

# CS & IT ENGINEERING

## Theory of Computation

Turing Machine : Recursively Enumerable

Turing Machine (Part 3)



Lecture No. 3



By- DEVA Sir

## TOPICS TO BE COVERED

01 Closure Properties

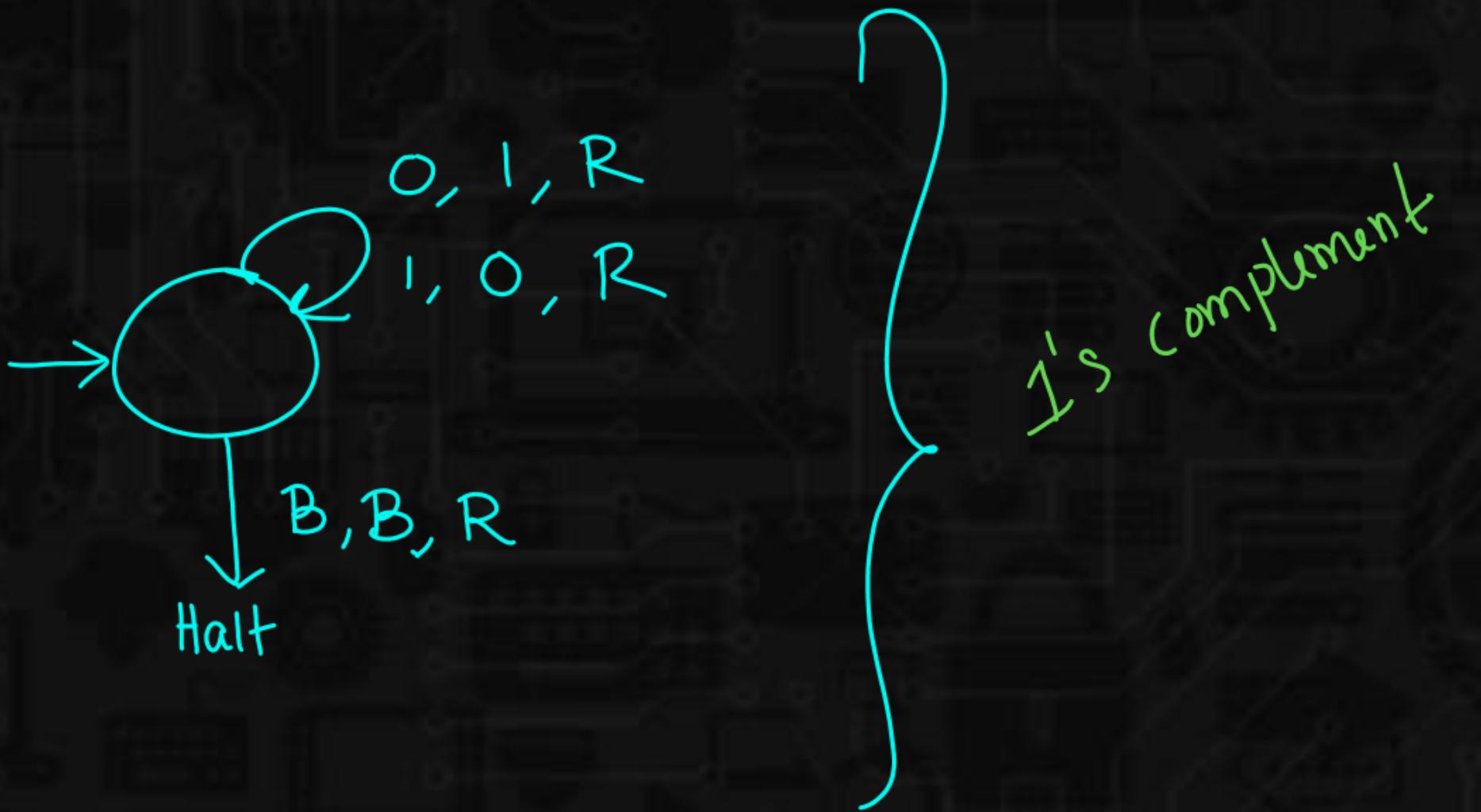
02 Terminology: Decidable

03 Terminology: Undecidable

04

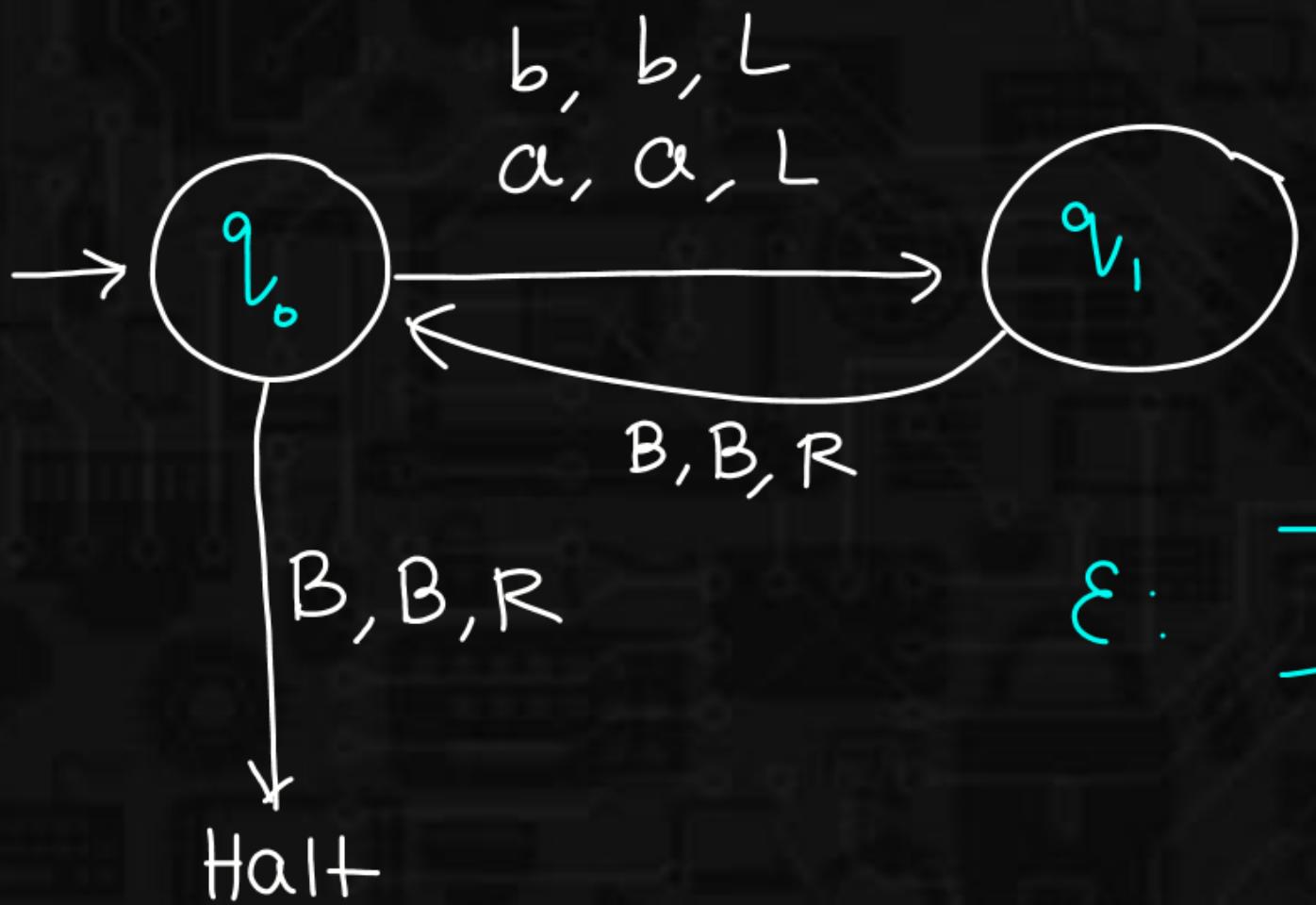
# Practice

I)

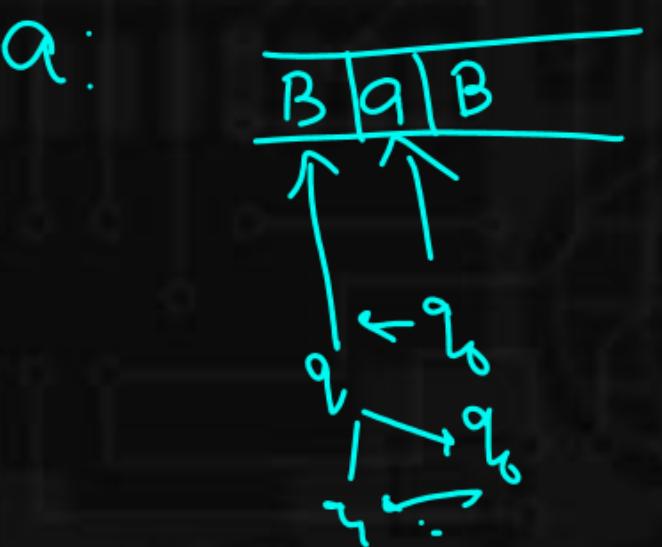
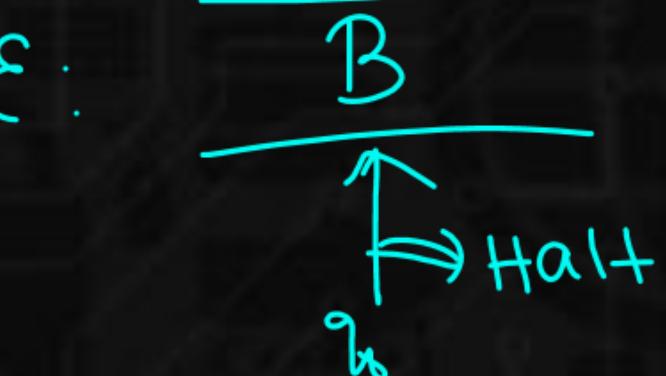


P  
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II)



$(a+b)^+$  doesn't halt  
 $\epsilon$  halts



a ↪

b ↪

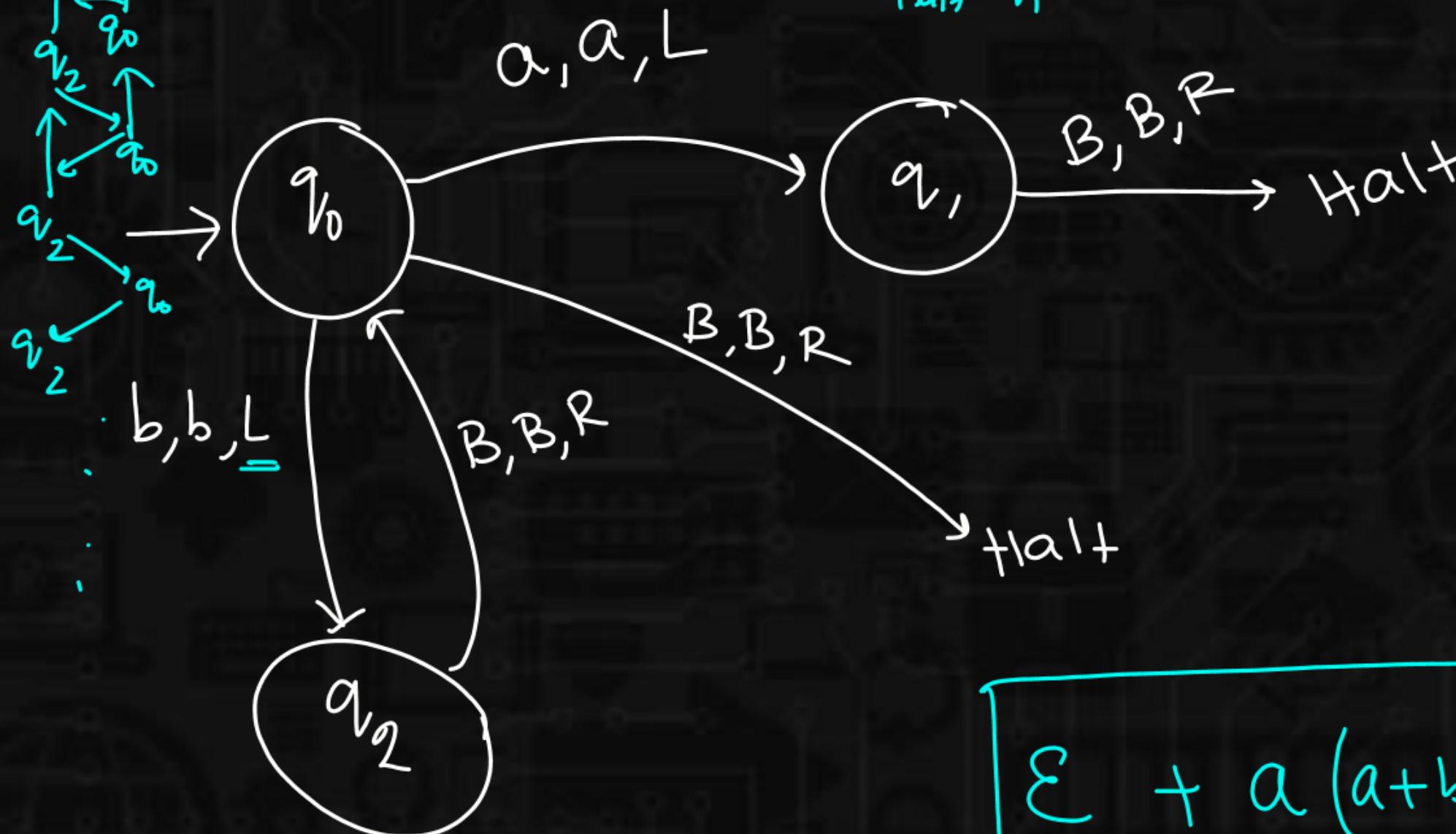
aa ↪

ab ↪

ba ↪

bb ↪

III)  $BbabB$



$\overline{BabB}$

Halt

$q_0$

$q_1$

Halt

$q_2$

$q_0$

$q_2$

$q_0$

$q_1$

$q_2$

$q_0$

$q_1$

$q_2$

$(a+b)^*$

↓

$\epsilon \rightarrow \text{Halt}$

$a \rightarrow a$

$b \rightarrow b$

$a^a \rightarrow a^a$

$a^b \rightarrow a^b$

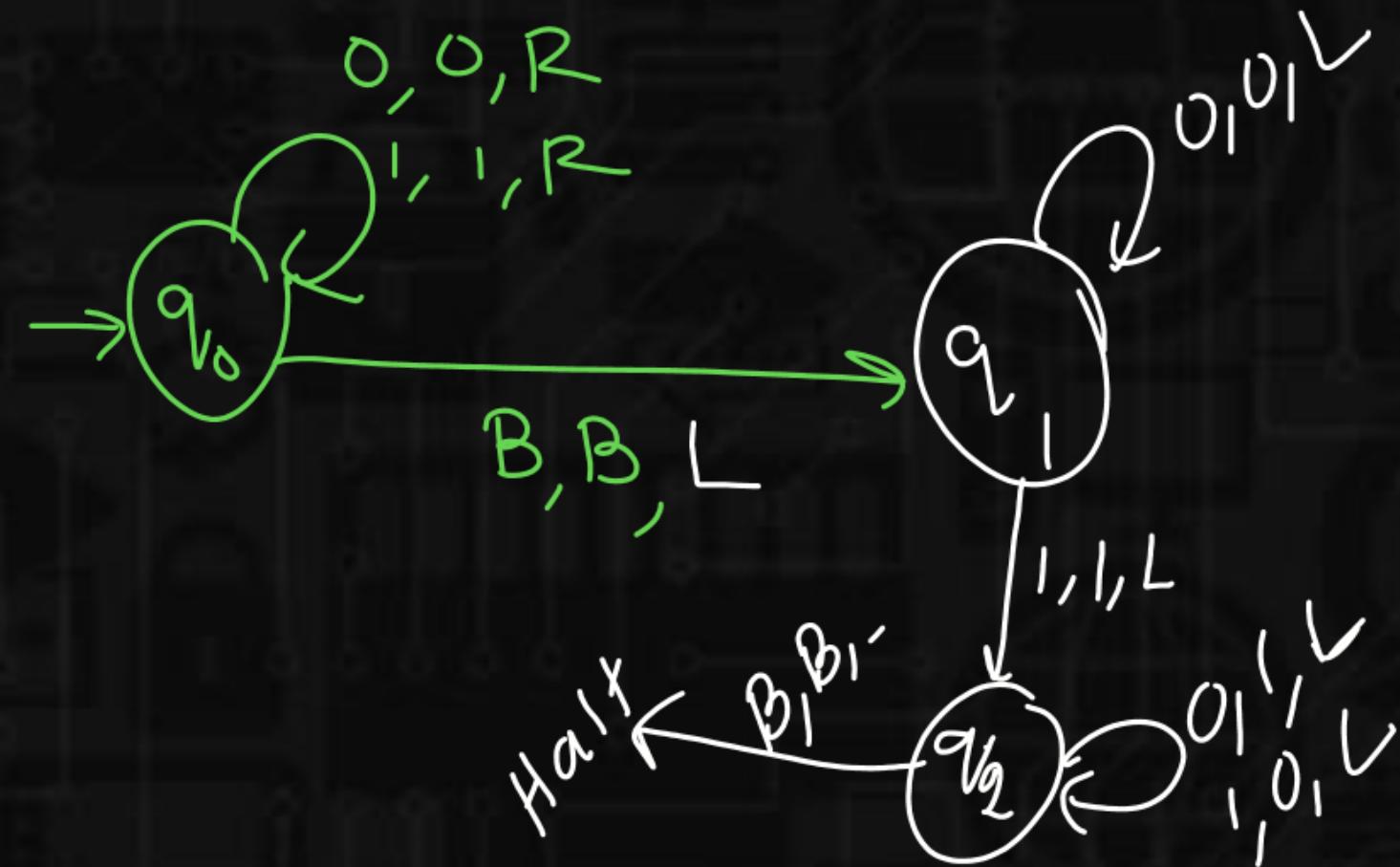
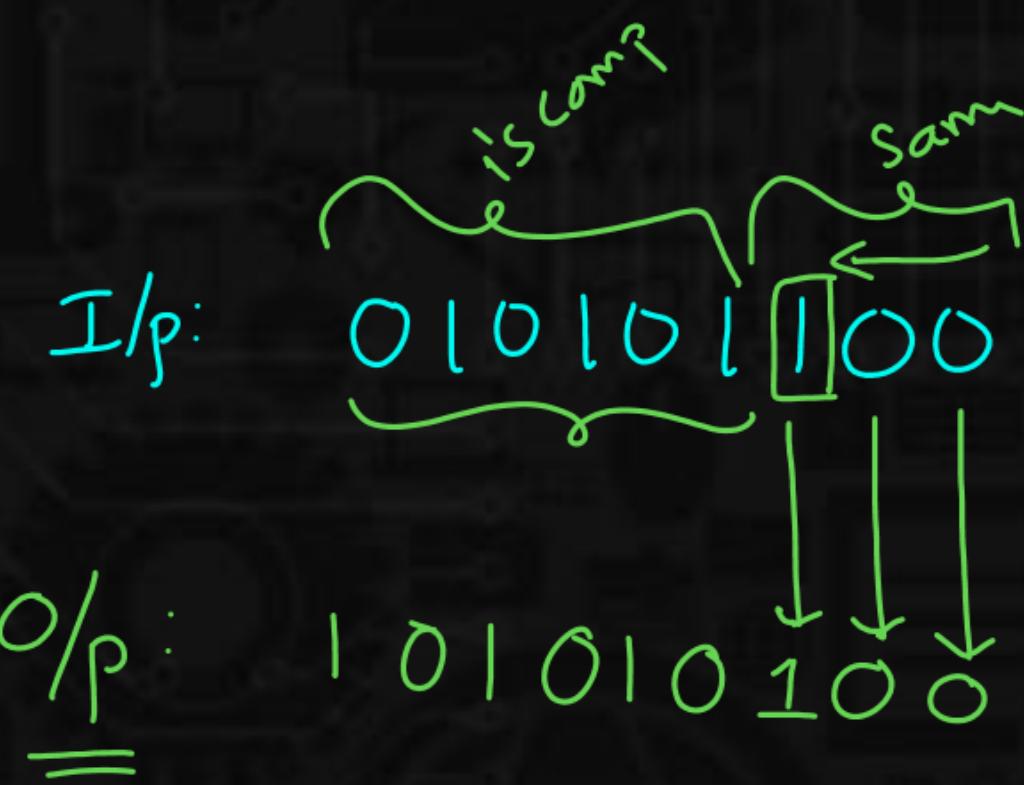
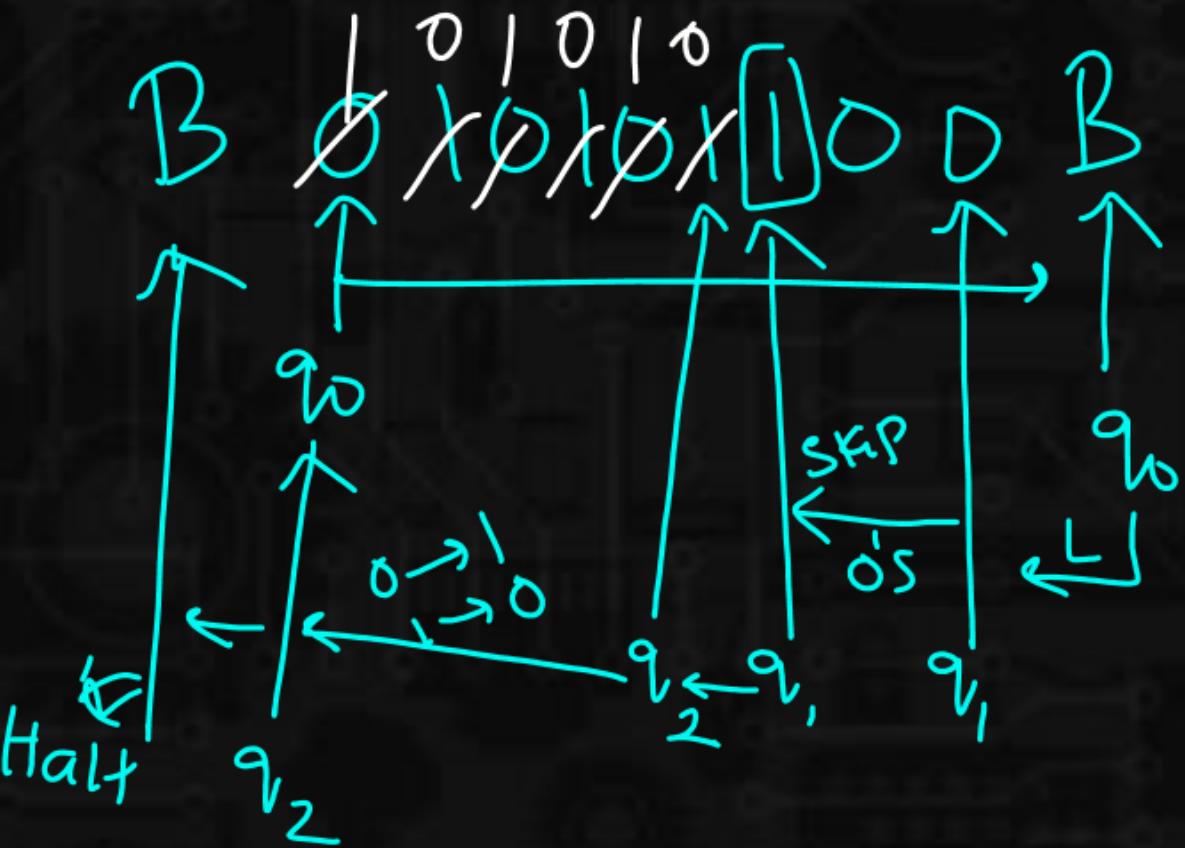
$b^a \rightarrow b^a$

$b^b \rightarrow b^b$

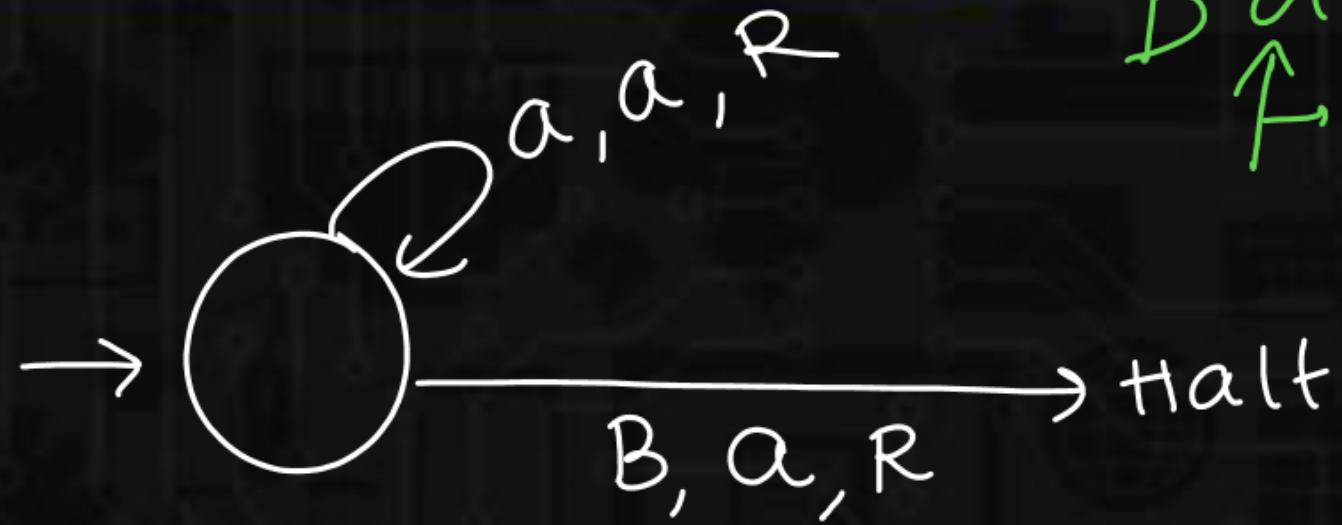
$\boxed{\epsilon + a(a+b)^* \rightarrow \text{Halt}}$

$\boxed{b(a+b)^* \Rightarrow \text{Never halts}}$

## IV) 2's complement



V)



$B \xrightarrow{a} a \xrightarrow{a} a \xrightarrow{a} B \xrightarrow{B} B \xrightarrow{B}$

A) Increment

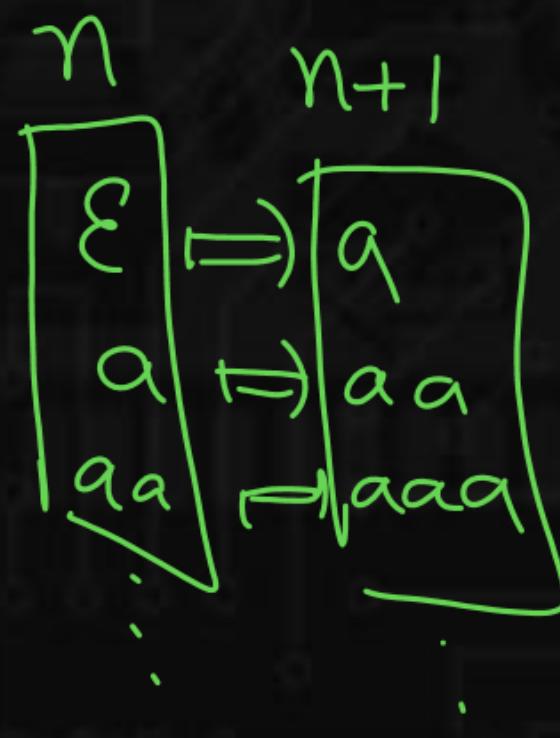
B) Decrement

~~C) Accepts  $a^*$~~

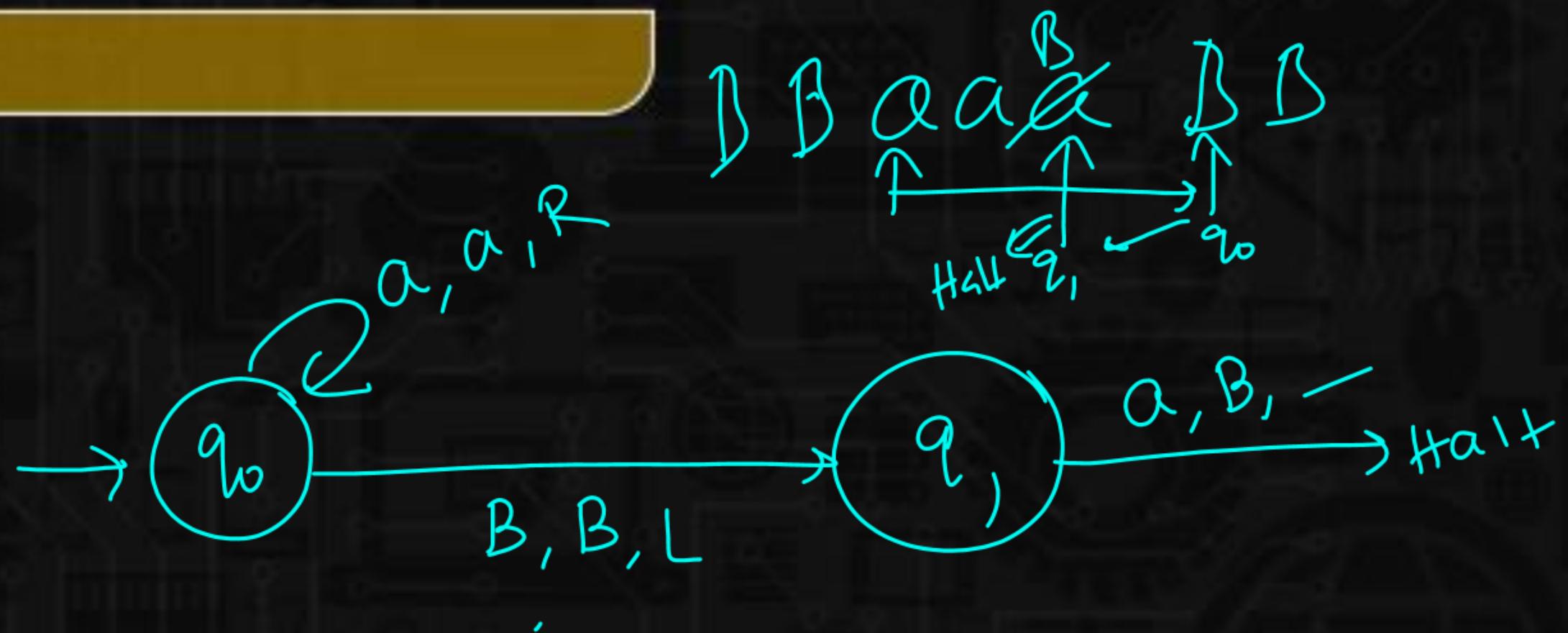
D) None

Unary Increment:

$$f(x) = x + 1$$



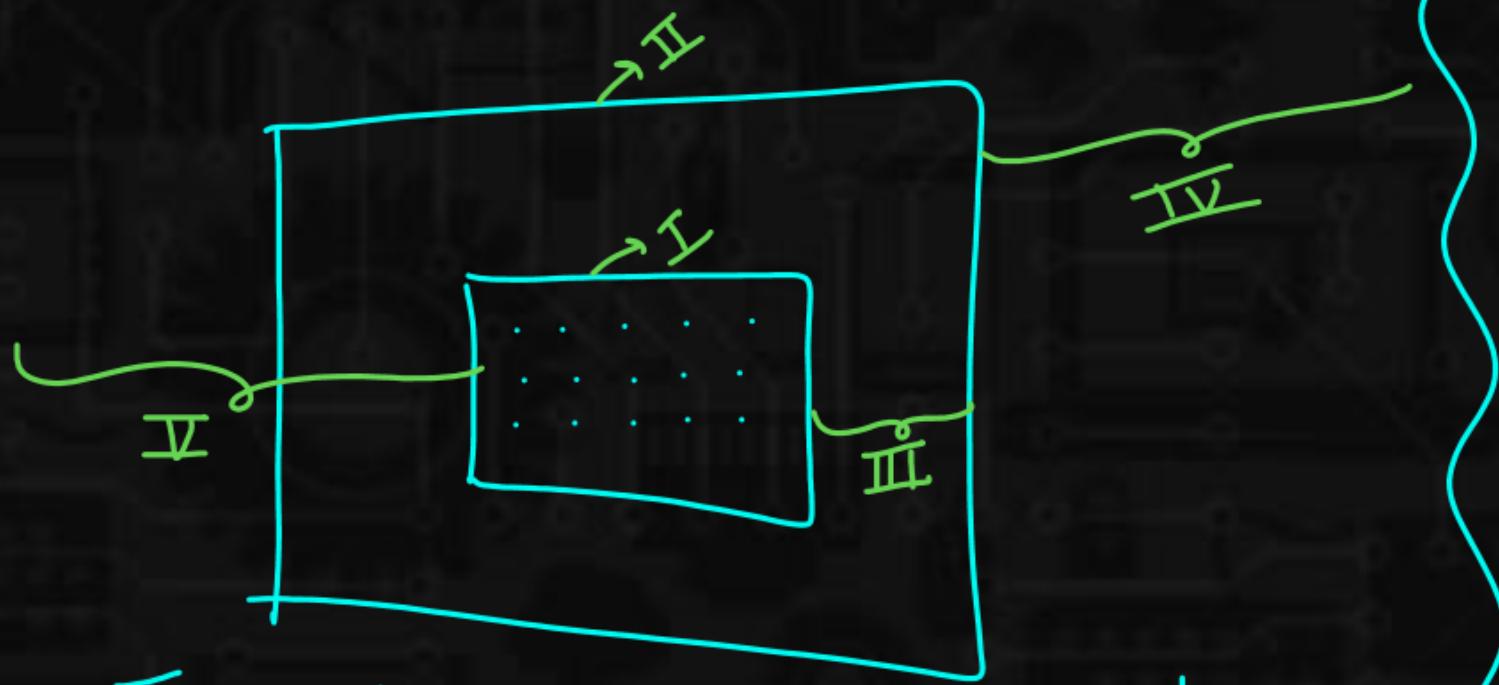
VI)



Decrement in unary

$\delta$	$a$	$B$
$\rightarrow q_0$	$(q_0, a, R)$	$(q_1, B, L)$
$q_1$	$(\text{Halt}, B, -)$	-

$f(x) = x - 1 ; x > 0$



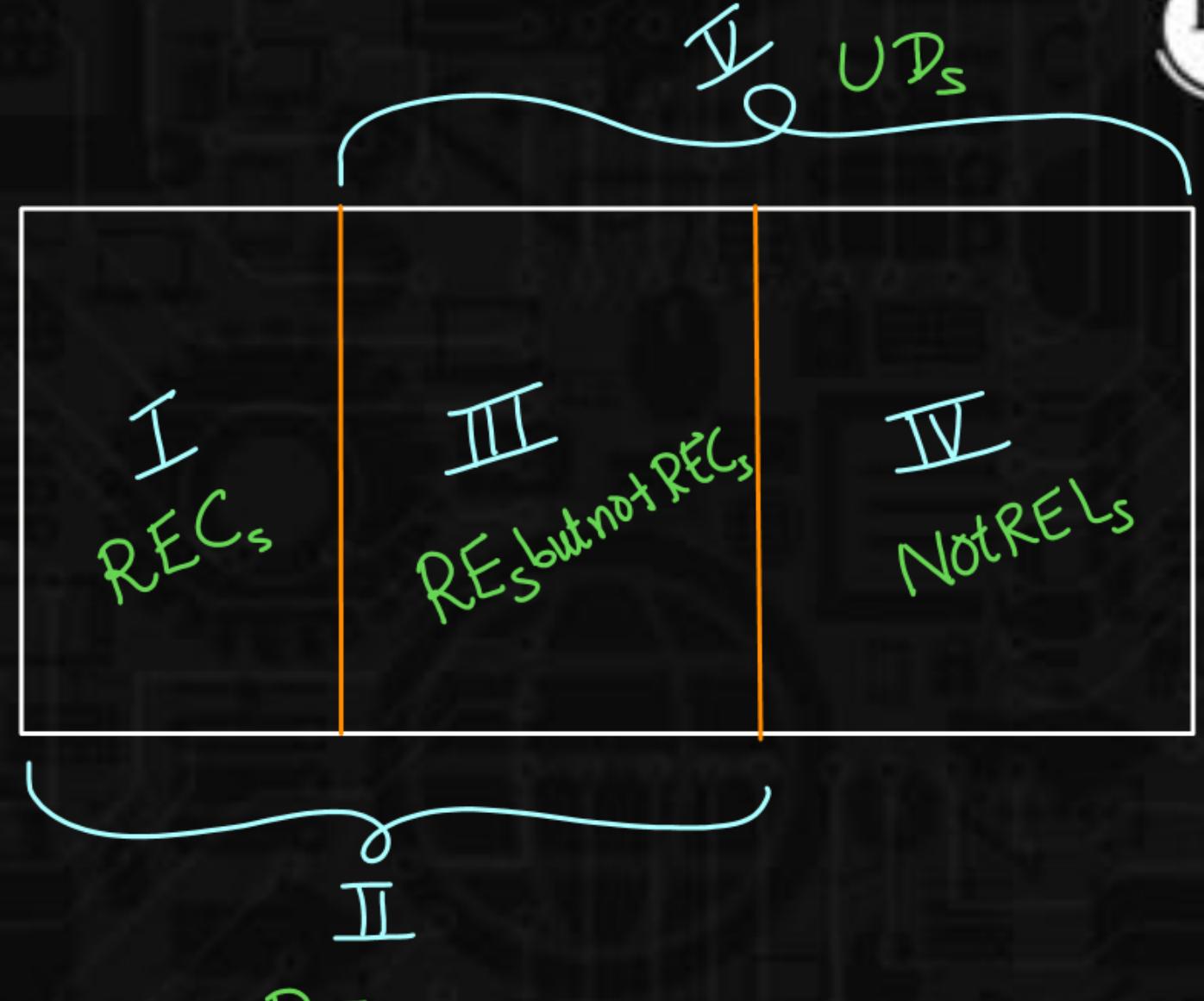
V: Set of all non recursive languages

I: Set of all recursive languages

II:  $I \cup III$  = Set of all REL<sub>S</sub>

III:  $II - I$  = Set of all REL<sub>S</sub> but not recursive

IV: Set of all not REL<sub>S</sub>



$a^*$  $(a+b)^*$  $a^* b^*$  $a^n b^n$  $a^n b^n c^n$  $WW^R | w \in \{a, b\}^*$  $w \# w |$ 

CSVS

 $a^n \# a^n \# a^n$  $a^{n!}$  $a^{2^n}$ prime  
 $a$  $a^n$ 

CSLS

# Language

→ Set of all strings over  $\Sigma$

## Examples:

- Set of symbols
  - Set of grammars
  - Set of languages
  - Set of DFAs
  - Set of TMs
  - Set of objects
- Set of Regs
  - Set of DCFLs
  - Set of finite languages

$$L = \{ \square, \square, \square, \dots \}$$

Symbol / String / Language / Regular / Grammar / DFA | NFA | PDA | NPDA | TM | - - -

TM  $\equiv$  Program

HTM  $\equiv$  Halting program  $\equiv$  Algorithm

# "closure properties" for REGS

[Not closed:

$\subseteq$ ,  $f$ ,  $h$ ,  $f_{\text{finite}}$ ,  $\text{Inf}(v, \cap, \cup, \subseteq, f)$

P  
W

- ① Union
- ② Intersection
- ③ Complement

- ④ Difference

- ⑤ Concatenation

- ⑥ Reversal

- ⑦ Kleene Star

- ⑧ Kleene plus

- ⑨ SubSet

- ⑩ Prefix

- ⑪ Suffix

- ⑫ SubString

- ⑬  $f(L)$

- ⑭  $h(L)$

- ⑮  $\epsilon$ -free  $h(L)$

- ⑯  $h^{-1}(L)$

- ⑰
- ⑱
- ⑲
- ⑳
- ㉑

- ㉒ Finite Substitution

- ㉓
- ㉔
- ㉕
- ㉖
- ㉗
- ㉘
- ㉙
- ㉚

- Infinite ( $\cup, \cap, \neg, \cdot, \subseteq, f$ )

# "closure properties" for RELs

Not closed:

$\bar{L}$ ,  $L_1 - L_2$ ,  $\subseteq$ , FinDiff, Inv( $\cup, \cap, \neg, \subseteq, f$ )

P  
W

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⑯  $h^{-1}(L)$

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FinItX ( $\cup, \cap, \neg, \subseteq, f$ )

Infinite ( $\cup, \cap, \neg, \subseteq, f$ )

$\text{REC}_s \times \boxed{f}_h \subseteq \text{Inf} \dots$   
fh  
Finite substitution

$\text{REL}_s \times i_{l_1, l_2} \subseteq \text{Inf} \dots$   
c  
FinDiff

## ① Union

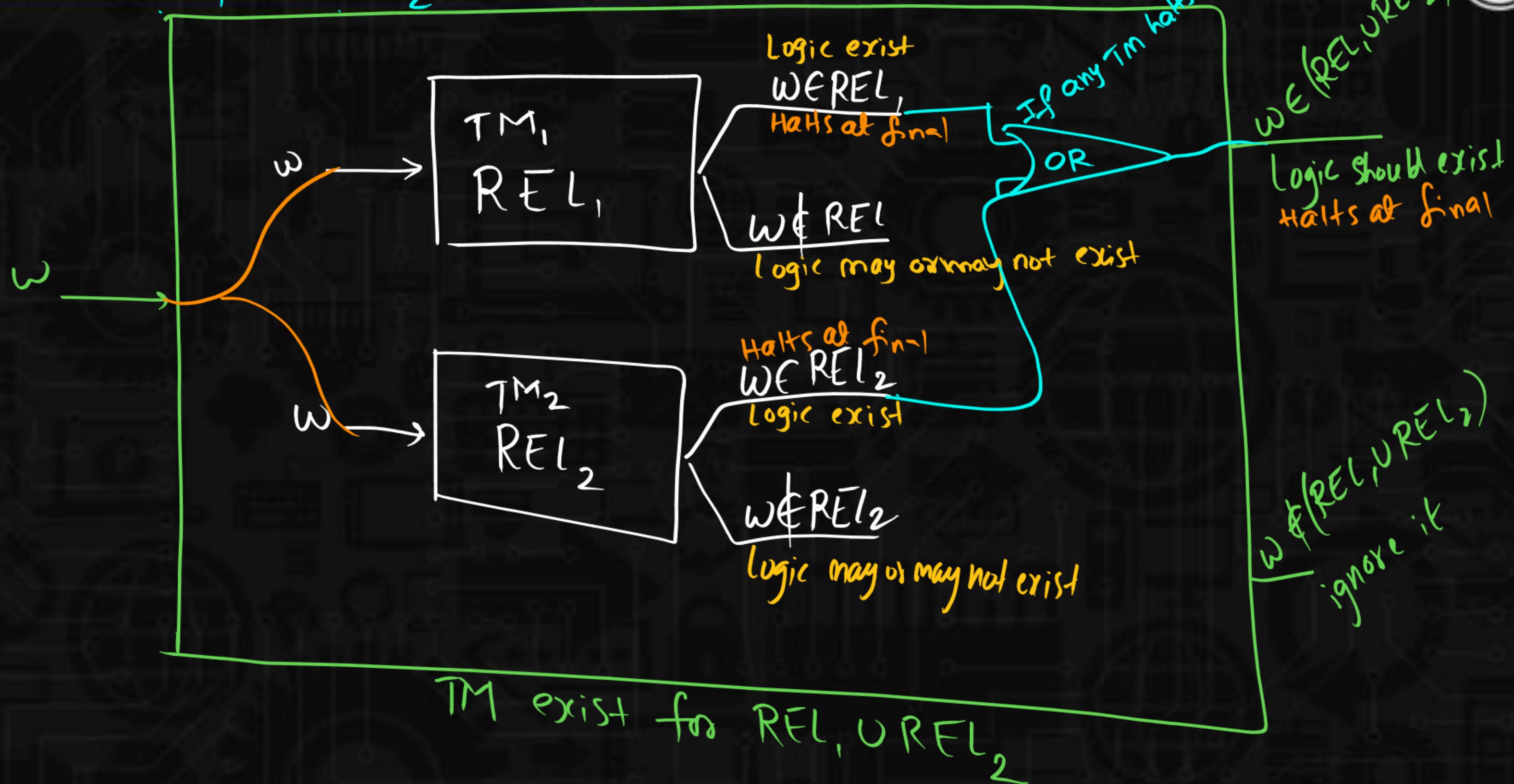
→ closed for RELs  
→ closed for Recursive languages (REC<sub>s</sub>)

$$\text{REL}_1 \cup \text{REL}_2 \Rightarrow \text{REL}$$

$$\text{REC}_1 \cup \text{REC}_2 \Rightarrow \text{REC}$$

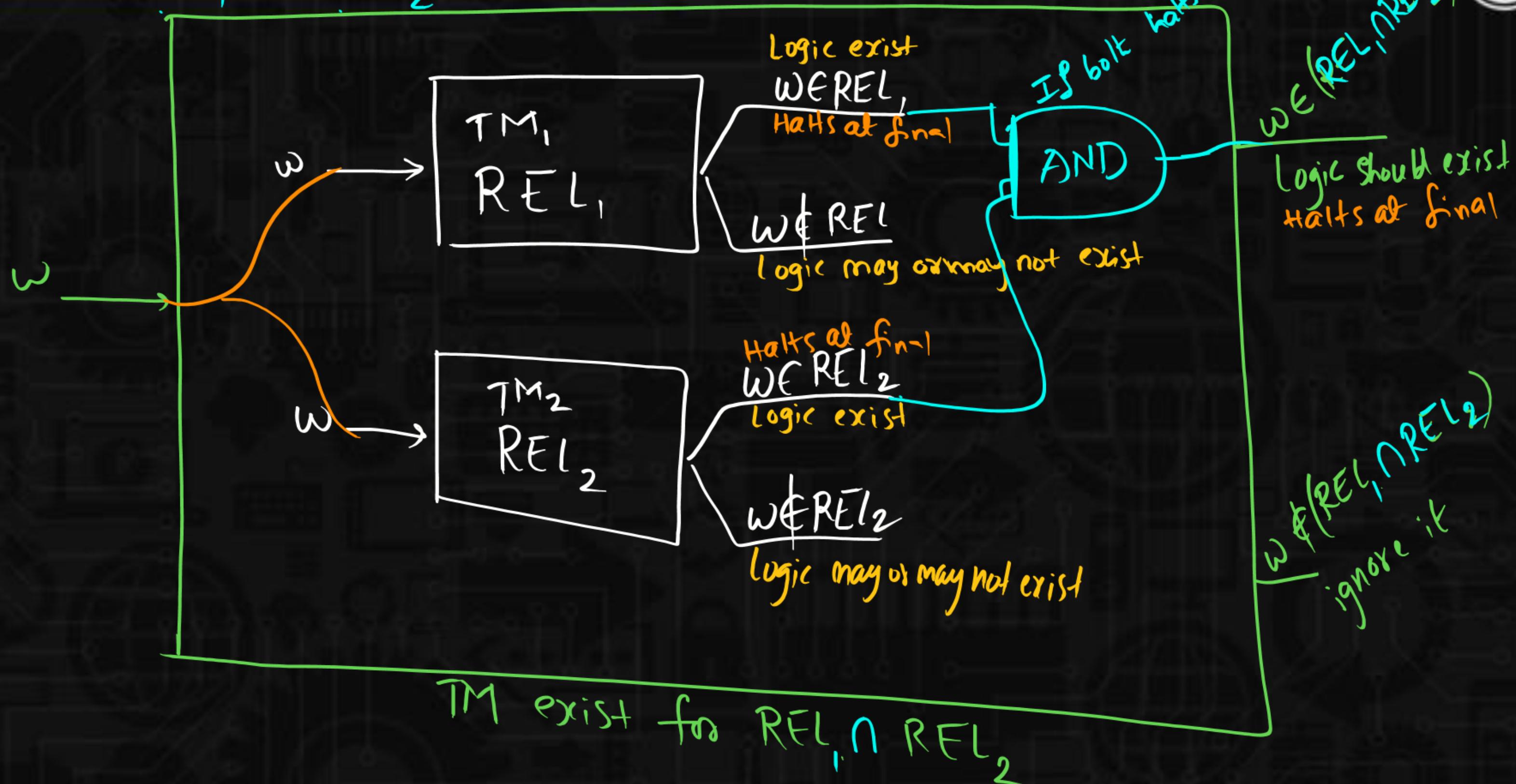
$\text{REL}, \text{UREL}_2 \Rightarrow \text{REL}$

P  
W



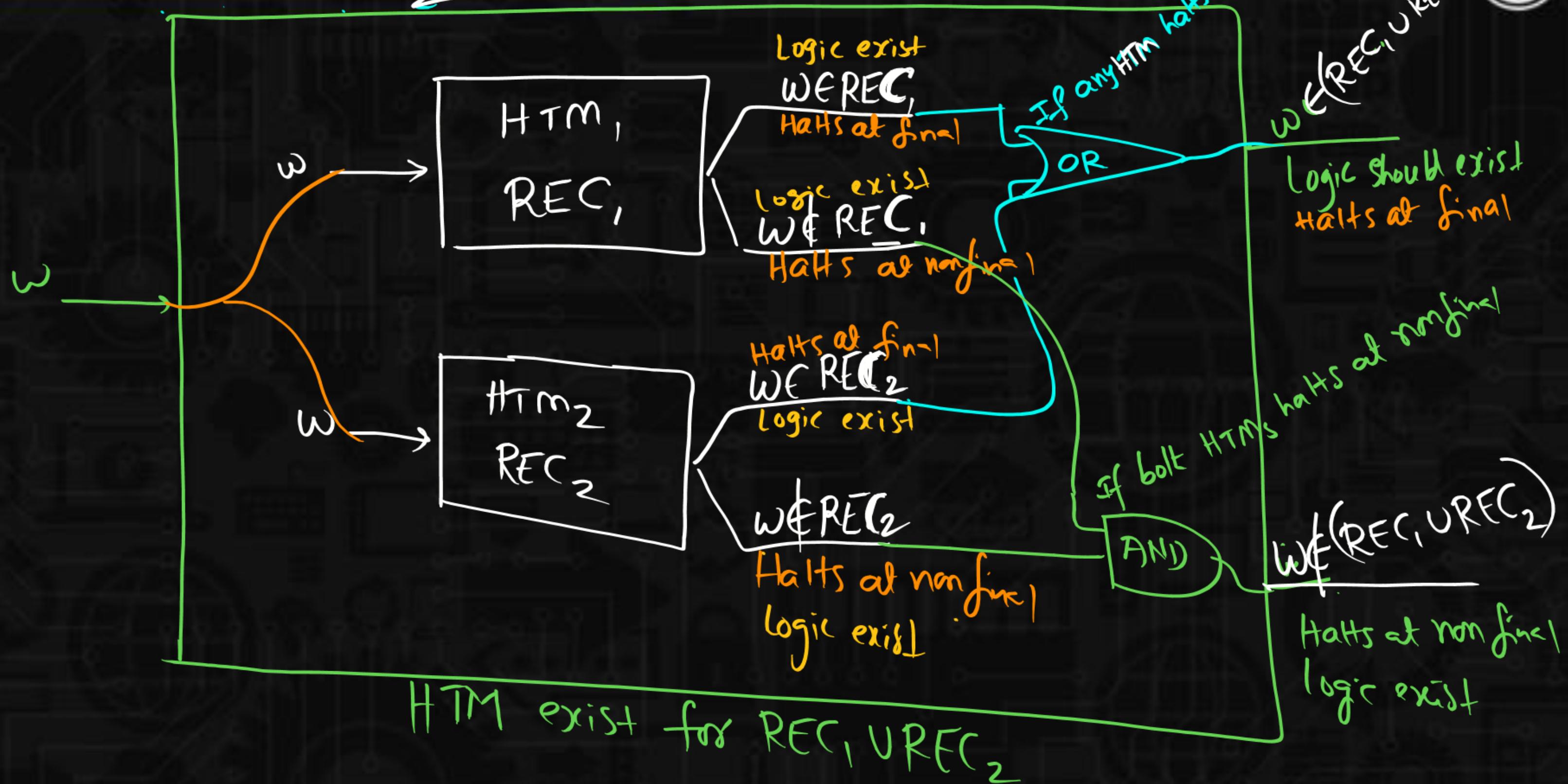
$\text{REL}_1 \cap \text{REL}_2 \Rightarrow \text{REL}$

P  
W



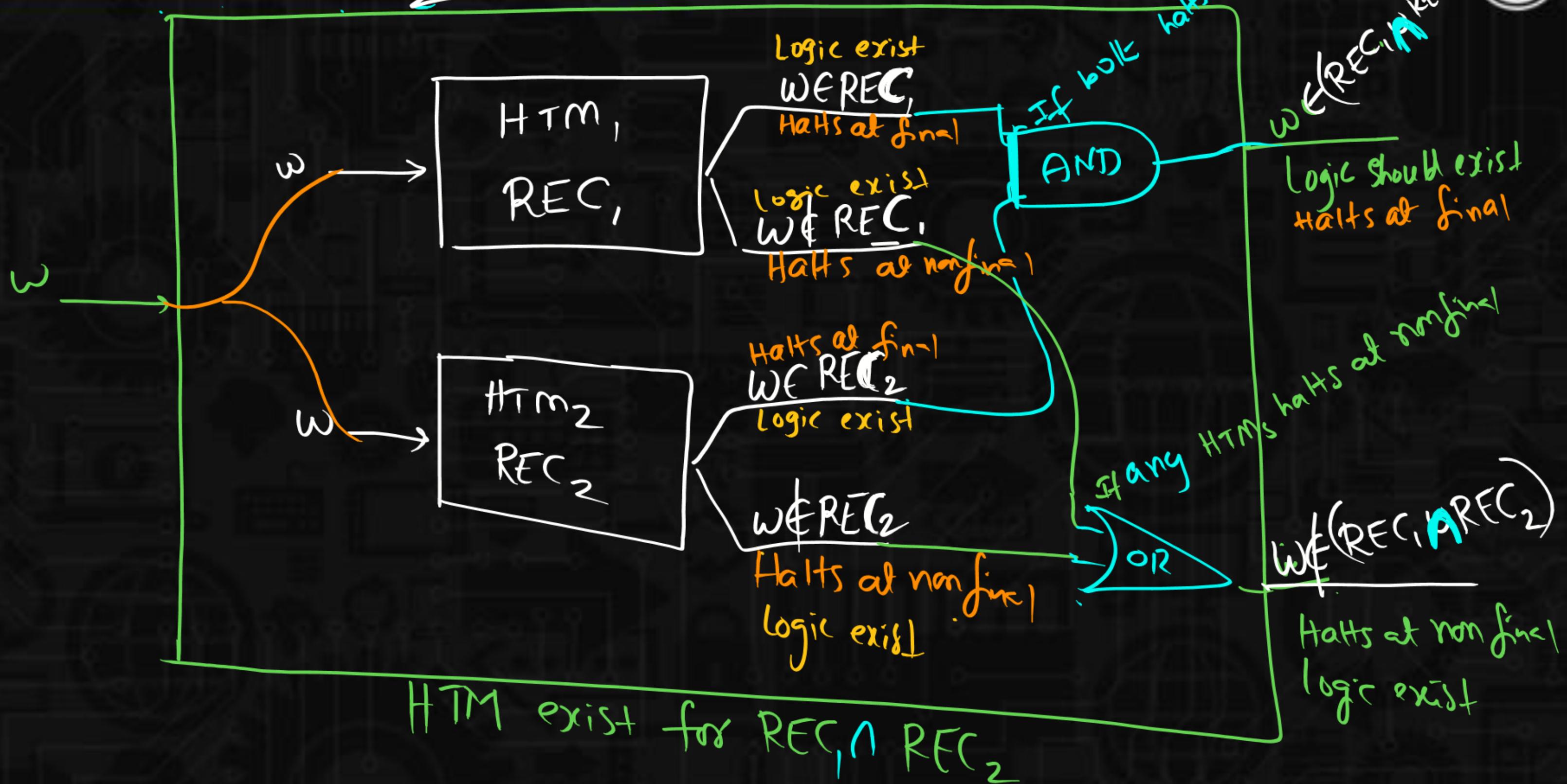
$\text{REC}_1 \cup \text{REC}_2 \Rightarrow \text{REC}$

P  
W



$\text{REC}_1 \cap \text{REC}_2 \Rightarrow \text{REC}$

P  
W



① Union

↳ closed for RELs & REC<sub>s</sub>

② Intersection

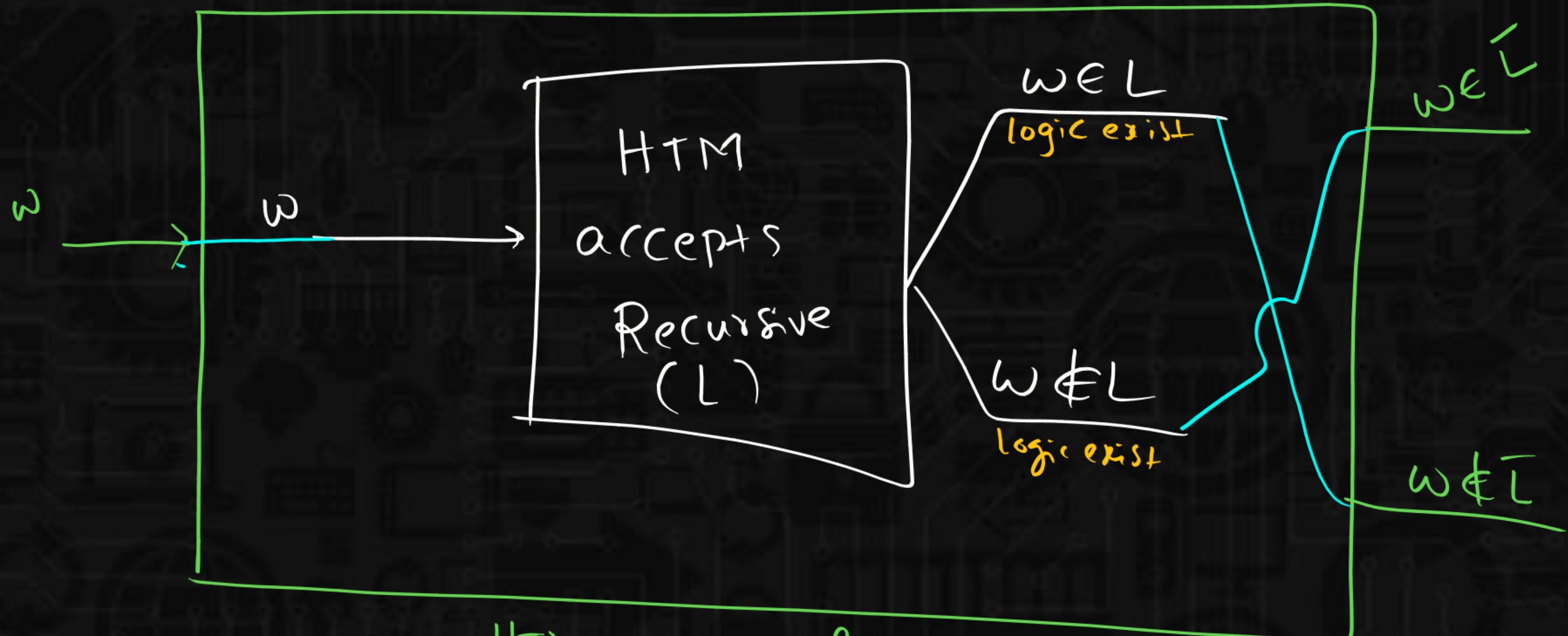


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③ Complement

→ closed for REC<sub>s</sub>  
→ Not closed for RELs

# Complement of Recursive language is Recursive

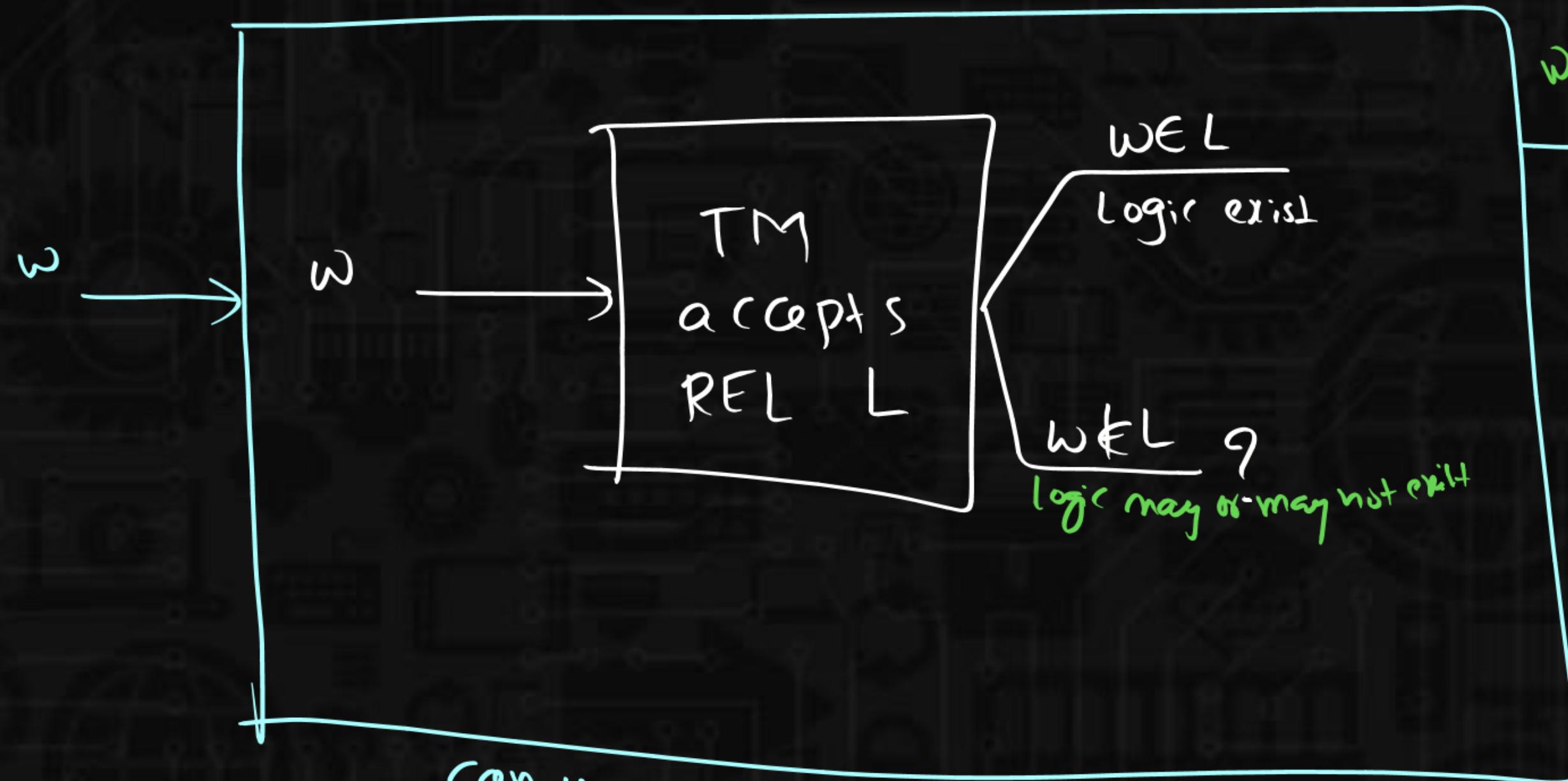
P  
W



HTM exist for complement of L

Complement of REL is may or may not be REL  
(either Recursive or Not REL)

P  
W



Can you prove TM exist for  $\bar{L}$  ?

We don't know  
we L logic may or may not exist  
Can you show logic for valid strings in L

REL  $\Rightarrow$  either "Recursive" OR "Not REL"

P  
W

If  $L$  is Recursive then  $\bar{L}$  is Recursive

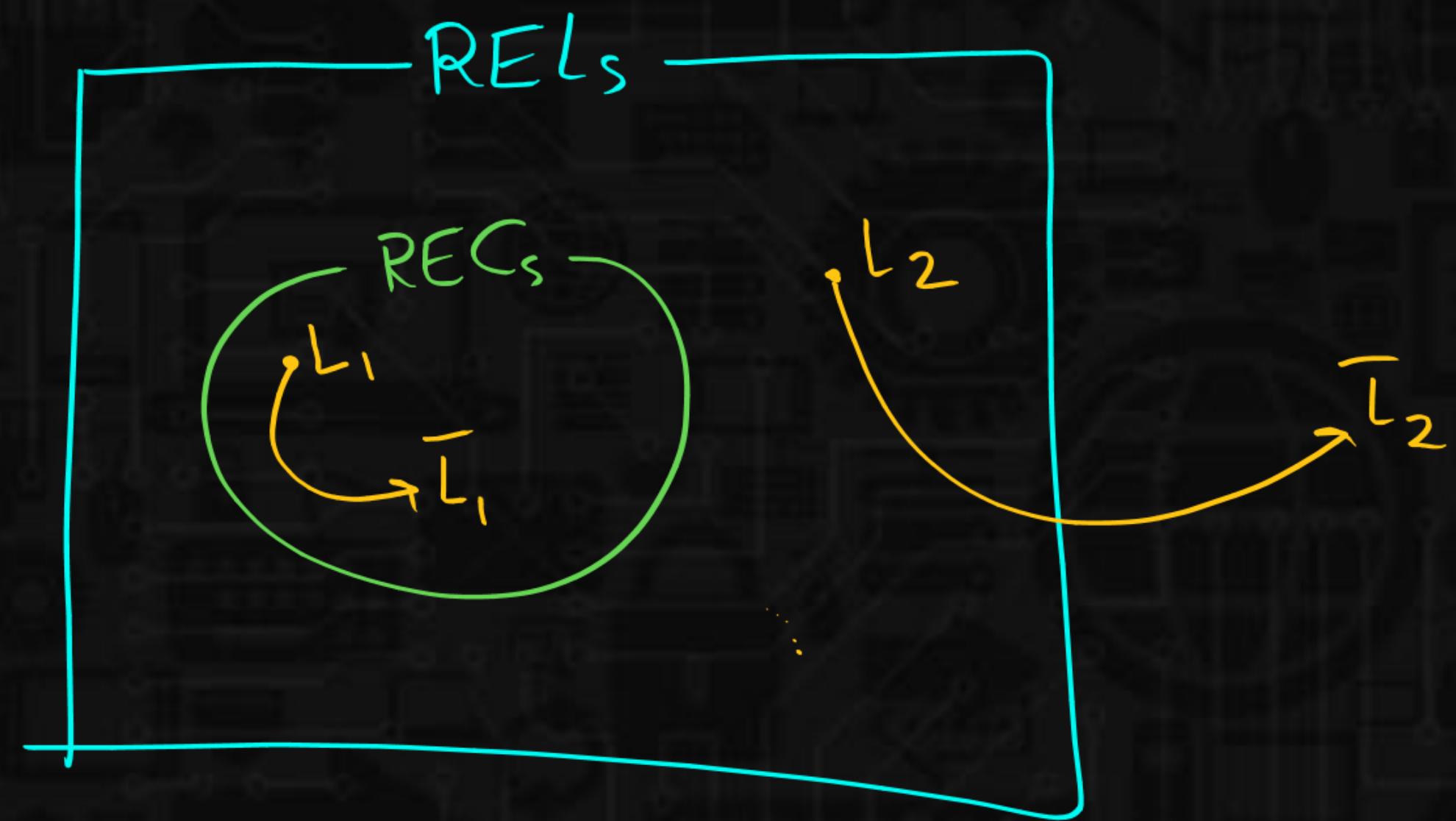
If  $L$  is "REL but not recursive"

then  $\bar{L}$  is Not REL

I)  $\overline{\text{REL}}$   $\Rightarrow$  May or may not be REL

II)  $\overline{\text{REL}}$   $\Rightarrow$  either "Recursive" OR "Not REL"

III)  $\overline{\text{REL}}$   $\Rightarrow$  Never be "REL but not Recursive"



- ①  $\text{REL}, \text{UREL}_2 \Rightarrow \text{REL}$
- ②  $\text{REC}, \text{UREC}_2 \Rightarrow \text{REC}$
- ③  $\text{REC} \cup \text{REL} \Rightarrow \text{REL}$
- ④  $\text{REC} \cup \text{Reg} \Rightarrow \text{REC}$
- ⑤  $\text{REC} \cup \text{CFL} \Rightarrow \text{REC}$
- ⑥  $\text{REC} \cup \text{CSL} \Rightarrow \text{REC}$
- ⑦  $\text{REL} \cup \text{Reg} \Rightarrow \text{REL}$
- ⑧  $\text{REL} \cup \text{CFL} \Rightarrow \text{REL}$
- ⑨  $\text{REL} \cup \text{CSL} \Rightarrow \text{REL}$

- \* \* \*
- ⑩  $\text{REL}, \neg \text{REL}_2 \Rightarrow \text{REL}$
  - ⑪  $\text{REC}, \neg \text{REC}_2 \Rightarrow \text{REC}$
  - ⑫  $\text{REC} \cap \text{REL} \Rightarrow \text{REL}$
  - ⑬  $\text{REC} \cap \text{Reg} \Rightarrow \text{REC}$
  - ⑭  $\text{REC} \cap \text{CFL} \Rightarrow \text{REC}$
  - ⑮  $\text{REC} \cap \text{CSL} \Rightarrow \text{REC}$
  - ⑯  $\text{REL} \cap \text{Reg} \Rightarrow \text{REL}$
  - ⑰  $\text{REL} \cap \text{CFL} \Rightarrow \text{REL}$
  - ⑱  $\text{REL} \cap \text{CSL} \Rightarrow \text{REL}$

\*\*\*

$$\boxed{DCFL \cap CFL \Rightarrow CSL}$$

$$CFL \cap CSL \Rightarrow CSL$$

$$DCFL \cap DCFL \Rightarrow CSL$$

$$CFL \cap CFL \Rightarrow CSL$$

Rec  $\Rightarrow$  Rec

REL  $\Leftrightarrow$  either Recursive OR Not REL

REL but not Rec  $\Rightarrow$  Not REL

→ closure properties ✓

Complement  
Important

closed

