



Stacks

* Linear Data Structures

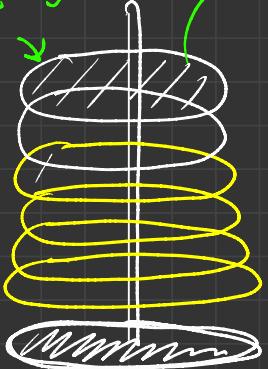
- ↳ Arrays
- ↳ ArrayLists
- ↳ Linked List
- ↳ Queues

* Non Linear Data Structures

- ↳ HashMap | HashSet
- ↳ tree - trees - heaps
- ↳ graph

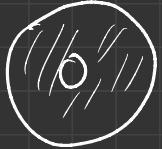
Stacks

last CD that added → topmost CD



LIFO

{ last in first out }





Stack (LIFO)

push(int v): Add a value v to the top of the
stack

Stack

pop(): Remove the topmost Ele

peek(): able to view the topmost data.

size(): size of the stack

```
class Stack
```

```
{ ArrayList<Integer> list;
```

```
int size;
```

```
Stack()
```

```
{
```

```
list = new ArrayList();
```

```
size = 0;
```

```
}
```

```
void push (int v) TC: O(1)
```

```
{
```

```
list.add (v);
```

```
size++;
```

```
}
```

```
int pop() TC: O(1)
```

```
{ if (size > 0)
```

```
{ int ele = list.remove(size - 1);
```

```
size--;
```

```
return ele;
```

```
} else
```

```
return -1;
```

```
int peek() TC: O(1)
```

```
{ if (size > 0)
```

```
{ return list.get(size - 1);
```

```
} else
```

```
return -1;
```

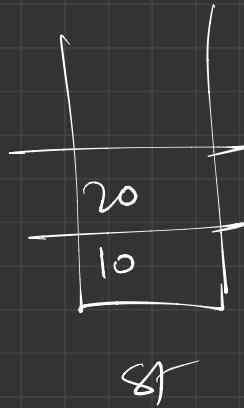
```
}
```

Practical Examples of Stack

- ↳ Recursion
- =====
- ↳ Memory management
- ↳ Cache
- ↳ Android Back

Stack < Integer > st = new Stack<>();
↓ ↓
class reference
variable

st.push(10);
st.push(20);
st.push(30);
print(st.peek()); → 30
st.pop(); → 30
print(st.peek()); → 20
size(); → 2



Q. Extra Brackets

String Sto = " $(a+b)$ ";  No

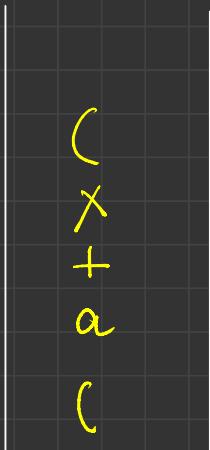
" $((a+b))$ ";  Yes

" $(a+b) + (c+d+(e*f))()$ ";  Yes

" $((a)+(b))$ ";  No

$$Stack = " \left(a + (b * d + f - (m) + n - o) \times (z) () \right) "$$

last open is first closed!



Stack

```

public boolean ExtraBrackets(String exp) {
    // Write your code here
    Stack<Character> st = new Stack<>();

    for (int i = 0; i < exp.length(); i++) {
        char ch = exp.charAt(i);

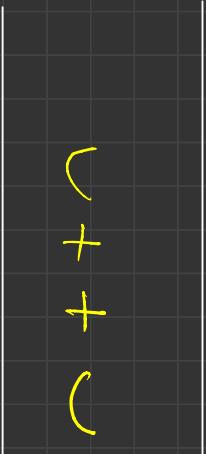
        if (ch != ')') {
            st.push(ch);
        } else {
            if (st.peek() == '(') {
                return true;
            } else {
                // remove the exp in between
                while (st.peek() != '(') {
                    st.pop();
                }
            }
        }

        // as this pair holds a exp in between, that means a valid bracket
        // so remove the corresponding opening bracket
        st.pop();
    }

    // if no extra bracket was found
    return false;
}

```

$$exp = ((a) + (b) + ((c+d)))$$



TC: $O(N)$
SC: $O(N)$

St.

$N \rightarrow push()$

$N \rightarrow pop()$

$O(N) + O(N) \approx \underline{O(N)}$

Next Greater Element On Right

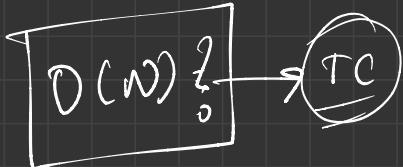
$$\text{ans}[] = \{ 0, 1, 2, 3, 4, 5, 6, 7, 8, 9 \}$$

$$\text{ngers}[] = \{ 6, 7, 2, 7, -1, 4, 8, 2, 5, -1 \}$$



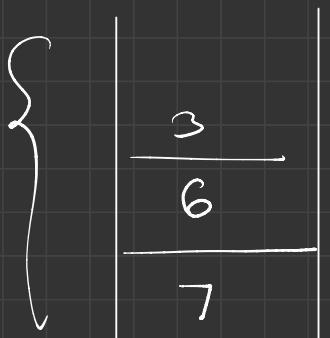
Brute force

$$\begin{aligned} \text{TC: } & O(N^2) \\ \text{SC: } & O(1) \end{aligned}$$



$arr[] = \{ \cancel{3}, \cancel{6}, \cancel{1}, \cancel{2}, \cancel{7}, \cancel{3}, \cancel{4}, \cancel{1}, \cancel{2}, \cancel{8} \}$

$nger[] = \{ 6, 7, 2, 7, -1, 4, 8, 2, 8, -1 \}$



Stack \rightarrow all the potential nge

```

public static long[] nextLargerElement(long[] arr, int n) {
    //Write code here and print output
    Stack<Long> st = new Stack<>();
    long[] nger = new long[n];

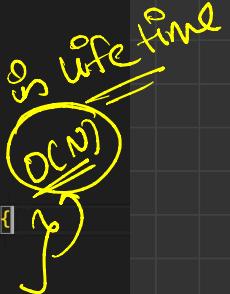
    // traverse from right to left
    // because I want info about potential nge
    for (int i = n - 1; i >= 0; i--) {
        // remove smaller potential nge from stack
        // as I'm covering for them
        while (st.size() > 0 && st.peek() <= arr[i]) {
            st.pop();
        }

        if (st.size() == 0) {
            nger[i] = -1;
        } else {
            nger[i] = st.peek();
        }

        st.push(arr[i]);
    }

    return nger;
}

```



$\downarrow \leftarrow \downarrow \leftarrow \times$

$arr[] = \{ 6, 8, 0, 1, 3 \}$

$\{ 8, -1, 1, 3, -1 \}$

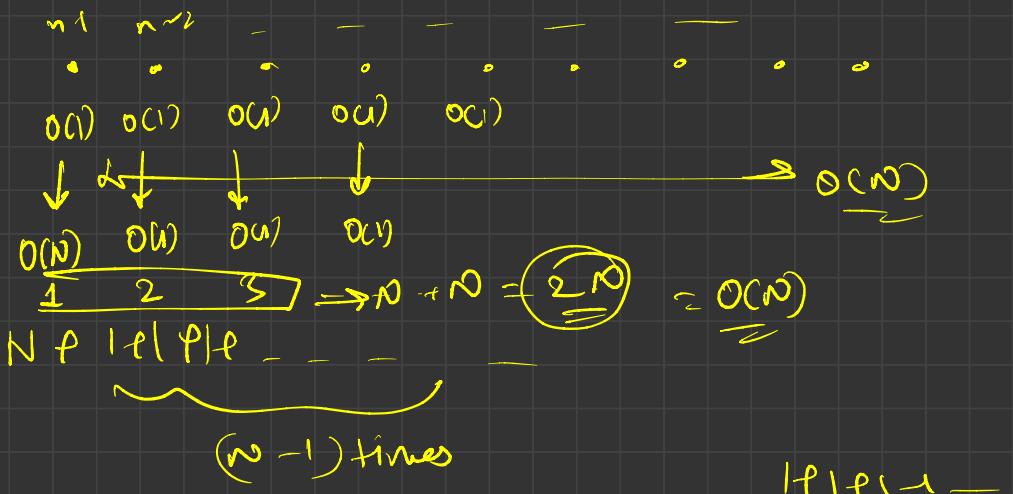
6
8

St

↑
↓
Atomized O(N)

TC: O(N)

SC: O(N)

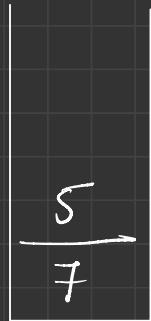


$$N + \underbrace{(N-1)}_{(N-1) \text{ times}} = (2N-1)$$

$O(N)$

~~ans[] = { 3, 6, 1, 2, 7, 3, 4, 1, 2, 8 }~~

~~ngers[] = { 6, 7, 2, 7, -1, 4, 8, 2, 5, -1 }~~



Stack → pool of people looking for ngers

```

public static long[] nextLargerElement(long[] arr, int n) {
    // Write code here and print output
    Stack<Integer> st = new Stack<>();
    long[] nger = new long[n];

    // traverse from left to right
    // because I want to create a pool of people looking for nge
    for (int i = 0; i < n; i++) {
        long ele = arr[i];

        // I can be a nge
        // go and see in pool, who all interested
        while (st.size() > 0 && arr[st.peek()] < ele) {
            nger[st.peek()] = ele;
            st.pop();
        }

        // I'm also looking for my nge
        st.push(i);
    }

    // people still looking for nge, will have it as -1
    while (st.size() > 0) {
        nger[st.peek()] = -1;
        st.pop();
    }

    return nger;
}

```

~~1~~ ~~2~~ ~~3~~ ~~4~~ ~~5~~ ~~6~~ ~~7~~ ~~8~~ ~~9~~ ↓
~~3~~, 6, 1, 2, 7, 3, 4, 1, 2, 8 }

~~0~~ ~~1~~ ~~2~~ ~~3~~ ~~4~~ ~~5~~ ~~6~~ ~~7~~ ~~8~~ ~~9~~
nger[] = { 6, 7, 2, 7, -1, 4, 5, 2, 5, -1 }

TC : O(N)
 SC : O(N)

—

st

Stock Span Problem

{ 0, 1, 2, 3, 4, 5, 6 }
100, 80, 60, 70, 60, 75, 85 }

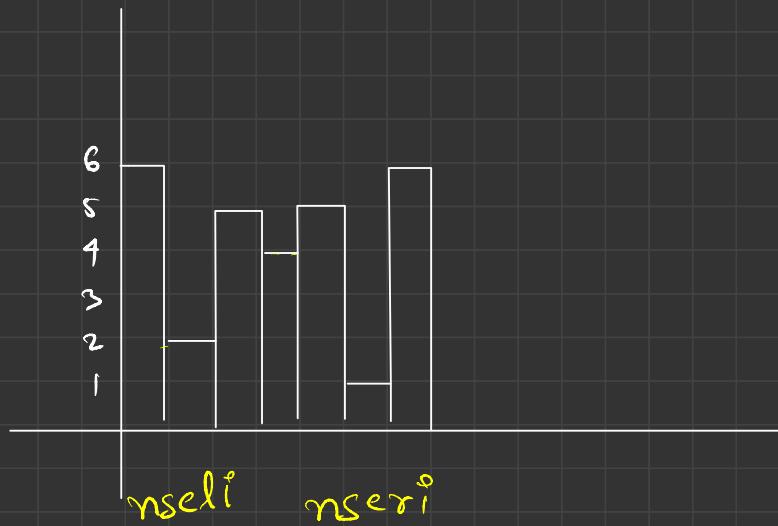
-1, 0, 1, 1, 3, 1, 0

Stopping Index

disk

$0 - (-1) \quad 1 - 0 \quad 2 - 1 \quad 3 - 1 \quad 4 - 3 \quad 5 - 1 \quad 6 - 0$
= 1 = 1 = 1 = 2 = 1 = 4 = 6 → Aws

Largest Area Histogram



maxArea = ~~10~~ ~~12~~ ✓

$$hi = \{ 6, 2, 5, 4, 8, 1, 6 \}$$

$$nsele = \{ -1, -1, 1, 1, 3, -1, 5 \}$$

$$nseri = \{ 1, 8, 3, 8, 5, 7, 7 \}$$

$$\text{width} = \{ 1, 5, 1, 3, 1, 7, 1 \}$$

Area

$$\underline{\overline{\max^n}}$$

- {
 - ① get NSELI
 - ② get NSERI
 - ③ get width right - left - 1

Area = width \times height

- ④ Store Max Area