



## Binary Search (Algorithm)

int[] arr = { 0 1 2 3 4 5 6 7 }  
1, 3, 7, 10, 11, 14, 20, 40 } target = 14

→  
sorted

### Brute force

```
for (int i = 0; i < n; i++)  
{  
    if (arr[i] == target)  
        return i;  
}
```

Linear Search

TC:  $O(n)$   
SC:  $O(1)$

what is the lowest floor from which, throwing brick will break

1500

1000 Bricks

↳ Brute force

① Optimized Approach

$$\log_2(1000) = \log_2 10^3 = 3 \times \log_2 10$$

$$\approx \sqrt[3]{10} \approx 2.1$$

Bricks

→ 500<sup>th</sup> floor (Breaks)

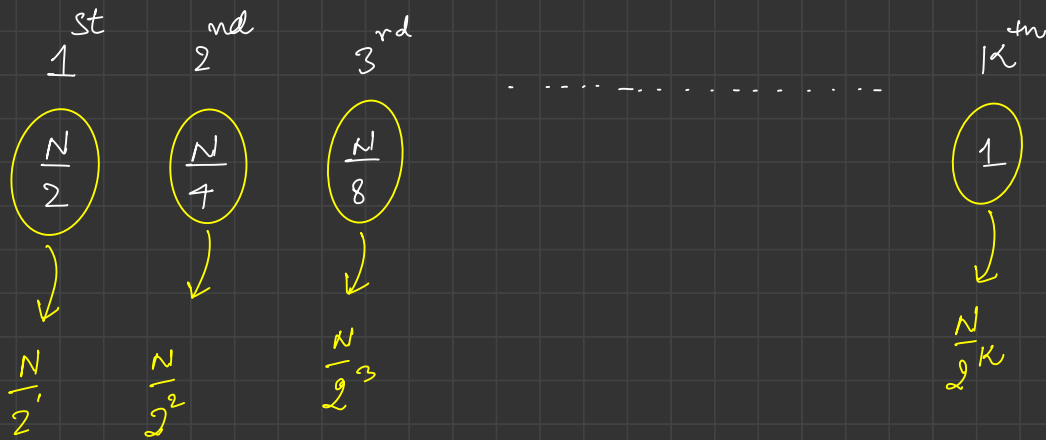
→ By using first Brick, I eliminated 500 floors

→ 375<sup>th</sup> floor (Not Break) → By using second Brick, I eliminated 250 more floors

→ By using 3<sup>rd</sup> Brick, I eliminate 125 more floors.

→ 250<sup>th</sup> floor (Not Break)

0



$$\frac{N}{2^K} = 1$$

$$N = 2^K$$

taking  $\log_2$  Both side

$$\log_2 N = \log_2 2^K$$

$$\log_2 N = K \log_2 2 \rightarrow 1$$

$$K = \log_2 N$$

→ Bricks are required!

int[] arr = { 0 1 2 3 4 5 6 7 } sorted Array  
1, 3, 7, 10, 11, 14, 20, 40 } target = 11

↑ mid  
si  
↑  
ei

while (si <= ei)

{ if (target == arr[mid])  
return mid  
else if (target > arr[mid])  
si = mid + 1;  
else  
ei = mid - 1;

TC:  $O(\log N)$   
SC:  $O(1)$

Binary Search

# ~~\*\*\*\*~~ Binary Search

↳ always implemented over a sorted region

→ Sorted region

→ TC:  $O(\log_2 N)$  (Expected)

99%  
→ Binary Search

Search insert position (ceil value)

arr[] = { 1, 3, 7, 10, 11, 20, 40 }

key = 2

Brute force

{ if (arr[i] == key) return i;  
else if (arr[i] > key) return i;

TC:  $O(N)$

SC:  $O(1)$

$arr[] = \{ \overset{0}{1}, \overset{1}{3}, \overset{2}{7}, \overset{3}{10}, \overset{4}{11}, \overset{5}{20}, \overset{6}{40} \}$   
 $key = 2$   
 $\uparrow$   $\uparrow$   $\uparrow$   
 $si$   $mid$   $ei$

Case 1

$arr[mid] == key$

Case 2

$arr[mid] > key$

$ans = mid$

$ei = mid - 1;$

Case 3

$arr[mid] < key$

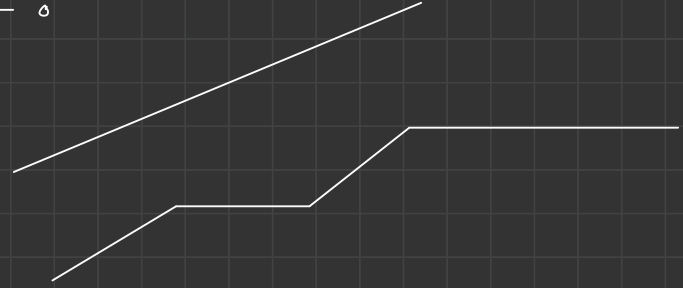
$si = mid + 1;$

$ans = arr.length$   
 $\downarrow$   
cell value  $\rightarrow$  (default value)



find first and last pos. of a element.

inc. array  
non-decreasing array



int[] arr = { 0 1 2 3 4 5 6 7 8 9 10 11 12  
1, 2, 2, 2, 2, 2, 3, 4, 4, 10, 20, 20, 30 }

ele = 2

fo = 1    lo = 5

Brute force

TC:  $O(N)$

SC:  $O(1)$

int[] arr = { 0 1 2 3 4 5 6 7 8 9 10 11 12  
 1, 2, 2, 2, 2, 2, 3, 4, 4, 10, 20, 30, 30 }

ele = 2

↑    ↑  
 ei   si

first Occurrence    mid

~~f0 = 1~~

case 1 :    arr[mid] == ele    }  
              f0 = mid;  
              ei = mid - 1;

case 2 :    arr[mid] > ele  
              ei = mid - 1;

case 3 :    arr[mid] < ele  
              si = mid + 1;

## Search in a 2D-Matrix (2D Sorted Matrix)

Row [ ] [ ] col

	0	1	2	3
0	1	2	3	4
1	5	6	7	8
2	9	10	11	12
3	13	14	15	16
4	17	18	19	20

target = 14

(5 x 4)

---