

## Binary addition & subtraction

\* One advantage of 1's complement & 2's complement no. system is that subtraction is just an addition. No extra subtractor ckt is required. If you have designed an adder, that is sufficient for subtraction as well.

\* Addition we have talked about.

$$\begin{array}{r} B: 0100 \\ 10010 \\ \hline 01100 \\ \hline 00110 \end{array}$$

$$\begin{array}{r} B: 1100 \\ 00011 \\ \hline 00100 \\ \hline \textcircled{1}1111 \end{array}$$

## Subtraction using 1's complement

\* say  $A - B$

$(-B)$  represent by 1's complement.

$$A + (-B)$$

$$= A + B's \text{ 1's complement}$$

→ If carry is there, add it to result

→ The result is +ve

→ If carry is not there, the result is -ve & is in 1's complement form.

eg.  $6-2$ ,

$$\begin{array}{r} 6: 0110 \\ -2: 1101 \\ \hline 10011 \\ \text{Carry } 1 \rightarrow \\ \hline 0100 \\ = +4 \end{array}$$

$n=4$  (4 bit representation)

$$\begin{array}{r} 2: 0010 \\ -2: 1101 \end{array}$$

(the result is -ve)

eg.  $3-5$

$$\begin{array}{r} 3: 0011 \\ -5: 1010 \\ \hline 1101 \end{array}$$

$$5: 0101$$

$= -2$  in 1's complement form

(the result is -ve)

Ex:

$5-21$  ,  $3-1$  ,  $13-15$

Subtraction using 2's complement

\* say,  $A-B$

$$= A + (-B)$$

$$= A + 2's \text{ comp. of } B$$

\* ignore carry if present, and

\* carry not there, result is -ve, and already in 2's complement form!

eg.

$$\begin{array}{r} 6 : 0110 \\ -2 : 1110 \\ \hline 10100 \end{array}$$

discard ←

$$= +4$$

$$\begin{array}{r} 2 : 0010 \\ 2's \text{ com. of } 2 \\ : 11011 \\ = 1110 \end{array}$$

pos. no. since carry is coming out.

eg. 3-5

$$\begin{array}{r} 3 : 0011 \\ -5 : 1011 \\ \hline 1110 \end{array}$$

$$= 2's \text{ comp. of } -2$$

neg. no. since carry is not there

**try!**

$$3-13, \quad 4-12, \quad 16-14$$

\* Overflow happens when both the no.s are having same sign.

2's complement

$$\begin{array}{r} +6 : 0110 \\ +5 : 0101 \\ \hline 1011 \end{array}$$

MSB is 1 : -ve

4 bit 2's complement 11 cannot be represented

This is overflow.  
out of Range eg. (-7 to +8) is range.



## BCD and gray code

- \* For less complex mathematical operation, we donot need ~~decimal~~ <sup>binary</sup> decimal-binary conversion. Binary operations are faster.
- \* decimal to binary & vice versa are complex.

### Binary Coded decimal (BCD)

→ In BCD, every decimal digit is represented by 4 bit binary equivalent.

eg.  $238_{10} = 0010 \ 0011 \ 1000$  in BCD

$13_{10} = 0001 \ 0011$  in BCD

$1210_{10} = 0001 \ 0010 \ 0001 \ 0000$

→ Conversion is simple.

→ Six 4-bit combinations are not valid in BCD. those are 1010, 1011, 1100, 1101, 1110, 1111.

Decimal	Binary
0	0000
1	0001
2	0010
3	0011
4	0100
5	0101
6	0110
7	0111
8	1000
9	1001

## Addition of BCD no.!

(14/8/2019)

\* Add 0110 to the

→ if one of the nibbles have invalid combination

→ if there is a carry from previous nibble

eg. 23 + 46

$$\begin{array}{r} 0010 \ 0011 \\ + \ 0100 \ 0110 \\ \hline 0110 \ 1001 \end{array} = 69$$

→ valid BCD.

→ no carry.

} → no need to add 0110

eg. 23 + 48

$$\begin{array}{r} 0010 \ 0011 \\ 0100 \ 1000 \\ \hline 0110 \ 1011 \\ + \ 0110 \\ \hline 0111 \ 0001 \end{array} = 71$$

→ correction step is done at this nibble because the nibble is invalid

→ both nibble invalid means add 0110 to each nibble

try: 23 + 49, 30 + 49, 59 + 99

eg. 28 + 39

$$\begin{array}{r} 0010 \ 1000 \\ 0011 \ 1001 \\ \hline 0110 \ 0001 \\ + \ 0110 \\ \hline 0110 \ 0111 \end{array} = 67$$

## Gray Codes

- Successive bits no.s differ by single bit change.
- Non weighted binary code.
- Also called cyclic code.
- Gray codes are used to reduce dynamic power consumption as it minimizes switching. Also
- reduce error in conversion

Binary	Gray
000	000
001	001
010	011
011	010
100	110
⋮	⋮

### Binary to Gray:

$$\begin{array}{rcl} B: & 0 & 1 & 1 \\ & \downarrow & \oplus & \oplus \\ G: & 0 & 1 & 0 \end{array}$$

$$\begin{array}{rcl} B: & B_2 & B_1 & B_0 \\ G: & G_2 & G_1 & G_0 \end{array}$$

$$G_2 = B_2$$

$$G_1 = B_2 \oplus B_1$$

$$G_0 = B_1 \oplus B_0$$

### Gray to Binary:

$$G: \quad G_2 \quad G_1 \quad G_0$$

$$B: \quad B_2 \quad B_1 \quad B_0$$

$$B_2 = G_2, \quad B_1 = G_1 \oplus B_2, \quad B_0 = G_0 \oplus B_1$$

$$\begin{array}{rcl}
 G: & 0 & 10 \\
 & \downarrow & \oplus \oplus \\
 B: & 0 & 11
 \end{array}$$

try:  $(10110)_2 = (?)_G$

$$(0110)_2 = (?)_G$$

$$(11001)_G = (?)_B$$

$$(011010)_G = (?)_B$$

→ Gray code is called self-reflecting code,  
 eg.  $\begin{array}{cc} 00 \\ 01 \\ 10 \\ 11 \end{array}$  → same thing if you want to represent  
 using 3 bits ↴

0	00
0	01
0	10
0	11
<hr/>	
1	11
1	10
1	01
1	00

mirror image  
 self reflecting

similarly try 3 bit representation to 4 bit representation.



## ASCII code

- American standard code for Information exchange.
- Its a code for representing 128 english characters as numbers, with each letter assigned a no. from 0 to 127.
- Most computers use ASCII code to represent text, which makes it possible to transfer data from one computer to another.
- Standard ASCII character set uses just 7 bits for each character, while some ch. set uses 8 bits.
- Also used in telecommunication.

→	A → 065	a → 097
	B → 066	b → 098
	⋮	⋮
	Z → 090	z → 122

0 → 048  
1 → 049  
9 → 057

& so on.



## EBCDIC

(extended binary coded decimal interchange code)  
8-bit character encoding used mainly on IBM computers.

8 bit means  $(0 - (2^8 - 1))$  possible combinations.

A  $\rightarrow$  193

a  $\rightarrow$  129

0  $\rightarrow$  240

B  $\rightarrow$  194

b  $\rightarrow$  130

1  $\rightarrow$  241

$\vdots$

$\vdots$

$\vdots$