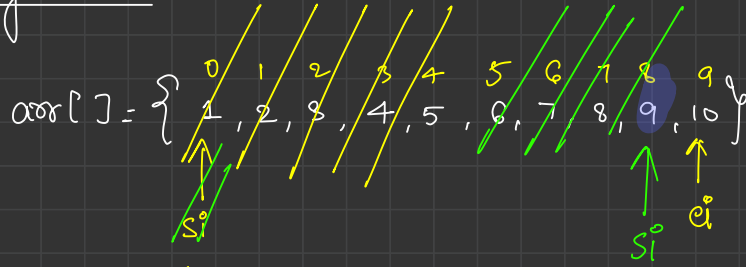
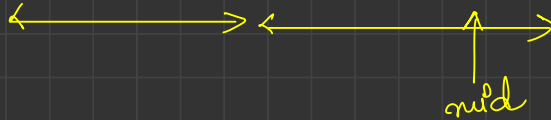




# Binary Search



target = 9



length of array

$$\begin{cases} TC: O(\log N) \\ SC: O(1) \end{cases}$$

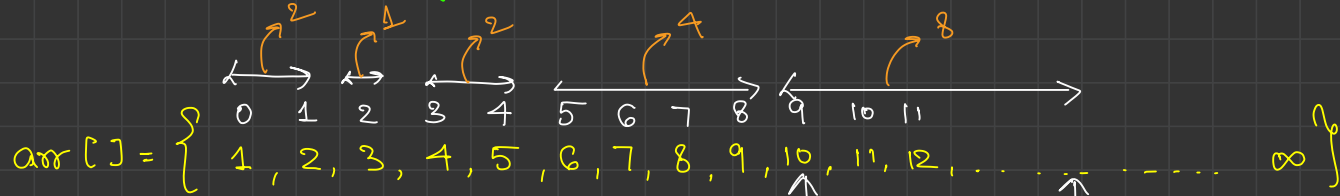
- ① define range when target can be find.
- ② get the mid point
- ③ take a decision to eliminate half of the range
- ④ and move towards the part where your answer lies.

# Binary Search over an infinite sorted array

target = 7

int[] arr = { 0, 1, 2, 3, 4, 5, 6, 7, ... }  $\infty$

NOTE: We don't know length of array, but all the ele. are sorted in inc. order



target = 10

$si = 0, ei = 1 \rightarrow [1, 2]$

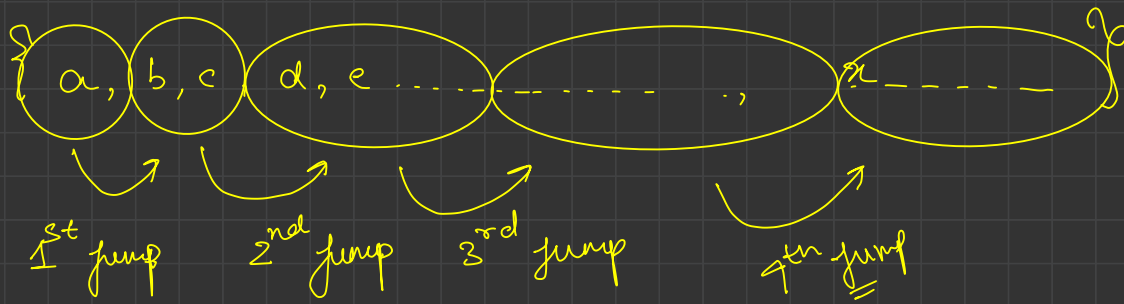
$si = 2, ei = 2 \rightarrow [3, 3]$   
 $si = 3, ei = 4 \rightarrow [4, 5]$   
 $si = 5, ei = 8$

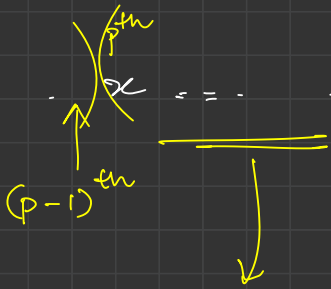
$\left. \begin{array}{l} si = ei + 1; \\ ei = 2 * ei; \end{array} \right\}$

$si = 9, ei = 16$

All the ele from  $si$  to  $ei$  are sorted.

Range  $[arr[si], arr[ei]]$

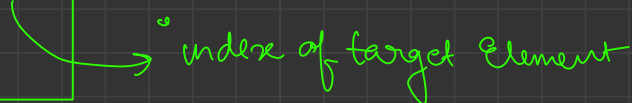


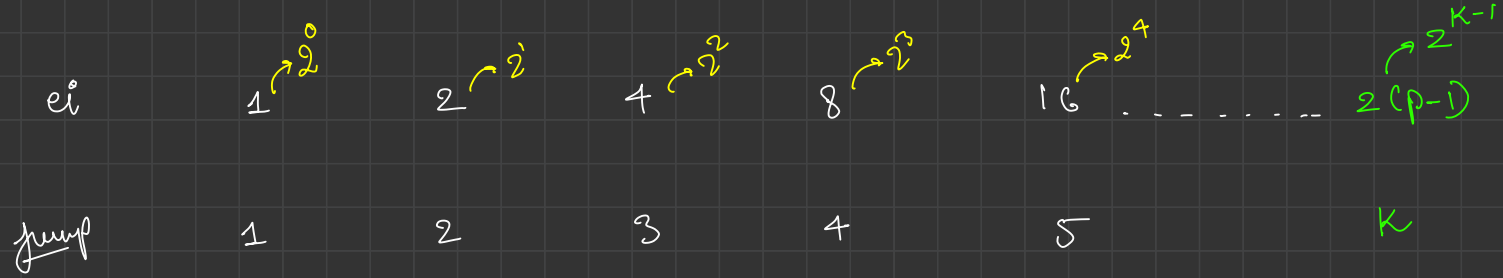
arr[] = { 0 1 2 3  
1, 2, 3, 4, 5, 6, ... .. }  


$$s_i = p$$

$$e_i = 2 \times (p-1)$$

$$\begin{aligned} \text{Number of Elements in this range} &= e_i - s_i + 1 \\ &= 2p - 2 - p + 1 \\ &= (p - 1) \rightarrow \underline{\text{Elements}} \end{aligned}$$

TC:  $O(\log_{f_2} P)$   
 of BS.  index of target element



$$2^{(p-1)} = 2^{K-1}$$

$$(p-1) = \frac{2^{K-1}}{2}$$

$$(p-1) = 2^{K-2}$$

take  $\log_2$  Both Sides

$$\log_2 (p-1) = (K-2) \log_2 2 \rightarrow 1$$

$$\log_2 (p-1) + 2 = K$$

TC for K jumps

$$= O(K)$$

$$\approx O(\log_2 P)$$

TC: to find range =  $O(\log_2 l)$

TC: to find in range =  $O(\log_2 l)$

Total TC: =  $2O(\log_2 l)$

$= O(\log_2 l)$

→ index of target  
element

```
int find (int[] arr, int target)
```

```
{
```

```
    if (target < arr[0])  
        return -1;
```

```
    int si = 0;
```

```
    int ei = 1;
```

```
    while (arr[ei] < target)
```

```
{
```

```
        si = ei + 1;
```

```
        ei = 2 * ei;
```

```
}
```

```
    int index = BinarySearch(arr, si, ei, target);
```

```
    return index;
```

```
}
```

TC:  $O(\log_2 l)$

SC:  $O(1)$



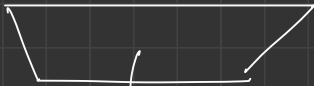
# Capacity to Ship Packages within B days

$A[] = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$

days = 5

1 day

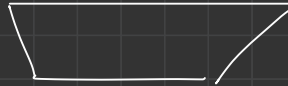
$\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$



55kg

10 days

$\{1\}$



22kg

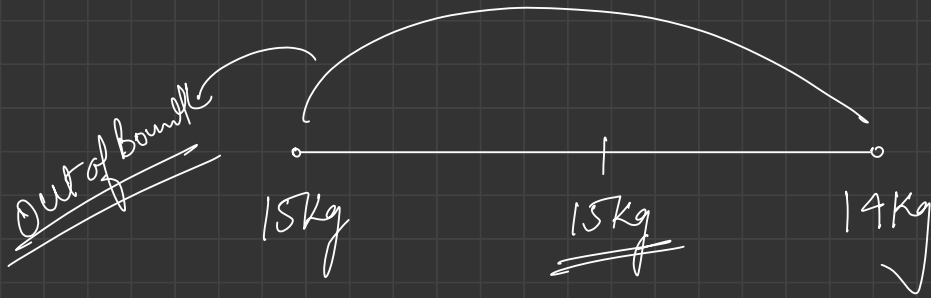


⋮

$\{10\}$



10kg



~~ans = 32~~  
~~20~~  
~~15~~  
15

$A[] = \{ 1, 2, 3, 4, 5, 6, 7, 8, 9, 10 \}$   
 $\uparrow \uparrow \uparrow \uparrow \uparrow \uparrow \uparrow \uparrow \uparrow \uparrow$

B = 5 days

capacity = 15 Kg

day 1  $\rightarrow 1 + 2 + 3 + 4 + 5$

day 2  $\rightarrow 6 + 7$

day 3  $\rightarrow 8$

day 4  $\rightarrow 9$

day 5  $\rightarrow 10$

}

## Minimum limit of balls in a bag

$$\text{arr}[ ] = \{ 2, \textcircled{4}, 8, 2 \}$$

$$\text{maxOpt} = \cancel{4} \cancel{3} \cancel{2} 1$$

$$\rightarrow (\overset{\checkmark}{1}, 3) (2, 2)$$

$$\{ 2, 1, 3, \textcircled{8}, 2 \}$$

$$\rightarrow (4, 4) (\overset{\checkmark}{3}, 5) (2, 6) (1, 7)$$

$$\{ 2, 1, 3, 3, 5, 2 \}$$

$$\rightarrow (2, 3) (\overset{\checkmark}{1}, 4)$$

$$\{ 2, 1, 3, 3, 1, \textcircled{4}, 2 \}$$

$$\rightarrow (1, 3) (\overset{\checkmark}{2}, 2)$$

$$\{ 2, 1, 3, 3, 1, 2, 2, 2 \} \rightarrow \underline{\underline{\text{max} = 3}}$$



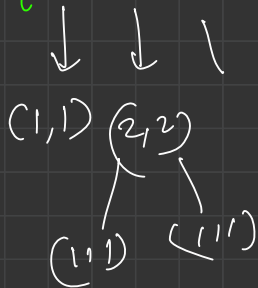
$$\text{pairs} = \cancel{2}$$



$$\underline{\underline{\text{opt} = 4}}$$

$$\text{arr}[] = \{2, 4, 8, 2\}$$

$$\text{maxBalls} = 1$$



$$\text{opt} = \cancel{7} \cancel{3} 4$$

$\{2, (4), 8, 2\}$

$\downarrow$   
 $(2, 2)$

maxDpt = 4

$\{2, 2, 2, (8), 2\}$

$\downarrow$   
 $(4, 4)$

$\{2, 2, 2, (4), 4, 2\}$

$\downarrow$   
 $(2, 2)$

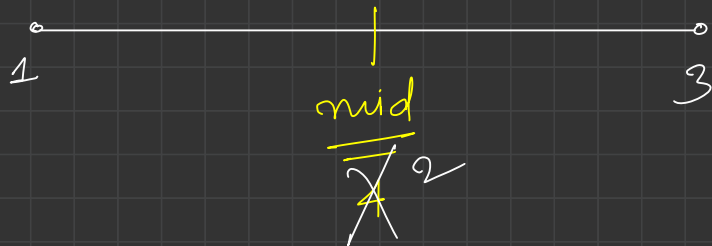
$\{2, 2, 2, 2, 2, 2, 2, 2\}$

$\swarrow$   
penalty = 2

$\{2, 2, 2, 2, 2, (4), 2\}$

$\downarrow$   
 $(2, 2)$

$\{2, 4, 8, 2\}$



parts = 4

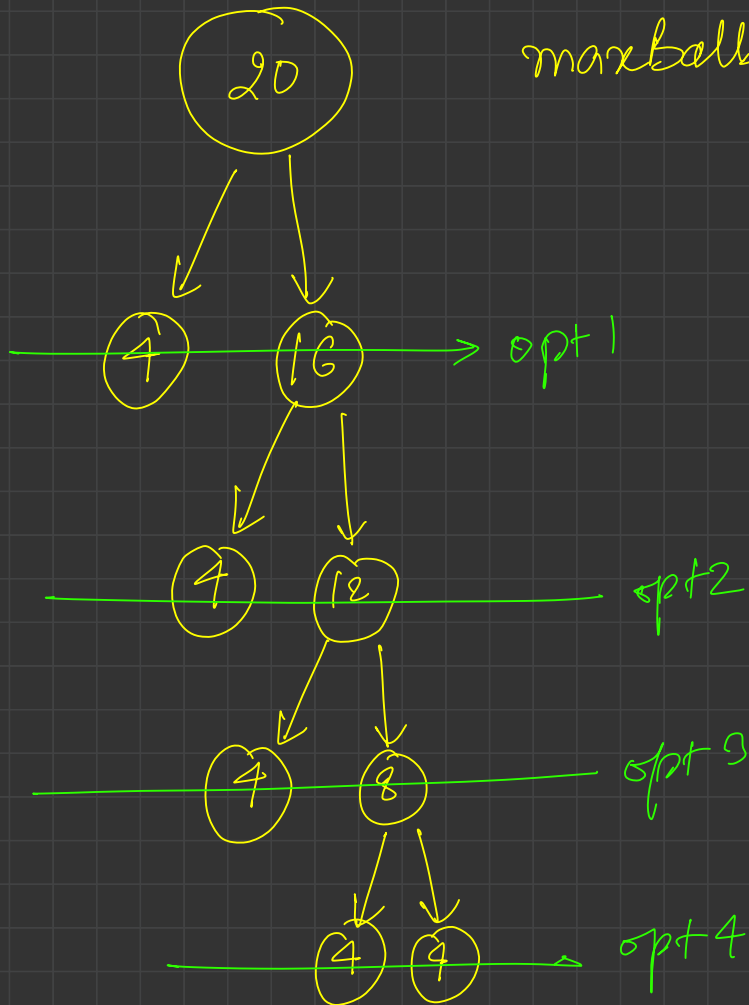
$\{2, 4, 8, 2\}$   
Four arrows point upwards from below the numbers 2, 4, 8, and 2 respectively.

max balls = 4

opt = 1

$$\left( \frac{8}{4} - 1 \right) = \underline{\underline{1}}$$

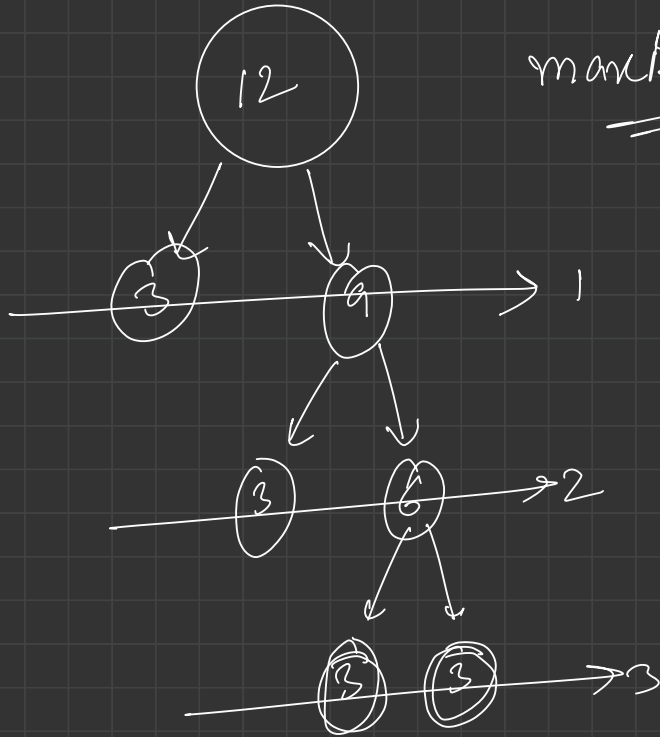
maxBalls = 4



opt

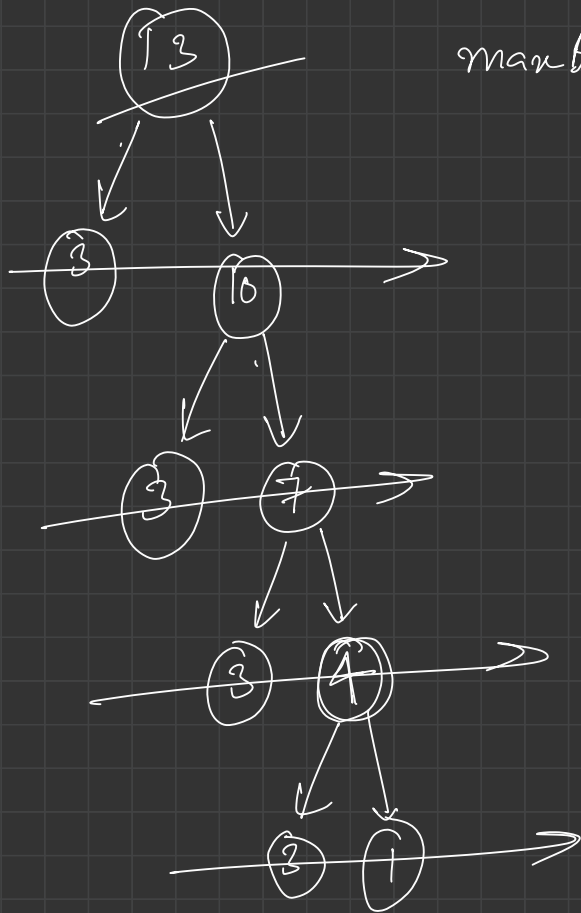
$$\left( \frac{\text{Balls In Bag}}{\text{max Balls}} - 1 \right)$$





maxballs = 3

$$\left( \frac{12}{3} - 1 \right) \text{opt}$$



max balls = 3

Number of opt

if (Balls % maxballs == 0)

{ opt =  $\frac{\text{Balls}}{\text{maxballs}} - 1$  ; }

else

{ opt =  $\frac{\text{Balls}}{\text{maxballs}}$  ; }

