

TUTORIAL: 3

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1. int LinearSearch (int arr, int n, int key)
    for (i = 0 to n-1)
        if (arr[i] == key)
            return i
    return -1;

```

```

2. insertion sort:
    iterative:
        for j = 2 to n,
            key = a[j]    i = j-1;
            while i > 0 and a[i] > key
                a[i+1] = a[i];
                i = i-1;
            a[i+1] = key;

```

3. Complexity for sort algorithms are:

1. Bubble sort

Best =  $O(n^2)$   
 Worst =  $O(n^2)$   
 Avg =  $O(n^2)$

2. Selection sort

Best =  $O(n^2)$   
 Worst =  $O(n^2)$   
 Avg =  $O(n^2)$

Insertion sort

Best =  $O(n)$   
 Worst =  $O(n^2)$   
 Avg =  $O(n^2)$

merge sort

Best =  $O(n \log n)$   
 Worst =  $O(n \log n)$   
 Avg =  $O(n \log n)$

Quick sortBest =  $O(n \log n)$ Worst =  $O(n^2)$ Avg =  $O(n \log n)$ Counting sortBest =  $O(n+k)$ Worst =  $O(n+k)$ Avg =  $O(n+k)$ 7 Heap sortBest =  $O(n \log n)$ Avg =  $O(n \log n)$ Worst =  $O(n \log n)$ 4 Iterative pseudo codeBinary search

int binarysearch (int arr, int l, int r, int x)

{  
    while (l <= r)

{

int m = (l + r) / 2;

if (arr[m] == x)

return m;

if (arr[m] &lt; x)

l = m + 1;

} else

r = m - 1;

}

return -1;

}

Time ComplexityBest Case  $\rightarrow O(1)$ Worst Case  $\rightarrow O(\log n)$ Avg Case  $\rightarrow O(\log n)$



5. Inplace Sorting: Bubble, Selection, Insertion, heap, Quick sort.

Online sort: self; insertion

Stable sort  $\rightarrow$  Bubble, insertion, merge, Counting

6. Quick sort is mostly used in practical because it is the fastest general purpose sort. In most practical situation quick sort is the method of choice if stability is important and space is available merge sort might be best.

7.  $T(n) = T(n/2) + 1$  put  $n = n/2$

$$T(n/2) = T(n/4) + 1$$

$$T(n) = T(n/4) + 1 + 1$$

$$T(n/4) = T(n/8) + 1$$

$$T(n) = T(n/8) + 1 + 1 + 1$$

$$T(n) = T(n/2^k) + k$$

$$\frac{n}{2^k} = 1 \Rightarrow 2^k = n, \Rightarrow k = \log_2 n$$

$$T(n) = T\left(\frac{n}{2^{\log_2 n}}\right) + \log_2 n$$

$$T(1) + \log_2 n$$

$$1 + \log_2 n = T(n) = O(\log_2 n)$$

11. Best Case of Quick sort.

$$T(n) = 2T(n/2) + n$$

$$a = 2 \quad b = 2 \quad F(n) = n$$

$$= n (\log_b a) \Rightarrow n \log_2 2 = n$$

$$F(n) = n$$

$$F(n) = a (n \log_b a \log_k n)$$

$$F(n) = \Theta(n \log_k n) = \Theta(n \log n)$$

Worst case of Quick sort

$$T(n) = T(n-1) + n \quad n = n-1,$$

$$T(n-1) = T(n-2) + n-1,$$

$$T(n) = T(n-2) + (n-1) + n$$

$$\text{For } n = n-2$$

$$T(n-2) = T(n-3) + n-2$$

$$T(n) = T(n-3) + (n-2) + (n-1) + n,$$

$$T(n) = T(n-k) + (n-(k-1)) + (n-(k-2))$$

$$T(k) = T(n-k)$$

$$n = k+1,$$

$$k = n-1$$

$$T(n) = 1 + 2 + 3 + \dots + n.$$

$$T(n) = \frac{n(n+1)}{2} = O(n^2)$$

~~T(n)~~

$$T(n) = \underline{\underline{O(n^2)}}$$



Best case of merge sort:

$$T(n) = 2T(n/2) + n$$

$$a=2, b=2, f(n)=n,$$

$$n \log_b a = n \log_2 2^2 = n$$

$$f(n) = O(n \log_b a \cdot \log_k n)$$

$$f(n) = O(n \log_k n)$$

$$f(n) = O(n \log n),$$

Same as worst case:

Simultaneously is Best Case Complexities of merge sort and quick sort both are same

$$T(n) = O(n \log n) \text{ But worst case}$$

Complexity of both the sorting are difficult. merge sort =  $O(n \log n)$  &  
Quick sort =  $O(n^2)$