### COMPUTER SCIENCE



Database Management System

FD's & Normalization



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11 Keys Concept

12 Finding Candidate keys





RDBMS Concept

Row (Tube | Record)

Relation (Table)

Column (Attoible | field)

Degree | Arity: # Attoible

Corelinality: # Tubles

Relational Schema: STUDENT (RULINO, Name Branch, CGPA)
Relational Instance extension: Set of Record

### FD (Functional Dependency)

X > if tix=to:x then tix = tox must be some.

Trivial FD

Won Tovial FD

Semi Non Trivial FD

Super key

Attribute closure

A-B, B-HC

@ R does not Bunctionally determine C.

Condude (1) Rule out the FD Bossed on the Table.
(2) Trivial FD are always blid.

### **Keys Concept**

SUPER KEY (Assume 6 C.K) Minimal Gudidate key 1 select og Remain CK Scrondal

except Pk

(5cr

### **Keys Concept**

Cornelidate key: Minimal of Suber key IB Any (Proper Subset) of (Suber key) is also Super key than that Prober Subset is Called Candidate key (2 30 om)

RIARCOE) [AB-C, C-D, B-EA] (ABCDE) Prime | Cey = (B) Parper subset 15 is Candidate

RIABODEI (AB-)C, C+D, B>E) (ABCDE) AR is suber lay Proposer SUDSET  $(A)^{\dagger} = (A)$ Prime low - (A,B) (B) - (BE) Non Prime | Non Pay = [CIDIE]

AB is Candidate key

· Bis Candidate key.

By Any Super Set of B

Bis Super key.

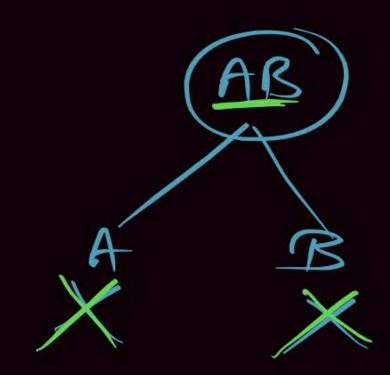
Any Super Set of Super key is also super key.

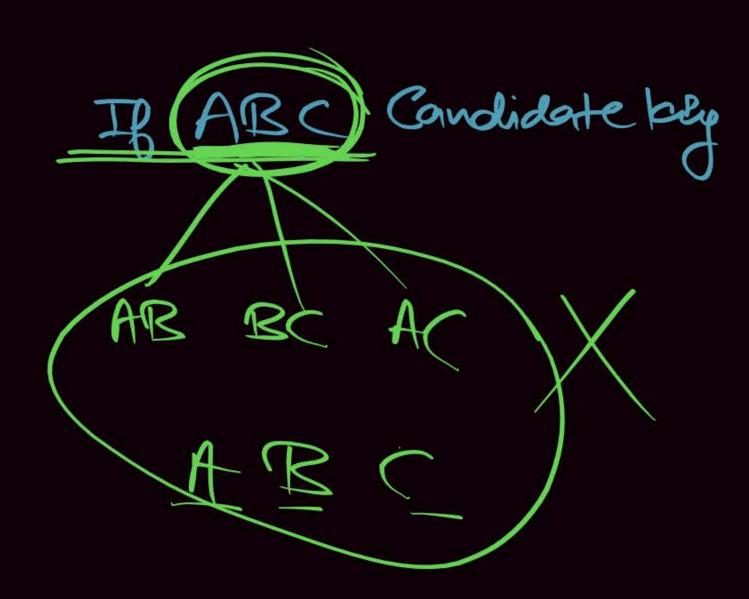


BBBCACE, BACDE

Suber Key

## If AB is Candidate ky





### **Keys Concept**

Suber key.
Ly Any Suber set of suber key is also suber key.

Comolidate key: minimal ab Super key

· Every Goodidate key is a Subset key.

(Note) But every Super key is Not a Candidate key.

Recourse Candidate key is a Minimal of Subser key

Prime key Attribute: Set at Attributes that Present/belongs to Any Some Candidate |cext.

# Non Prime Non key Attribute: Set of Attribute

that Not Present/

Not belongs to Any Candidate key

### Finding Candidate key:

AB is Candidate key.

The Attorbute which Is Not Present in Right Hond Side (R.M.S) that Attrobute Must be Pregent in Candidate (cy

### Finding Multiple candidate key:

Procedure:

First Final Any One Candidate key, then that Attoibute (which Present is called Prime Attoibute.

Prime Attribute) than multiple

Candidate key are there.

D-13 Prime key Attribute = (B, Assume (Assume)

B is Candidate key.  $\mathcal{D} \longrightarrow \mathcal{B}$  $\begin{array}{cccc}
D & \longrightarrow B \\
DE & \longrightarrow B \\
DE & \longrightarrow B
\end{array}$ 



#### R(ABCDEF) $\{A \rightarrow B, B \rightarrow C, D \rightarrow CEF\}$



#### Find candidate keys for the relation R?

$$(A)^{+} = (ABC)$$

$$(D)^{+} = (DCEF)$$

$$(AD)^{+} = (ABCDEF)$$

No Multiple Candidate key Only One C.K



R(ABCDE) {AB 
$$\rightarrow$$
 C, C  $\rightarrow$  D, D  $\rightarrow$  E, B  $\rightarrow$ A, C  $\rightarrow$  B}



$$\subseteq \to \mathbb{B}$$

### R(ABCD) $\{A \rightarrow B, B \rightarrow C, C \rightarrow D, D \rightarrow A\}$

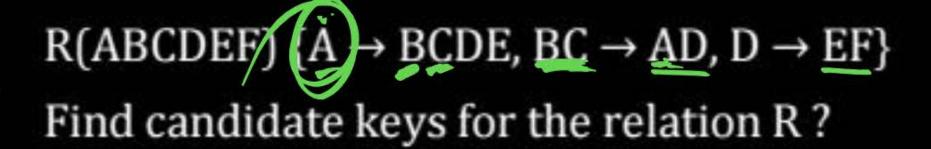


#### Find candidate keys for the relation R?

$$D \rightarrow A$$

$$(D)^{+} - (DARC)$$











### R(ABCDEF) $\{A \rightarrow BCDE, BC \rightarrow AD, D \rightarrow EF\}$





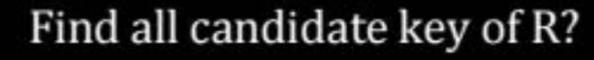
$$(B)^{\dagger} = (B)$$

$$(B)^{\dagger}=(B)$$
 $(C)^{\dagger}=(C)$ 
 $(B)^{\dagger}=(B)$ 
 $(B)^{\dagger}=(B)$ 
 $(C)^{\dagger}=(C)$ 
 $(C)^{\dagger}=(C)$ 
 $(C)^{\dagger}=(C)$ 
 $(C)^{\dagger}=(C)$ 



### R(ABCD) F: $\{AB \rightarrow C, B \rightarrow D, C \rightarrow B, D \rightarrow B\}$







$$(B)^{+} = (B)$$

$$\frac{D}{A} = [ADBc]$$

$$\frac{D}{A} = [A]$$

$$\frac{D}{A} = [A]$$

$$\frac{D}{A} = [A]$$

ABValready taken

ABValready taken

already taken





Consider the following relational schema R(ABCDEF) with  $\bigcup$  functional dependency {AB  $\rightarrow$  C, C  $\rightarrow$  D, D  $\rightarrow$  E, E  $\rightarrow$  F, F  $\rightarrow$  B} The number of candidate keys for relation R?



 $R(ABCDE) : \{AB \rightarrow C, BC \rightarrow D\}$ 



Find Candidate keys for the Relation R?

# Any Doubt?





