COMPUTER SCIENCE



Database Management System

FD's & Normalization

Lecture_08

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Finding Super keys

Lossy and Lossless Join



- 1 RDBMS Concept
- 2) FD Concept & its type
- 3 Attribute closure
- (4) SUPER KEY
- 6 Coundidate key 6 Finding Multiple C.K

- 3 Membership set
- (8) Equality blood FD Set
- 9) Minimal Cover
- (10) Closure of FD Set

Any Super Set of Key (Candidak Key) is also
Super Key



Finding Number of super keys

- - (i) With only Condidate key As?
 - (ii) With only Candidate key As, Az?
 - (iii) With Only Candidate key A, Az, AzAy?
 - (in) With Only Candidak key AIA2, A2 A3?
- " (v) with only Candidak key A, A2, A3?

With only Candidate key AI

AI AZ AL AZ AZ

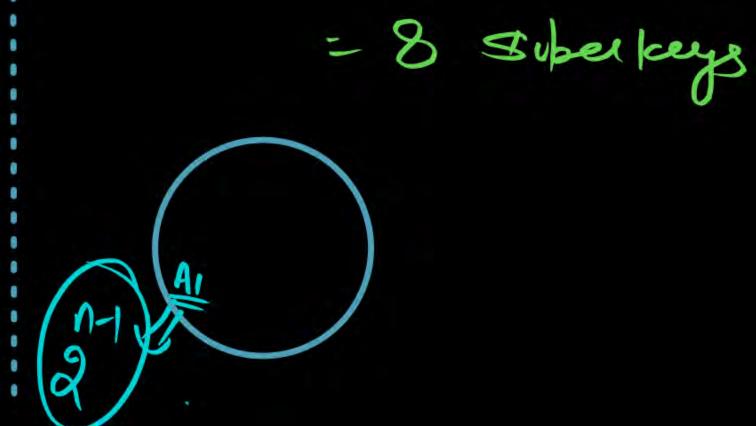


(Ai) Az Az Ay As An

Super key



RIABCD) C.K:A





n=4 Attribute R(ABCD) C.K:A

7 8 Super Keys

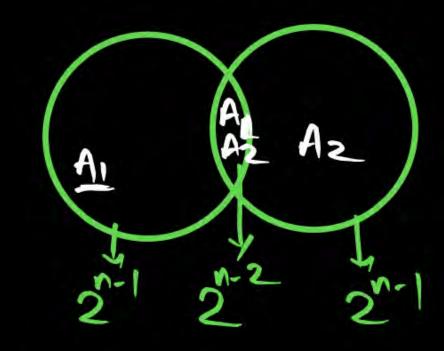
R(ABCD) with CK: A



With only Candidate key A, Az?



AL AZ AZ AM AT AI AZ AZ AM AT AZ AI AZ AY AZ AI AZ AI AZ AI AZ AI AZ AI AZ AI AZ



A1 A2 A3 An A5 ... An :

Az Ar Az An Az. . . An

n(AUB) = n(A) +n(B) - n (A NB)

Superbey = 2 + 2 - 2

.

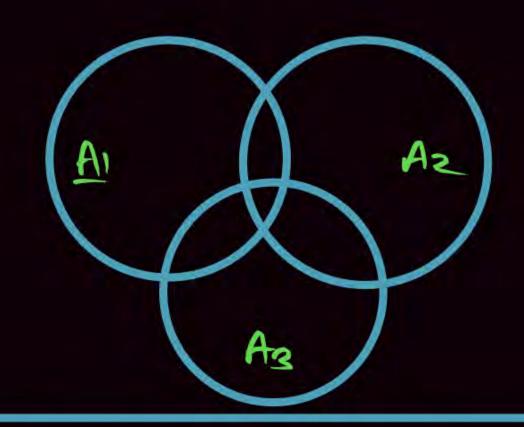




With Buly Candidate key AIAz, Az Az? # Super

super = n-2 n-3 keys = 2+2-2

(SUS) with only Candidate by A, A2, A3 ?



Super keys =
$$\frac{n-1}{2} + \frac{n-1}{2} + \frac{n-1}{2} - \frac{n-2}{2} - \frac{n-2}{2} + \frac{n-3}{2}$$

Finding Number of super keys

(QI) R(ABCDE) With CK: A? (Q.2) R(ABCDE) with CK A, B? (Q.3) R(ABCDE) with C.K AR, CD? (Q4) RIABCDE) With CK AB, BC? QS) R(ABCDE) With CK A, B, C ? (0.6) R(ABCDE) With CK: ABC, DE (07) R(ABCDE) With CK A, BCD?

Seri) R(ABCDE) C.k: A # subser keys = 2⁵⁻¹ = 2⁴ = 16 Avs Sum3) C.k: AB, CD 25-2+25-2-2-1 23+23-21 8+8-2 =14 Avg

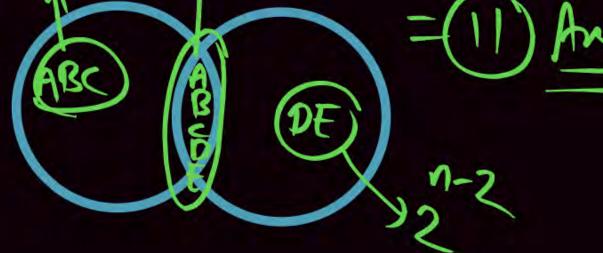
(2) Ck(A,B,C) 2-1+2-1+2-1 -2-2-2-2-2 + 5-3

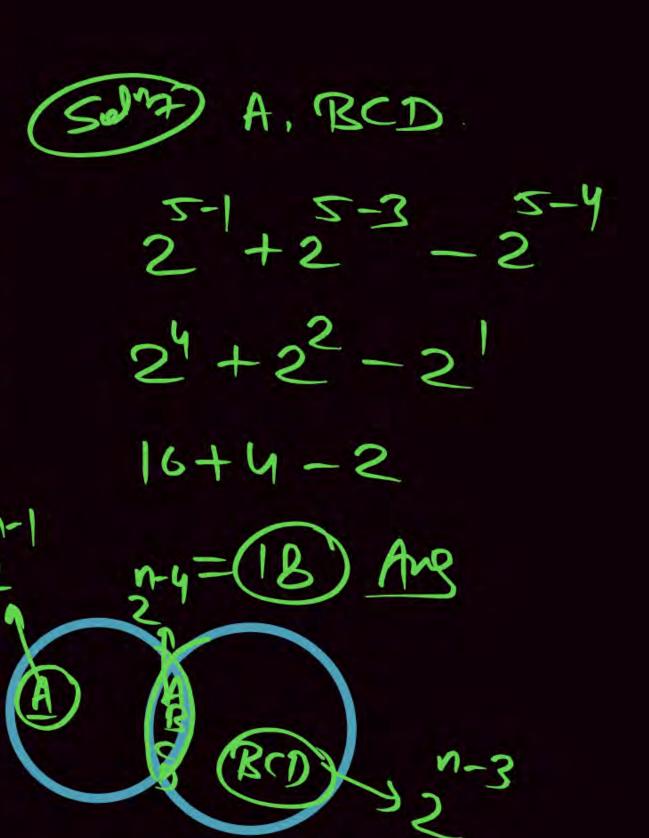
 $2^{1}+2^{1}-2^{3}$ $2^{1}+2^{1}-2^{3}$ =16+16-8=24 Ang

AB, BC (3) 2-2+2-2 2-12 8+8-4 -(12) Ang

 $= \frac{32+21+21}{-2^{3}-2^{3}-2^{3}+2}$ $= \frac{32-4}{(28)} \text{ Ang}$

Sulvep =
$$5-3$$
 $5-2$ $5-3$ | Icerys = $2+2-2$





(GATES) R (EFFGH) With C.K E?

$$\begin{array}{l}
\text{(GW)} \\
\text{(SW)} \\$$

R(EFGH) CK:E

RIEFGH) 8 Super keys E, EF, EG, EH EFG. EGH, EFH, EFGH. 00) JEH 0 JEG 1 1 7 EGH 100 JEF O 1 JEF 1 DIEFG



R(ABCDF) With Candidate key AB

Suber Keys

5.2 R(ARCDE)

AB, ABC, ABD, ABE,

ABCD, ABDE, ABCE, ABCDE.

(QI) R(ABCDE) CK: A, BC

subset =
$$2 + 2 - 2$$

= $2^{4} + 2^{3} - 2^{2}$
= $2^{4} + 2^{3} - 2^{2}$
= $2^{6} + 8 - 4$
= $2^{6} + 8 - 4$
= $2^{6} + 8 - 4$

(RI) R(ABCDE) C.K: A, BC.

II nd Method:

(Q2) R(ABCDE) CK: [A,B,CD)

Siper rigs!

$$ABCD \Rightarrow 2^3 = 8$$
 $ABCACAD, ABCD$
 $ABCD$
 $ABCACAD, ABCD$
 $ABCACA, ACD, ABCD$
 $ABCACA, ABCA, ACD, ABCD$
 $ABCACA, ABCACA, ACD, ABCD$
 $ABCACA, ACD, ABCD$
 $ABCACA, ABCACA, ABCACA, ABCD$
 $ABCACA, ABCACA, ABCACA, ABCD$
 $ABCACA, ABCACA, ABCACA, ABCD$
 $ABCACA, ABCACA, A$

then # Subel lays?

Here n:5

Total Max - # S.K. = 2 -1

Total Maximum Number of Super Keys = 2 - 1

n: # 06 Attribute

(under the Assumption is every Single Attrobute is Candidate key)

Maximum Number of Candidate key = nc 12/2/

n: # af Attailure

$$[2.1] = 2$$

$$[2.9] = 2$$

R(ABCDEF) then bind (Maximum #(C.K)? (n:6 Attribute)

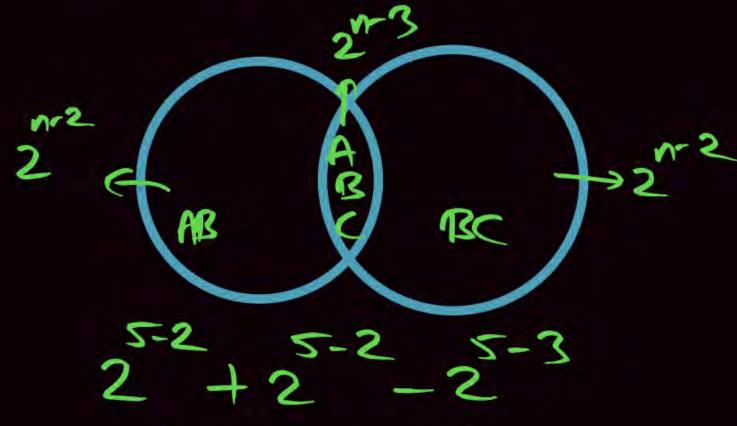
(i) IB every Single Attorbute born a key: 60: every 2 Attribute born a key: 602: (15) (C) 20 (20) (20) 3 Attribute " (iii) 4 Attribute " (i) 5 Attroibute " " (VI) 6 Attribute " "

R(ABCDEF)

- (i) 6C1: A, B.C. D, E, F.
- (ii) 6C2: AB, AC, AD, AE, AF, BC, BD, BE, BF,
 CD, CE, CF, DE, DF, EF = (15)
 - 6C3: ABC, ACD, ADE, AEF....

(in) every 6 Attribute born a Ck: ABCDEF = 0

QI) R(ABCDE) C.K[AB, BC]



II method

R (AMCDF)

$$\frac{8}{8} = \frac{3}{2} = 8$$

$$\frac{3}{8} = \frac{2}{3} = \frac{3}{4}$$

12 guper Iceys (Q.) R(ABCDF) C.K: [AB, CD]

Venn Diagram (AB. CD)

Super = 2 + 2 - 2 = 2 + 2 - 2 = 2 + 2 - 2 = 8 + 8 - 2

= 14 Super Keys.

RIABODE)

AB) CDE = 2 = 8

CD E

2' = 2

AB CD

ARC CDE

ARD CDAT

ABE

ARCD CDB

ABDE CDAE
ABCE CDBE

To suberkey

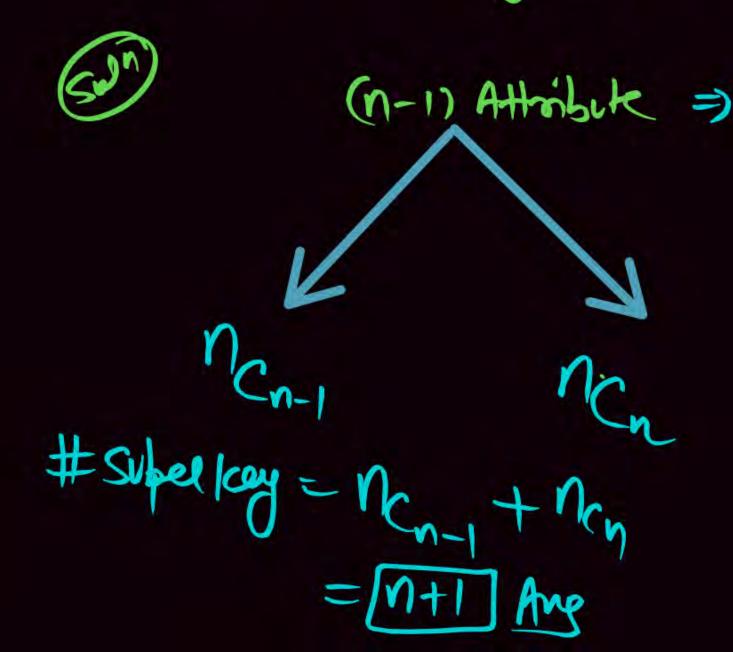
14/

CK (AB, CD)

= 14 Ang

Other Approach.

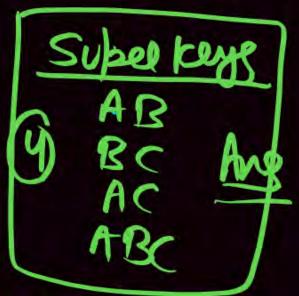
A Relation R with n Attorbutes, is every (n-1) attorbutes from a Candidate key, then Total Number of Suber key.



(n-1) Attribute => Subserkey = nCn-1 + Any Subserset & n-1 & R(ABC)

3CK AC

11:3 Attoibute



(e)
$$R(ABCD)$$
 (n= 4)

Every (n-1) = every

 $(y-1)=3$
 $(y-1)=3$

9+1 =) (5 S.K) Ang

ARC

ACD

ABD

RCD

ABCD

Of A Relation R with n Attributes, is every (n-1) attributes from a Canadidate key, then Total Number of Super key. (n-1) Attribute => Subsetting = ncn-1 + Any Subset Set

Super key

COSET Other Approach. (2) A Relation R with n Attributes, is every (n-2) attributes from a Candidate key, then Total Number of Super key. (n-2) Attroibute =) Subset key = non-2 + Any Subset Set ncn-2 ncn-1 # Super key = $n_{n-2} + n_{n-1} + n_{n}$ Ans

Con Ray Nan Nan

#Super =
$$5c_3 + 5c_4 + 5c_5$$

reggs = $10 + 5 + 1$
= (16)

CREET Other Approach

3 A Relation R with n Attorbutes, is every Single attorbutes from a Candidate key. Then Total Number of Suber key.



#Super = $nc_1+nc_2+nc_3+...nc_n$

Q R(PBRST) CK: [P. T, BR, RS).

Total Maximum # Suber | Cey = 2-1 > 2-1

=(31 Super Keys

Not a key (Never, key) R S S

or or or or

Condidate bey: Minimal

OR, RS CK

R Never Served Candidate Key

