


CS & IT ENGINEERING

Theory of Computation

Lecture No.- 02

A man with a beard and mustache, wearing a black polo shirt, stands with his arms crossed in front of a bookshelf. He is wearing a black watch on his left wrist.

Malleham Devasane Sir

Topics to be Covered



Topic

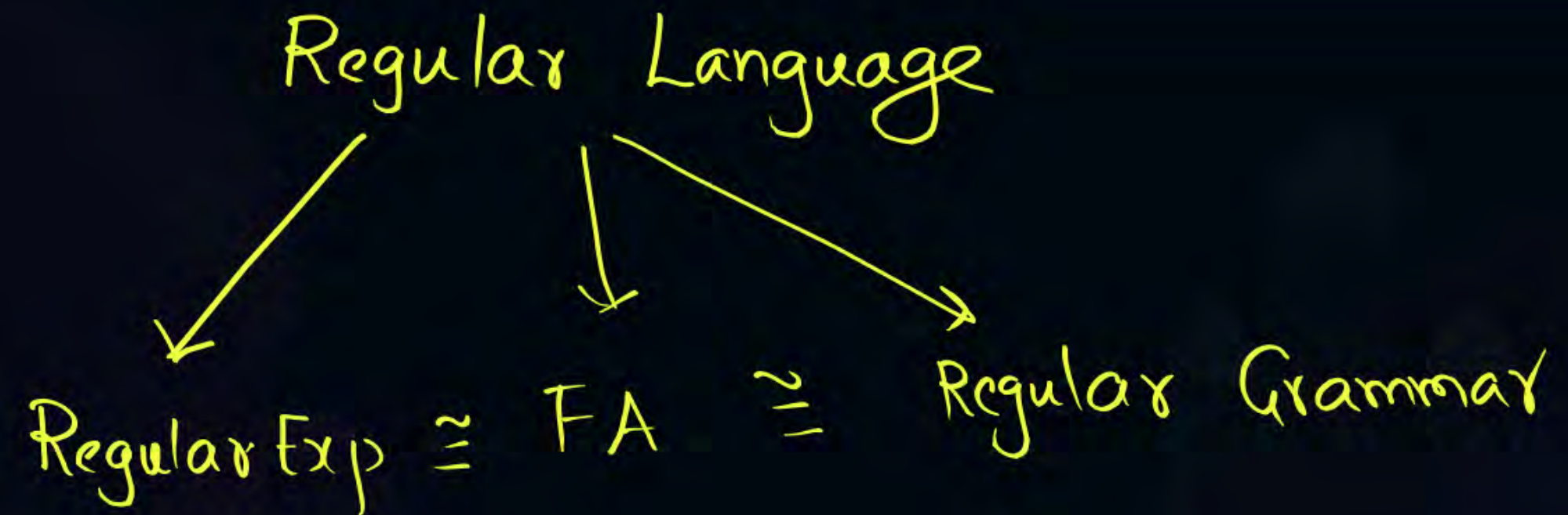
Regular Expression

Topic

Finite Automata



Regular Expression & Finite Automata





Regular Exp & FA : MCQ

$$L = \{\epsilon, b^*, aa, bab^*ab^*, \dots\}$$

Q15. Consider the following regular language L:

$$L = \{w \mid \text{number of a's (w) mod } 3 \neq 1 \text{ where } w \in (a, b)^*\}$$

DEVA SIR PW

Which one of the following represents above language L?

~~A~~

$$(\epsilon + b^*ab^*ab^* + ab^*)(a + b)^* = (a + b)^*$$

~~B~~

$$[(b^*ab^*ab^*ab^*)^* + b^*](b^*ab^*ab^*)$$

ϵ aa

~~C~~

$$[(b^*ab^*ab^*ab^*)^* + (b^* + \epsilon + b^*ab^*ab^*)]$$

$3k$ 0 a's 2 a's

D

$$[b^*ab^*ab^*ab^*]^* + b^*](\epsilon + b^*ab^*ab^*)]$$

$\epsilon \text{ a } \epsilon \text{ a } \epsilon \text{ a } b$ ϵ

5 a's
Not possible

$$n_a(w) \% 3 = 0 \text{ or } 2$$

$$\# \text{a's} = 0, 2, 3, 5, 6, 8, \dots$$

$$n_a(w) \% 3 \neq 1$$

$$\# \text{a's} \neq 1, 4, 7, 10, \dots$$

$$\boxed{\#a's \% 3 = 0 \text{ or } 2}$$

$$\Sigma = \{a, b\}$$

$$= \left(b^* a b^* a b^* a b^* \right)^* \left(b^* + b^* a b^* a b^* \right)$$

$$= \left(\left(b^* a b^* a b^* a b^* \right)^* + b^* \right) \left(\epsilon + b^* a b^* a b^* \right)$$



Regular Exp & FA : MCQ



Q16. Which of the following represents set of all strings starts and ends with different symbols over a's and b's?

$$a \Sigma^* b + b \Sigma^* a$$



$a(a + b)^* b$



$\underline{a}(a + b)^* \underline{a} + \underline{b}(a + b)^* \underline{b}$



$a(a + b)^* b + b(a + b)^* a$



None of these



Regular Exp & FA : MCQ



Q17. Which of the following is TRUE?

~~A~~

$$\underbrace{1(11)^*}_{\text{Min}=1} = \underbrace{\left(1(11)^*\right)^*}_{\text{Min}=\epsilon}$$

~~B~~

$$\underbrace{(11)^*}_{\substack{\text{1 not possible} \\ \epsilon, 11}} = \underbrace{\left(1(11)^*\right)^*}_{\epsilon, 1, 11, 111, \dots}$$

☒ C

$$\underbrace{\left(1(11)^*\right)^*}_{\epsilon, 1, 11, 111, \dots} = \underbrace{1^*}_{\epsilon, 1, 11, 111, \dots}$$

~~D~~

$$\underbrace{(11)^*}_{\text{even}} = \underbrace{1^*}_{\text{All}}$$



Regular Exp & FA : MCQ



Q18. Consider the following regular expression given below:

$$R_1 = (01 + (1 + 01)0)^* (1 + 01)$$

$$R_2 = (01)^* (1 + 01) (0(01)^* (0 + 01))^*$$

Which of the following is correct about R_1 and R_2 ?

- A** String "0110" generated by R_2 but not R_1 .
- B** String "0110" generated by R_1 but not R_2 .
- C** Both expression generates the same language.
- D** None of these



Regular Exp & FA : MCQ

Q19. Suppose the ~~length~~^{Size} of language $|L_1| = 5$ and $|L_2| = 4$ then, which of the following is correct?

$$|L_1 L_2| \geq 8$$

$$L_1 = \{a, b, c, d, e\}$$

$$L_2 = \{f, g, h, i\}$$

$$L_1 L_2 = \{ \boxed{a} \boxed{f}, \boxed{a} \boxed{g}, \boxed{a} \boxed{h}, \boxed{a} \boxed{i}, \boxed{b} \boxed{f}, \dots \}$$

20 strings

$$8 \leq |L_1 L_2| \leq 20$$

- A** $|L_1 \cdot L_2| \geq 5$
- B** $|L_1 \cdot L_2| \geq 20$
- C** $|L_1 \cdot L_2| \geq 9$
- D** $|L_1 \cdot L_2| \leq \underline{\underline{20}}$



Regular Exp & FA : MCQ



Q20. Choose correct statement.

A

$$\overbrace{R + \phi}^R = \overbrace{R \cdot \phi}^{\phi} \quad \times$$

B

$$\underbrace{R + R}_R = \underbrace{R \cdot R}_{R^2} \quad \times$$

C

$$R + \epsilon = \underbrace{R + \phi}_R \quad \times$$

D

$$\underbrace{R + R}_R = \underbrace{R \cdot \epsilon}_R$$



Regular Exp & FA : MCQ



Q20. ~~Choose correct statement.~~

Which of the following is possible for some R ?

☒ **A** $R + \phi = R \cdot \phi$ possible when $R = \phi$

☒ **B** $R + R = R \cdot R$ possible when $R = \epsilon$

☒ **C** $R + \epsilon = R + \phi$

☒ **D** $R + R = R \cdot \epsilon$

$$\underbrace{\phi + \phi}_{\phi} = \underbrace{\phi \cdot \phi}_{\phi}$$



Regular Exp & FA : NAT

Q21. Consider the following regular expressions R:

$$R = (\underbrace{ab^* + ba^*})^* (\underbrace{ba^* + ab^*}) = (a+b)^+ \Rightarrow 2 \text{ states}$$

Number of states are needed to design a DFA for above expression R is ____.

H.W.:

$$R = (ab^+ + ba^+)^* (ba^* + ab^*)$$

ϵ , a, b, aa, ab, ba, bb, ...
x, ✓, ✓, x, ✓, ✓, x, ...



Regular Exp & FA : MCQ

Q22. Consider the following deterministic finite automata (DFA):

In above DFA {S} is starting state and F_1 and F_2 are final states.

Which of the following language belong to above DFA?

~~A~~

$L = \{ \overbrace{wxw^R}^{a \ b \ a} \mid w, x \in (a, b)^+ \} = axa + bxb$

B

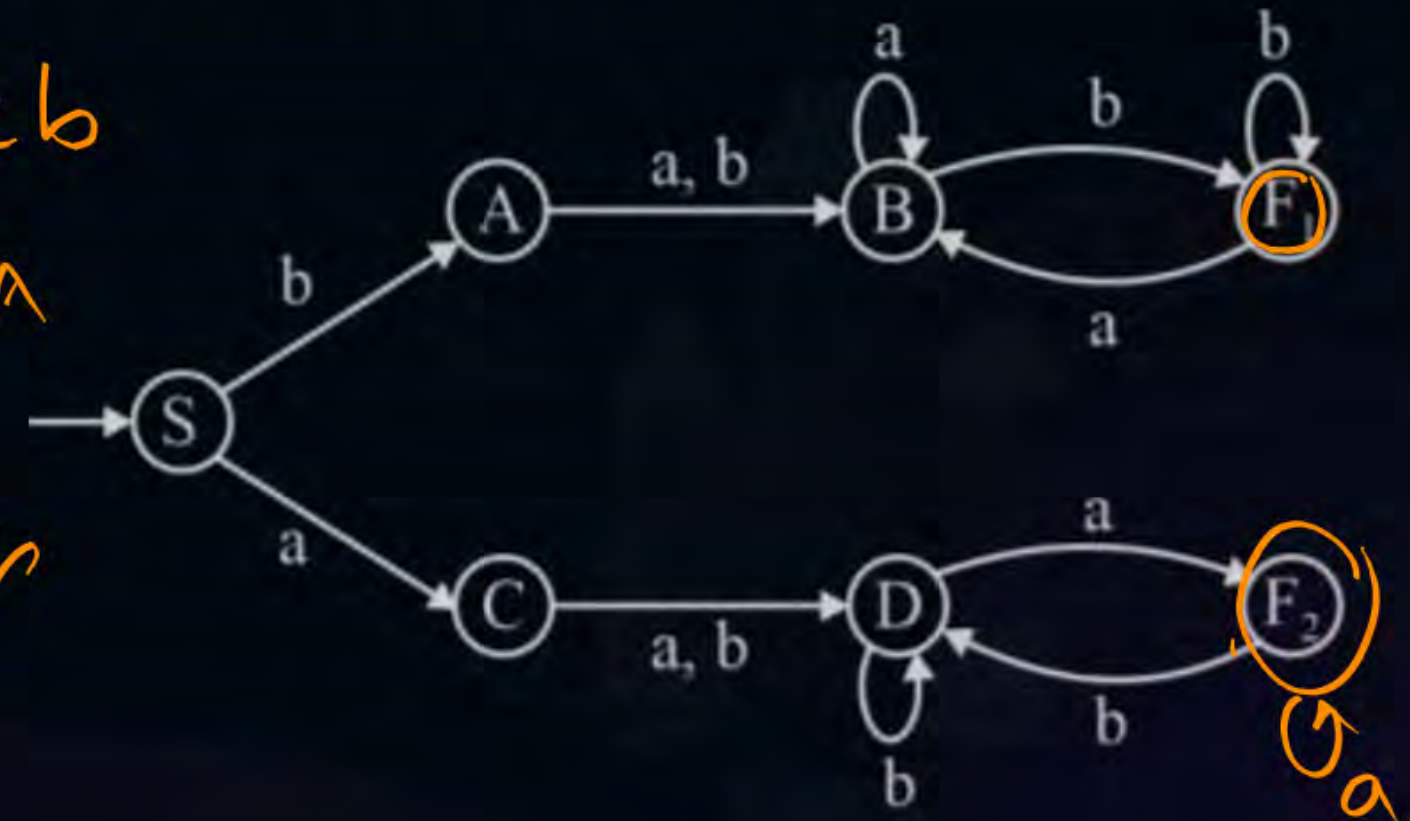
$L = \{ ww^R x \mid w, x \in (a, b)^+ \} \Rightarrow \text{No DFA}$

C

$L = \{ wxw^R \mid w, x \in (a, b)^* \} = (a+b)^*$
 $\rightarrow \text{No DFA}$

D

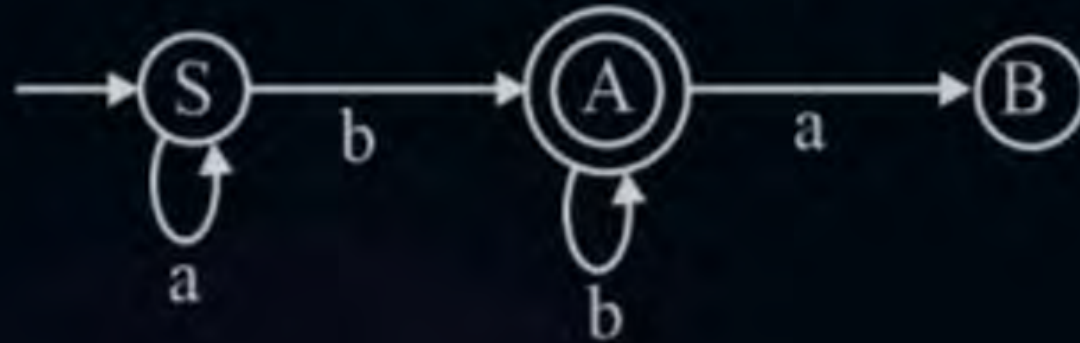
None of these.





Regular Exp & FA : MSQ

Q23. Consider the following finite automaton F:



$$= a^* b b^*$$

In above finite automata S is starting state and A is final state. Which of the following is/are correct regular expression for above finite automata?

~~A~~

$a^* b b^* a$

~~B~~

$a^* b^+$

~~C~~

$(b + a a^* b) + (b + a a^* b) b b^*$

~~D~~

$a^* b b^*$

$$\begin{aligned} & \underline{b + a a^* b + \underbrace{b b b^*}_{b b b} + \underbrace{a a^* b b b^*}_{a^+ b b^+}} \\ & \checkmark \end{aligned}$$

$$\begin{aligned} & a^0 b^+ \checkmark \\ & a^+ b^+ \end{aligned}$$

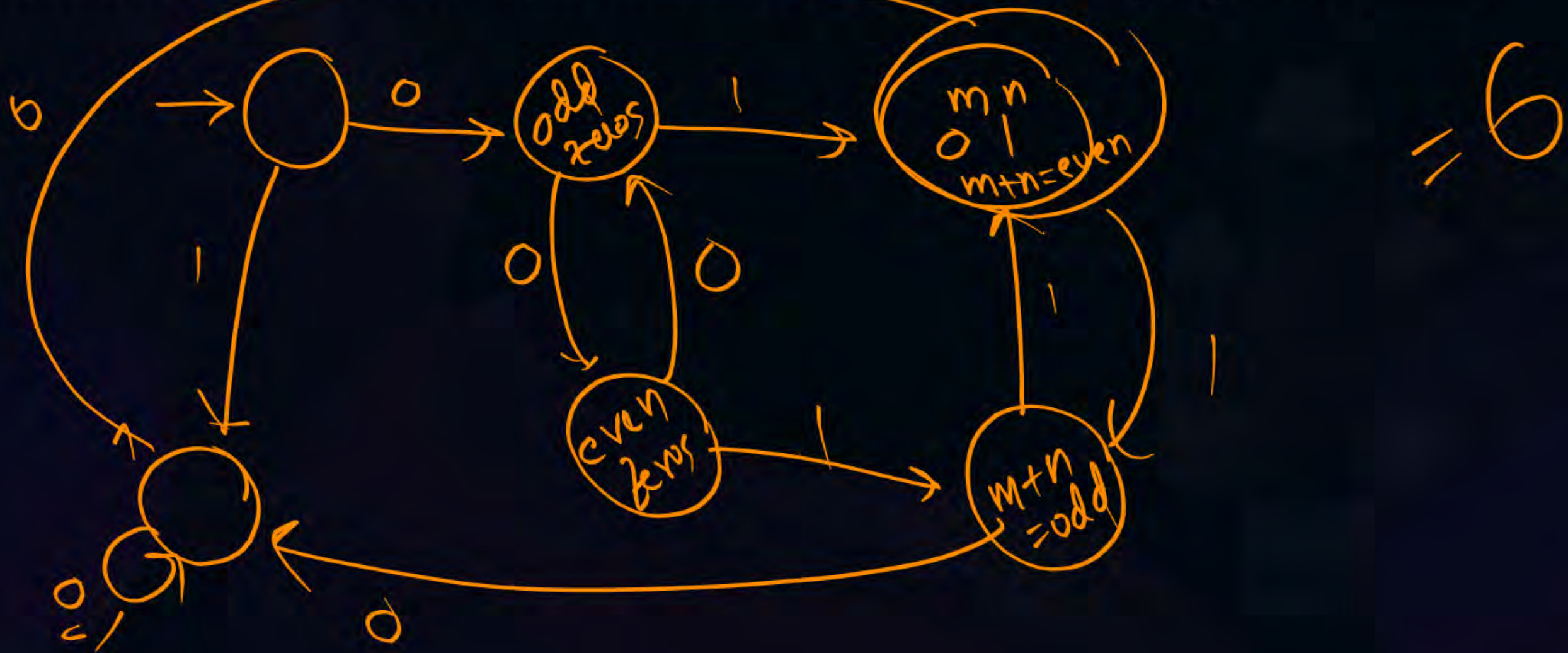


Regular Exp & FA : NAT

Q24. Consider the language L on $\Sigma = \{0, 1\}$:

$$L = \{0^m 1^n \mid m+n = \text{even and } m, n \geq 1\} = \{01, 0011, 0111, 0001, \dots\}$$

The minimum number of states needed for DFA of language L is ____.





Regular Exp & FA : NAT

Q25. Consider the following regular expression R:

$$R = (00 + 11 + (01 + 10) (00 + 11)^* (01 + 10))^*$$

Minimum number of states are needed in DFA for above regular expression R is

— $L = \{w \mid w \in \{0,1\}^*, n_0(w) = \underbrace{\text{even}}_{\text{div by 2}}, n_1(w) = \underbrace{\text{even}}_{\text{div by 2}}\}$





Regular Exp & FA : MSQ

Q26. Consider the Language L:

$$L = \{w \mid |w| \% 99 \leq 27 \text{ where, } w \in \{a, b\}^*\}$$

What is the regular expression for above language L?

$$|w| \% 3 \leq 1$$

$$(a+b)^3 (\epsilon + a+b)$$

A

$$((a+b)^{27})^* (a+b)^{99} \quad |w| = 27k + 99, k \geq 0$$

B

$$((a+b)^{99})^* (a+b)^{27} \quad |w| = 99k + 27, k \geq 0$$

C

$$((\epsilon + a + b)^{99})^* (\epsilon + a + b)^{27} = (a+b)^* \quad |w| \geq 0$$

D

$$\underbrace{((a+b)^{99})^*}_{\epsilon} \underbrace{(\epsilon + a + b)^{27}}_{\epsilon} \quad |w| = 99k + m, k \geq 0, m \leq 27$$

$$|W| \% 99 \leq 27$$

$$|W| \% 99 = 27$$

remainder

$$\left(\sum^{99} \right)^* \sum^{27} (a+b)^{99} (a+b)^{27}$$

$$\left(\sum^{99} \right)^* (\epsilon + \Sigma)^{27}$$

$$|W| \% 99 = 0$$

OR

$$|W| \% 99 = 1$$

OR

$$|W| \% 99 = 2$$

OR

$$|W| \% 99 = 3$$

OR

$$|W| \% 99 = 27$$

$$|W| \% 99 = 0$$

$$\left(\sum^{99} \right)^*$$

multiple of 99

$$|W| \% 99 = 1$$

$$|W| = 99K + 1$$

$$\left(\sum^{99} \right)^* \Sigma$$

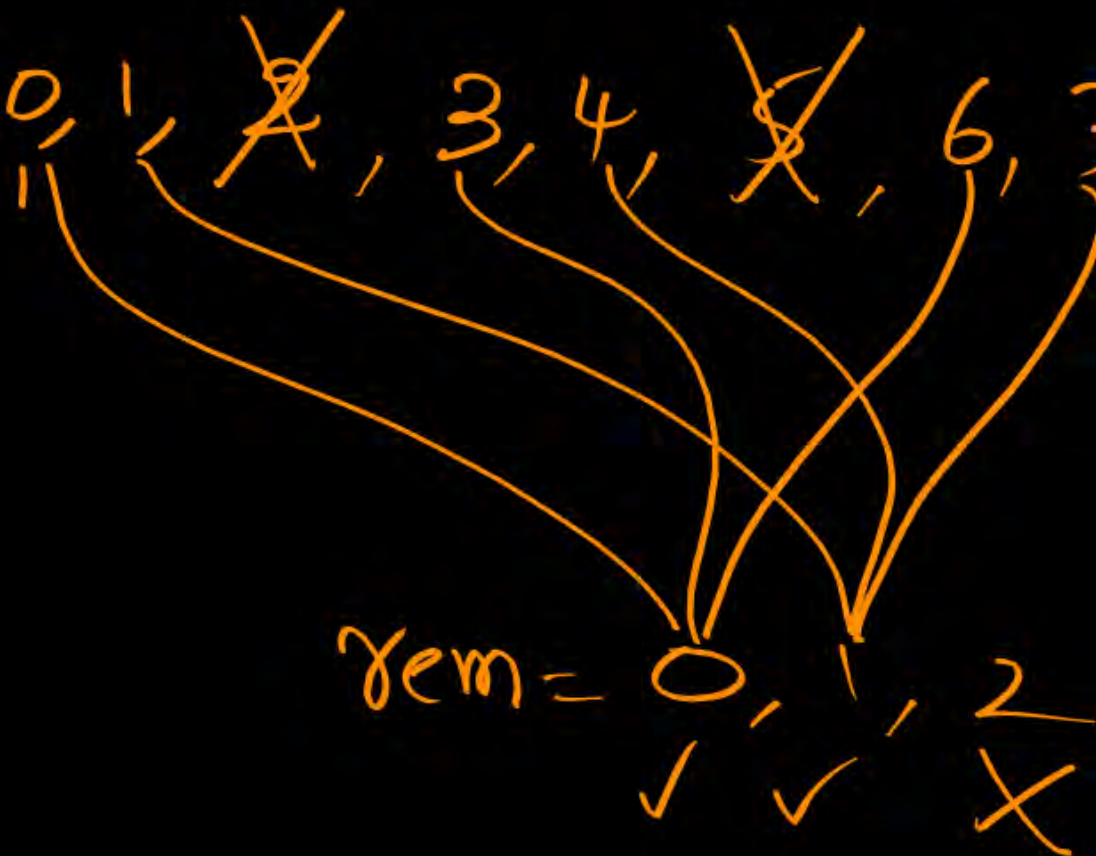
$$|W| \% 99 \leq 1$$

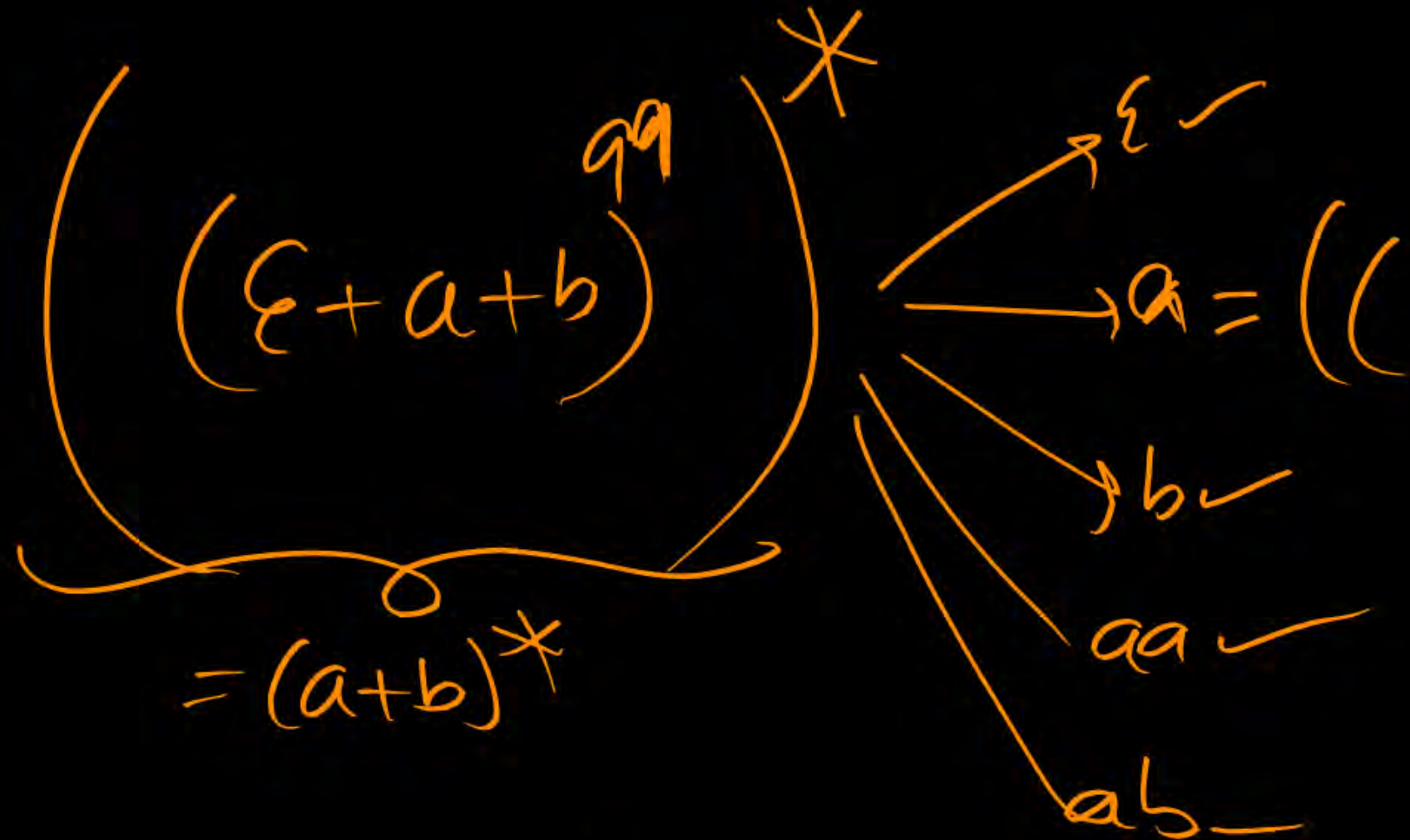
$$\left(\sum^{99} \right)^* (\epsilon + \Sigma)$$

$$|w| \% 3 \leq 1$$

$|w| = 0, 1, \cancel{2}, 3, 4, \cancel{5}, 6, 7, \cancel{8}, \dots$

rem = 0, 1, 2
 ✓ ✓ ✗





$$\begin{pmatrix} aa \\ \end{pmatrix} = \underbrace{(a) (\epsilon)}_{a}$$





Regular Exp & FA : MCQ

Q27. Which of the following regular expression represents the no two consecutive zeros ending with 1?

☒ A

$(01 + 10 + 11 + 1)^*$ $\rightarrow \epsilon \checkmark$
 $\rightarrow 0 \times$

☒ B

$(01 + 1)^+$ $\rightarrow \epsilon \times$

☒ C

$(01 + 10 + 11 + 1)^* 0$ $\rightarrow \epsilon \times$

☒ D

None of these

$L = \{ \epsilon, 0, 1, 00, 01, 10, 11, \dots \}$

$000, 001, \dots$



Regular Exp & FA : MCQ

Q28. Consider the following regular expression R.

$$R = \epsilon + 0(10)^*(\epsilon + 1) + 1(01)^*(\epsilon + 0)$$

which one of the following languages over the alphabet $\{0,1\}$ is described by the above regular expression R?

$$\Sigma^* 00 \Sigma^* + \Sigma^* 11 \Sigma^*$$

~~A~~

Set of all binary strings having either 00 or 11 as a substring.

→ ϵ is invalid

~~B~~

Set of all binary strings not having 00 as a substring.

$\epsilon, 0, 1, 01, 10, 11$

~~C~~

Set of all binary strings not having 11 as substring.

→ 00 should be in R
 $00 \notin R$ if this option is correct

~~D~~

Set of all binary strings neither having 00 nor 11 as a substring.



Regular Exp & FA : MCQ



Q29. Suppose $L_1 = 0^*$

$$L_2 = 10$$

$$L_3 = \{1^m 0^m \mid m \geq 0\}$$

$$L_4 = 1^*$$

If $L_5 = ((L_2/L_1) - L_4) - \bar{L}_3$ Then, the language L_5 will be:

A

ϕ

B

$\{10\}$

C

1^*

D

$\{\epsilon\}$



Regular Exp & FA : NAT



Q30. The number of states in the minimum sized DFA that accepts the language defined by the regular expression $(00 + 111)^*$ is_____.

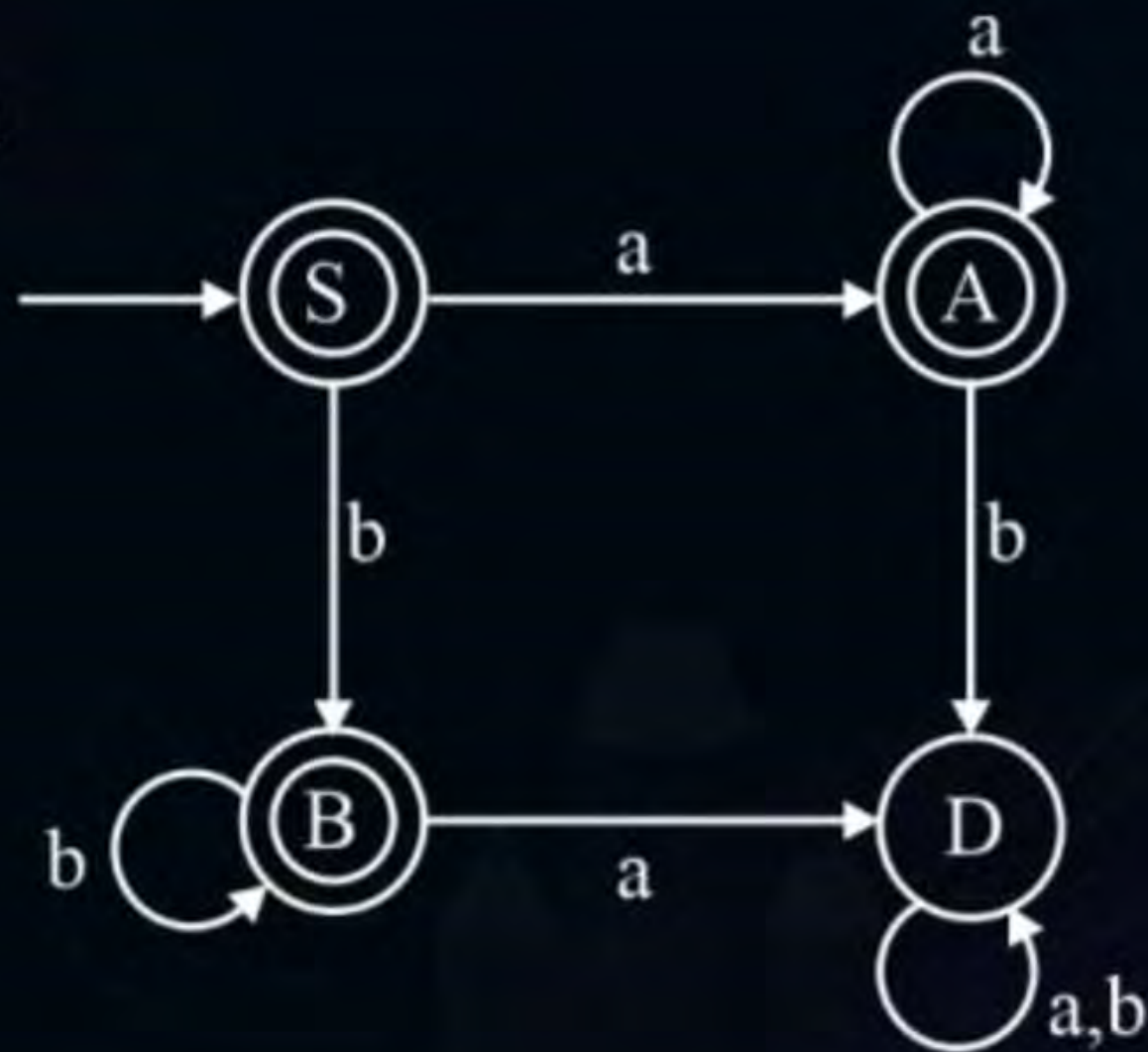


Regular Exp & FA : MSQ



Q31. Which of the following regular expression is equivalent to the finite automaton?

- A** $(a + b)^* a + (a + b)^* b$
- B** $a^* + b^*$
- C** $\epsilon + a^+ + b^+$
- D** $\epsilon + aa^* + bb^*$





Regular Exp & FA : NAT



Q32. Consider grammar G:

G:

$S \rightarrow aSa \mid a \mid b \mid \epsilon$

Let $L = \{w \mid w \in L(G) \text{ and } |w| = 14\}$

Then how many strings are possible in L? ____.



Regular Exp & FA : MCQ



Q33. Consider two languages L_1 and L_2 on $\Sigma = \{a, b\}$.

$L_1 = \{aa, ab\}$ and $L_2 = \{aa, ab, abab\}$ then which of the following is true?

A $L_1^* \subset L_2^*$

B $L_2^* \subset L_1^*$

C $L_1^* = L_2^*$

D $(L_1 \cup L_2)^* = (a+b)^*$



Regular Exp & FA : MSQ



Q34. Which of the following is/are correct regular expression for $L = \{\text{starting and ending with } a\}$?

A

$a(a + b)^* a$

B

$a(a + b)^* a + a$

C

$a(a + bb^*a)^*$

D

$a(a + bb^*a + aa)^*$

THANK - YOU