

# CS & IT ENGINEERING

## Compiler Design

Syntax Directed Translations

Lecture No. 2



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SDTs

→ Definitions

→ Evaluations

# SDTs Definitions :

- ① L-attributed Grammar
- ② S-attributed Grammar



## L-attributed SDT

- ① Computation depends on parent / left siblings / children.

$S \rightarrow ABC$

$\left\{ \begin{array}{l} A.x = S.y; \\ B.x = A.x + S.x; \\ C.x = A.x + 1; \\ S.x = B.x * 2 \end{array} \right.$

- ② Translation can be anywhere

## S-attributed SDT



- ① computation depends on only children [Synthesized attributes only]

- ② Every translation should appear only at end of production.

$E \rightarrow AB \{E.x = A.x + B.y\}$

Attribute type  $\begin{cases} \rightarrow \text{Inherited} \\ \rightarrow \text{Synthesized} \end{cases}$

SDT type  
(definition)  $\begin{cases} \rightarrow \text{L-attribute Grammar} \\ \rightarrow \text{S-attribute Grammar} \end{cases}$



Find Type/Definition of SDT.



①

$S \rightarrow a \{ S.x = a.val \}$



S-attributed Grammar  
(also, L-attributed)

②

$S \rightarrow Aa \{ A.x = S.y \}$

$S \rightarrow Bb \{ S.y = B.x \}$

$A \rightarrow a \{ A.x = a.val \}$

$B \rightarrow b \{ B.y = b.val \}$



L-attributed  
but not S-attributed  
 $x$  is neither inherited nor synthesized  
 $y$  is synthesized attribute

③

$$E \rightarrow E_1 + E_2 \quad \{ E.val = E_1.val + E_2.val \}$$

$$E \rightarrow a \quad \{ E.val = a.val \}$$

→ S-attributed Grammar  
(L-attributed)

→ val is synthesized attribute



④

$S \rightarrow \boxed{A}B \quad \{ A.x = \underbrace{B.x}_{\text{Right Sibling}} \}$

$A \rightarrow a \quad \{ \quad \}$

$B \rightarrow b \quad \{ B.x = b.val \}$

$\Rightarrow$  Not L-attributed  
SDT

(neither  
L-attributed

nor  
S-attributed SDT)



⑤

$D \rightarrow T L$

$T \rightarrow \text{int}$

$T \rightarrow \text{float}$

$L \rightarrow \text{id}$

$\{ L.type = \overset{\text{Left sibling}}{T.type} \}$

$\{ T.type = \underset{\text{child}}{\text{int}} \}$

$\{ T.type = \underset{\text{child}}{\text{float}} \}$

$\{ \}$

This SDT is  
L-attribute  
but  
not S-attribute

⑥

$S \rightarrow \{S.x = \underbrace{a.val}_{\text{child}}\} a$   
 translation is not at end

$\Rightarrow$  Not S-attributed  
 $\Rightarrow$  It is L-attributed

Attributes not computed

⑦

$E \rightarrow a \{ \text{print}(a.val) \}$

$\Rightarrow$  S-attributed SDT  
 (L-attributed SDT)

⑧

$E \rightarrow \{ \text{print}(a.val) \} a$

$\Rightarrow$  only L-attributed SDT



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$i$  is inherited attribute  
 $s$  is synthesized attribute

$$A \rightarrow P.Q$$

$$\begin{aligned} P.i &= A.i + 2; \\ Q.i &= P.i + A.i; \\ A.s &= \underbrace{P.s + Q.s}_{\text{children}}; \end{aligned}$$

parent

left sibling

parent

only L-attribute Definition

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$i$  is inherited

$$A \rightarrow XY$$

$$\begin{aligned} X.i &= A.i + Y.s, \\ Y.i &= X.s + A.i, \end{aligned}$$

parent

right sibling

left sib

parent

Not 2-attribute (Not S-attribute)



Every S-attributed grammar is L-attributed.  
 Every not L-attributed grammar is not S-attributed

## L-attributed Grammars

### S-attributed Grammars

• ① : ③  
 : ⑦

• ⑧

• ②

• ⑤

• ⑥

• ⑨

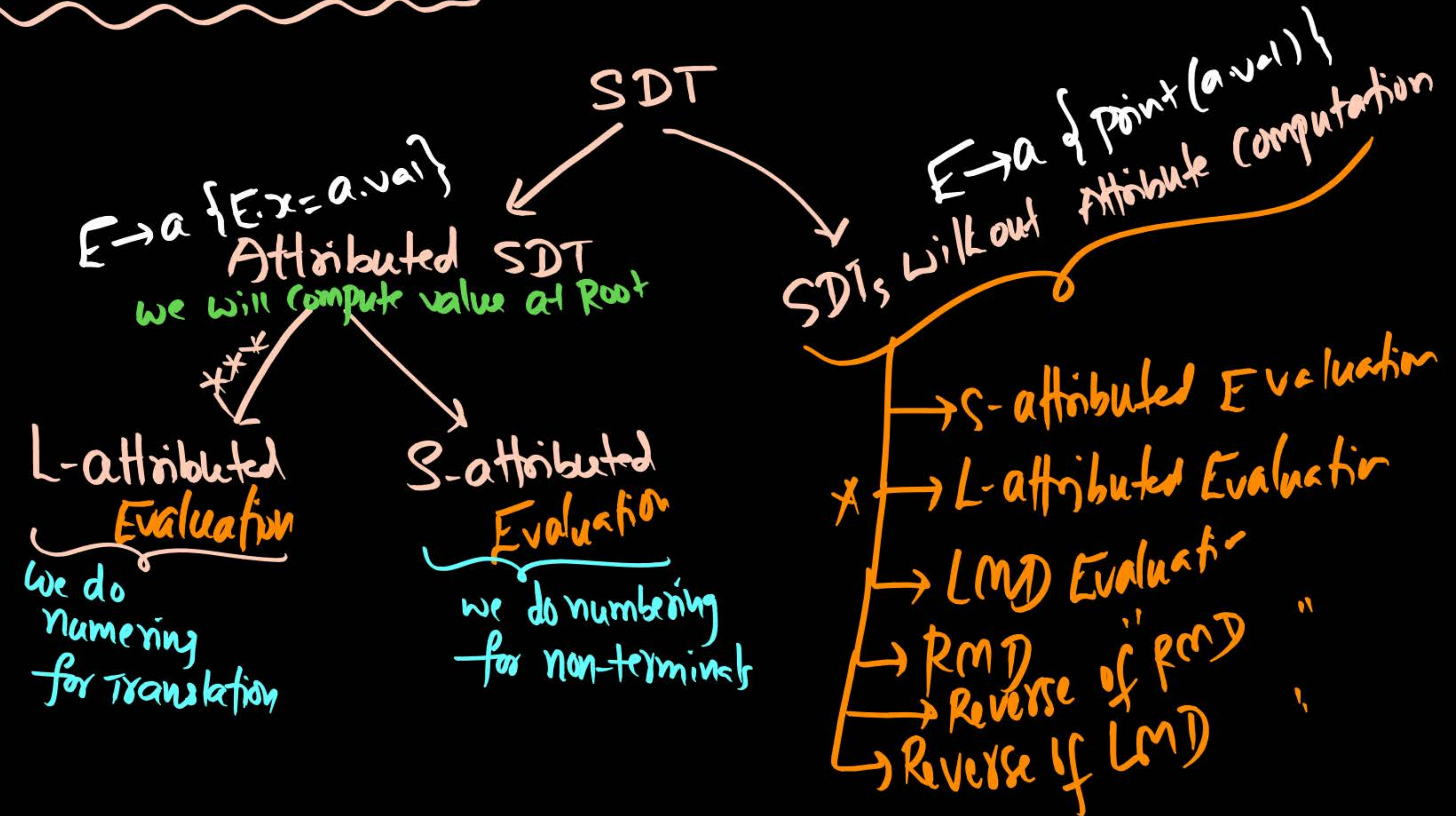
• ④

Not L-attributed  
 SDTs

• ⑩



# Evaluations of SDT:





①  $E \rightarrow E_1 + T \{ E.val = E_1.val * T.val \}$

$E \rightarrow a \{ E.val = a.val \}$

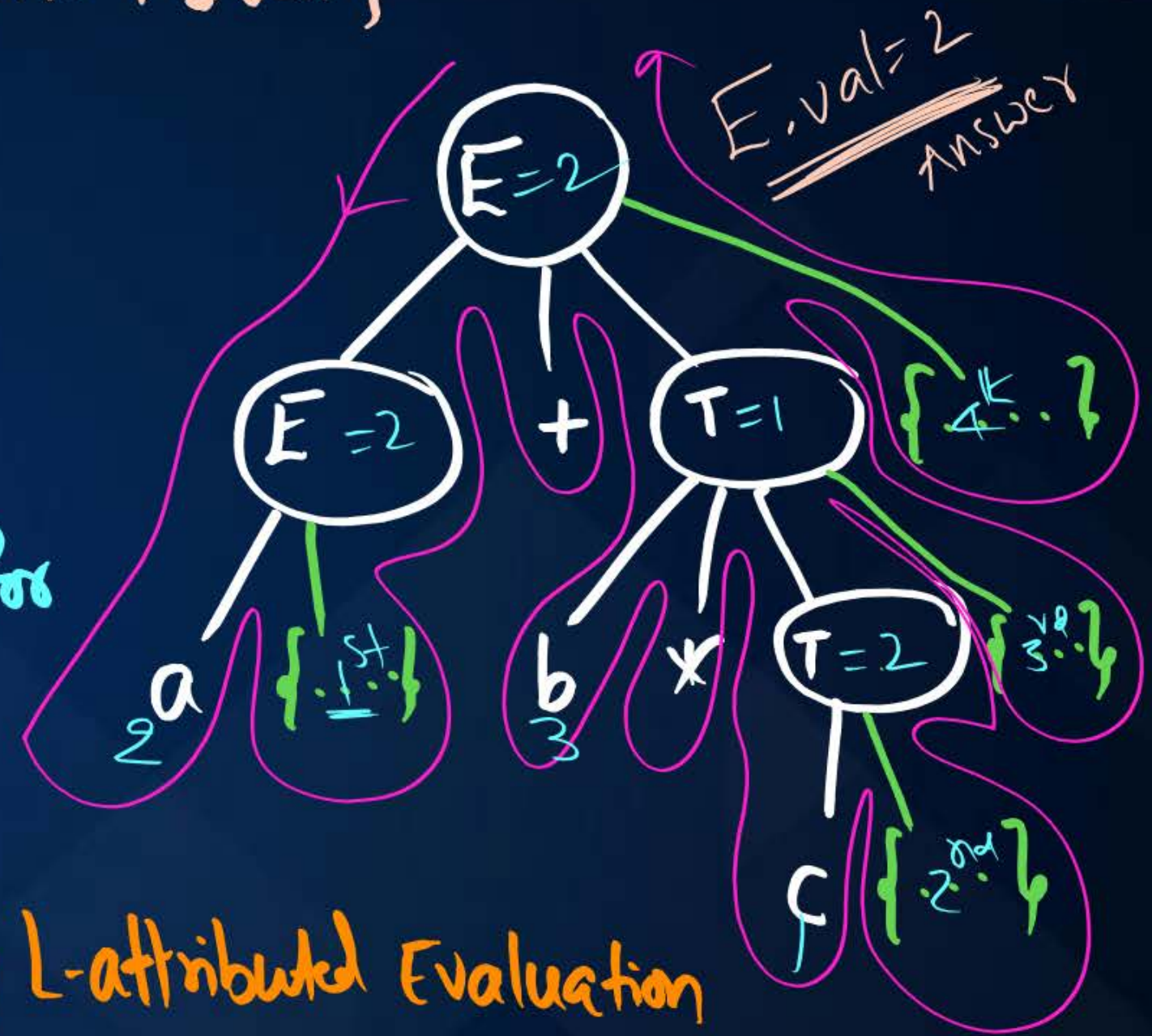
$T \rightarrow b * T_1 \{ T.val = b - T_1.val \}$

$T \rightarrow c \{ T.val = c.val + 1 \}$

Find attribute value at the Root for  
input  $2 + 3 * 1$

Note: This SDT is both L-attribute & S-attribute.

Answer = 2



Using L-attribute Evaluation

Depth Left and Follow Inorder

(All translations evaluate from Left to Right)

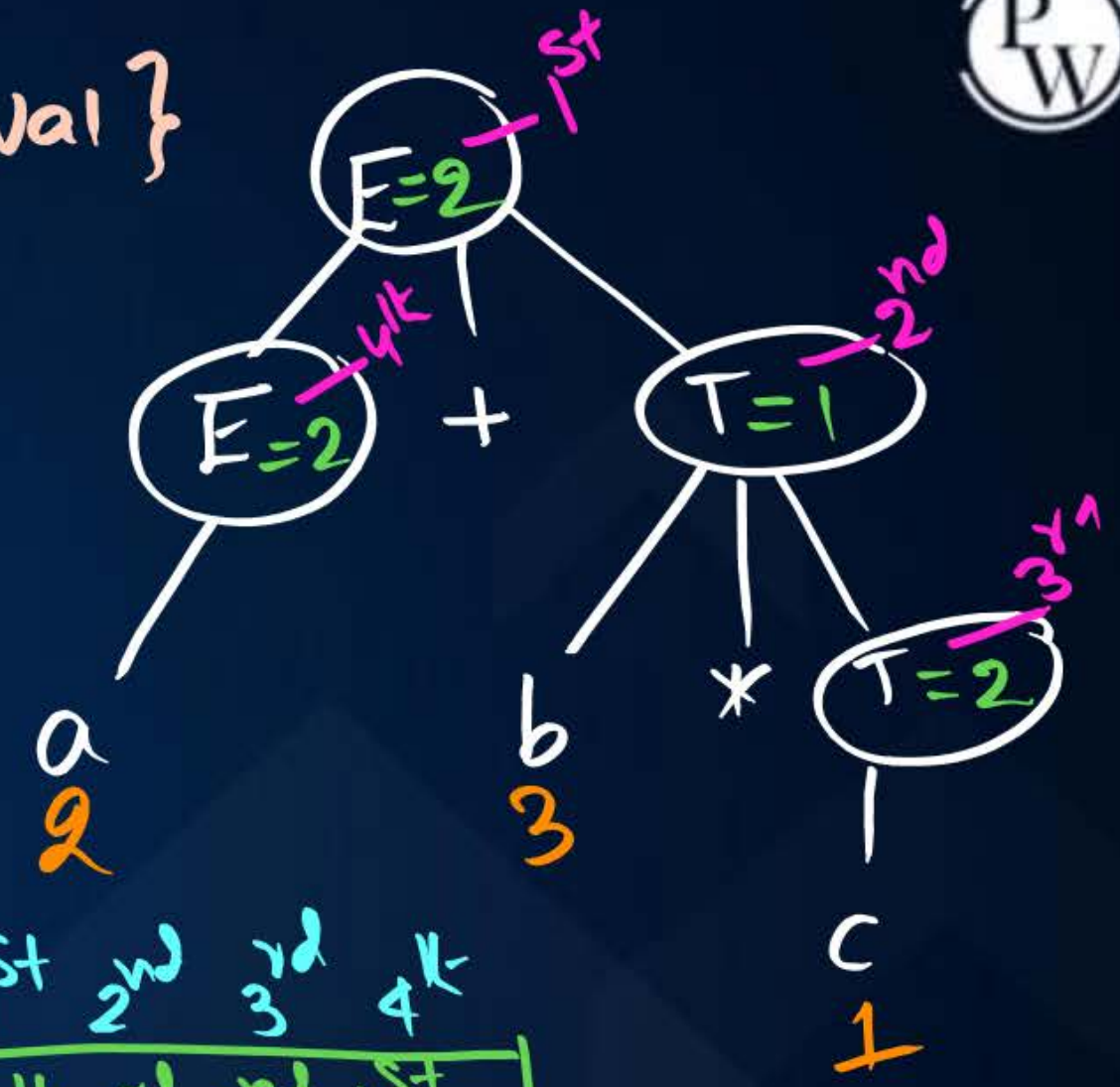


①  $E \rightarrow E_1 + T \{ E.val = E_1.val + T.val \}$

$E \rightarrow a \{ E.val = a.val \}$

$T \rightarrow b * T_1 \{ T.val = b + T_1.val \}$

$T \rightarrow c \{ T.val = c.val + 1 \}$



Find attribute value at the Root for

input  $2 + 3 * 1$ :

RMD order: 1<sup>st</sup> 2<sup>nd</sup> 3<sup>rd</sup> 4<sup>th</sup>

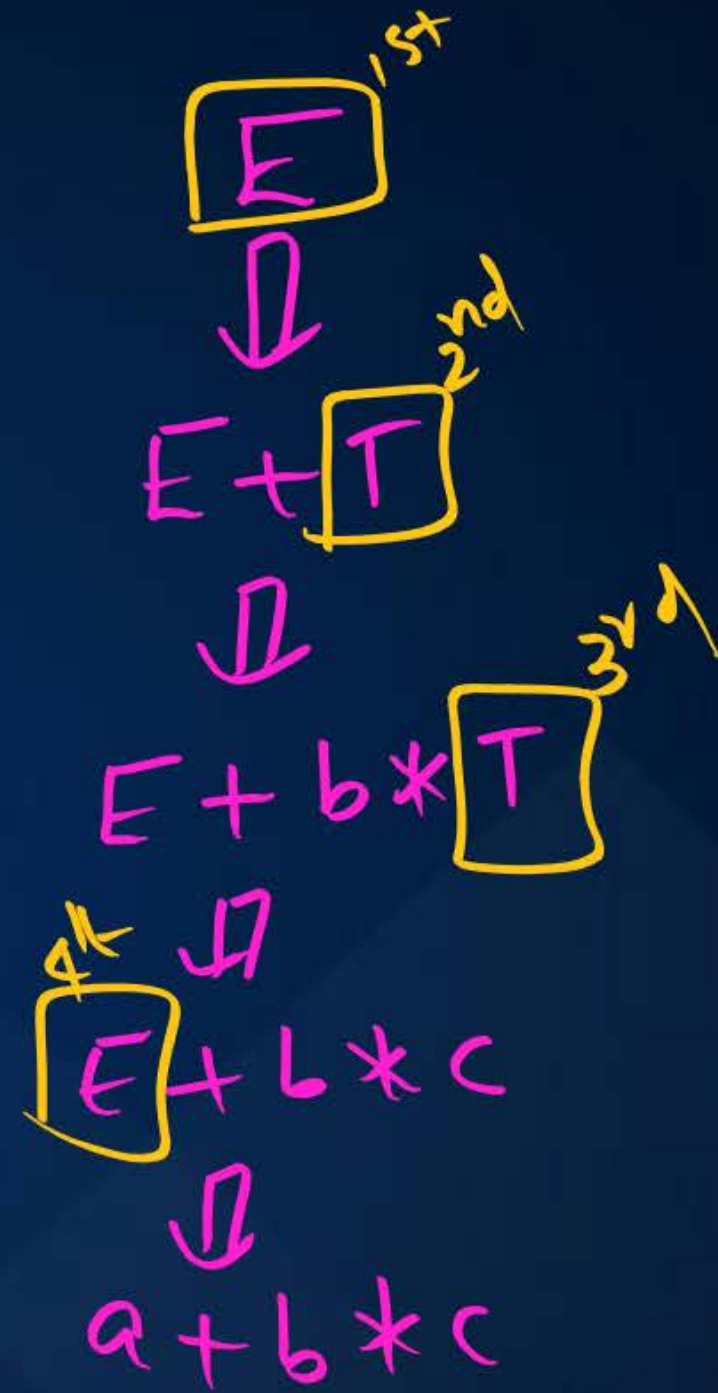
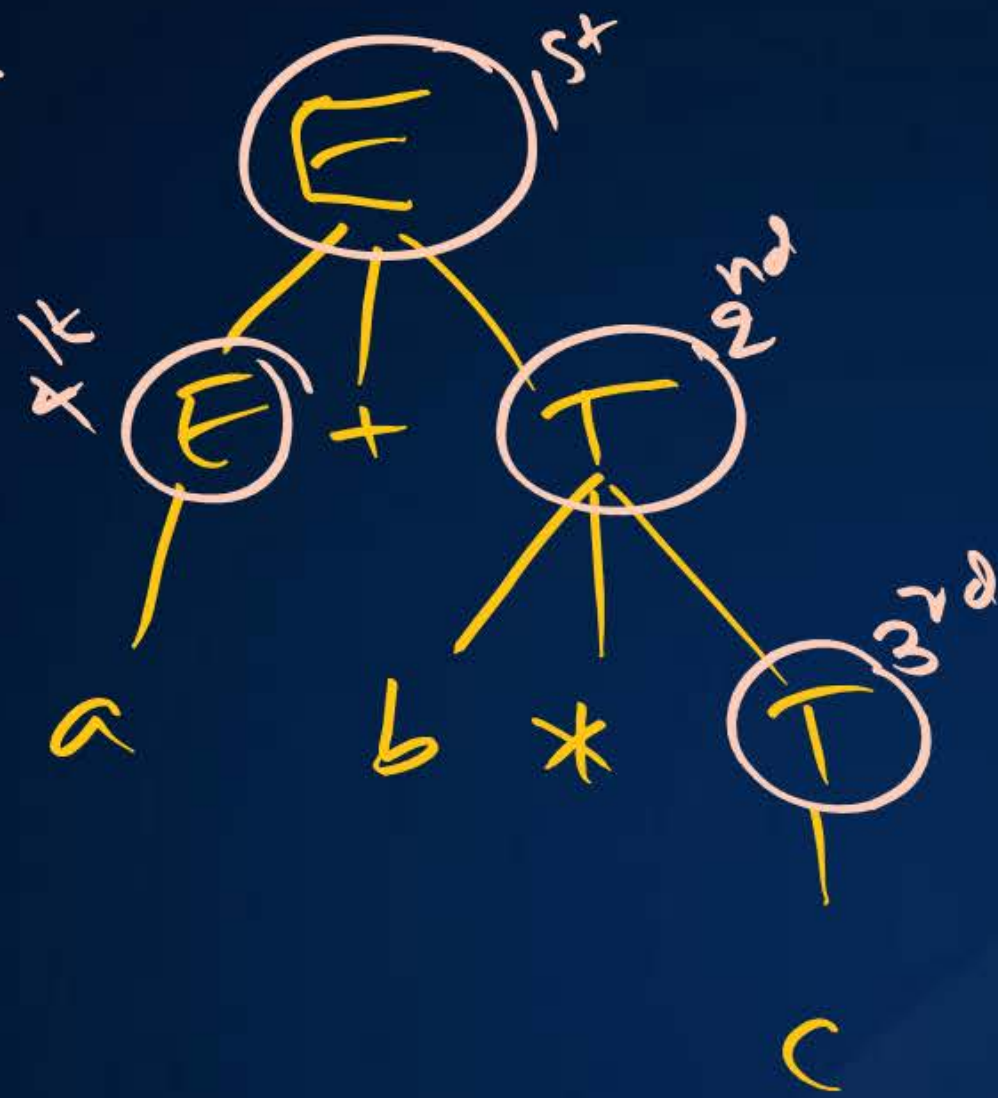
Reverse of RMD: 4<sup>th</sup> 3<sup>rd</sup> 2<sup>nd</sup> 1<sup>st</sup>

Note: This SDT is both L-attribute & S-attribute.

↓  
Answer = 2

⇒ Using S-attributed Evaluation  
(Bottom up parsing Approach)  
(Reverse of RMD)

AND numbers:





Note:  
 $\Rightarrow$

If no evaluation technique mentioned then

'Use L-attributed'

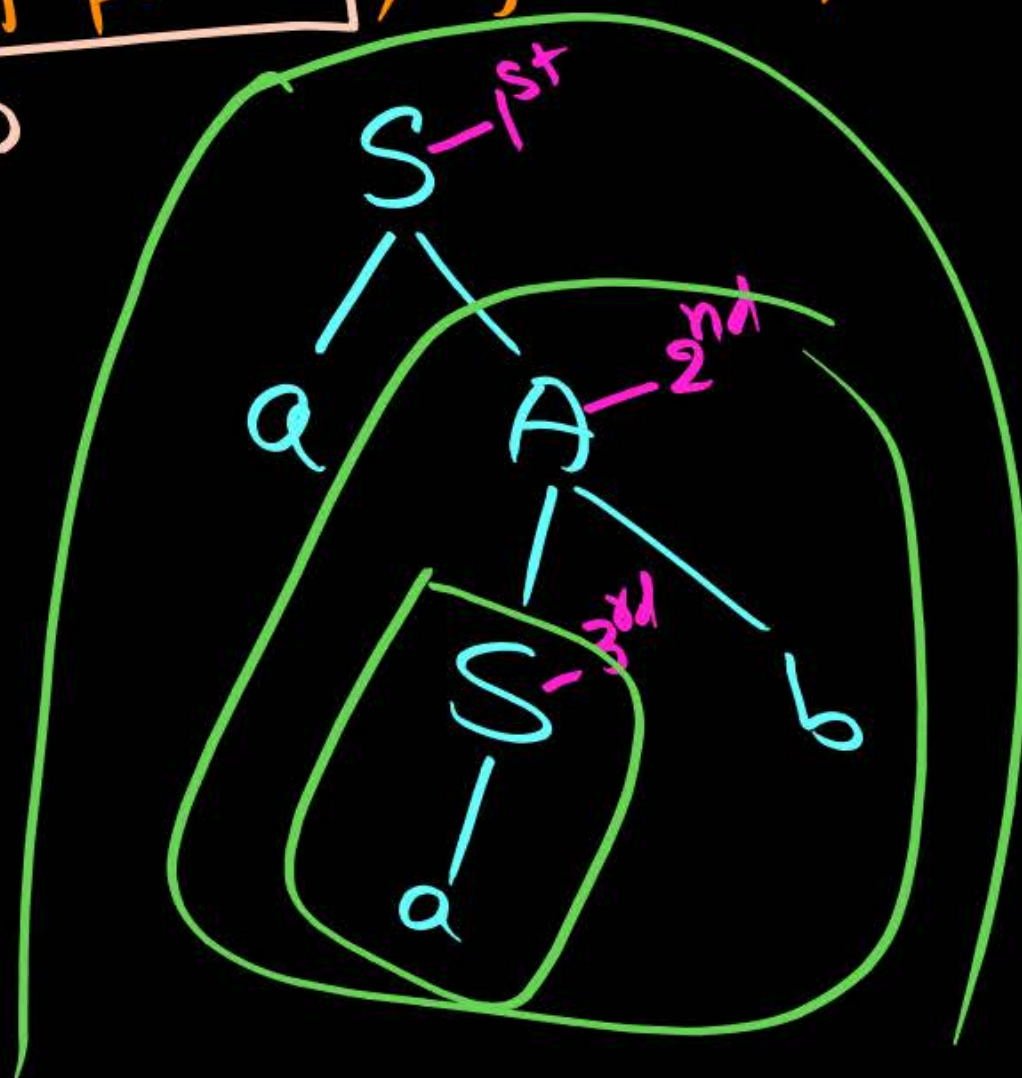


②

$S \rightarrow aA$  {print 1}  
 $S \rightarrow a$  {print 2}  
 $A \rightarrow Sb$  {print 3}

Using Bottom up parser, find o/p for "aab".

Reverse of RMD



RMD order: 1<sup>st</sup> 2<sup>nd</sup> 3<sup>rd</sup>

Reverse of RMD: 3<sup>rd</sup> 2<sup>nd</sup> 1<sup>st</sup>

o/p: 2 3 1



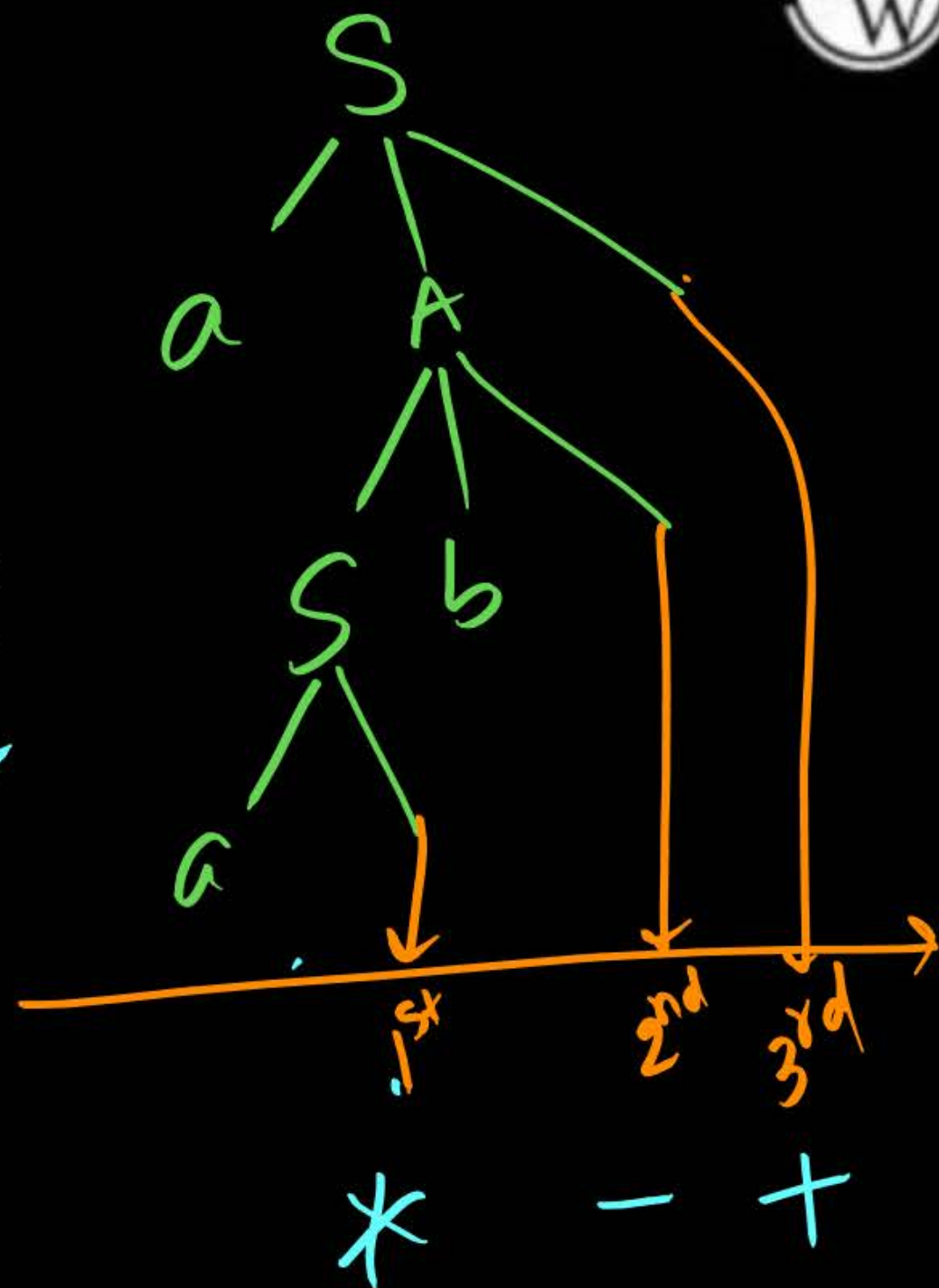
③

$S \rightarrow aA$  {print '+'}  
 $S \rightarrow a$  {print '\*'}  
 $A \rightarrow Sb$  {print '-'}

Find o/p for aab using

- I) Not mentioned any evaluation
- II) Using l-attributed evaluation
- III) Using S-attributed / Bottom up parsing / LR parser / SR parser / LALR / SLR / CLR
- IV) Using top-down parsing / LL(1) / predictive parsing
- V) Using RMD
- VI) Using Reverse of LMD

same o/p  
\* - +





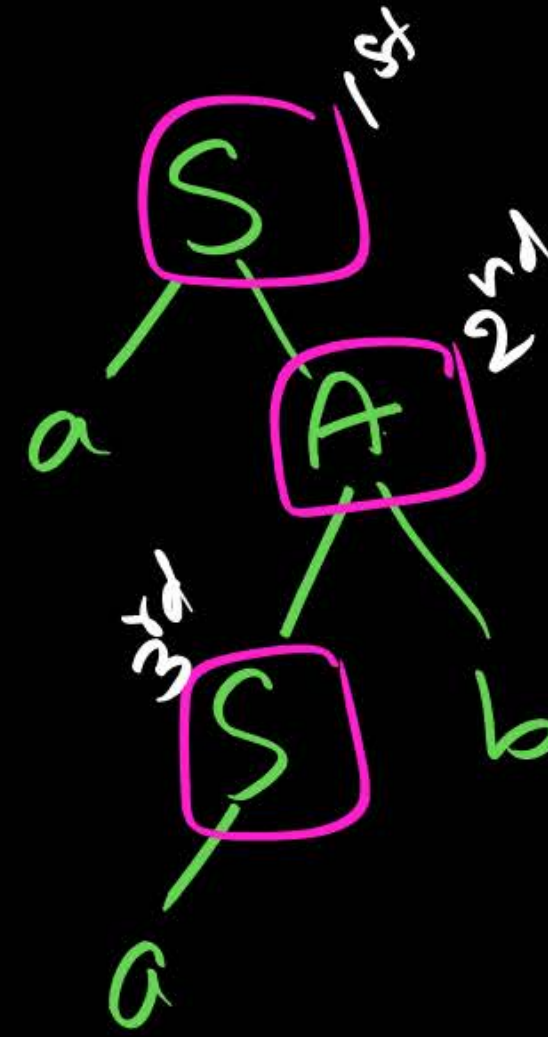
③

$S \rightarrow aA$  {print '+'}  
 $S \rightarrow a$  {print '\*'}  
 $A \rightarrow Sb$  {print '-'}

Find o/p for aab using

- I) Not mentioned any evaluation
- II) Using l-attributed evaluation
- ☒ III) Using S-attributed / Bottom up parsing / LR parser / SR parser / LALR / SLR / CLR
- IV) Using top-down parsing / LL(1) / predictive parsing
- ☒ V) Using RMD  $\Rightarrow + - *$
- VI) Using Reverse of LMD

same o/p  $\rightarrow * - +$



RMD: 1<sup>st</sup> 2<sup>nd</sup> 3<sup>rd</sup>  
 + - \*

Reverse of RMD: \* - +

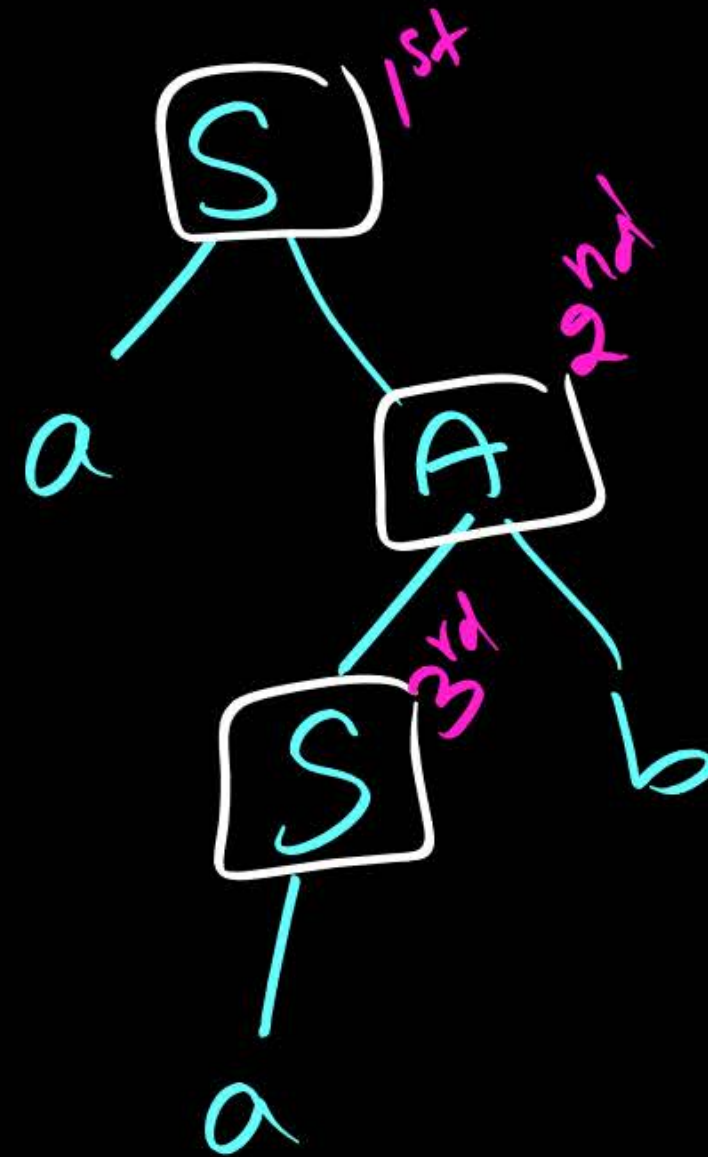


③

$S \rightarrow aA$  {print '+'}  
 $S \rightarrow a$  {print '\*'}  
 $A \rightarrow Sb$  {print '-'}

Find o/p for aab using

- I) Not mentioned any evaluation.
- II) Using l-attributed evaluation
- III) Using S-attributed / Bottom up parsing / LR parser / SR parser / LALR / SLR / CLR
- ~~IV)~~ Using TOP-down parsing / LL(1) / predictive parsing
- V) Using RMD
- ~~VI)~~ Using Reverse of LMD  $\Rightarrow * - +$



LMD: 

1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>
+	-	*

Reverse of LMD: \* - +



④

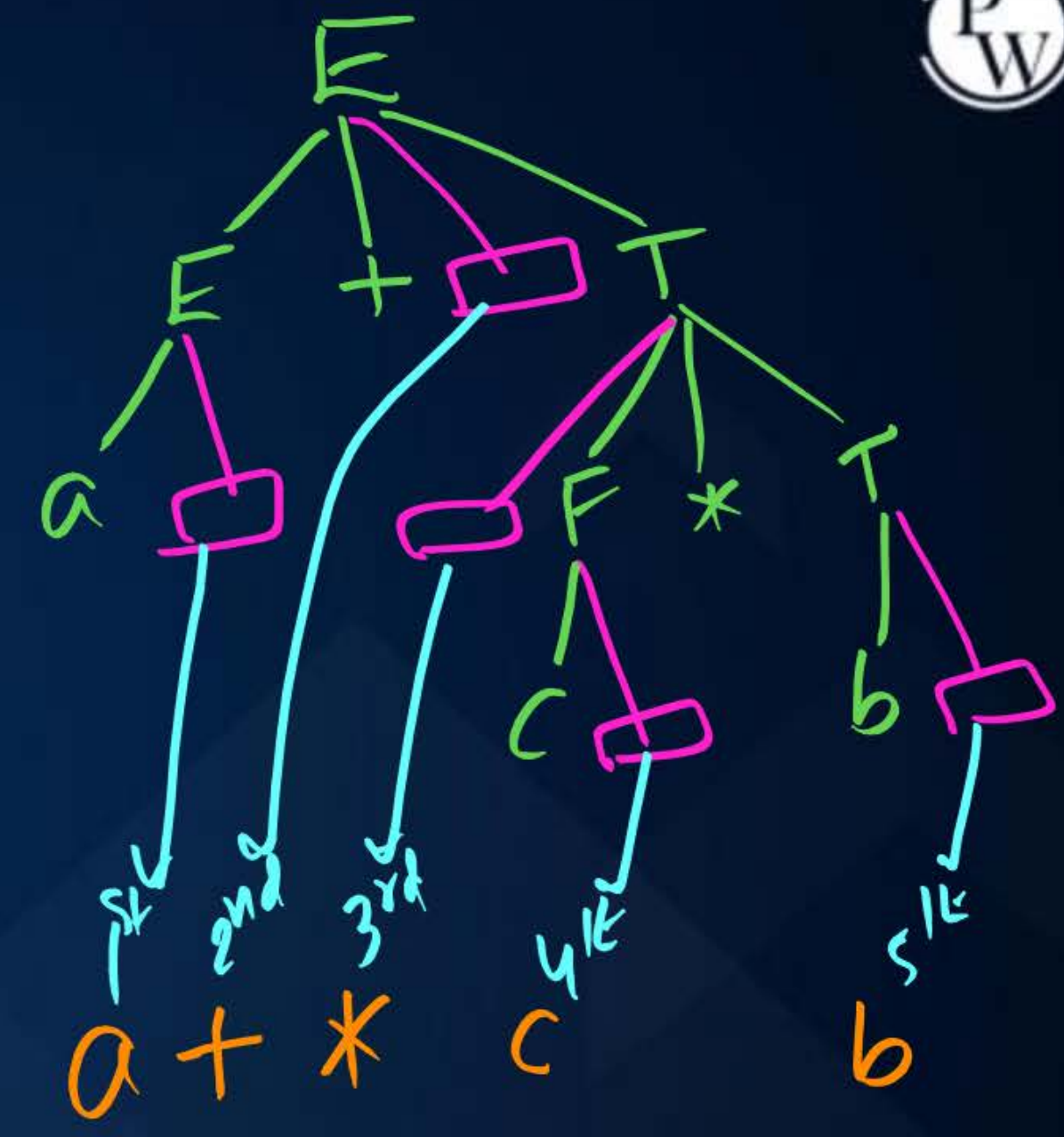
$$E \rightarrow E + \{ \text{print } + \} T$$

$$E \rightarrow a \{ \text{print } a \}$$

$$T \rightarrow \{ \text{print } * \} F * T$$

$$T \rightarrow b \{ \text{print } b \}$$

$$F \rightarrow c \{ \text{print } c \}$$



Find o/p for Input =  $a + c * b$  using:

- \* I) L-attributed evaluation  $\Rightarrow a + * c b$
- II) Bottom up parsing
- III) TOP-down parsing



④

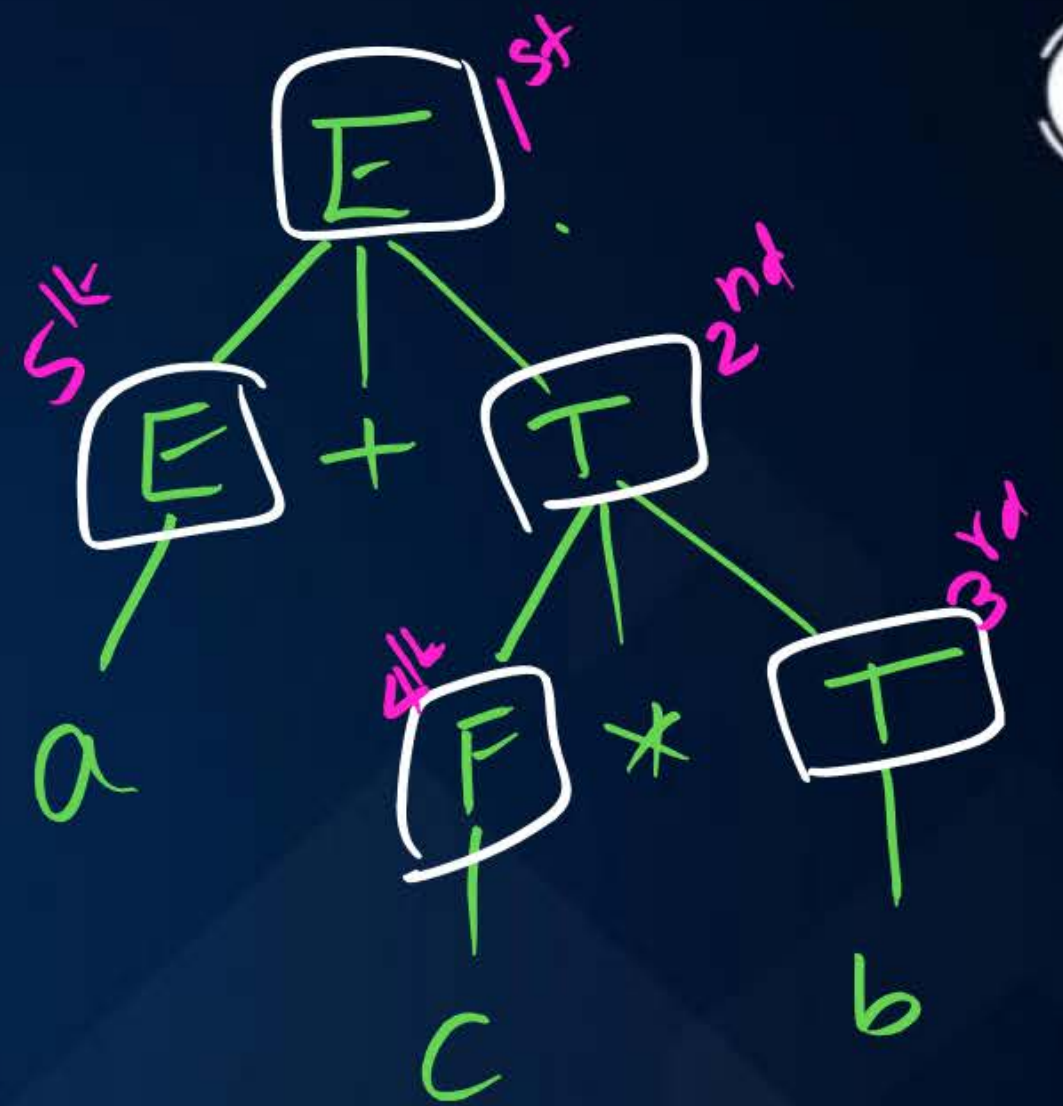
$$E \rightarrow E + \{ \text{print } + \} T$$

$$E \rightarrow a \{ \text{print } a \}$$

$$T \rightarrow \{ \text{print } * \} F * T$$

$$T \rightarrow b \{ \text{print } b \}$$

$$F \rightarrow c \{ \text{print } c \}$$



Find o/p for Input =  $a + c * b$ . using:

- \* I) L-attributed evaluation
- II) Bottom up parsing  $\Rightarrow a c b * +$
- III) TOP-down parsing

Reverse of RMD: 5<sup>th</sup> 4<sup>th</sup> 3<sup>rd</sup> 2<sup>nd</sup> 1<sup>st</sup>  
 $a \ c \ b \ * \ +$

④

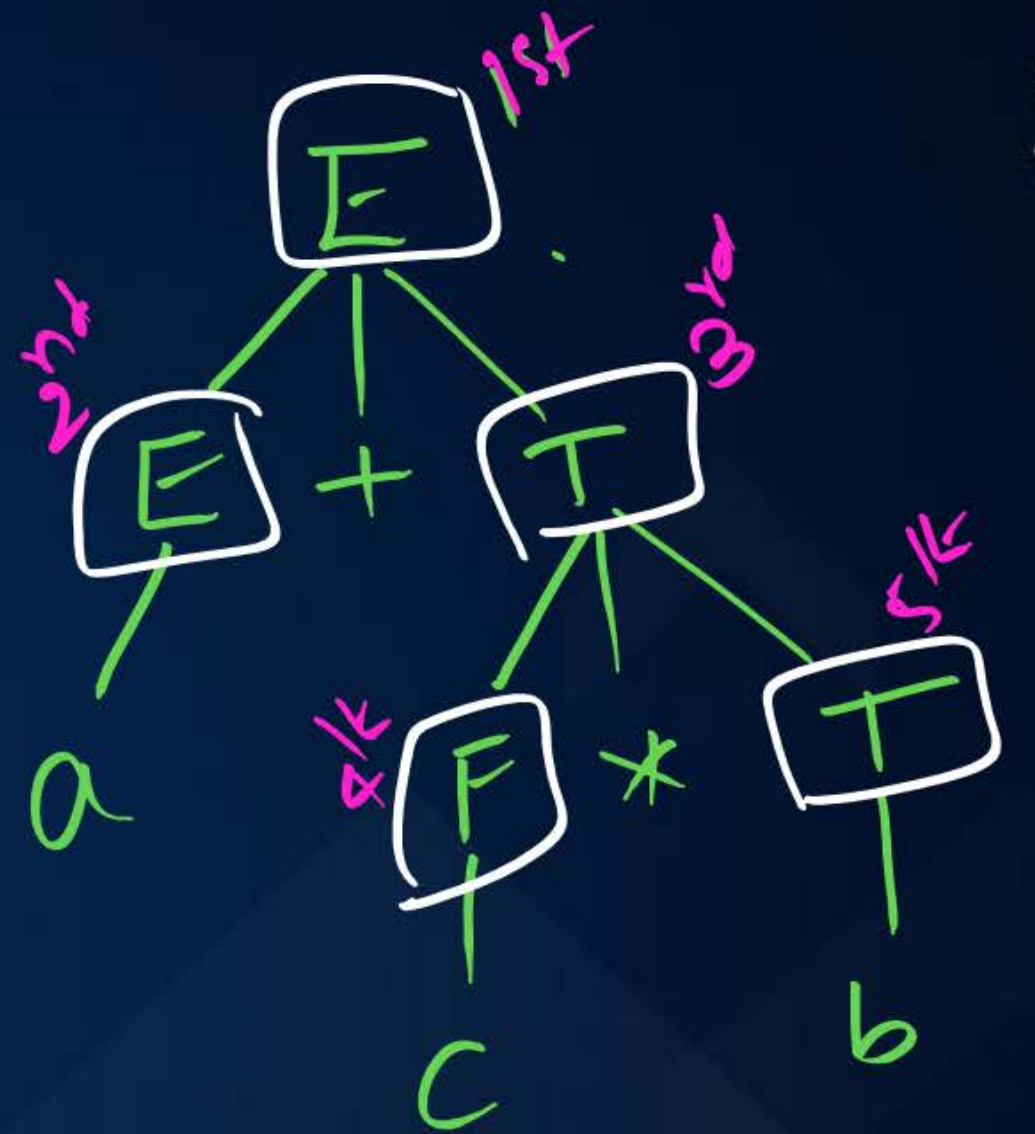
$$E \rightarrow E + \{ \text{print } + \} T$$

$$E \rightarrow a \{ \text{print } a \}$$

$$T \rightarrow \{ \text{print } * \} F * T$$

$$T \rightarrow b \{ \text{print } b \}$$

$$F \rightarrow c \{ \text{print } c \}$$



Find o/p for Input =  $a + c * b$ . using:

\* I) L-attributed evaluation

II) Bottom up parsing

III) TOP-down parsing  $\Rightarrow + a * c b$

LMD: 1<sup>st</sup> 2<sup>nd</sup> 3<sup>rd</sup> 4<sup>th</sup> 5<sup>th</sup>  
 + a \* c b



5  
H.W.

$$E \rightarrow \{ \text{print } + \} E + T$$

$$E \rightarrow \{ \text{print } a \} a$$

$$T \rightarrow \{ \text{print } * \} F * T$$

$$T \rightarrow b \{ \text{print } b \}$$

$$F \rightarrow c \{ \text{print } c \}$$

Find o/p for Input =  $a + c * b$ . using:

- \* I) L-attributed evaluation
- II) Bottom up parsing
- III) TOP-down parsing

