CS & IT

ENGINERING

Database Management System

FD's & Normalization

DPP 01 Discussion Notes



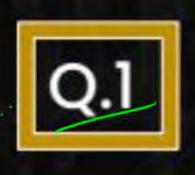




TOPICS TO BE COVERED

01 Question

02 Discussion



According to RDBMS rules, choose the correct statement from the following.





A relation in RDBMS can have multiple attributes.



A relation in RDBMS is a set of rows and columns.



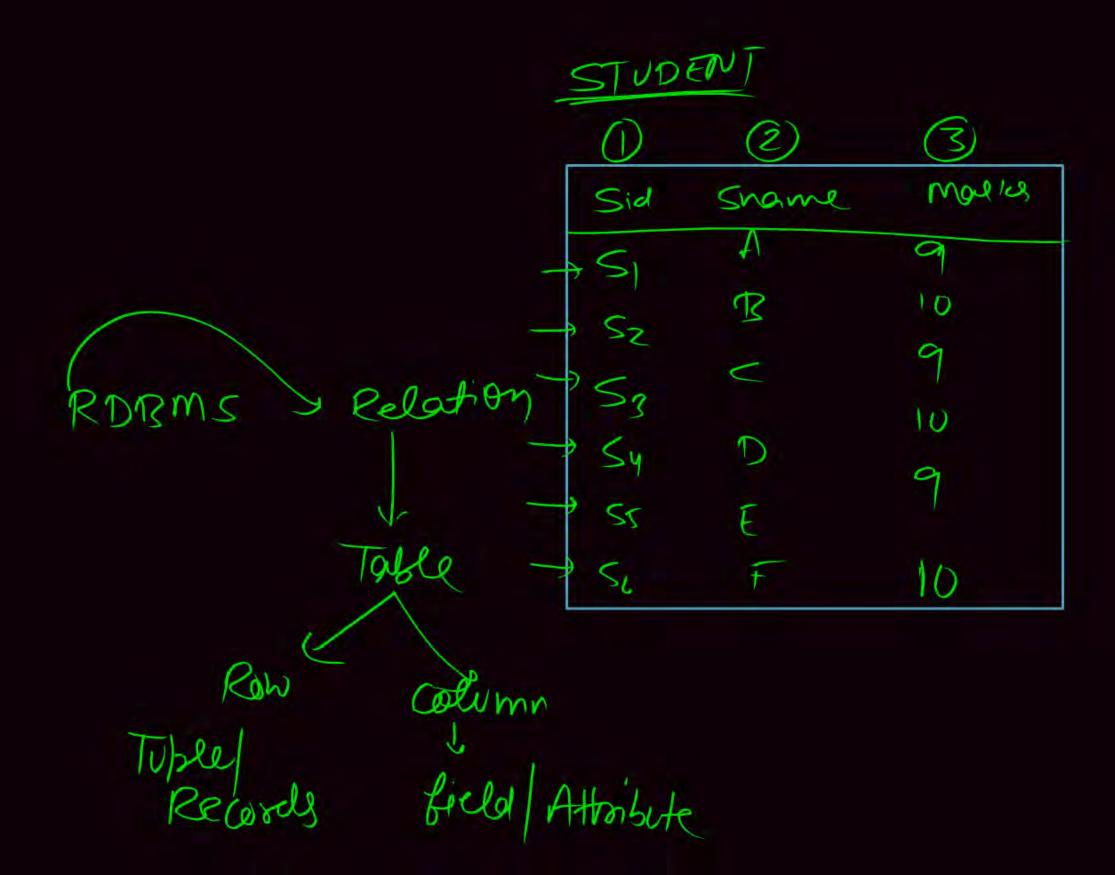
A tuple in a relation can have multiple values for an attribute.

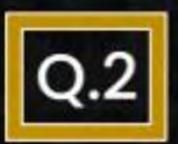


All of the above



Multivalued



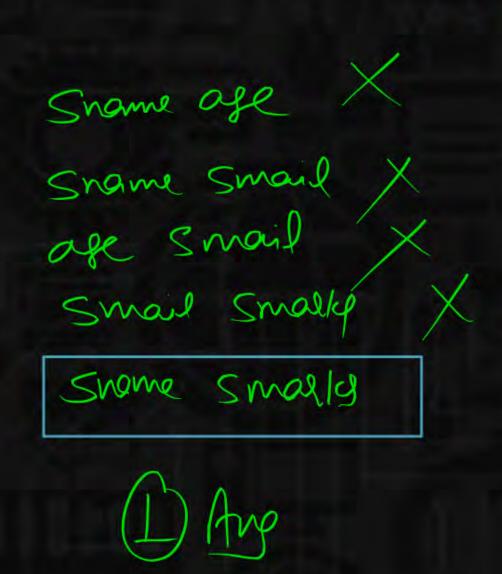


Consider the student relation shown below with schema stud (Sname, Sage, Smail, Smarks), [NAT]

id (W)

Stud

Sname	Sage	Smail	Smarks
Rohit	28	R@pw.live	68
Kanika	25	K@pw.live	75
Pankaj	25	K@pw.live	75
Rohit	28	R@pw.live	88
Anjali	26	A@pw.live	75



For the above given instance how many 2-set of attributes can determine a row uniquely?

Consider a relation schema R(A, B, C, D, E, F, H) with the given

Functional dependency set:

 $\{A \rightarrow BC, C \rightarrow AD, DE \rightarrow F, C \rightarrow F\}$

The attribute closure that contains all the attributes of the relation R is?



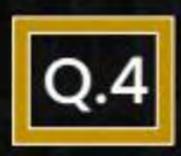




AEH+ [AFH] - [AFHBCDF] - (ABCD FFH)

All of the above





(G)

Consider the below relation schema Stud (Rid, name, course, mail, phone) with FD set as:



Rid → {Rid} -> Toivial FD

Rid → {name, mail} → Non Trivial FD

Semi course -> {course, phone}

phone → {phone} → Tovial

mail → {Rid, course} -> Non Tokial FD (5)

Servir Non Trivial: 1

name → {phone, mail, course} The number of non-trivial FD's in the given FD set is/are?



Tovial FD X -y is tovial

IR X = y

 $AB \rightarrow B$ $AB \rightarrow AB$

Non Trivial FD

X-y is Non Torvial

ibb XMy = \$ A Must Satisfy FD Defination

> A -> B Rall -> Branch

Semi Non Toivial FD

X & y 4 × Ny + o

A -> AB

Jeither x \$ y

exny + 6

Q.5

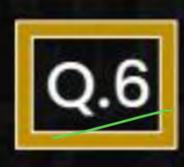
Consider the following set of FD's:

$$\{V \rightarrow W, W \rightarrow XZ, X \rightarrow YZ\}$$
 for relation

R(V, W, X, Y, Z)

Then the attribute closure of YZ⁺ contains how many elements?





For the given FD set: $\{P \to QT, Q \to SU, V \to U\}$ of a relation R(P, Q, T, S, U, V). Find the set of attributes that is Super key but not a [MCQ]



A PTQ
$$[PTQ]^{\dagger} = [PTQSU]$$

B. $[PV]$ $[PV]^{\dagger} = [PVQTSU] \Rightarrow [PQTSUV]$

PQV $[QV]^{\dagger} = [QVSU]$

QV $[QV]^{\dagger} = [QVSU]$

Keys Supel Set of Subel Keys is Super key getting all attribute of Relational Super key

Schena R then X is super key Candidate key: Minimal of Super Key

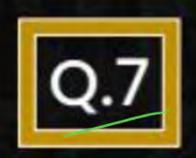
AB is c.k

(Note) Every Condidate key is a Subselley Ret Every Subser key is Not Candidate key

C.K

P.K are
Super
Key

Super key -> Candidate by (4 C.K Minimal I Select as & Poimally key Remaining 3CK AK/ Secondeen key



In a schema with attribute X, Y, Z, W, V, the following set of functional dependencies are given:



$$\{Y \rightarrow X, Y \rightarrow Z, ZW \rightarrow V, X \rightarrow W, V \rightarrow X\}.$$

Which of the following FD is not implied by the above set?





$$YX \rightarrow ZW$$

$$\underline{\mathbf{YX}} \to \underline{\mathbf{ZW}} \qquad (\underline{\mathbf{YX}}^{\dagger} - (\underline{\mathbf{YX}} \times \underline{\mathbf{Z}} \, \underline{\mathbf{W}} \, \underline{\mathbf{V}})$$



$$XV \rightarrow YZ$$



$$ZW \rightarrow V$$



$$XV \rightarrow XW$$

$$\frac{ZW \to V}{XV \to XW} \left(\frac{ZW}{Y} \right)^{+} = \left(\frac{ZW \times X}{X} \right)^{+} = \left(\frac$$

F [A>B, B>c]

Check ADC ? (A) = [ABC]

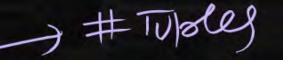
X-y Logically implied from the closure of X

(X3t-[...y.] determines y.



Choose the correct statement from the following.







The cardinality is defined as the number of attributes in a relation.



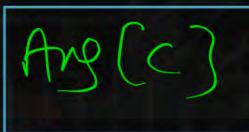
Degree of the relation is the number of tuples in the relation.

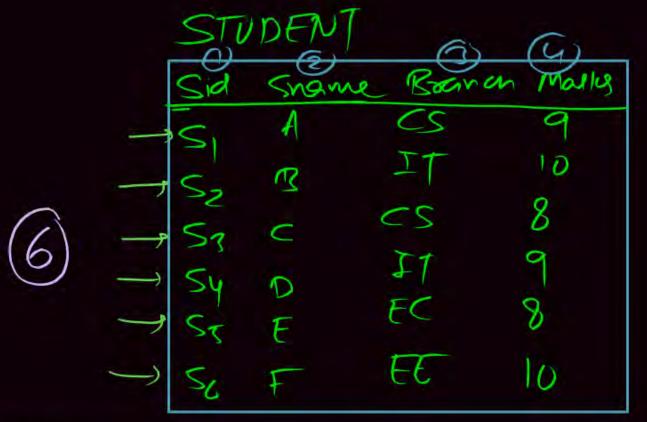


Relation instance is the set of tuples of a relation at a particular instance of time.



All of the above





Depoee

Arity: Number of Attorbute

(4)

Relational schema: STUPENT (Sid Sname Brand Marly) (6)

Relational Instance: Set of Records (Snapshot



