An Introduction to CUDA using IPT

1. Introduction to IPT

The Interactive Parallelization Tool (IPT) is a tool for semi-automatically parallelizing existing serial programs written in C/C++. It supports parallelization using the three most popular parallel programming paradigms: OpenMP, MPI, and CUDA. The users interactively specify which regions or loops to parallelize in a serial program. On the basis of the specifications provided by the users, IPT can parallelize several types of C/C++ applications. It therefore frees the users from the burden of learning the low-level syntax of the different parallel programming paradigms, and in manually reengineering their serial code for creating parallel versions. In order to do so, IPT walks the users through a series of questions related to the parallelization of the input C/C++ programs and generates parallelized versions as output.

2. Introduction to GPU Programming

Graphic Processing Units (GPUs) are computing devices primarily designed for processing graphic and video inputs, and relaying the output to a display. GPUs are therefore designed to handle high throughput, where throughput is the number of tasks completed per unit of time. In comparison to a Central Processing Unit (CPU) that could have fewer but faster cores, a GPU could be having hundreds of cores that can handle thousands of threads. Therefore, even though a GPU's clock is slower than that of a CPU, the GPU is able to execute a greater number of commands per clock-cycle.

GPUs can be used for general-purpose programming and are used for running parallelized applications in several areas, such as physical simulations or machine learning. Nvidia and AMD are the two popular GPU manufacturers.

Compute Unified Device Architecture (CUDA) is a paradigm for parallelizing C and C++ programs by utilizing a system's GPU for general computing. In the early 2000's, the CUDA Toolkit was developed by Nvidia to provide programmers with a library and compiler for parallel programming on Nvidia GPUs.

In the CUDA programming model there is a **host,** and one or more **devices**. The CPU is the host, and the GPU is the device. The GPU acts as an accelerator to the CPU, which is the main processor or the host. The host and the device have separate memory. Therefore, the programmers may need to allocate memory on both the host and device, and explicitly write code to copy or move the data between them as needed. Similarly, one needs to demarcate the functions that run on the device from those running on the host. A function that runs on the device, and is called from the host, is called a **kernel**. The GPU responds to the kernels launched by the host (or the CPU). Only the host can launch kernels. Kernels are labeled with the identifier `__global__` to indicate to the CUDA compiler that this function will be called by the CPU to run on the GPU. In addition to global functions, there are also device functions that are demarcated with the identifier `__device__`. These functions can only be called from kernels or other device functions. The host cannot call device functions. Similarly, the host also has functions that can only be run by the host. For example the `main` function is always run from the host. Device and host functions are split at the compile-time. Host functions are processed by the standard compiler (such as `gcc`), and the device functions are processed by the `nvcc` compiler from the CUDA toolkit.

When a kernel is launched, the instructions are executed in a **thread**. The GPU can utilize thousands of threads to complete a task. Each thread will execute the same set of instructions from the kernel. Threads are run in **warps** which are groups of 32 threads that execute simultaneously. Each thread has its own local memory that is private to the thread, as well as an identification number. There is a three-dimension arragement of threads on a GPU: x, y, and z are the dimensions. This dimensionality can be utilized when parallelizing problems with multiple dimensions such as performing calculations on a 2-D matrix.

A group of threads that share memory is called a **block**. Threads running in the same block can access the same address space in memory. Depending on the GPU, the maximum number of threads in a block is either 512 or 1024. Blocks are assigned to run on **stream multiprocessors**, or **SMs** for short. SMs are hardware constructs that have registers, cache, and memory. SMs also have 2 warp schedulers. When a SM executes the threads in a block, it runs them in groups called warps. As mentioned before, warps are groups of up 32 threads that run simultaneously. If a block contains 64 threads, the SM will run them in two warps. If a block has 100 threads, it will run 4 warps. The first 3 warps will all contain 32 threads, and the last warp would have 4 threads and 28 dummy threads.

Blocks are themselves organized in a three-dimensional structure called a grid. Blocks also extend in the x, y, and z dimensions (this kernel dimensionality is referred to by variables with `dim3` type and will be covered in detail at a later point) and have their own identification number. When launching a kernel, the programmer specifies the number of blocks to run in each dimension, and also specified the number of threads in each block per dimension. There is no way to specify what SM a block is assigned to, or the order that the SM will run the blocks in. As a result, blocks must be independent from one another and must be able to be run in any order.

Indexing is another aspect of CUDA grid dimensionality. Often times kernels will be written based on the thread index. As both threads and blocks have identification numbers in the x, y, and z dimensions, CUDA has the following built-in variables related to identification and dimensions: `uint3 threadIdx`, `uint3 blockIdx`, `dim3 blockDim`, `dim3 gridDim`. The `uint3` data type is a coordinate variable in CUDA based on the typical `int` data type. These variables can be used to

find either aggregate properties about programs (for example, how many kernels are running) or as indices to control thread actions. For example, in order to specify that a kernel only runs on threads in the z dimension, one can add an if statement that checks the thread is a z-dimension thread.

3. CUDA Program Structure

Many CUDA programs follow the same general structure since there is a defined relationship between the host and device. The CPU is responsible for most of the overhead for running a kernel such as moving data from the host to the device, launching kernels, and copying the results from the device back to the host. To launch a kernel, the CPU will allocate storage on the GPU before copying data from the CPU to the GPU. In order to allocate memory and copy the data, often the functions `cudaMalloc` (for allocating memory on the device), `cudaMemcpy` (for copying data from the host to the device, or vice versa), and `cudaFree` (for releasing used memory on the device) are used.

Executing a kernel is a concurrent operation in CUDA unless specified, therefore the CPU will continue executing instructions while the kernel is running. However, a new kernel operation cannot start until the previous one has concluded.

Figure 1 shows the general structure of a CUDA program, in which only one kernel is called. A kernel is launched using the syntax shown in Figure 2.

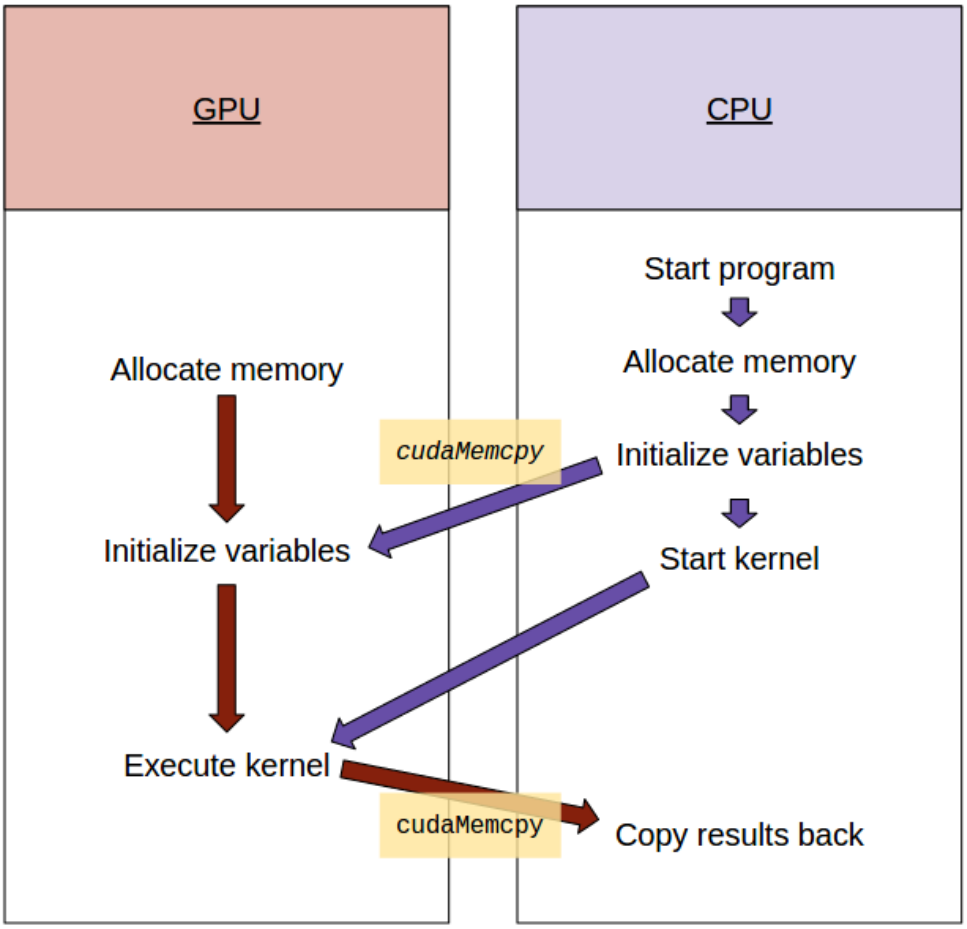

Figure 1: General Strcuture of a CUDA program

```
dim3 gridSize(M,1,1);

dim3 blockSize(N,1,1);

kernelName<<< gridSize, blockSize >>>(kernel Parameters);
```

Figure 2: A C code snippet showing the syntax for kernel launch

The three pairs of chevrons shown in Figure 2 are used to set the execution parameters. In Figure 2, the gridSize represents the number of blocks and the blockSize represents the number of threads in each block.

The parameters inside the chevrons are of **`dim3`** type and are useful for setting the dimensionality for kernel calls. In Figure 2, `dim3` was initialized with all three parameters but the user could choose to provide only 1 or 2 parameters depending on the problem being solved. Any parameter that is not initialized defaults to 1. To determine the total number of threads that a kernel is using, one must consider the number of threads in each dimension as well as the number of blocks in each dimension. The total number of threads is equal to the product of the number of blocks and the number of threads per block. To determine the number of threads in a block, multiply the number of threads in each dimension together for the thread `dim3` object. In the example code snippet in Figure 2, the total number of threads in each block is N*1*1 = N. The same process is repeated to find the total number of blocks in the grid M*1*1=M. The total number of threads used here is: M*N.

In summary, the program structure shown in Figure 1, should have the components listed in Figure 3. Note that, as mentioned before, because the GPU and CPU have separate memory, the GPU results obtained in the kernel must be returned to the CPU for further processing. This is done via `cudaMemcpy()` as well. Finally, the allocated GPU memory can be freed via `cudaFree()`.

1. CPU allocates storage on GPU (`cudaMalloc()`)
2. CPU copies data from CPU to GPU (`cudaMemcpy()`)
3. CPU launches kernel on GPU to perform some computation
4. CPU copies results back from GPU (`cudaMemcpy()`)

Figure 3: General components of the CUDA program represented in Figure 1

4. Generating the First CUDA Program with IPT

The C program shown in figure 4 can be used to compute the sum of all integers in the interval [0, 10000]. We will use this program as a test case for demonstarting the process of using IPT for generating CUDA code. The steps for parallelizing the code with IPT are shown in Figure 5. The Generated code is shown in Figure 6. The user may need to modify the number of threads and blocks that are specified in the generated code.

```
#include <stdio.h>
int main(int argc, char *argv[]){
```

```
   int i, sum,upToVal;
   upToVal = 10001;
   sum = 0;
   for(i=0; i< upToVal; i++){
      sum = sum +i;
   }
   printf("\nSum of first %d numbers is: %d\n\n",upToVal, sum);
   return 0;
}
```

Figure 4: C program to compute the sum of integers in the interval [0, 10000]

```
Please select a parallel programming model from the following available
options:
1. MPI
2. OpenMP
3. CUDA
3

Please enter the function in which you wish to insert the kernel call(or
parallelize the for-loop).
Please choose the function that you want to parallelize from the list below
1 : main
1

Would you like
1. For-loop
2. TBD
1

for (i = 0; i < upToVal; i++) {
  sum = sum + i;
}
Is this the for loop you are looking for?(y/n)
y
Do you want to perform reduction on any variables ? (Y/N)
y

Please enter the variables to reduce ([ format:  1,2,3 etc. ] with 1 is for
the first variable, 2 is for the second variable and so on).
 Possible variables to reduce are:
1. sum type is int
2. i type is int
1

Please enter the reduction operation for variable [sum]. Possible reduction
operations are:
1. Sum
2. Product
1
Do you want to perform reduction on any arrays ? (Y/N)
n
```

Figure 5: IPT prompts and user responses (bold)

```
#include <stdio.h>
void __global__ kernel0(int sum[],int upToVal, int device_M , int device_N){
int i =  blockIdx.x * blockDim.x + threadIdx.x;
int  print_statement_deleted_here=0;
sum[i] = 0;
{
sum[i] = sum[i] + i;
}
__syncthreads();
__syncthreads();
}

int main(int argc,char *argv[])
{
  int total_sum = 0;
  int *host_sum;
  int *device_sum;
  int i;
  int sum;
  int upToVal;
  upToVal = 10001;
  sum = 0;
//Inserting code for memory allocation grid size and block size calculation
host_sum= (int*)malloc((1)*((int) ((upToVal - 0 ) / (1)))*sizeof(int));
cudaMalloc((void **) &device_sum,( 1)*((int) ((upToVal - 0 ) /
(1)))*sizeof(int));
//Please note this is the section wherein the number of blocks and threads
are calculated.  To change the number of threads alter the dimBlock whereas
to change the number of blocks alter the dimGrid
int D_rows = 1 + ( ((int) ((upToVal - 0 ) / (1)) > 1024 ) ? (int) ((upToVal
- 0 ) / (1))/1024 : (int) ((upToVal - 0 ) / (1)) );
int D_cols = ((int) ((upToVal - 0 ) / (1)) > 1024 ) ? 1024 : 1;
dim3 dimGrid(D_rows,1);
dim3 dimBlock(D_cols,1);
kernel0<<<dimGrid,dimBlock>>>(device_sum,upToVal,1,(int) ((upToVal - 0 ) /
(1)));
/*
  int IPT_function_replace;
*/
//Copying from Device to Host
cudaMemcpy(host_sum,device_sum,((int) ((upToVal - 0 ) / (1)))*sizeof(int),
cudaMemcpyDeviceToHost);
 //code for variable reduction
 for(long row = 0; row < 1; ++row){for(long col = 0; col < (int) ((upToVal -
0 ) / (1)); ++col)      {  total_sum+= host_sum[row*(int) ((upToVal - 0 ) /
(1))+ col]; }  }
sum+= total_sum;
//  Ending Parallelization
  printf("\nSum of first %d numbers is: %d\n\n",upToVal,sum);
  return 0;
}
```

Figure 6: CUDA code generated by IPT – sum of numbers

5. Parallelizing the Code for Finding Sum of Arrays

The program shown in Figure 7 computes the sum of two arrays (A, B) and saves the results in array C. The steps to parallelize this code with IPT are shown in Figure 8, and the generated code is shown in Figure 9.

```
#include <stdio.h>
int main() {
  int lm = 10000;
  int A[lm];
  int B[lm];
  int C[lm];
  for (int j=0; j < lm; j++) {
      A[j] = j;
      B[j] = j;
  }
  // Addition
  for (int q=0; q < lm; q++) {
      C[q] = A[q] + B[q];
  }
  // Checker
  for (int s=0; s < lm; s++) {
      if (C[s] != 2*s) {
      printf("failed at row %d\n", s);
      }
  }
}
```

Figure 7: Serial C program to compute the sum of arrays A and B

```
Please select a parallel programming model from the following available
options:
1. MPI
2. OpenMP
3. CUDA
3

Please enter the function in which you wish to insert the kernel call(or
parallelize the for-loop).
Please choose the function that you want to parallelize from the list below
1 : main
1

Would you like
1. For-loop
2. TBD
1

for (int j = 0; j < lm; j++) {
  A[j] = j;
  B[j] = j;
}
Is this the for loop you are looking for?(y/n)
```

```
n

OK - will find the next loop if available.

// Addition
for (int q = 0; q < lm; q++) {
  C[q] = A[q] + B[q];
}
Is this the for loop you are looking for?(y/n)
y
cannot automatically find the number of iteration for the loop, please
specify the number of iteration:
10001
Do you want to perform reduction on any variables ? (Y/N)
n
Do you want to perform reduction on any arrays ? (Y/N)
n
Is the following array [ C ]
1. Input , 2. Output 3. Input/Output 4. Neither Input nor Output
2
Is the following array [ A ]
1. Input , 2. Output 3. Input/Output 4. Neither Input nor Output
1
Is the following array [ B ]
1. Input , 2. Output 3. Input/Output 4. Neither Input nor Output
1
```

Figure 8: IPT prompts and user responses (bold)

```
#include <stdio.h>
__global__  void kernel0(int * C,int * A,int * B,int lm, int device_M , int
device_N){
int q =  blockIdx.x * blockDim.x + threadIdx.x;
int  print_statement_deleted_here=0;
{
C[q] = A[q] + B[q];
}
__syncthreads();
}

int main()
{
  int *device_B;
  int *device_A;
  int *device_C;
  int lm = 10000;
  int A[lm];
  int B[lm];
  int C[lm];
  for (int j = 0; j < lm; j++) {
      A[j] = j;
      B[j] = j;
  }
// Addition
//Please note this is the section wherein the number of blocks and threads
```

```
are calculated.  To change the number of threads alter the dimBlock whereas
to change the number of blocks alter the dimGrid
int D_rows = 1 + ( (10001 > 1024 ) ? 10001/1024 : 10001 );
int D_cols = (10001 > 1024 ) ? 1024 : 1;
dim3 dimGrid(D_rows,1);
dim3 dimBlock(D_cols,1);
cudaMalloc((void **) &device_C,(lm)*sizeof(int));
cudaMalloc((void **) &device_A,(lm)*sizeof(int));
cudaMemcpy(device_A,A,(lm)*sizeof(int),cudaMemcpyHostToDevice);
cudaMalloc((void **) &device_B,(lm)*sizeof(int));
cudaMemcpy(device_B,B,(lm)*sizeof(int),cudaMemcpyHostToDevice);
kernel0<<<dimGrid,dimBlock>>>(device_C,device_A,device_B,lm,1,10001);
/*
  int IPT_function_replace;
*/
// Checker
cudaMemcpy(C,device_C,(lm)*sizeof(int), cudaMemcpyDeviceToHost);
cudaFree(device_C);
cudaFree(device_A);
cudaFree(device_B);
  for (int s = 0; s < lm; s++) {
      if (C[s] != 2 * s) {
      printf("failed at row %d\n",s);
      }
  }
}
```

Figure 9: CUDA code generated by IPT – sum of arrays

6. Parallelizing the Circuit Satisfiability Code

The program shown in Figure 10 is the serial version of the circuit satisfiability problem (https://people.sc.fsu.edu/~jburkardt/c_src/satisfy/satisfy.html).

This code simulated a logical ciruit can computes the number of combinations of 0s and 1s that are provided to the circuit for which the circuit produce an output of 1. The steps to parallelize this code with IPT are shown in Figure 11, and the generated code is shown in Figure 12. Note that only the number of solutions found is printed by the CUDA version.

```
# include <stdlib.h>
# include <stdio.h>
# include <time.h>

int main ( int argc, char *argv[] );
int circuit_value ( int n, int bvec[] );
void i4_to_bvec ( int i4, int n, int bvec[] );
void timestamp ( void );

/******************************************************************************
***/
```

```
int main ( int argc, char *argv[] )

/******************************************************************************
***/
/*
  Purpose:

      MAIN is the main program for SATISFY.

  Licensing:

      This code is distributed under the GNU LGPL license.

  Modified:

      20 March 2009

  Author:

      John Burkardt

  Reference:

      Michael Quinn,
      Parallel Programming in C with MPI and OpenMP,
      McGraw-Hill, 2004,
      ISBN13: 978-0071232654,
      LC: QA76.73.C15.Q55.
*/
{
# define N 23

  int bvec[N];
  int i;
  int ihi;
  int j;
  int n = N;
  int solution_num;
  int value;

  printf ( "\n" );
  timestamp ( );
  printf ( "\n" );
  printf ( "SATISFY\n" );
  printf ( "  C version\n" );
  printf ( "  We have a logical function of N logical arguments.\n" );
  printf ( "  We do an exhaustive search of all 2^N possibilities,\n" );
  printf ( "  seeking those inputs that make the function TRUE.\n" );
/*
  Compute the number of binary vectors to check.
*/
  ihi = 1;
  for ( i = 1; i <= n; i++ )
  {
      ihi = ihi * 2;
  }
```

```
  for ( i = 1; i <= n/2; i++ )
  {
      printf("\ntest\n");
  }
  printf ( "\n" );
  printf ( "  The number of logical variables is N = %d\n", n );
  printf ( "  The number of input vectors to check is %d\n", ihi );
  printf ( "\n" );
  printf ( "   #        Index ---------Input Values-----------------------
\n" );
  printf ( "\n" );
/*
  Check every possible input vector.
*/
  solution_num = 0;

  for ( i = 0; i < ihi; i++ )
  {
      i4_to_bvec ( i, n, bvec );

      value = circuit_value ( n, bvec );

      if ( value == 1 )
      {
      solution_num = solution_num + 1;

      printf ( "  %2d  %10d:  ", solution_num, i );
      for ( j = 0; j < n; j++ )
      {
      printf ( " %d", bvec[j] );
      }
      printf ( "\n" );
      }
  }
  //added this for testing. If the already parallelized loop is not the last
one, then there is no error, else segementation fault
  //for ( i = 1; i <= n/2; i++ )
  //{
  //  printf("\ntestagain\n");
  //}

//  Report.

  printf ( "\n" );
  printf ( "  Number of solutions found was %d\n", solution_num );
/*
  Shut down.
*/
  printf ( "\n" );
  printf ( "SATISFY\n" );
  printf ( "  Normal end of execution.\n" );
  printf ( "\n" );
  timestamp ( );

  return 0;
```

```
# undef N
}
/******************************************************************************
***/

int circuit_value ( int n, int bvec[] )

/******************************************************************************
***/
/*
  Purpose:

      CIRCUIT_VALUE returns the value of a circuit for a given input set.

  Licensing:

      This code is distributed under the GNU LGPL license.

  Modified:

      20 March 2009

  Author:

      John Burkardt

  Reference:

      Michael Quinn,
      Parallel Programming in C with MPI and OpenMP,
      McGraw-Hill, 2004,
      ISBN13: 978-0071232654,
      LC: QA76.73.C15.Q55.

  Parameters:

      Input, int N, the length of the input vector.

      Input, int BVEC[N], the binary inputs.

      Output, int CIRCUIT_VALUE, the output of the circuit.
*/
{
  int value;

  value =
      (  bvec[0]  ||  bvec[1]  )
      && ( !bvec[1]  || !bvec[3]  )
      && (  bvec[2]  ||  bvec[3]  )
      && ( !bvec[3]  || !bvec[4]  )
      && (  bvec[4]  || !bvec[5]  )
      && (  bvec[5]  || !bvec[6]  )
      && (  bvec[5]  ||  bvec[6]  )
      && (  bvec[6]  || !bvec[15] )
      && (  bvec[7]  || !bvec[8]  )
      && ( !bvec[7]  || !bvec[13] )
```

```
      && (  bvec[8]  ||  bvec[9]  )
      && (  bvec[8]  || !bvec[9]  )
      && ( !bvec[9]  || !bvec[10] )
      && (  bvec[9]  ||  bvec[11] )
      && (  bvec[10] ||  bvec[11] )
      && (  bvec[12] ||  bvec[13] )
      && (  bvec[13] || !bvec[14] )
      && (  bvec[14] ||  bvec[15] )
      && (  bvec[14] ||  bvec[16] )
      && (  bvec[17] ||  bvec[1]  )
      && (  bvec[18] || !bvec[0]  )
      && (  bvec[19] ||  bvec[1]  )
      && (  bvec[19] || !bvec[18] )
      && ( !bvec[19] || !bvec[9]  )
      && (  bvec[0]  ||  bvec[17] )
      && ( !bvec[1]  ||  bvec[20] )
      && ( !bvec[21] ||  bvec[20] )
      && ( !bvec[22] ||  bvec[20] )
      && ( !bvec[21] || !bvec[20] )
      && (  bvec[22] || !bvec[20] );

  return value;
}
/******************************************************************************
***/

void i4_to_bvec ( int i4, int n, int bvec[] )

/******************************************************************************
***/
/*
  Purpose:

      I4_TO_BVEC converts an integer into a binary vector.

  Licensing:

      This code is distributed under the GNU LGPL license.

  Modified:

      20 March 2009

  Author:

      John Burkardt

  Parameters:

      Input, int I4, the integer.

      Input, int N, the dimension of the vector.

      Output, int BVEC[N], the vector of binary remainders.
*/
{
```

```
  int i;

  for ( i = n - 1; 0 <= i; i-- )
  {
      bvec[i] = i4 % 2;
      i4 = i4 / 2;
  }

  return;
}
/******************************************************************************
***/

void timestamp ( void )

/******************************************************************************
***/
/*
  Purpose:

      TIMESTAMP prints the current YMDHMS date as a time stamp.

  Example:

      31 May 2001 09:45:54 AM

  Licensing:

      This code is distributed under the GNU LGPL license.

  Modified:

      24 September 2003

  Author:

      John Burkardt

  Parameters:

      None
*/
{
# define TIME_SIZE 40

  static char time_buffer[TIME_SIZE];
  //const struct tm *tm;
  //size_t len;
  //time_t now;

  //now = time ( NULL );
  //tm = localtime ( &now );

  //len = strftime ( time_buffer, TIME_SIZE, "%d %B %Y %I:%M:%S %p", tm );

  //printf ( "%s\n", time_buffer );
```

```
  return;
# undef TIME_SIZE
}
```

Figure 10: Serial C program for the circuit-satisfiability problem

```
Please select a parallel programming model from the following available
options:
1. MPI
2. OpenMP
3. CUDA
3

Please enter the function in which you wish to insert the kernel call(or
parallelize the for-loop).
Please choose the function that you want to parallelize from the list below
1 : main
2 : circuit_value
3 : i4_to_bvec
4 : timestamp
1

Would you like
1. For-loop
2. TBD
1

for (i = 1; i <= n; i++) {
  ihi = ihi * 2;
}
Is this the for loop you are looking for?(y/n)
n

OK - will find the next loop if available.

for (i = 1; i <= n / 2; i++) {
  printf("\ntest\n");
}
Is this the for loop you are looking for?(y/n)
n

OK - will find the next loop if available.
Note: With your response, you will be selecting or declining the
parallelization of the outermost for-loop in the code region shown below. If
instead of the outermost for-loop, there are any inner for-loops in this
code region that you are interested in parallelizing, then, you will be able
to select those at a later stage.

for (i = 0; i < ihi; i++) {
  i4_to_bvec(i,n,bvec);
  value = circuit_value(n,bvec);
  if (value == 1) {
      solution_num = solution_num + 1;
      printf("  %2d  %10d:  ",solution_num,i);
```

```
      for (j = 0; j < n; j++) {
      printf(" %d",bvec[j]);
      }
      printf("\n");
  }
}
Is this the for loop you are looking for?(y/n)
y
Do you want to perform reduction on any variables ? (Y/N)
y

Please enter the variables to reduce ([ format:  1,2,3 etc. ] with 1 is for
the first variable, 2 is for the second variable and so on).
 Possible variables to reduce are:
1. i type is int
2. n type is int
3. value type is int
4. solution_num type is int
5. print_statement_deleted_here type is int
4

Please enter the reduction operation for variable [solution_num]. Possible
reduction operations are:
1. Sum
2. Product
1
Do you want to perform reduction on any arrays ? (Y/N)
n
Is the following array [ bvec ]
1. Input , 2. Output 3. Input/Output 4. Neither Input nor Output
4
```

Figure 11: IPT prompts and user responses (bold)

```
# include <stdlib.h>
# include <stdio.h>
# include <time.h>
int main(int argc,char *argv[]);
int circuit_value(int n,int bvec[]);
void i4_to_bvec(int i4,int n,int bvec[]);
void timestamp();
/******************************************************************************
***/
__device__ void i4_to_bvec_rose(int i4,int n,int bvec[]){
int i;
for(i = n - 1;0 <= i;i--) {bvec[i] = i4 % 2;i4 = i4 / 2;}
return ;
}
__device__ int circuit_value_rose(int n,int bvec[]){
int value;
value =(bvec[0] || bvec[1]) &&(!bvec[1] || !bvec[3]) &&(bvec[2] || bvec[3])
&&(!bvec[3] || !bvec[4]) &&(bvec[4] || !bvec[5]) &&(bvec[5] || !bvec[6])
&&(bvec[5] || bvec[6]) &&(bvec[6] || !bvec[15]) &&(bvec[7] || !bvec[8])
&&(!bvec[7] || !bvec[13]) &&(bvec[8] || bvec[9]) &&(bvec[8] || !bvec[9])
&&(!bvec[9] || !bvec[10]) &&(bvec[9] || bvec[11]) &&(bvec[10] || bvec[11])
```

```
&&(bvec[12] || bvec[13]) &&(bvec[13] || !bvec[14]) &&(bvec[14] || bvec[15])
&&(bvec[14] || bvec[16]) &&(bvec[17] || bvec[1]) &&(bvec[18] || !bvec[0])
&&(bvec[19] || bvec[1]) &&(bvec[19] || !bvec[18]) &&(!bvec[19] || !bvec[9])
&&(bvec[0] || bvec[17]) &&(!bvec[1] || bvec[20]) &&(!bvec[21] || bvec[20])
&&(!bvec[22] || bvec[20]) &&(!bvec[21] || !bvec[20]) &&(bvec[22] ||
!bvec[20]);
return value;
}
__global__  void kernel0(int solution_num[],int ihi,int n,int value,int j,
int device_M , int device_N){
int i =  blockIdx.x * blockDim.x + threadIdx.x;
int  print_statement_deleted_here=0;
int bvec[23];
solution_num[i] = 0;
{
i4_to_bvec_rose(i,n,bvec);
value = circuit_value_rose(n,bvec);
if(value == 1) {solution_num[i] = solution_num[i] +
1;print_statement_deleted_here;for(j = 0;j < n;j++)
{print_statement_deleted_here;}print_statement_deleted_here;}
}
__syncthreads();
__syncthreads();
}

int main(int argc,char *argv[])
/******************************************************************************
***/
/*
  Purpose:
      MAIN is the main program for SATISFY.
  Licensing:
      This code is distributed under the GNU LGPL license.
  Modified:
      20 March 2009
  Author:
      John Burkardt
  Reference:
      Michael Quinn,
      Parallel Programming in C with MPI and OpenMP,
      McGraw-Hill, 2004,
      ISBN13: 978-0071232654,
      LC: QA76.73.C15.Q55.
*/
{
  int total_solution_num = 0;
  int *host_solution_num;
  int *device_solution_num;
# define N 23
  int bvec[23];
  int i;
  int ihi;
  int j;
  int n = 23;
  int solution_num;
  int value;
```

```
  printf("\n");
  timestamp();
  printf("\n");
  printf("SATISFY\n");
  printf("  C version\n");
  printf("  We have a logical function of N logical arguments.\n");
  printf("  We do an exhaustive search of all 2^N possibilities,\n");
  printf("  seeking those inputs that make the function TRUE.\n");
/*
  Compute the number of binary vectors to check.
*/
  ihi = 1;
  for (i = 1; i <= n; i++) {
      ihi = ihi * 2;
  }
  for (i = 1; i <= n / 2; i++) {
      printf("\ntest\n");
  }
  printf("\n");
  printf("  The number of logical variables is N = %d\n",n);
  printf("  The number of input vectors to check is %d\n",ihi);
  printf("\n");
  printf("   #    Index ---------Input Values-----------------------\n");
  printf("\n");
/*
  Check every possible input vector.
*/
  solution_num = 0;
//Inserting code for memory allocation grid size and block size calculation
host_solution_num= (int*)malloc((1)*((int) ((ihi - 0 ) / (1)))*sizeof(int));
cudaMalloc((void **) &device_solution_num,( 1)*((int) ((ihi - 0 ) /
(1)))*sizeof(int));
//Please note this is the section wherein the number of blocks and threads
are calculated.  To change the number of threads alter the dimBlock whereas
to change the number of blocks alter the dimGrid
int D_rows = ((int) ((ihi - 0 ) / (1)) > 1024 ) ? (int) ((ihi - 0 ) /
(1))/1024 : (int) ((ihi - 0 ) / (1));
int D_cols = ((int) ((ihi - 0 ) / (1)) > 1024 ) ? 1024 : 1;
dim3 dimGrid(D_rows,1);
dim3 dimBlock(D_cols,1);
kernel0<<<dimGrid,dimBlock>>>(device_solution_num,ihi,n,value,j,1,(int)
((ihi - 0 ) / (1)));
/*
  int IPT_function_replace;
*/
//added this for testing. If the already parallelized loop is not the last
one, then there is no error, else segementation fault
//for ( i = 1; i <= n/2; i++ )
//{
//  printf("\ntestagain\n");
//}
//  Report.
//Copying from Device to Host
cudaMemcpy(host_solution_num,device_solution_num,((int) ((ihi - 0 ) /
(1)))*sizeof(int), cudaMemcpyDeviceToHost);
 //code for variable reduction
```

```
 for(long row = 0; row < 1; ++row){for(long col = 0; col < (int) ((ihi - 0 )
/ (1)); ++col)    {  total_solution_num+= host_solution_num[row*(int) ((ihi
- 0 ) / (1))+ col];     }  }
solution_num+= total_solution_num;
//  Ending Parallelization
  printf("\n");
  printf("  Number of solutions found was %d\n",solution_num);
/*
  Shut down.
*/
  printf("\n");
  printf("SATISFY\n");
  printf("  Normal end of execution.\n");
  printf("\n");
  timestamp();
  return 0;
# undef N
}
/******************************************************************************
***/

int circuit_value(int n,int bvec[])
/******************************************************************************
***/
/*
  Purpose:
      CIRCUIT_VALUE returns the value of a circuit for a given input set.
  Licensing:
      This code is distributed under the GNU LGPL license.
  Modified:
      20 March 2009
  Author:
      John Burkardt
  Reference:
      Michael Quinn,
      Parallel Programming in C with MPI and OpenMP,
      McGraw-Hill, 2004,
      ISBN13: 978-0071232654,
      LC: QA76.73.C15.Q55.
  Parameters:
      Input, int N, the length of the input vector.
      Input, int BVEC[N], the binary inputs.
      Output, int CIRCUIT_VALUE, the output of the circuit.
*/
{
  int value;
  value = (bvec[0] || bvec[1]) && (!bvec[1] || !bvec[3]) && (bvec[2] ||
bvec[3]) && (!bvec[3] || !bvec[4]) && (bvec[4] || !bvec[5]) && (bvec[5] ||
!bvec[6]) && (bvec[5] || bvec[6]) && (bvec[6] || !bvec[15]) && (bvec[7] ||
!bvec[8]) && (!bvec[7] || !bvec[13]) && (bvec[8] || bvec[9]) && (bvec[8] ||
!bvec[9]) && (!bvec[9] || !bvec[10]) && (bvec[9] || bvec[11]) && (bvec[10]
|| bvec[11]) && (bvec[12] || bvec[13]) && (bvec[13] || !bvec[14]) &&
(bvec[14] || bvec[15]) && (bvec[14] || bvec[16]) && (bvec[17] || bvec[1]) &&
(bvec[18] || !bvec[0]) && (bvec[19] || bvec[1]) && (bvec[19] || !bvec[18])
&& (!bvec[19] || !bvec[9]) && (bvec[0] || bvec[17]) && (!bvec[1] ||
bvec[20]) && (!bvec[21] || bvec[20]) && (!bvec[22] || bvec[20]) &&
```

```
(!bvec[21] || !bvec[20]) && (bvec[22] || !bvec[20]);
  return value;
}
/******************************************************************************
***/

void i4_to_bvec(int i4,int n,int bvec[])
/******************************************************************************
***/
/*
  Purpose:
      I4_TO_BVEC converts an integer into a binary vector.
  Licensing:
      This code is distributed under the GNU LGPL license.
  Modified:
      20 March 2009
  Author:
      John Burkardt
  Parameters:
      Input, int I4, the integer.
      Input, int N, the dimension of the vector.
      Output, int BVEC[N], the vector of binary remainders.
*/
{
  int i;
  for (i = n - 1; 0 <= i; i--) {
      bvec[i] = i4 % 2;
      i4 = i4 / 2;
  }
  return ;
}
/******************************************************************************
***/

void timestamp()
/******************************************************************************
***/
/*
  Purpose:
      TIMESTAMP prints the current YMDHMS date as a time stamp.
  Example:
      31 May 2001 09:45:54 AM
  Licensing:
      This code is distributed under the GNU LGPL license.
  Modified:
      24 September 2003
  Author:
      John Burkardt
  Parameters:
      None
*/
{
# define TIME_SIZE 40
//  static char time_buffer[40];
//  const struct tm *tm;
//  size_t len;
```

```
//  time_t now;
//  now = time(((void *)0));
//  tm = (localtime((&now)));
//  len = strftime(time_buffer,40,"%d %B %Y %I:%M:%S %p",tm);
//  printf("%s\n",time_buffer);
  return ;
# undef TIME_SIZE
}
```

Figure 12: Generated CUDA code – circuit-staisfiability

7. Parallelizing Matrix Multiplication

The serial C++ program shown in Figure 13 calculates the product of two matrixes. The steps to parallelize this code are shown in Figire 14, and the generated CUDA code is shown in Figure 15. This code can be run as-is without any modifcations.

```
#include <iostream>
#include <stdio.h>
#include <cstdlib>

using namespace std;

int main(int argc, char** argv){

  //int ** mult = (int**)malloc (sizeof(int)*M);
  //int *a = (int*) malloc (sizeof(int)*M);
  //int *b = (int*) malloc (sizeof(int)*K);
  const int M = 100;
  const int N = 100;
  const int K = 100;

  int mult[M][N];
  int a[M][K];
  int b[K][N];
  //a = &a_t;
  //b = &b_t;
  //mult = &mult_t;
  cout << "here\n";
  for (int i = 0; i < M; ++i) {
      for ( int k = 0; k < K; k++) {
         a[i][k] = i;
      }
  }
  for (int k = 0; k < K; ++k) {

      for ( int j = 0; j < N; j++) {
      b[k][j] = j;
      }
  }
  for (int i = 0; i < M; ++i) {
```

```
      for ( int j = 0; j < N; j++) {
      mult[i][j] = 0;
      }
  }

  for(int i = 0; i < M; ++i) {
      for(int j = 0; j < N; ++j ) {
            for(int k = 0; k < K; ++k)
            {
                  mult[i][j] += a[i][k] * b[k][j];
            }
    }
}
  for (int i = 0; i < M; ++i) {
      for ( int j = 0; j < N; j++) {
                  cout << mult[i][j] << " ";
      }
    cout << "\n";
  }
  return 0;

}
```

Figure 13: C++ program to compute the multiplication of matrices A and B

```
Please select a parallel programming model from the following available
options:
1. MPI
2. OpenMP
3. CUDA
3

Please enter the function in which you wish to insert the kernel call(or
parallelize the for-loop).
Please choose the function that you want to parallelize from the list below
1 : main
1

Would you like
1. For-loop
2. TBD
1
Note: With your response, you will be selecting or declining the
parallelization of the outermost for-loop in the code region shown below. If
instead of the outermost for-loop, there are any inner for-loops in this
code region that you are interested in parallelizing, then, you will be able
to select those at a later stage.

for (int i = 0; i < M; ++i) {
  for (int k = 0; k < K; k++) {
      a[i][k] = i;
  }
}
Is this the for loop you are looking for?(y/n)
n
```

```
OK - will find the next loop if available.

for (int k = 0; k < K; k++) {
  a[i][k] = i;
}
Is this the for loop you are looking for?(y/n)
n

OK - will find the next loop if available.
Note: With your response, you will be selecting or declining the
parallelization of the outermost for-loop in the code region shown below. If
instead of the outermost for-loop, there are any inner for-loops in this
code region that you are interested in parallelizing, then, you will be able
to select those at a later stage.

for (int k = 0; k < K; ++k) {
  for (int j = 0; j < N; j++) {
      b[k][j] = j;
  }
}
Is this the for loop you are looking for?(y/n)
n

OK - will find the next loop if available.

for (int j = 0; j < N; j++) {
  b[k][j] = j;
}
Is this the for loop you are looking for?(y/n)
n

OK - will find the next loop if available.
Note: With your response, you will be selecting or declining the
parallelization of the outermost for-loop in the code region shown below. If
instead of the outermost for-loop, there are any inner for-loops in this
code region that you are interested in parallelizing, then, you will be able
to select those at a later stage.

for (int i = 0; i < M; ++i) {
  for (int j = 0; j < N; j++) {
      mult[i][j] = 0;
  }
}
Is this the for loop you are looking for?(y/n)
n

OK - will find the next loop if available.

for (int j = 0; j < N; j++) {
  mult[i][j] = 0;
}
Is this the for loop you are looking for?(y/n)
n

OK - will find the next loop if available.
```

```
Note: With your response, you will be selecting or declining the
parallelization of the outermost for-loop in the code region shown below. If
instead of the outermost for-loop, there are any inner for-loops in this
code region that you are interested in parallelizing, then, you will be able
to select those at a later stage.

for (int i = 0; i < M; ++i) {
  for (int j = 0; j < N; ++j) {
      for (int k = 0; k < K; ++k) {
      mult[i][j] += a[i][k] * b[k][j];
      }
  }
}
Is this the for loop you are looking for?(y/n)
y
cannot automatically find the number of iteration for the loop, please
specify the number of iteration:
100
Do you want to perform reduction on any variables ? (Y/N)
n
Do you want to perform reduction on any arrays ? (Y/N)
n
Is the following array [ mult ]
1. Input , 2. Output 3. Input/Output 4. Neither Input nor Output
3
Is the following array [ a ]
1. Input , 2. Output 3. Input/Output 4. Neither Input nor Output
1
Is the following array [ b ]
1. Input , 2. Output 3. Input/Output 4. Neither Input nor Output
1
```

Figure 14: IPT prompts and user responses (bold)

```
#include <iostream>
#include <stdio.h>
#include <cstdlib>
using namespace std;
__global__  void kernel0(int * mult,int * a,int * b,const int M,const int
N,const int K, int device_M , int device_N){
int i =  blockIdx.x * blockDim.x + threadIdx.x;
int  print_statement_deleted_here=0;
{
for(int j = 0;j < N;++j) {for(int k = 0;k < K;++k) {mult[j + i *(1 * 100)]
+= a[k + i *(1 * 100)] * b[j + k *(1 * 100)];}}
}
__syncthreads();
}

int main(int argc,char **argv)
{
  int *device_b;
  int *device_a;
  int *device_mult;
//int ** mult = (int**)malloc (sizeof(int)*M);
```

```
//int *a = (int*) malloc (sizeof(int)*M);
//int *b = (int*) malloc (sizeof(int)*K);
  const int M = 100;
  const int N = 100;
  const int K = 100;
  int mult[100][100];
  int a[100][100];
  int b[100][100];
//a = &a_t;
//b = &b_t;
//mult = &mult_t;
  cout<<"here\n";
  for (int i = 0; i < M; ++i) {
      for (int k = 0; k < K; k++) {
      a[i][k] = i;
      }
  }
  for (int k = 0; k < K; ++k) {
      for (int j = 0; j < N; j++) {
      b[k][j] = j;
      }
  }
  for (int i = 0; i < M; ++i) {
      for (int j = 0; j < N; j++) {
      mult[i][j] = 0;
      }
  }
//Please note this is the section wherein the number of blocks and threads
are calculated.  To change the number of threads alter the dimBlock whereas
to change the number of blocks alter the dimGrid
int D_rows = (100 > 1024 ) ? 100/1024 : 100;
int D_cols = (100 > 1024 ) ? 1024 : 1;
dim3 dimGrid(D_rows,1);
dim3 dimBlock(D_cols,1);
cudaMalloc((void **) &device_mult,(100)*(100)*sizeof(int));
for(int rose_i0 = 0; rose_i0 < 100;rose_i0++)
cudaMemcpy(device_mult + rose_i0*100,mult[
rose_i0],(100)*sizeof(int),cudaMemcpyHostToDevice);
cudaMalloc((void **) &device_a,(100)*(100)*sizeof(int));
for(int rose_i0 = 0; rose_i0 < 100;rose_i0++)
cudaMemcpy(device_a + rose_i0*100,a[
rose_i0],(100)*sizeof(int),cudaMemcpyHostToDevice);
cudaMalloc((void **) &device_b,(100)*(100)*sizeof(int));
for(int rose_i0 = 0; rose_i0 < 100;rose_i0++)
cudaMemcpy(device_b + rose_i0*100,b[
rose_i0],(100)*sizeof(int),cudaMemcpyHostToDevice);
kernel0<<<dimGrid,dimBlock>>>(device_mult,device_a,device_b,M,N,K,1,100);
/*
  int IPT_function_replace;
*/
for(int rose_i0 = 0; rose_i0 < 100;rose_i0++)
cudaMemcpy(mult[ rose_i0],device_mult + rose_i0*100,(100)*sizeof(int),
cudaMemcpyDeviceToHost);
cudaFree(device_mult);
cudaFree(device_a);
cudaFree(device_b);
```

```
  for (int i = 0; i < M; ++i) {
      for (int j = 0; j < N; j++) {
      (cout << mult[i][j])<<" ";
      }
      cout<<"\n";
  }
  return 0;
}
```

Figure 15: Generated CUDA Code

8. Parallelizing a Molecular Dynamics Code Using IPT

The Molecular Dynamics (MD) simulation application is used in multiple domains (*e.g.*, biochemistry, biophysics, and material science) for studying the movement of atoms and molecules. Typically, the simulation is run for a fixed period of time during which the atoms and molecules are allowed to interact with each other [1]. During the interaction, the atoms and molecules exert force on each other, in a constrained or an unconstrained manner, thus giving a dynamic view of their interaction over a period of time.

The MD simulation application that we are considering here is designed to follow the path of particles that exert force on each other and are not constrained by any walls [2]. Therefore, after colliding, the particles move past each other. This MD code uses the velocity Verlet time integration scheme and the particles in the simulation interact with a central pair potential [2].

The compute-intensive steps in this application are related to calculating the force and energies (potential energy and kinetic energy) in each time-step, as well as updating the values of the positions, velocities, and accelerations of the particles (*i.e.*, the atoms and molecules) in the simulation. The size of the system (*i.e.,* the number of particles in the simulation), the length of time between the evaluation of energies and force (*i.e.*, the size of the integration step), and the time duration for which the interaction of particles needs to be studied impact the overall runtime of the simulation. In real-world scenarios, often there are several thousand particles in an MD simulation and studying their movement using a serial implementation of the code can take a large amount of time. Therefore, parallel computing is used to reduce the overall run-time of such simulations.

Serial Version of the MD Code and its Parallelization

Let us consider the serial version of the code for the MD simulation application. The complete code for this application is available at [2]. After profiling the code with gprof, we can see that the function named compute is the most time-consuming function and is a good candidate for parallelization. The program spends about 98.9% of its time in this function. A code snippet of the compute function is shown in Figure 1.

```
1. for ( k = 0; k < np; k++ ){
2.     //Compute the potential energy (pe) and forces (f).
3.     for ( i = 0; i < nd; i++ ){
4.       f[i+k*nd] = 0.0;
5.     }
6.     for ( j = 0; j < np; j++ ){
7.       if ( k != j ){
8.          d = dist ( nd, pos+k*nd, pos+j*nd, rij );
9.          if ( d < PI2 ) {
```

```
10.              d2 = d;
11.            } else{
12.              d2 = PI2;}
13.            pe = pe + 0.5 * pow ( sin ( d2 ), 2 );
14.            for ( i = 0; i < nd; i++ ){
15.              f[i+k*nd] = f[i+k*nd] - rij[i] * sin ( 2.0 * d2 ) / d; }
16.         }
17.        }
18.        //Compute the kinetic energy (ke).
19.        for ( i = 0; i < nd; i++ ) {
20.            ke = ke + vel[i+k*nd] * vel[i+k*nd]; }
21.  }
```

Figure 16: Snippet of the compute function – serial version of the MD code

The number of iterations of the for-loop beginning at line # 1 of Figure 16 is equal to the number of particles (`np`) in the simulation. The potential energy (`pe`), force (`f`), and kinetic energy (`ke`) of each particle are calculated in each iteration, and their values across all the iterations are added together. The number of spatial dimensions for the simulation is represented by `nd` and the distance between the particles is represented by d in the code snippet shown in Figure 16. The values of the elements in the displacement vector `rij` are calculated for each particle in the function named `dist` in each iteration of the for-loop beginning at line # 6 of the for-loop. The variable PI2 shown in the code snippet in Figure 1 represents the value of the constant `pi` divided by 2, that is, (3.14/2). Except the loop-variables that are of type integer, all other values in this example are of type double.

There are no dependencies between the iterations of the for-loop at line # 1 of the code in Figure 16. Hence, this for-loop can be parallelized by distributing the computations in the loop across multiple threads or processes. Once all the threads or processes have finished their share of computations and have their local results ready, we can combine those results in a meaningful manner to obtain a global result. This global result should match the results of running the application in the serial mode as closely as possible (some rounding off errors may be permissible). Combining the locally computed values of variables while applying a mathematical operation (*e.g.*, sum or multiplication) is referred to as **reduction**.

In the MD example, if the for-loop at line # 1 of Figure 16 is parallelized, the values of the variables `pe` and `ke` will be computed by each thread or process participating in the computation for a certain number of iterations. However, these values should be collected together and added using reduction to obtain the same result as one would get without parallelization.

Generating the CUDA version of the MD Code using IPT

The steps for generating the CUDA version of the MD code using IPT are shown in Figure 17, and the questions posed by IPT during the process of generating the CUDA version of the MD application are shown in boldface. The details related to the meaning of the questions are available in the IPT documentation.

```
$IPT md.c

NOTE: We currently support only C and C++ programs.

Please select a parallel programming model from the following available options:
1. MPI
2. OpenMP
3. CUDA
3

Please choose the function in which you wish to insert the kernel call (or parallelize
the for-loop).
1 : main
2 : compute
3 : cpu_time
4 : dist
5 : initialize
6 : r8mat_uniform_ab
7 : timestamp
8 : update
2

Would you like to parallelize
1. For-loop
2. TBD
1

Note: With your response, you will be selecting or declining the parallelization of the
outermost for-loop in the code region shown below. If instead of the outermost for-loop,
there are any inner for-loops in this code region that you are interested in
parallelizing, then, you will be able to select those at a later stage.

for (k = 0; k < np; k++) {
  //Compute the potential energy and forces.
  for (i = 0; i < nd; i++) {
    f[i + (k * nd)] = 0.0;
  }
  for (j = 0; j < np; j++) {
    if (k != j) {
      d = dist(nd,(pos + (k * nd)),(pos + (j * nd)),rij);
      //Attribute half of the potential energy to particle J.
      if (d < PI2) {
        d2 = d;
      }
      else {
        d2 = PI2;
      }
      pe = (pe + (0.5 * pow(sin(d2),2)));
      for (i = 0; i < nd; i++) {
        f[i + (k * nd)] = (f[i + (k * nd)] - ((rij[i] * sin((2.0 * d2))) / d));
      }
    }
  }
  //Compute the kinetic energy.
  for (i = 0; i < nd; i++) {
    ke = (ke + (vel[i + (k * nd)] * vel[i + (k * nd)]));
  }
}
Is this the for loop you are looking for?(y/n)
y

Do you want to perform reduction on any variables? (Y/N)
y

Please enter the variables to reduce ([format:  1,2,3 etc. ] where 1 is for the first
variable, 2 is for the second variable and so on). Possible variables to reduce are:
1. nd type is int
2. k type is int
3. np type is int
4. d type is double
5. PI2 type is double
```

The loop shown in Figure 1 is found by IPT

`ke` and `pe` are selected for reduction

```
6. d2 type is double
7. pe type is double
8. ke type is double
7,8

Please enter the reduction operation for variable [pe]. Possible reduction operations
are:
1. Sum
2. Product
1

Please enter the reduction operation for variable [ke]. Possible reduction operations
are:
1. Sum
2. Product
1

Is the following array [ f ]
1. Input , 2. Output 3. Input/Output 4. Neither Input nor Output
3
Please specify the dimensions for this variable (format: [ D1,D2,D3, etc. ] with D1 is
the first dimensions, D2 is the second dimensions and so on). Enter 1 if this is not an
array.
np*nd

Is the following array [ pos ]
1. Input , 2. Output 3. Input/Output 4. Neither Input nor Output
1
Please specify the dimensions for this variable (format: [ D1,D2,D3, etc. ] with D1 is
the first dimensions, D2 is the second dimensions and so on). Enter 1 if this is not an
array.
np*nd

Is the following array [ rij ]
1. Input , 2. Output 3. Input/Output 4. Neither Input nor Output
4

Is the following array [ vel ]
1. Input , 2. Output 3. Input/Output 4. Neither Input nor Output
1
Please specify the dimensions for this variable (format: [ D1,D2,D3, etc. ] with D1 is
the first dimensions, D2 is the second dimensions and so on). Enter 1 if this is not an
array.
nd*np

Running Consistency Tests
```

Reduction with sum operation

Checking if the data needs to be copied to the device (GPU) and back to the host CPU

Figure 17: Steps for generating the CUDA version of the MD code using IPT

After the user follows the steps shown in Figure 17, IPT generates an output file named **rose_md_CUDA.cu** in the current working directory. A snippet of the CUDA code generated by IPT is shown in Figure 18 and the code inserted by IPT is highlighted in boldface. The complete code is provided in figure 19. As can be noticed from the CUDA code in figure 19, IPT reengineers the serial code to create and insert the kernel and device functions, do memory management, perform the reduction operation, and calculate the number of threads using which the CUDA program can be run. It replaces the for-loop selected for parallelization with the call to the CUDA kernel. It also inserts comments in the code to inform the user about the purpose of the changes made to their original code.

```
//Inserting code for memory allocation grid size and block size calculation
host_pe= (double*)malloc((1)*((int) ((np - 0 ) / (1)))*sizeof(double));
cudaMalloc((void **) &device_pe,( 1)*((int) ((np - 0 ) / (1)))*sizeof(double));
```

```
//Inserting code for memory allocation grid size and block size calculation
host_ke= (double*)malloc((1)*((int) ((np - 0 ) / (1)))*sizeof(double));
cudaMalloc((void **) &device_ke,( 1)*((int) ((np - 0 ) / (1)))*sizeof(double));
```

IPT inserts code for memory management on the GPU & calculates the number of threads in each block of a grid using which the code will run on the GPU

```
//Please note this is the section wherein the number of blocks and threads are calculated.  To
change the number of threads alter the dimBlock whereas to change the number of blocks alter the
dimGrid
int D_rows = ((int) ((np - 0 ) / (1)) > 1024 ) ? (int) ((np - 0 ) / (1))/1024 : (int) ((np - 0 )
/ (1));
int D_cols = ((int) ((np - 0 ) / (1)) > 1024 ) ? 1024 : 1;
dim3 dimGrid(D_rows,1);
dim3 dimBlock(D_cols,1);

cudaMalloc((void **) &device_f,(np*nd)*sizeof(double));
cudaMemcpy(device_f,f,(np*nd)*sizeof(double),cudaMemcpyHostToDevice);
cudaMalloc((void **) &device_pos,(np*nd)*sizeof(double));
cudaMemcpy(device_pos,pos,(np*nd)*sizeof(double),cudaMemcpyHostToDevice);
cudaMalloc((void **) &device_vel,(nd*np)*sizeof(double));
cudaMemcpy(device_vel,vel,(nd*np)*sizeof(double),cudaMemcpyHostToDevice);

kernel0<<<dimGrid,dimBlock>>>(device_pe,device_ke,device_f,device_pos,device_vel,np,i,nd,j,d,PI2,
d2,1,(int) ((np - 0 ) / (1)));
```

The loop shown in Figure 1 is replaced by the call to a kernel named **kernel0**

```
/*
  int IPT_function_replace;
*/

//Copying from Device to Host
cudaMemcpy(host_pe,device_pe,((int) ((np - 0 ) / (1)))*sizeof(double), cudaMemcpyDeviceToHost);
```

IPT inserts code for reduction

```
//code for variable reduction
 for(long row = 0; row< 1; ++row){
    for(long col = 0; col <(int)((np - 0 ) / (1));
++col){
      total_pe+= host_pe[row*(int) ((np - 0 ) / (1))+ col];
    }
 }
 pe+= total_pe;

//Copying from Device to Host
cudaMemcpy(host_ke,device_ke,((int) ((np - 0 ) / (1)))*sizeof(double), cudaMemcpyDeviceToHost);
 //code for variable reduction
 for(long row = 0; row < 1; ++row){
    for(long col = 0; col < (int) ((np - 0 ) / (1)); ++col){
       total_ke+= host_ke[row*(int) ((np - 0 ) / (1))+ col];
  }
 }
 ke+= total_ke;
 cudaMemcpy(f,device_f,(np*nd)*sizeof(double), cudaMemcpyDeviceToHost);
 cudaFree(device_f);
 cudaFree(device_pos);
 cudaFree(device_vel);
```

Figure 18: Code snippet of the CUDA version of the MD code generated using IPT

```
$cat rose_md_CUDA.cu

# include <stdlib.h>
# include <stdio.h>
# include <time.h>
# include <math.h>
int main(int argc,char *argv[]);
void compute(int np,int nd,double pos[],double vel[],double mass,double
```

```
f[],double *pot,double *kin);
double cpu_time();
double dist(int nd,double r1[],double r2[],double dr[]);
void initialize(int np,int nd,double pos[],double vel[],double acc[]);
void r8mat_uniform_ab(int m,int n,double a,double b,int *seed,double r[]);
void timestamp();
void update(int np,int nd,double pos[],double vel[],double f[],double
acc[],double mass,double dt);
/**************************************************************************
*/
__device__ double dist_rose(int nd,double r1[],double r2[],double dr[]){
double d;
int i;
d = 0.0;
for(i = 0;i < nd;i++) {dr[i] = r1[i] - r2[i];d = d + dr[i] * dr[i];}
d = sqrt(d);
return d;
}
void __global__ kernel0(double pe[],double ke[],double * f,double * pos,double
* vel,int np,int i,int nd,int j,double d,double PI2,double d2, int device_M ,
int device_N){
int k =  blockIdx.x * blockDim.x + threadIdx.x;
int  print_statement_deleted_here=0;
double rij[3];
pe[k] = 0;
ke[k] = 0;
{
for(i = 0;i < nd;i++) {f[i + k * nd] = 0.0;}
for(j = 0;j < np;j++) {if(k != j) {d = dist_rose(nd,pos + k * nd,pos + j *
nd,rij);if(d < PI2) {d2 = d;}else {d2 = PI2;}pe[k] = pe[k] + 0.5 *
pow((sin(d2)),2);for(i = 0;i < nd;i++) {f[i + k * nd] = f[i + k * nd] - rij[i]
* sin(2.0 * d2) / d;}}}
for(i = 0;i < nd;i++) {ke[k] = ke[k] + vel[i + k * nd] * vel[i + k * nd];}
}
__syncthreads();
__syncthreads();
}

int main(int argc,char *argv[])
/**************************************************************************
*/
/*
  Purpose:
    MAIN is the main program for MD.
  Discussion:
    MD implements a simple molecular dynamics simulation.
    The velocity Verlet time integration scheme is used.
    The particles interact with a central pair potential.
    This program is based on a FORTRAN90 program by Bill Magro.
  Usage:
    md nd np step_num dt
    where:
    * nd is the spatial dimension (2 or 3);
    * np is the number of particles (500, for instance);
    * step_num is the number of time steps (500, for instance).
    * dt is the time step (0.1 for instance)
  Licensing:
```

```
    This code is distributed under the GNU LGPL license.
  Modified:
    27 December 2014
  Author:
    John Burkardt.
*/
{
  double *acc;
  double ctime;
  double dt;
  double e0;
  double *force;
  int i;
  int id;
  double kinetic;
  double mass = 1.0;
  int nd;
  int np;
  double *pos;
  double potential;
  int step;
  int step_num;
  int step_print;
  int step_print_index;
  int step_print_num;
  double *vel;
  timestamp();
  printf("\n");
  printf("MD\n");
  printf("  C version\n");
  printf("  A molecular dynamics program.\n");
/*
  Get the spatial dimension.
*/
  if (1 < argc) {
    nd = atoi(argv[1]);
  }
   else {
    printf("\n");
    printf("  Enter ND, the spatial dimension (2 or 3).\n");
    scanf("%d",&nd);
  }
//
//  Get the number of particles.
//
  if (2 < argc) {
    np = atoi(argv[2]);
  }
   else {
    printf("\n");
    printf("  Enter NP, the number of particles (500, for instance).\n");
    scanf("%d",&np);
  }
//
//  Get the number of time steps.
//
  if (3 < argc) {
```

```
    step_num = atoi(argv[3]);
  }
   else {
    printf("\n");
    printf("  Enter ND, the number of time steps (500 or 1000, for
instance).\n");
    scanf("%d",&step_num);
  }
//
//  Get the time steps.
//
  if (4 < argc) {
    dt = atof(argv[4]);
  }
   else {
    printf("\n");
    printf("  Enter DT, the size of the time step (0.1, for instance).\n");
    scanf("%g",&dt);
  }
/*
  Report.
*/
  printf("\n");
  printf("  ND, the spatial dimension, is %d\n",nd);
  printf("  NP, the number of particles in the simulation, is %d\n",np);
  printf("  STEP_NUM, the number of time steps, is %d\n",step_num);
  printf("  DT, the size of each time step, is %f\n",dt);
/*
  Allocate memory.
*/
  acc = ((double *)(malloc((nd * np) * sizeof(double ))));
  force = ((double *)(malloc((nd * np) * sizeof(double ))));
  pos = ((double *)(malloc((nd * np) * sizeof(double ))));
  vel = ((double *)(malloc((nd * np) * sizeof(double ))));
/*
  This is the main time stepping loop:
    Compute forces and energies,
    Update positions, velocities, accelerations.
*/
  printf("\n");
  printf("  At each step, we report the potential and kinetic energies.\n");
  printf("  The sum of these energies should be a constant.\n");
  printf("  As an accuracy check, we also print the relative error\n");
  printf("  in the total energy.\n");
  printf("\n");
  printf("      Step      Potential       Kinetic        (P+K-E0)/E0\n");
  printf("                Energy P        Energy K       Relative Energy
Error\n");
  printf("\n");
  step_print = 0;
  step_print_index = 0;
  step_print_num = 10;
  ctime = cpu_time();
  for (step = 0; step <= step_num; step++) {
    if (step == 0) {
      initialize(np,nd,pos,vel,acc);
    }
```

```
     else {
       update(np,nd,pos,vel,force,acc,mass,dt);
     }
     compute(np,nd,pos,vel,mass,force,&potential,&kinetic);
     if (step == 0) {
       e0 = potential + kinetic;
     }
     if (step == step_print) {
       printf("  %8d  %14f  %14f  %14e\n",step,potential,kinetic,(potential +
kinetic - e0) / e0);
       step_print_index = step_print_index + 1;
       step_print = step_print_index * step_num / step_print_num;
//printf("\nforce = ");
/*for (int ttt = 0; ttt < nd*np; ttt++) {
          printf("%lf ",force[ttt]);
       }*/
// printf("\n\n");
     }
   }
/*
  Report timing.
*/
  ctime = cpu_time() - ctime;
  printf("\n");
  printf("  Elapsed cpu time: %f seconds.\n",ctime);
/*
  Free memory.
*/
  free(acc);
  free(force);
  free(pos);
  free(vel);
/*
  Terminate.
*/
  printf("\n");
  printf("MD\n");
  printf("  Normal end of execution.\n");
  printf("\n");
  timestamp();
  return 0;
}
/******************************************************************************
*/

void compute(int np,int nd,double pos[],double vel[],double mass,double
f[],double *pot,double *kin)
/******************************************************************************
*/
/*
  Purpose:
    COMPUTE computes the forces and energies.
  Discussion:
    The computation of forces and energies is fully parallel.
    The potential function V(X) is a harmonic well which smoothly
    saturates to a maximum value at PI/2:
      v(x) = ( sin ( min ( x, PI/2 ) ) )^2
```

```
    The derivative of the potential is:
      dv(x) = 2.0 * sin ( min ( x, PI/2 ) ) * cos ( min ( x, PI/2 ) )
            = sin ( 2.0 * min ( x, PI/2 ) )
  Licensing:
    This code is distributed under the GNU LGPL license.
  Modified:
    21 November 2007
  Author:
    John Burkardt.
  Parameters:
    Input, int NP, the number of particles.
    Input, int ND, the number of spatial dimensions.
    Input, double POS[ND*NP], the positions.
    Input, double VEL[ND*NP], the velocities.
    Input, double MASS, the mass of each particle.
    Output, double F[ND*NP], the forces.
    Output, double *POT, the total potential energy.
    Output, double *KIN, the total kinetic energy.
*/
{
  double *device_vel;
  double *device_pos;
  double *device_f;
  double total_ke = 0;
  double *host_ke;
  double *device_ke;
  double total_pe = 0;
  double *host_pe;
  double *device_pe;
  double d;
  double d2;
  int i;
  int j;
  int k;
  double ke;
  double pe;
  double PI2 = 3.141592653589793 / 2.0;
  double rij[3];
  pe = 0.0;
  ke = 0.0;
//Inserting code for memory allocation grid size and block size calculation
host_pe= (double*)malloc((1)*((int) ((np - 0 ) / (1)))*sizeof(double));
cudaMalloc((void **) &device_pe,( 1)*((int) ((np - 0 ) /
(1)))*sizeof(double));
//Inserting code for memory allocation grid size and block size calculation
host_ke= (double*)malloc((1)*((int) ((np - 0 ) / (1)))*sizeof(double));
cudaMalloc((void **) &device_ke,( 1)*((int) ((np - 0 ) /
(1)))*sizeof(double));
//Please note this is the section wherein the number of blocks and threads are
calculated.  To change the number of threads alter the dimBlock whereas to
change the number of blocks alter the dimGrid
int D_rows = ((int) ((np - 0 ) / (1)) > 1024 ) ? (int) ((np - 0 ) / (1))/1024
: (int) ((np - 0 ) / (1));
int D_cols = ((int) ((np - 0 ) / (1)) > 1024 ) ? 1024 : 1;
dim3 dimGrid(D_rows,1);
dim3 dimBlock(D_cols,1);
cudaMalloc((void **) &device_f,(np*nd)*sizeof(double));
```

```
cudaMemcpy(device_f,f,(np*nd)*sizeof(double),cudaMemcpyHostToDevice);
cudaMalloc((void **) &device_pos,(np*nd)*sizeof(double));
cudaMemcpy(device_pos,pos,(np*nd)*sizeof(double),cudaMemcpyHostToDevice);
cudaMalloc((void **) &device_vel,(np*nd)*sizeof(double));
cudaMemcpy(device_vel,vel,(np*nd)*sizeof(double),cudaMemcpyHostToDevice);
kernel0<<<dimGrid,dimBlock>>>(device_pe,device_ke,device_f,device_pos,device_v
el,np,i,nd,j,d,PI2,d2,1,(int) ((np - 0 ) / (1)));
/*
  int IPT_function_replace;
*/
//Copying from Device to Host
cudaMemcpy(host_pe,device_pe,((int) ((np - 0 ) / (1)))*sizeof(double),
cudaMemcpyDeviceToHost);
 //code for variable reduction
 for(long row = 0; row < 1; ++row){for(long col = 0; col < (int) ((np - 0 ) /
(1)); ++col)    {  total_pe+= host_pe[row*(int) ((np - 0 ) / (1))+ col];    }
}
pe+= total_pe;
//  Ending Parallelization
//Copying from Device to Host
cudaMemcpy(host_ke,device_ke,((int) ((np - 0 ) / (1)))*sizeof(double),
cudaMemcpyDeviceToHost);
 //code for variable reduction
 for(long row = 0; row < 1; ++row){for(long col = 0; col < (int) ((np - 0 ) /
(1)); ++col)    {  total_ke+= host_ke[row*(int) ((np - 0 ) / (1))+ col];    }
}
ke+= total_ke;
//  Ending Parallelization
cudaMemcpy(f,device_f,(np*nd)*sizeof(double), cudaMemcpyDeviceToHost);
cudaFree(device_f);
cudaFree(device_pos);
cudaFree(device_vel);
  ke = ke * 0.5 * mass;
   *pot = pe;
   *kin = ke;
  return ;
}
/*****************************************************************************
**/

double cpu_time()
/*****************************************************************************
**/
/*
  Purpose:

    CPU_TIME reports the total CPU time for a program.
  Licensing:
    This code is distributed under the GNU LGPL license.
  Modified:
    27 September 2005
  Author:
    John Burkardt
  Parameters:
    Output, double CPU_TIME, the current total elapsed CPU time in second.
*/
{
```

```
  double value;
  value = ((double )(clock())) / ((double )10000001);
  return value;
}
/******************************************************************************
*/

double dist(int nd,double r1[],double r2[],double dr[])
/******************************************************************************
*/
/*
  Purpose:
    DIST computes the displacement (and its norm) between two particles.
  Licensing:
    This code is distributed under the GNU LGPL license.
  Modified:
    21 November 2007
  Author:
    John Burkardt.
  Parameters:
    Input, int ND, the number of spatial dimensions.
    Input, double R1[ND], R2[ND], the positions of the particles.
    Output, double DR[ND], the displacement vector.
    Output, double D, the Euclidean norm of the displacement.
*/
{
  double d;
  int i;
  d = 0.0;
  for (i = 0; i < nd; i++) {
    dr[i] = r1[i] - r2[i];
    d = d + dr[i] * dr[i];
  }
  d = sqrt(d);
  return d;
}
/******************************************************************************
*/

void initialize(int np,int nd,double pos[],double vel[],double acc[])
/******************************************************************************
*/
/*
  Purpose:
    INITIALIZE initializes the positions, velocities, and accelerations.
  Licensing:
    This code is distributed under the GNU LGPL license.
  Modified:
    26 December 2014
  Author:
    John Burkardt.
  Parameters:
    Input, int NP, the number of particles.
    Input, int ND, the number of spatial dimensions.
    Output, double POS[ND*NP], the positions.
    Output, double VEL[ND*NP], the velocities.
    Output, double ACC[ND*NP], the accelerations.
```

```
*/
{
  int i;
  int j;
  int seed;
/*
  Set positions.
*/
  seed = 123456789;
  r8mat_uniform_ab(nd,np,0.0,10.0,&seed,pos);
/*
  Set velocities.
*/
  for (j = 0; j < np; j++) {
    for (i = 0; i < nd; i++) {
      vel[i + j * nd] = 0.0;
    }
  }
/*
  Set accelerations.
*/
  for (j = 0; j < np; j++) {
    for (i = 0; i < nd; i++) {
      acc[i + j * nd] = 0.0;
    }
  }
  return ;
}
/******************************************************************************
*/

void r8mat_uniform_ab(int m,int n,double a,double b,int *seed,double r[])
/******************************************************************************
*/
/*
  Purpose:
    R8MAT_UNIFORM_AB returns a scaled pseudorandom R8MAT.
  Discussion:
    This routine implements the recursion
      seed = 16807 * seed mod ( 2^31 - 1 )
      unif = seed / ( 2^31 - 1 )
    The integer arithmetic never requires more than 32 bits,
    including a sign bit.
  Licensing:
    This code is distributed under the GNU LGPL license.
  Modified:
    03 October 2005
  Author:
    John Burkardt
  Reference:
    Paul Bratley, Bennett Fox, Linus Schrage,
    A Guide to Simulation,
    Second Edition,
    Springer, 1987,
    ISBN: 0387964673,
    LC: QA76.9.C65.B73.
    Bennett Fox,
```

```
    Algorithm 647:
    Implementation and Relative Efficiency of Quasirandom
    Sequence Generators,
    ACM Transactions on Mathematical Software,
    Volume 12, Number 4, December 1986, pages 362-376.
    Pierre L'Ecuyer,
    Random Number Generation,
    in Handbook of Simulation,
    edited by Jerry Banks,
    Wiley, 1998,
    ISBN: 0471134031,
    LC: T57.62.H37.
    Peter Lewis, Allen Goodman, James Miller,
    A Pseudo-Random Number Generator for the System/360,
    IBM Systems Journal,
    Volume 8, Number 2, 1969, pages 136-143.
  Parameters:
    Input, int M, N, the number of rows and columns.
    Input, double A, B, the limits of the pseudorandom values.
    Input/output, int *SEED, the "seed" value.  Normally, this
    value should not be 0.  On output, SEED has
    been updated.
    Output, double R[M*N], a matrix of pseudorandom values.
*/
{
  int i;
  const int i4_huge = 2147483647;
  int j;
  int k;
  if ( *seed == 0) {
    fprintf(stderr,"\n");
    fprintf(stderr,"R8MAT_UNIFORM_AB - Fatal error!\n");
    fprintf(stderr,"  Input value of SEED = 0.\n");
    exit(1);
  }
  for (j = 0; j < n; j++) {
    for (i = 0; i < m; i++) {
      k =  *seed / 127773;
       *seed = 16807 * ( *seed - k * 127773) - k * 2836;
      if ( *seed < 0) {
         *seed =  *seed + i4_huge;
      }
      r[i + j * m] = a + (b - a) * ((double )( *seed)) * 4.656612875E-10;
    }
  }
  return ;
}
/*****************************************************************************
*/

void timestamp()
/*****************************************************************************
*/
/*
  Purpose:
    TIMESTAMP prints the current YMDHMS date as a time stamp.
  Example:
```

```
    31 May 2001 09:45:54 AM
  Licensing:
    This code is distributed under the GNU LGPL license.
  Modified:
    24 September 2003
  Author:
    John Burkardt
  Parameters:
    None
*/
{
# define TIME_SIZE 40
  static char time_buffer[TIME_SIZE];
  const struct tm *tm;
  size_t len;
  time_t now;
  now = time ( NULL );
  tm = localtime ( &now );
  len = strftime ( time_buffer, TIME_SIZE, "%d %B %Y %I:%M:%S %p", tm );
  printf ( "%s\n", time_buffer );
  return ;
# undef TIME_SIZE
}
/******************************************************************************
*/

void update(int np,int nd,double pos[],double vel[],double f[],double
acc[],double mass,double dt)
/******************************************************************************
*/
/*
  Purpose:
    UPDATE updates positions, velocities and accelerations.
  Discussion:
    The time integration is fully parallel.
    A velocity Verlet algorithm is used for the updating.
    x(t+dt) = x(t) + v(t) * dt + 0.5 * a(t) * dt * dt
    v(t+dt) = v(t) + 0.5 * ( a(t) + a(t+dt) ) * dt
    a(t+dt) = f(t) / m
  Licensing:
    This code is distributed under the GNU LGPL license.
  Modified:
    21 November 2007
  Author:
    John Burkardt.
  Parameters:
    Input, int NP, the number of particles.
    Input, int ND, the number of spatial dimensions.
    Input/output, double POS[ND*NP], the positions.
    Input/output, double VEL[ND*NP], the velocities.
    Input, double F[ND*NP], the forces.
    Input/output, double ACC[ND*NP], the accelerations.
    Input, double MASS, the mass.
    Input, double DT, the time step.
*/
{
  int i;
```

```
  int j;
  double rmass;
  rmass = 1.0 / mass;
  for (j = 0; j < np; j++) {
    for (i = 0; i < nd; i++) {
      pos[i + j * nd] = pos[i + j * nd] + vel[i + j * nd] * dt + 0.5 * acc[i +
j * nd] * dt * dt;
      vel[i + j * nd] = vel[i + j * nd] + 0.5 * dt * (f[i + j * nd] * rmass +
acc[i + j * nd]);
      acc[i + j * nd] = f[i + j * nd] * rmass;
    }
  }
  return ;
}
```

Figure 19: CUDA version of the MD code generated using IPT

Compiling and running the serial version

The serial version of the MD application can be compiled using the command:

```
$ icc -o md md.c
```

The executable "md" of the serial version can be run as follows on a compute node:

```
./md 2 1000 1000 0.01
```

In the aforementioned command, we provided 4 arguments that specify:

```
ND, the spatial dimension, which is 2
NP, the number of particles in the simulation, which is 1000
STEP_NUM, the number of time steps, which is 1000
DT, the size of each time step, is which 0.010000
```

Compiling and running the CUDA version

To compile and run the CUDA code generated by IPT, you would need access to a GPU. If you are using the Lonestar5 system, then you may need to load the "cuda" module.

The following commands show the steps for loading the "cuda" module, and compiling the code using the `nvcc` command:

```
$ module load cuda
$ nvcc -o rose_md_cuda rose_md_CUDA.cu
```

Then run the executable on the `gpu` queue on Lonestar 5:

```
$ nvcc -o rose_md_cuda rose_md_CUDA.cu
```

The executable can be run as follows on the compute node in the "gpu" queue:

```
./rose_md_cuda 2 1000 1000 0.01
```

9. CUDA Code Optimizations

A key design point that will help program optimization is to write kernels that look like serial programs and can be run on only one thread. While it is computationally inexpensive to launch and run thousands of threads on a GPU, interdependencies between threads will slow-down a program. It is best to design the kernel like a generalized serial program that is run on a single thread.

Data transfer between the CPU and GPU is essential but this process can be costly in terms of time and resources. Therefore, it is important to minimize the data transfers between the CPU and GPU. This means that the ratio of computation to communication should be kept as high as possible. This can be achieved by ensuring that the computations performed by the GPU are not trivial and are worth running on the device.

A related factor to consider when making performance modifications is maximizing arithmetic density. Arithmetic density is the ratio of math:memory time. This ratio can be increased by maximizing work and computation operations per thread or minimizing memory time spent on memory access per thread. Time can be minimized on memory access using faster memory for frequently accessed data. Local memory is the fastest to access, followed by shared, with global being the slowest. Another way to decrease the time spent on memory access is to use coalesced memory for global memory access. Time used for memory access will be minimal when the threads read from and write to contiguous memory locations.

Similarly, when launching the kernels in a thread, it is also important to be mindful of thread divergence. Thread divergence occurs when some threads may execute a statement that others do not due to logical branching. Divergence is most often caused by if-statements and loops. Divergence hurts performance because of implicit barriers that prevent threads from moving on. Consider a group of 32 threads running in a warp with identification numbers from 0-31. If these threads are running on a kernel with a loop that runs as many times as the square thread's identification number, then threads with lower identification number will finish much faster than the threads with higher identification numbers. The thread with identification number 0 will run through the code once, while thread with identification number 32 will run the code 1024 times. All threads in the warp will wait until the last thread finishes. While this example is rather extreme it demonstrates how thread divergence can impact the performance.

10. Common Operations

CUDA programs allow a few parallel communication models. These operations are gather, scatter, map, stencil, and transpose. Gather is a many-to-one operation where threads read multiple elements to produce a single output. Scatter is a one-to-many operation where an output is written to many memory locations. Map is a one-to-one pattern where read and write operations only occur at a specific location. Stencil is a specific type of gather operation where the output is dependent upon the neighboring data points. Transpose is a reordering tool to change how data is arranged in memory. It can be useful to simplify data structures for calculations such as flattening matrices into arrays.

11. Synchronization

Apart from communication between threads, CUDA also offers the possibility of synchronizing them. CUDA threads running in a kernel can be programmed to do so until reaching a barrier - called by `__syncthreads()`, where they are all forced to wait until all threads reach the barrier. Barriers help in preventing threads from reading from or writing to memory locations before the program is ready for them. However one must be wary of using too many synchronization barriers because it does slow-down the program since some threads will become idle in the waiting processes.

Similarly, there is an implicit barrier between two kernels. This means that if a program calls kernel A and then kernel B, the program will finish kernel A before moving on to kernel B. If the programmer wants to add in an explicit barrier, `__syncthreads()` can be used. Recall that threads are run in groups of 32 called warps. In order to synchronize the threads in a warp, one needs to call `__syncthreads()`. Note that `__syncthreads()` will only synchronize threads in the same block and does not synchronize at the grid or device level.

Another way to control threads is to use atomic operations. Atomic operations help prevent memory collisions where threads read and write from the same locations at the same time. When memory collisions occur, data can be corrupted resulting in incorrect results. Atomic operations serialize memory accesses so that only one thread can access for reading or writing at a time. While atomic actions are useful for preventing memory collisions, serialization in this manner will decrease program speed.

12. References

[1] Molecular Dynamics (MD) Simulation Application. Website accessed on August 10, 2017. https://en.wikipedia.org/wiki/Molecular_dynamics

[2] Molecular Dynamics (MD) code. Website accessed on August 10, 2017. https://people.sc.fsu.edu/~jburkardt/c_src/md/md.html