

# CS & IT ENGINEERING

Theory of Computation

Finite Automata:

NFA / DFA-4

Lecture No. 8



By- DEVA Sir



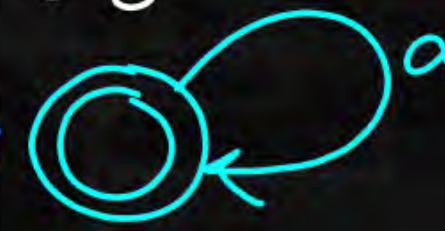
# TOPICS TO BE COVERED

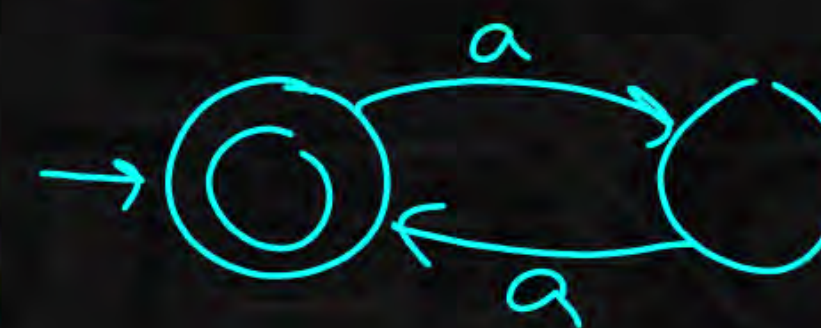
## 01 Remaining DFA Models

I, II, III ✓

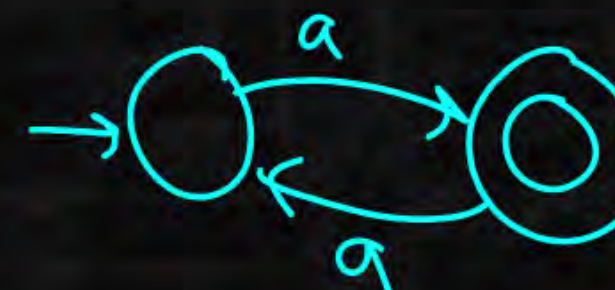
IV  
↓

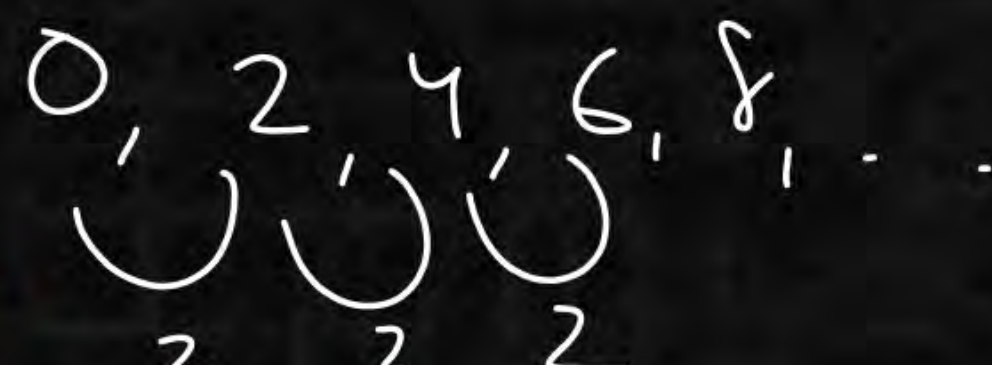


(27)  $a^* = \{a^n \mid n \geq 0\} \rightarrow$  

$L =$  (28)  $(aa)^* = \{a^{2n} \mid n \geq 0\}$   
 $= \{\epsilon, a^2, a^4, a^6, \dots\}$  

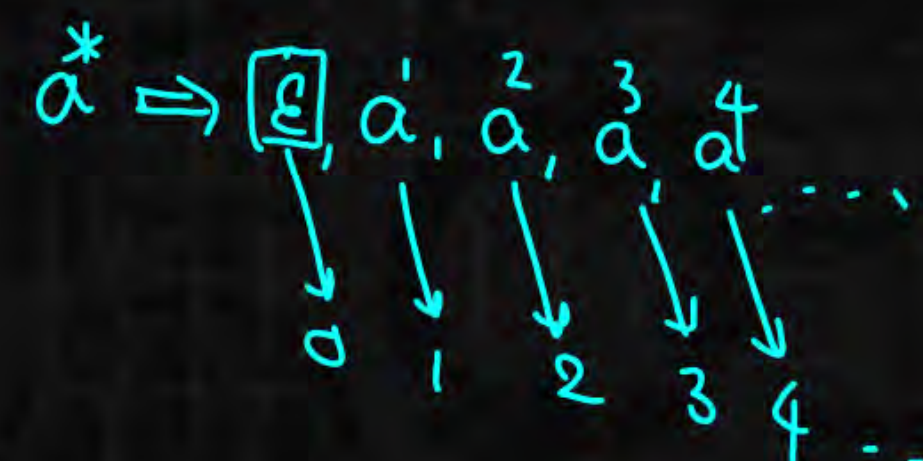
A.P.  
 $0, 1, 2, 3, \dots$

$\bar{L} =$  (29)  $a(aa)^* = \{a^{2n+1} \mid n \geq 0\}$   
 $= \{a, a^3, a^5, a^7, \dots\}$  

$0, 2, 4, 6, 8, \dots$   


(30)  $\{a^{3n+2} \mid n \geq 0\}$

(31)  $\{a^{3n+5} \mid n \geq 0\}$

$a^* \Rightarrow \boxed{\epsilon}, a^1, a^2, a^3, a^4, \dots$   


$a(aa)^*$

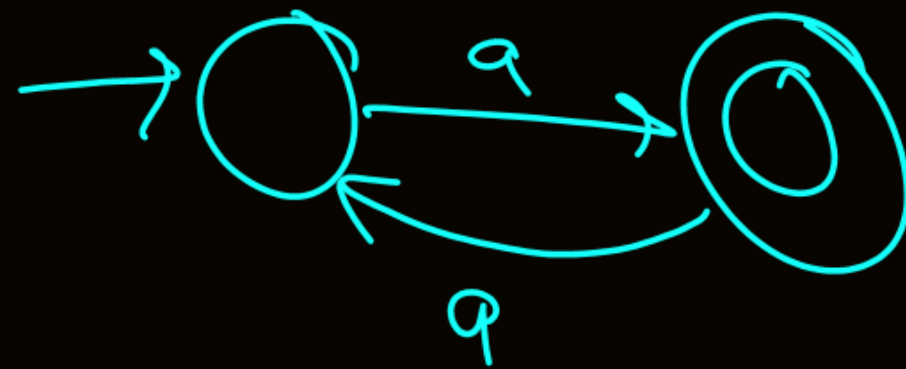


$$a(aa)^0 = a$$

$$a(aa)^1 = a^3$$

$$a(aa)^2 = a(aa)^1(aa)^1 \\ = a^5$$

$$L = \{a, a^3, a^5, a^7, \dots\}$$



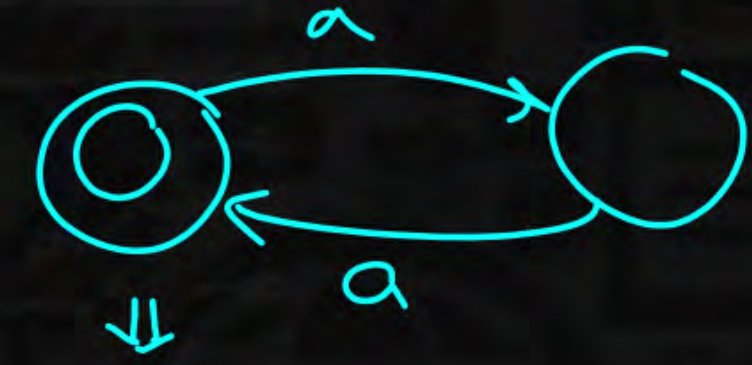
$$\underbrace{(aa)^*}_{\text{even } a's} + \underbrace{a(aa)^*}_{\text{odd } a's} = \underbrace{a^*}_{\text{All } a's}$$



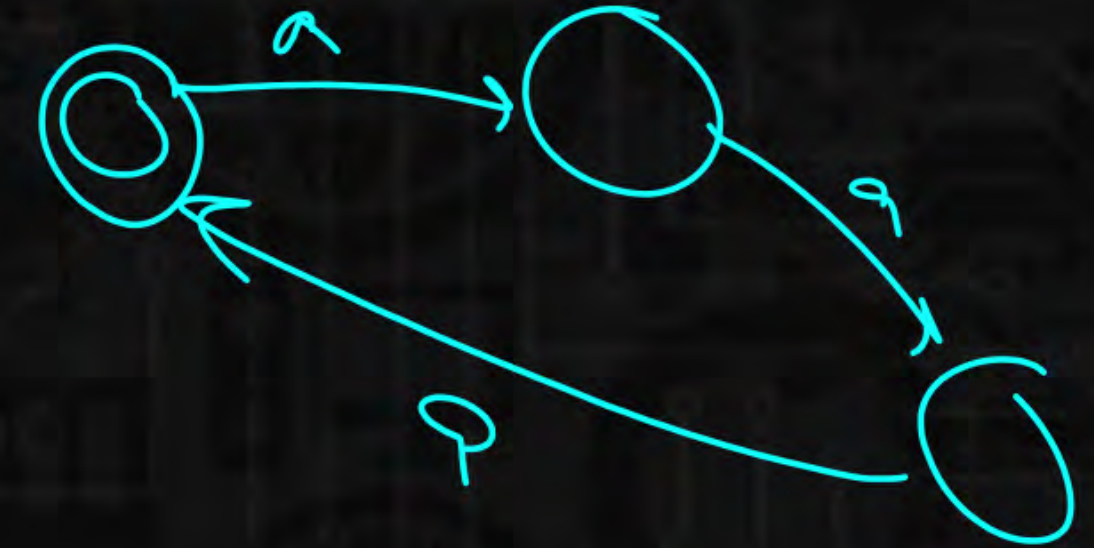
1



2



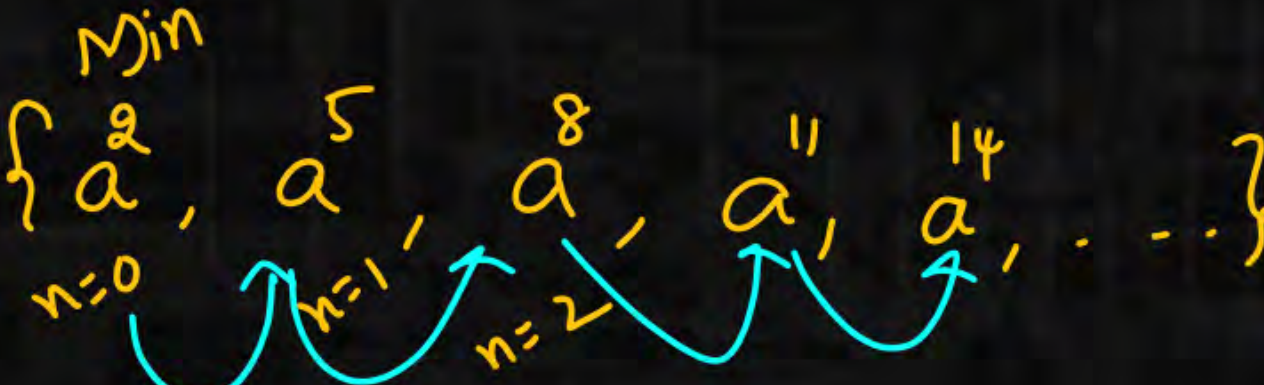
3

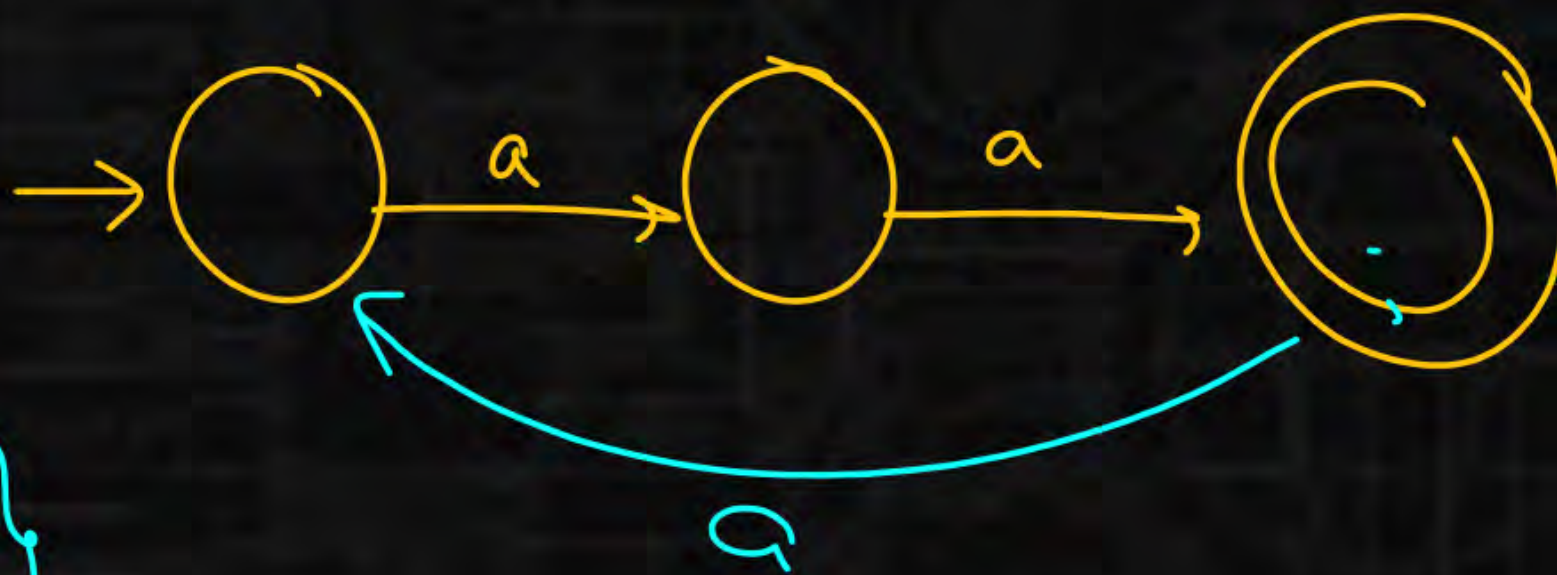


(30)  $\{a^{3n+2} \mid n \geq 0\} = \{a^2, a^5, a^8, a^{11}, a^{14}, \dots\}$

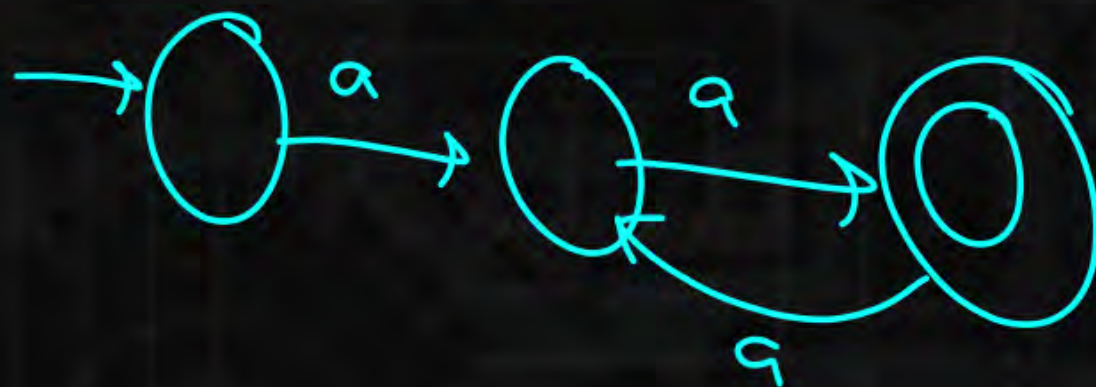
Min

$n=0$   $n=1$   $n=2$





(32)  $\{a^{2n+2} \mid n \geq 0\}$

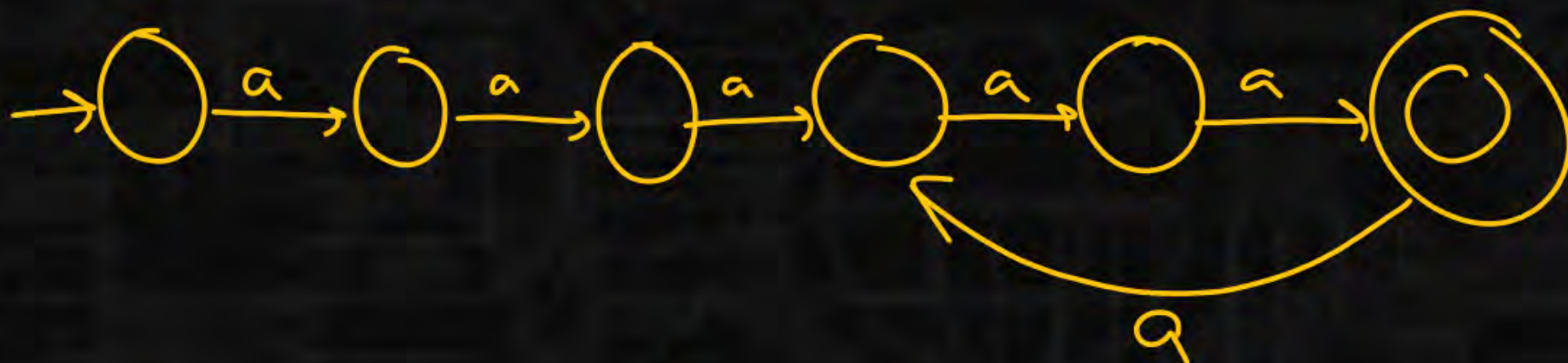




$$(31) \{a^{3n+5} \mid n \geq 0\} = \{a^5, a^8, a^{11}, a^{14}, \dots\}$$

$\underbrace{a^5}_{3a's} \quad \underbrace{a^8}_{3a's} \quad \underbrace{a^{11}}_{3a's} \quad \underbrace{a^{14}}_{3a's} \quad \dots$

$$\begin{aligned} a^{3n+5} &= a^{3n} \cdot a^5 \\ &= (aaa)^* a^5 \end{aligned}$$



$$\begin{array}{cccccc} \epsilon & a & a^2 & a^3 & a^4 & a^5 \\ & & & a^6 & a^7 & a^8 \\ & & & a^9 & a^{10} & a^{11} \\ & & & a^{3n+3} & a^{3n+4} & a^{3n+5} \end{array} \quad n \geq 0$$

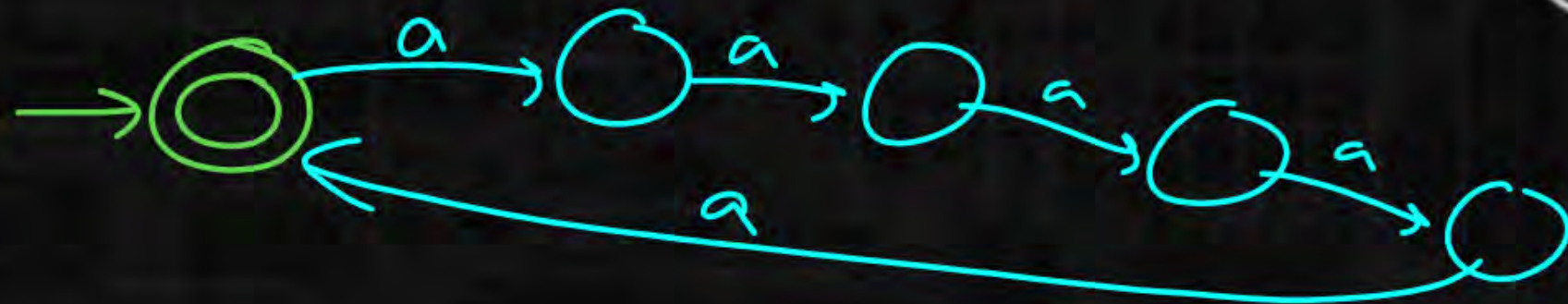


$$s_n = s_{n+0}$$

$k_1 + k_2$   
 $a$   
 $k_1 > k_2$

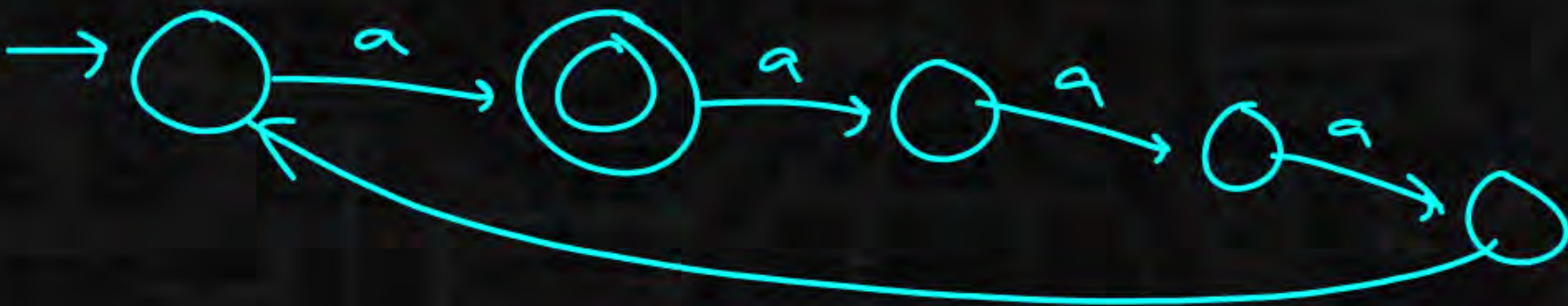
(32)  $\{a^{s_n} \mid n \geq 0\} \Rightarrow$

5 states



(33)  $\{a^{s_{n+1}} \mid n \geq 0\} \Rightarrow$

5 states



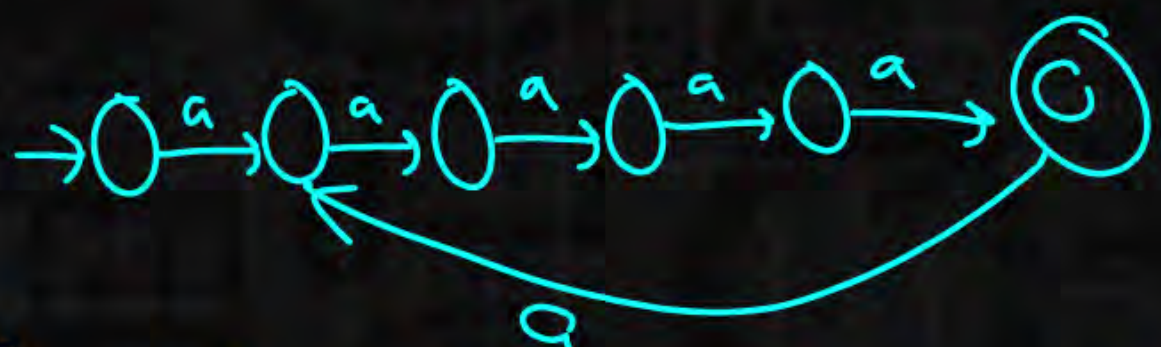
(34)  $\{a^{s_{n+4}} \mid n \geq 0\}$

5 states

$k_1 + k_2$   
 $a$   
 $k_1 \leq k_2$

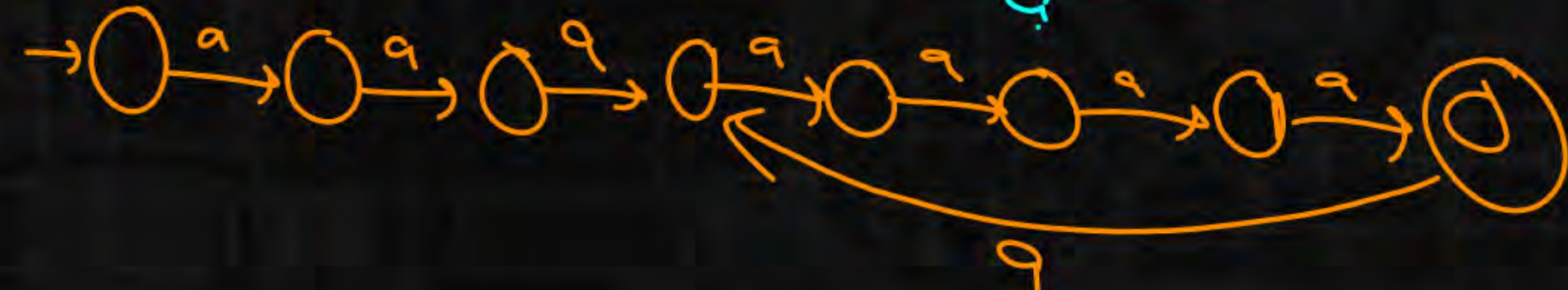
(35)  $\{a^{s_{n+5}} \mid n \geq 0\} = \{a^5, a^{10}, a^{15}, \dots\}$

6 states



(36)  $\{a^{s_{n+7}} \mid n \geq 0\}$

8 states





$$(37) \quad \{a^{2n} \mid n \geq 1\} = \{a^2, a^4, \dots\} = \{a^{2n+2} \mid n \geq 0\} \Rightarrow 3 \text{ states}$$

$$(38) \quad \{a^{3n+2} \mid n \geq 1\} = \{a^{3n+5} \mid n \geq 0\} \Rightarrow 6 \text{ states}$$

$$(39) \quad \{a^{2n+1} \mid n \geq 1\} = \{a^{2n+3} \mid n \geq 0\} \Rightarrow 4 \text{ states}$$



$$\{ a^{[K_1]n + [K_2]} \mid n \geq 0 \}$$

If  $K_1 > K_2$

$K_2$  acts as remainder



$K_1$  states in min DFA

If  $K_1 \leq K_2$

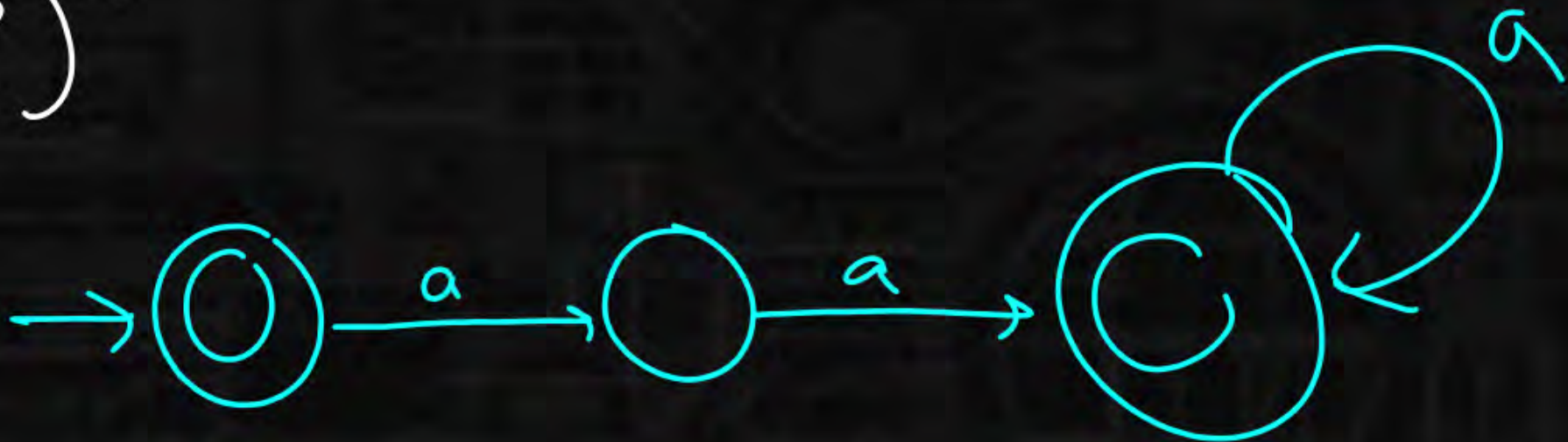
$\nexists$  ( $K_2$  will not be remainder)

$K_2 + 1$  States

(40)

$$(aa+aaa)^* = \{\epsilon, \overset{\text{A.P. begins here}}{a^2}, a^3, a^4, a^5, \dots\}$$

$$(a^2 + a^3)^*$$



$$L = \{a^n \mid n \neq 1\}$$


$$L = \{a^n \mid n=0 \text{ or } n \geq 2\}$$

$$= \epsilon + aaa^*$$

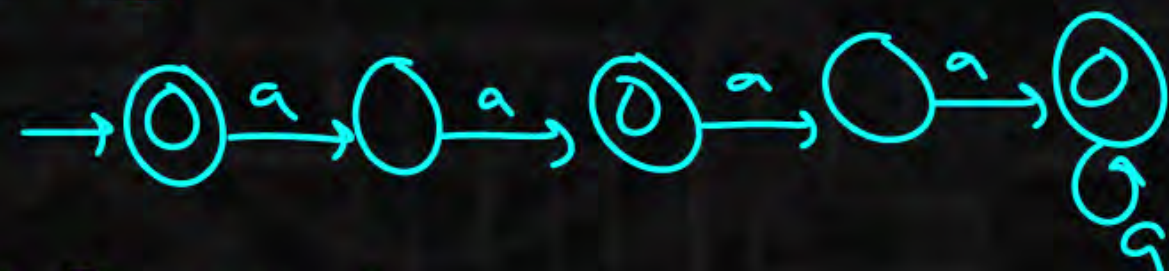
$$= \epsilon + aa^2$$



(41)  $(\underline{a} + \underline{aaa})^* = a^* \rightarrow$    $\Rightarrow 1 \text{ state}$

(42)  $(\underline{aa} + \underline{aaaa})^* = (aa)^* \Rightarrow 2 \text{ states}$     
 A.P. begins

(43)  $(aa + \underline{aaaaaa})^* = \{ \epsilon, a^2, \boxed{a^4}, a^5, a^6, a^7, a^8, a^9, \dots \} \Rightarrow 5 \text{ states}$

(44)  $(aaa + \underline{aaaaa})^*$     
  $= \{ \epsilon, a^3, a^4, \boxed{a^6}, a^7, a^8, a^9, \dots \} \Rightarrow 7 \text{ states}$



$$(aaa + aaaa)^* = (a^3 + a^4)^*$$

$$= (a^3 + a^4)^0 + (a^3 + a^4)^1 + (a^3 + a^4)^2 + \dots$$

$$\checkmark \epsilon \leftarrow ( )^0$$

$$\times a \leftarrow \times$$

$$\times a^2 \leftarrow \times$$

$$\checkmark a^3 \leftarrow (a^3 + a^4)^1$$

$$\checkmark a^4 \leftarrow (a^3 + a^4)^1$$

$$\times a^5 \leftarrow \times$$

$$\checkmark a^6 \checkmark$$

$$\checkmark a^7 \checkmark$$

$$\vdots$$

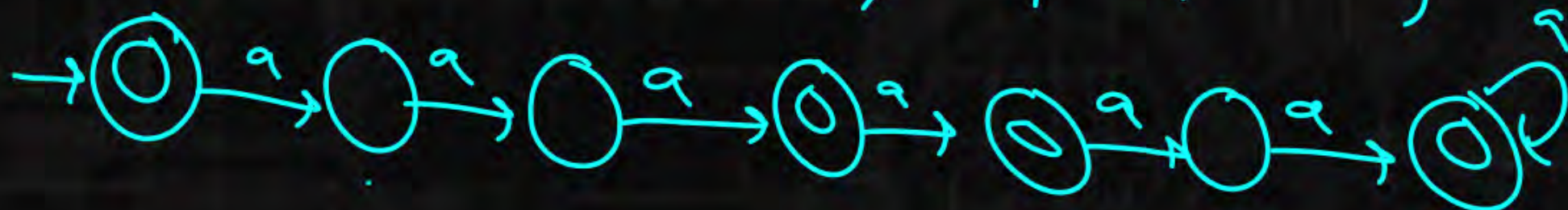
$$a^3 \cdot a^3 = a^6$$

$$a^3 \cdot a^4 = a^7$$

$$a^4 \cdot a^3 = a^7$$

$$a^4 \cdot a^4 = a^8$$

$$= \{ \epsilon, a^3, a^4, a^6, a^7, a^8, a^9, a^{10}, \dots \}$$





check given number is even or not.

0, 1, 2, 3, 4, 5, 6, 7

$n$

```
main()
{
    int n;
    scanf("%d", &n);
    if (n % 2 == 0)
        printf("even");
    else
        printf("not even");
}
```

even

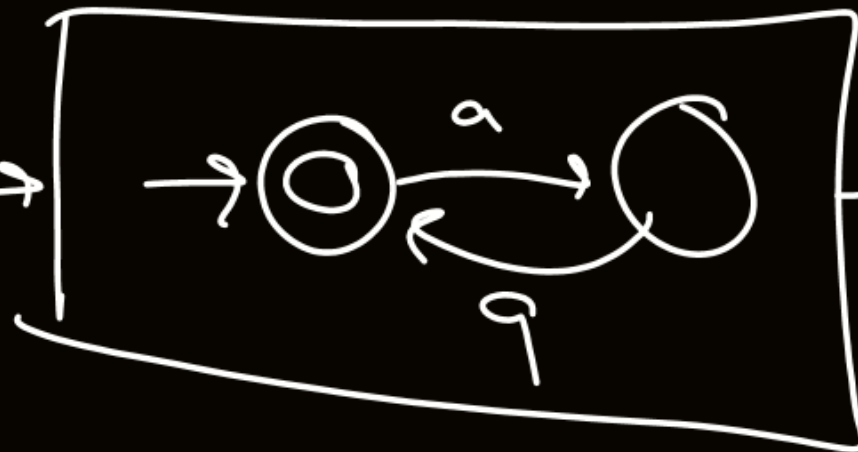
not even

$L = (aa)^*$  over  $\Sigma = \{a\}$

$\epsilon$

$a$

$aa$



$\epsilon$  is even

$a$  is not even

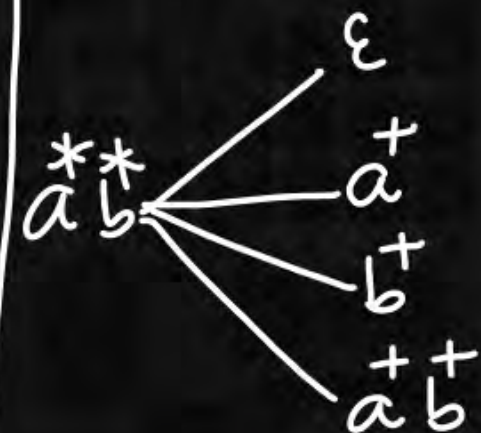
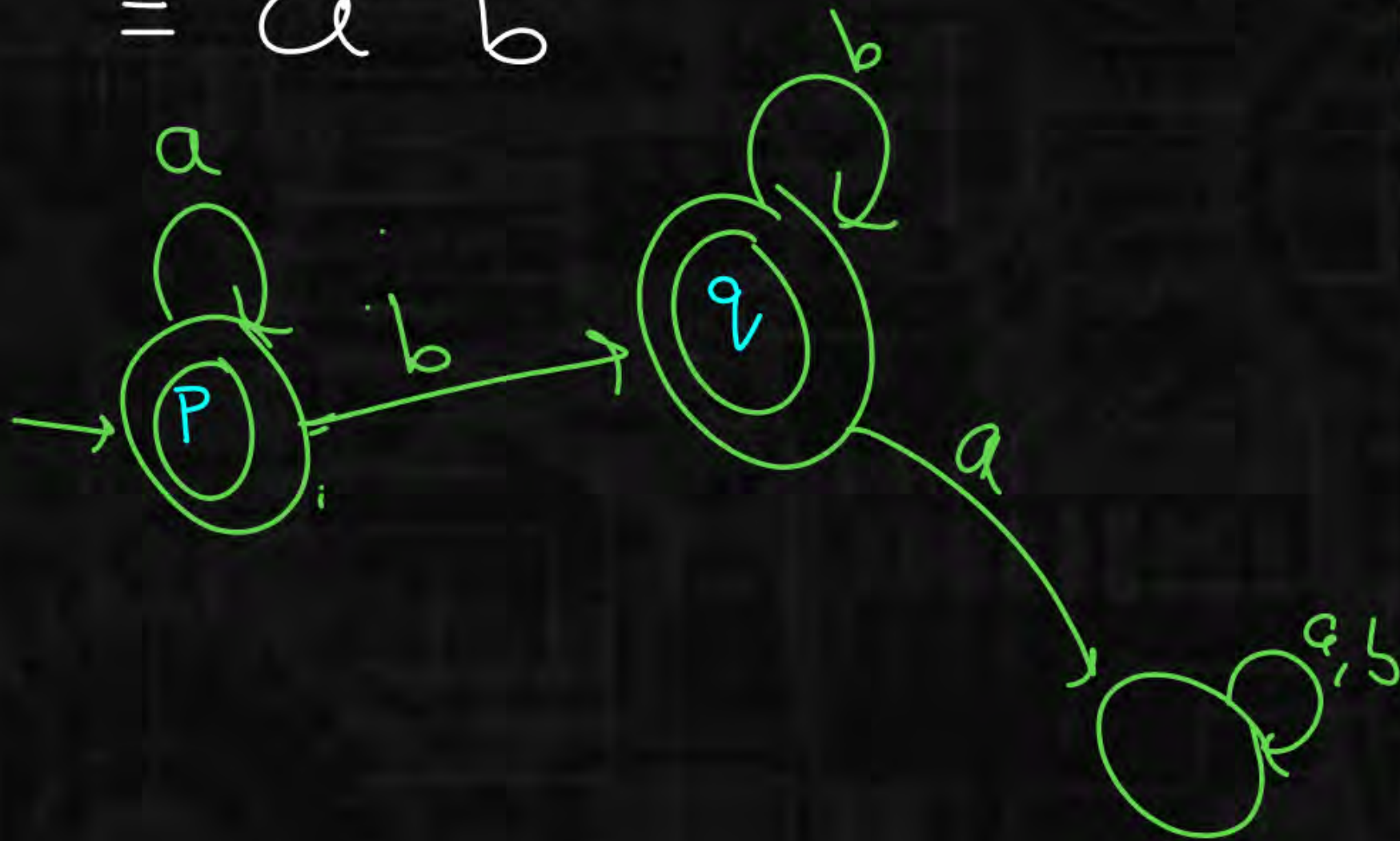
$aa$  is even

(45)

$$L = \{ \underline{a^m} b^n \mid m, n \geq 0 \}$$

Any no. of a's followed by Any no. of b's

$$= a^* b^*$$

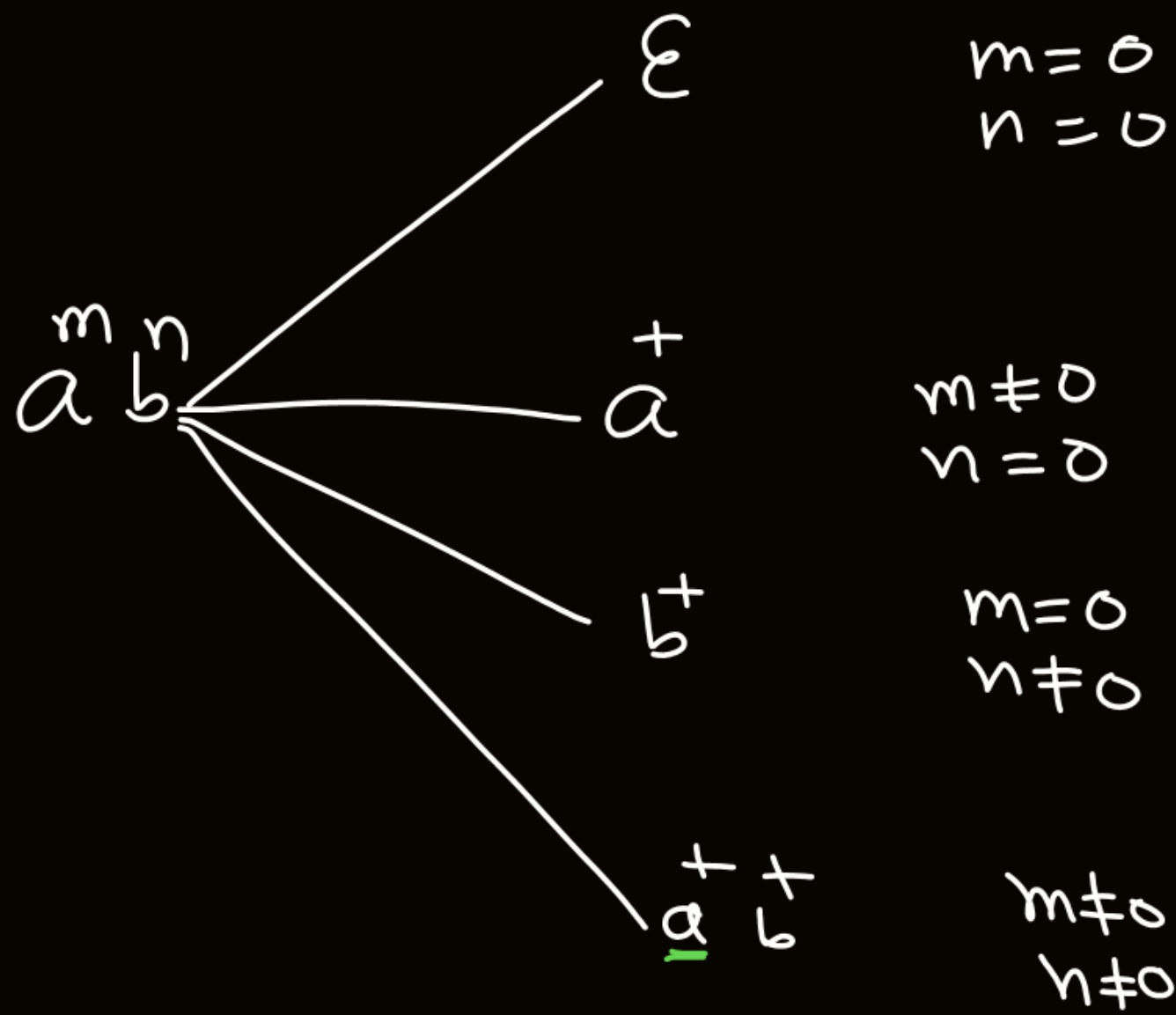


$b$  never comes before  $a$   
 $a$  never comes after  $b$



$$\begin{aligned}
 P &= a^* \\
 q &= a^* b^* \\
 \hline
 P + q &= a^* + a^* b^* \\
 &= a^* (\epsilon + b^+) \\
 &= a^* b^*
 \end{aligned}$$





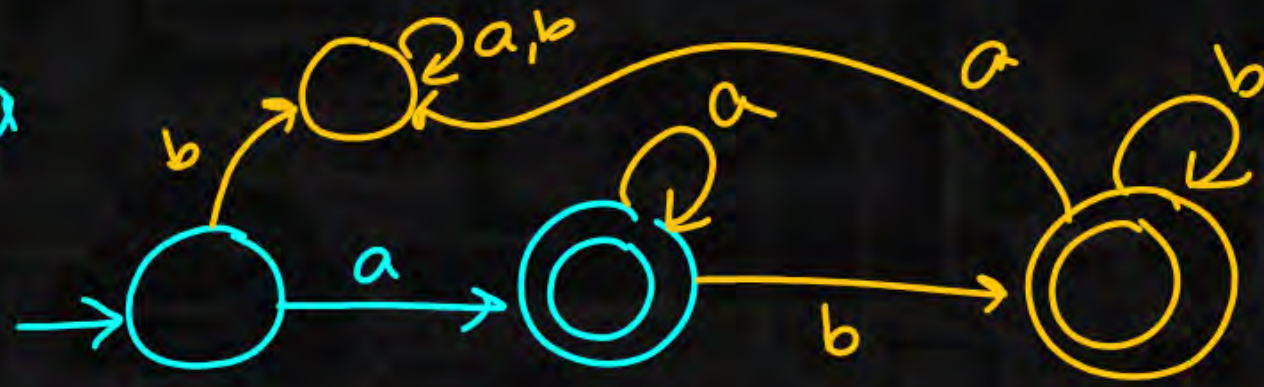
# DFA Construction



46

$a^+b^*$

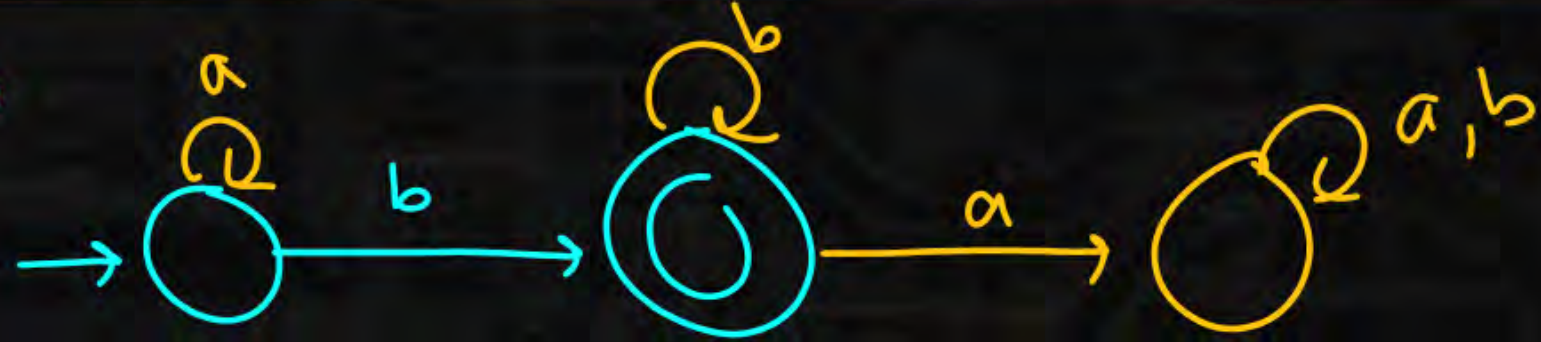
Min = a



47

$a^*b^+$   
 $= a^*bb^*$

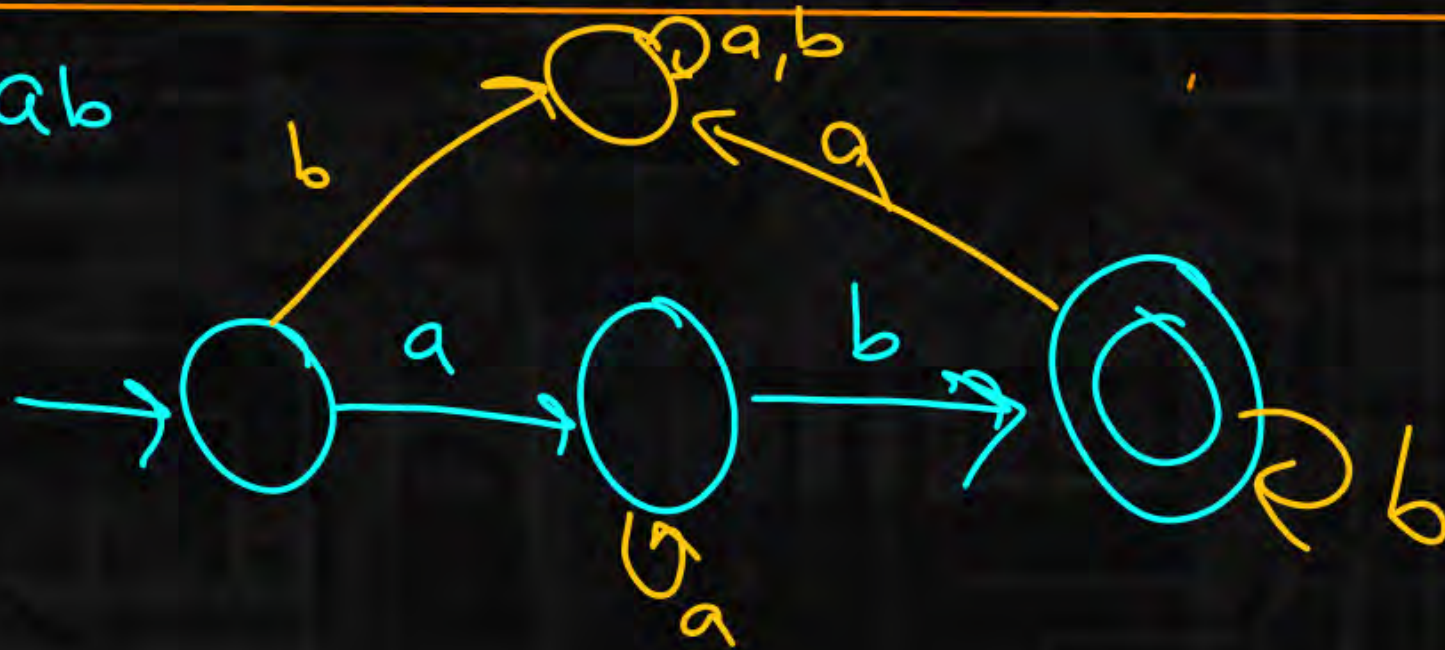
min = b



48

$a^+b^+$   
 $= aa^*bb^*$

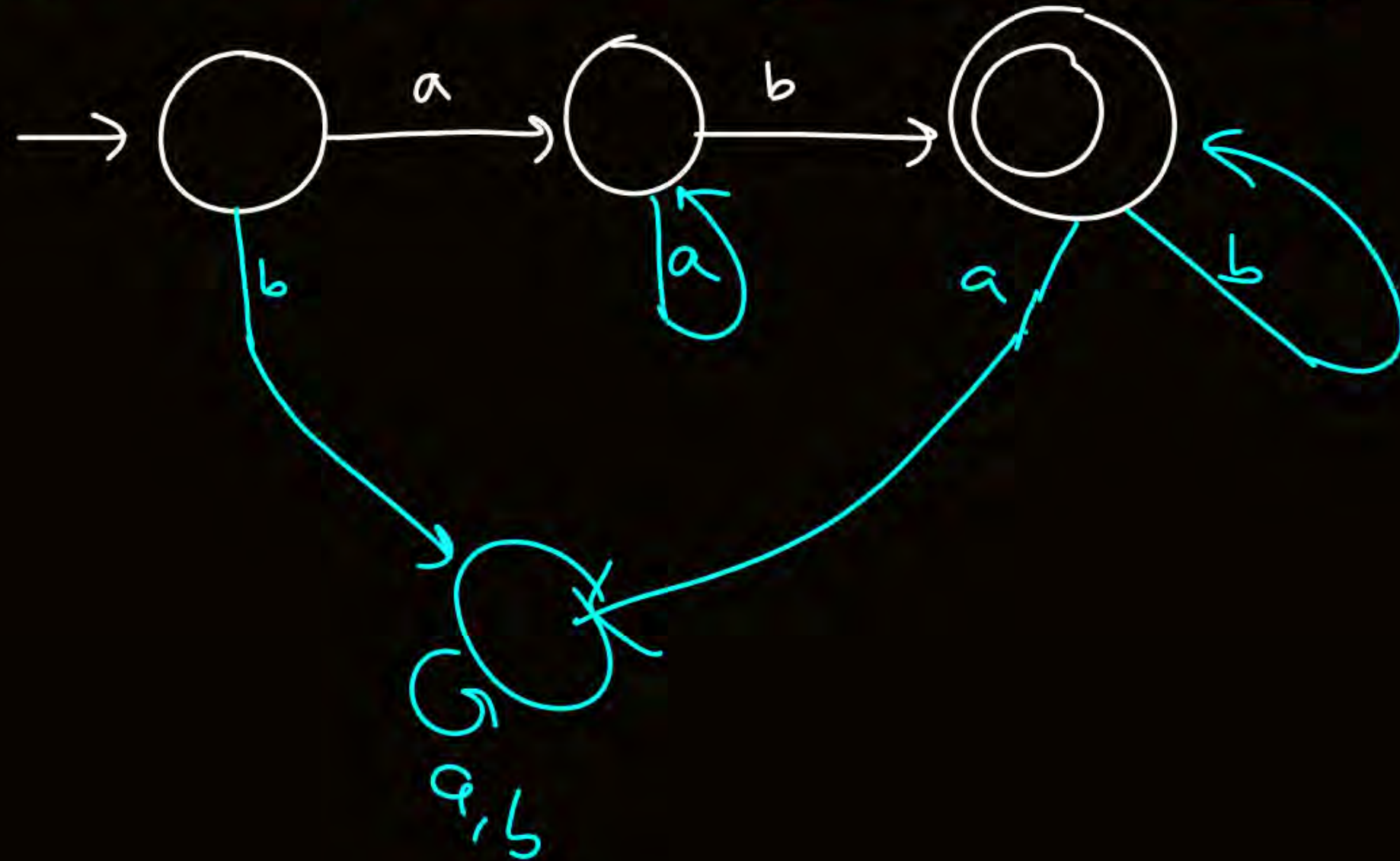
min = ab





$$a^+ b^+ = \left\{ \underbrace{ab}_{\text{min}}, aab, abb, aabb, \dots \right\}$$

Atleast one a followed atleast one b

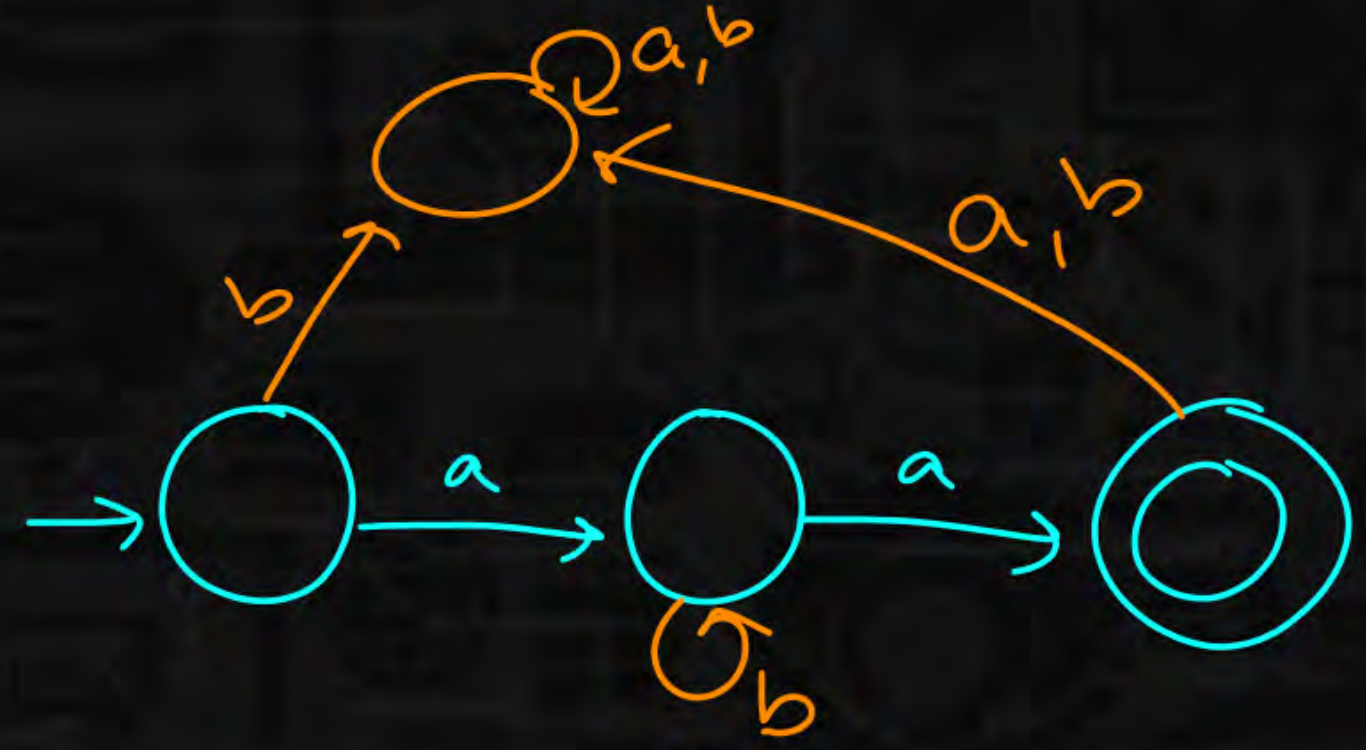




(49)

$a b^* a$

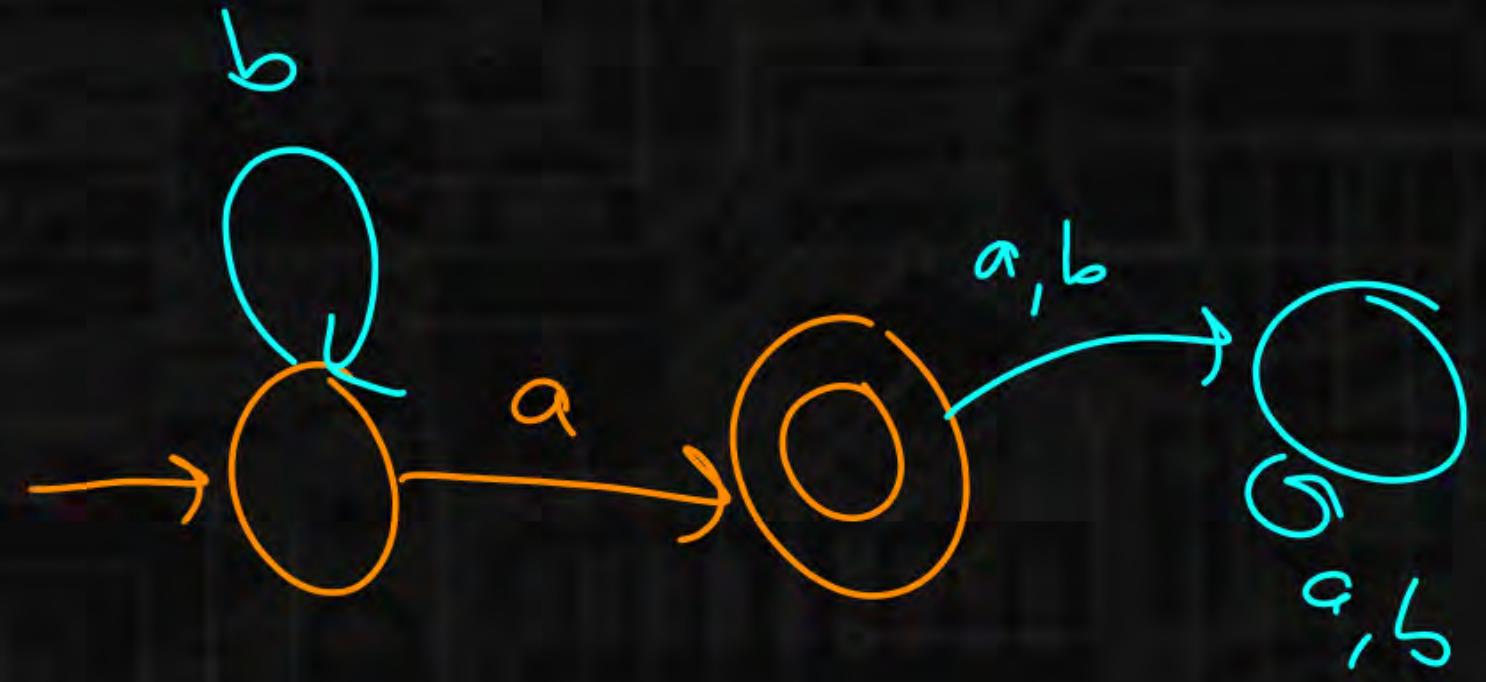
min = aa



(50)

$b^* a$

min = a





H.W.

$$(51) \quad b^* a^*$$

$$(52) \quad b^+ a^*$$

$$(53) \quad b^* a^+$$

$$(54) \quad b^+ a^+$$

$$\Sigma = \{a, b\}$$

$$(55) \quad a^* b^* c^* = \{a^m b^n c^k \mid m, n, k \geq 0\}$$

$$(56) \quad a^* b^+ c^*$$

$$(57) \quad a^* b^* c^+$$

$$(58) \quad a^+ b^* c^*$$

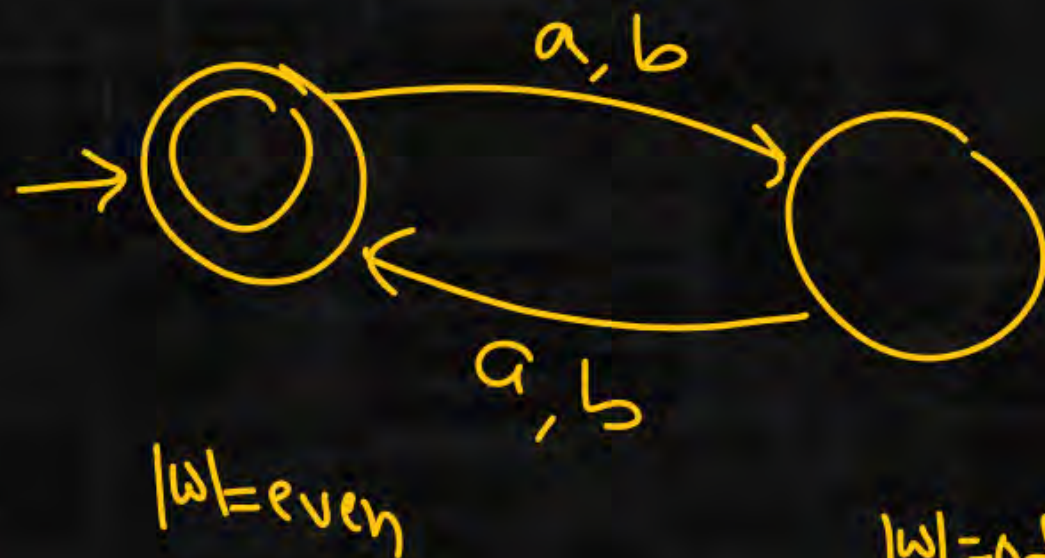
$$(59) \quad a^+ b^+ c^*$$

$$(60) \quad a^+ b^+ c^+$$

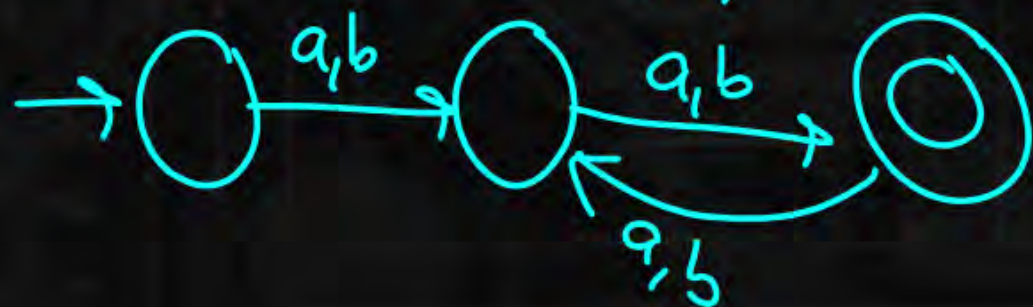


(61)  $\{w \mid w \in \{a, b\}^*, |w| \text{ is divisible by } 2\}$

$$= [(a+b)^2]^*$$



(62)  $\{w \mid w \in \{a, b\}^+, |w| \text{ is divisible by } 2\}$



$$= [(a+b)^2]^+$$

I: special

II:  $|w| = k$   
+ + + + +

III:  $n_a(w) = k$   
+ + + + +

IV: L over 1 symbol,  
but forms A.P.

V: sequence



$|w|$  is divisible by 2

$|w|$  is even

$$|w| \% 2 = 0$$

$$\text{rem}(|w|/2) = 0$$

$$|w| \cong_2 0$$

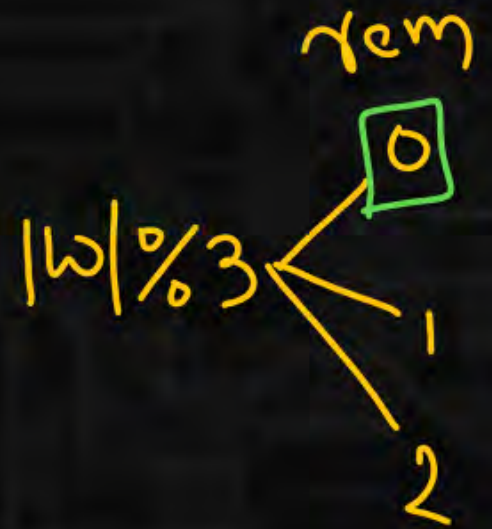
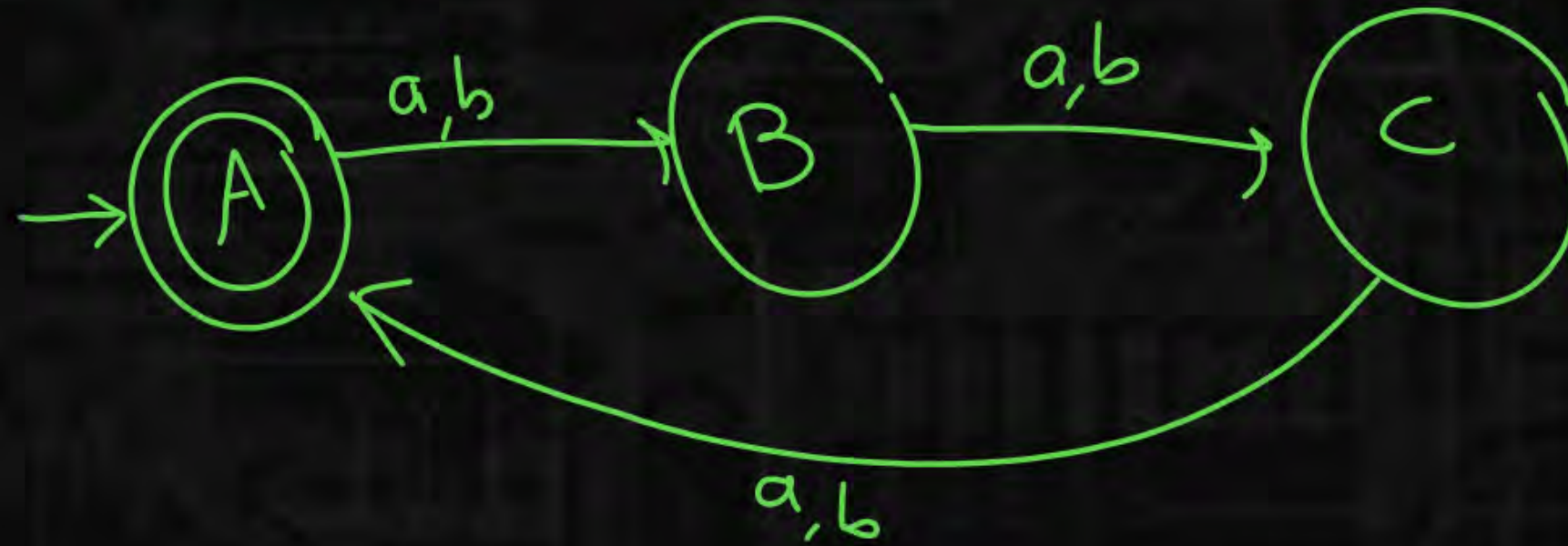
$$|w| = 0 \bmod 2$$

$$|w| = 2n, n \geq 0$$

$$(63) \{w \mid w \in \{a,b\}^*, \underbrace{|w| = \text{odd}}_{|w| \% 2 = 1}\} = \overline{(61)}$$

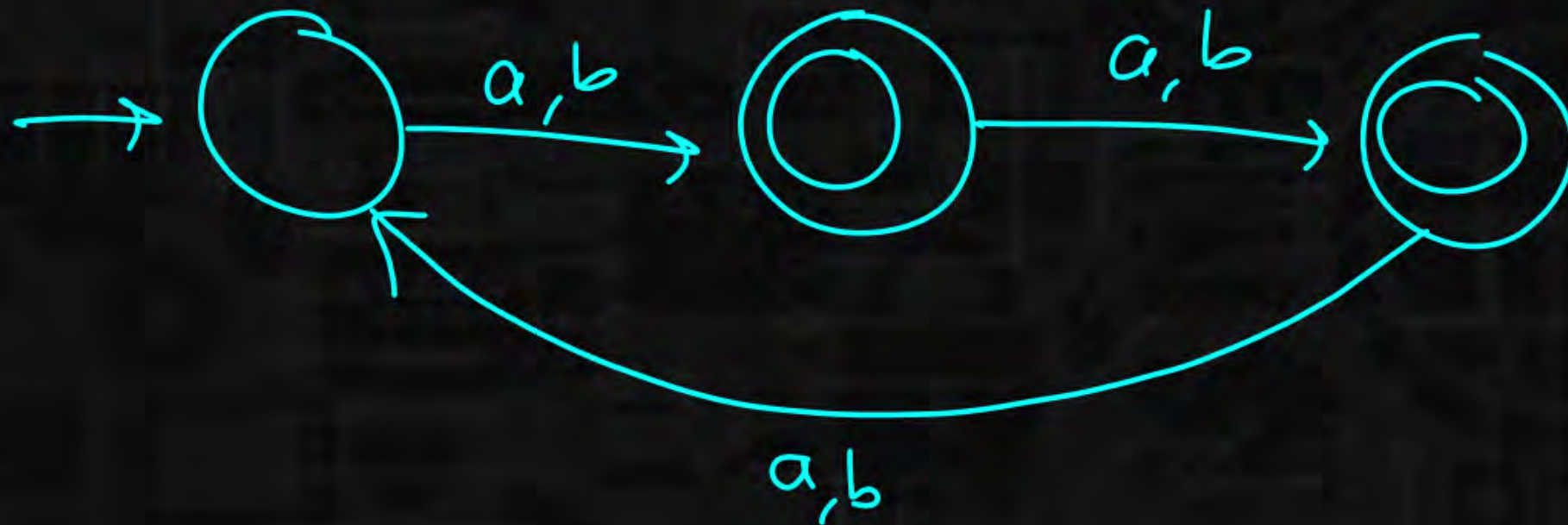


$$(64) \{w \mid w \in \{a,b\}^*, |w| \text{ is divisible by } 3\}$$





(65)  $\{w \mid w \in \{a,b\}^*, |w| \text{ is not divisible by } 3\}$   
( $|w| \% 3 \neq 0$ )



$$|w| = 3n$$

$$|w| = 3n + 1$$

$$|w| = 3n + 2$$

$$(66) \quad \{w \mid w \in \{a,b\}^*, |w| = \underbrace{100n + 5}_{\substack{K_1 > K_2 \\ \text{remainder}}}, n \geq 0\}$$

$\hookrightarrow K_1$  states  
100 states

$$(67) \quad \{w \mid w \in \{a,b\}^*, |w| = \underbrace{100n + 123}_{K_1 \leq K_2}, n \geq 0\}$$

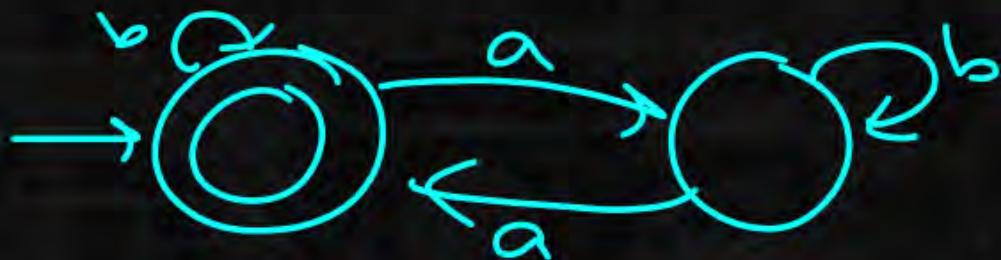
$\hookrightarrow 124$  states  
 $(K_2 + 1)$



## Model-VII [Remainder based &amp; no. of a's]

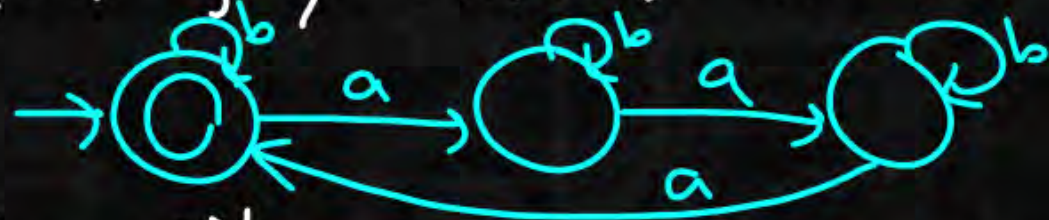
(68)  $\{w \mid w \in \{a, b\}^*, n_a(w) \text{ is divisible by } 2\}$

→ 2 states



(69)  $\{w \mid w \in \{a, b\}^*, n_a(w) \text{ is divisible by } 3\}$

→ 3 states



In general

(70)  $\{w \mid w \in \{a, b\}^*, n_a(w) = k_1n + k_2, n \geq 0\}$

Case I: If  $k_1 > k_2$  then  $k_1$  states

Case II: If  $k_1 \leq k_2$  then  $k_2 + 1$  States



Home Work

min=a  
3 states

71

$$L = \underline{a}(a+b)^* \Rightarrow \{w \mid w \in \{a, b\}^*, w \text{ starts with 'a'}\}$$

min=aa  
4 states

72

$$L = \underline{aa}(a+b)^* \Rightarrow \{w \mid w \in \{a, b\}^*, w \text{ starts with 'aa'}\}$$

min=a  
2 states

73

$$L = (a+b)^*a \Rightarrow \text{Set of all strings ending with 'a'}$$

min=aa  
3 states

74

$$L = (a+b)^*aa \Rightarrow \text{Set of all strings ending with 'aa'}$$

min=a  
2 states

75

$$L = (a+b)^*a(a+b)^* \Rightarrow \{w_1aw_2 \mid w_1, w_2 \in \{a, b\}^*\}$$

min=aa  
3 states

76

$$L = (a+b)^*aa(a+b)^* \Rightarrow \{w \mid w \in \{a, b\}^*, w \text{ contains 'aa' as substring}\}$$

min string + 2

min string + 1



exactly  $\Rightarrow$  Dead state required

At least  $\Rightarrow$  Dead state  
is not required

At most

Start

end

containing

Dead state is not required

$$\Sigma = \{\epsilon, a\}$$
$$\downarrow \quad \downarrow$$
$$= \{x, y\}$$

$$(\epsilon + a)^* = (x + y)^*$$

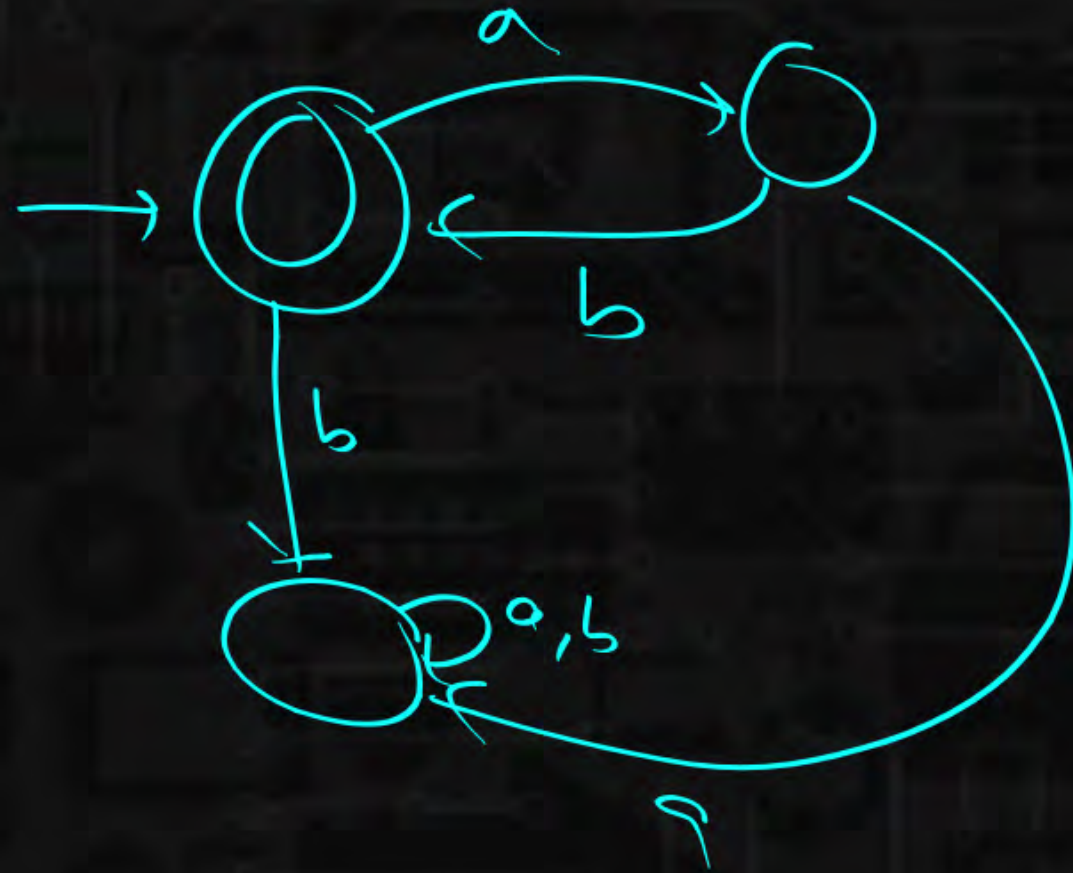
→ In this question, it is i/p symbol

$$(\epsilon a)^* \neq a^*$$

$$(xy)^*$$



$$L = (ab)^*$$



# Summary



model - I

⋮

model - VII

model - VIII

⇒ easy  
DFA is easy  
takes less time to derive

⇒ Answer might be easy  
DFA construction takes time



