# CS & IT ENGINEERING

Theory of Computation

PDA: closure properties

Lecture No. 06



By- DEVA Sir



TOPICS TO BE COVERED



### Closure properties



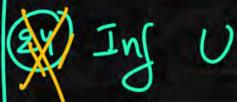
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La Finite languages
- Infinite languages
 -> Regular languager
```

# closure properties for CFLs

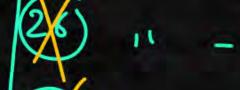








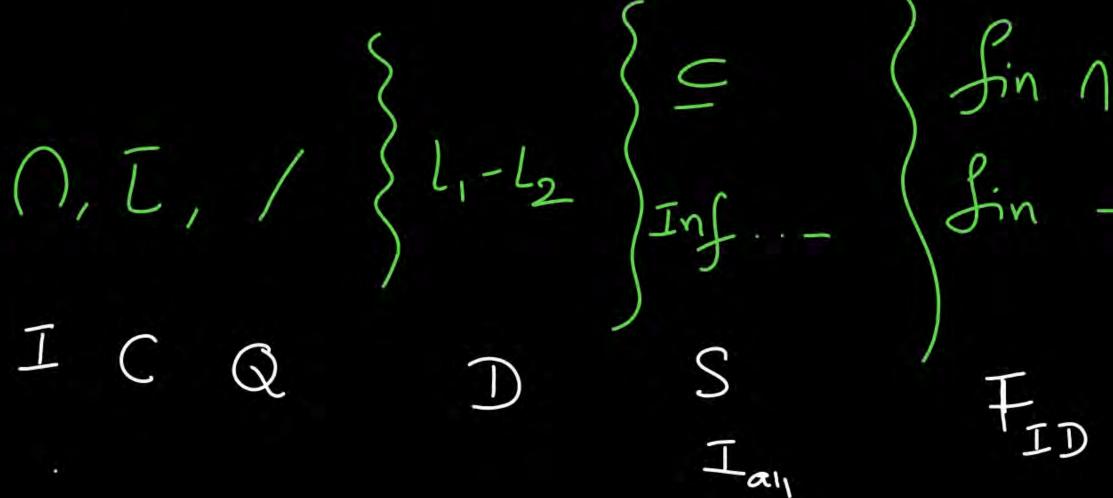








Not closed for CFLs:







# closure properties for DIFLS

- 1 L, Ul2
- (2) LINL2
- 3 I
  - (4) L1-L2
  - (5) L1. L2
  - (G) Rev
  - 7 1
  - 8 Lt

- 9 Subset (L)
- (L) prefix (L)
  - (i) Suffix (L)
  - (12) Substring(L)
  - (B) f(L)
- (14) h(L)
- (15) E- free h(L)

- (L)
  - (F) L1/L2
- 18 Finite U
- (19) "
- (20) 11 -
- 2) 11.
- 2 II S
  - (25) 1/ f

- 24) Ins U
- 25) 11
- 26) 11 -
- (2<del>)</del> "•
- (25) 11 S
  - 29 11 f

closed for DCF 13:



I, pref, finc,

C
P
In F,

# Regular closures for CFls

Pw

CFL

## Regular closures for DCFLs



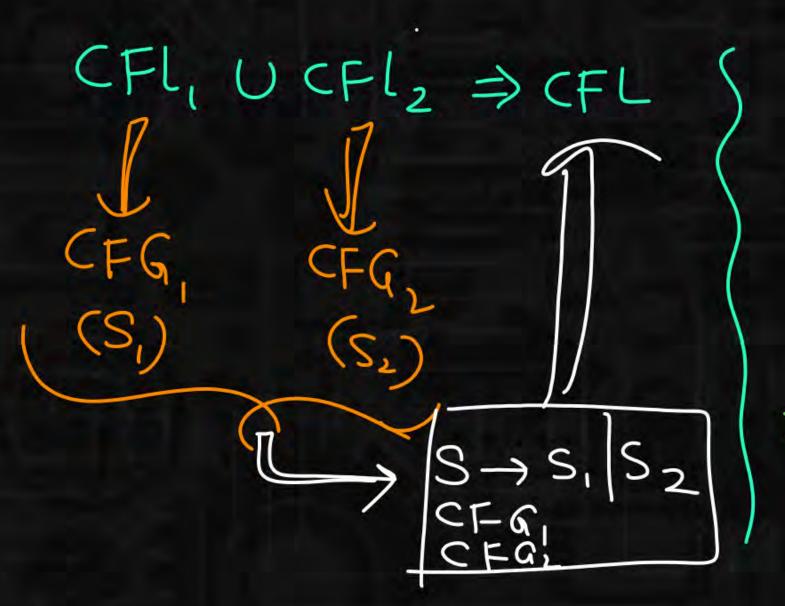
\* DCYL

#### Union



-> closed for CFLs

-> Not closed for DCFL,



D(FL, UDCFL2 F) Need not be DCFL (Blueys (FL) \* 11) { a 6 0 0 0 da 5 2 = not DUFL

Pw

- (I) CFL, UCFL, => CFL
- (2) DCFL, UDCFL2 => CFL
- 3 CFL UDGFL = CFL
- (4) CFL U Reg => CFL
- (5) DOFL U Reg = ) DOFL
- (6) CFL U Fin =) CFL
- DIFL U Fin =) JUFL

- (8) CFL UING F) Inf
- 9 DOFL UING A) INS

Z. Jak



- (1) (a+b)\* U {a"b"} → (a+b)\*

  Reg
- (2) dais U dais ) CFL but not DCFL
- (3) 中 (1) (2) 一人(2)
- (4) da b U da b = CFL byt not DCFL
  - 可有的的人有意识的一种

## 2 Intersection

Pw

Hy Not closed for CFls

July 11 11 for DCFls

CFL, MCFl<sub>2</sub> Deed not be CFL

DCFL, NDCFL2 (Always CSL)

Example:

(Always C

(



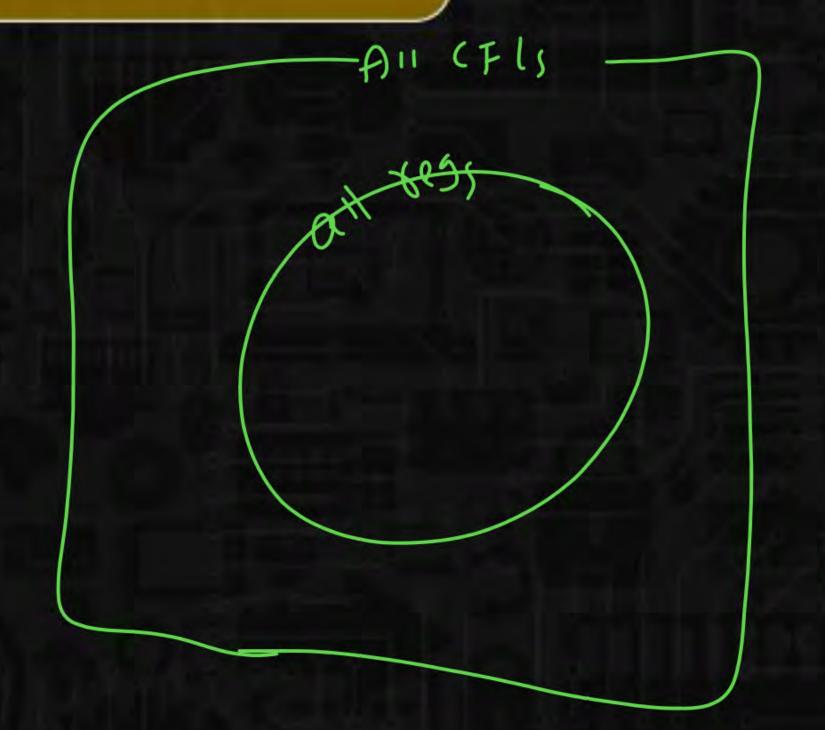






Jabn Nat Dable CFL N Reg Arked not be Set of CFL, NSUL of Reg, I) Set of Reg,





Set of all set of Resident Set of all Rep

3 Complement



In Not closed for CFLs

-> Not closed for DCFLs

CFL => Need not be CFL
(Always CSL)

Prof

L = 260 is CF(

L=abcn is not cfc

DCFL => DCFL

Soul :

Lis D(FL I DPDA I + - nf DPDA

L BDIFL



$$= \left\{ \begin{array}{l} \sum_{k=1}^{n} \sum_{k=1$$

Pw

app is (Fl

abais not (FL

(a+b+1) is (FL \$ 15 CFL



$$\frac{L_2}{L_2} = \frac{\{\omega \omega^R \mid \omega \in \{a, b\}^*\}}{|\omega \in \{a, b\}^*\}} \text{ is also } cfl$$

(4) D'ifference



> Not closed for CFLs

"" " Too DCFLs

1,-12 - 1, 1 T2

(5) Concatenation

-> closed for CFIs

-> Not closed for D(FIs

CFL, . CFl2 => Always (FL

DCFL, DCFL, =) Need not be DOFL aaa



Reversal -> closed for CFIs -> Not closed for DCFIs (CFL) => CFL

CELE) CEC Exent Leverse (Len)

DCFL) => Need not be octhed to the control of the c

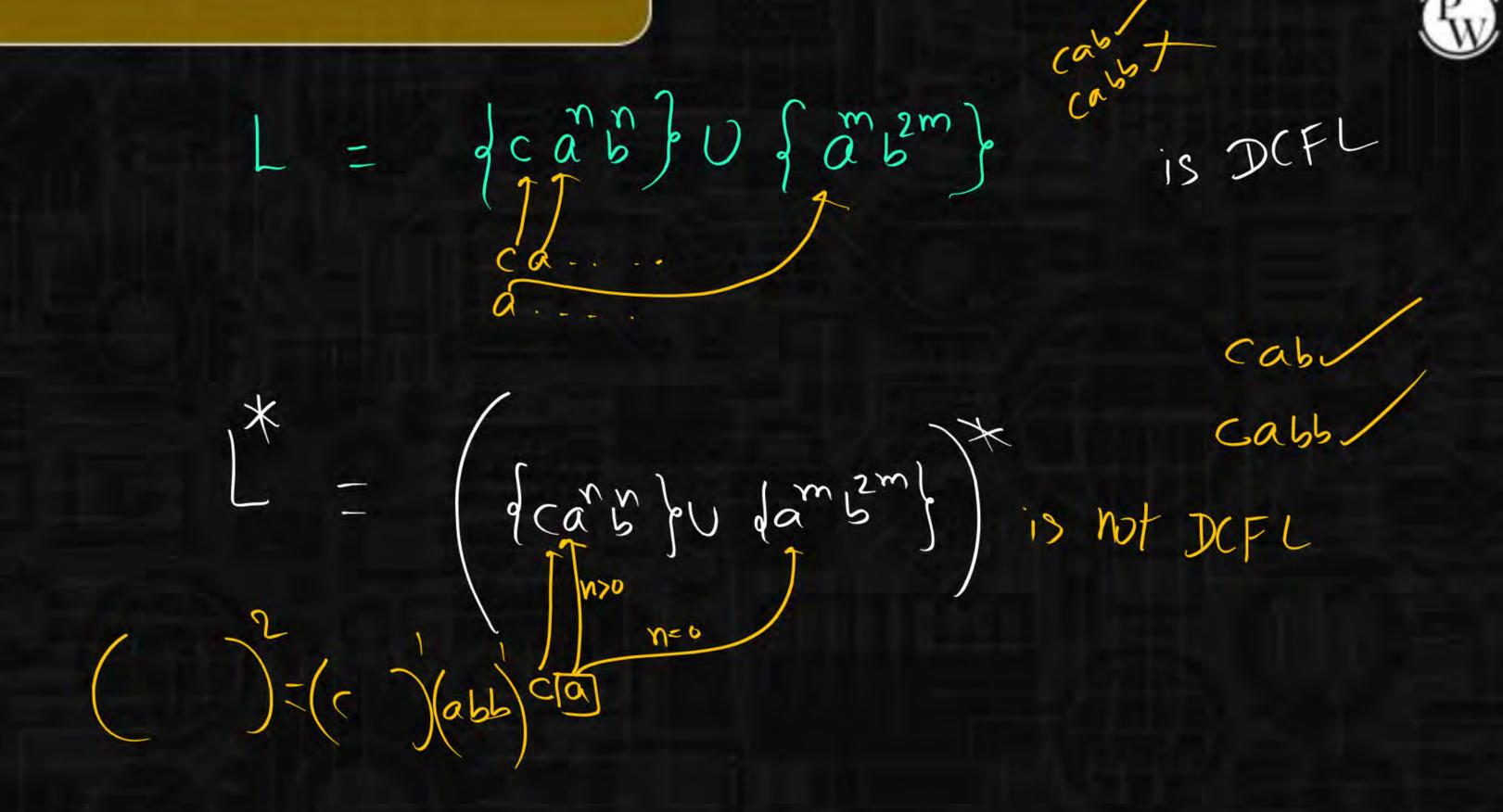
L={dcab} O {dab} is DCFL

Rue l'Eacfol Em m de is rot art

(7) Kleine star



(8) Kleene plus J





L=
$$\{ab\}$$
 is DIFL

 $\begin{array}{c}
\chi = \{ab\} \\
\chi = \{ab$ 

Subset Ve5s DCFLs mot closed CFLS CSLS Recs REL DCFL CFL CFL 554 ap apu φ not DOFC notry

DEPL

DCFL

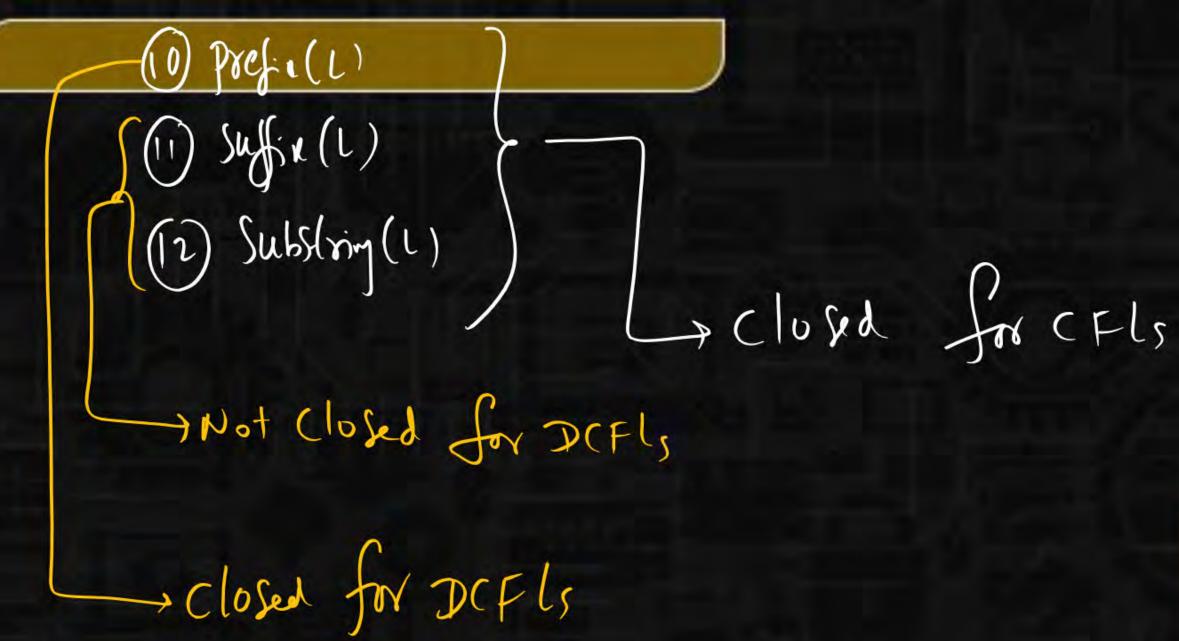




Subset of regular language is need not be seen

Subset of DCFL is need not be DCFL

Subset of CFL is need not be CFL





Prefix of DCFL =>DCFL
Complement of DCFL



Prof for Prefa of

L=ab CFL

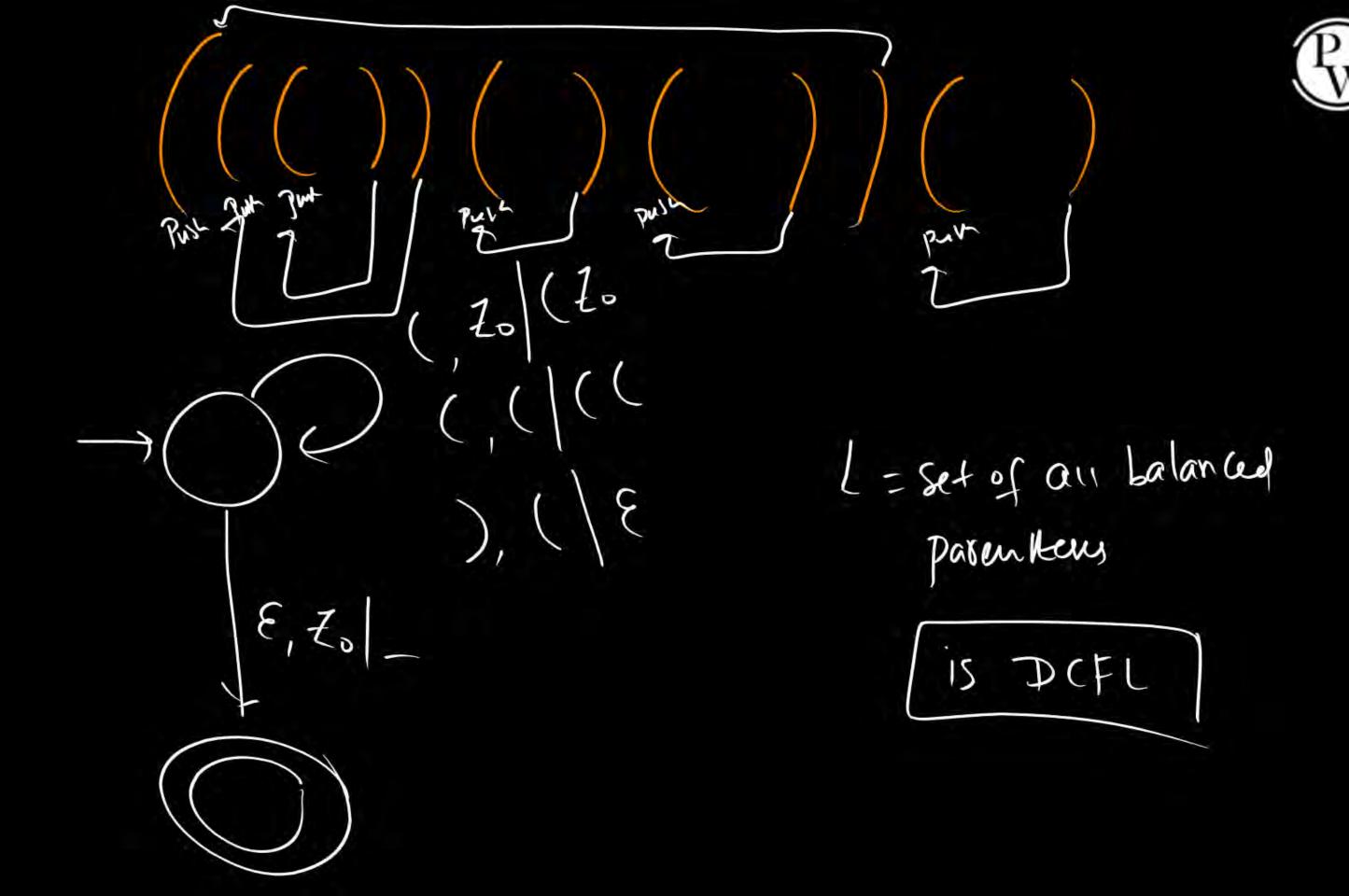
a a S a S b

breto (r)= on p/m=n/

Homomorphism & Substitution -> closed for EFLS
-> Not closed for DCFLS Tifcas Detr 1(c)=1E}

Homomorphism & Substitution a Not closh for DCFls Tifcaphylasp) 4(6)=(6) 3-9 B-aBbb/E







CFLs

| River Met Close1

O, I Diff = fin 0

Quotient Diff = fin 0

Sin 3:4

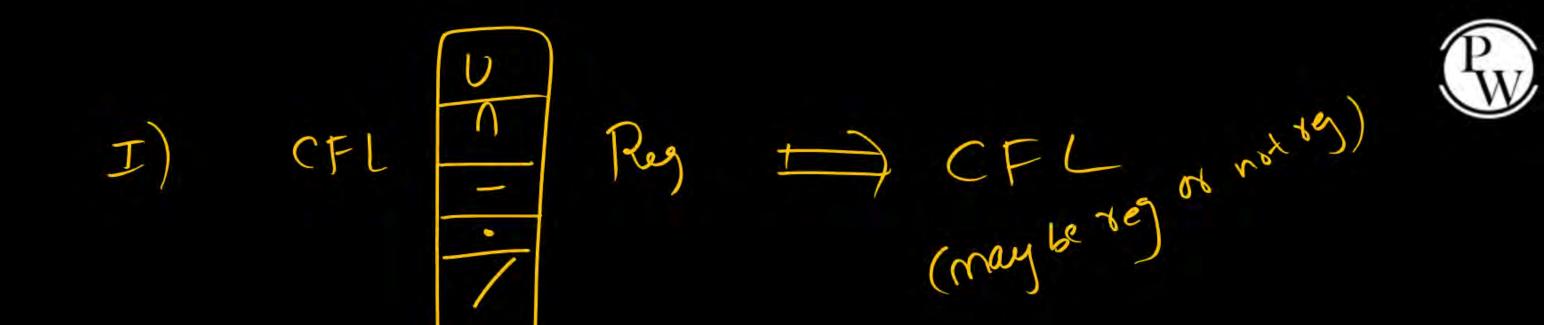
D(TL

Remember

Remember

Closed

T Pret | Fin s



T) Reg-CFL => need not be CFL

Reg-CFL => Reg n CFL => Reg n CSL => CSL

D(FL Py-D(FL #)D(FL Rey-D(FL #)D(FL Rey-D(FL #)D(FL FL Rey ND(FL Rey ND(FL

