



Hashing

It is a technique for searching purpose

① Linear Search

3	10	1	5	8	7	12	-10
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TC: $O(N)$ \leadsto Searching

② Binary Search \rightarrow {array should be sorted}

{ TC: $O(\log n)$ \leadsto Searching }

③ Hashing \rightarrow {TC: O(1) Searching}

{ ~~8~~, ~~3~~, ~~13~~, ~~6~~, ~~4~~, ~~10~~ }, 50

target = 4

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	...	50
A				3	4		6		8		10			13		...	50

search(4)

key

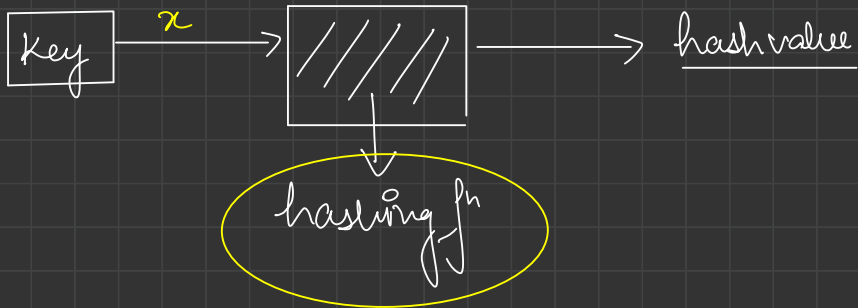
TC: O(1)

```
if (A[key] != null)
{
    return true;
}
else
    return false;
```

memory wasted

disadv

↳ high memory is used.
↳ hashing is introduced.



step1

$$g(x) = k$$

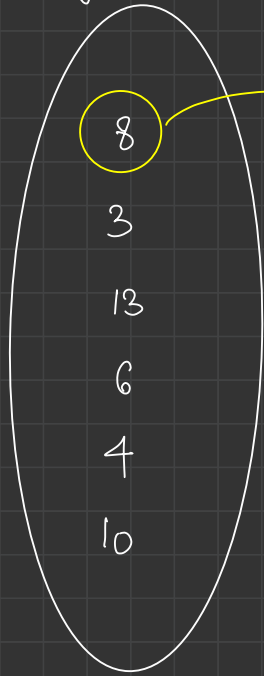
Two arrows point down from x and k to the words "key" and "integer value" respectively. A large curly brace is on the right side of these two words.

step2

$$h(k) = y$$

An arrow points down from k to the text "hashing fn". An arrow points from y to an oval containing the text "hash value".

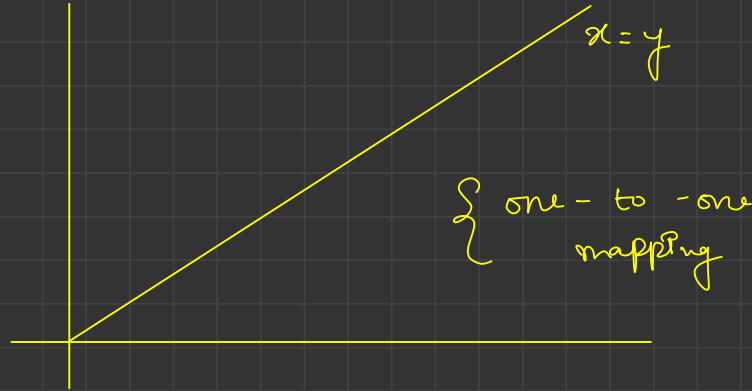
Key Space



hashing fⁿ $h(K) = y$ hash value →

$$\underline{h(x) = x}$$

$$h(8) = \textcircled{8} \rightarrow \underline{\text{hash value}}$$



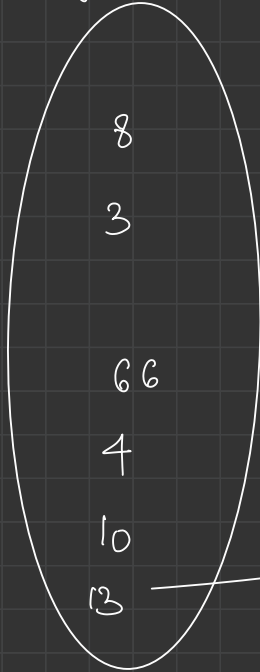
hash table

0	
1	
2	
3	
4	
5	
6	
7	
8	8
9	

Many - to - One Mapping

$$h(x) = x^0 / 10$$

key space



$$\underline{h(x) = x \% 10}$$

$$h(8) = 8 \% 10 = 8$$

$$h(3) = 3 \% 10 = 3$$

$$h(66) = 66 \% 10 = 6$$

$$h(4) = 4 \% 10 = 4$$

$$h(10) = 10 \% 10 = 0$$

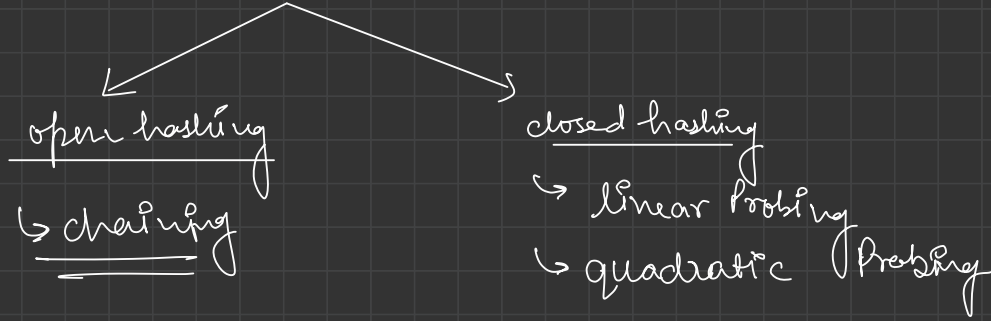
$$h(13) = 13 \% 10 = 3$$

hash table

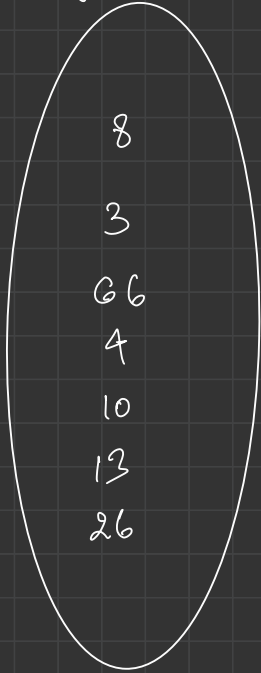
0	10
1	
2	
3	3
4	4
5	
6	66
7	
8	8
9	

collision

Methods to Remove Collision



Key Space



$$\underline{h(x) = x \% 10}$$

$$h(8) = 8 \% 10 = 8$$

$$h(3) = 3 \% 10 = 3$$

$$h(66) = 66 \% 10 = 6$$

$$h(4) = 4 \% 10 = 4$$

$$h(10) = 10 \% 10 = 0$$

$$h(13) = 13 \% 10 = 3$$

$$h(26) = 26 \% 10 = 6$$

Search(13)

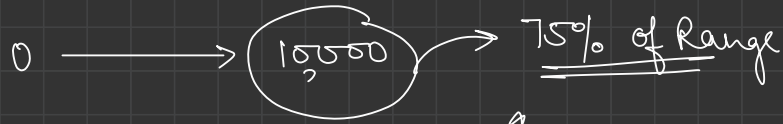
TC: O(1)

Searching

avg. time Complexity

hash table

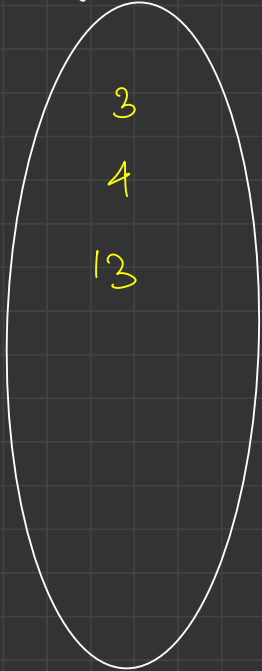
0	10	
1		
2		
3	3	→ 13
4	4	
5		
6	66	→ 26
7		
8	8	
9		



$h(n) = n \% \text{ size}$

linear probing

key space



$$h(x) = \{ h'(x) + f(i) \} \% \text{size}$$

$$h(x) = \{ h'(x) + f(i) \} \% 10$$

$$h'(x) = x \% 10 \quad \quad f(i) = i, 0, 1, 2, 3, \dots$$

$$\begin{aligned} h(3) &= \{ h'(3) + f(i) \} \% 10 \\ &= \{ 3 + 0 \} \% 10 = 3 \end{aligned}$$

$$h(4) = \{ 4 + 0 \} \% 10 = 4$$

$$h(13) = \{ 3 + 0 \} \% 10 = 3$$

$$\{ 3 + 1 \} \% 10 = 4$$

$$\{ 3 + 2 \} \% 10 = 5$$

disadv

clustering
issues

hash table

0	
1	
2	
3	3
4	4
5	13
6	
7	
8	
9	

Quadratic Probing

$$h(x) = \left\{ h'(x) + f(i) \right\} \% \text{ size}$$

$$h'(x) = x \% \text{ size}$$

$$f(i) = i^2, \underline{\underline{0, 1, 2, \dots}}$$

HashMap, HashSet



{ Hashing Algorithm }

TreeMap, TreeSet



{ Red-Black trees }

HashSet : collection of unique entity,

3, 13, 13, 4, 3, 10, 9, 9, 4

HashSet

3, 13, 4, 10, 9

Values are
in random order

{ TC: $O(1)$ searching
TC: $O(1)$ insertion

TreeSet

Values will be
inc order

3, 4, 9, 10, 13

{ TC: $O(\log N)$ searching
TC: $O(\log N)$ insertion

HashMap

→ key - value pairs

HashMap < Integer, String > map = new HashMap();

map.get(1) → "Accio"

TC: O(1) → Searching =

1.	"Accio"
2.	"Parsen"
3.	"Amuray"

Keys are stored in random order.

TreeMap

→ stores keys in asc. order

→ search TC: $O(\log N)$, insertion TC: $O(\log N)$

Break till 10:40pm°

```

// faith -> returns number of people managed by this emp
int rec (String emp, HashMap<String, ArrayList<String>> directReportee) {
    ① if (directReportee.containsKey(emp) == false) {
        return 0;
    }

    ② int cnt = 0;
    ③ for (String directs : directReportee.get(emp)) {
        cnt += rec(directs, directReportee) + 1;
    }
    ④ System.out.println(emp + " " + cnt);
    ⑤ return cnt;
}

```

