# **Branch: CSE & IT**

# **Database Management System**

# **Transaction & Concurrency Control**

**DPP 04** 

#### [MCQ]

- **1.** Consider the following statements:
  - P: If a schedule is allowed by 2PL, it guarantee conflict serializable.
  - Q: A conflict serializable schedule is always allowed by 2PL.

Which of the following is/are correct?

- (a) Ponly
- (b) Q only
- (c) Both P and Q
- (d) Neither P nor Q

#### [MSQ]

- 2. Consider the following schedule: -
  - S:  $R_1(P)$ ;  $W_1(P)$ ;  $R_2(P)$ ;  $R_2(Q)$ ;  $R_1(Q)$ ;  $W_1(Q)$ ;  $C_1$ ;  $C_2$ ; Which of the following is correct?
  - (a) S is conflict serializable.
  - (b) S is not conflict serializable.
  - (c) S is allowed by 2PL.
  - (d) S is not allowed by 2PL.

#### [MCQ]

- **3.** Consider the following statements:
  - P: A schedule allowed by 2PL is always strict recoverable
  - Q: A strict recoverable schedule is always allowed by 2PL.

Which of the following statement is/are CORRECT?

- (a) Q only
- (b) Neither P nor Q
- (c) Ponly
- (d) Both P and Q

#### [MSQ]

- **4.** Consider the following schedule
  - $S: R_x(A); \ W_y(A); \ W_y(B); \ R_z(B); \ R_z(C); \ W_x(C); \ C_x; \\ C_v; \ C_z;$

Which of the following is/are CORRECT?

- (a) The given schedule is conflict serializable.
- (b) The given schedule is allowed by 2PL.
- (c) The given schedule is not allowed by 2PL.
- (d) The given schedule is not conflict serializable.

#### [MSQ]

- **5.** Consider the following schedule
  - $S: W_2(P); \ W_1(P); \ W_2(Q); \ W_3(P); \ W_1(Q); \ W_3(Q); \ C_1; \\ C_2; \ C_3;$

**Batch: Hinglish** 

Which of the following is/are correct?

- (a) S is not allowed by 2PL.
- (b) S is allowed by 2PL.
- (c) The equivalent serial schedule to S is  $T_3 \rightarrow T_2 \rightarrow T_1$ .
- (d) The equivalent serial schedule to S is  $T_2 \rightarrow T_1 \rightarrow T_3$ .

### [MSQ]

- **6.** Consider the following schedule
  - S:  $W_2(P)$ ;  $W_1(P)$ ;  $W_3(P)$ ;  $W_2(Q)$ ;  $W_1(Q)$ ;  $W_3(Q)$ ;  $C_1$ ;  $C_2$ ;  $C_3$ ;

Which of the following is/are correct?

- (a) S is conflict serializable.
- (b) S is allowed by 2PL.
- (c) S is not allowed by 2PL.
- (d) Serial schedule equivalent to S is  $T_2 \rightarrow T_1 \rightarrow T_3$ .

# [MCQ]

- 7. Which of the following is FREE from deadlock?
  - (a) Basic 2PL.
  - (b) Strict 2PL.
  - (c) Conservative 2PL.
  - (d) Rigorous 2PL.

#### [NAT]

- **8.** Consider the following schedule (S):
  - $S: R_1(A); W_2(A); W_1(A); W_3(A); C_1; C_2; C_3;$
  - (i) S is conflict serializable.
  - (ii) S is not conflict serializable.
  - (iii) S is not allowed by 2PL.

How many statements is/are correct for above schedule(s)? \_\_\_\_.

#### [MCQ]

9. Consider the following transaction for schedule(S)S:

$T_1$	$T_2$	$T_3$
R(A)		
	W(A)	
	W(A) W(B)	
		R(B) R(C)
		R(C)
W(C)		

For schedule(S) which of the following is true?

- (a) Conflict serializable and allowed by 2PL.
- (b) Not conflict serializable and not allowed by 2PL.
- (c) Conflict serializable but not allowed by 2PL.
- (d) None of the above.

#### [MCQ]

- **10.** Consider the following schedule(S)
  - $S: \ W_2(A); W_1(A); W_2(B); W_3(A); W_1(B); W_3(B); C_1; \\ C_2; C_3;$

Which of the following is INCORRECT?

- (a) S is recoverable.
- (b) S is cascade recoverable.
- (c) S is conflict serializable
- (d) S is not view serializable.

## [MCQ]

- 11. Which of the following is true about 2PL?
  - (a) If schedule is allowed by 2PL then, there is no lost update problem.
  - (b) If schedule is allowed by 2PL then, there is no starvation.
  - (c) If schedule is allowed by 2PL then, there is possibility of deadlock.
  - (d) None of the above.

#### [MCQ]

**12.** Consider the following schedule(s);

$T_1$	R(A)		W(A)		С	
$T_2$		R(A)		W(A)		C

Which of the following is true about above schedule(S):

- (a) S is free from deadlock.
- (b) S is allowed by 2PL.
- (c) S is conflict serializable.
- (d) None of the above.

#### [MCQ]

- **13.** Consider the following statements:
  - $S_1$ : Every serial schedule is not strict 2PL schedule.
  - S<sub>2</sub>: If schedule allowed by 2PL then it's may / may not be allowed by strict 2PL.
  - S<sub>3</sub>: Every serial schedule may not be conflict serializable schedule.

Which of the following is INCORRECT?

- (a)  $S_1$  only.
- (b)  $S_2$  and  $S_3$  only.
- (c)  $S_1$  and  $S_3$  only.
- (d)  $S_1$ ,  $S_2$  and  $S_3$ .

# **Answer Key**

- 1. (a)
- 2. (b, d)
- **3. (b)**
- 4. (c, d)
- 5. (b, d)
- (a, c, d)
- 7. (c)

- 8. (2) 9. (b)
- 10. (d) 11. (c)
- **12.** (d)
- 13. (d)



# **Hints & Solutions**

1. (a)

> If a schedule is allowed by 2PL then, it is always conflict serializable. But vice-versa not true.

2. (b, d)

Precedence graph -



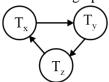
Not conflict serializable and not allowed by 2PL.

**(b) 3.** 

Bothe statements are INCORRECT.

(c, d)

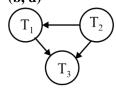
Precedence graph



→ Since a cycle exists, it is not conflict serializable

Hence, it is not allowed by 2PL.

5. (b, d)



→It is conflict serializable

2PL Testing-

$\mathbf{T_1}$	$T_2$	$T_3$
	lock x(P)	
	$W_2(P)$	
	lock x (Q)	
	unlock (P)	
lock x (p)		
$W_1(P)$		
	$W_2(Q)$	
	unlock (Q)	
lock X(Q)		
unlock (P)		
		lock X(P)
		W <sub>3</sub> (P)

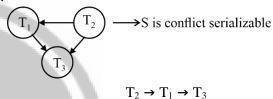
W <sub>1</sub> (Q) unlock (Q)	
	lock X(Q)
	$W_3(Q)$
	W <sub>3</sub> (Q) unlock (P) unlock (Q)
	unlock (Q)

It is allowed by 2PL

Equal serial schedule  $-T_2 \rightarrow T_1 \rightarrow T_3$  (based on point order)

(a, c, d)

Precedence graph -



2PL testing;

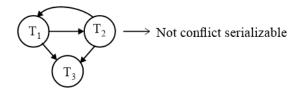
$\mathbf{T}_1$	$T_2$	$T_3$
	lock x(P)	
	$W_2(P)$	
	lock x (Q)	
	unlock (P)	
lock x (p)		
$W_1(P)$	Not possible	
	(unlock (Q)	
not possible		
(Lock X(Q))		
not possible		
(unlock (P)		denied
		Lock x (P)
		$W_3(P)$
	$W_2(Q)$	
	unlock(Q)	
$W_1(Q)$		
		W <sub>3</sub> (Q)
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Hence, the above schedule is not allowed by 2PL.

# 7. (c)

Conservative 2PL is free from deadlock and remaining are not free from deadlock.

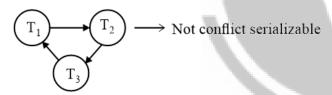
## 8. (2)



$\mathbf{T_1}$	$T_2$	$T_3$
X(A)		
R(A)		
	$X(A) \rightarrow denied$	
	W(A)	
W(A)		
		W(A)

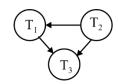
Not allowed by 2PL Only (2) statement are correct.

## 9. (b)



• Not allowed by 2PL because R<sub>3</sub>(C) and W<sub>1</sub>(C)willconflict for X(c)and S(A).

#### 10. (d)



- No cycle so conflict serializable.
- If schedule is conflict serializable then it will be view serializable also.
- Hence, option (d) is correct.

# 11 (c)

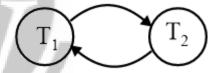
- 2PL condition not sufficient to avoid irrecoverable, cascading roll back and lost update problem.
- 2PL restriction may lead to starvation and deadlock.

Hence, option (c) is correct.

### 12. (d)

	$T_1$	$T_2$
	<b>▼</b> S(A) R(A)	
deadlock both wants to upgrade lock and write on da	¥X(A)	S(A) R(A) U(A) W(A)
item A		

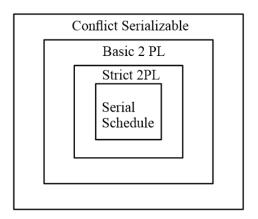
• Not allowed by 2PL



• S is not conflict serializable

Hence option (d) is correct.

# 13. (d)



All statements are incorrect.



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