

CS & IT ENGINEERING

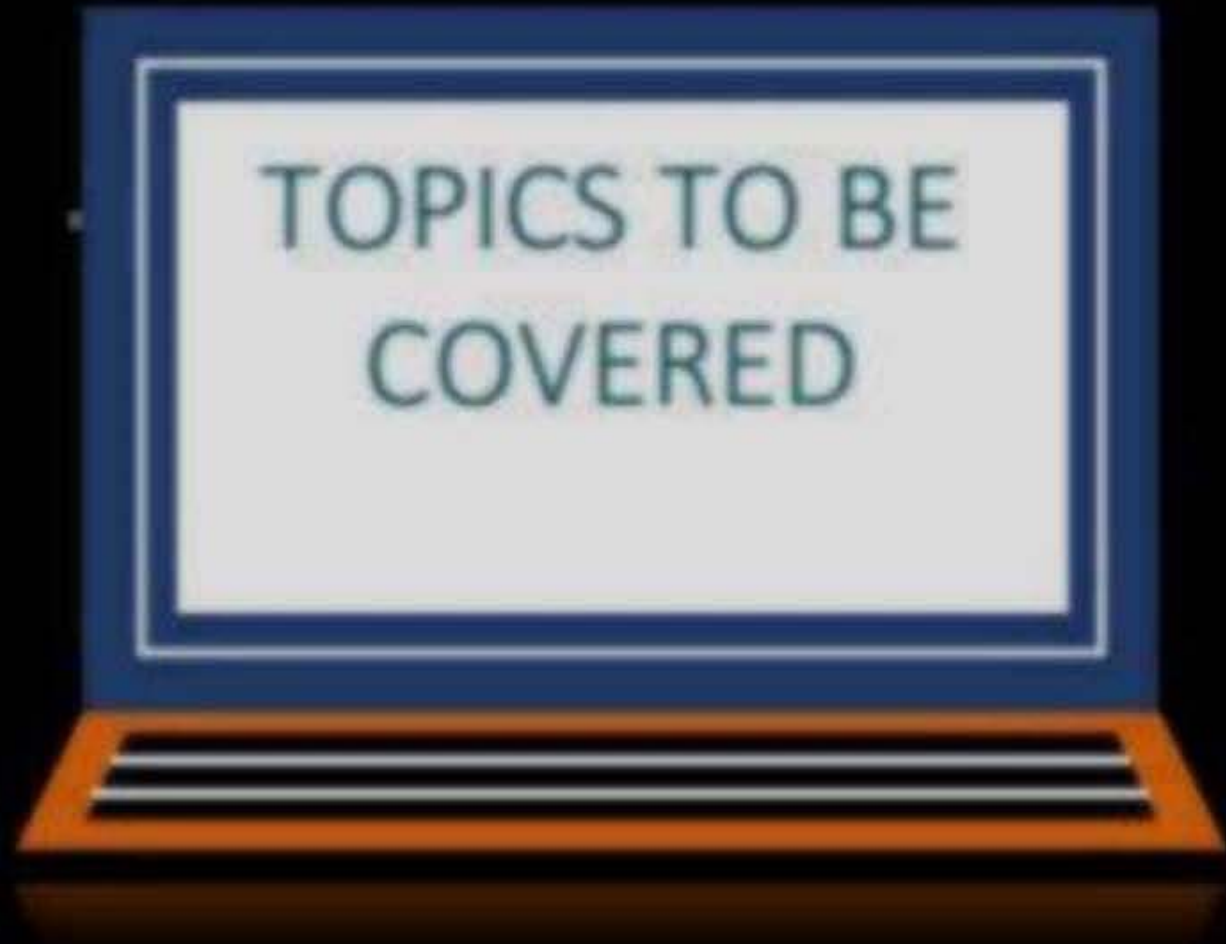
Compiler Design

Intermediate code and code optimization

Lecture No. 3



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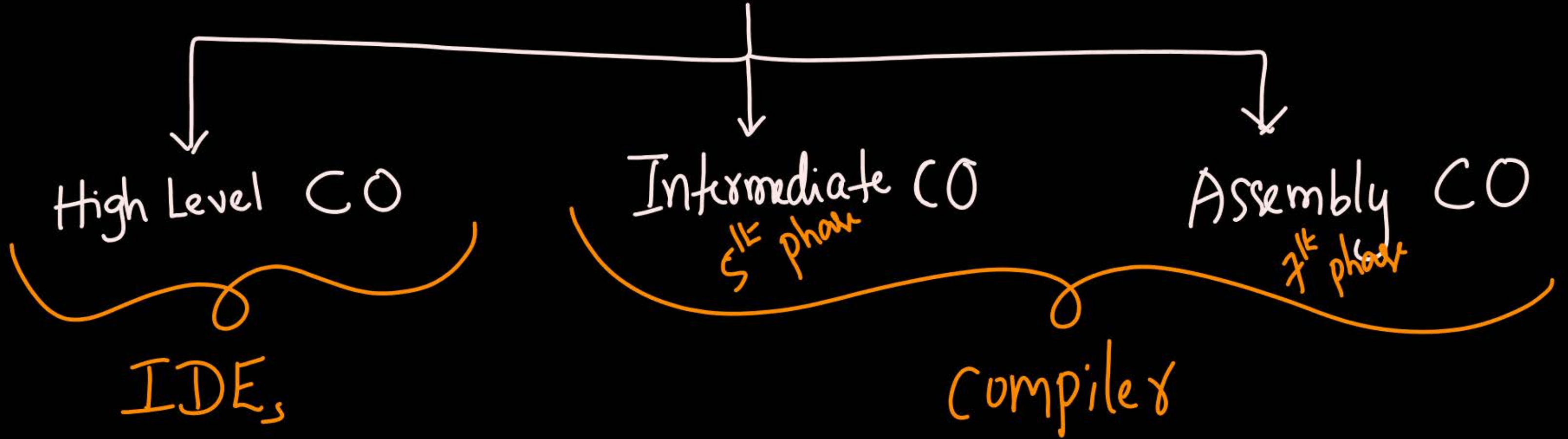
Code Optimization [CO]

- What is CO?
- CO Techniques
- Data Flow Analysis
 - Live Variable Analysis
 - Reaching Definition Analysis

Code Optimization

↳ It may save Space / Time

Code Optimization



Code Optimization

Local Optimizations

Global Optimizations

- Statement Level
- Basic Block level
- Loop Level
- Intra-procedural Level
- Inter-procedural Level

Code Optimization Techniques :



③ Common Sub expressions elimination [using DAG]

④ Dead code elimination

- Code motion
- Induction variables elimination

⑤ Loop Optimizations

- loop merging
- loop Unrolling

① Constant Folding



$$x = 2 * 3 + y$$

folding



$$x = 6 + y$$

$$x = 2 + 3 * y$$

Constant
Folding
not possible

②

Identity Simplification:

$$X = Y + 0 + Z$$



$$\underbrace{Y + 0 = Y}$$

Identity Law

$$X = Y + Z$$

$$X = \overbrace{Y * 1} + Z$$



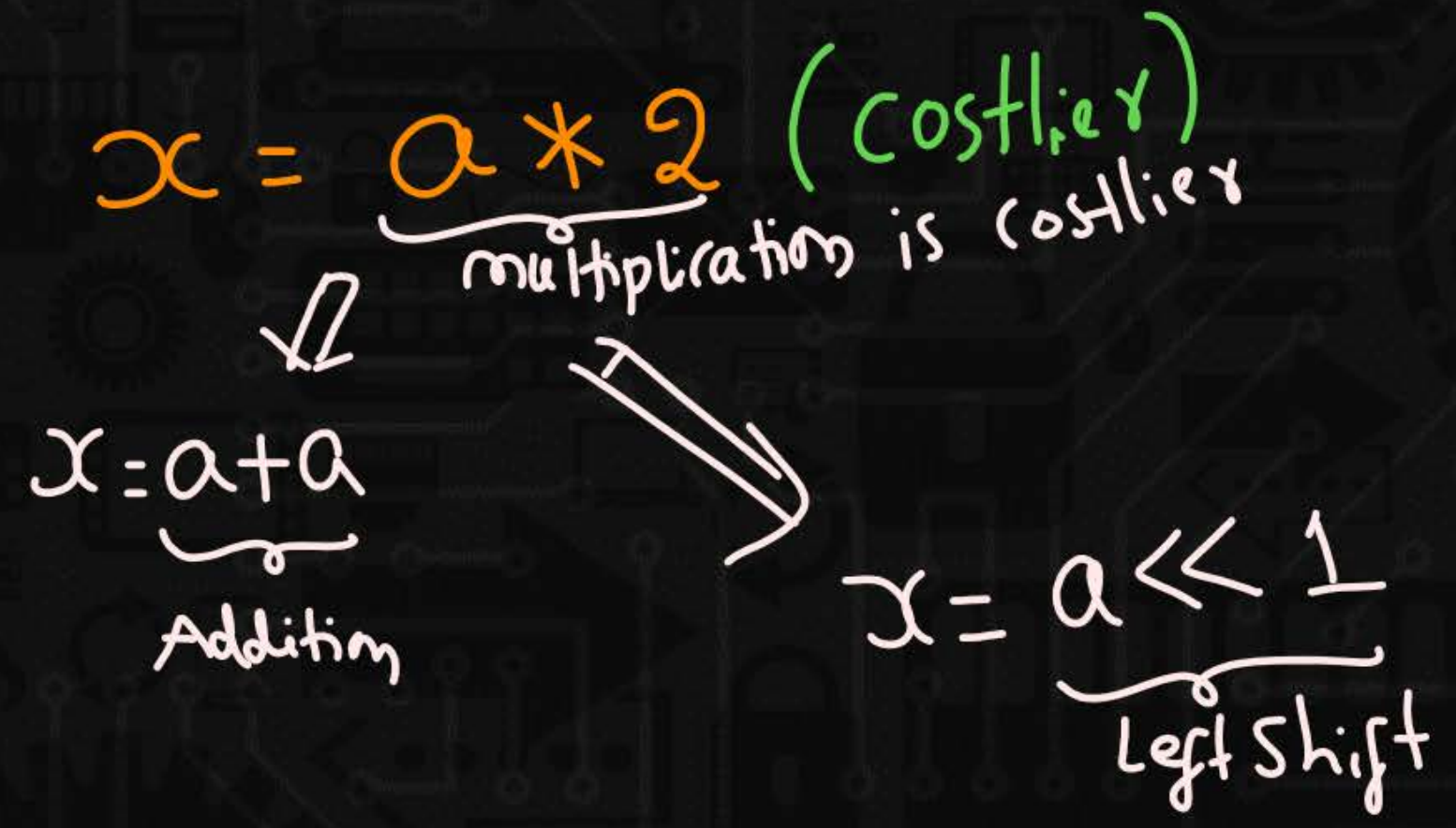
$$Y * 1 = Y$$

Identity Law *

$$X = Y + Z$$

***③ Strength Reduction

↳ "costlier" Instruction Can be replaced with "cheaper"



$$x = a * 8$$



$$x = a \ll 3$$

$$x = a / 8$$



$$x = a \gg 3$$

④ Cancellation



$$x = \cancel{a} + b * c - \cancel{a}$$



$$x = b * c$$

*** (5)

Copy propagation

→ We need DATA Flow Analysis (Live variable Analysis)

$$x = 5$$

$$y = x * a$$

$$z = x + y$$

↙ constant propagation

$$\begin{aligned} y &= 5 * a \\ z &= 5 + y \end{aligned}$$

$$x = b$$

$$y = x * a$$

$$z = x + y$$

↙ variable propagation

$$\begin{aligned} y &= b * a \\ z &= b + y \end{aligned}$$

***⑥ Common Sub-expression Elimination



- we can use DAG to eliminate common sub-exps
- we can use Available Expression Analysis (Data Flow Analysis)

$$x = (a+b) * (a+b)$$

$$\begin{aligned} t_1 &= a+b \\ x &= t_1 * t_1 \end{aligned}$$

⑦

Dead Code Elimination



↳ We can do live variable analysis
(Data Flow Analysis)

```
x = a + b  
y = a * c  
z = x + c  
print(z)
```

Dead code



```
x = a + b  
z = x + c  
print(z)
```


⑧ Loop optimizations



i) Code Motion

↳ Find loop invariant code and move outside loop

$n=0 \Rightarrow 2$ statements

$n=1 \Rightarrow 6$ "

$n=2 \Rightarrow 10$ "

$4n+2$ statements

```
for (i=0; i<n; i++)
```

```
{
```

```
    x = y;
```

loop Invariant Code

```
    a = a + i;
```

```
}
```

if (n > 0)

```
    x = y;
```

```
    for (i=0; i<n; i++) { a = a + i; }
```

$3n+3$

Statements

ii) Induction Variables Elimination

```
for(i=0; i<n; i++)  
{  
    x = i+a;  
    y = i*b;  
    z = i-c;  
}  
  
for (i=0, j=0; i<n; i++)  
{  
    K=0;  
    x = i+a;  
    y = j*b;  
    z = K-c;  
    K++;  
    j++;  
}
```

j and K are deleted

<u>Variables</u>	<u>Induction Variables</u>
i j K x y z a b c	i j K x y z

iii) Loop Merge/Loop Combine/Loop fusion

```
for (i=0; i<n; i++)  
{  
    A[i] = i+1;  
}  
for (j=0; j<n; j++)  
{  
    B[j] = j*5;  
}
```

⇒

```
for (i=0; i<n; i++)  
{  
    A[i] = i+1;  
    B[i] = i*5;  
}
```


iv) Loop Unrolling

```
for (i = 1; i <= 4n; i++)  
{  
    printf("gate");  
}
```

```
for (i = 1; i <= 2n; i++)  
{  
    printf("gate");  
    printf("gate");  
}
```

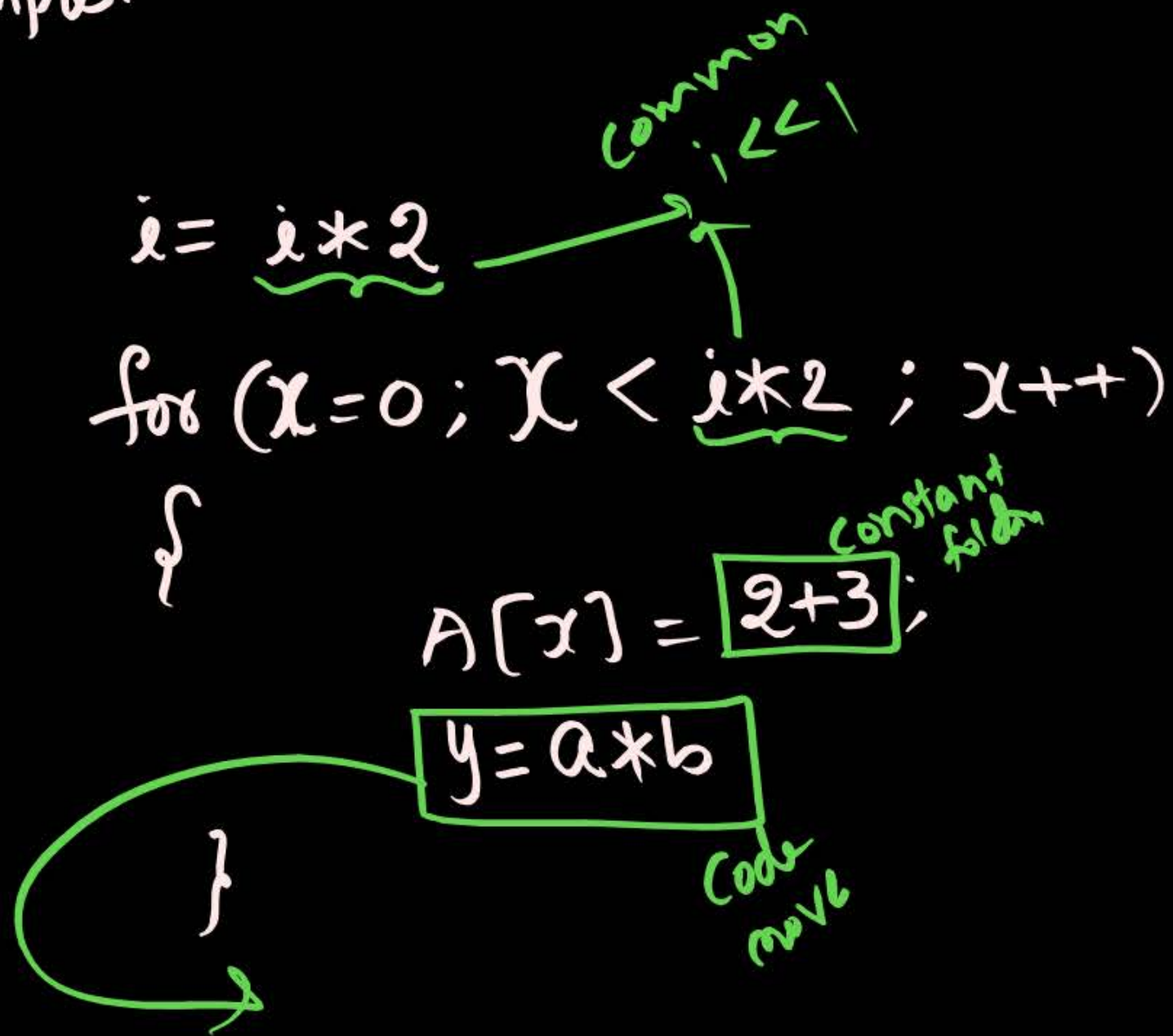
Double

Half

↓ one more time
loop unroll

```
for (i = 1; i <= n; i++)  
{  
    printf("gate");  
    printf("gate");  
    printf("gate");  
    printf("gate");  
}
```

Example:



- ☒ A) Code Motion
- ☒ B) Strength Reduction
- ☒ C) Constant Folding
- ☐ D) Copy propagation
- ☒ E) Common Sub exp elimination



Example:

After propagation, this code is dead code

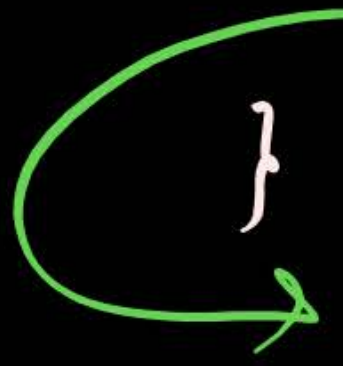
$i = 10;$

for ($x = 0; x < 10 * 2; x++$)
{

$A[x] = a;$

$y = a * b$

code move



☒ A) Code Motion

☐ B) Strength Reduction

☒ C) Constant Folding

☒ D) Copy propagation

☐ E) Common Subexp elimination

☒ F) Dead code Elimination



Data Flow Analysis



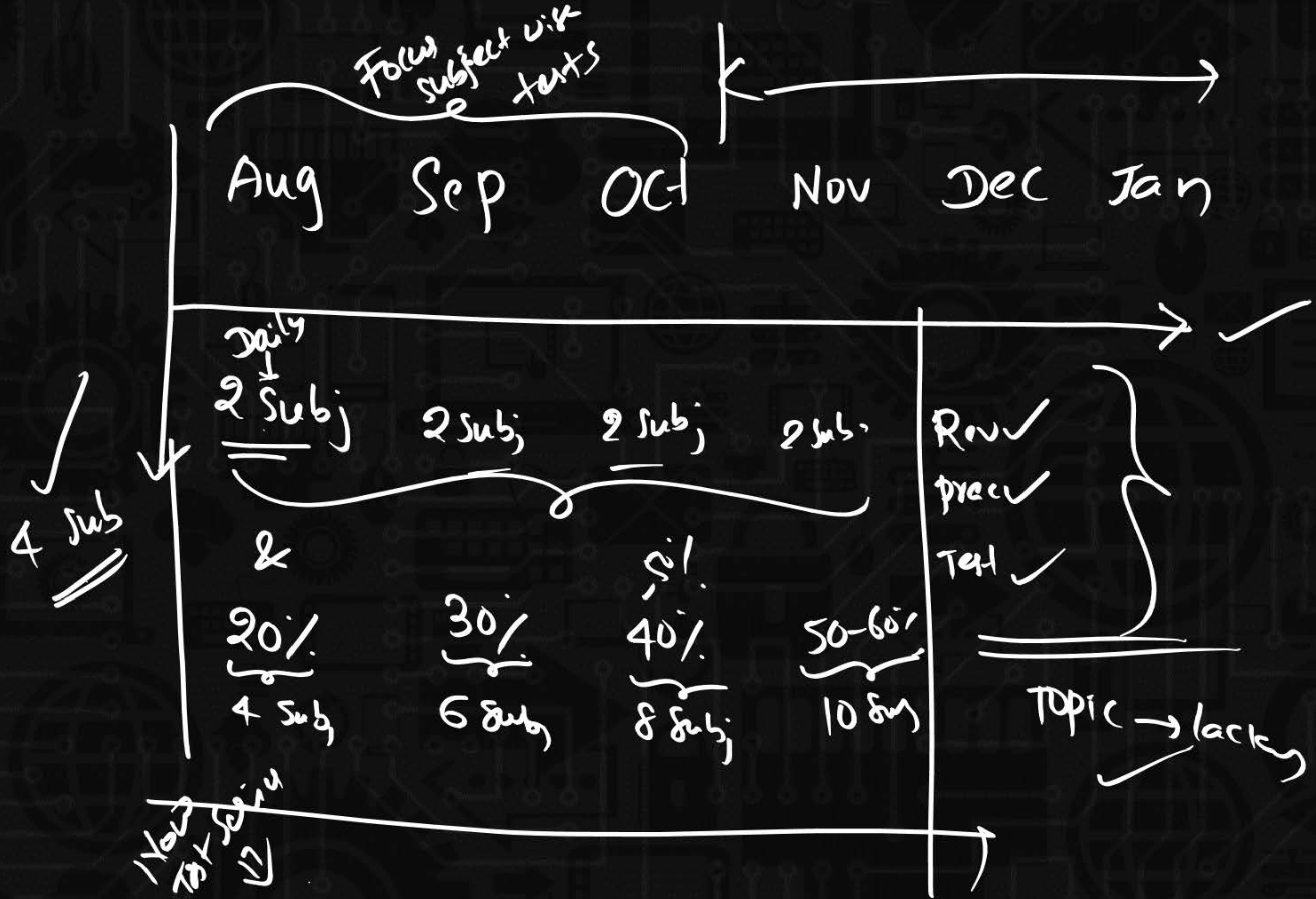
1) Forward Analysis

→ Reaching definitions Analysis

→ Available Expressions Analysis

2) Backward Analysis

*** → live variable Analysis



↳ Code optimization techniques

Next: Data Flow Analysis

