

Homework IV

Due on Nov. 14, 2021

Justify your answers with proper reasonings/proofs.

1. Suppose you are given an array $M[1\dots n, 1\dots n]$ of numbers, which may be positive, negative, or zero, and which are not necessarily integers. Describe an algorithm to find the largest sum of elements in any rectangular subarray of the form $M[i\dots i', j\dots j']$. Your algorithm should run in $O(n^3)$ time.
2. For a sequence of n days, you are given subsets S_1, S_2, \dots, S_n of $\{1, 2, \dots, k\}$. Think of S_i as the subset of people who are available for work on day i . You need to pick exactly one person from S_i for each of the days $i = 1, \dots, n$. For a person j , let Δ_j denote the $\sum_{i: j \in S_i} \frac{1}{|S_i|}$. This is the expected number of times j would be picked if we pick a random person from S_i on each of days $i = 1, \dots, n$. A selection of persons, one from each set S_i , is said to be good if each person j is picked at most $\lceil \Delta_j \rceil$ times. Show that such a selection is always possible, and give an efficient algorithm to find such a selection.
3. Let (u, v) be a directed edge in arbitrary flow network G . Prove that if there is a minimum (s, t) -cut (S, T) such that $u \in S$ and $v \in T$, then there is no minimum cut (S', T') such that $u \in T', v \in S'$. Note that by definition of cut, $s \in S, t \in T$, and similarly $s \in S', t \in T'$.
4. Let G be an undirected graph and s and t be two special vertices in it. Give an efficient algorithms to find the maximum number of node disjoint paths from s to t (a set of paths from s to t are said to be node-disjoint if no two of them share a vertex other than s or t).
5. Let G be an undirected graph. For a subset S of vertices, let $e(S)$ denote the number of edges which have both the end-points in S . Given a rational number α , we would like to find out if there is a subset S of vertices such that $\frac{e(S)}{|S|} \geq \alpha$. Give an efficient algorithm to solve this problem.