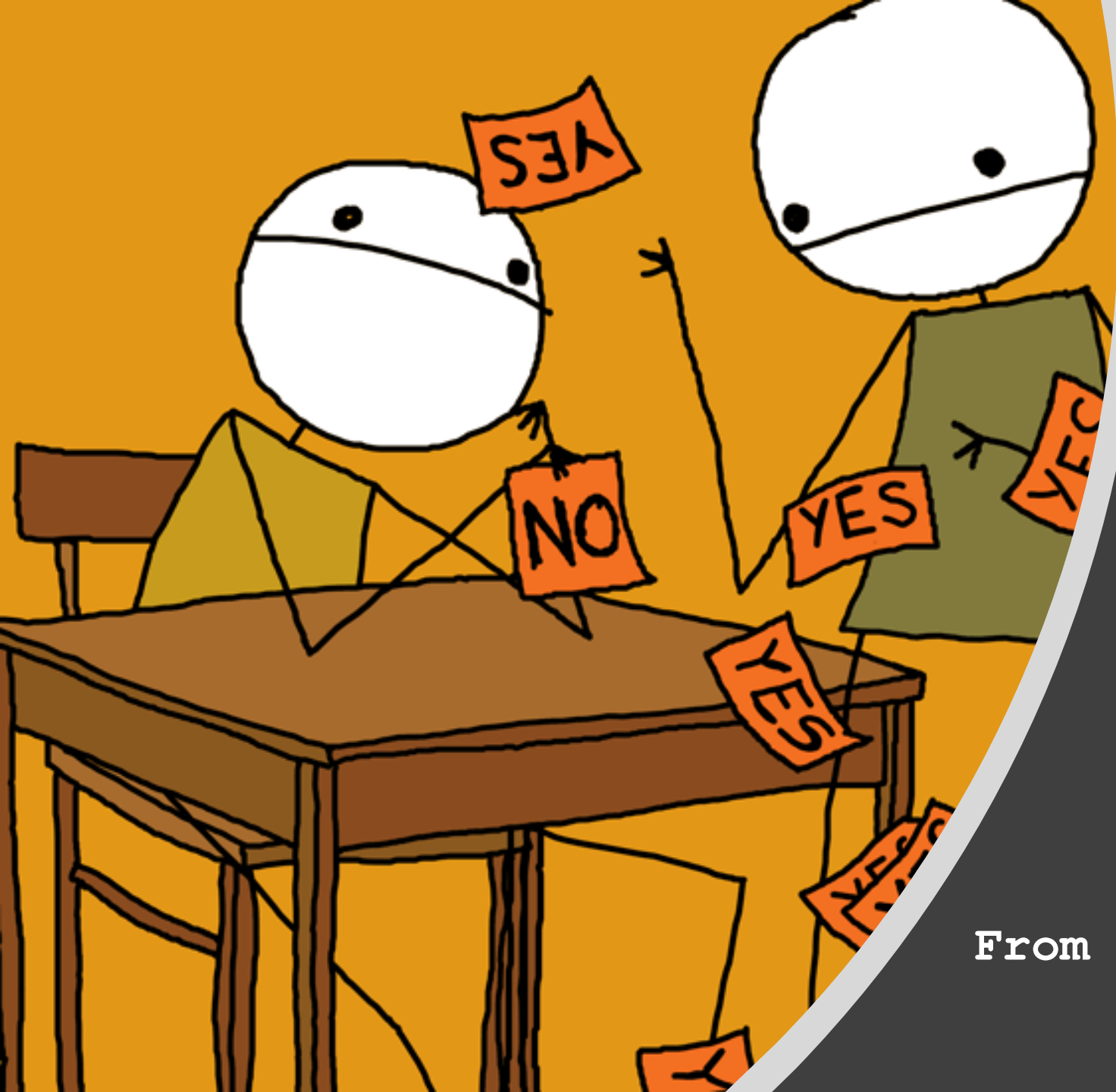


Decisions in MIPS Assembly Language



- All instructions we've seen so far allow us to manipulate data.
- To build a computer we must have the ability to make decisions.



Branches

From if-else/switch to assembly

Conditional Statement in HLL

```
// if-else in C/Java
if (condition) clause
if (condition) {
    clause1
}
else {
    clause2
}
```

labels



```
// C: Rewrite with goto
if (condition) goto L1
clause2
goto L2
L1: clause1
L2:
```

Same meaning in C

No **goto** in Java

Conditional Branches in MIPS

Branch if (registers are) equal: **beq** **reg1**, **reg2**, label

```
// C
if (reg1 == reg2)
    goto label1 ;
```

C to MIPS

```
# MIPS:
# go to label1 if $s1 == $s2
beq $s1 $s2 label1
```

Branch if (registers are) not equal: **bne** **reg1**, **reg2**, label

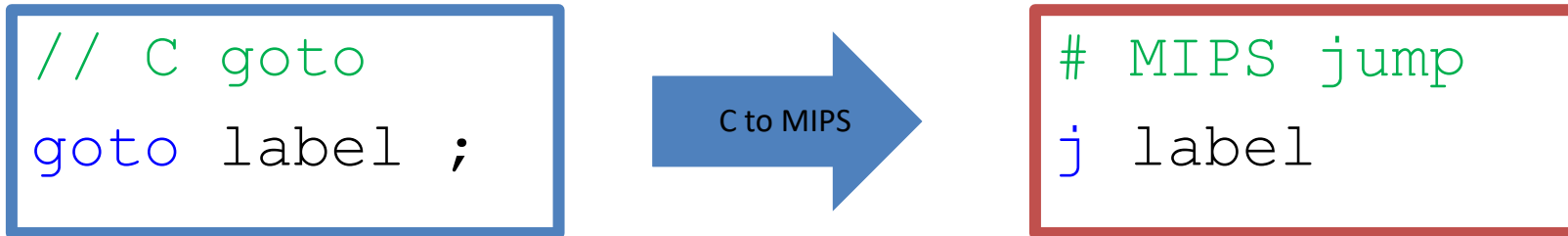
```
// C
if (reg1 != reg2)
    goto label1 ;
```

C to MIPS

```
# MIPS
# go to label1 if $s1 != $s2
bne $s1 $s2 label1
```

Unconditional Branch

- **Jump Instruction:** Jump directly to a label



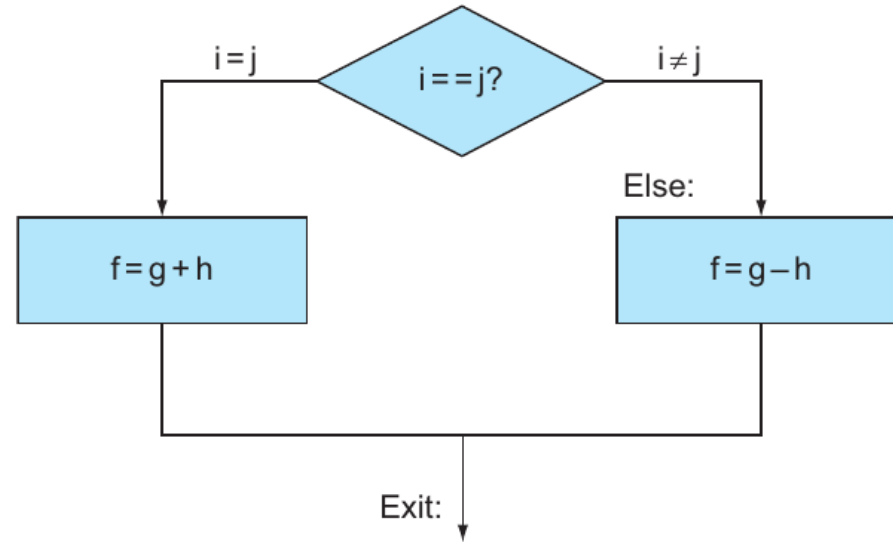
Technically, the following instruction is the same.

There is an important difference. We will see in MIPS representation!

```
# beq version
beq $0, $0, label
```

Conditional Statement in HLL

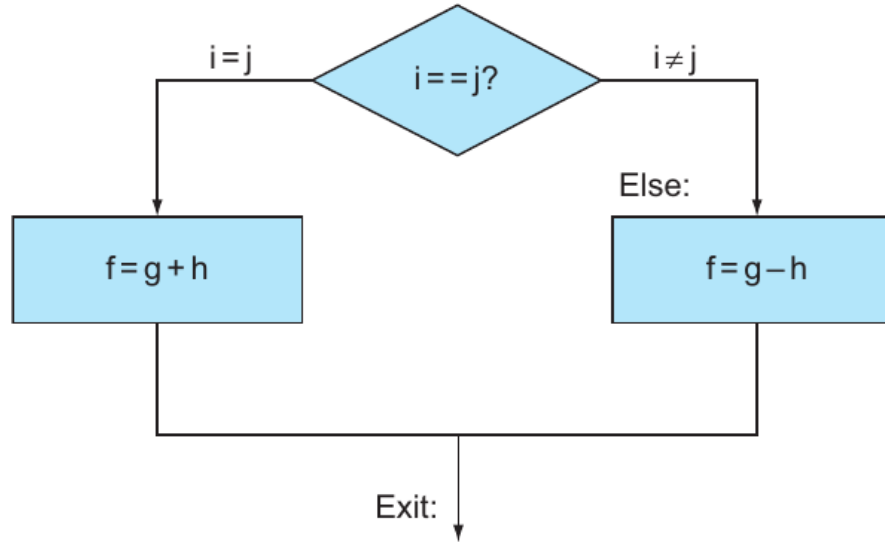
```
// C and Java  
if ( i == j ) {  
    f = g + h ;  
} else {  
    f = g - h ;  
}
```



Compiling *if-else* into MIPS

// C and Java

```
if ( i == j ) {  
    f = g + h ;  
} else {  
    f = g - h ;  
}
```



compiler automatically creates labels to handle decisions (branches).

Registers

\$s0	f
\$s1	g
\$s2	h
\$s3	i
\$s4	j

MIPS

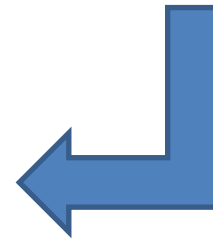
```
beq $s3, $s4, True    # branch i == j  
sub $s0, $s1, $s2     # f = g - h (false)  
j    Exit              # jump to Exit  
True: add $s0, $s1, $s2 # f = g + h (true)  
Exit:
```


The Switch Statement in HLL

Choose among four alternatives depending on whether *k* has the value 0, 1, 2 or 3.

```
// Rewrite it with if-else
if      (k==0) f = i + j ;
else if (k==1) f = g + h ;
else if (k==2) f = g - h ;
else if (k==3) f = i - j ;
```

```
// Switch Statement
switch (k) {
    case 0: f=i+j; break ;
    case 1: f=g+h; break ;
    case 2: f=g-h; break ;
    case 3: f=i-j; break ;
}
```



The switch-statement

- How would you code this in MIPS?

```
        bne    $s5, $0, L1      # if k != 0, branch to L1
        add    $s0, $s3, $s4    # if k == 0, f = i + j
        j      Exit            # end of case, exit
L1:     addi    $t0, $s5, -1      # $t0 = k-1
        bne    $t0, $0, L2      # if k != 1, branch to L2
        add    $s0, $s1, $s2    # if k == 1, f = g + h
        j      Exit            # end of case, exit
L2:     addi    $t0, $s5, -2      # $t0 = k - 2
        bne    $t0, $0, L3      # if k != 2, branch to L3
        sub    $s0, $s1, $s2    # if k == 2, f = g - h
        j      Exit            # end of case, exit
L3:     addi    $t0, $s5, -3      # $t0 = k - 3
        bne    $t0, $0, Exit    # if k != 3, branch to Exit
        sub    $s0, $s3, $s4    # if k == 3, f = i - j
Exit:
```



Loops

Q: How did the programmer die in the shower?

A: He read the shampoo bottle instructions: Lather. Rinse. Repeat.



Loops in C and Assembly

HLL has three types of loops: while, do-while, for. Each can be rewritten as the other



MIPS: There are multiple ways to write a loop with conditional branch

Loops in HLL: 3 ways

Example: Sum of Series

$\text{sum} = 1 + 2 + 3 + 4 + 5$

```
// while
int i = 1 ;
int N = 5 ;
int sum = 0 ;

while ( i<=N ) {
    sum += i ;
    i++ ;
}
```

```
// for
int i = 1 ;
int N = 5 ;
int sum = 0 ;

for (i=1 ; i<=N ; i++)
    sum += i ;
```

```
// do-while
int i = 1 ;
int N = 5 ;
int sum = 0 ;

do {
    sum += i ;
    i++ ;
} while (i<=N) ;
```

From do-while to goto

Example: Sum of Series

$\text{sum} = 1 + 2 + 3 + 4 + 5$

```
int i = 1 ;
int N = 5 ;
int sum = 0 ;
// do-while loop in C
do {
    sum = sum + i ;
    i = i + 1 ;
} while ( i != N ) ;
```

do-while to goto

```
int i = 1 ;
int N = 5 ;
int sum = 0 ;
// Rewrite it with goto in C
Loop: sum = sum + i ;
      i = i + 1 ;
if ( i != N )
    goto Loop ;
```

From **do-while** to MIPS assembly

// do-while loop in C

```
do {  
    sum = sum + i ;  
    i = i + 1 ;  
} while ( i != N ) ;
```

do-while to goto

// Rewrite it with goto in C

```
Loop: sum = sum + i ;  
      i = i + 1 ;  
if ( i != N )  
    goto Loop ;
```

Registers

\$s1	i
\$s2	N
\$s3	sum

MIPS code

```
Loop: add  $s3, $s3, $s1  # sum = sum + i  
      addi $s1, $s1, 1    # i = i + 1  
      bne  $s1, $s2, Loop # go to Loop if i != N
```

Inequalities

So far, we only test equalities. What about inequalities?

Inequalities in MIPS

- `beq` and `bne` only tested equalities

```
if ( i == j )  
if ( i != j )
```

C to MIPS

```
beq $s1 $s2 label1  
bne $s1 $s2 label1
```

- We need to test `<`, `<=`, `>`, `>=`

```
if ( i < j )  
if ( i <= j )  
if ( i > j )  
if ( i >= j )
```

C to MIPS

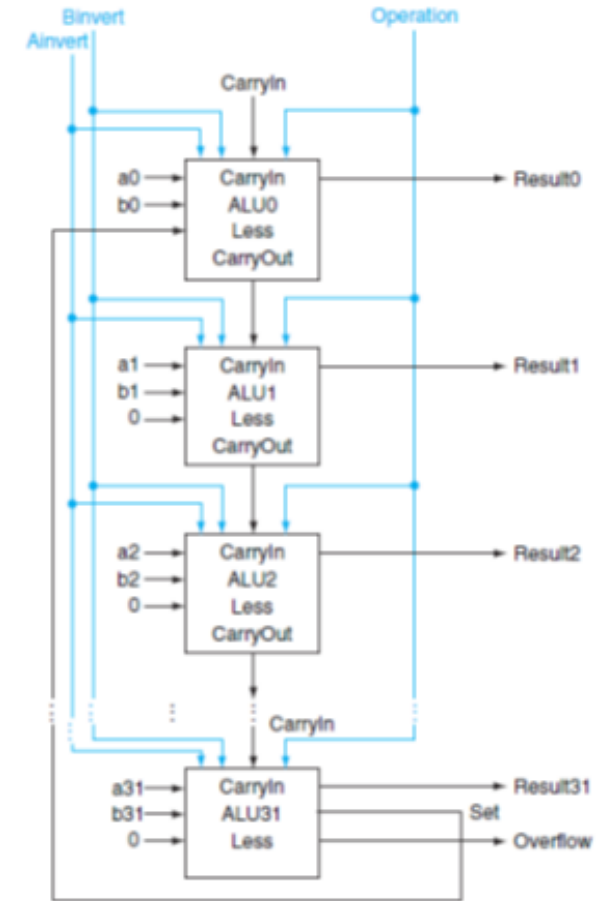
Inequalities in MIPS: `slt`

- Syntax:

```
slt reg1, reg2, reg3
```

- Compare `reg2` and `reg3`
- Place the result in `reg1`

```
// HLL style
if ( reg2 < reg3 )
    reg1 = 1 ;
else
    reg1 = 0 ;
```



Remember "**S**et on **L**ess **T**han" From ALU?

Inequalities in MIPS: from *goto* to MIPS



```
// C
if ( g < h )
    goto Less ;
```

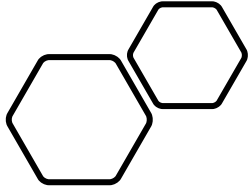
Registers

\$s0	g
\$s1	h
\$t0	

```
# MIPS: branch to Less if $s0 < $s1
slt $t0, $s0, $s1    # if $s0 < $s1 (g < h), $t0 = 1
bne $t0, $0, Less    # branch if $t0 != 0
```

\$0 always contains 0

bne and **beq** often use it for comparison after an **slt** instruction.



Inequalities in MIPS

We have now seen slt for $<$, what about $>$, \leq and \geq ?



MIPS philosophy: **Simpler is Better!** Can we implement them using just slt and beq/bne

Four Combinations of `slt` and `beq/bne`

```
slt $t0, $s0, $s1    # $t0 = 1 if $s0 < $s1 (g < h)
bne $t0, $0, Less     # if $t0 != 0, goto Less ( g < h )
```

```
slt $t0, $s0, $s1    # $t0 = 1 if $s0 < $s1 (g < h)
beq $t0, $0, Geq      # if $t0 == 0, goto Geq ( g >= h )
```

```
slt $t0, $s1, $s0    # $t0 = 1 if $s1 < $s0 (h < g)
bne $t0, $0, Gtr      # if $t0 != 0 goto Gtr ( g > h )
```

```
slt $t0, $s1, $s0    # $t0 = 1 if $s1 < $s0 (h < g)
beq $t0, $0, Leq      # if $t0 == 0, goto Leq ( g <= h )
```

Pseudo-instructions for Inequalities

Too complicated? Good News!

MARS translates pseudo-instructions into MIPS instructions

PSEUDOINSTRUCTION SET

NAME	MNEMONIC	OPERATION
Branch Less Than	blt	if($R[rs] < R[rt]$) PC = Label
Branch Greater Than	bgt	if($R[rs] > R[rt]$) PC = Label
Branch Less Than or Equal	ble	if($R[rs] \leq R[rt]$) PC = Label
Branch Greater Than or Equal	bge	if($R[rs] \geq R[rt]$) PC = Label
Load Immediate	li	$R[rd] = \text{immediate}$
Move	move	$R[rd] = R[rs]$



Inequalities with Immediates

Immediates in Inequalities

- **Syntax:**

slti **Result** **Source** **Immediate**

- Result = 1 if Source < Immediate, or 0 otherwise
- **slti** is the immediate version of **slt**

```
// C
if ( g >= 1 )
    goto Loop ;
```

```
# MIPS
slti $t0, $s0, 1      # $t0 = 1 if $s0 < 1
beq  $t0, $0, Loop    # goto Loop if $t0 == 0
```


Unsigned Immediates in Inequalities

- Syntax:

sltu **Result** **Source1** **Source2**

sltiu **Result** **Source** **Immediate**

- Set result to 1 or 0 depending on unsigned comparisons

```
# MIPS
sltu  $t0, $s0, $s1    # $t0 = 1 if $s0 < $s1
sltiu $t0, $s0, 5      # $t0 = 1 if $s0 < 5
```

Immediates in Inequalities



```
slt    $t0, $s0, $s1  
sltu   $t1, $s0, $s1
```

Assume

$\$s0 = 0xFFFF\ FFFA$

$\$s1 = 0x\ 0000\ FFFA$

What is value of \$t0, \$t1?

Review and More Information

- High-level languages
 - Conditional statement: `if-else`, `switch`
 - Loop: `while`, `do-while`, `for`
- MIPS uses **conditional branches**:
 - Equality: `beq`, `bne`
 - Inequality: `slt`, `slti`, `sltu`, `sltiu`
 - Jump: `j`
- Textbook Section 2.7
- **Try it out in MARS**