

BREAKING SAMSUNG'S ARM TRUSTZONE

Quarkslab

Black Hat USA 2019

OUR TEAM



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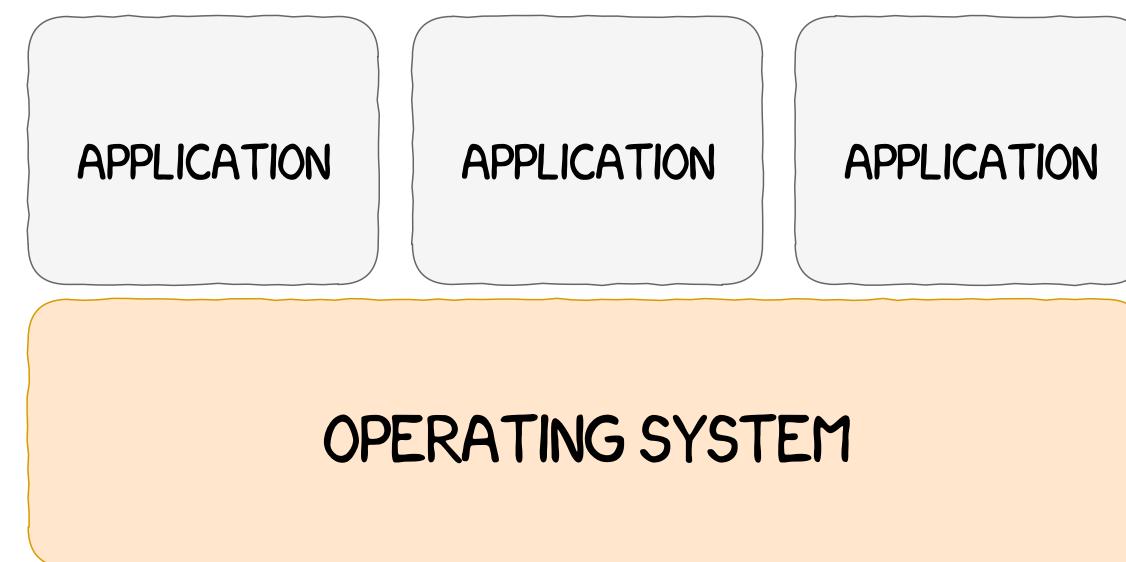
PRESENTATION OUTLINE

- Current state of embedded security
- Introduction to the ARM TrustZone technology
- Samsung's TrustZone Overview
- Trusted Components
- Vulnerability Research Tools
- Vulnerability Analysis
- Exploitation
- Post-Exploitation Demonstrations

CURRENT STATE OF EMBEDDED SECURITY

A LONG TIME AGO...

TRADITIONAL ARCHITECTURE



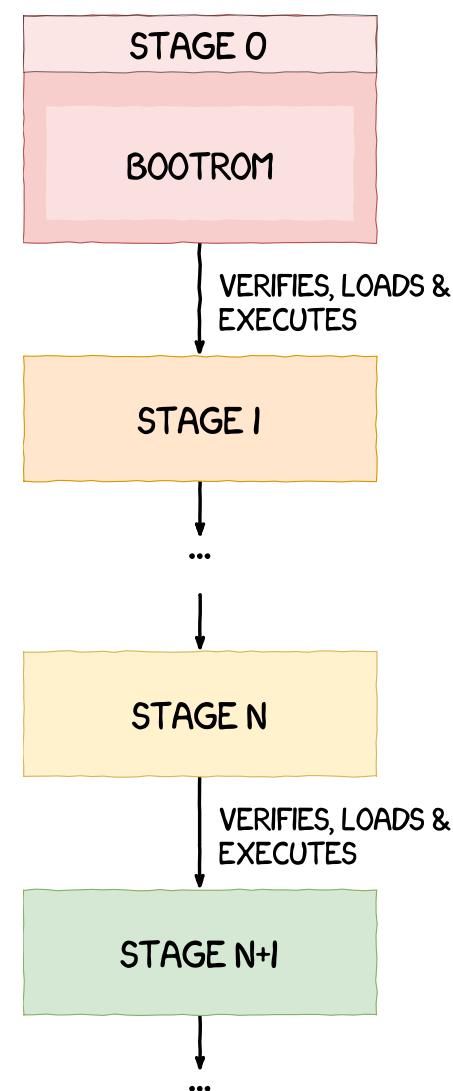
- Kernel unbreakable...?

HOW DO WE PROTECT OURSELVES...

- ... if the kernel is corrupted during the boot process?
- ... if the kernel is corrupted when the system is already running?

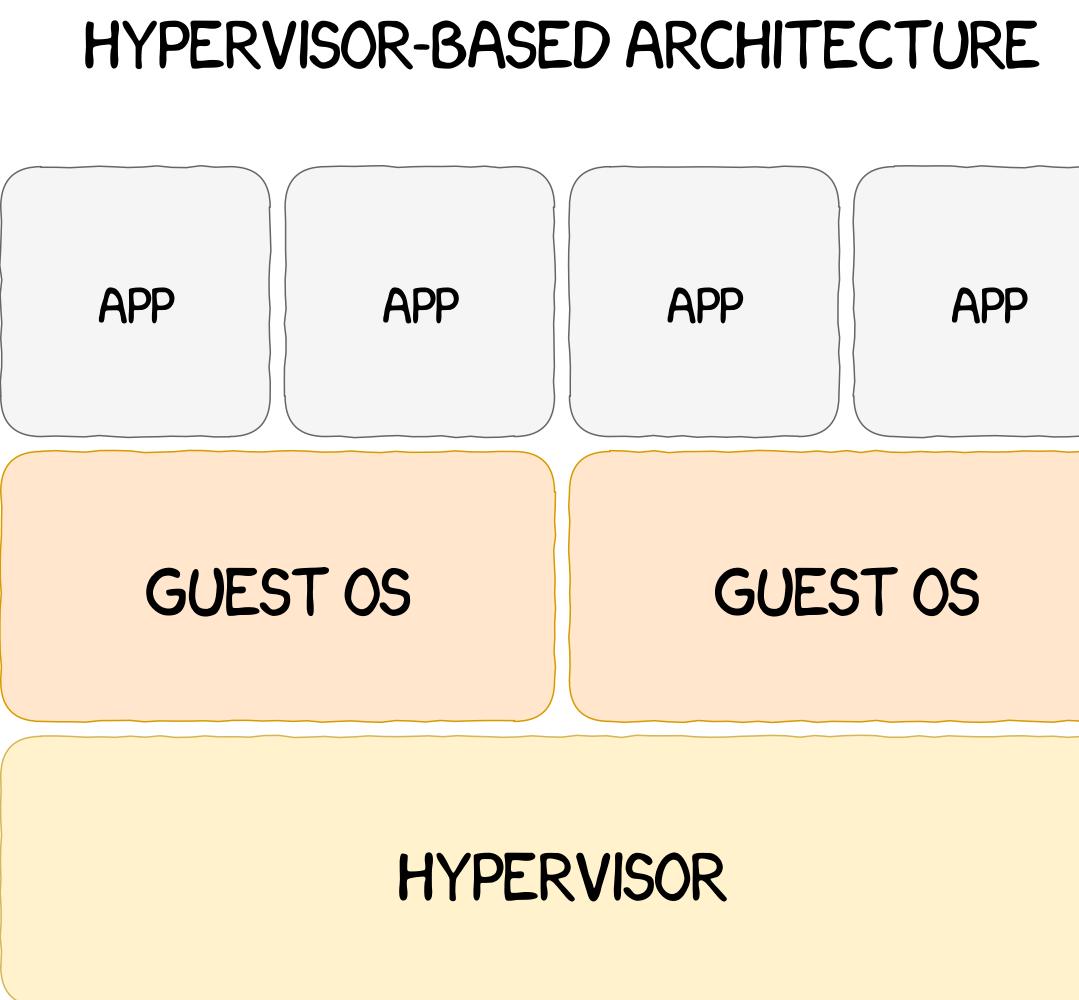
PROTECTION DURING THE BOOT PROCESS

Secure Boot



- Prevent the execution of untrusted or unauthorized code on end users devices

RUNTIME PROTECTION USING AN HYPERVISOR

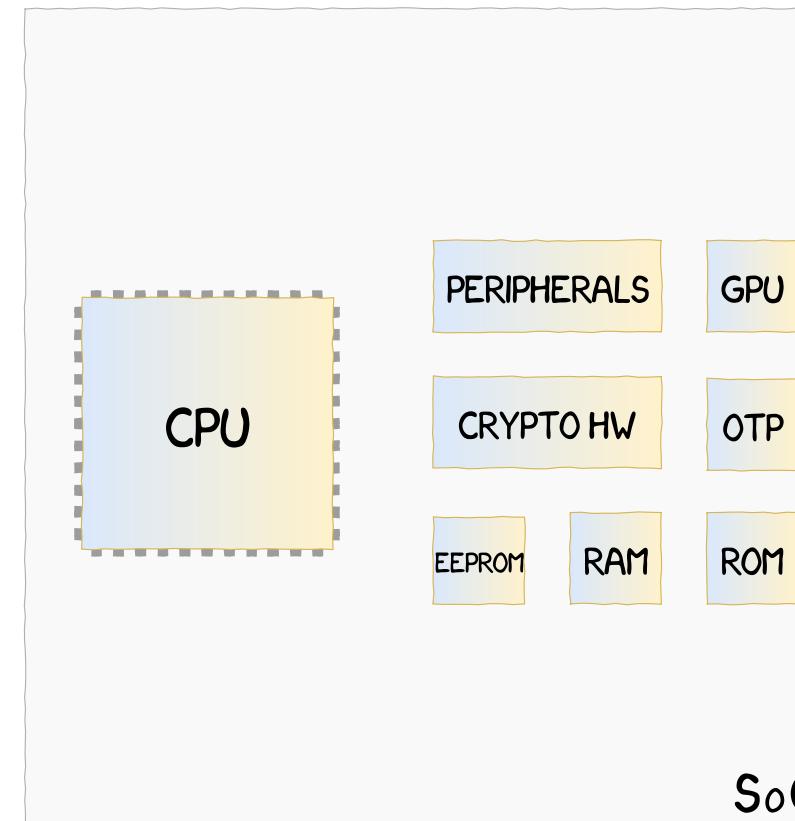


- Hypervisor based guest kernel protection
- **Problem:** VM escapes and hypervisor compromissions

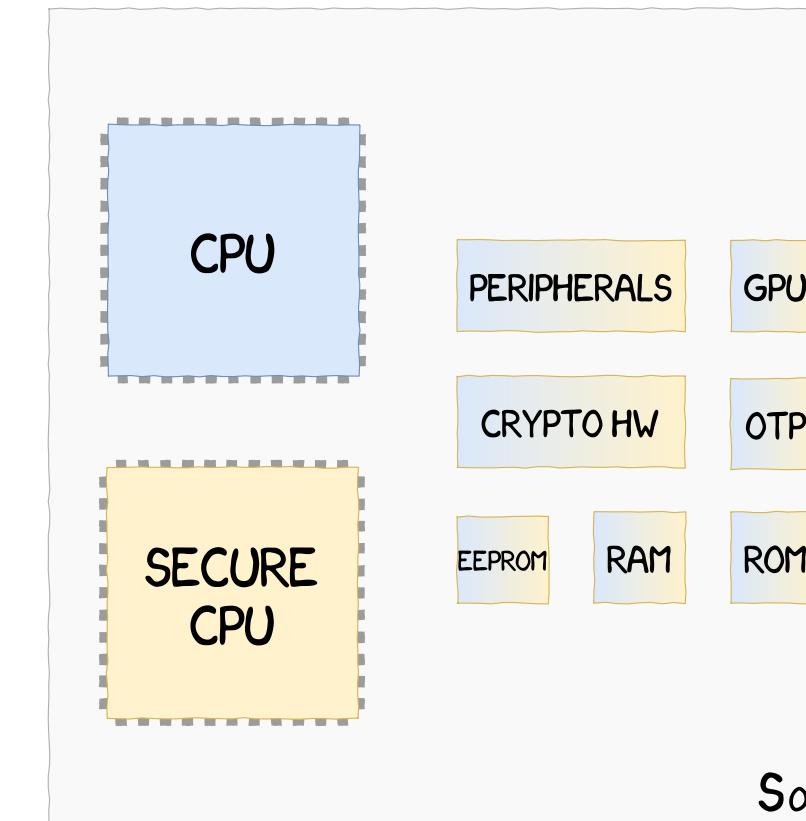
TRUSTED EXECUTION ENVIRONMENTS

Taken from *Le TEE, nouvelle ligne de défense dans les mobiles*, SSTIC 2013

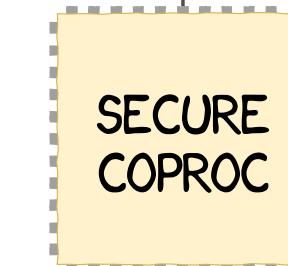
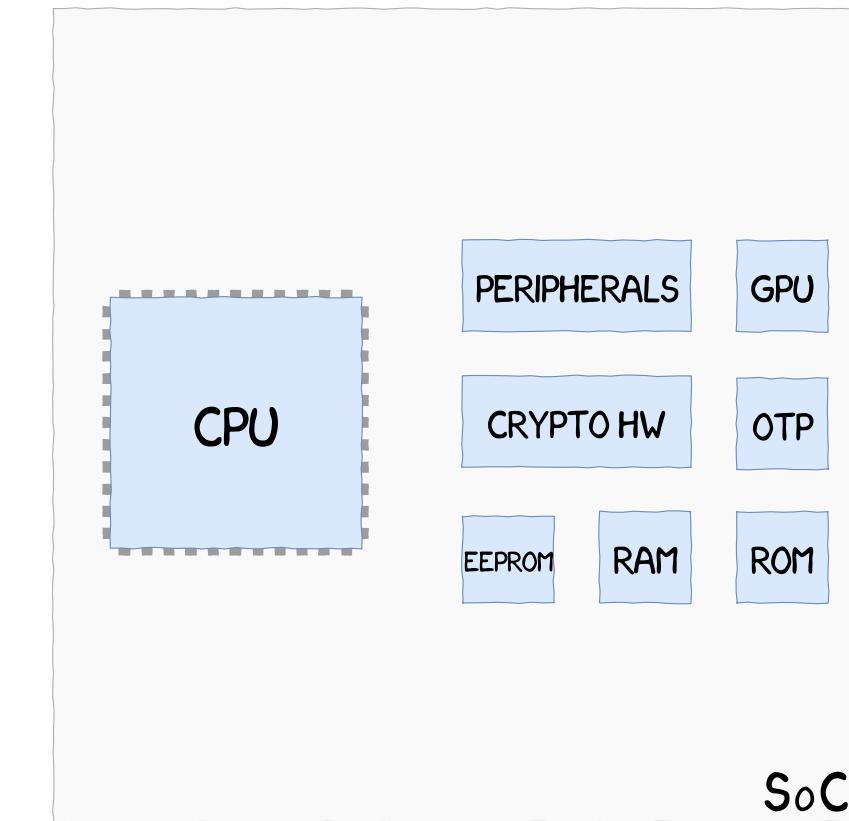
Virtual Processor (e.g. ARM TrustZone)



On-SoC Processor (e.g. Apple SEP)



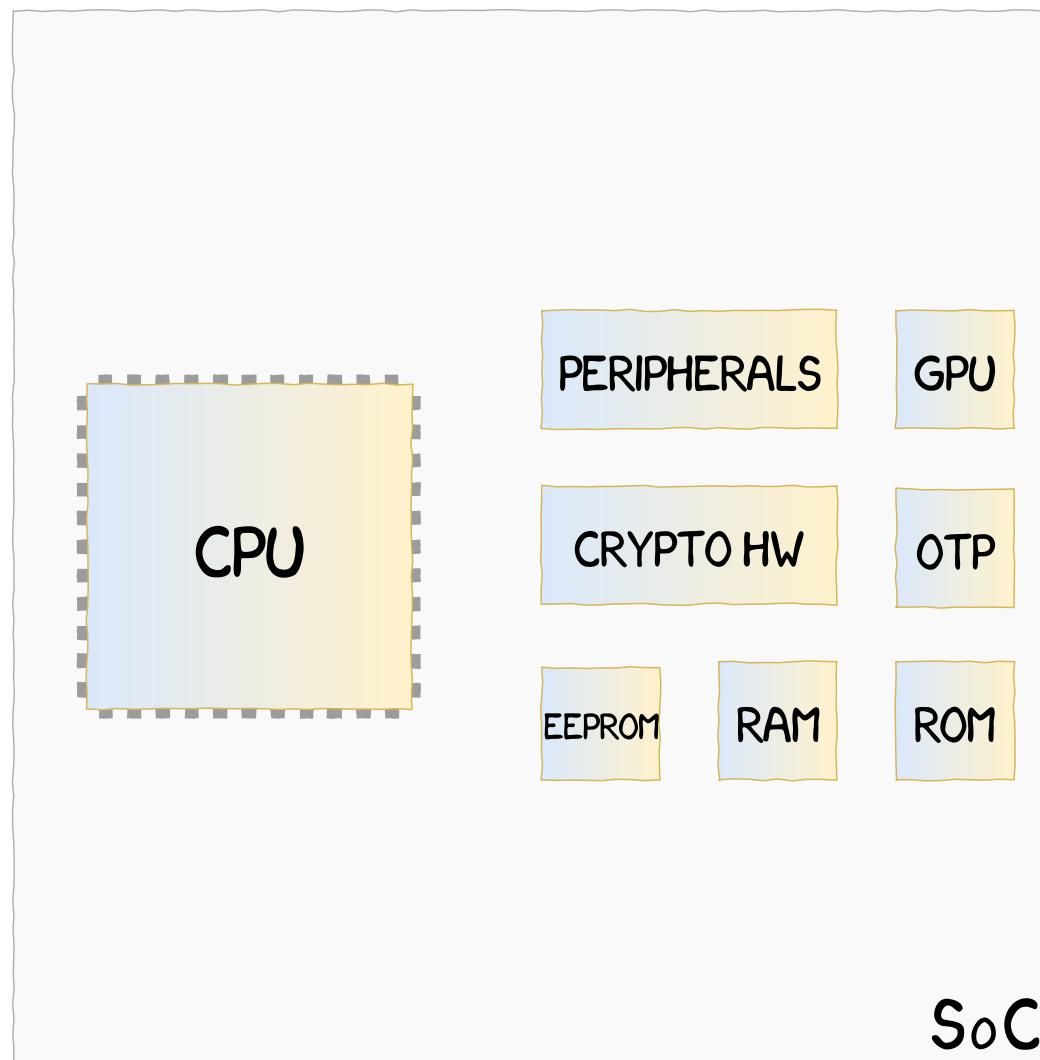
External Coprocessor (e.g. Google Titan M)



ARM TRUSTZONE TECHNOLOGY

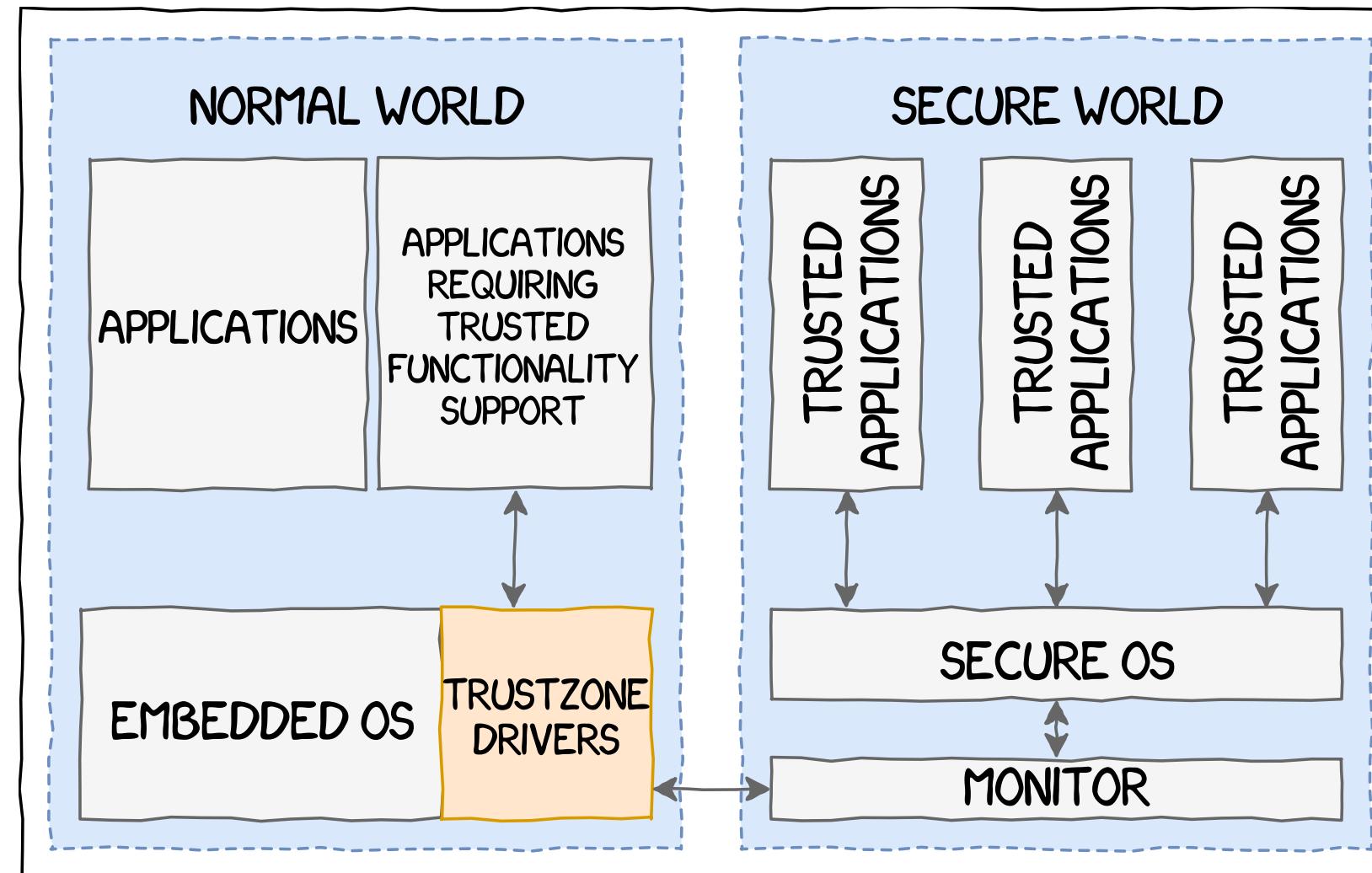
OVERVIEW

ARM TrustZone is a system-wide hardware isolation mechanism



- **Hardware architecture**
 - Partitioning of all the SoC's hardware and software resources
 - TZPC, TZASC, TZMA, etc.
- **Software architecture**
 - Software implementation used in secure state
 - Communications between secure and non-secure components

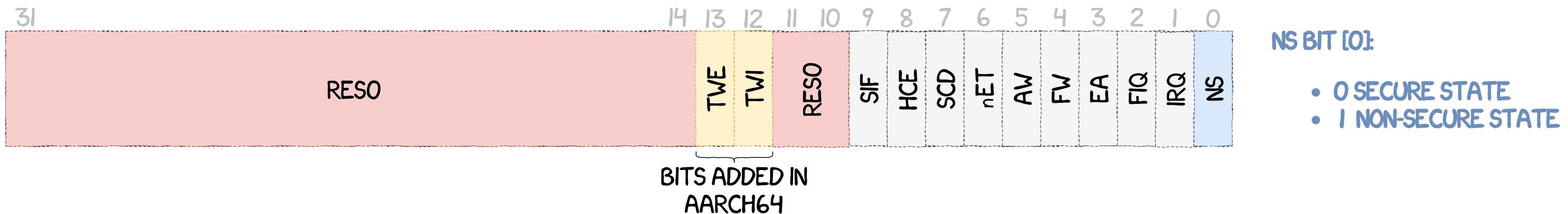
SECURE AND NON-SECURE WORLDS



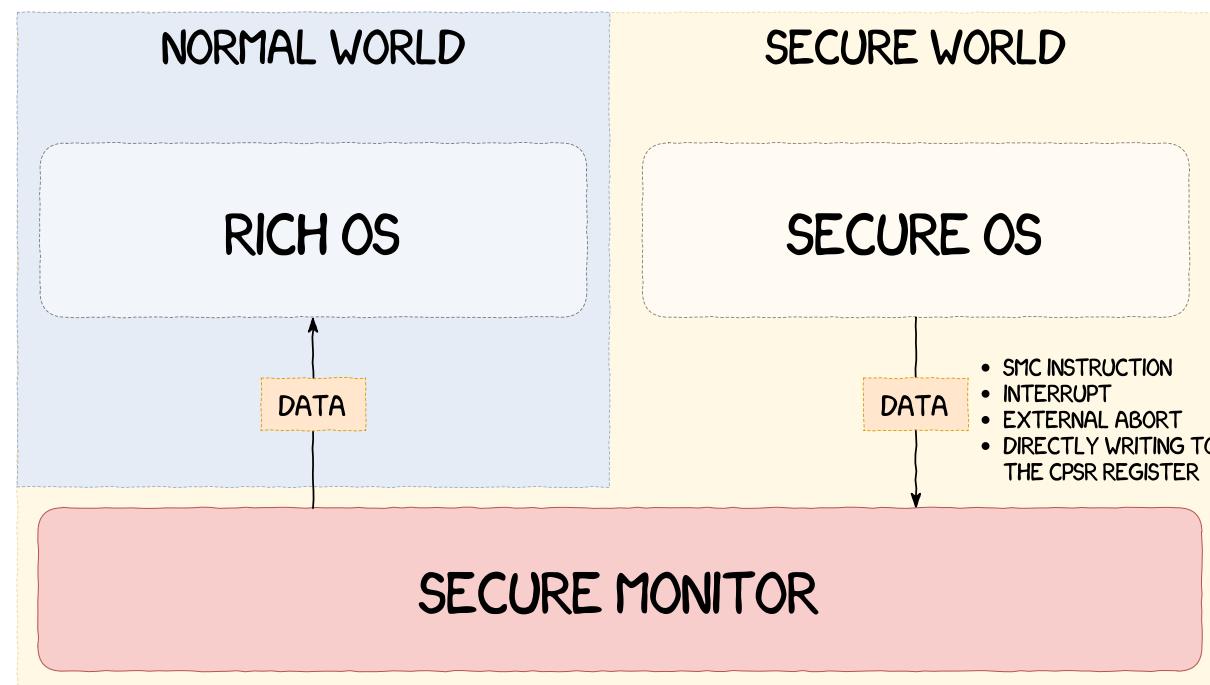
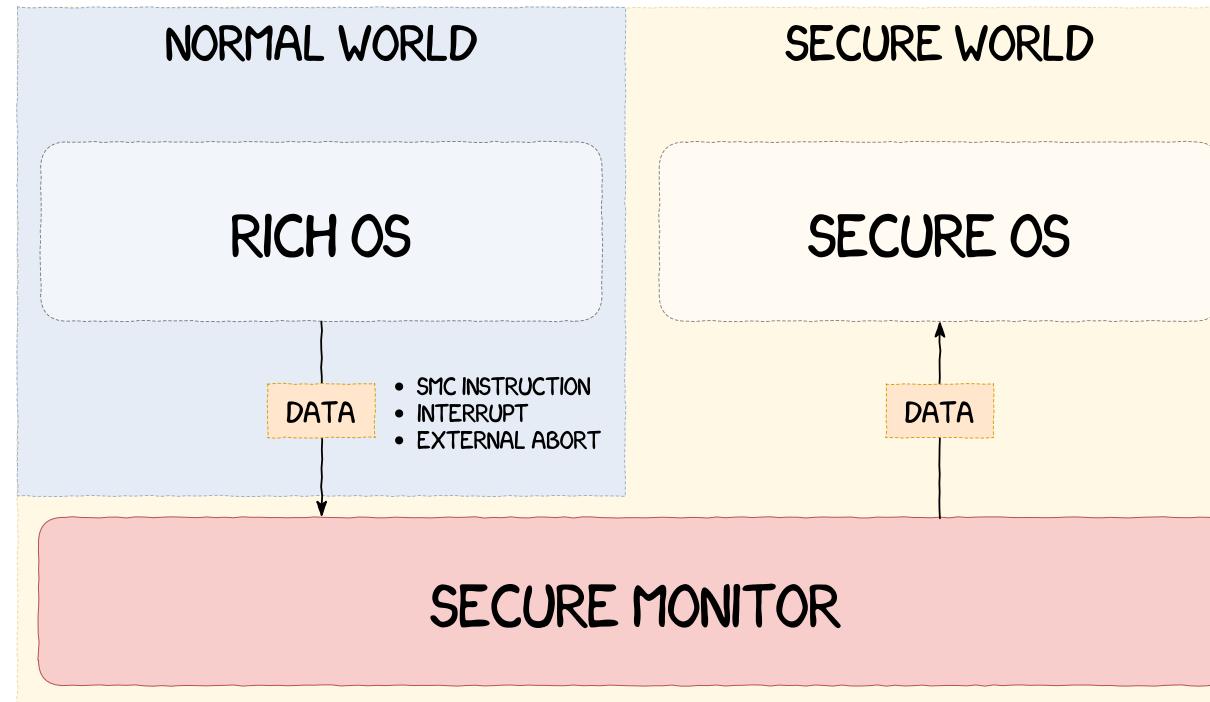
- **Secure World**
 - Runs trusted code
 - Performs sensitive operations
- **Normal World**
 - Considered as compromised by design
 - Performs non-sensitive operations

SECURE CONFIGURATION REGISTER

The **Secure** (or **Non-Secure**) state of the CPU is determined by the least significant bit of the **Secure Configuration Register** (SCR)



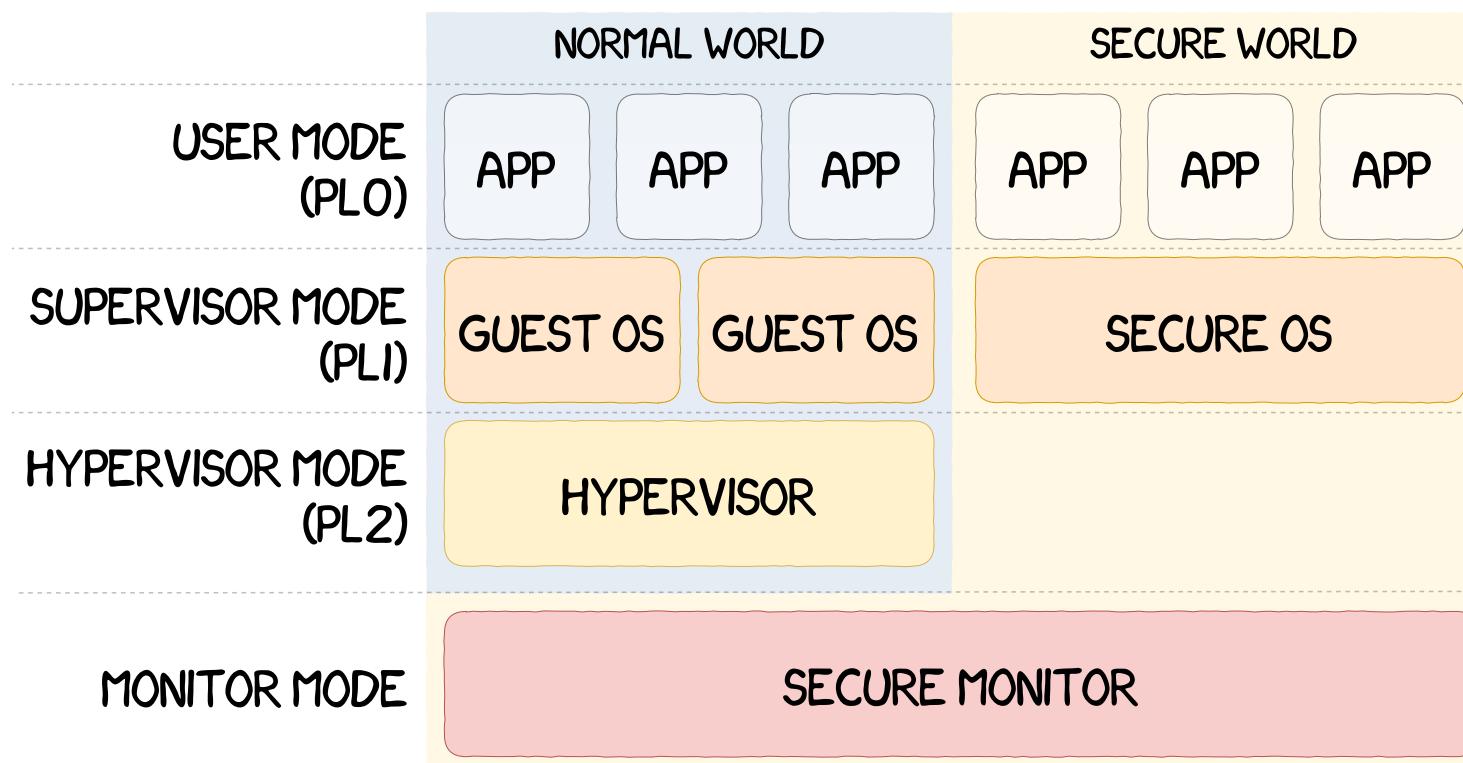
COMMUNICATING BETWEEN WORLDS



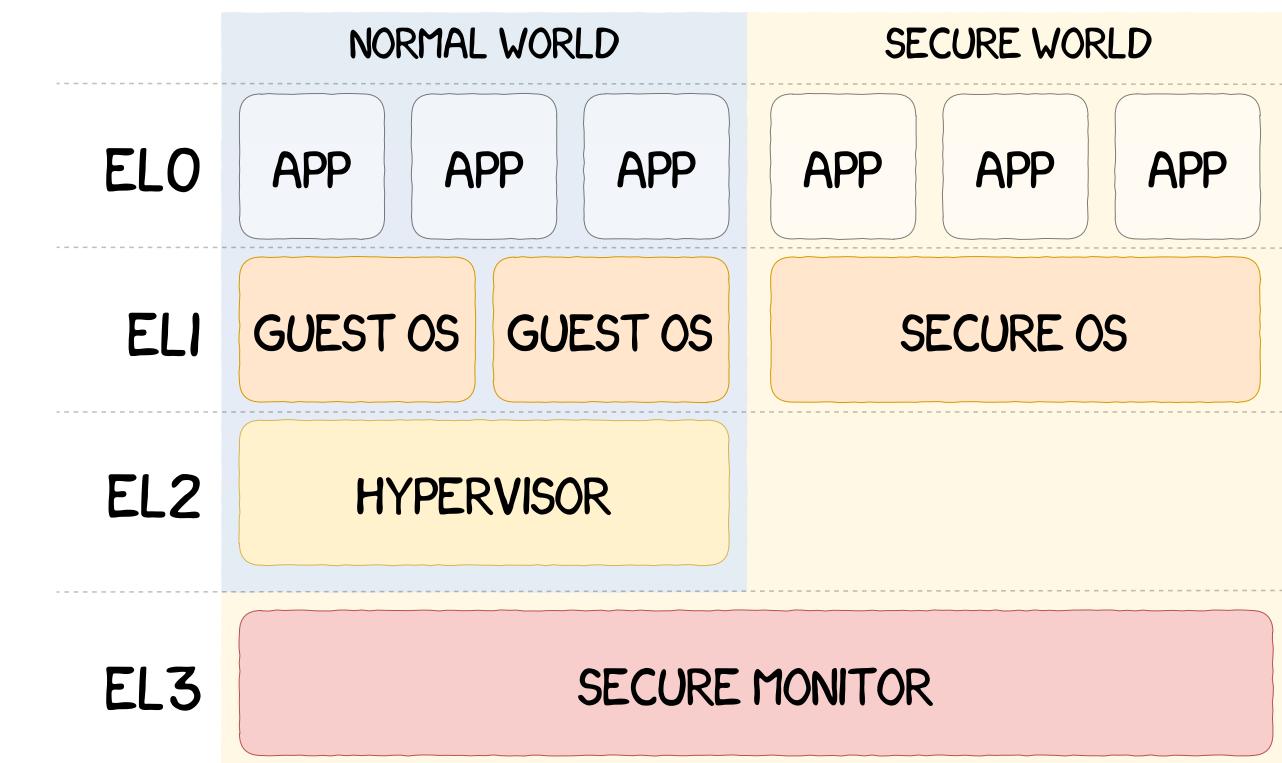
- Switches between worlds are performed by the **Secure Monitor**
 - Runs at the highest privilege level (**EL3** in ARMv8/**Monitor Mode** in ARMv7)
- Data exchanged through
 - Exceptions
 - Interruptions
 - Writing to the PSTATE/CPSR registers (privileged operation)

PRIVILEGES SEPARATION

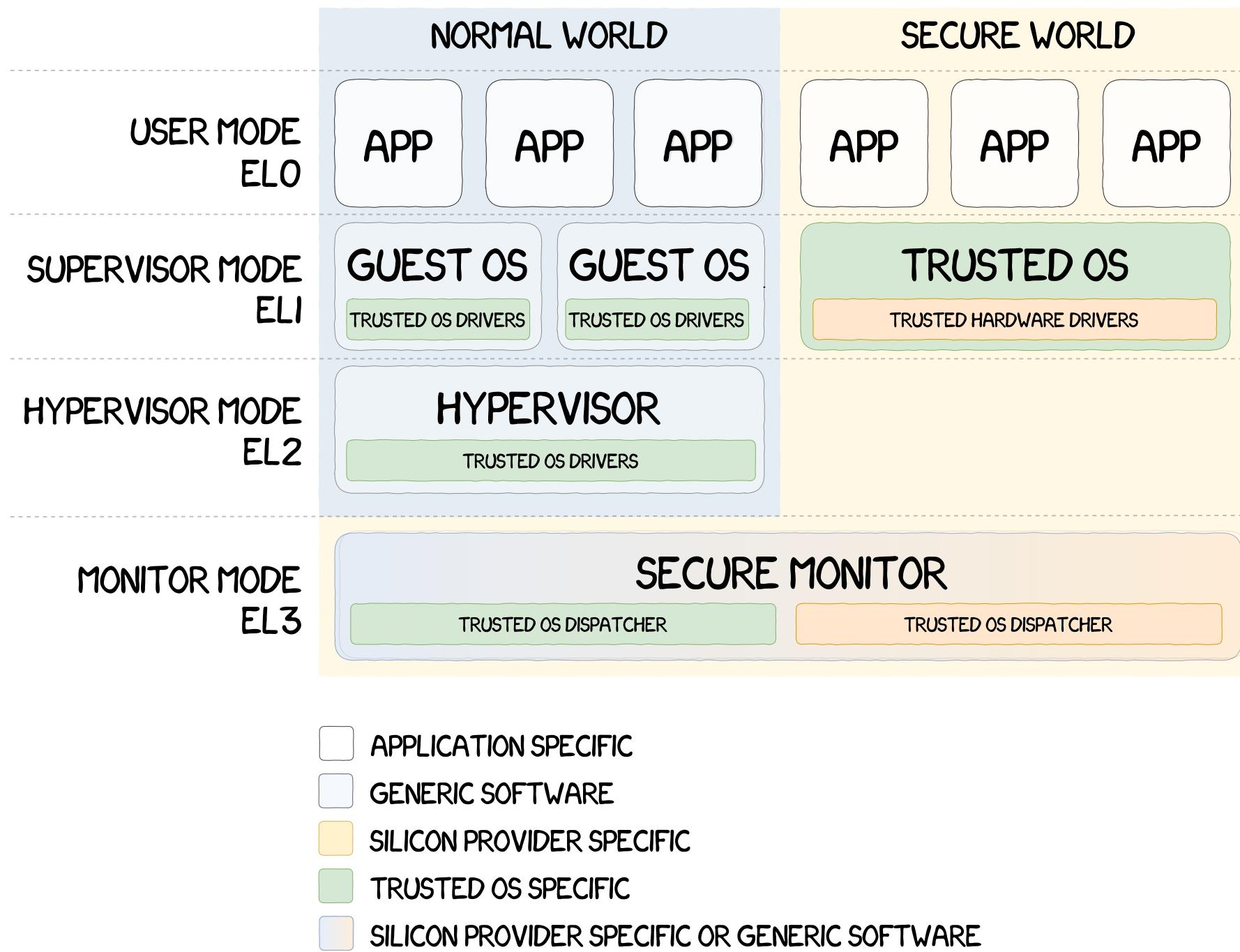
ARMv7 Privilege Levels



ARMv8 Exception Levels



THE TRUSTED COMPONENTS FRAGMENTATION ISSUE

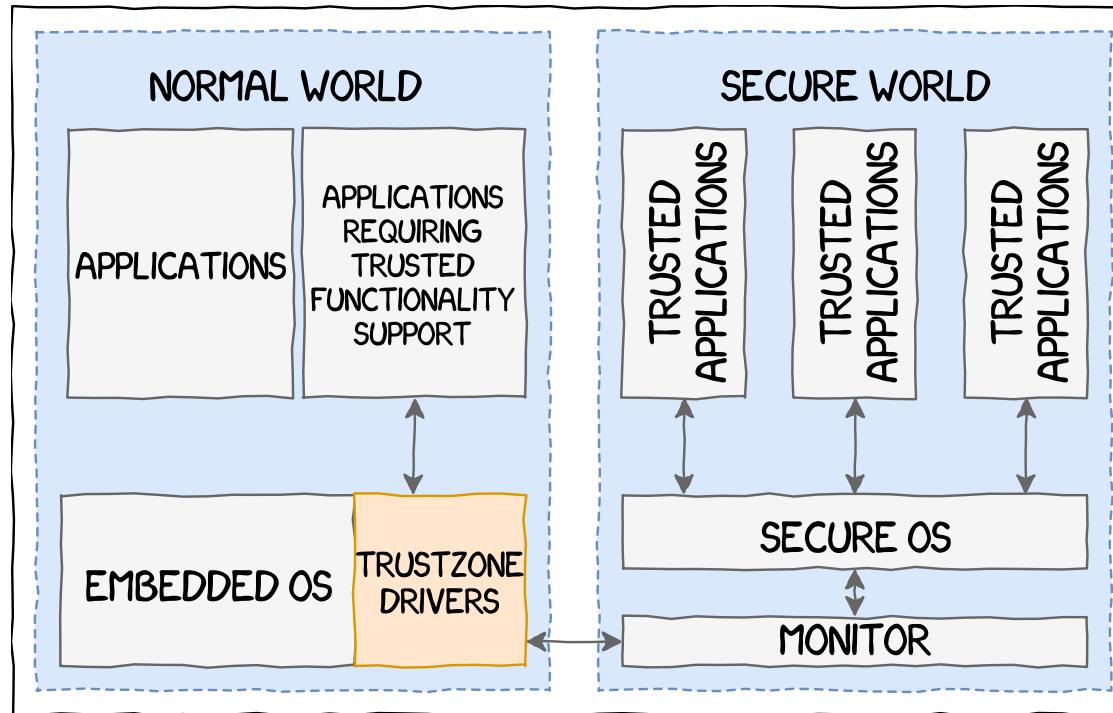


- **Privilege escalation by design**
 - No hardware isolation between S-EL1 and EL3
 - Access to all the physical memory
 - Will be fixed in ARMv8.4 with *Secure Partitions*
- **Fragmentation**
 - System developed by different vendors
 - Cooperation and mutual trust required

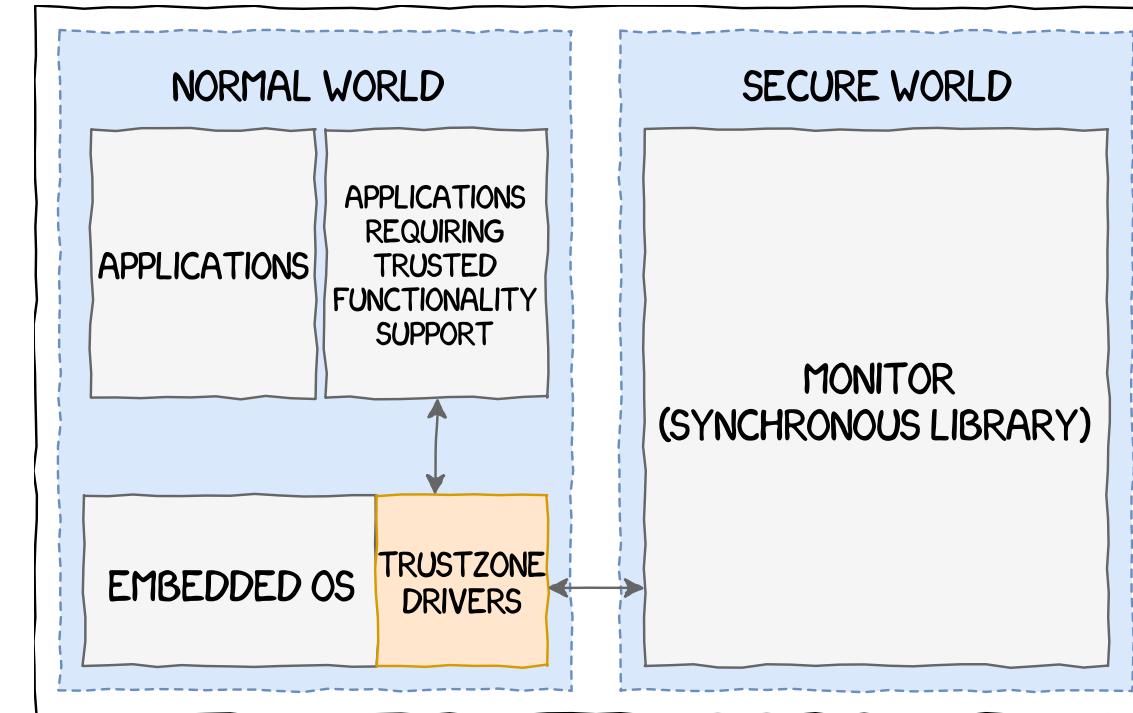
TRUSTZONE SOFTWARE ARCHITECTURE

Several implementations of the software stack running in TrustZone are possible

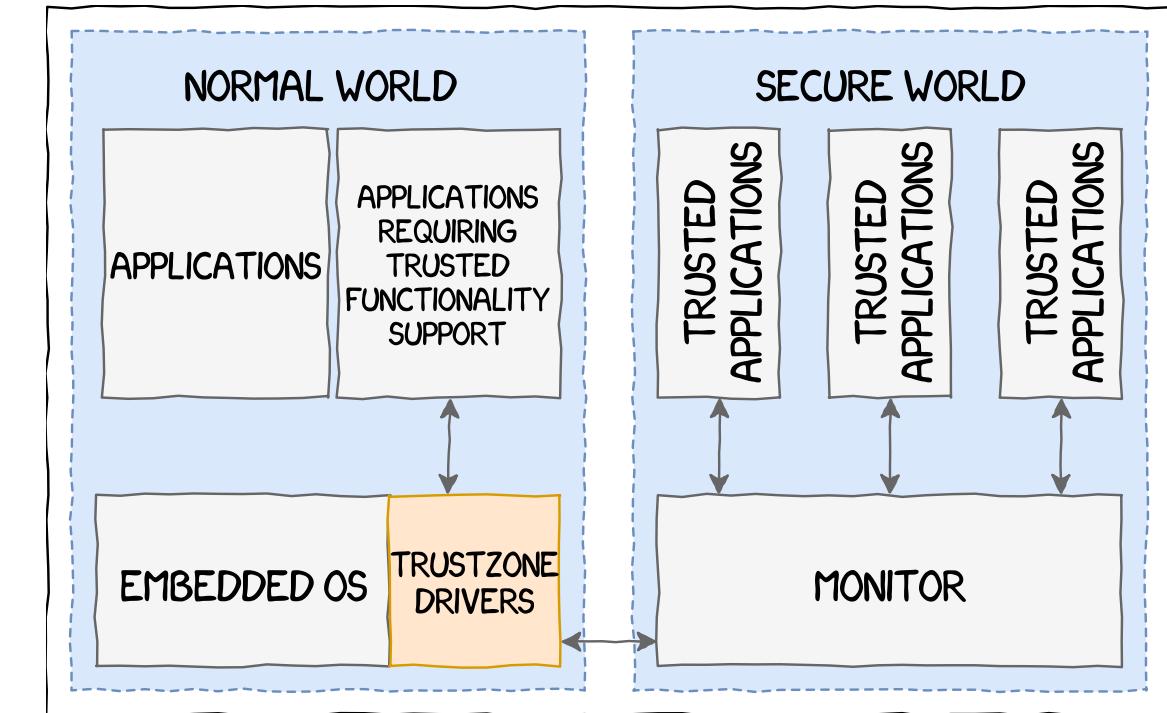
Operating System



Synchronous Library



Intermediate Option



TRUSTZONE'S USE CASES

- Accessing hardware-backed features:
 - Cryptographic engine
 - Credentials storage (Hardware-backed Keystore)
 - True random number generator
 - ...
- Digital Rights Management (by leveraging the cryptographic engine)
- Protecting and monitoring of the Normal World by the Secure World
 - **Example:** Samsung's Real-Time Kernel Protection (RKP) and Periodic Kernel Measurement (PKM)

SAMSUNG'S ARM TRUSTZONE

OVERVIEW

- **Samsung Devices**

- Use both Samsung's Exynos and Qualcomm's Snapdragon SoCs
 - The same phone models can have different SoCs depending on the country

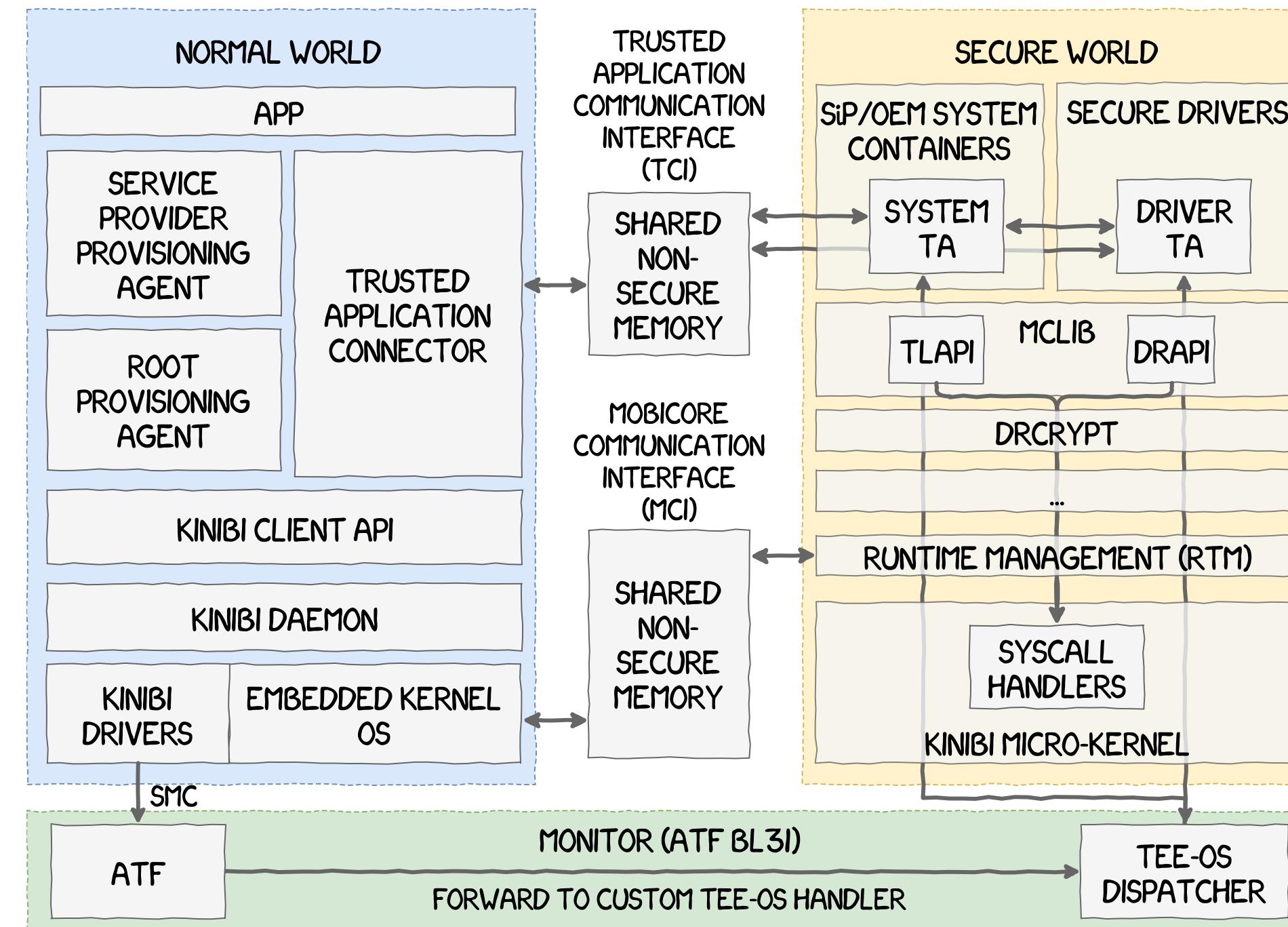
- **Samsung's TrustZone**

- Found only on Exynos SoCs
- First used in the Samsung Galaxy S3
- Trusted OS used:
 - **Kinibi** developed by Trustonic (Galaxy S3 to Galaxy S9)
 - **TEEGRIS** developed by Samsung (Galaxy S10)
 - Both are used in other models too
 - This talk will focus on **Kinibi**

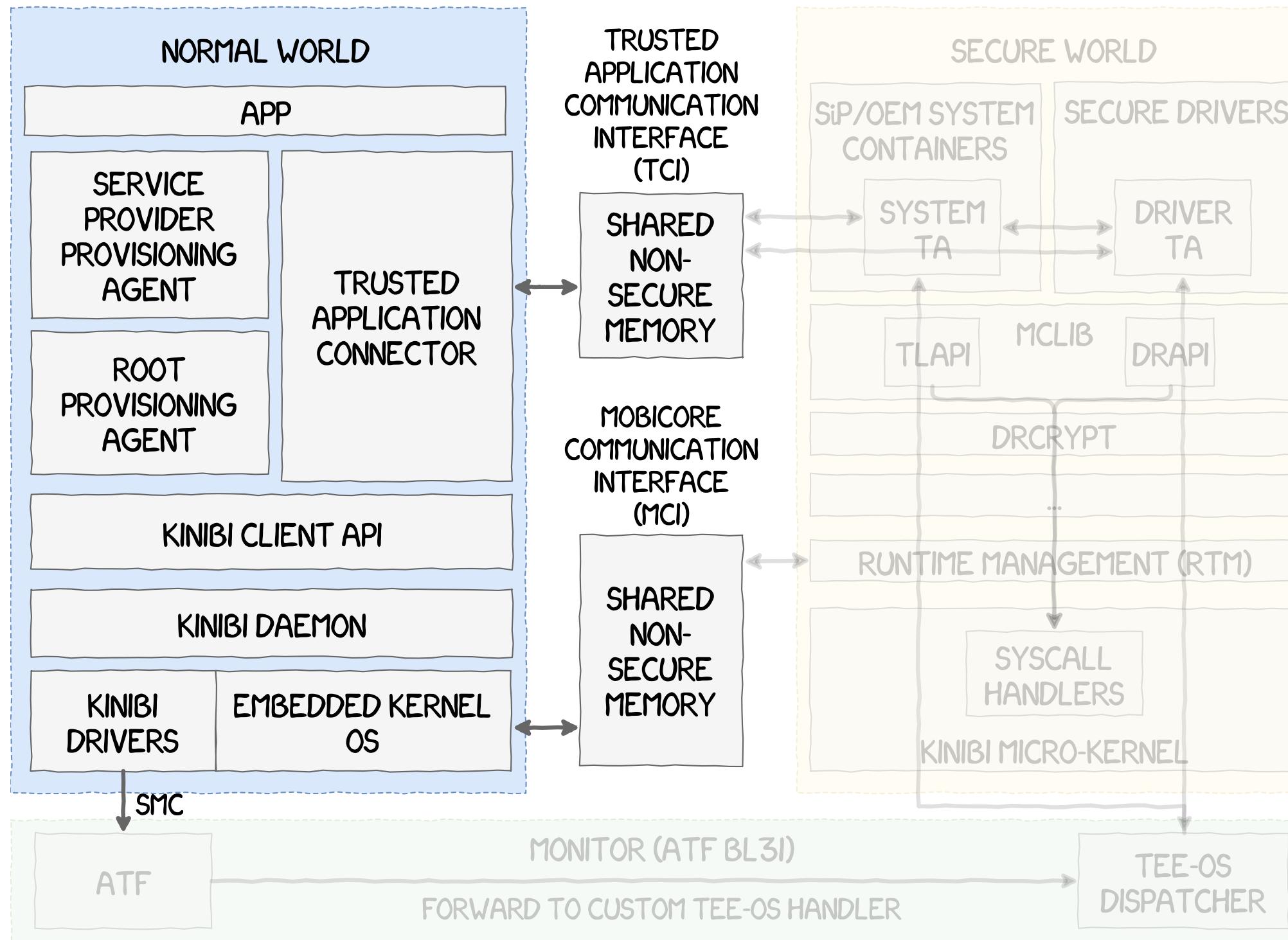
PREVIOUS WORKS

- Reverse Engineering Samsung S6 SBOOT (2-part article series) by **Fernand Lone Sang**
 - ARM Trusted Firmware usage on Samsung devices and extraction process from an OTA of the TEE-OS
- Unbox Your Phone (3-part article series) by **Daniel Komaromy**
 - Reverse-engineering of the Trusted OS and exploitation of vulnerabilities in trustlets
- Trust Issues: Exploiting TrustZone TEEs by **Gal Beniamini**
 - Security analysis of different Trusted Execution Environments

SAMSUNG'S TRUSTZONE ARCHITECTURE

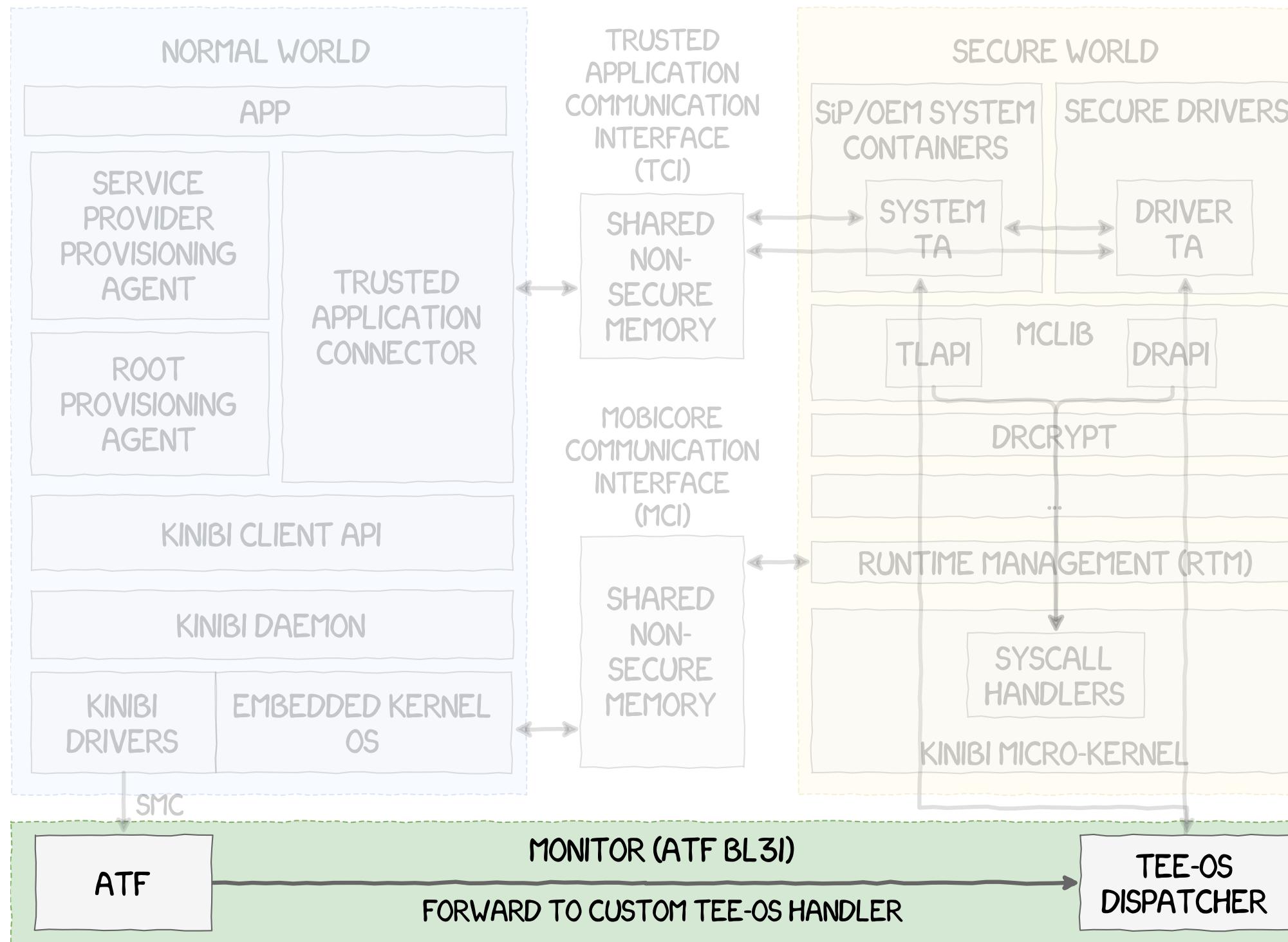


NORMAL WORLD COMPONENTS



- Drivers, daemons, libraries and interfaces used for communicating with the Secure World
- Communications pass through SMCs and shared memory buffers

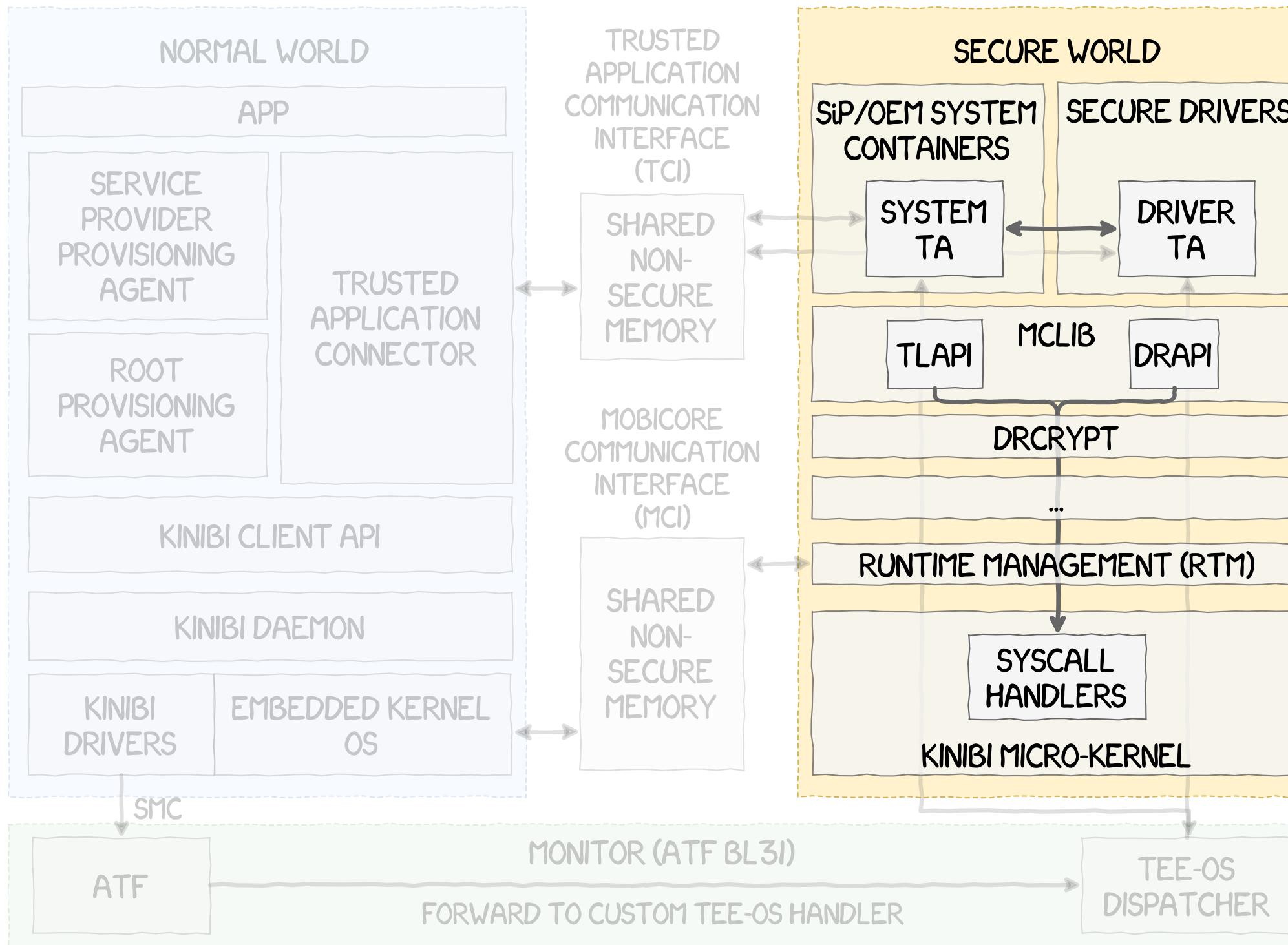
SECURE MONITOR



- **ARM Trusted Firmware**

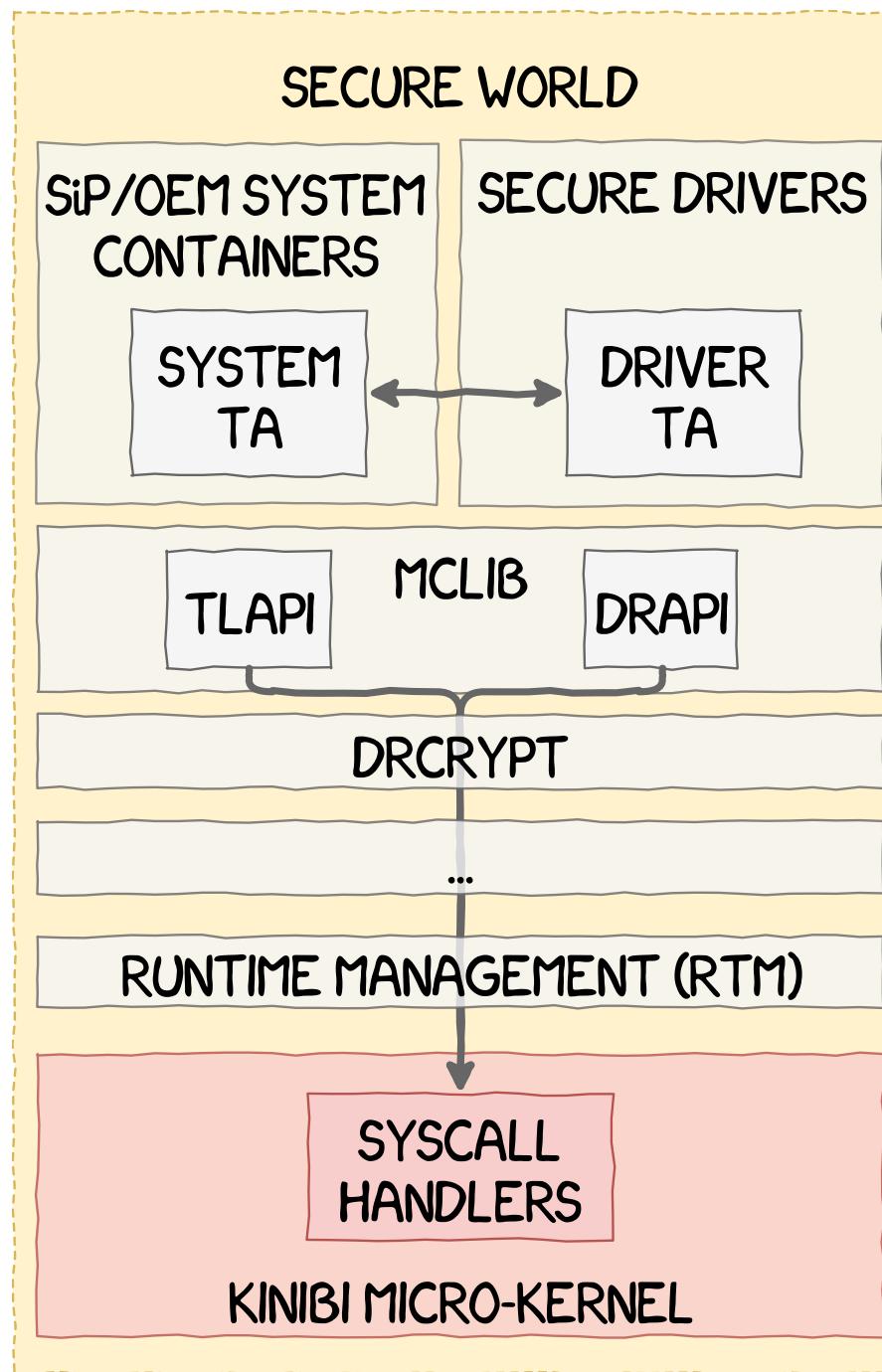
- Open-source reference implementation of Secure World software provided by ARM
- Contains a modular secure monitor implementation
- Custom SMC handlers, called **runtime services**, can be added to fit the vendors requirements
- **Example:** runtime services are used by Samsung to forward SMCs handled by Kinibi

SECURE WORLD COMPONENTS



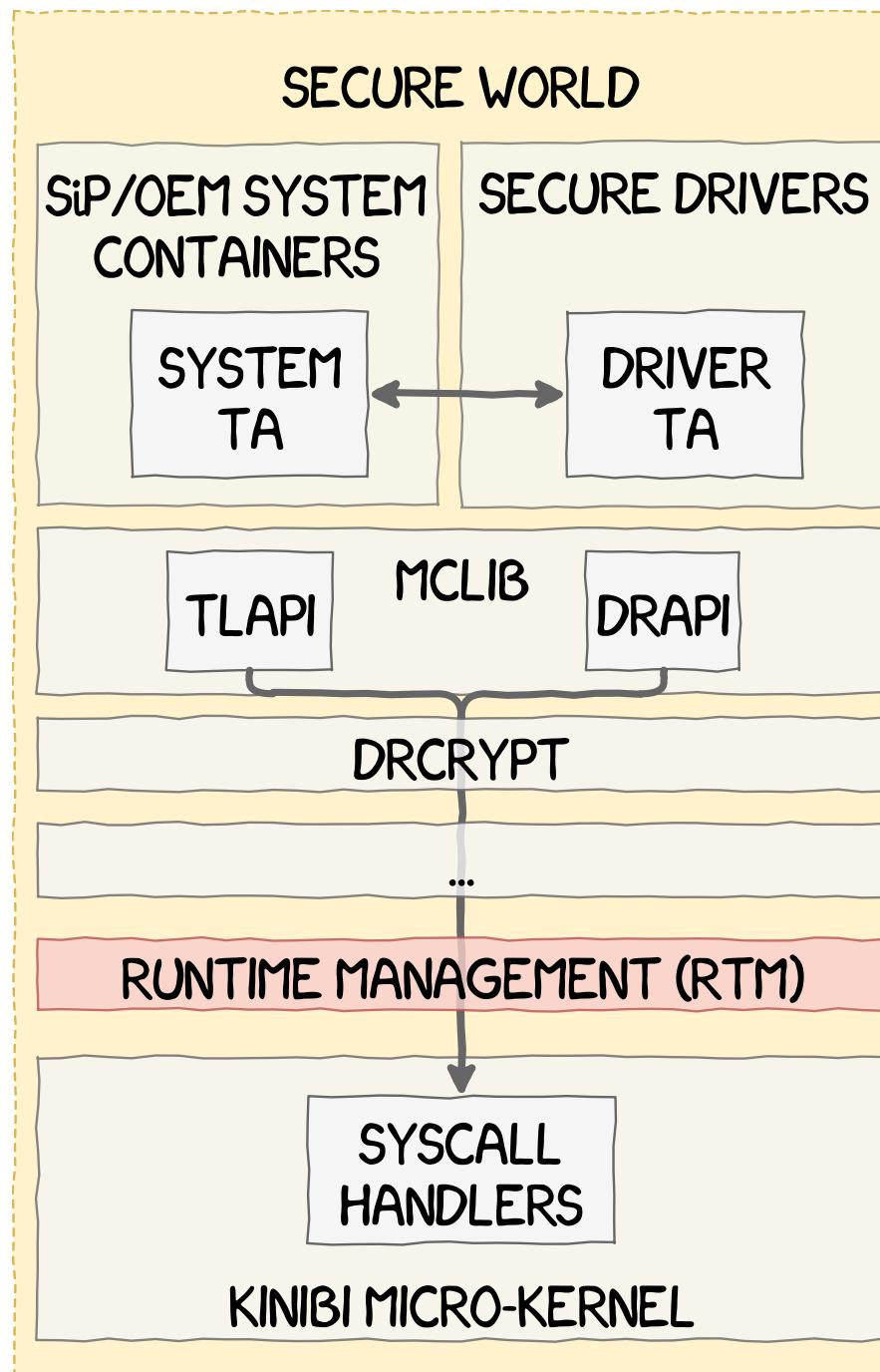
- Secure world based on a micro-kernel architecture

MTK: KINIBI'S MICRO KERNEL



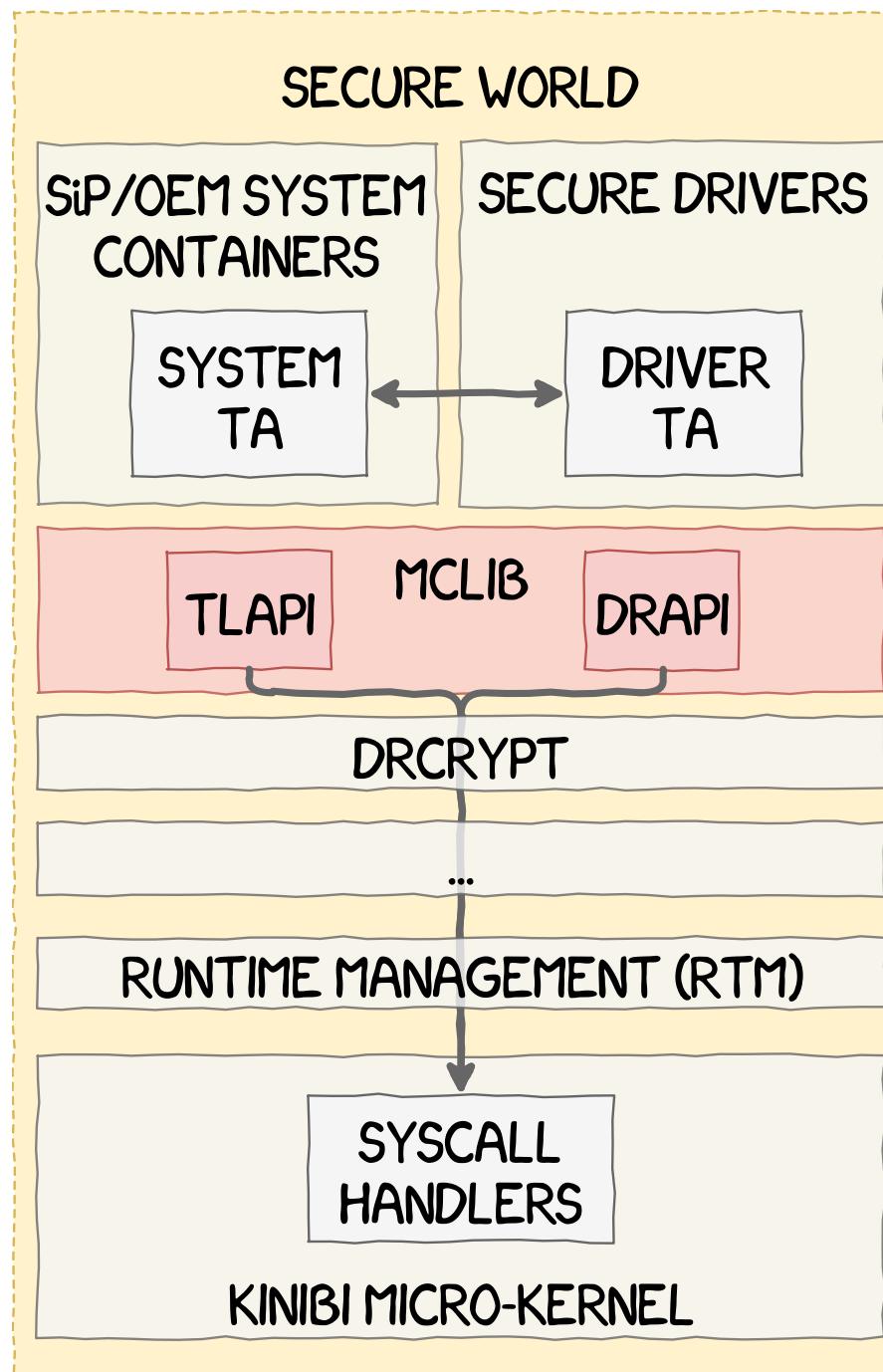
- Kinibi is a 32-bit OS developed by **Trustonic**
 - Used to be called *Mobicore* and *t-base*
- **MTK**: micro-kernel and only component running in S-EL1
- Provides syscalls (SVCs)
 - Memory mapping, process creation, SMCs, etc.
 - SVCs available depend on the privileges of the calling process
- Loads other components (embedded drivers, etc.) and especially **RTM**

RUN-TIME MANAGER



- Special Secure World trusted application equivalent to the init process on Linux
- **Main tasks**
 - starting and managing processes
 - notifying trustlets of incoming data from the NWd
- **Implements communication channels**
 - Inter-Process Communications
 - Mobicore Communication Interface (MCI)
 - A communication channel with the Normal World based on the Mobicore Control Protocol (MCP)

MCLIB: KINIBI'S STANDARD LIBRARY

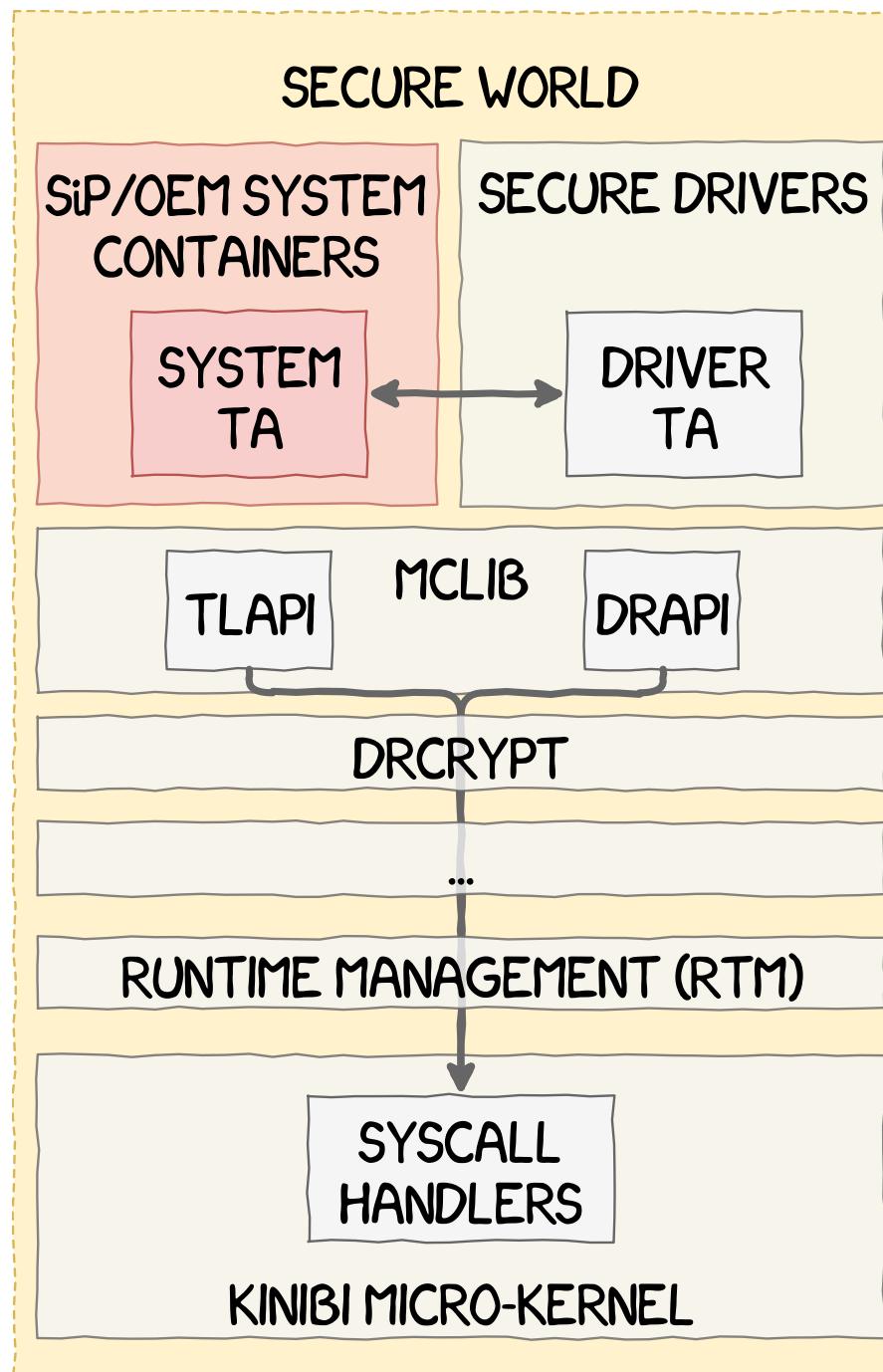


- Provides standard functions to Trusted Applications, Secure Drivers and RTM
- Separated into two APIs:
 - **TlApi:** set of functions used by trusted applications
 - **DraPI:** set of functions used by secure drivers
- Useful during exploitation to find gadgets
- **TlApi call example**

```

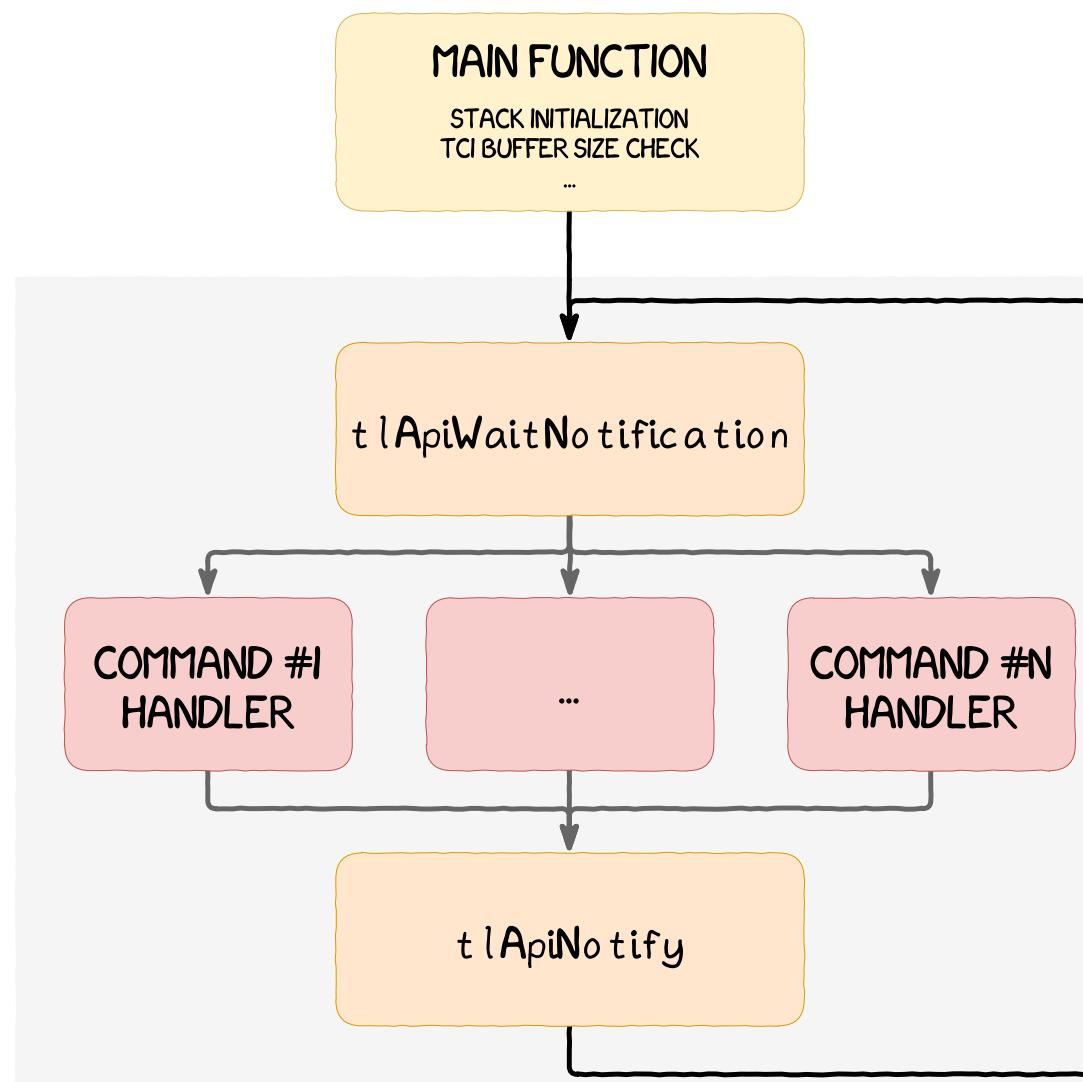
; _DWORD tlApiWaitNotification(_DWORD timeout)
MOV.W      R1, #0x1000
LDR.W      R2, [R1, #(tlApiLibEntry - 0x1000)]
MOV        R1, R0
MOVS     R0, #6
BX       R2
  
```

TRUSTED APPLICATIONS



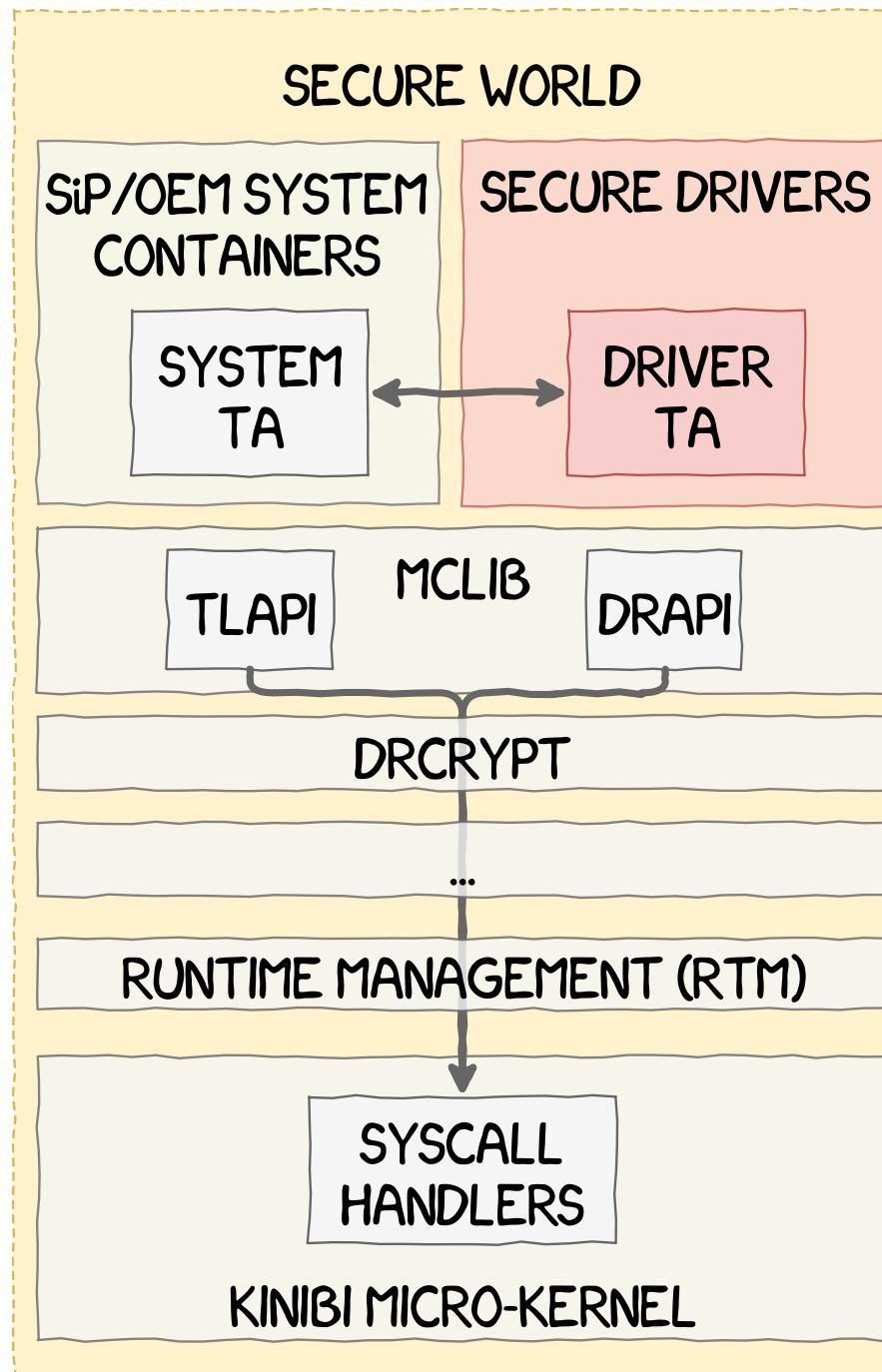
- Secure World equivalent of regular applications in the Normal World (run at S-EL0)
- Allow trusted third-parties to extend the functionalities of the TEE-OS
 - Trusted UI, DRM, storage of secrets, etc.
- Signed binaries loaded directly from the Normal World (so are SDs)

TRUSTED APPLICATIONS LIFE-CYCLE



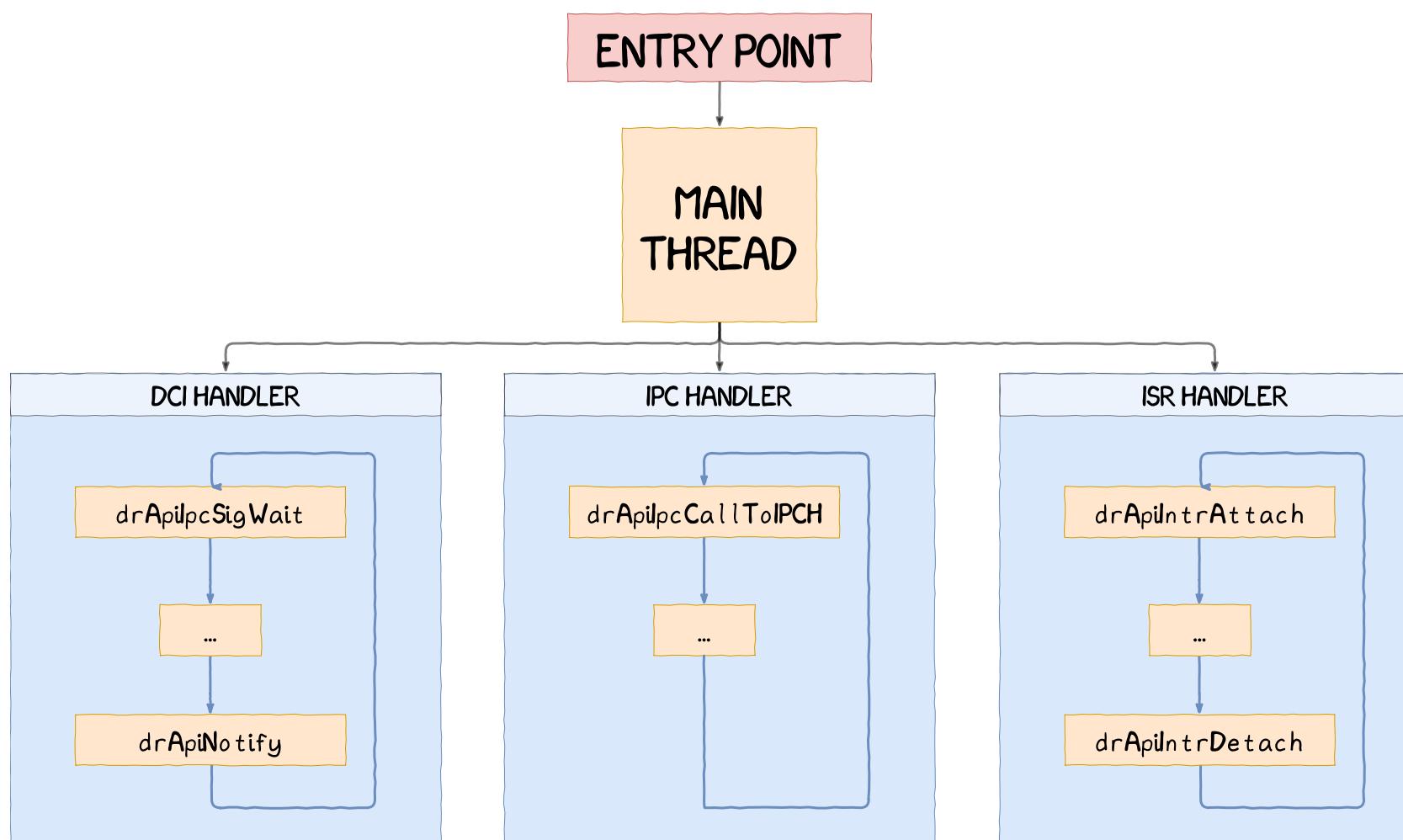
- Communications with the Normal World made through world-shared memory (named TCI buffer by Trustonic)
 - The TCI buffer contains commands to be handled by the trustlet
 - TCI buffer contains commands to be handled by the trustlet
-
- **Notifications**
 - **tlApiWaitNotification**
 - **tlApiNotify**

SECURE DRIVERS



- Special type of Trusted Applications
- Run at S-EL0 but have higher software-defined privileges
- Have access to a richer set of API and syscalls
- Are used by trustlets as an interface to access physical memory and reach secure peripherals in a controlled manner
- Communications with TAs made through IPCs and shared memory

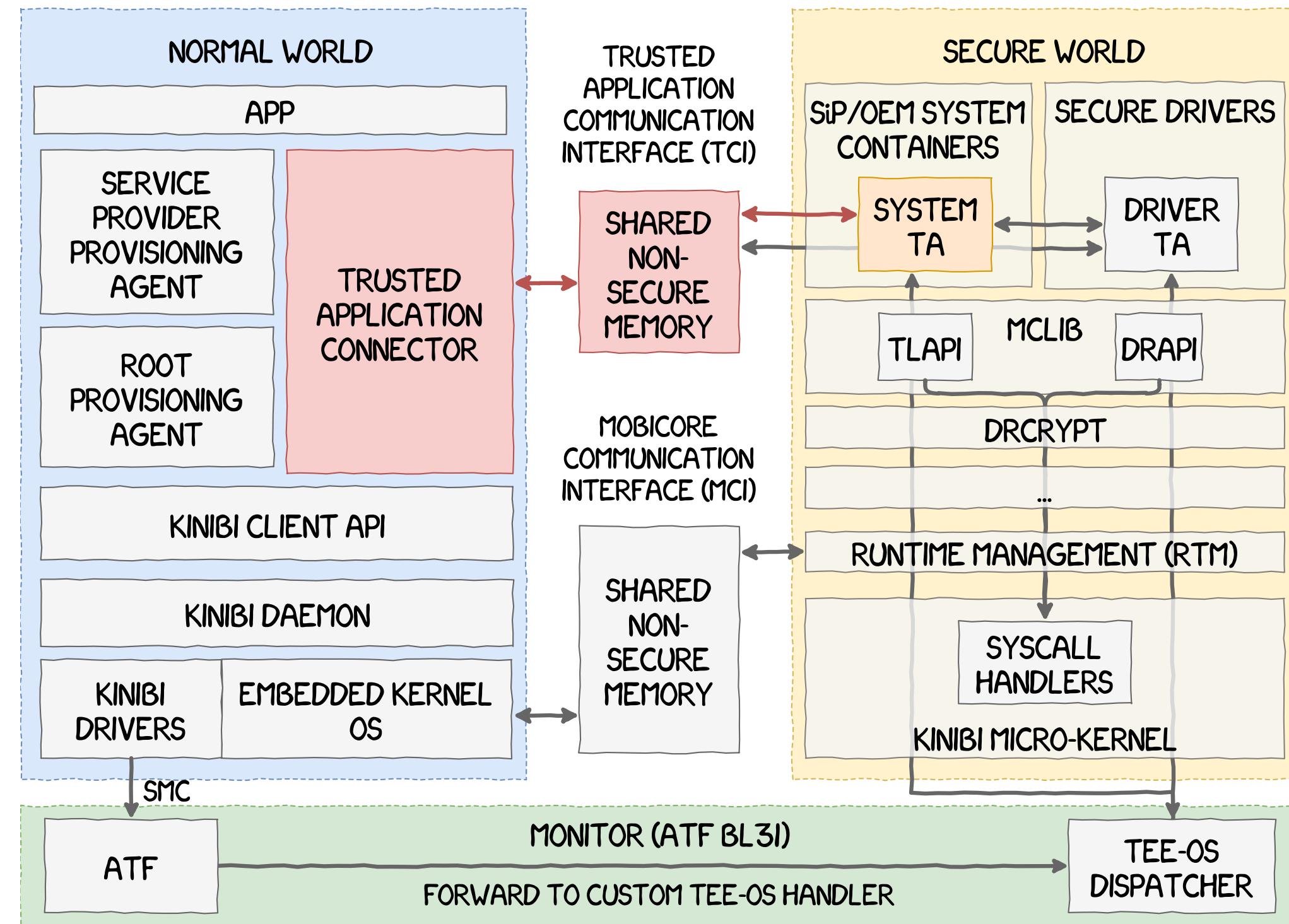
SECURE DRIVERS LIFE-CYCLE



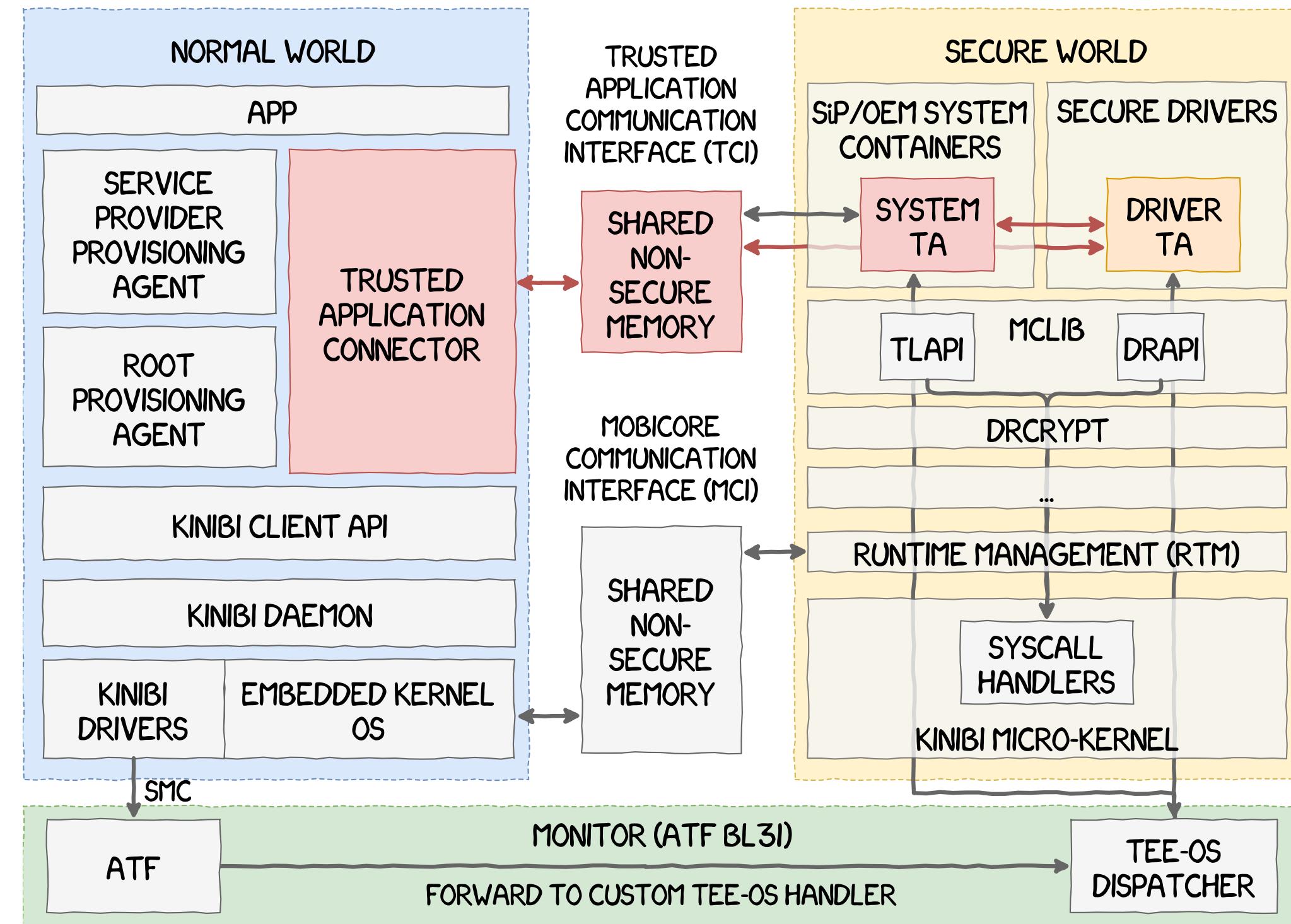
- **Multi-threaded application**
 - **DCI:** Normal World communications
 - **IPC:** trustlet communications
- **Trustlet interactions**
 - Retrieves IPC data by mapping the entire trustlet
 - Notifications using `drApiIpcCallToIPCH`

VULNERABILITY RESEARCH TOOLS

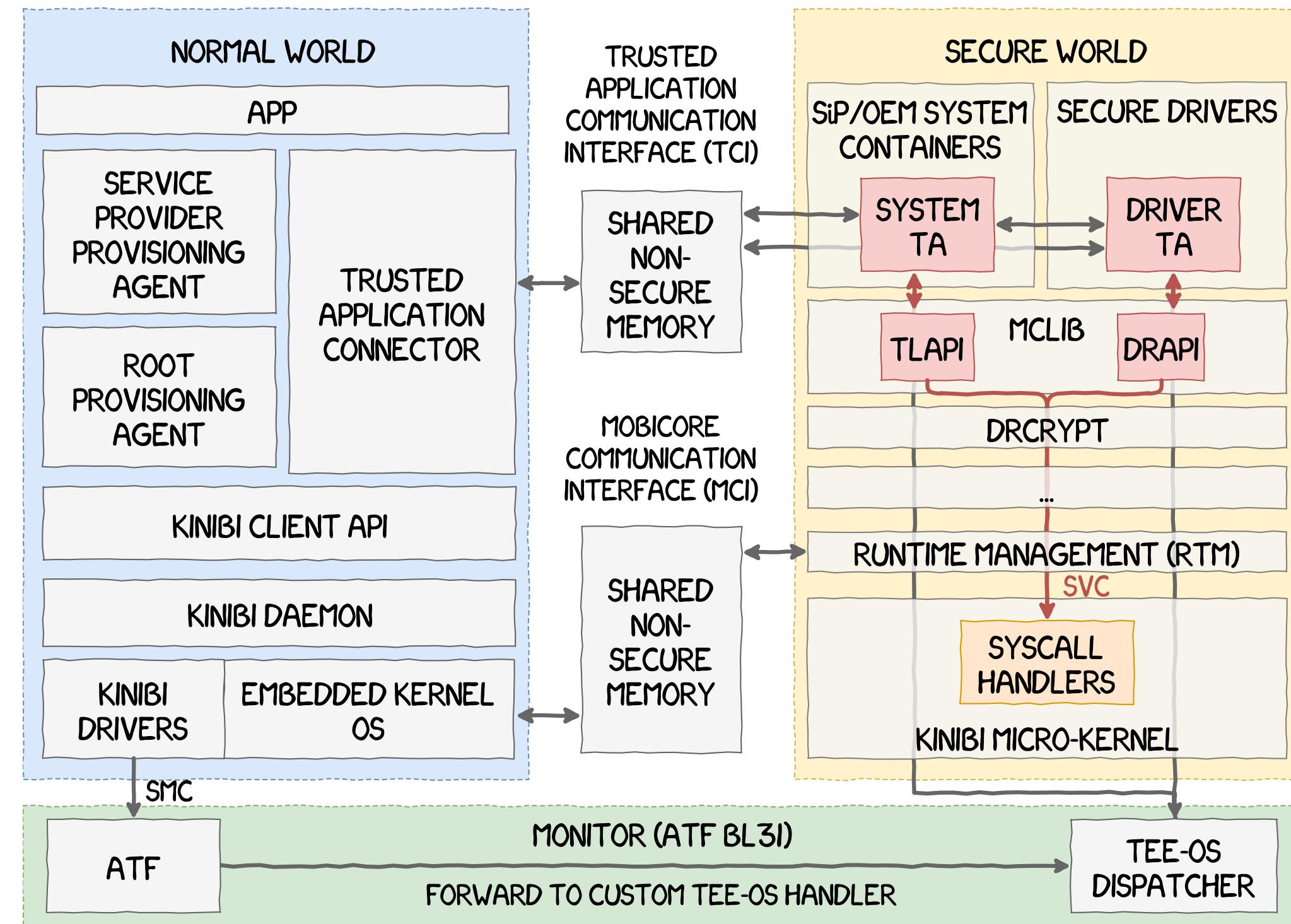
ATTACK SURFACE



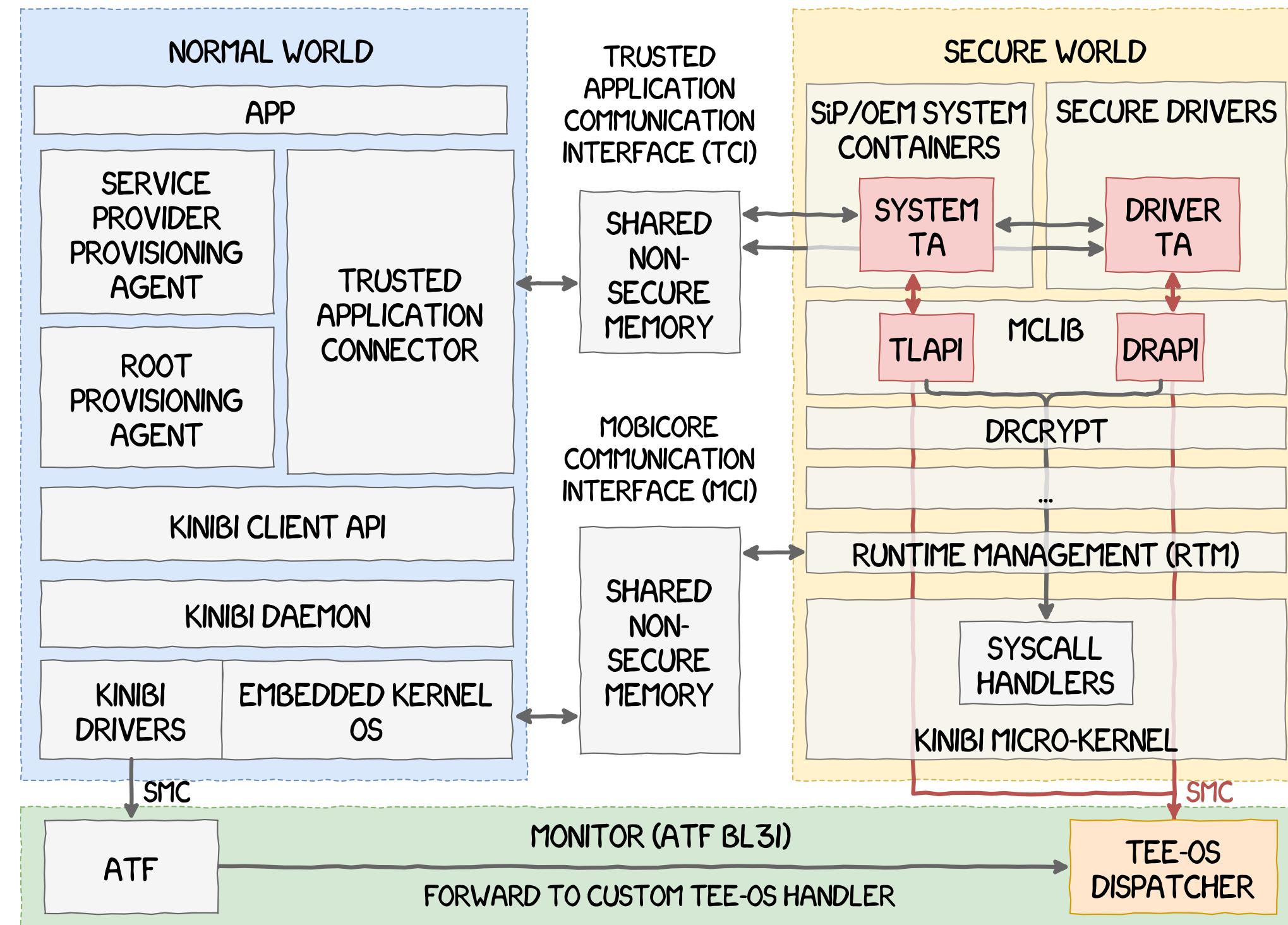
ATTACK SURFACE



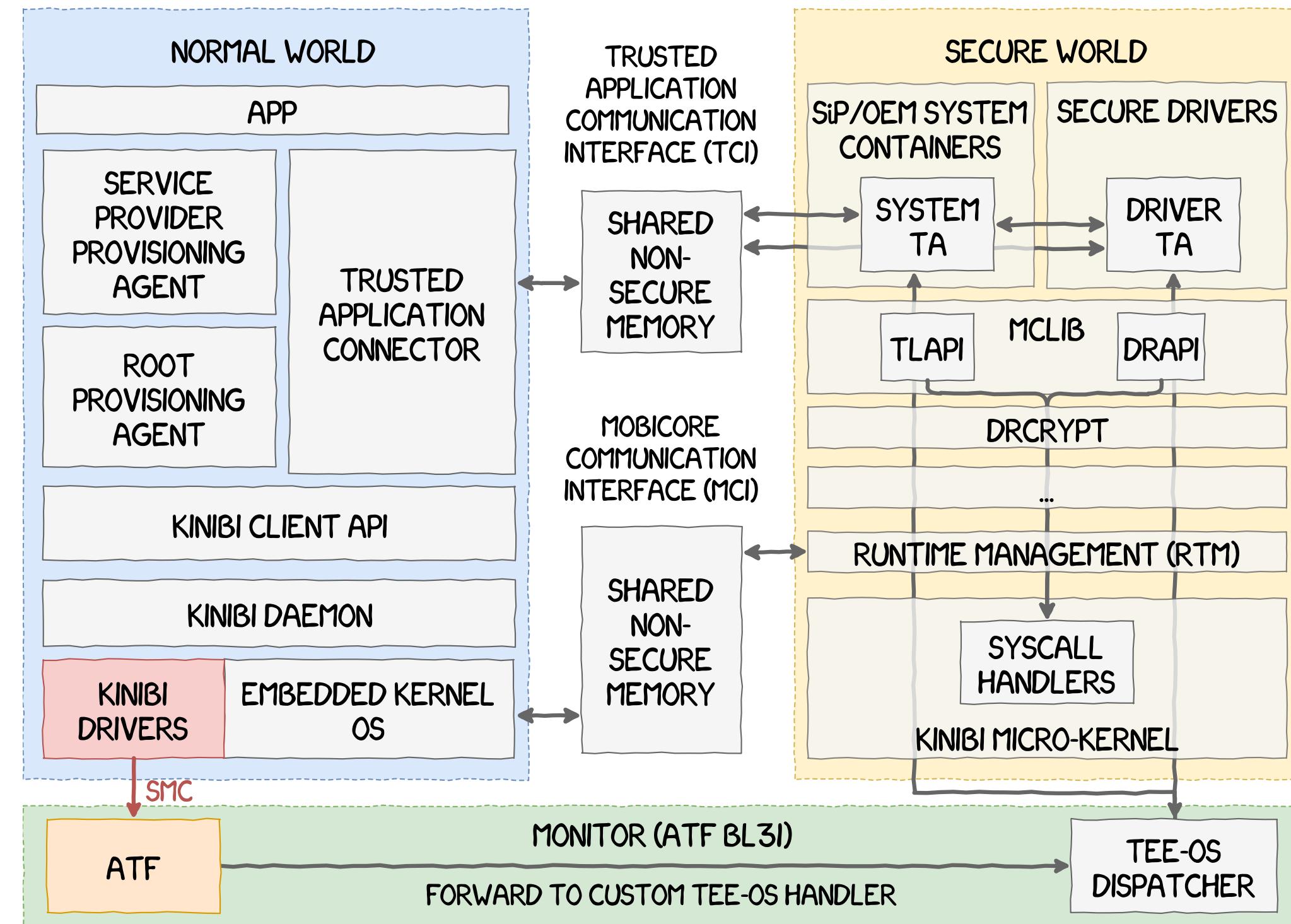
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ATTACK SURFACE



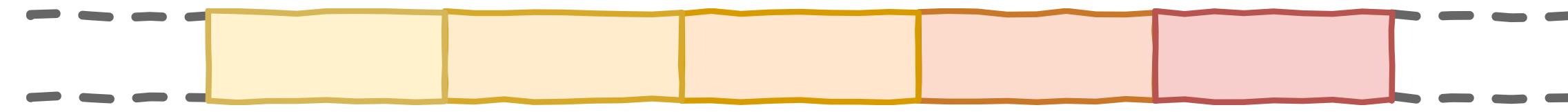
ATTACK SURFACE



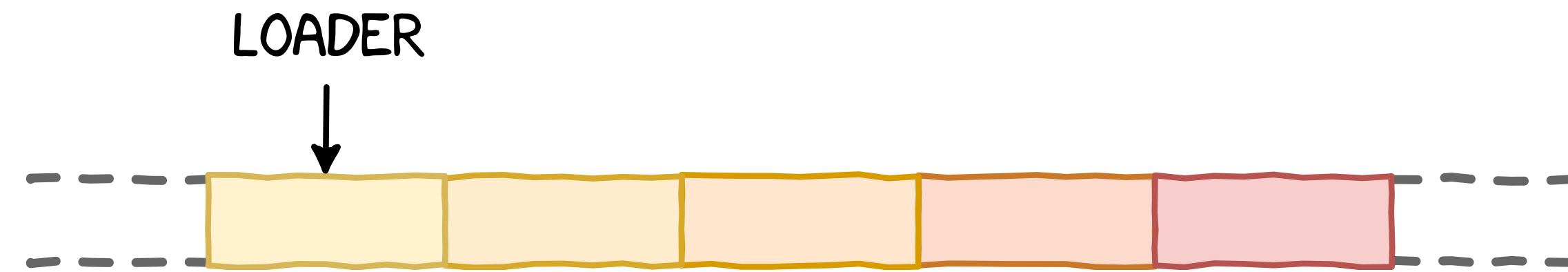
ATTACK SURFACE

- Must be reachable from the Normal World
- ATF is open-source, probably heavily reviewed
- Trusted Applications are low-hanging fruits

OUR JOURNEY IN 5 STEPS



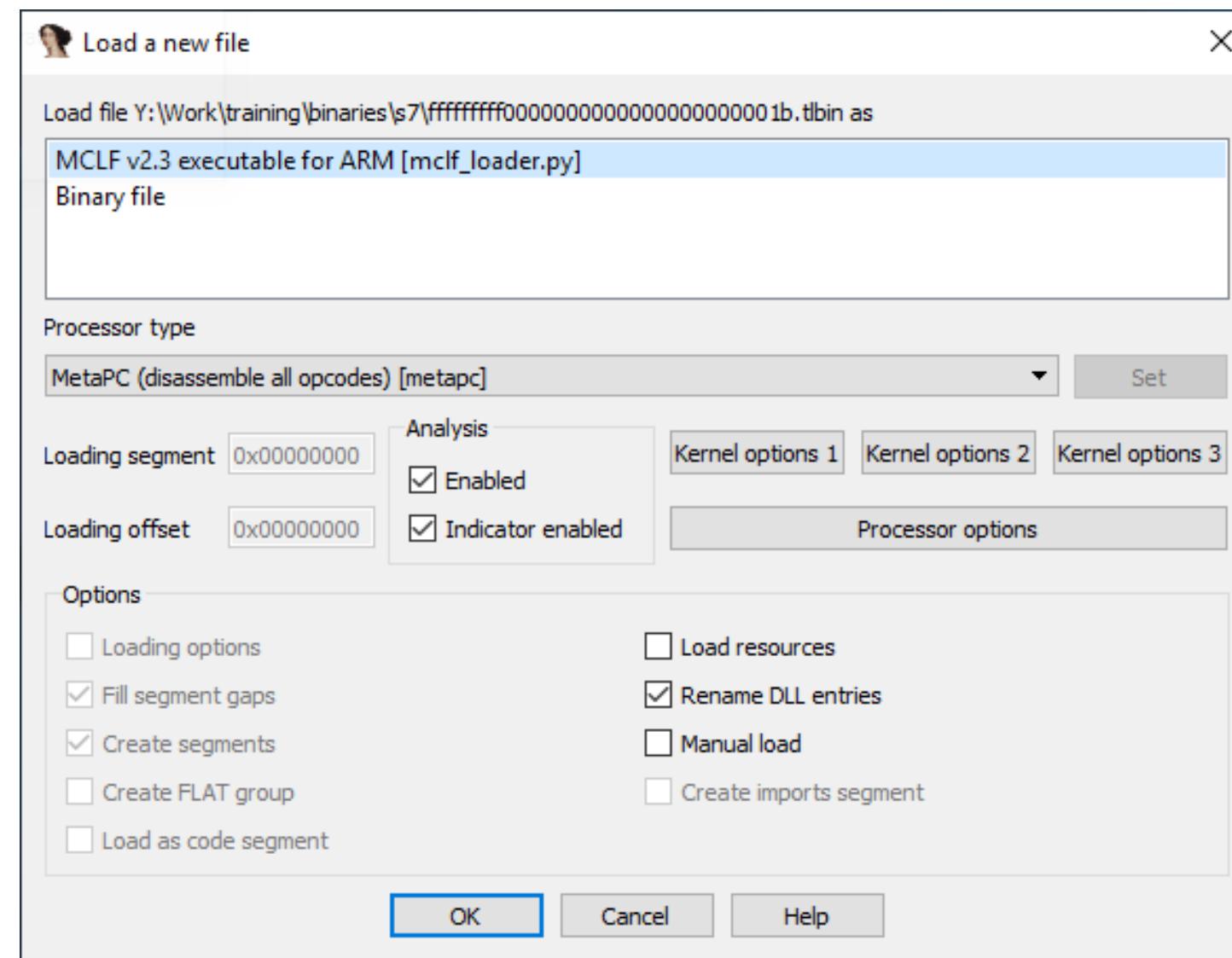
STEP #1 - LOADING INTO IDA/GHIDRA



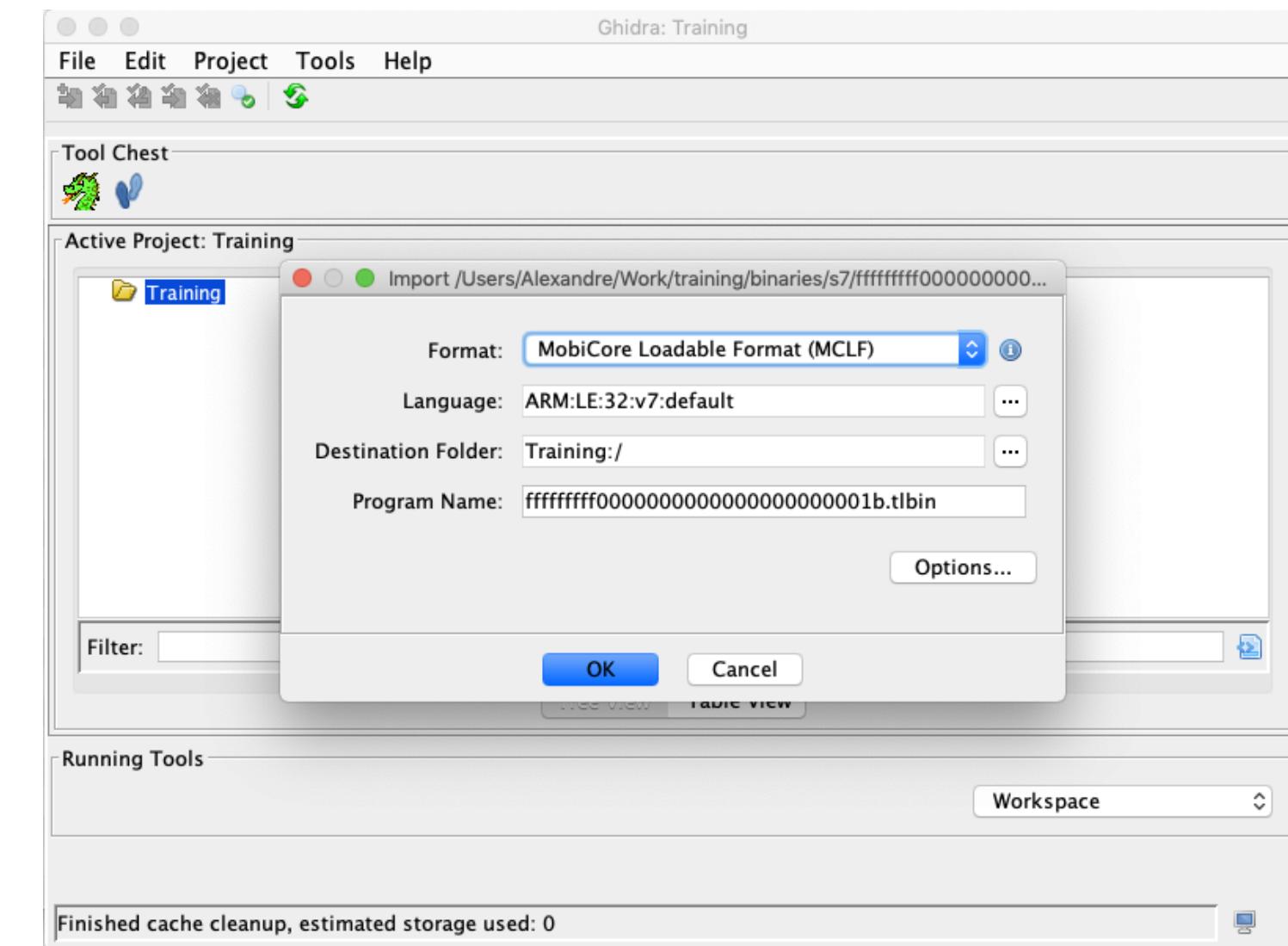
The screenshot shows a GitHub repository page for `Trustonic / trustonic-tee-user-space`. The repository has 5 stars and 5 forks. The current branch is `master`, and the file being viewed is `mcLoadFormat.h` located at `trustonic-tee-user-space / common / MobiCore / inc`. The commit history shows a single commit by `t-user` with the message `NWd from branches/rel_t-sdk-r7:12090`, dated `on 17 Feb 2015`. The file contains 246 lines (209 sloc) and is 11 KB in size. The code itself is a copyright notice for TRUSTONIC LIMITED, version 2013-2015, with redistribution conditions.

```
1  /*
2  * Copyright (c) 2013-2015 TRUSTONIC LIMITED
3  * All rights reserved.
4  *
5  * Redistribution and use in source and binary forms, with or without
6  * modification, are permitted provided that the following conditions are met:
7  *
8  * 1. Redistributions of source code must retain the above copyright notice,
9  *    this list of conditions and the following disclaimer.
10 *
```

- Proprietary File Format - MobiCore Loadable Format (MCLF)

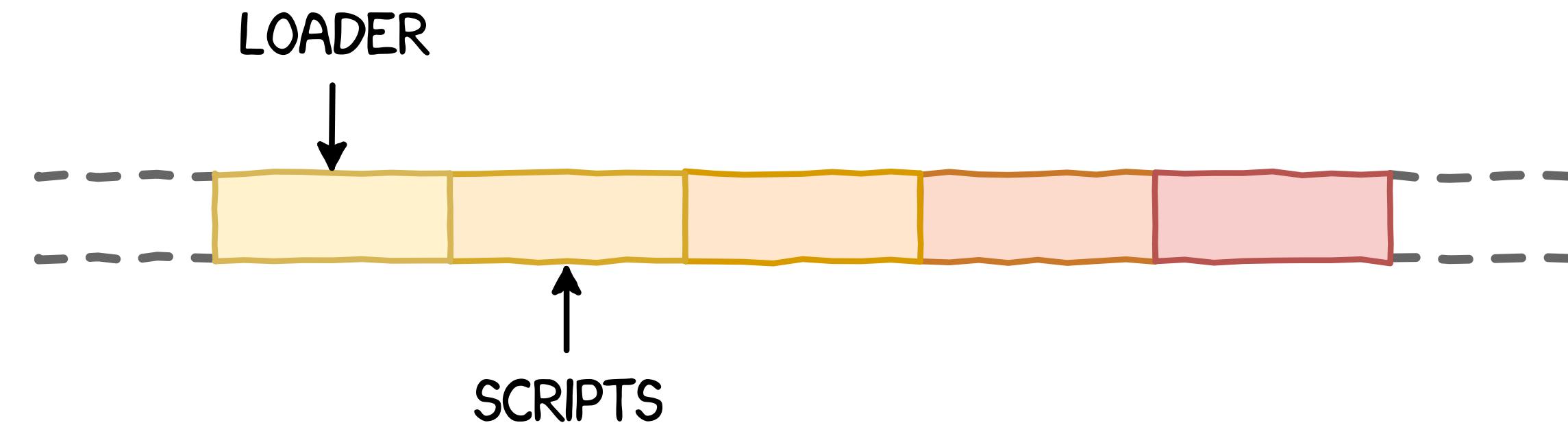


- mclf-ida-loader

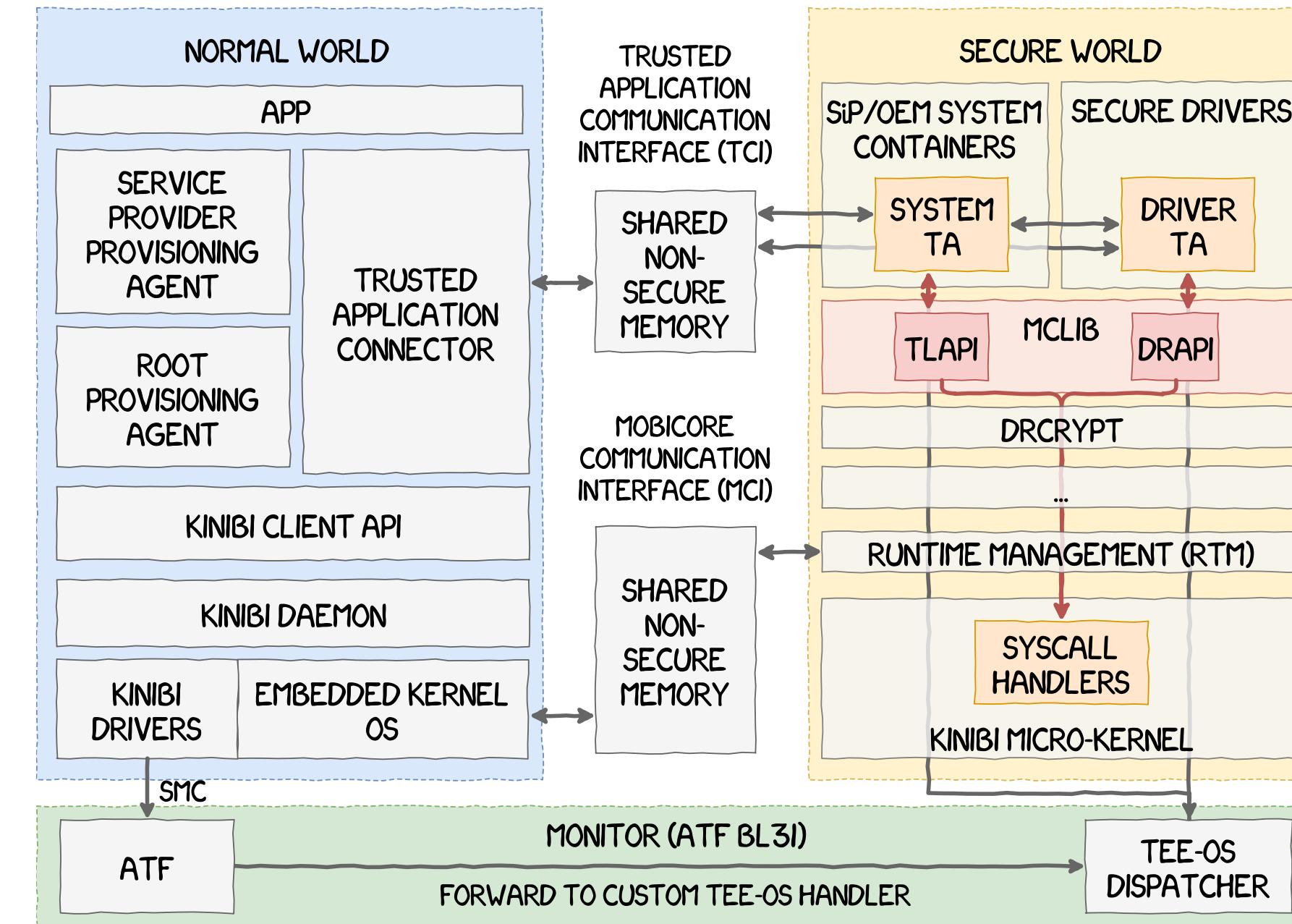


- mclf-ghidra-loader

STEP #2 - IDENTIFYING FUNCTIONS



MCLIB - STANDARD LIBRARY



- Renames tlApi/drAPI functions
- Sets the functions prototypes

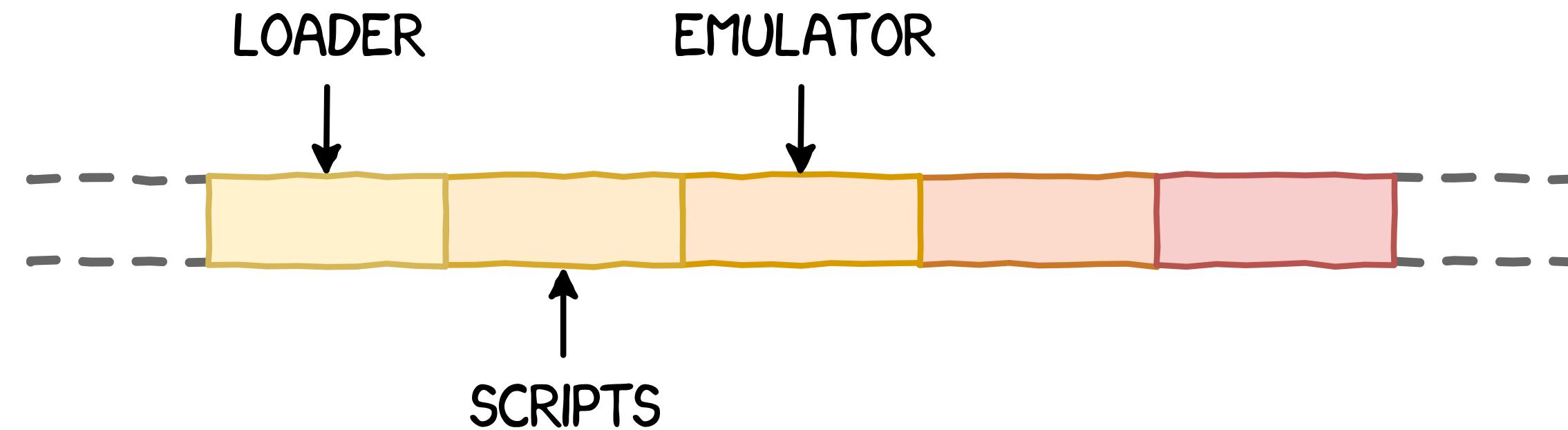
```
.text:00023EA8 ; ===== S U B R O U T I N E =====
.text:00023EA8
.text:00023EA8
.text:00023EA8 sub_23EA8          ; CODE XREF: _entry-4A80↑p
.text:00023EA8    LDR      R1, =dword_1000
.text:00023EAA    LDR.W   R2, [R1,#(tlApiLibEntry - 0x1000)]
.text:00023EAE    MOV     R1, R0
.text:00023EB0    MOVS    R0, #6
.text:00023EB2    BX     R2
.text:00023EB2 ; End of function sub_23EA8
.text:00023EB2
.text:00023EB2 ; -----
.text:00023EB4    DCD dword_1000      ; DATA XREF: sub_23EA8↑r
.text:00023EB8
.text:00023EB8 ; ===== S U B R O U T I N E =====
.text:00023EB8
.text:00023EB8 sub_23EB8          ; CODE XREF: sub_26CC+1A↑p
.text:00023EB8    ; sub_347C+C↑p
.text:00023EB8    LDR      R3, =dword_1000
.text:00023EBA    PUSH    {R4,R5}
.text:00023EBC    LDR.W   R4, [R3,#(tlApiLibEntry - 0x1000)]
.text:00023EC0    MOV     R3, R2
.text:00023EC2    MOV     R2, R1
.text:00023EC4    MOV     R1, R0
.text:00023EC6    MOV     R12, R4
.text:00023EC8    POP     {R4,R5}
.text:00023ECA    MOVS    R0, #0xE
.text:00023ECC    BX     R12
.text:00023ECC ; End of function sub_23EB8
```

Before

```
.text:00023EA8 ; ===== S U B R O U T I N E =====
.text:00023EA8
.text:00023EA8
.text:00023EA8 ; _DWORD __cdecl tlApiWaitNotification(_DWORD timeout)
.text:00023EA8 tlApiWaitNotification           ; CODE XREF: _entry-4A80↑p
.text:00023EA8    LDR      R1, =dword_1000
.text:00023EAA    LDR.W   R2, [R1,#(tlApiLibEntry - 0x1000)]
.text:00023EAE    MOV     R1, R0
.text:00023EB0    MOVS    R0, #6
.text:00023EB2    BX     R2
.text:00023EB2 ; End of function tlApiWaitNotification
.text:00023EB2
.text:00023EB2 ; -----
.text:00023EB4    off_23EB4      DCD dword_1000      ; DATA XREF: tlApiWaitNotification↑r
.text:00023EB8
.text:00023EB8 ; ===== S U B R O U T I N E =====
.text:00023EB8
.text:00023EB8 ; _DWORD __cdecl tlApiRandomGenerateData(_DWORD alg, _DWORD randomBuffer, _DWORD randomLen)
.text:00023EB8 tlApiRandomGenerateData           ; CODE XREF: sub_26CC+1A↑p
.text:00023EB8    ; sub_347C+C↑p
.text:00023EB8    LDR      R3, =dword_1000
.text:00023EBA    PUSH    {R4,R5}
.text:00023EBC    LDR.W   R4, [R3,#(tlApiLibEntry - 0x1000)]
.text:00023EC0    MOV     R3, R2
.text:00023EC2    MOV     R2, R1
.text:00023EC4    MOV     R1, R0
.text:00023EC6    MOV     R12, R4
.text:00023EC8    POP     {R4,R5}
.text:00023ECA    MOVS    R0, #0xE
.text:00023ECC    BX     R12
.text:00023ECC ; End of function tlApiRandomGenerateData
```

After

STEP #3 - MANUALLY FINDING VULNERABILITIES



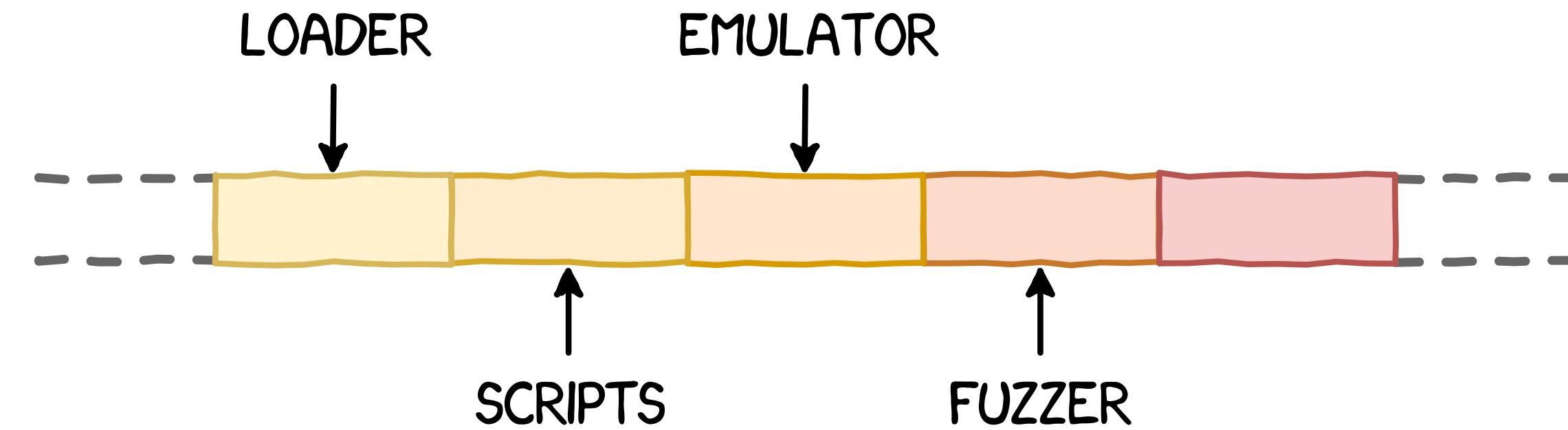
TRUSTLETS EMULATOR

- Based on Unicorn (external project)
- Split into simple tasks:
 - Loading the MCLF binary
 - Mapping the shared memory buffer
 - Hooking the McLib functions

TRUSTLETS EMULATOR

```
python emulator.py *41.tlbin cmd1.bin --tci 0x40100 -v
[+] Binary is a trustlet
[+] Trustlet size = 0x1ba4c
[+] Mapping text section at 0x00001000 with a size of 0x4874
[+] Mapping data section at 0x00007000 with a size of 0x168
[+] Mapping BSS section at 0x00007168 with a size of 0x17070
[+] Mapping region at 0x07d00000 (0x1 bytes)
[+] Mapping TCI buffer at 0x00100000 with a size of 0x40100
[i] drApiLogvPrintf(u'ICCC:Trustlet ICCC::Starting\n')
[+] Loading input data
[i] drApiLogvPrintf(u'TL ICCC: we got a command: 1\n')
[i] drApiLogvPrintf(u'ICCC: Initialize failed - tamper fuse set\n')
[i] drApiLogvPrintf(u'ICCC: Measurements result ret = 65548, ret hex = 1000c\n')
[i] drApiLogvPrintf(u'iccc: ICCC save data@\n')
[i] drApiLogvPrintf(u'Iccc_phys_read failed\n')
[i] drApiLogvPrintf(u'ICCC: check magic failed\n')
[i] drApiLogvPrintf(u'End of ICCC_Init, ret=1000c\n')
[i] drApiLogvPrintf(u'ICCC: Error writing Trustboot flag\n')
[+] tlApiNotify: Quitting!
```

STEP #4 - FINDING VULNERABILITIES AUTOMATICALLY



TRUSTLETS FUZZER

- Based on AFL_Uncorn (internal project)
 - Interfaces the fuzzer AFL with Unicorn
 - Usability and performance improvements
 - 100% of the code is written in Python!

TRUSTLETS FUZZER

process timing	overall results
run time : 0 days, 0 hrs, 4 min, 6 sec	cycles done : 0
last new path : 0 days, 0 hrs, 0 min, 3 sec	total paths : 65
last uniq crash : 0 days, 0 hrs, 0 min, 0 sec	uniq crashes : 21
last uniq hang : none seen yet	uniq hangs : 0
cycle progress	map coverage
now processing : 0 (0.00%)	map density : 0.02% / 1.79%
paths timed out : 0 (0.00%)	count coverage : 1.17 bits/tuple
stage progress	findings in depth
now trying : havoc	favored paths : 1 (1.54%)
stage execs : 4253/6528 (65.15%)	new edges on : 65 (100.00%)
total execs : 5407	total crashes : 88 (21 unique)
exec speed : 21.94/sec (slow!)	total tmouts : 0 (0 unique)
fuzzing strategy yields	path geometry
bit flips : 2/32, 1/31, 2/29	levels : 2
byte flips : 1/4, 0/3, 0/1	pending : 65
arithmetics : 10/224, 0/204, 0/68	pend fav : 1
known ints : 1/8, 0/18, 0/10	own finds : 64
dictionary : 0/0, 0/0, 0/0	imported : n/a
havoc : 0/0, 0/0	stability : 100.00%
trim : 50.00%/1, 0.00%	[cpu000: 6%]

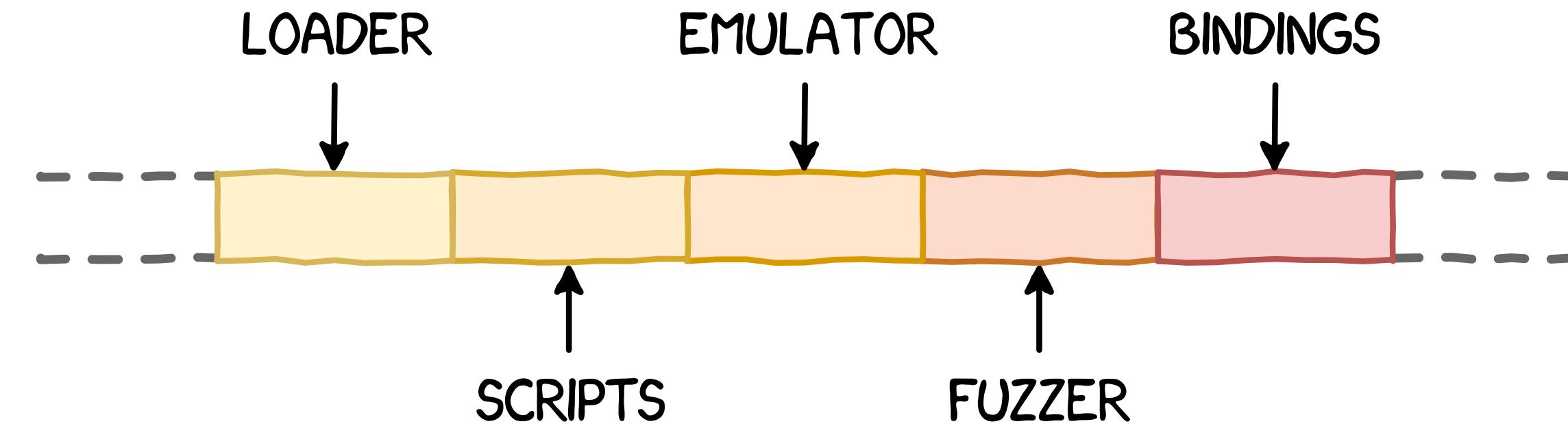
TRUSTLETS SYMBOLIC EXECUTOR

- Based on Manticore by Trail of Bits
- Uses very simple strategies:
 - Mark the shared memory buffer symbolic
 - Explore all the paths of the trustlet
 - Check reads or writes to memory
 - Ask the solver for an invalid address

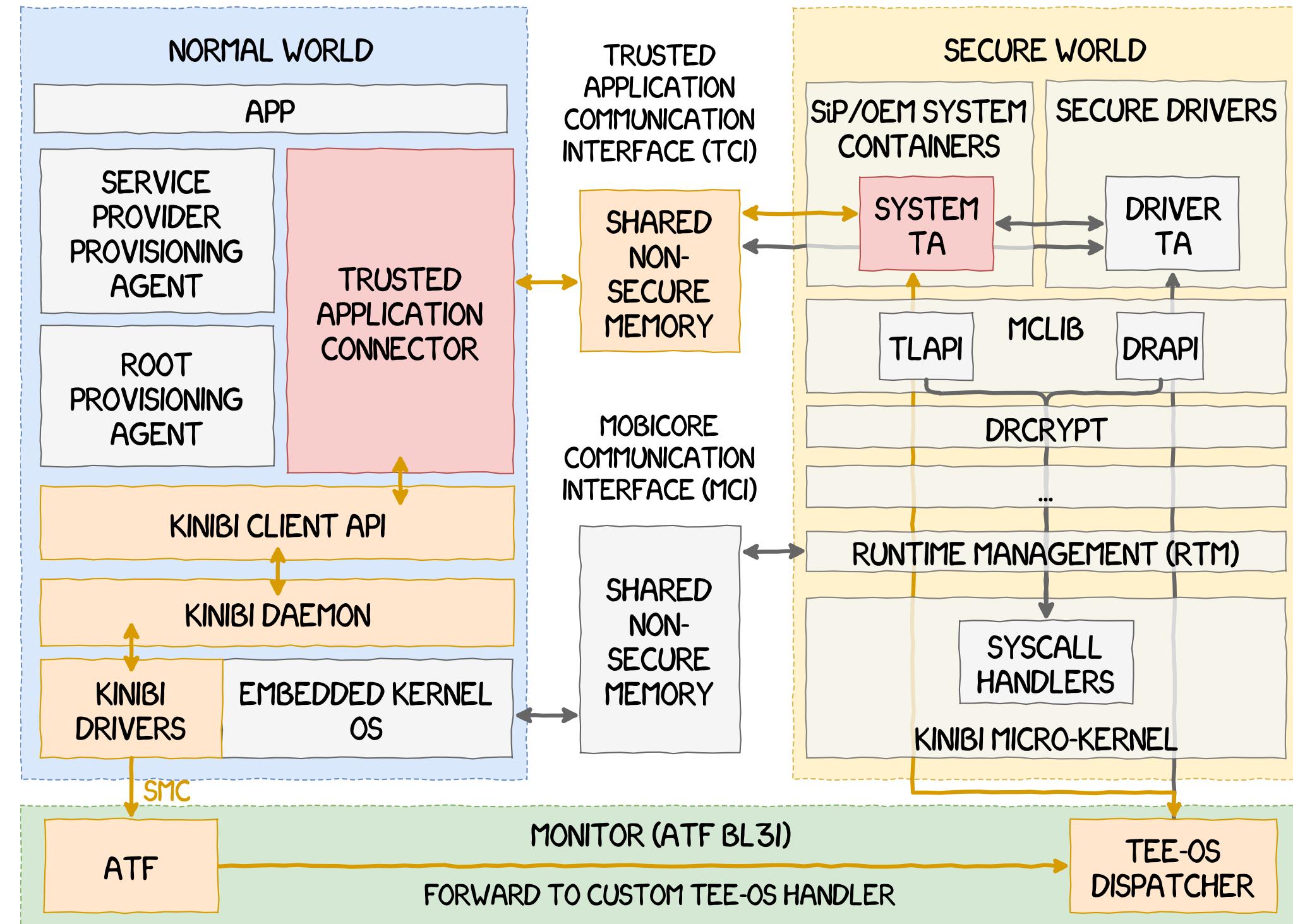
CRASH EXAMPLE

```
Command line:  
'./tainter.py -s 1036 ffffffff000000000000000000000005.tlbin.elf -t -c coverage.txt'  
Status:  
    Invalid symbolic memory access (mode:r)  
  
===== PROC: 00 =====  
Memory:  
000000000001000-000000000008000  r x 00000094 ffffffff000000000000000000000005.tlbin.elf  
000000000009000-0000000000024000  rw 00006dc8 ffffffff000000000000000000000005.tlbin.elf  
000000000100000-0000000000101000  rw 00000000  
0000000007d0000-0000000007d01000  rx 00000000 CPU:  
INSTRUCTION: 0x0000000000059ec:          pld      [r1, #0x80]  
APSR: 0x000000060000000  
R0 : 0x000000000009aac  
R1 : <BitVecExtract at 7f2571dbdeb8-T>  
R10: 0x000000000000000  
R11: 0x000000000000000  
R12: 0x000000000000000  
R13: 0x0000000000023a28
```

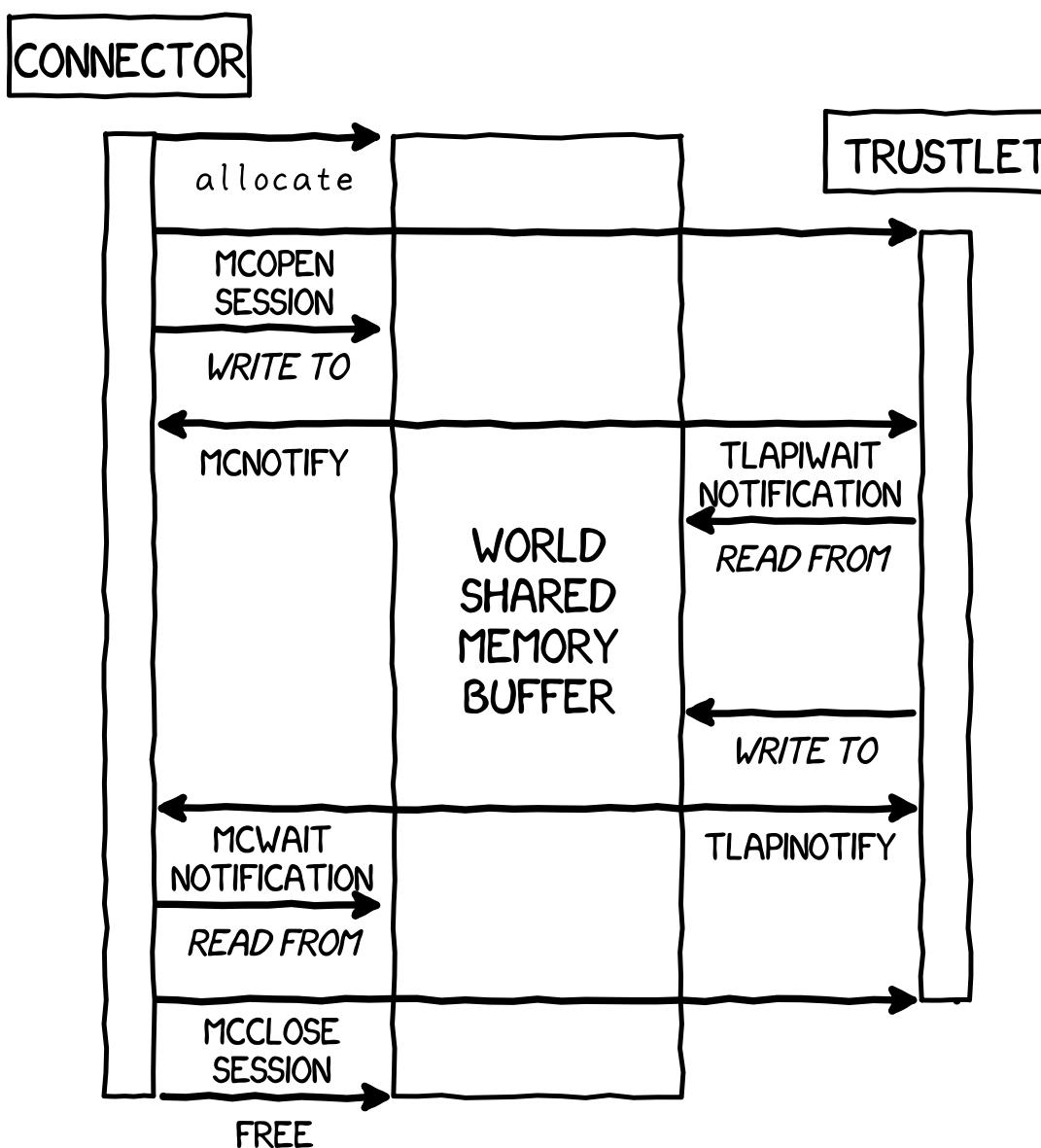
STEP #5 - EXPLOITING THE VULNERABILITIES



SOFTWARE STACK



CLIENT API



PYTHON BINDINGS

- Writing C is tedious, writing Python is a lot easier
- Bindings of the mcClient API called `pymcclient`
- Provides various utilities: hexdump, (dis)assemble, etc.
- Provides a command interpreter which is based on IPython

SCRIPT EXAMPLE

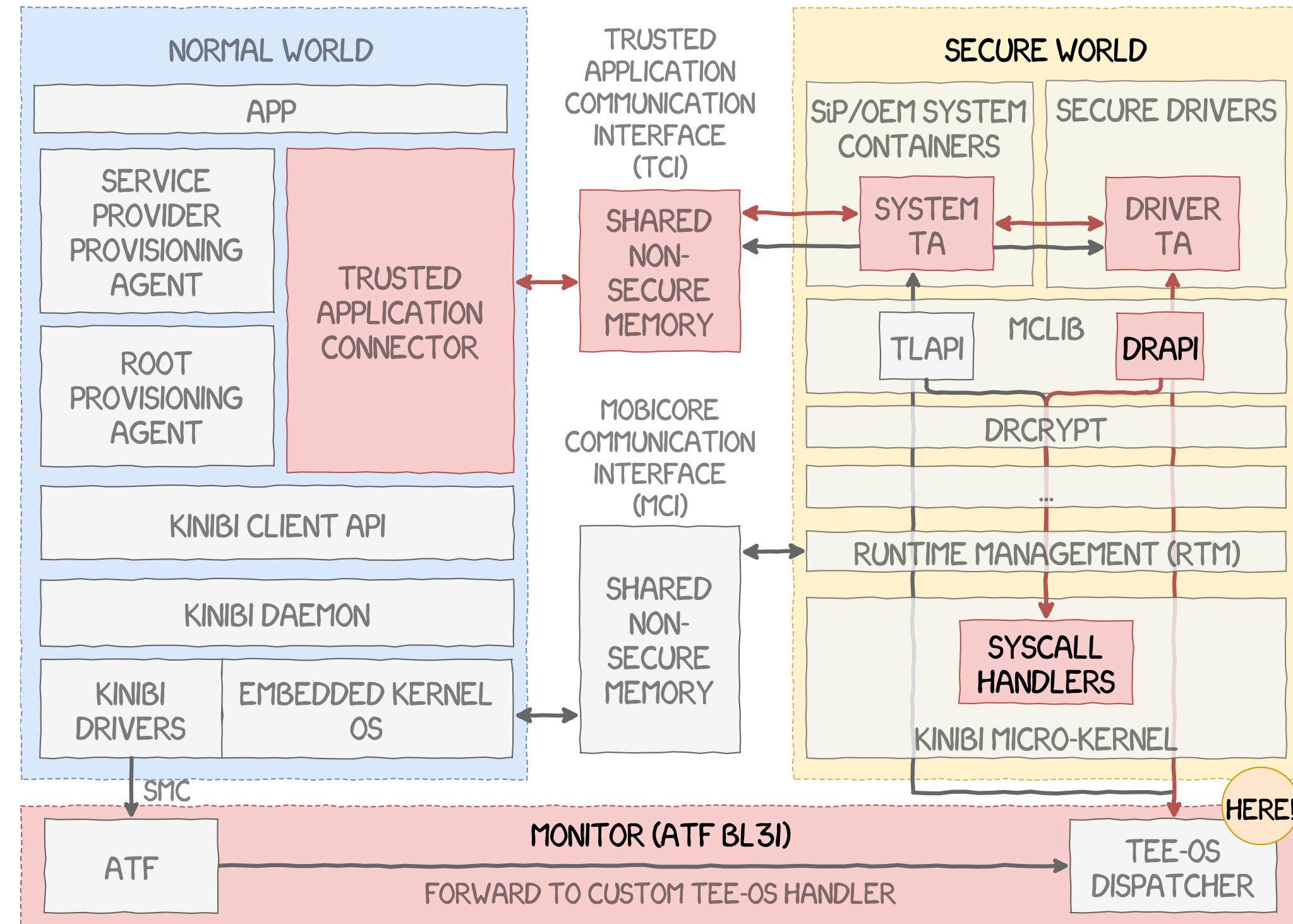
```
with Device(DEVICE_ID) as dev:  
    with dev.buffer(TCI_BUFFER_SIZE) as tci:  
        with open(TRUSTLET_FILE, "rb") as fd:  
            buf = fd.read()  
  
        with Trustlet(dev, tci, buf) as app:  
            tci.seek(0)  
            tci.write_dword(1)  
  
            app.notify()  
            app.wait_notification()  
  
            tci.seek(0)  
            print(tci.read_dword())
```

VULNERABILITY ANALYSIS & EXPLOITATION

OVERVIEW

- **Target:**
 - Samsung Galaxy S7 running Android 7.0
- **Main goal:**
 - Obtaining code execution in EL3
- **Prerequisites:**
 - Being part of the radio group
 - Being able to write files somewhere on the device

ATTACK PLAN

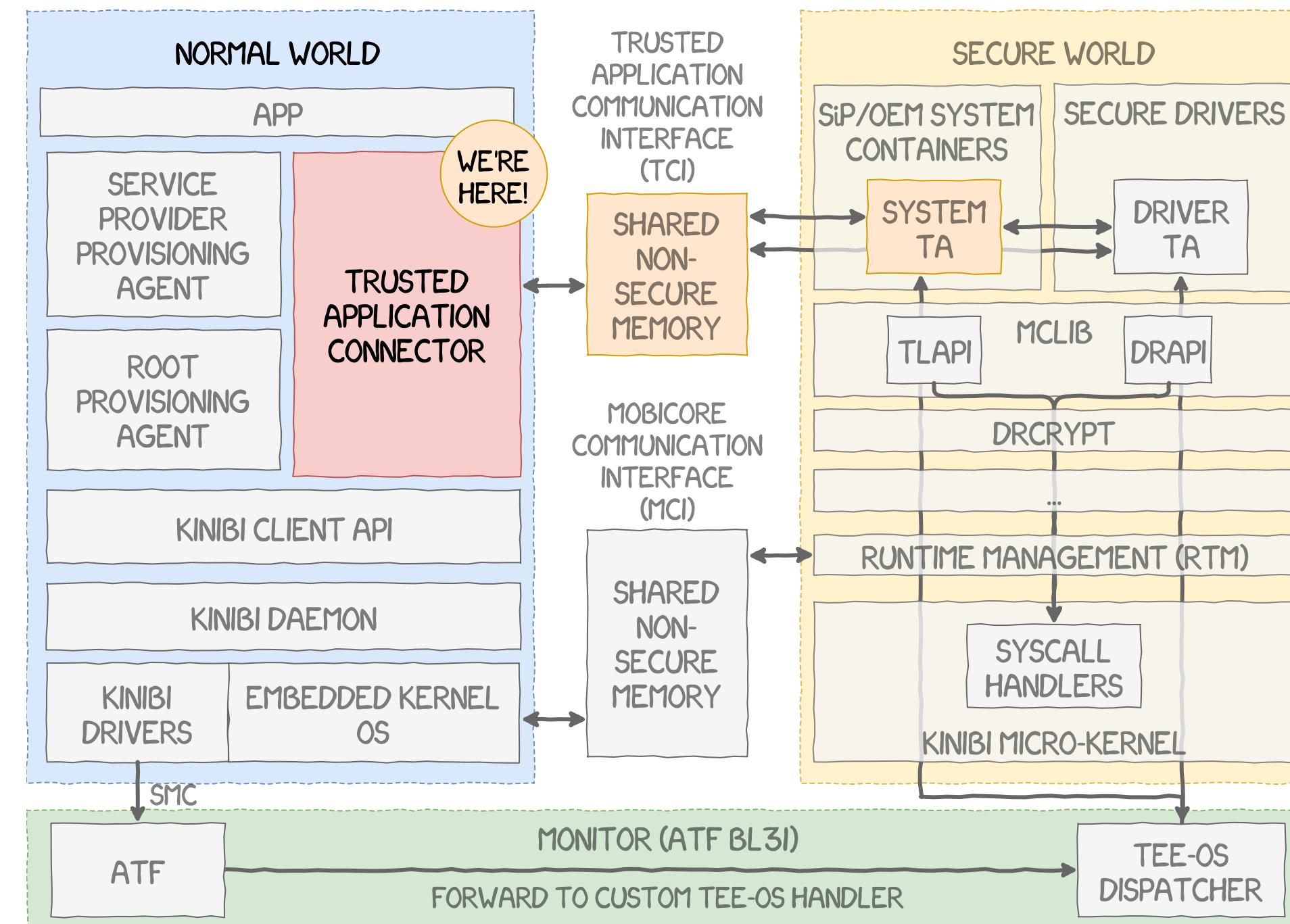


SOFTWARE MITIGATIONS

Model	XN bit	Canary	ASLR	PIE
S6	✓	✗	✗	✗
S7	✓	✗	✗	✗
S8	✓	✗	✗	✗
S9	✓	✓	✗	✗

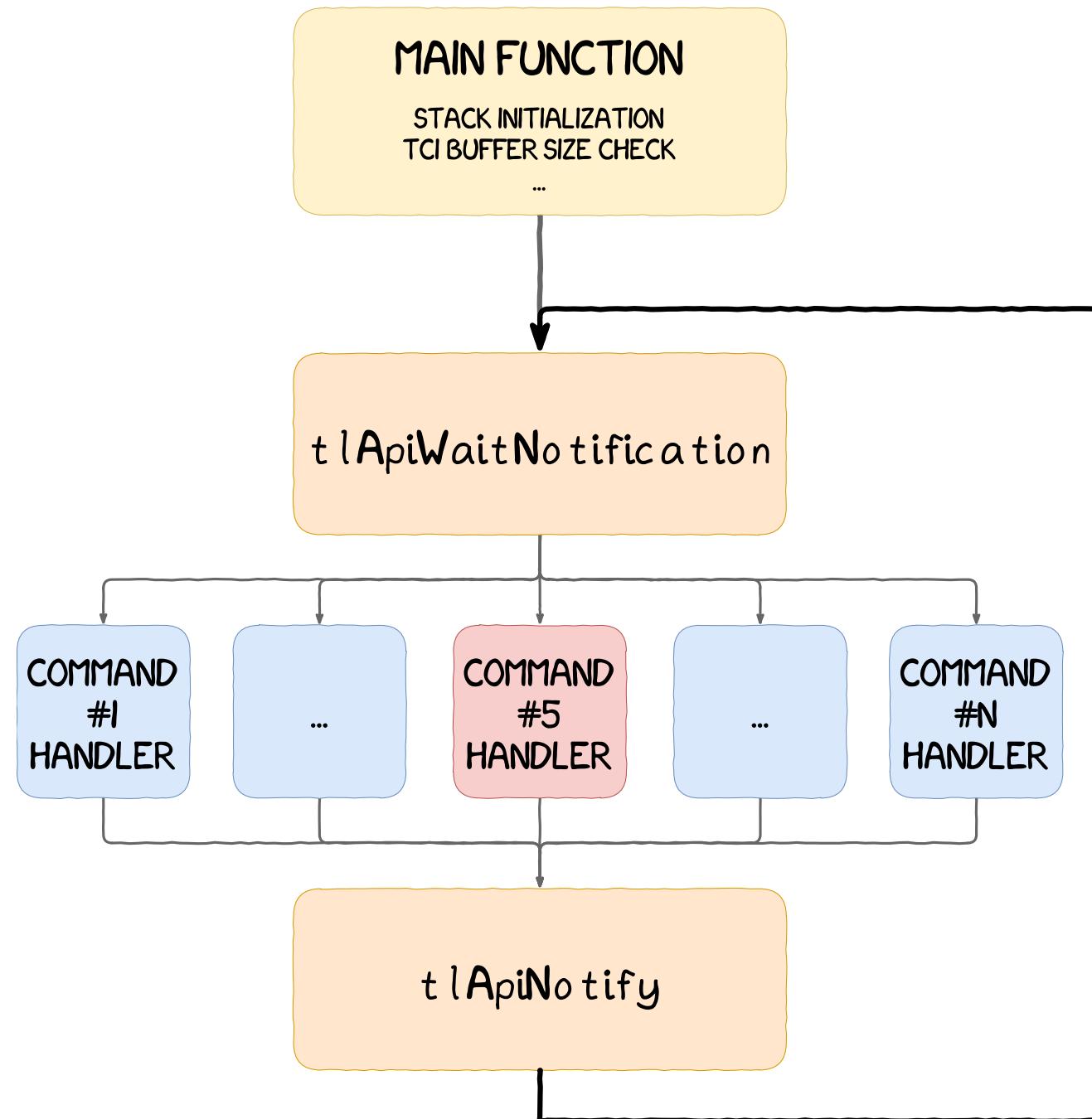
ATTACKING A TRUSTED APPLICATION

Overview



ATTACKING A TRUSTED APPLICATION

SEM Trustlet Vulnerability



Stack-based buffer overflow in the handler of the command ID #5

Call to `memcpy` at the beginning of the 5th command handler

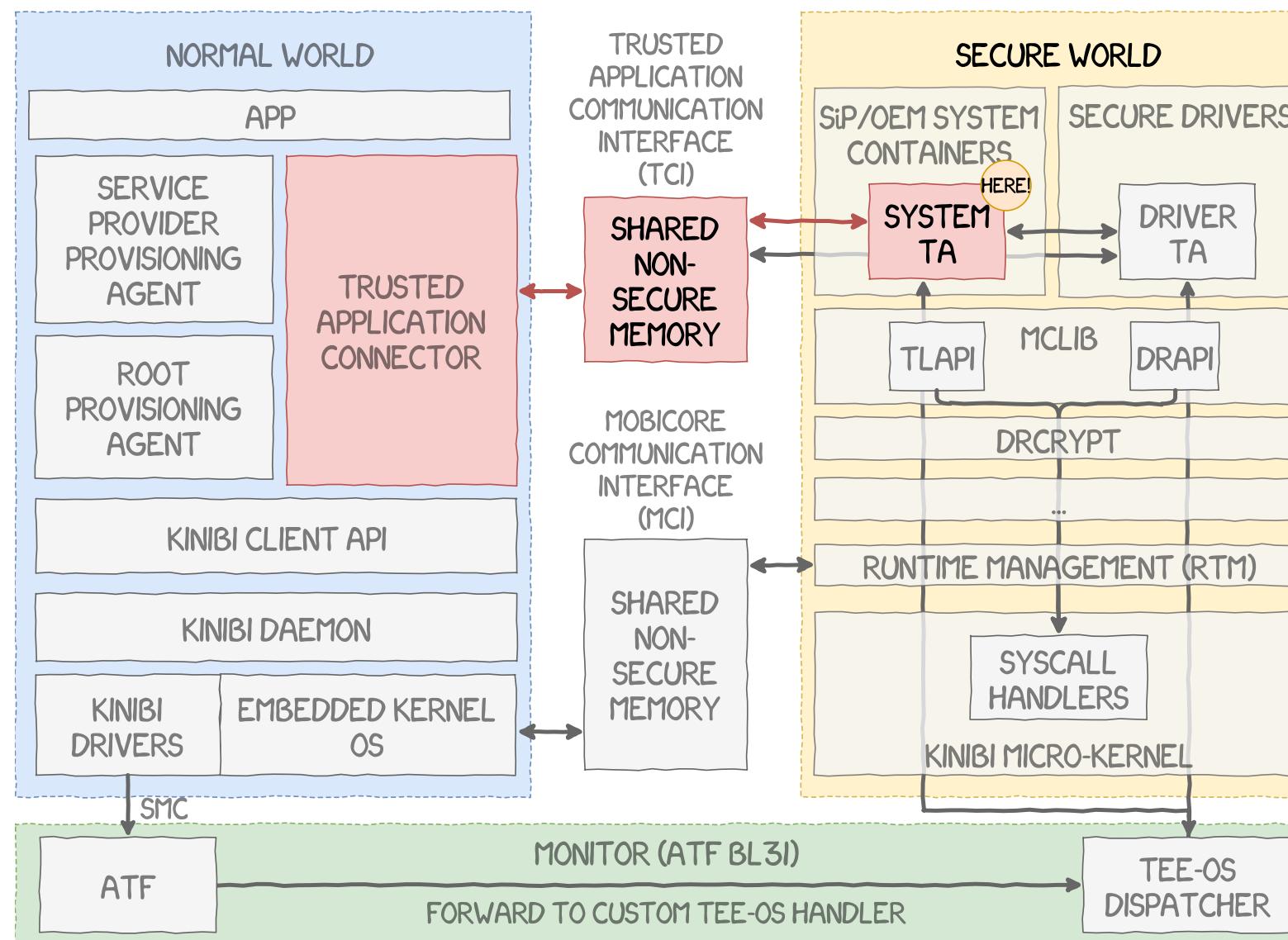
Before this call, the registers are set as follow:

- R0 = SP+0x4F8-0xF0, the destination buffer
- R1 = `tci_buffer + 0x8`, the source buffer
- R2 = `*(tci_buffer + 0x16808)`, the length of the buffer

.text:00020FB2	ADD.W	R1, R0, #0x16000
.text:00020FB6	MOV	R4, R1
.text:00020FB8	LDR.W	R2, [R1,#0x808]
.text:00020FBC	ADD.W	R1, R0, #8
.text:00020FC0	ADD.W	R0, SP, #0x4F8+var_F0
.text:00020FC4	BLX	memcpy_aligned

ATTACKING A TRUSTED APPLICATION

Exploitation Results



- Code execution in **S-ELO**
- **It is now possible to:**
 - Communicate with **Secure Drivers**
 - Make some syscalls (e.g. print characters, get system information, etc.)
- **Next target:** Secure Driver

ATTACKING A SECURE DRIVER

VALIDATOR Secure Driver Vulnerability

A vulnerability was found in the **VALIDATOR secure driver**

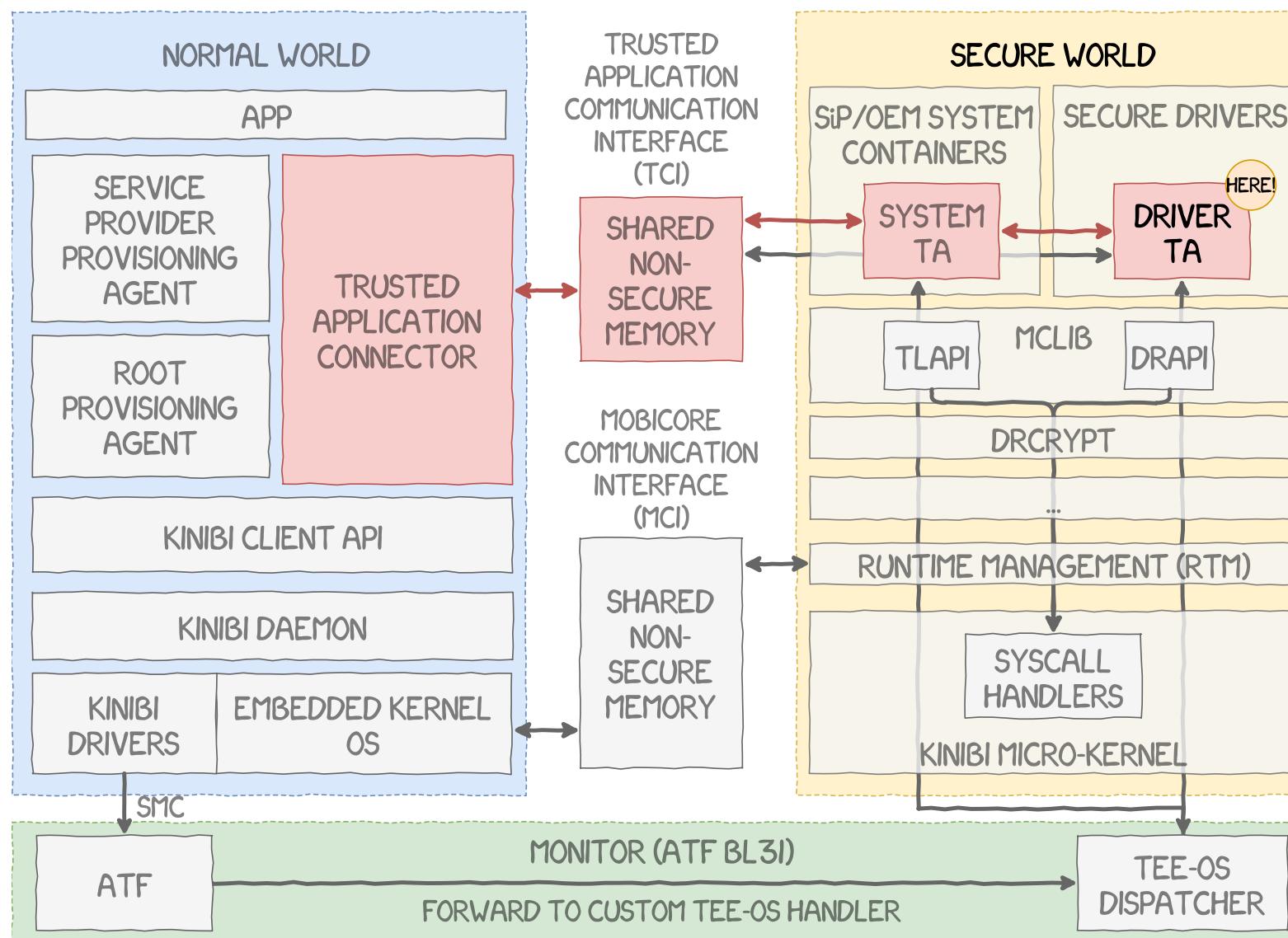
Stack-based buffer overflow in the handler of the command ID #15

Equivalent to the one found in the trustlet (i.e. `memcpy` in the stack and a user-controlled size)

```
.text:00001362          MOVS      R2, #0x37 ; '7'  
.text:00001364          MOV       R1, R4  
.text:00001366          ADDS      R0, R6, #1  
.text:00001368          BLX      memcpy
```

ATTACKING A SECURE DRIVER

Exploitation Results



- Code execution in **S-EL0** but with higher privileges
- **It is now possible to:**
 - Communicate with the **RunTime Manager** (or RTM, an init-like process)
 - Access more syscalls (e.g. map physical memory, create threads, make SMCs, etc.)
- **Next target:** Trusted OS & Monitor

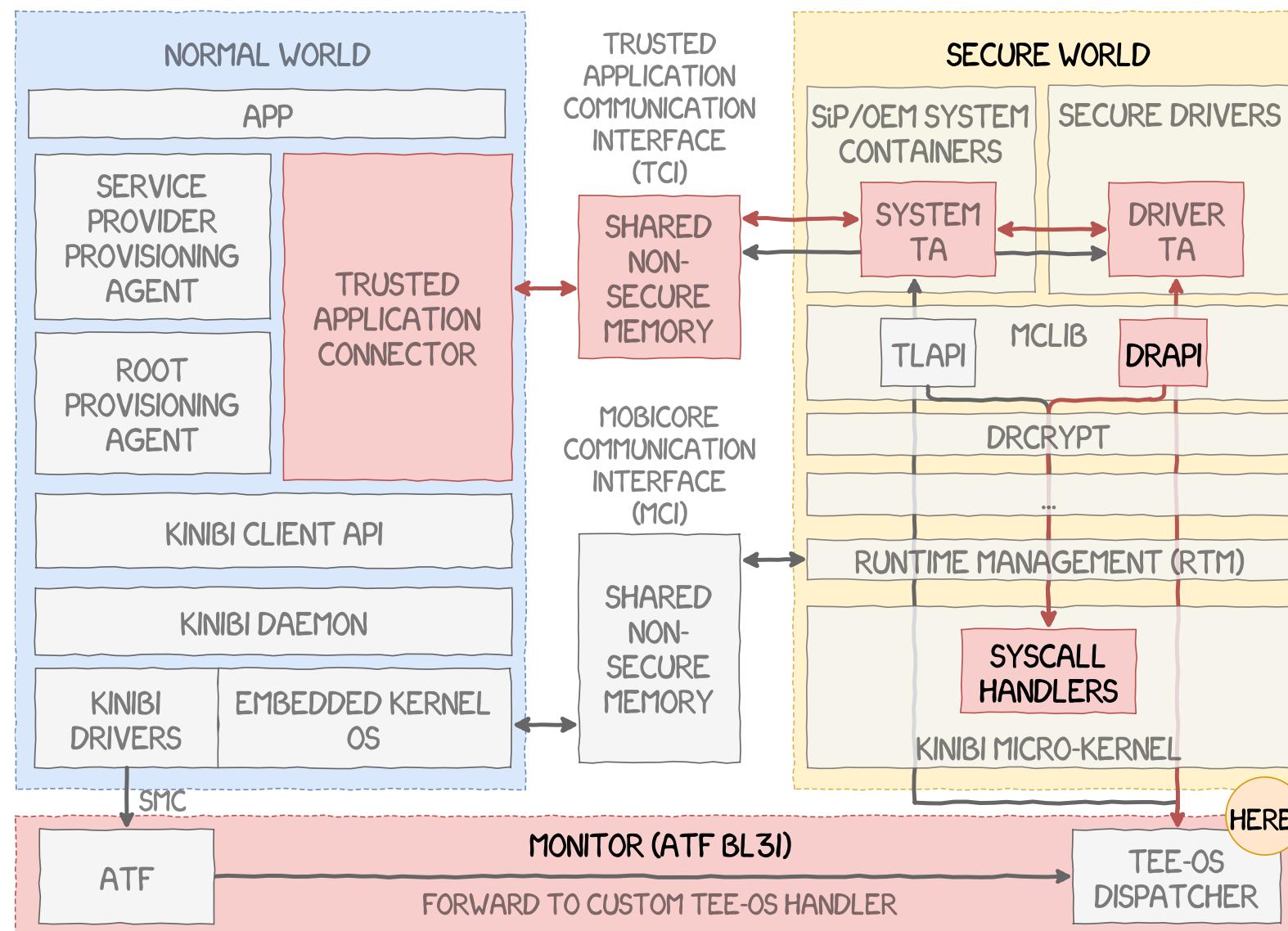
ATTACKING KINIBI AND THE MONITOR

Vulnerability Analysis

- **mmap**: secure and non-secure physical memory mapping syscall
- **Vulnerability**
 - Monitor mapped at `0x2022000`
 - Can be mapped using `mmap` to modify an SMC
 - Calling the hijacked SMC allows code execution in EL3
- **Patch**
 - Fixed in the newest versions by using a blacklist

ATTACKING KINIBI AND THE MONITOR

Exploitation Results



- Code execution in **EL3**
- Now possible to do anything we want!

POST-EXPLOITATION

TRUSTPWN FRAMEWORK

- Based on the previous vulnerabilities
- **Internals**
 - Uses the EL3 vulnerability to have arbitrary access to Kinibi
 - Adds a SVC and a drApi function to execute code in S-EL1 "natively"
 - SVCs and drApi functions are referenced in pointer arrays
- **Usage**
 - Read or write memory arbitrarily
 - Execute code in S-EL1 and EL3

DEMO

Finding the Master Key in the Monitor

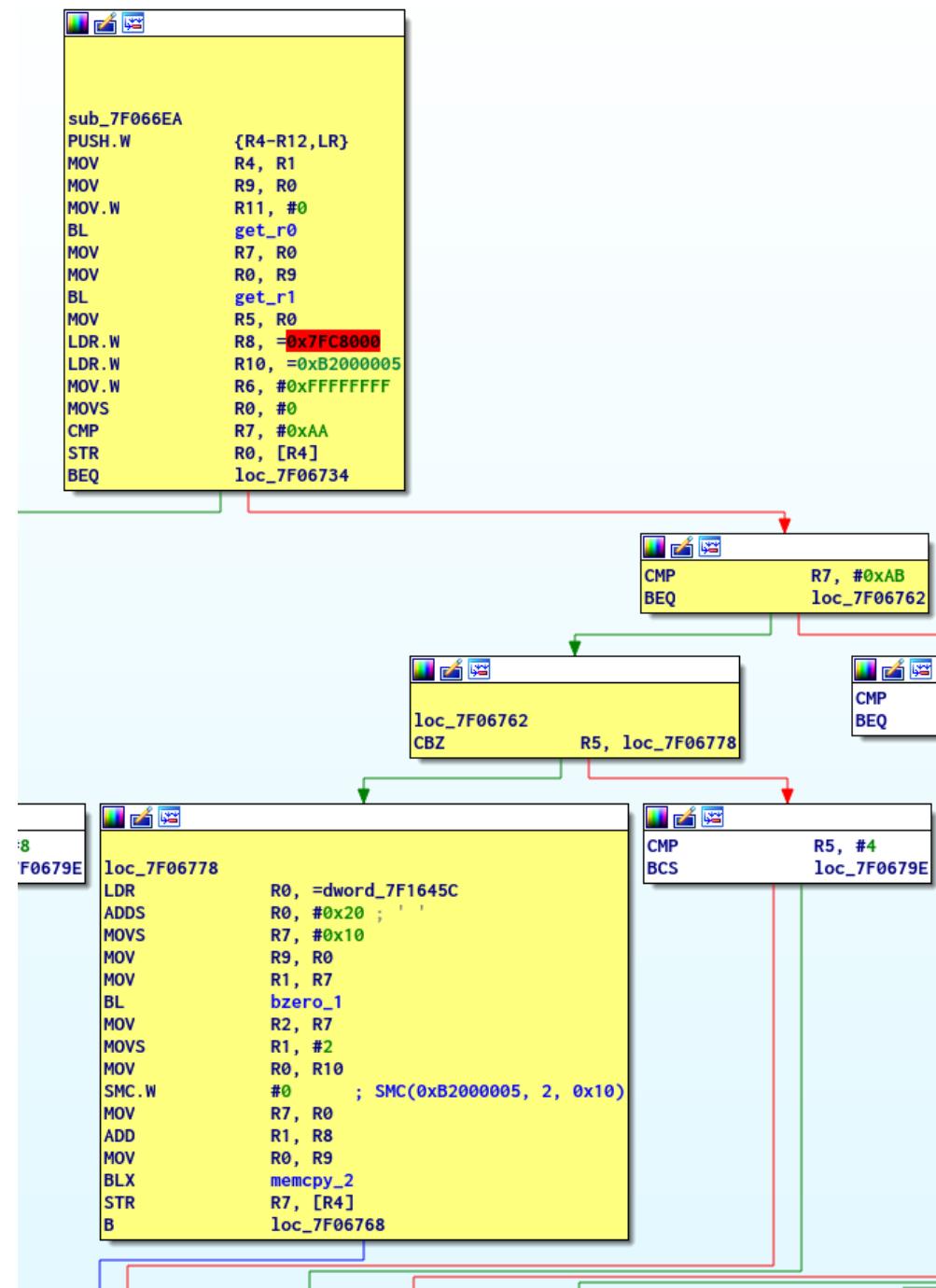
FINDING THE MASTER KEY IN THE MONITOR

DrApi Reversing

- Reversing the crypto-driver **drcrypto** (found embedded in Kinibi)
- **DrApi 0x1030**
 - Takes four possible command IDs (0xAA, 0xAB, 0xAC, 0xAD)
 - The interesting one is 0xAB
 - Wrapper around SMC 0xB2000005
- **SMC 0xB2000005**
 - SMC arguments:
 - **R0:** SMC ID
 - **R0:** command number (four possible values [0-3])
 - **R1:** number of bytes to read
 - Reads 0x10 bytes of the master key at 0x101E4000

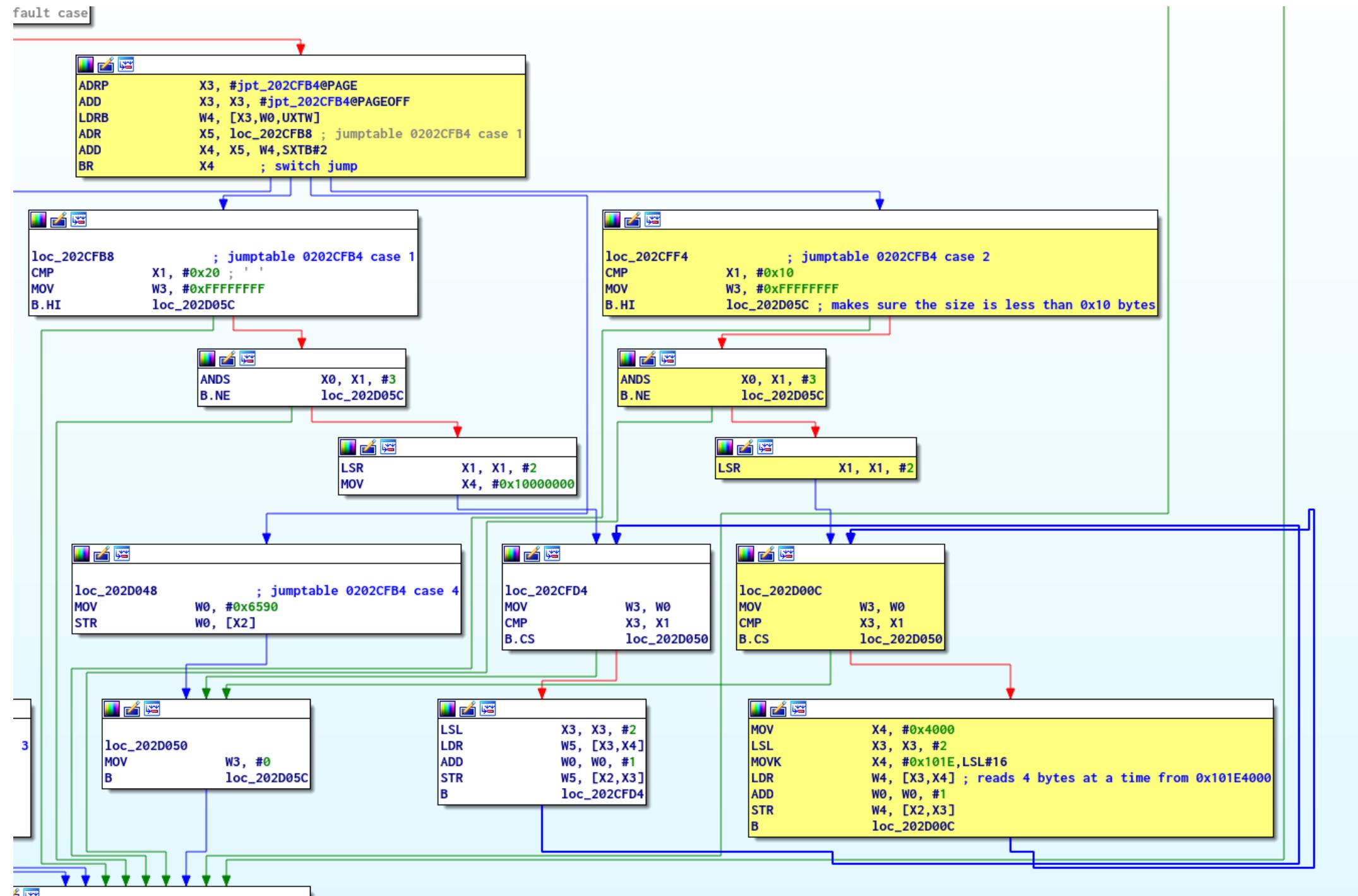
FINDING THE MASTER KEY IN THE MONITOR

DrApi Function



FINDING THE MASTER KEY IN THE MONITOR

SMC Function



DEMO

Bypassing Signature Checks

BYPASSING SIGNATURE CHECKS

Methodology

- Reversing RTM
- Finding the SHA-256 of the public key corresponding to the private key used to sign TAs and SDs
- Signature is verified using `tlApiSignatureVerify`
- Patch the checks and load your own TA or SD

FINDING THE MASTER KEY IN THE MONITOR

RTM Verifications

- **First check**

ROM: 00006E62	BL	tlApiSignatureVerify
ROM: 00006E66	LDR	R4, =0x40B00009
ROM: 00006E68	ADDS	R4, R4, #6
ROM: 00006E6A	CBNZ	R0, loc_6E7A
ROM: 00006E6C	LDRB.W	R0, [SP, #0xC0+var_44]

- **Second check**

ROM: 000073E0	BL	tlApiSignatureVerify
ROM: 000073E4	CBNZ	R0, loc_73FA
ROM: 000073E6	LDRB.W	R0, [SP, #0x1C0+var_48]

DEMO

Trusted-OS Instrumentation

TRUSTED-OS INSTRUMENTATION

Methodology

- Handles ARMv7 and Thumb
- Based on the **Undefined Instruction** exception
- **Undefined Instruction** handler is replaced by our own code
- Patch an instruction with the ARM undefined instruction UDF
0xNNNN
- When a breakpoint triggers the current context of the CPU is saved
 - Current context is saved
 - Overwritten instruction is executed

THANK YOU!