



# CS & IT ENGINEERING

## Computer Network

1500 Series

Lecture No.- 09



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# Recap of Previous Lecture



Topic

Questions Practice

Topic





# Topics to be Covered



Topic

Common Data Questions

Topic



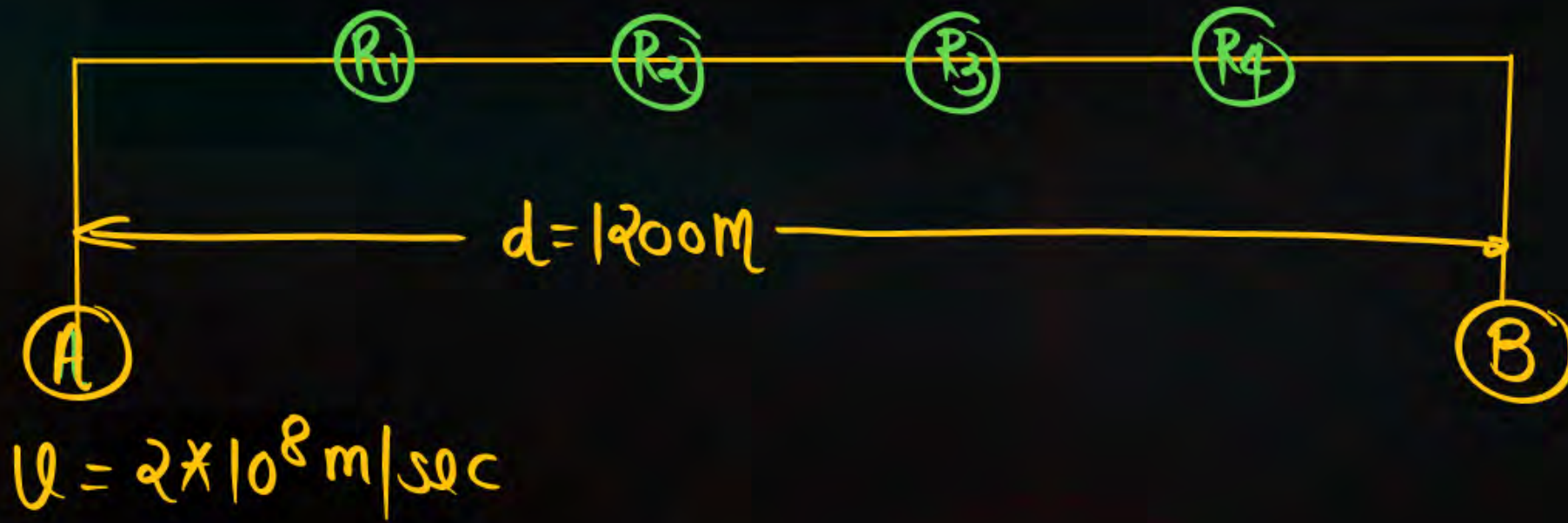
Suppose two nodes, A and B, are attached to opposite ends of an 1200m cable, and that they each have one frame of 1,500 bits (including all headers and preambles) to send to each other. Both nodes attempt to transmit at time t=0. Suppose there are four repeaters between A and B, each inserting a 40-bit delay. Assume the transmission rate is 100 Mbps, and CSMA/CD with backoff intervals of multiples of 512 bits times is used. After the collision, A draws K=0 and B draws K=1 in the exponential backoff protocol. Ignore the jam signal in this case.

#Q. What is the one-way propagation delay (including repeater delays) between A and B in <sup>micro</sup>seconds? Assume the signal propagation speed is  $2 \times 10^8$  m/sec.

#Q. At what time (in <sup>micro</sup>seconds) is A's packet completely delivered at B?

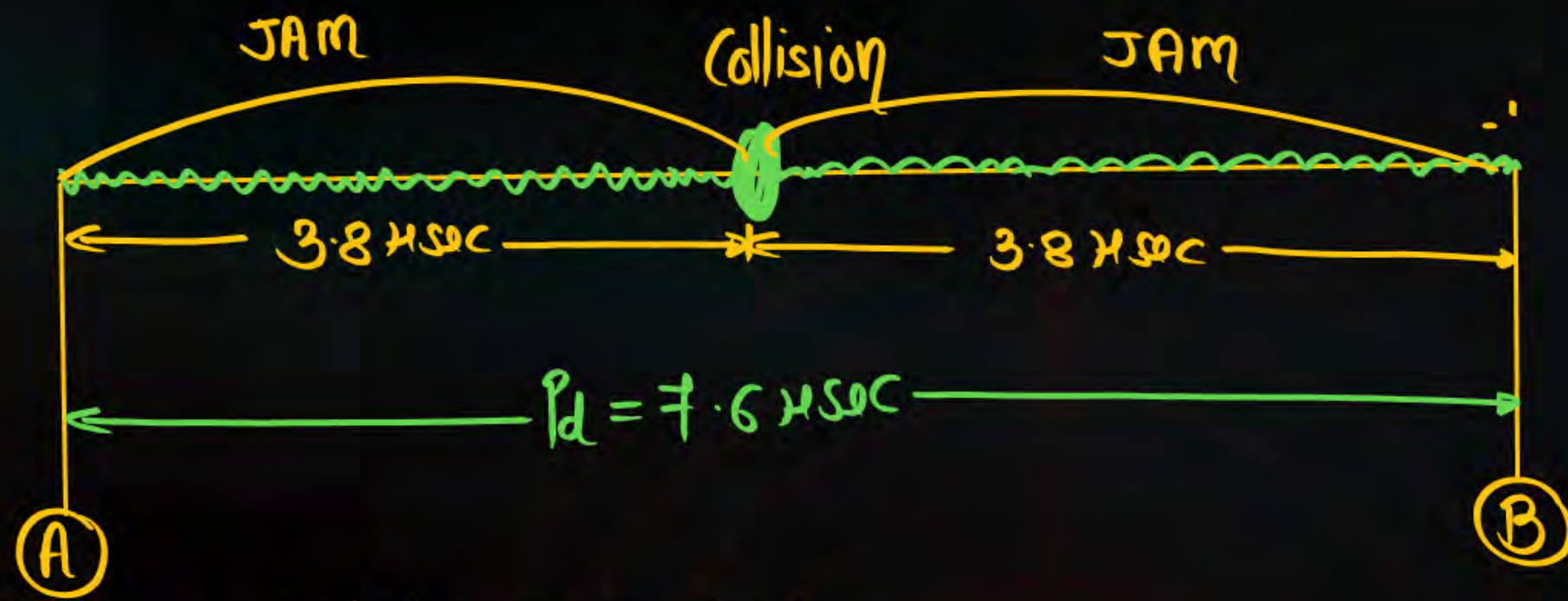
#Q. At what time (in microsec) 'A' retransmission 1<sup>st</sup> bit will reach at 'B'  
 $15.2 + 7.6 = 22.8$





$$\begin{aligned}
 P_d (\text{including Repeater delay}) &= \frac{d}{u} + \text{Repeater delay} \\
 &= \frac{1200 \text{ m}}{2 \times 10^8 \text{ m/sec}} + \frac{4 \times 40 \text{ bit}}{100 \times 10^6 \text{ bps/sec}} \\
 &= 6 \times 10^{-6} \text{ sec} + 1.6 \times 10^{-6} \text{ sec} \\
 &= 6 \mu\text{sec} + 1.6 \mu\text{sec} = 7.6 \mu\text{sec}
 \end{aligned}$$





At  $t=0$  Both A and B starts

At  $t=3.8 \mu\text{sec} \rightarrow$  Both 'A' and 'B' data collide

At  $t=7.6 \mu\text{sec} \rightarrow$  Both 'A' and 'B' detect the collision

$$K=0$$

$$WT = P_d = 7.6 \mu\text{sec}$$

At  $t=7.6 + 7.6 = 15.2 \mu\text{sec}$  'A' start retransmitting the data

$$K=1$$

$$WT = K \times \text{slot duration}$$

$$WT = 1 \times 512 \text{ bits}$$

$$T_d = \frac{\text{PKt size}}{\text{Bandwidth}}$$

$$= \frac{1500 \text{ bits}}{100 \times 10^6 \text{ bits/sec}}$$

$$= 15 \times 10^{-6} \text{ sec} = 15 \mu\text{sec}$$

At  $t = 15.2 \mu\text{sec} + 15 \mu\text{sec} = 30.2$  'A' Finished its transmission

At  $t = 30.2 + 7.6 \mu\text{sec} = 37.8 \mu\text{sec}$  'A's' Packet completely delivered at 'B'



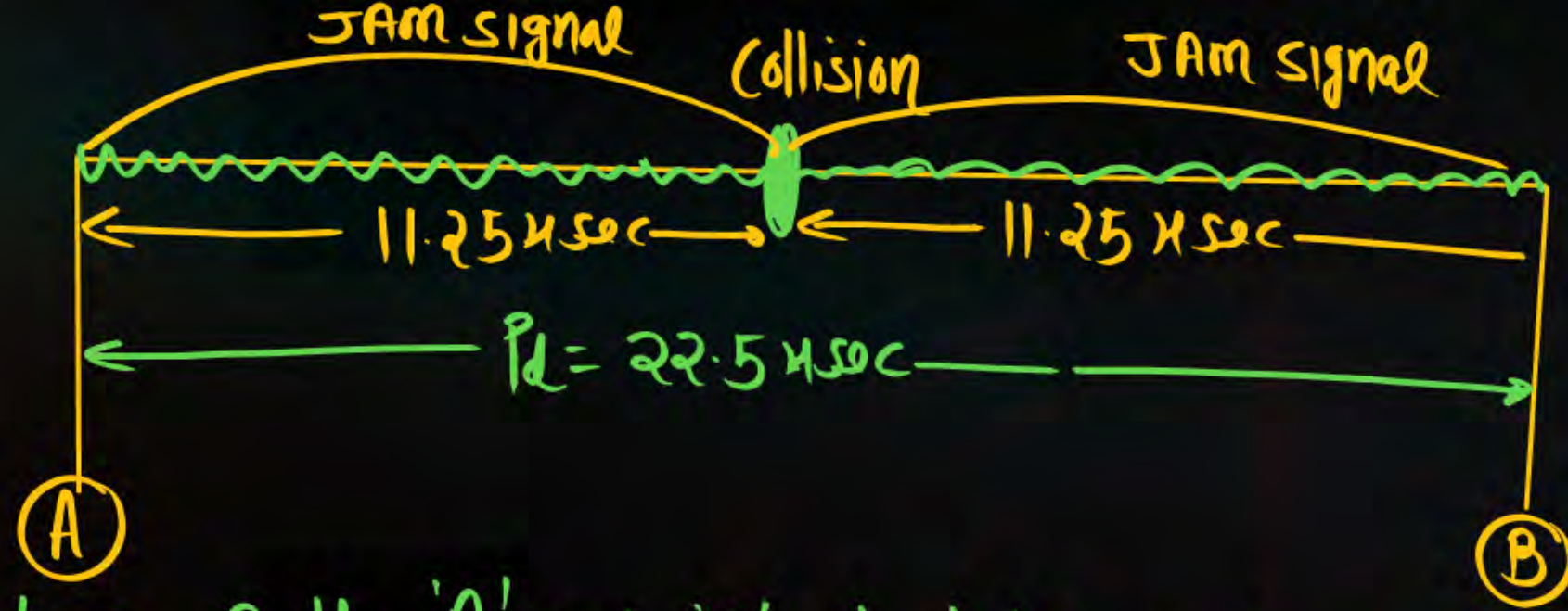
Suppose nodes A and B are on the same Ethernet bus, and the propagation delay between the two nodes is 225-bit times. Suppose both A and B send frames of 1,500 bits (including all headers and preambles) exactly at the same time, the frames collide, and then A and B choose different values of K ( $A = 0$  and  $B = 1$ ) in the CSMA/CD algorithm. Assume the transmission rate is 10 Mbps; CSMA/CD with backoff intervals of multiples of 512-bits is used. If a node detect collision, then it sends a 48-bit jam signal of inform other nodes. Assuming no other node is active and transmission time of a data frame is negligible

#Q. At what time (in microseconds) does A begin retransmission?

\_\_\_\_\_

#Q. At what time (in microseconds) A's packet completely delivered at B? \_\_\_\_\_





$$P_d = \frac{225 \text{ bit}}{10 \times 10^6 \text{ bits/sec}} = 22.5 \times 10^{-6} \text{ sec} = 22.5 \mu\text{sec}$$

At  $t=0$  Both 'A' and 'B' start transmitting data

At  $t=11.2 \mu\text{sec}$  Both 'A' and 'B' data Pkt collide

At  $t=11.25 + 4.8 + 11.25 = 27.3$  Both 'A' and 'B' detect collision

$K=0$

$WT = P_d = 22.5 \mu\text{sec}$

At  $t=27.3 + 22.5 = 49.8 \mu\text{sec}$  'A' start its retransmission

$K=1$

$$\begin{aligned} T_d(\text{JAM signal}) &= \frac{\text{JAM signal size}}{\text{Bandwidth}} \\ &= \frac{48 \text{ bits}}{10 \times 10^6 \text{ bits/sec}} \\ &= 4.8 \times 10^{-6} \text{ sec} \\ &= 4.8 \mu\text{sec} \end{aligned}$$

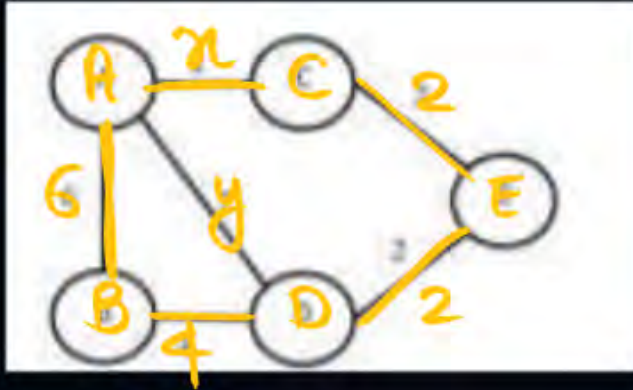


At  $t = 49.8 + 150 = 199.8 \mu\text{sec}$  'A' Finish its transmission  $T_d(\text{PKT}) = \frac{\text{PKT size}}{\text{Bandwidth}}$

At  $t = 199.8 + 22.5 = 222.3$  'A's PKT completely delivered at 'B'  $= \frac{1500 \text{ bits}}{10 \times 10^6 \text{ bits/sec}}$   
 $= 150 \times 10^{-6} \text{ sec} = 150 \mu\text{sec}$



Consider the following network



$$\cancel{x+y} < \cancel{x+2+2}$$

$$y < 4$$

Which of the following condition must be satisfied to ensure that traffic from B to C will always Flow through node A?

- (a)  $x > 4$
- (b)  $y + x < 6$
- ☒ (c)  $y + x < 4$
- (d)  $x < 4$



The network uses a Distance Vector Routing protocol to compute the distances and next hops between different node pairs. Given the initial distance vector table after first round of distance vector exchange

Information Stored at Node	Distance to Reach Node						
	A	B	C	D	E	F	G
A	0	2	7	$\infty$	$\infty$	3	$\infty$
B	2	0	4	$\infty$	$\infty$	$\infty$	$\infty$
C	7	4	0	$\infty$	2	$\infty$	$\infty$
D	$\infty$	$\infty$	$\infty$	0	10	1	$\infty$
E	$\infty$	$\infty$	2	10	0	$\infty$	2
F	3	$\infty$	$\infty$	1	$\infty$	0	1
G	$\infty$	$\infty$	$\infty$	$\infty$	2	1	0

Each distance vector is the distance of the best known path at that instance to nodes, A to G, where the distance to itself is 0. Also, all links are symmetric and the cost is identical in both directions. In each round, all nodes exchange their distance vectors with their respective neighbours. Then all nodes update their distance vectors. In between two rounds, any change in cost of a link will cause the two incident nodes to change only that entry in their distance vectors. The distance vector table entries of node A after second round of distance vector exchange with its neighbours is

	A	B	C	D	E	F	G
(a) ✓	0	2	7	4	$\infty$	3	4
(b) ✓	0	2	6	4	9	3	4
(c)	0	2	6	4	$\infty$	3	4
(d)	0	2	6	4	8	3	4





A Recd DV From B, C, F

From B

2
0
4
8
8
8
8

AB=2

0+2=2

4+2=6

$\infty+2=\infty$

$\infty+2=\infty$

$\infty+6=\infty$

From C

7
4
0
8
2
8
8

AC=7

4+7=11

0+7=7

~~8~~+7= $\infty$

2+7=9

$\infty+7=\infty$

$\infty+7=\infty$

From F

3
8
8
1
8
0
1

AF=3

$\infty+3=\infty$

$\infty+3=\infty$

1+3=4

$\infty+3=\infty$

0+3=3

1+3=4

New Routing table of 'A'

Des.	Dis.	NH
A	0	A
B	2	B
C	6	B
D	4	F
E	9	C
F	3	F
G	4	F



How many statements about link state routing protocol is/are TRUE?

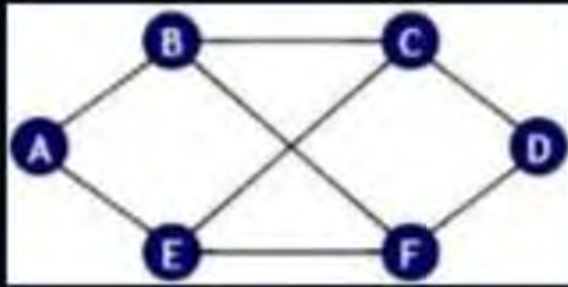
1. Routing messages from a router should be forwarded to all other routers in its area. (True)
2. Routers build a complete picture of the whole network and compute shortest LS paths locally. (True)
3. Suffers from the "count to infinity" problem. DVR suffers From count to infinity Not
4. Does not prevent routing loops. DVR does Not Prevent Routing loops & Not LSR (False) LSR (False)
5. Typically uses Dijkstra's shortest path algorithm. \_\_\_\_\_ (True)



## [MCQ]



Consider the following network, which utilize the Distance vector routing:



F New Routing table

Dest.	Dis.	NH
A	9	B
B	6	B
C	8	
D		
E		
F	0	F

Router F has the following routing table:

A:  $\infty$ , B: 6, C:  $\infty$ , D: 3, E: 5, F: 0

The following vectors have just come in to router F:

From B: A: 3, B: 0, C: 8, D: 12, E: 6, F: 2

From D: A: 16, B: 12, C: 6, D: 0, E: 9, F: 10

From E: A: 7, B: 6, C: 3, D: 9, E: 0, F: 4

What is F's new routing table?

~~(a) A:  $\infty$ , B: 6, C:  $\infty$ , D: 3, E: 5, F: 0~~

~~(b) A: 9, B: 6, C: 9, D: 3, E: 5, F: 8~~

(c) A: 9, B: 6, C: 8, D: 3, E: 5, F: 0

(d) A: 9, B: 6, C: 9, D: 3, E: 5, F: 0

At F

F Rcvd DV From B, D, E

From B

3  
0  
8  
12  
6  
2

FB=6

8+6=14

From D

16  
12  
6  
0  
9  
10

FD=3

6+3=9

From E

7  
6  
3  
9  
0  
4

FE=5

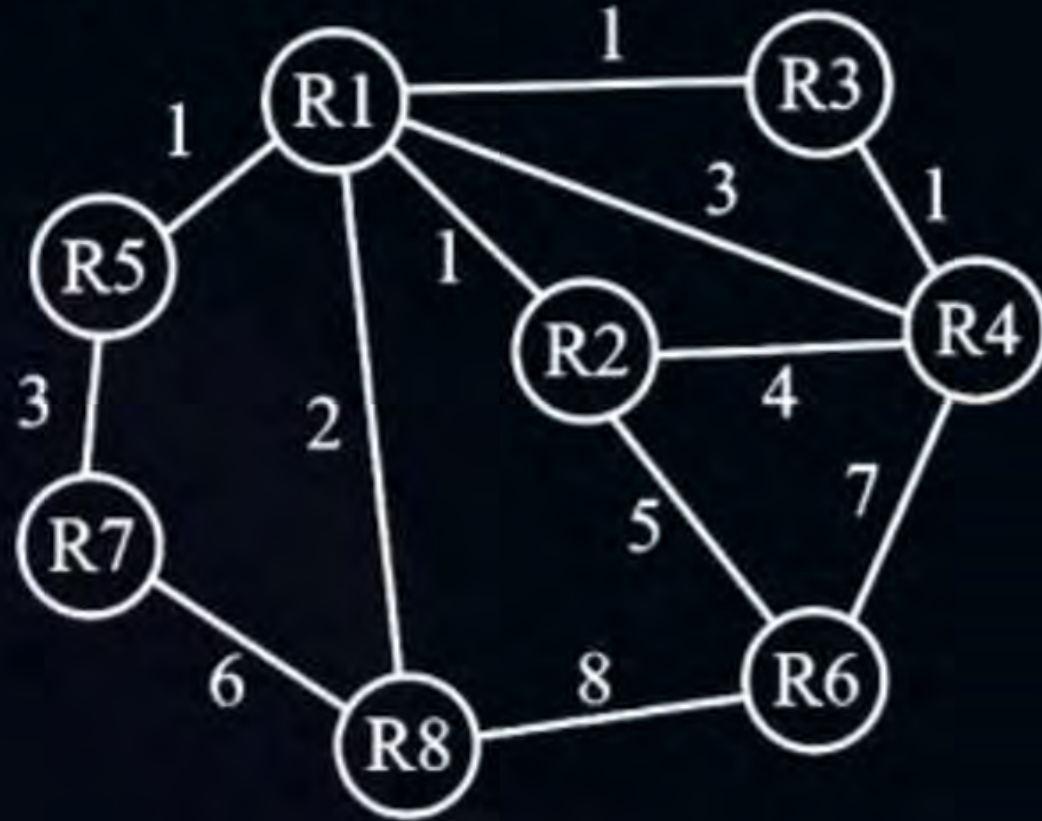
3+5=8



## [MCQ]



Consider a network with 8 routers R1 to R8 connected with links having weight shown in the following diagram:



Routing table of R2

Dest	Dis	NH
R1	1	
R2	0	
R3	2	
R4		
R5		
R6		
R7		
R8		

All the routers use the distance vector based routing algorithm to update their routing tables. Each router starts with its routing table initialized to contain an entry for each neighbour with the weight of the respective connecting link. After all the routing tables stabilize, then what will be the routing table at router R2?



A

R1	R2	R3	R4	R5	R6	R7	R8
1	0	5	4	2	5	5	3

B

R1	R2	R3	R4	R5	R6	R7	R8
1	0	2	3	2	5	5	3

C

R1	R2	R3	R4	R5	R6	R7	R8
1	0	5	4	2	5	4	3

D

R1	R2	R3	R4	R5	R6	R7	R8
1	0	4	5	2	5	4	3



Consider the following statements about the protocols:

- ~~S<sub>1</sub>~~: HTTP and FTP both are In-band protocol
- ~~S<sub>2</sub>~~: HTTP is stateful and FTP is stateless protocol
- ✓ S<sub>3</sub>: HTTP is stateless but FTP is stateful protocol
- ~~S<sub>4</sub>~~: HTTP and FTP both are out of band protocol

Which of the above statements are correct?

- ☐ A S<sub>1</sub>
- ☐ B S<sub>2</sub>
- ☒ C S<sub>3</sub>
- ☐ D S<sub>4</sub>

stateless	stateFull
DNS	POP
SMTP	IMAP
HTTP	FTP

In-Band	out of Band
DNS, SMTP, HTTP, IMAP POP	FTP





## 2 mins Summary



Topic

One

**Common Data Questions**

Topic

Two

Topic

Three

Topic

Four

Topic

Five



**THANK - YOU**