## CS & IT ENGINEERING



Routing Protocols
Lecture No-01



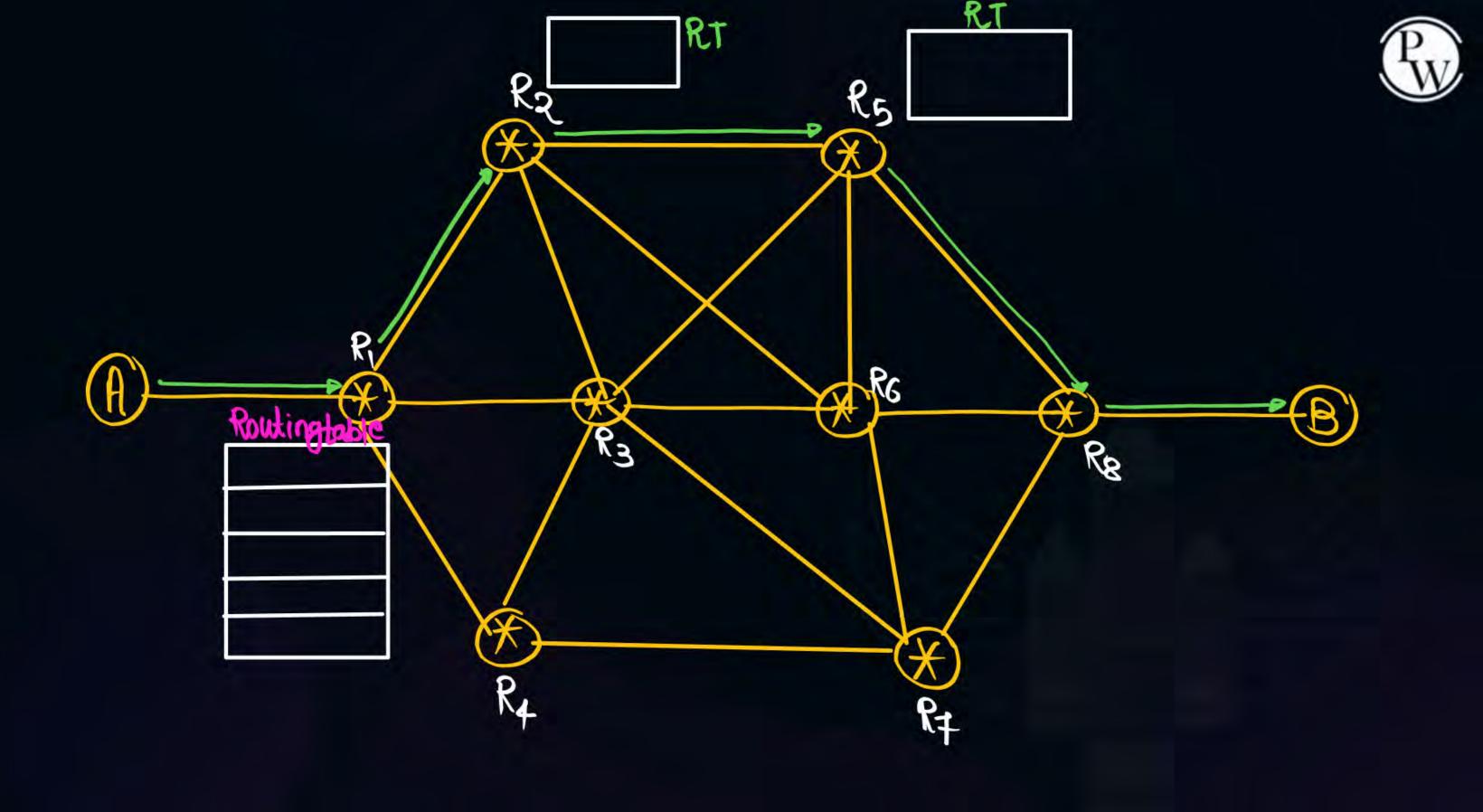
By-Ankit Doyla Sir





## TOPICS TO BE COVERED

Routing Algorithms

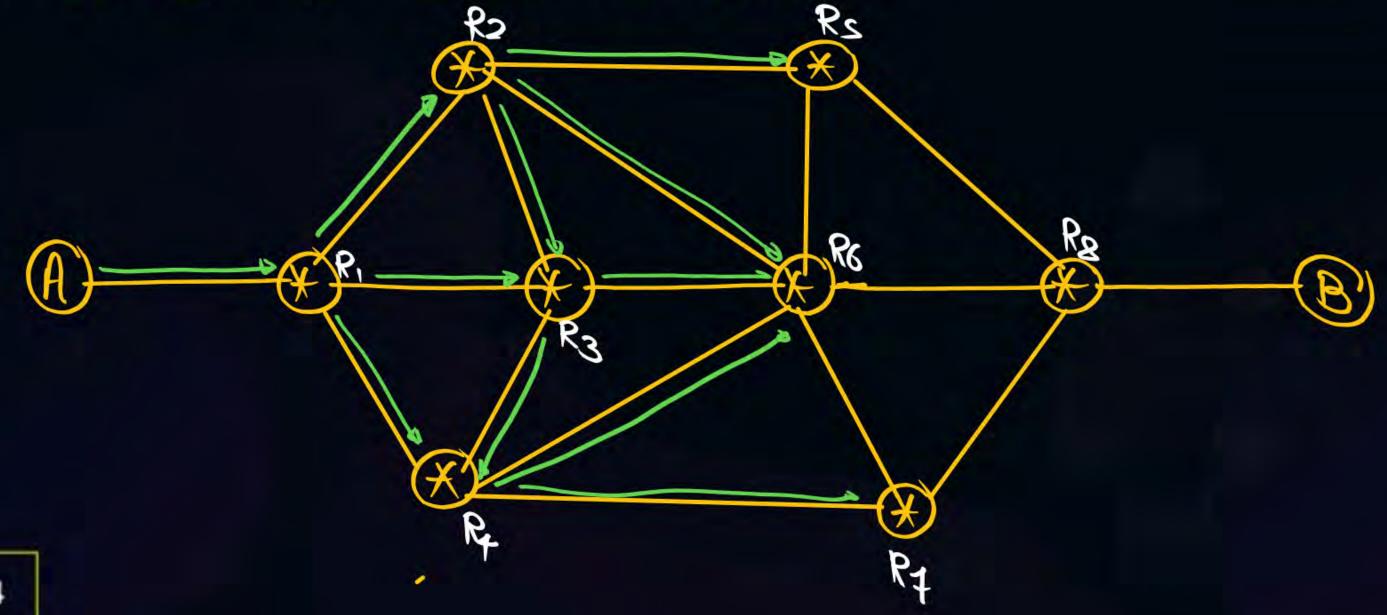




### **Topic: Flooding**



Flooding is simple computer network Routing Algorithm in which every incoming packet is sent through every outgoing link except one it arrived on.





## **Topic: Advantage of Flooding**



- No Routing is required
- 2. Shortest path is always guarantee i.e. the packet arrives at the destination first
- It is highly Reliable, if one path is down then the packet reach at the destination by choosing another path.



## **Topic: Disadvantage of Flooding**

- Traffic is very high
- 2. Many duplicate packets received by receiver



### **Topic: Advantage of Routing**



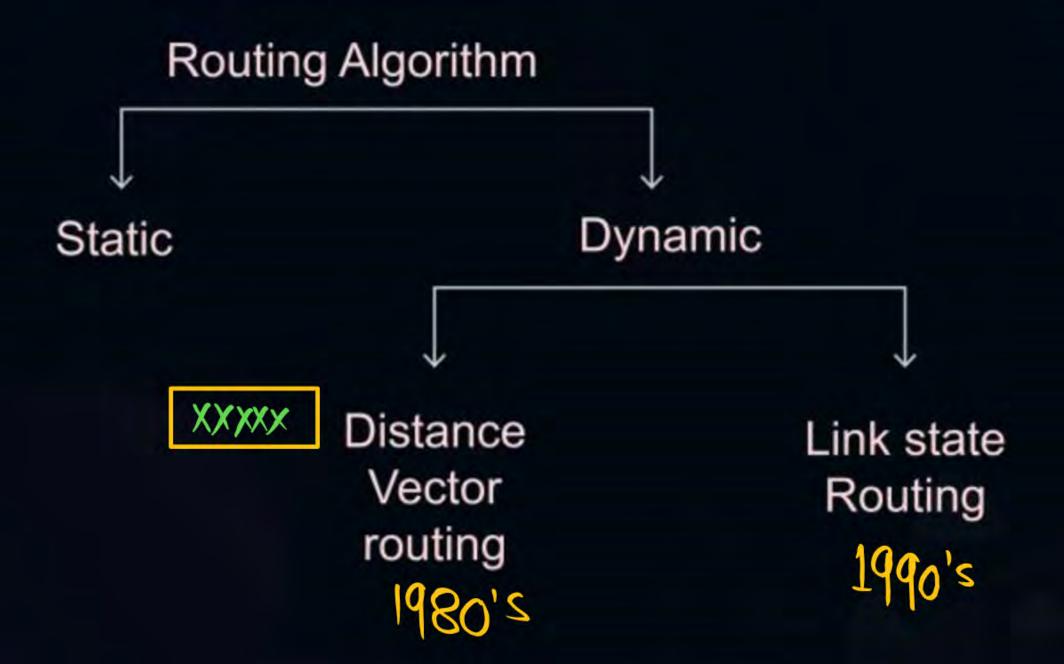
- Traffic is very less.
- No duplicate packet received by receiver.



### **Topic: Disadvantage of Routing**

- Routing table is required.
- Chosen path may be down so it is not highly Reliable.
- Shortest path is depends on the algorithm and some algorithms fails to find the shortest path.





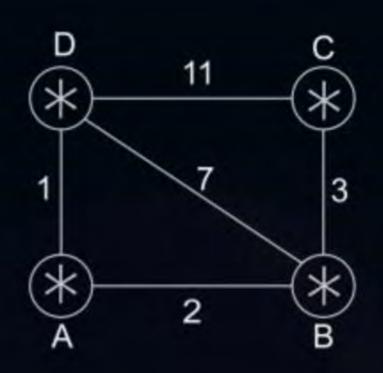


## Distance Vector Routing



## **Topic: Distance Vector Routing: (1980)**





#### **Routing Table**

Destination	Distance	Next Hop

#### Step 1:



Prepare the Routing table at every Router Based on the local knowledge.

Dest.	Dis.	NH
A	1	A
8		8
C	11	C
D	0	٥

Dest.	Dis.	NH
A	0	A
B	2	B
C	8	_
D	1	D



Dest.	Dis.	NH
A	8	
0	3	8
C	O	C
D	11	D

Dest.	Dis.	NH
A	2	A
B	0	B
C	3	C
0	17	D



#### Step 2:

At A

A Receive Distance vector from B, D

At B

B Receive Distance vector from A, C, D

At C

C Receive Distance vector from B, D

At D

D Receive distance from A, B, C

A Received Distance Vector From B, D

A Received Distance Vector From B

From B

From B

$$A \rightarrow B = \min 2 \quad A \rightarrow B + B \rightarrow B$$
 $A \rightarrow B = \min 2 \quad A \rightarrow B + B \rightarrow B$ 
 $A \rightarrow B = \min 2 \quad A \rightarrow B + B \rightarrow B$ 
 $A \rightarrow B = \min 2 \quad A \rightarrow B + B \rightarrow B$ 
 $A \rightarrow B = \min 2 \quad A \rightarrow B + B \rightarrow B$ 



Destination	Distance	NH
A	0	A
8	2	B
C	5	B
D	1	D

$$A \rightarrow C = \min \left\{ \begin{array}{l} A \xrightarrow{2} B + B \xrightarrow{3} c \\ A \rightarrow D + D \xrightarrow{1} c \end{array} \right.$$

$$A \rightarrow c = min 2 5 12) = 5$$

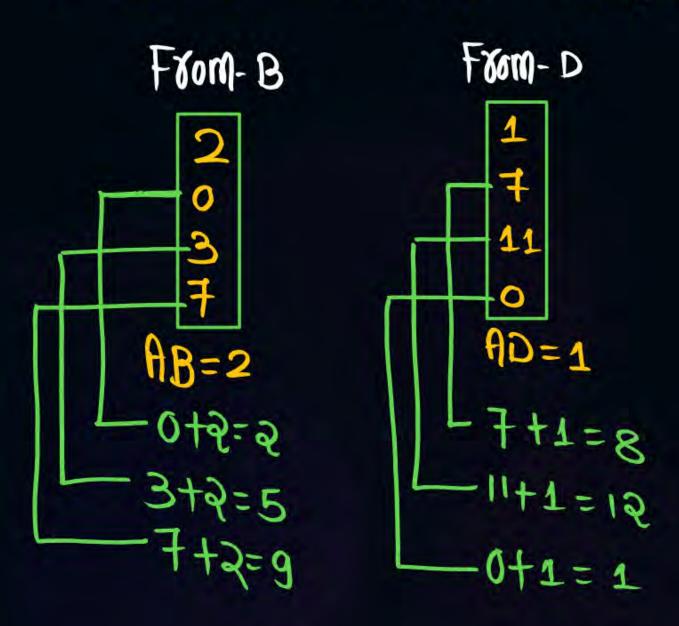
$$A \rightarrow D = \min_{0 \le 1} A \xrightarrow{2} B + B \rightarrow D$$

$$A \rightarrow D + D \rightarrow D$$

$$A \rightarrow D = \min_{0 \le 1} 29,1) = 1$$

## Shootcut (AD Rule)

At A Received Distance vector From B, D

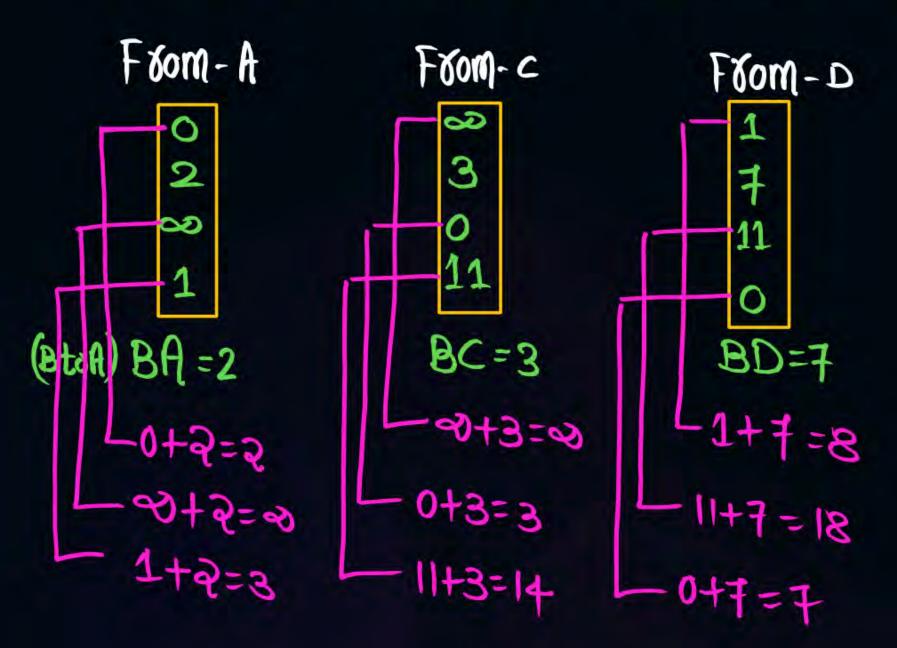




## A- New Routing table

Destination	Distance	ИН
A	0	A
B	2	8
C	5	B
D	1	0

## At B B Received Distance Vector From A,C,D





## B New Pouting table

Dastination	Distance	NH
A	2	A
8	0	В
C	3	C
0	3	A



#### At C: New Routing table

Dest.	Dis	NH
Α	5	В
В	3	В
С	0	С
D	(10)	В

At D: New Routing table

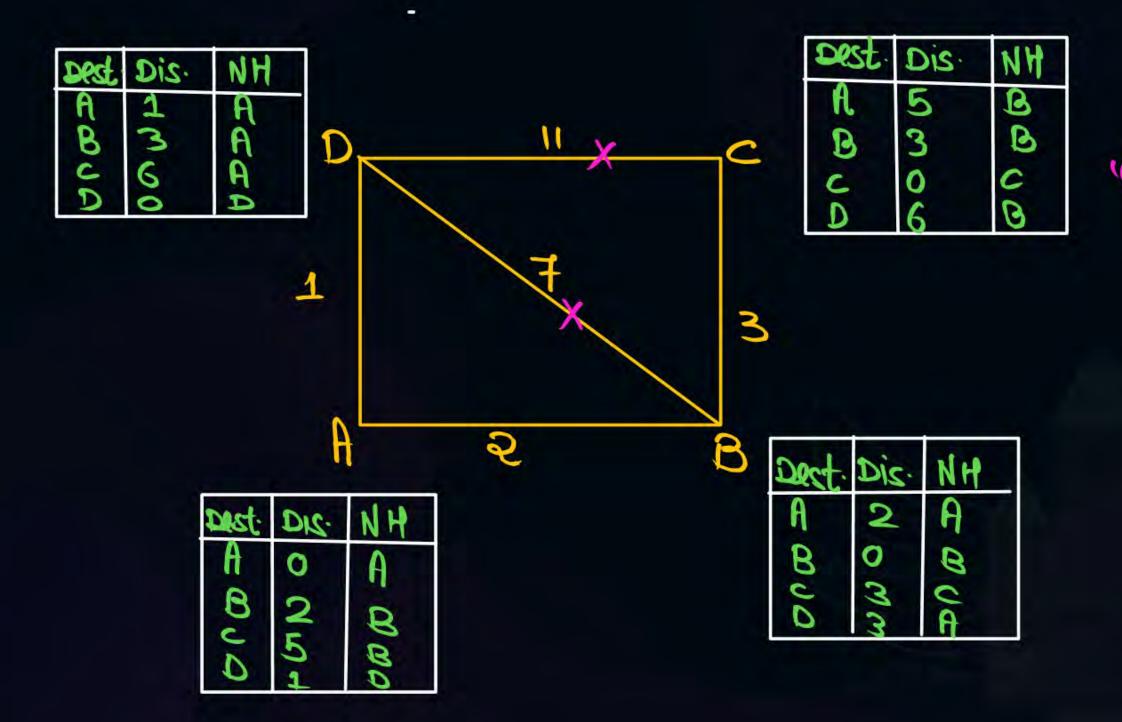
Dest.	Dis	NH
Α	1	Α
В	3	Α
С	10	В
D	0	D

## Step 3: At A A Received Distance Vector From



## Final Routing table in one step





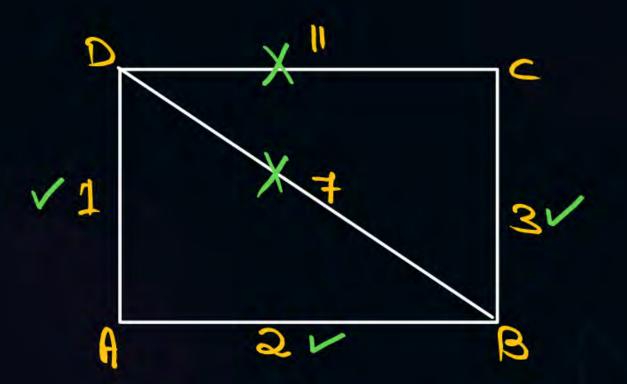
edges which are Not included in any Routing table"











edges which are Not included



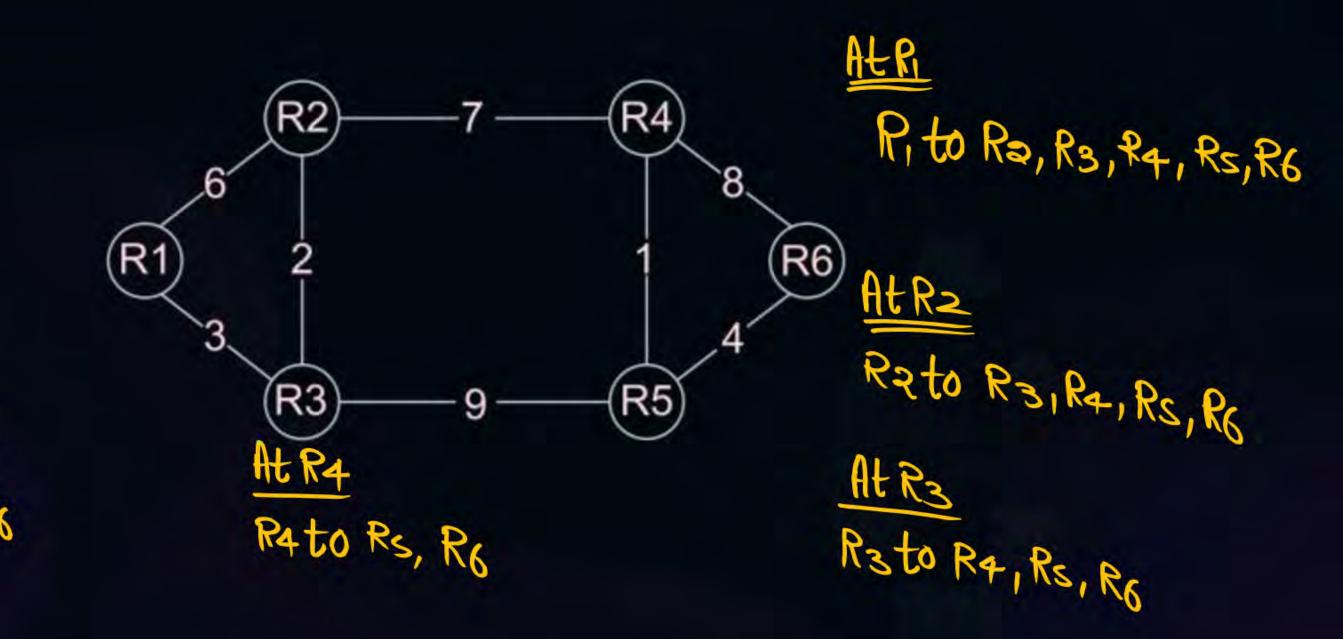
# Problem Solving On DVR

#### Gate 2010





Consider a network with 6 routers R1 to R6 connected with links having weights as shown in the following diagram



At RS
Roto Ro



## **Topic: Problem Solving on DVR**



## Common data Guestion

Q. All the routers use the distance vector-based routing algorithm to update their routing tables. Each router starts with its routing table initialized to contain an entry for each neighbor with the weight of the respective connecting link. After all the routing tables stabilize, how many links in the network will never be used for carrying any data?

A. 4

B. 3

C. 2

D. 1



## **Topic: Problem Solving on DVR**



Q. Suppose the weights of all unused links in the previous question are changed to 2 and the distance vector algorithm is used again until all routing tables stabilize. How many links will now remain unused?

A. 0

B. 1

C. 2

D. 3



## **Topic: Problem Solving on DVR**



Q. Consider the following subnet. If distance vector routing is used, how many link can never by used after all the routing table. Are stabilized?

