Hashing → technique used for searching purpose.

- Linear Search → TC: O(N)
- Binary Search → TC: $O(\log_2 N)$

TC: O(1) → Searching.

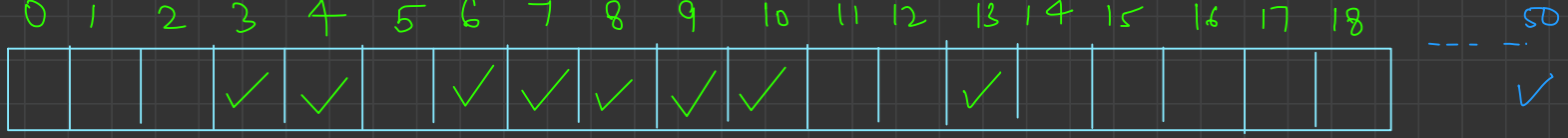
eg: 8, 3, 13, 6, 4, 10, 9, 7, 50 (user input)

Key = 13

search(int val)

→ TC: O(1)

Boolean arr[]



a lot of empty space

Memory wastage!

To overcome issue, hashing was introduced.

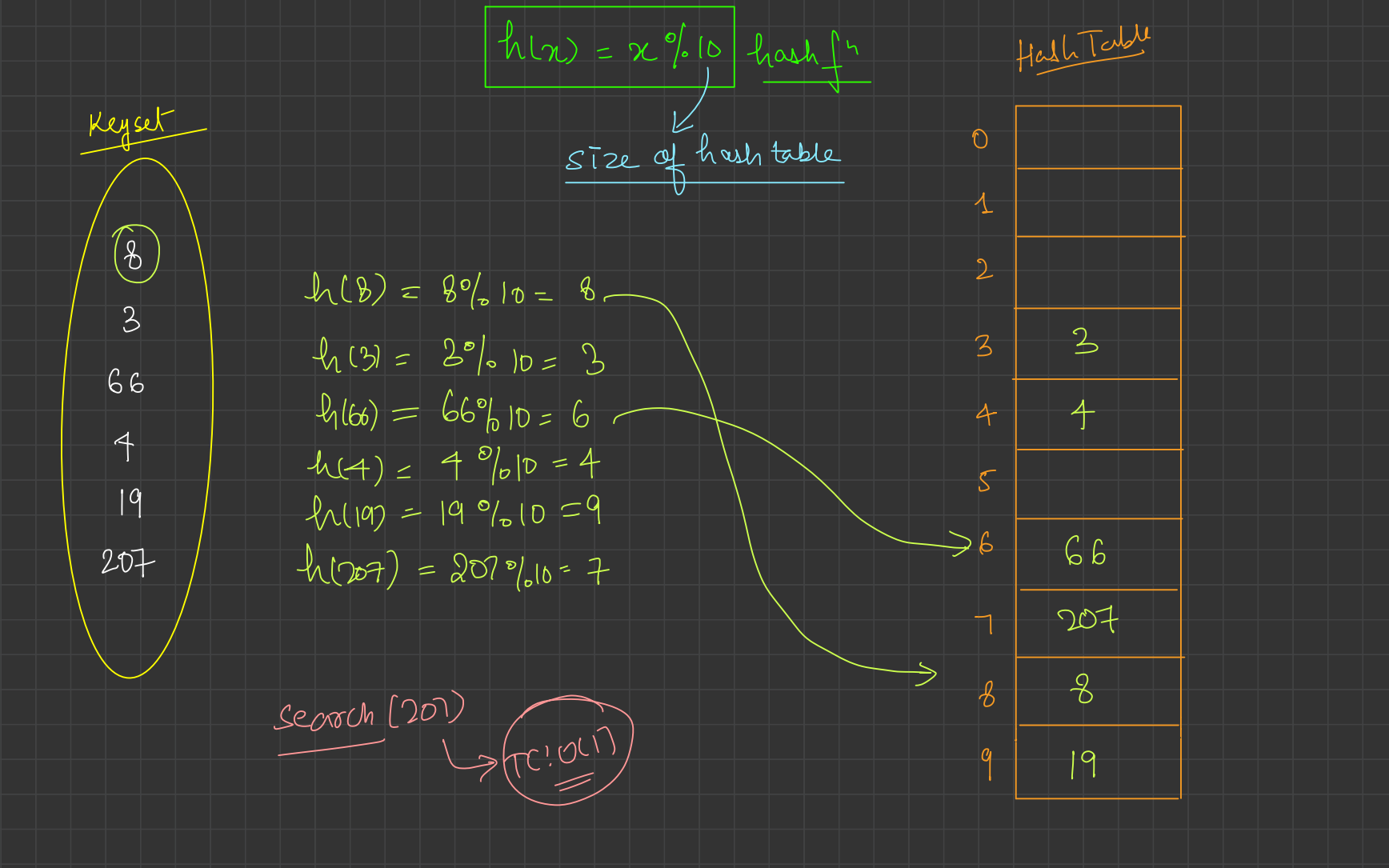
TC: O(1)

```
Search(13)
{
    if (arr[13] == true)
        return true;
    else
        return false;
}
```

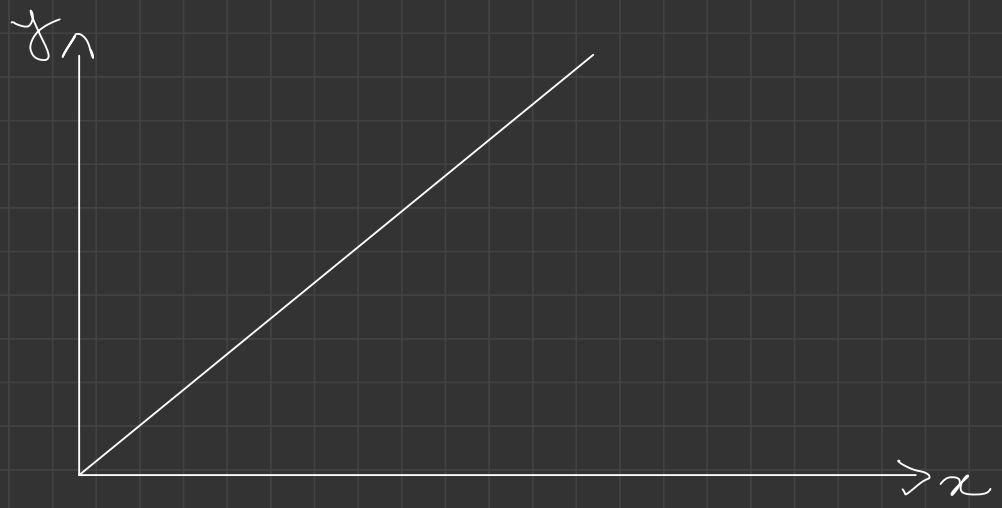
hashing fn (2 step process)

Step 1 $y(x) = k$ → integer value

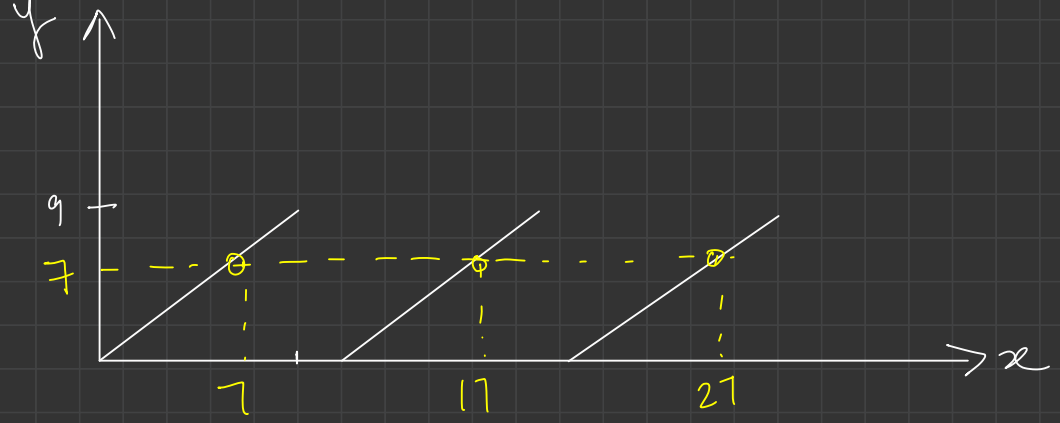
Step 2 $h(k) = y$ → hashed value


h(x) = x % 10
hash fn
size of hash table
Hash Table
Keyset
8
3
66
4
19
207
h(8) = 8 % 10 = 8
h(3) = 3 % 10 = 3
h(66) = 66 % 10 = 6
h(4) = 4 % 10 = 4
h(19) = 19 % 10 = 9
h(207) = 207 % 10 = 7
0
1
2
3
3
4
4
5
6
66
7
207
8
8
9
19
search (207)
TC: O(1)

One to One Mapping



Many to one



$h(x) = x \% 10$ hash fn

Keyset

8
3
66
4
19
207
33

$h(8) = 8 \% 10 = 8$

$h(3) = 3 \% 10 = 3$

$h(66) = 66 \% 10 = 6$

$h(4) = 4 \% 10 = 4$

$h(19) = 19 \% 10 = 9$

$h(207) = 207 \% 10 = 7$

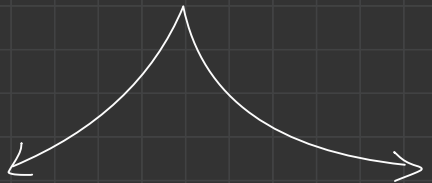
$h(33) = 33 \% 10 = 3$

Hash Table

0	
1	
2	
3	3
4	4
5	
6	66
7	207
8	8
9	19

Collision

Methods to Remove Collision



Closed hashing

- linear Probing
- quadratic Probing

Open hashing

- chaining

Chaining

$h(x) = x \% 10$ hash fn

AL<LL> Hash Table

Keyset

8, 3, 66, 4, 19, 207, 33, 23

$h(8) = 8 \% 10 = 8$

$h(3) = 3 \% 10 = 3$

$h(66) = 66 \% 10 = 6$

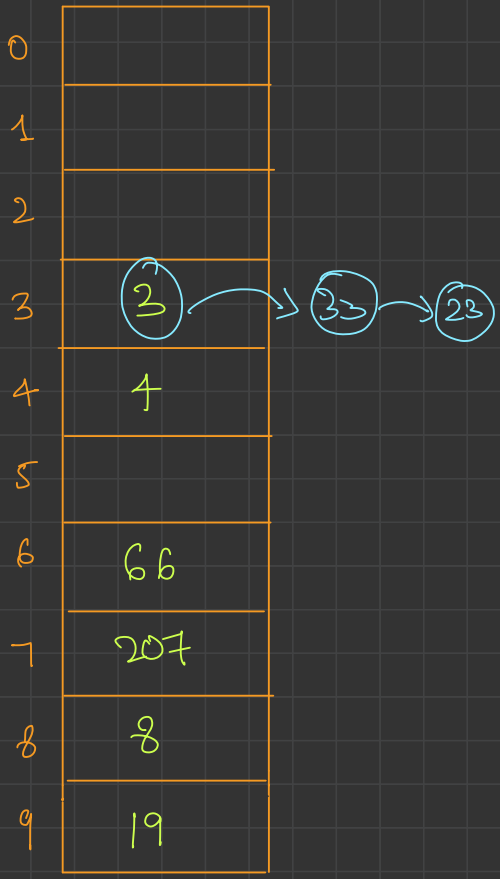
$h(4) = 4 \% 10 = 4$

$h(19) = 19 \% 10 = 9$

$h(207) = 207 \% 10 = 7$

$h(33) = 33 \% 10 = 3$

$h(23) = 3$



Searching (33) → 3

hash fn

$h(x) = x \% \text{loadfactor}$

↓

75% of range of values given

0 – 100

↳ lf = 75

Linear Probing

$$h'(x) = [h(x) + f(i)] \% \text{size} \quad (\text{size} \to 10)$$

$h(x) = x \% 10$

$f(i) = i, \quad i \in R[0, \infty)$

Keyset: 8, 3, 66, 33, 4, 23

$h'(8) = [h(8) + f(0)] \% 10$

$= (8 + 0) \% 10 = 8$

$h'(3) = [h(3) + f(0)] \% 10$

$= (3 + 0) \% 10 = 3$

$h'(66) = [h(66) + f(0)] \% 10$

$= (6 + 0) \% 10 = 6$

$h'(33) = [h(33) + f(0)] \% 10 = (3 + 0) \% 10 = 3$

$h'(33) = [h(33) + f(1)] \% 10 = (3 + 1) \% 10 = 4$

$h'(4) = [h(4) + f(0)] \% 10 = (4 + 0) \% 10 = 4$

fail $\longrightarrow$ 5

Hash Table

0	
1	
2	
3	3
4	33
5	4
6	66
7	23
8	8
9	

Linear hashing

↳ disadv : clustering issue !

Quad Probing

$$h_1'(n) = \left[h_1(n) + f(i) \right] \% \text{ size}$$

$h_1(n)$ - n % size

$f(i) = i^2$, $i \in R$

2 - Data Structures

1. HashSet
2. HashMap

Based On Hash Theory

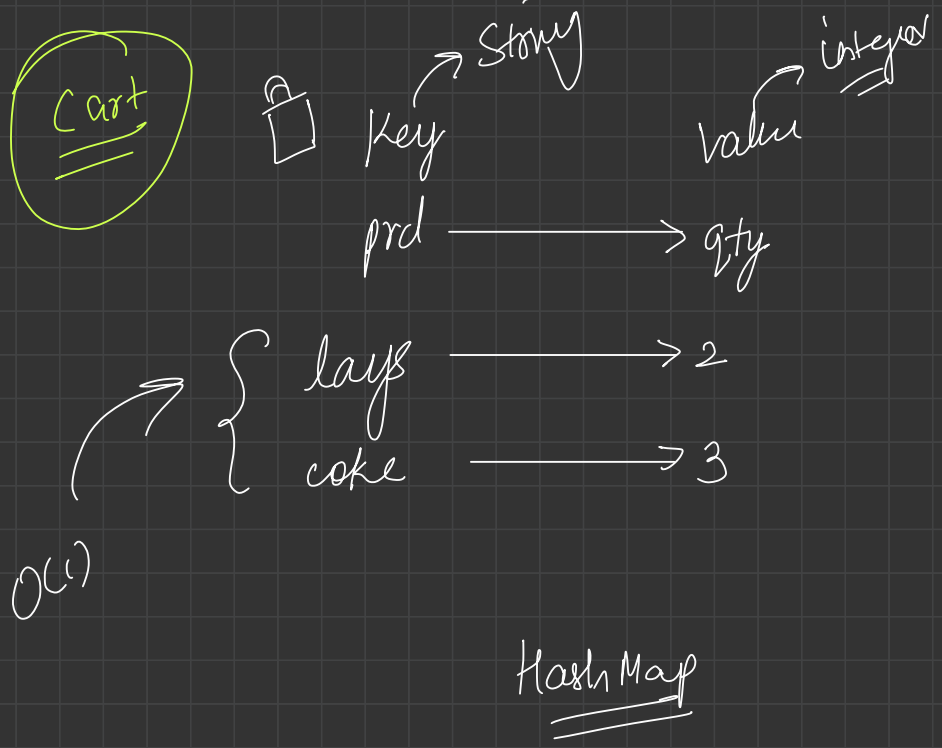
TC: O(1)

Searching

HashSet : Stores unique entity

HashMap

↳ <key, value>



HashMap

TreeMap
TreeSet
→ here key are in ordered format

Searching
↳ T.C: O(log N)

based on
Red Black Tree
↳ self Balancing Tree