

## School of CET

# **System Software and Compiler lab**

# **Assignment No.1**

### TY BTech CSE

**Assignment Title:** Design of Pass 1 of Two Pass Assembler.

**Aim:** Design suitable data structure & implement pass 1 of Two Pass Assembler pseudo machine.

**Objective:** Design suitable data structure & implement pass 1 of Two Pass Assembler pseudo machine. Subset should consist of a few instructions from each category & few assembler directive.

### **Theory:**

### Assembler

The assembler is a program for converting instructions written in low-level assembly code into relocatable machine code and generating along information for the loader.

It generates instructions by evaluating the mnemonics (symbols) in operation field and find the value of symbol and literals to produce machine code. Now, if assembler do all this work in one scan then it is called single pass assembler, otherwise if it does in multiple scans then called multiple pass assembler. Here assembler divide these tasks in two passes:

#### • Pass-1:

- 1. Define symbols and literals and remember them in symbol table and literal table respectively.
- 2. Keep track of location counter
- 3. Process pseudo-operations

#### • Pass-2:

1. Generate object code by converting symbolic op-code into respective numeric op-code

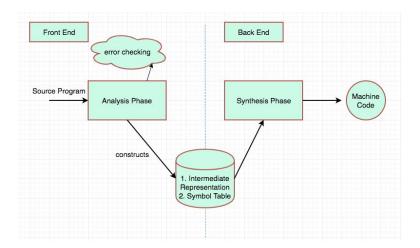
2. Generate data for literals and look for values of symbols

# Design specification of an Assember

<u>Analysis phase</u>: reads the source program and splits it into multiple tokens and constructs the intermediate representation of the source program. And also checks and indicates the syntax and semantic errors of a source program. It collects information about the source program and prepares the symbol table. Symbol table will be used all over the compilation process. This is also called as the front end of a compiler.

<u>Synthesis phase</u>: It will get the analysis phase input(intermediate representation and symbol table) and produces the targeted **machine level code**.

This is also called as the **back end** of a compiler.



### Algorithm

Algorithm for Pass I

- 1. locentr: =0; (default value)
- 2. While next statement is not an END statement

(a) If label is present then this-label: = symbol in label field; Enter (this-label, loccntr) in SYMTAB.

- (b) If a START or ORIGIN statement then
  - *loccntr:* = value specified in operand field;
- (c) If an EQU statement then
  - (i) this-addr:= value of <address spec>;
  - (ii) Correct the symtab entry for this-label to (this-label, this-addr).
- (d) If a declaration statement then
  - (i) *code:* = code of the declaration statement;
  - (ii) *size:* = size of memory area required by DC/DS.
  - (iii) locentr := locentr + size;
  - (iv) Generate IC '(DL, code).
  - (f) If an imperative statement then
    - (i) *code:* = machine opcode from OPTAB;
    - (ii) *locentr:* = *locentr* + instruction length from OPTAB;
    - (iii) If operand is a symbol then

this-entry: = SYMTAB entry number of operands;
Generate IC '(IS, code) (S, this-entry)';

- 3. (Processing of END statement)
- (b) Generate IC
- (c) Go to Pass II.

## **Listing & Error Handling:**

Input: Assembly Language Program. /Intermediate code generated by PASS I

### **Output:**

1. Mnemonic Table

Mnemonic	Op-code	Length

2. Symbol Table

Symbols/ Labels	Address

3. Intermediate Form (After Pass I)/Final Output

Address (LC value)	Op-code	Operand 1	Operand 2
		(Temp)	(Temp)

**Conclusion:** The function of Pass I in assembler are studied along with errors coming in each pass.

**Platform:** Linux (JAVA)

```
import java.io.*;
import java.util.*;
class Condition {
class Register {
class Operator {
```

```
class Symbol {
class Tuple<X, Y> {
public class assembler rg {
        // static ArrayList<Symbol> symbolTable = new
ArrayList<Symbol>(25);
          static ArrayList<String> intermed code = new
ArrayList<String>(25);
    static ArrayList<Tuple<String, String>> literalTable = new
ArrayList<Tuple<String, String>>(25);
    static ArrayList<Tuple<String, String>> symbolTable = new
ArrayList<Tuple<String, String>>(25);
                                                         new
ArrayList<Integer>(25);
```

```
static ArrayList<Operator>
ArrayList<Operator>(25);
ArrayList<Register>(25);
ArrayList<Condition>(25);
```

```
System.out.println("----");
   } catch (NumberFormatException e) {
public static void main(String args[]) throws IOException {
            BufferedReader br = new BufferedReader(new
```

```
System.out.println("################");
Arrays.toString(splitted));
splitted[2]);
                                  Tuple<String, String> lt = new
Tuple<String, String>(splitted[2], "Address not given");
                              literalTable.add(lt);
```

```
Tuple<String, String> st = new
Tuple<String, String>(splitted[2], "Address not given");
symbolTable.size(); k++) {
symbolTable.get(k);
                                    intermed code.add("(S," + k +
                                             String[] newArray =
Arrays.copyOfRange(splitted, 1, splitted.length);
                                  Tuple<String, String> s = new
Tuple<String, String>(splitted[0], line number + "");
```

```
literal table
Tuple<String, String>(splitted[3], "Address not given");
                                  Tuple<String, String> st = new
Tuple<String, String>(splitted[3], "Address not given");
Arrays.copyOfRange(splitted, 1, splitted.length);
                  Operator x = null;
                   Iterator<Operator> i = aList.iterator();
```

```
Register r = null;// name,mc code
h()-1));
1))) {
splitted[2].contains("=")) {
k++) {
literalTable.get(k);
```

```
!splitted[2].contains("+")) {
k++) {
splitted[2].contains("+"))// ORIGIN L1 +3
                         for (int gg = 0; gg < symbolTable.size();</pre>
gg++) {
Integer.parseInt(s.y);
Integer.parseInt(splitted[2].substring(1)) - 1;
```

```
i++) {
                           Operator x = null;
                                         Iterator<Operator> hh =
aList.iterator();
                                   // System.out.println(x.name+"
intermed code.add("(C,"+literalTable.get(i).x.substring(2,3)+")")
begins for literal table;
```

```
Operator x = null;
                                         Iterator<Operator> hh =
aList.iterator();
                                   // System.out.println(x.name+"
                                    intermed code.add("(" + x.cls
```

```
START 200
MOVER AREG, ='5'
MOVEM AREG, X
L1 MOVER BREG, ='2'
ORIGIN L1 +3
LTORG
X DS 1
END
###################
##################
201: [MOVER, AREG,, ='5']
splitted[2]:='5'
----LITERAL TABLE-----
```

```
##################
202: [MOVEM, AREG,, X]
splitted[2]:X
###################
203: [L1, MOVER, BREG,, ='2']
```

```
(C,200)
###################
204: [ORIGIN, L1, +3]
splitted[2]:+3
```

```
(C,205)
##################
206: [LTORG]
```

```
(C,5)
207: [X, DS, 1]
splitted[2]:1
```

```
(DL,2)
##################
208: [END]
ggggggggg
(C,200)
```