



# POORNIMA FOUNDATION

## DETAILED LECTURE NOTES

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Services provided by DLL

- framing
- error control
- flow control (one is fast  $\leftarrow$  other is slow)
- Physical addressing
- Access control

- This is speed matching mechanism.

- It tells the sender how much data it can transmit before it ~~can~~ must wait for an acknowledgement from the receiver.

flow control  
protocols

Noiseless  
Channels

• stop & wait

Noisy  
Channel

- stop & wait ARQ

- Go back N ARQ  
- selective Repeat ARQ

Sliding window  
protocol.

### 1) Stop & wait

- is DLL protocol for transmission of frames over noiseless channels.
- It provides unidirectional data transmission with flow control facilities but without error control facilities.
- The idea is straightforward
- After transmitting one frame, the sender waits for an acknowledgement before transmitting the next frame.

### Sender side

Rule 1 - send one data packet at a time

Rule 2 send the next packet only after receiving ACK for the previous

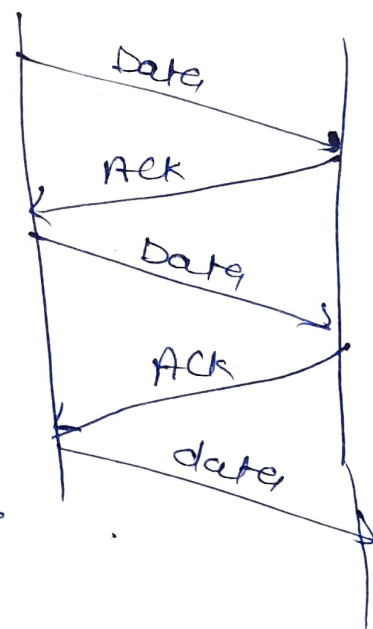
### Receiver side

Rule 1 Receive & consume data packet

Rule 2 After consuming packet, Ack need to be sent (flow control)

### Problems

1) due to lost data  
sender wait for infinite amount of time for ack  
Receiver for data

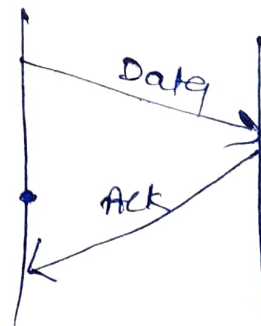


2) Due to lost ACK  
sender waits for ack for infinite amount of time

3) Due to delayed ACK (data)

After timeout, a delayed ack is received ~~might~~

delayed ack might be wrongly considered as ack of some other data packet.





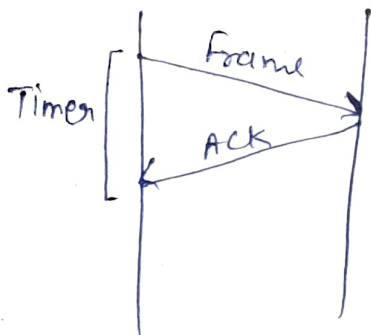
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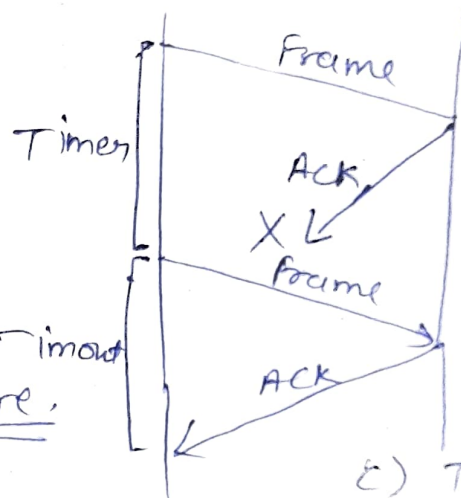
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stop and wait ARQ (Automatic Repeat Request)  
protocol.

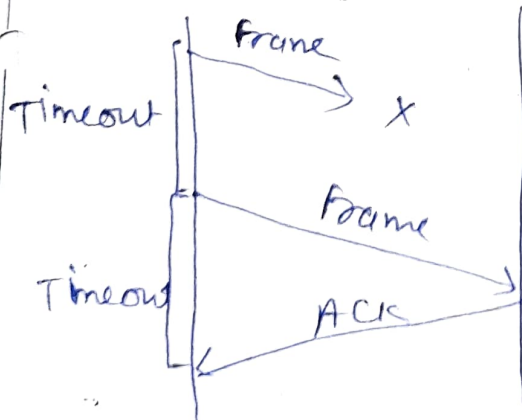
= stop & wait + Timeout timer +  
sequence No.



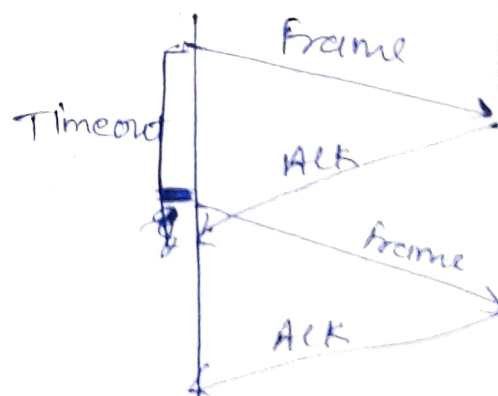
a) ACK is received  
before timeout expires.



c) The ACK is lost



b) The original frame is  
lost



d) The time out fires too soon

ACK is late. ACK received  
after timeout expired. It retransmit again.