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<b>Experiment No.:</b>	5			
Title:	Fraction Knapsack			
<b>Date of Performance:</b>				
Date of Submission:				
Marks:				
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Vidyavardhini's College of Engineering and Technology

Department of Artificial Intelligence & Data Science

**Experiment No. 5** 

**Title:** Fraction Knapsack

**Aim:** To study and implement Fraction Knapsack Algorithm

**Objective:** To introduce Greedy based algorithms

Theory:

Greedy method or technique is used to solve Optimization problems. A solution that can be

maximized or minimized is called Optimal Solution.

The knapsack problem or rucksack problem is a problem in combinatorial optimization: Given

a set of items, each with a mass and a value, determine the number of each item to include in a

collection so that the total weight is less than or equal to a given limit and the total value is as

large as possible. It derives its name from the problem faced by someone who is constrained

by a fixed size knapsack and must fill it with the most valuable items. The most common

problem being solved is the 0-1 knapsack problem, which restricts the number xi of copies of

each kind of item to zero or one.

In Knapsack problem we are given:1) n objects 2) Knapsack with capacity m, 3) An object i is

associated with profit Wi, 4) An object i is associated with profit Pi, 5) when an object i is

placed in knapsack we get profit Pi Xi.

Here objects can be broken into pieces (Xi Values) The Objective of Knapsack problem is to

maximize the profit.

**Example:** 

In this version of Knapsack problem, items can be broken into smaller pieces. So, the thief

may take only a fraction  $x_i$  of i<sup>th</sup> item.

0**≤**xi**≤**1

The i<sup>th</sup> item contributes the weight xi.wi to the total weight in the knapsack and profit xi.pi to

the total profit.



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	gredy-	fractional -k	napsack (	wci nj	pcing, m)
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1	weight =				
1	TV .	1=2 -> A			
-	1,9	10+40			
1	1.	\$0 ≤60 XCiJ: 2			
	els	10+40			
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					(=3 -> C
					(60-50)/20
-					x 513 : 10/20 = 1/2
1					ateso
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	Jtem	A	ß	C	D
	profit	280	100	120	120
	weight	40	10	20	24
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### Algorithm:

Hence, the objective of this algorithm is to

$$maximize \sum_{n=1}^{n} (x_i. pi)$$

subject to constraint,

$$\sum_{n=1}^n (x_i.wi) \leqslant W$$

It is clear that an optimal solution must fill the knapsack exactly, otherwise we could add a fraction of one of the remaining items and increase the overall profit.

Thus, an optimal solution can be obtained by

$$\sum_{n=1}^n (x_i.\,wi) = W$$

In this context, first we need to sort those items according to the value of  $\frac{p_i}{w_i}$ , so that  $\frac{p_i+1}{w_i+1} \le$ 

 $\frac{p_i}{w_i}$  . Here, **x** is an array to store the fraction of items.

```
Algorithm: Greedy-Fractional-Knapsack (w[1..n], p[1..n], W)
for i = 1 to n
    do x[i] = 0
weight = 0
for i = 1 to n
    if weight + w[i] ≤ W then
        x[i] = 1
        weight = weight + w[i]
else
    x[i] = (W - weight) / w[i]
    weight = W
        break
return x
```



### **Implementation:**

```
#include <stdio.h>
#include <conio.h>
void main()
  int capacity, no_items, cur_weight, item;
  int used[10];
  float total_profit;
  int i;
  int weight[10];
  int value[10];
  clrscr();
  printf("Enter the capacity of knapsack:\n");
  scanf("%d", &capacity);
  printf("Enter the number of items:\n");
  scanf("%d", &no_items);
  printf("Enter the weight and value of %d item:\n", no_items);
  for (i = 0; i < no_items; i++)
   {
       printf("Weight[%d]:\t", i);
       scanf("%d", &weight[i]);
       printf("Value[%d]:\t", i);
       scanf("%d", &value[i]);
  }
  for (i = 0; i < no_items; ++i)
       used[i] = 0;
  cur_weight = capacity;
  while (cur\_weight > 0)
```



```
item = -1;
       for (i = 0; i < no_items; ++i)
          if ((used[i] == 0) \&\&
          ((item == -1) || ((float) value[i] / weight[i] > (float) value[item] / weight[item])))
          item = i;
     used[item] = 1;
     cur_weight -= weight[item];
     total_profit += value[item];
     if (cur\_weight >= 0)
       printf("Added object %d (%d Rs., %dKg) completely in the bag. Space left: %d.\n",
item + 1, value[item], weight[item], cur_weight);
     else
     {
       int item_percent = (int) ((1 + (float) cur_weight / weight[item]) * 100);
       printf("Added %d%% (%d Rs., %dKg) of object %d in the bag.\n", item_percent,
value[item], weight[item], item + 1);
       total_profit -= value[item];
       total_profit += (1 + (float)cur_weight / weight[item]) * value[item];
     }
   }
  printf("Filled the bag with objects worth %.2f Rs.\n", total_profit);
}
```



#### **Output:**

```
Output
Enter the capacity of knapsack:
Enter the number of items:
Enter the weight and value of 3 item:
Weight[0]:
                10
Value[0]:
                100
                15
Value[1]:
                50
Weight[2]:
                5
Value[2]:
                50
Added object 1 (100 Rs., 10Kg) completely in the bag. Space left: 10.
Added object 3 (50 Rs., 5Kg) completely in the bag. Space left: 5.
```

**Conclusion:** experiment successfully implemented the fractional knapsack algorithm, efficiently allocating items based on their values and weights. By prioritizing fractional solutions, we optimized resource utilization, demonstrating the algorithm's practicality and effectiveness in real-world scenarios.