

CS & IT ENGINEERING

COMPUTER NETWORKS

TCP & UDP

Lecture No-9



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TOPICS TO
BE
COVERED



Error control in TCP

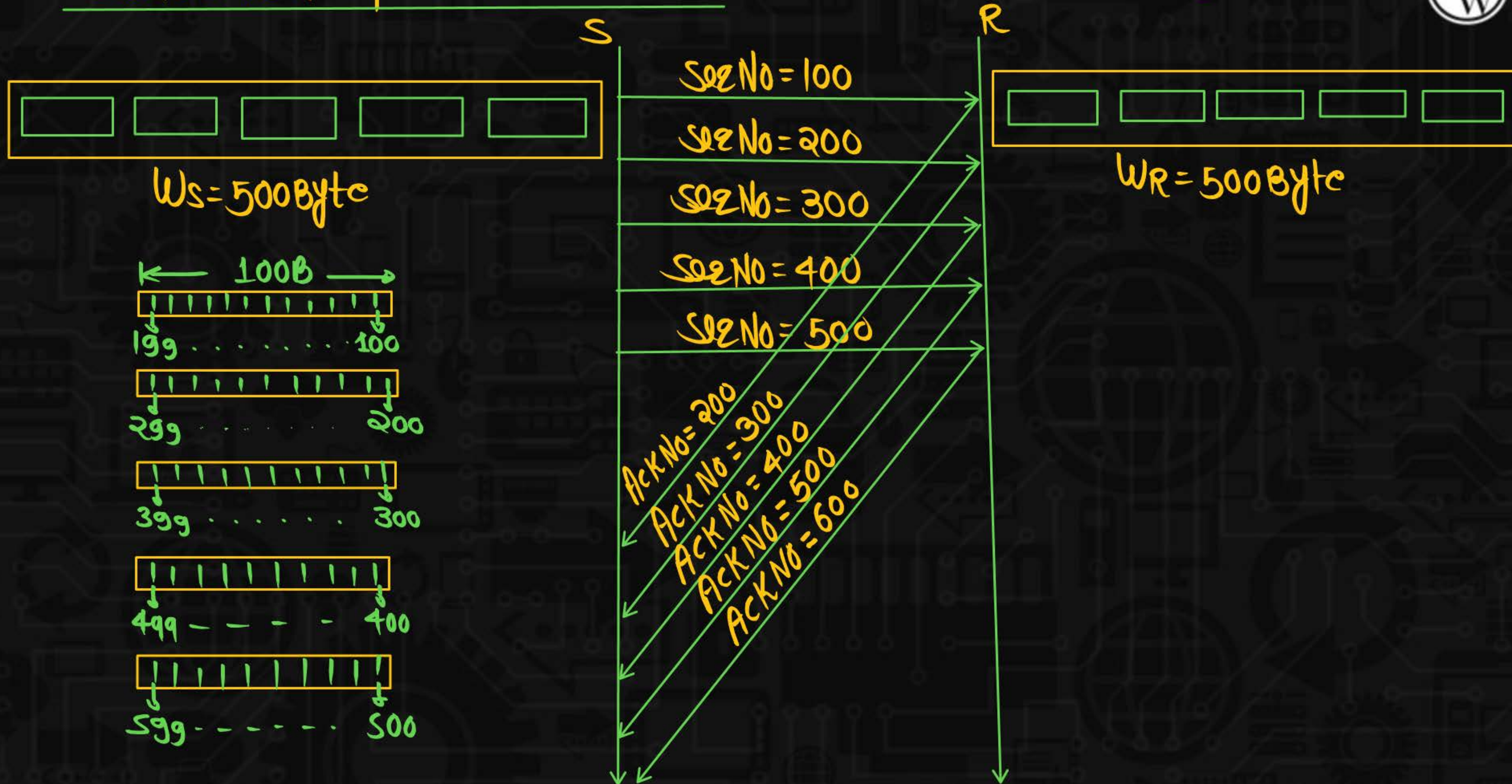
Error control in TCP:

- > TCP can use both selective and cumulative acknowledgement.
- > Receiver may choose to send independent ACK or cumulative ACK
- > TCP uses a combination of selective repeat and GO-Back-N protocol for error control and flow control.
- > In TCP sender window size = receiver window size.
- > In TCP out of order packets are accepted by the receiver.

When ever receiver receives an out of order packet ,it accept that packet but send an acknowledgement for the expected packet.

- > Out of order segments are never delivered to the process.
- > TCP guarantee that data are delivered to process in order.

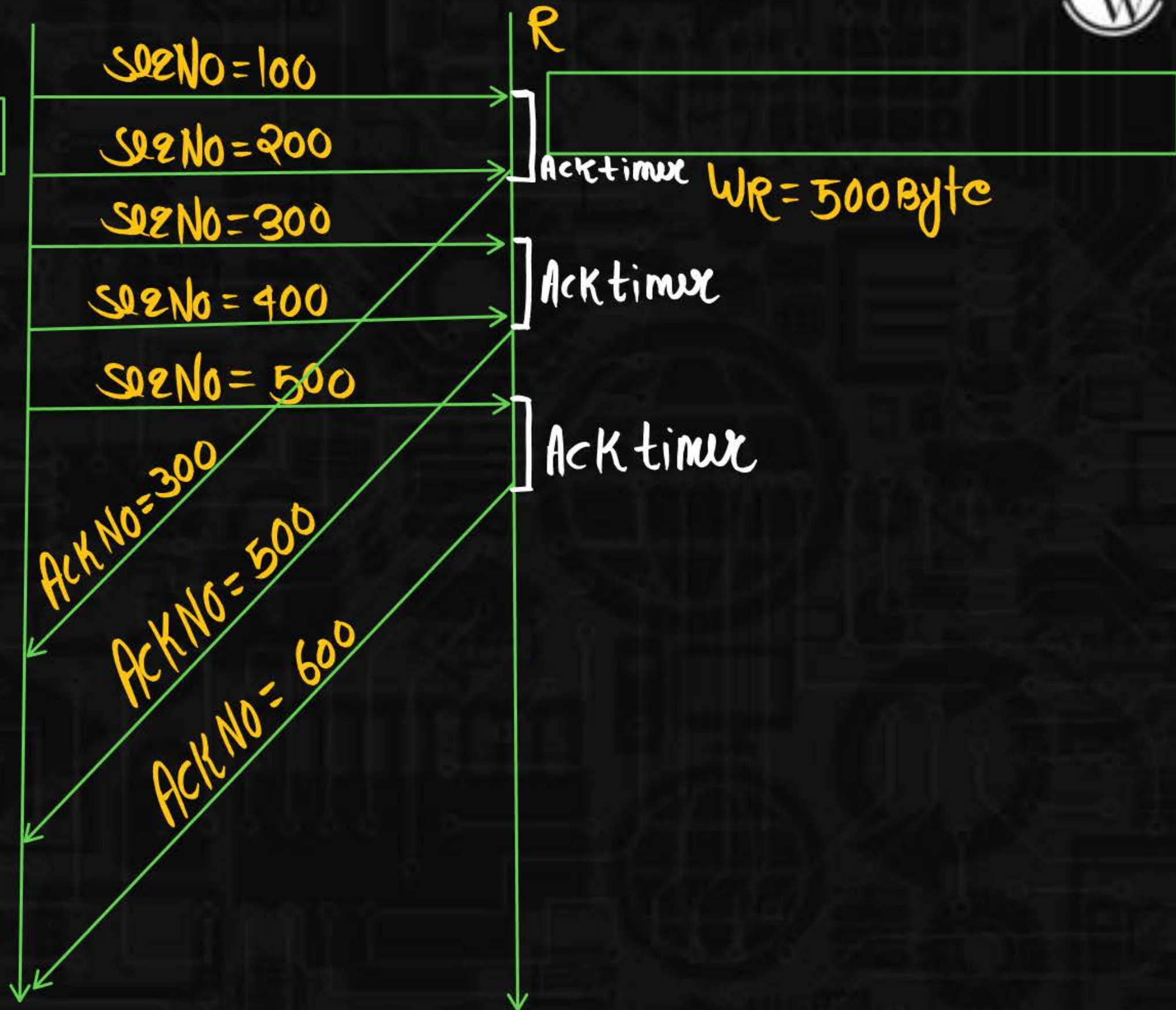
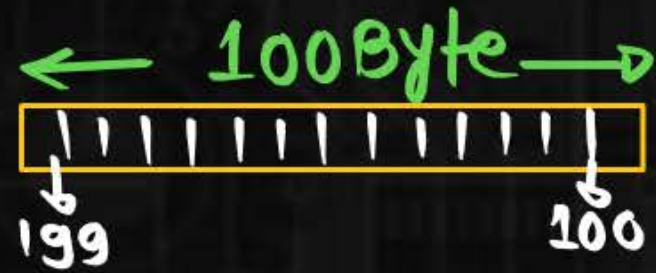
Selective Ack/Independent Ack



Cumulative Ack



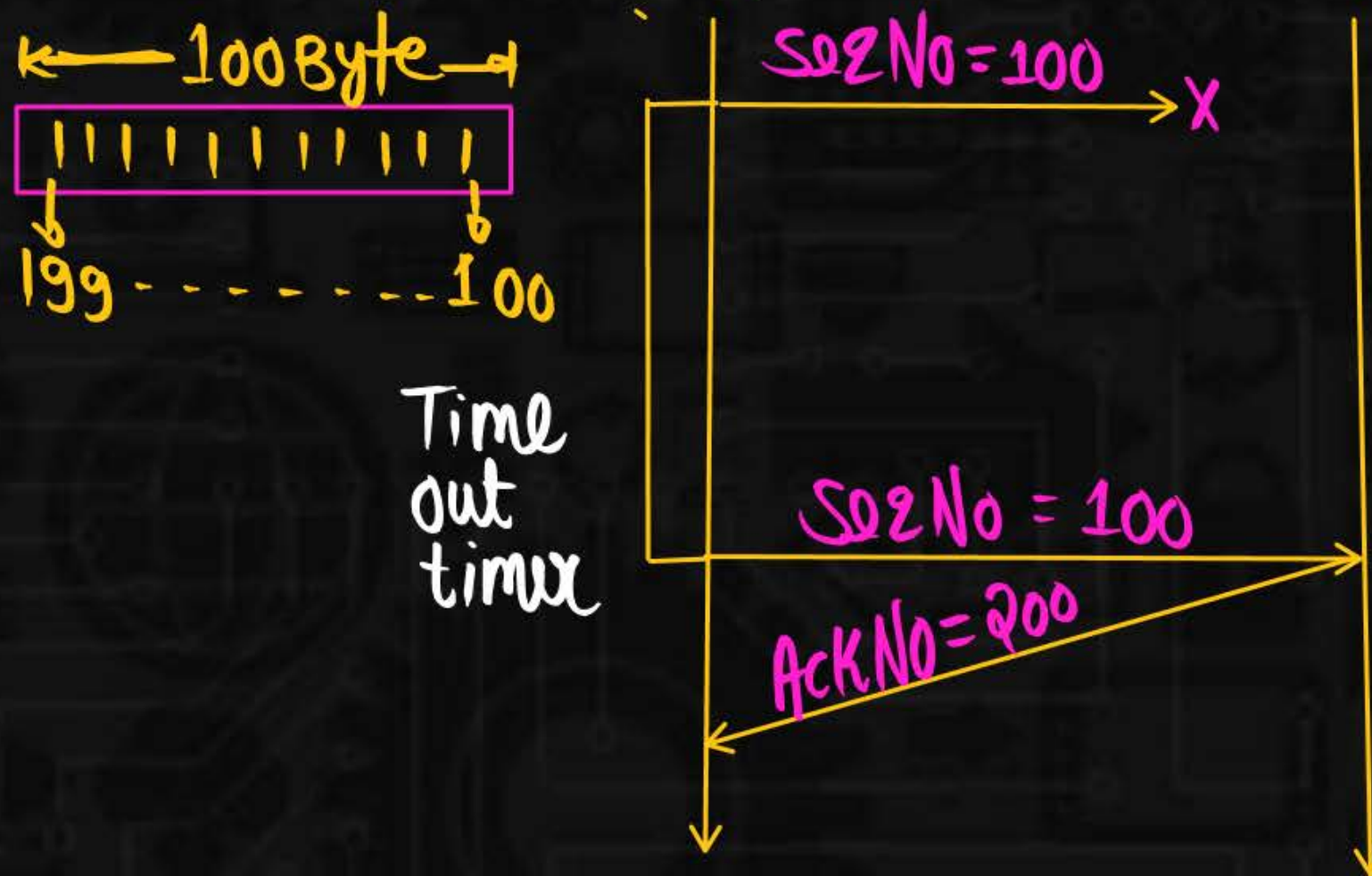
Ws = 500 Byte



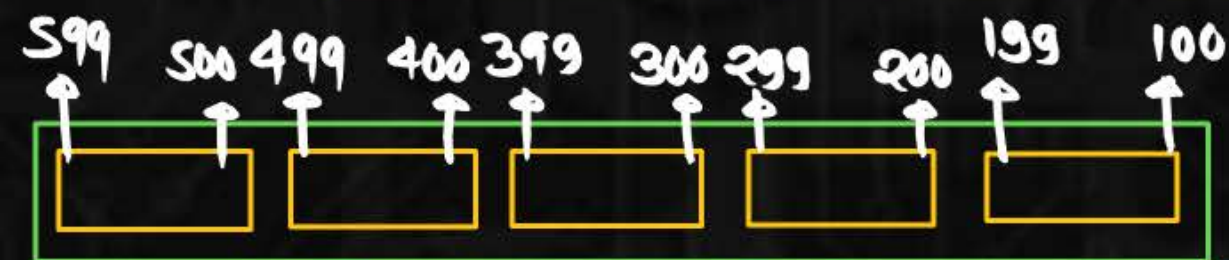
Retransmission in TCP

- (i) Retransmission after time out timer
- (ii) Retransmission after 3 duplicate Ack

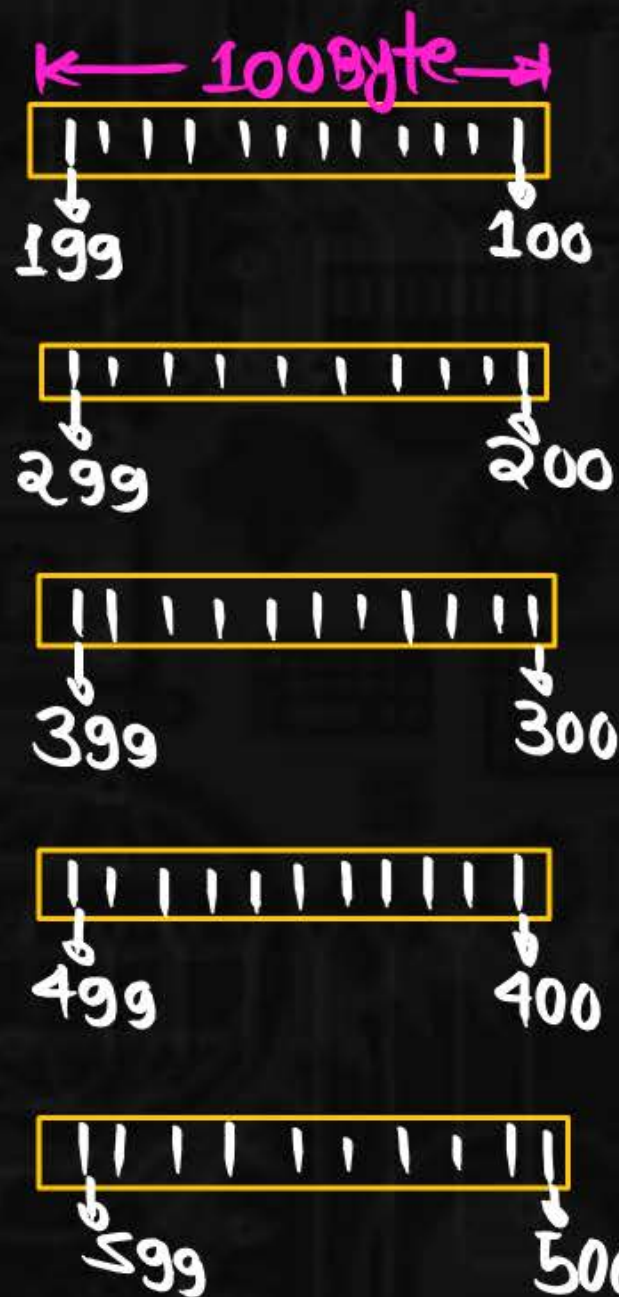
(i) Retransmission after time out timer



(ii) Retransmission after 3 duplicate Ack [Fast Retransmission]



WS = 500 Byte



SeqNo = 100

SeqNo = 200 → X

SeqNo = 300

SeqNo = 400

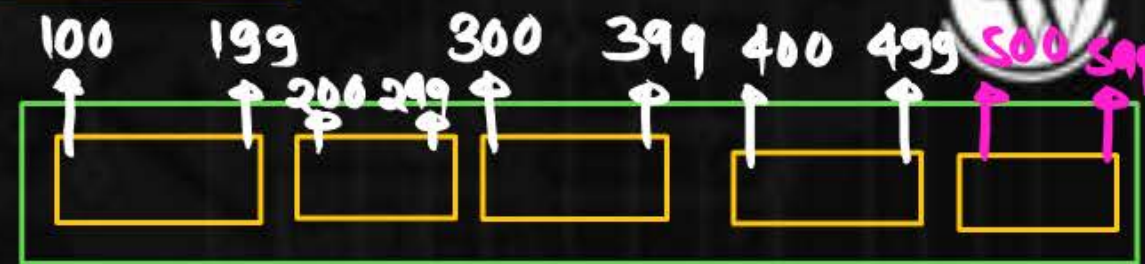
SeqNo = 500

AckNo = 200

AckNo = 200
AckNo = 200
AckNo = 200

SeqNo = 200

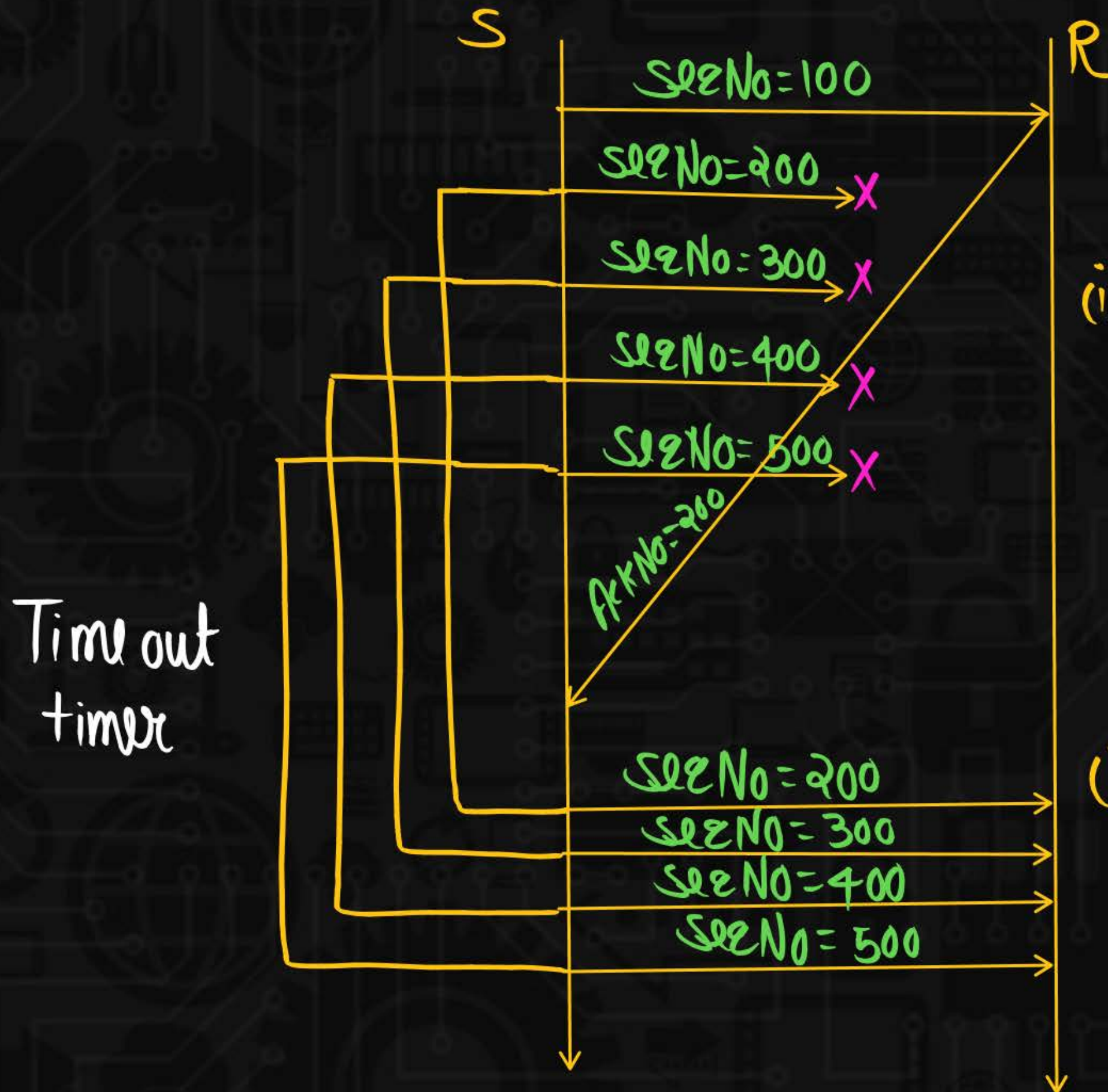
AckNo = 600



WR = 500 Byte

3 duplicate Ack

↓
Indicate mild congestion condition



Note

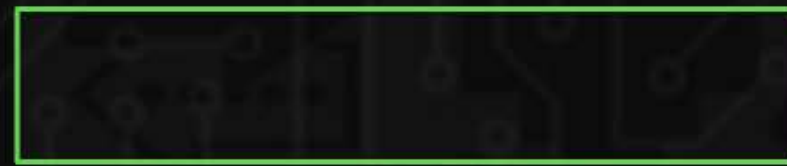
- (i) If 3 duplicate Ack Not possible then we use Time out timer concept For retransmitting the Lost Packet
- (ii) Time out timer indicate severe congestion condition

Lost Acknowledgement

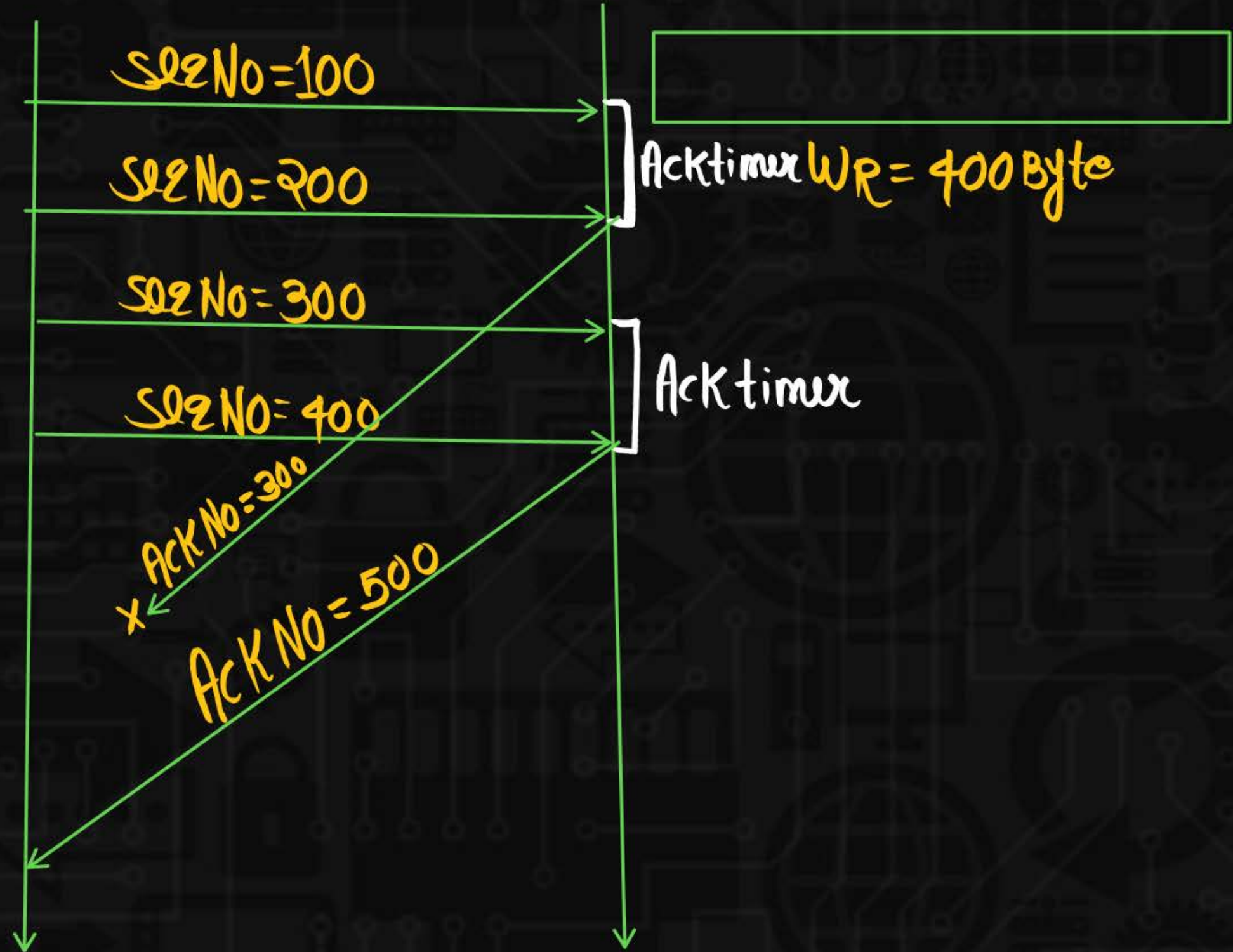


- (i) Automatically corrected Lost Ack
- (ii) Lost Ack corrected by Resending the segment

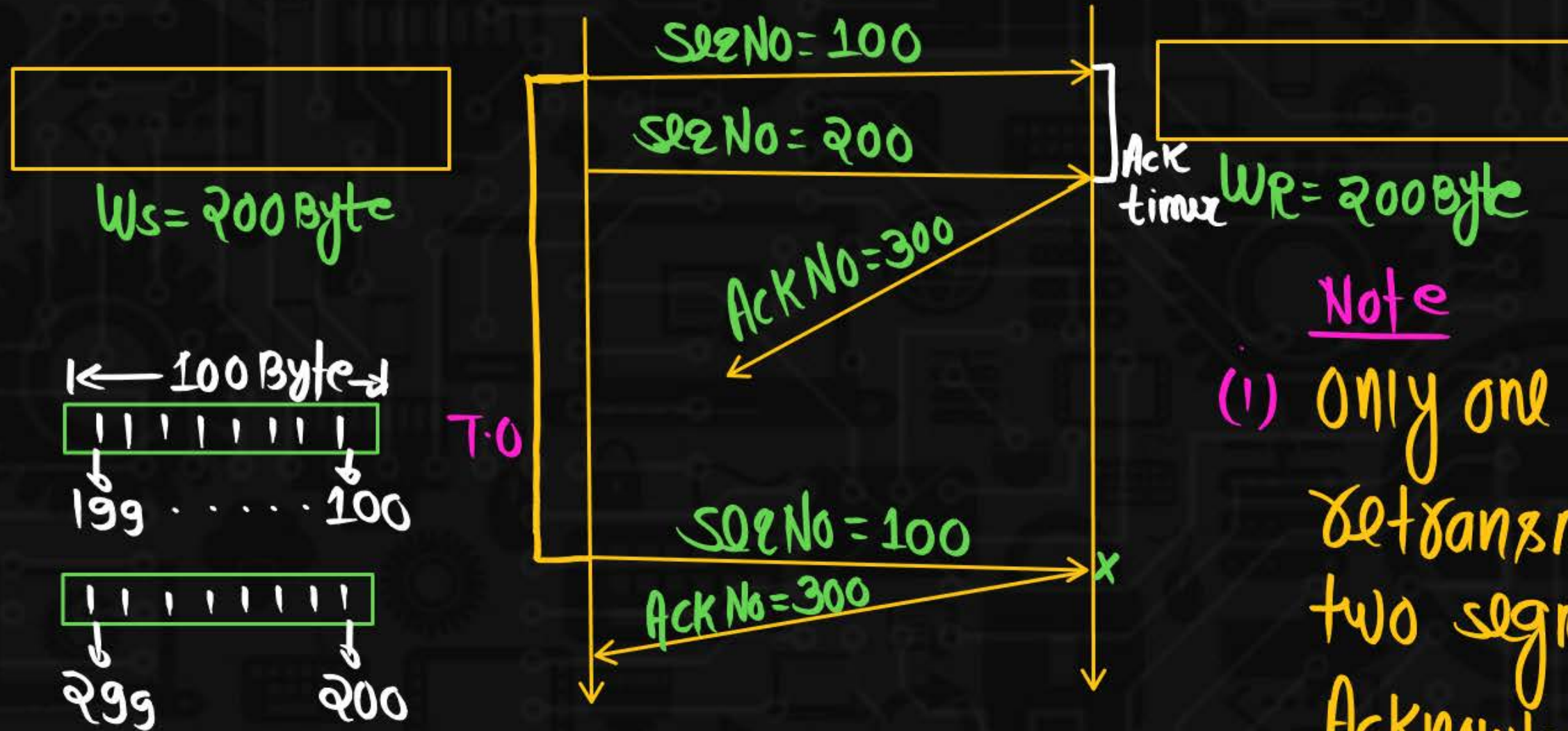
(i) Automatically corrected lost Ack



WS=400Byte



(ii) Lost Ack Corrected by Resending the segment



Note

(i) Only one segment is retransmitted Although two segments are Not Acknowledge

