COMPUTER SCIENCE & I.T



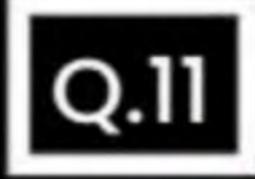
Process Synchronization/ Coordination 01







- IPC & Synchronization
- Types of Synchronization



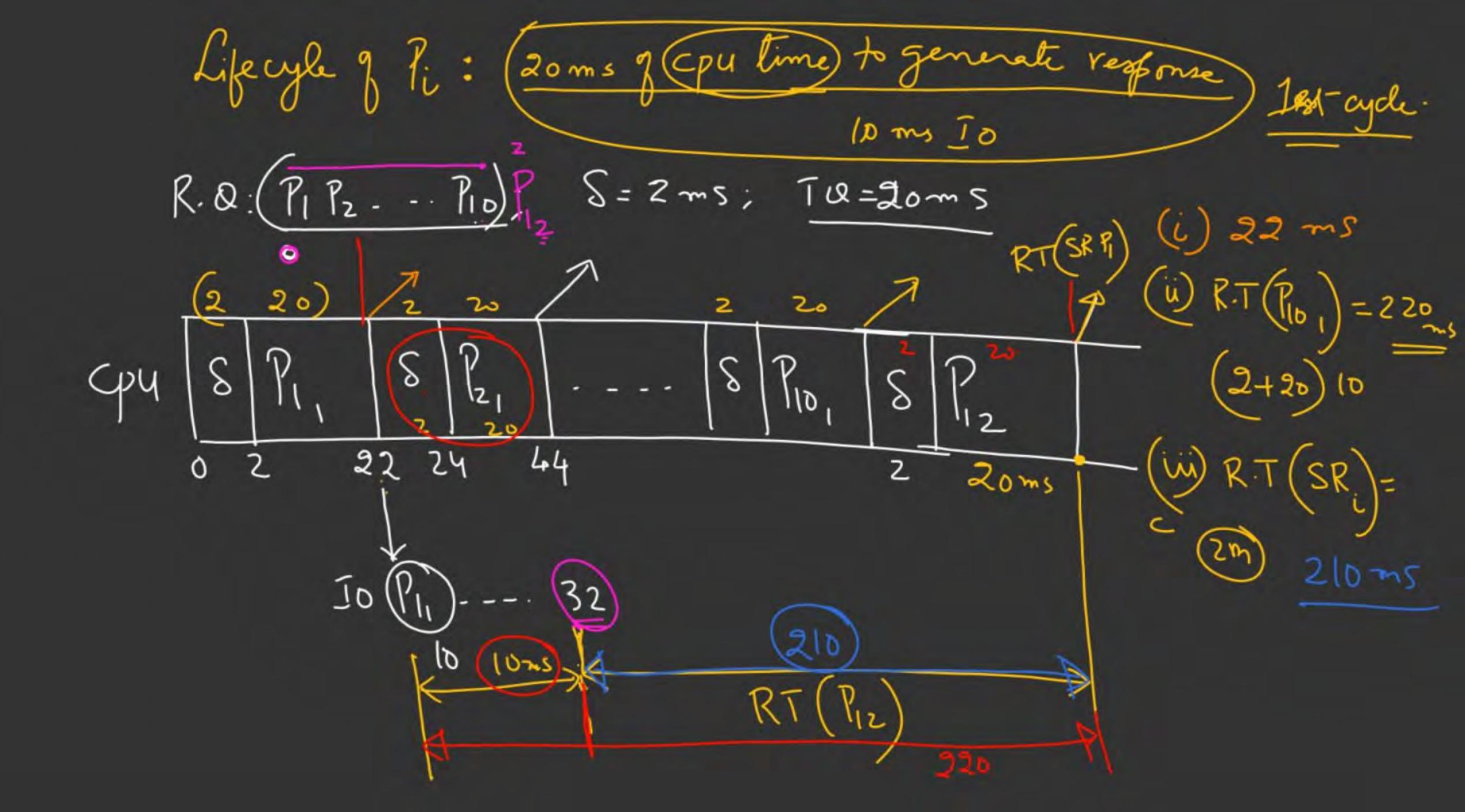




Consider a System using Round Robin Scheduling with 10 (W) Processes all arriving at the time 0. Each Process is associated with 20 identical Request. Each Process request consumes 20 ms of CPU time after which it spends 10 ms of time on I/O, thereafter, initiates subsequent Request. Assuming Scheduling Overhead of 2 ms and Time Quantum of 20 ms, Calculate

/ i. Response time of the 1st request of the 1st Process ii. Response time of the 1st request of the last Process iii. Response time of the subsequent request of any Process.

Assignment: Repeat the Problem
y TQ=10ms

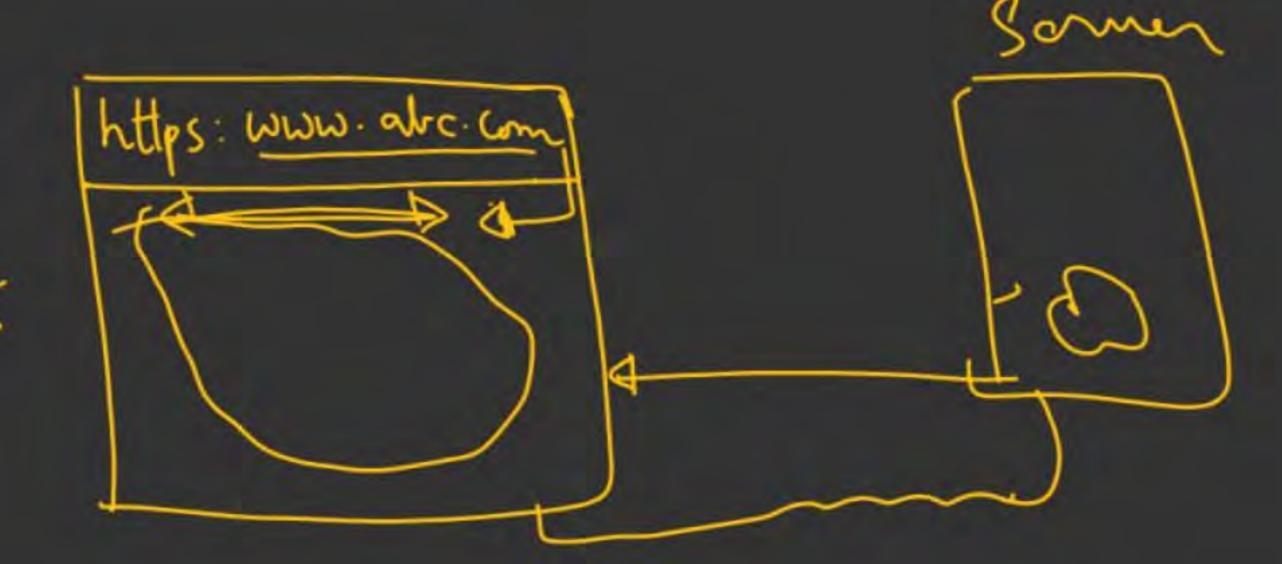


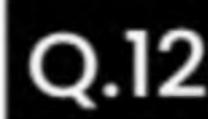
Response Jime: The Jime at which the request is Submitted by the process, to the time at which it generales the terponse (Initial results);

Scenario -1

Scenemo:

Browser







Consider a System using RR Scheduling with TQ of 'Q' seconds & CPU Scheduling overhead is 'S' seconds.

Each Process on an average run for 'T' seconds before blocking on I/O. Give a formula for CPU efficiency for each of the following conditions.

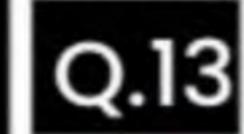
1.
$$Q = \infty$$

2.
$$Q > T$$

3.
$$S < Q < T$$

4.
$$Q = S$$

5.
$$Q \approx 0$$





Consider a System using Preemptive Priority based

scheduling with dynamically changing priorities. On its arrival a Process is assigned a priority of zero and Running Process Priority increases at the rate of ' β' and Priority of the Processes in the ready Q increases at the rate of ' α '. By dynamically changing the values of α and β one can achieve different Scheduling disciplines among the Processes. What discipline will be followed for the following conditions.

1.
$$\beta > \alpha > 0$$

2. $\alpha < \beta < 0$

$$2. \alpha < \beta < 0$$

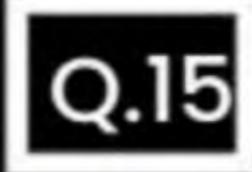


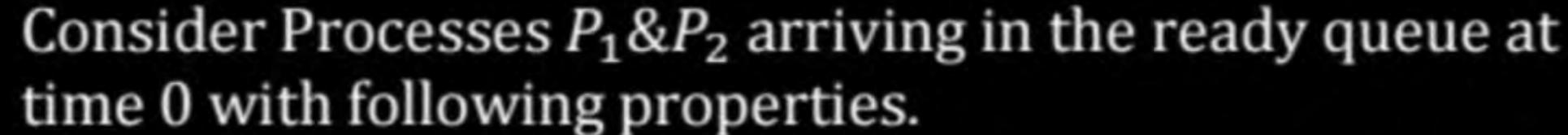




Consider a System using Preemptive Priority based scheduling with dynamically changing priorities. On its arrival a Process is assigned a priority of zero and Running Process Priority increases at the rate of ' \beta ' and Priority of the Processes in the ready Q increases at the rate of ' α '. By dynamically changing the values of α and β one can achieve different Scheduling disciplines among the Processes. What discipline will be followed for the following conditions.

- 1. $\beta > \alpha > 0$ 2. $\alpha < \beta < 0$



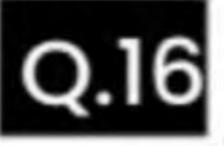




- i) P1 needs a total of 12 units of CPU time and 20 units of I/O time. After every 3 units of CPU time P1 spends 5 units on I/O.
 - ii) P2 needs a total of 15 units of CPU time and no L/O. P2 arrives just after P1.

Compute the Completion times of $P_1 \& P_2$ using the following scheduling techniques:

- 1.SRTF
- 2.Round Robin with Time Quanta = 4 units

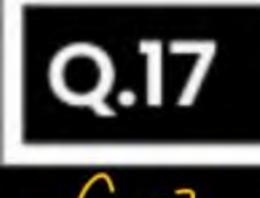


Three processes A, B and C each execute a loop of 100 iterations. In each iteration of the loop, a process performs a single computation that requires to CPU milliseconds and then initiates a single I/O operation that lasts for t_{io} milliseconds. It is assumed that the computer where the processes execute has sufficient number of I/O devices and the OS of the computer assigns different I/O devices to each process. Also the scheduling overhead of the OS is negligible. The processes have the following characteristics:

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	V	W

Process Id	$\mathbf{t}_{\mathbf{c}}$	t_{i0}
A	100 ms	500 ms
В	350 ms	500 ms
C	200 ms	500 ms

The processes A, B, and C are started at times 0,5 and 10 milliseconds respectively in a pure time-sharing system (round robin scheduling) that uses a time slice of 50 milliseconds. The time in milliseconds at which process C would complete its first I/O operation is



Which of the following statements (s) is/are correct in the context of CPU Scheduling?





The goal is to only maximize CPU utilization and minimize X throughput



Α

Turnaround time includes waiting time



Implementing preemptive scheduling needs hardware support



Round-robin policy can be used even when the CPU time required by each of the processes is not known Apriori.

8) Multi-Level Queue Scheduling Starvation: Processes in Lower Q's (having boson Priorits) a lot of Searithing Time; Simple Ready &' Multi-Level (Feedback) Multi-Level Q' Process_types gator. -> All processes will FCMD FUND have to use only one Scheduling FCFS BEND. technique





