifting maybe faster than multiplication!
 - a good compiler usually notices when C code multiplies by a power of 2 and compiles it to a shift instruction:

```
a = a * 8 ; // C
```



```
sll $s0, $s0, 3 # MIPS
```

- Likewise, shift right to divide by powers of 2
 - Use `sra` but watch out for negative numbers as the result is rounded down

```
b = b / 2 ; // C
```



```
sra $s1, $s1, 1 # MIPS
```

with \$s1 = -3 answer is -2

Why it's important

Fast multiplication by 2

Fast division by 2

Regular multiplication and division require n clock cycles, where n is the number of steps.

Shifting runs in a single clock cycle.

Shift Arithmetic Instructions

Example: $\$t0 = 1001_2$ What happen after the operations?

```
sra $t1, $t0, 1  
srl $t2, $t0, 1  
sll $t3, $t0, 1
```

Shift Arithmetic Instructions

Example: $\$t0 = 1001_2$ What happen after the operations?

```
sra $t1, $t0, 1  
srl $t2, $t0, 1  
sll $t3, $t0, 1
```

$\$t0 = 1001_2 = -7_{10}$

$\$t1 = 1100_2 = -4_{10}$

$\$t2 = 0100_2 = 4_{10}$

$\$t3 = 0010_2 = 2_{10}$

Use Shift to Extract Information

- Suppose we want to isolate byte 0 (rightmost 8 bits) of a word stored in \$t0

```
andi $t0, $t0, 0xFF
```

- Suppose we want to isolate byte 1 (bit 15 to bit 8) of a word stored in \$t0.

```
andi $t0, $t0, 0xFF00
```

- How do we “extract” the information?



Use Shift to Extract Information


```
sll $t0, $t0, 16  
srl $t0, $t0, 24
```

0001 0010 0011 0100 0101 0110 0111 1000

0101 0110 0111 1000 0000 0000 0000 0000

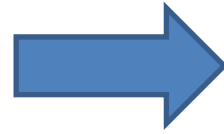
0101 0110 0111 1000 0000 0000 0000 0000

0000 0000 0000 0000 0000 0000 0101 0110



Application: Pixel Data for Images

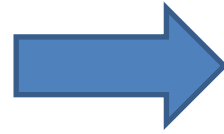
Example: Packing pixel data for images



- Suppose each pixel of an image has
 - *Red* r $[0 \sim 255]$
 - *Green* g $[0 \sim 255]$
 - *Blue* b $[0 \sim 255]$
 - *Alpha* a , transparency value $[0 \sim 255]$

Each of them can be
represented by a byte (8-bit)

Example: Packing pixel data for images



- Instead of using 4 registers for r, g, b, and a, we **pack** each of 8 bits into a 32-bit integer

```
// C code
int packARGB ( int r, int g, int b, int a ) {
    return a << 24 | r << 16 | g << 8 | b ;
}
```

Example: Packing pixel data for images

```
// C code
int packARGB ( int r, int g, int b, int a ) {
    return a << 24 | r << 16 | g << 8 | b ;
}
```

```
# MIPS
sll    $t0, $a3, 24      # shift a to the left
or     $v0, $t0, $zero   # combine a with $v0
sll    $t0, $a0, 16      # shift r to the left
or     $v0, $t0, $v0     # combine r with $v0
sll    $t0, $a1, 8       # shift g to the left
or     $v0, $t0, $v0     # combine g with $v0
or     $v0, $a2, $v0     # combine b with $v0
jr     $ra               # return to $ra
```


Review

- Logical and Shift Instructions operate on bits individually, unlike arithmetic, which operate on entire word
- Use Logical and Shift Instructions to isolate fields, either by masking or by shifting back and forth
- New Instructions:
 - Logical: `and, andi, or, ori, nor`
 - Shift: `sll, srl, sra`
- Practice: try writing MIPS functions
- Textbook 2.6