

CS & IT ENGINEERING



Data structures &
Programming

Linked List


Lec- 06



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TOPICS TO BE
COVERED



Linked List -6

✓ The following C function takes a S.L.L P as an argument.

```
int f(struct item *p){
```

```
return ((P == NULL) || (P->next == NULL)) ||
```

the func.

returns 1

if and only if

$$\left((P \rightarrow \text{data} \leq P \rightarrow \text{next} \rightarrow \text{data}) \right) \&\&$$
$$f(p \rightarrow \text{next}))$$

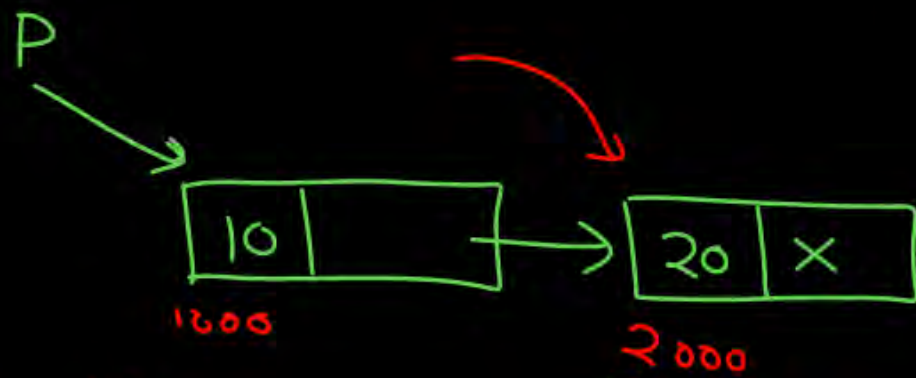
- A) The list is empty or exactly 1 ele.
- B) The ele in list are sorted in non-dec. order.
- C) The " " " " " " -ing "
- D) Not all the ele. in the list have same data.

✓ The following C function takes a S.L.L P as an argument.

```
int f(struct item *P){
```

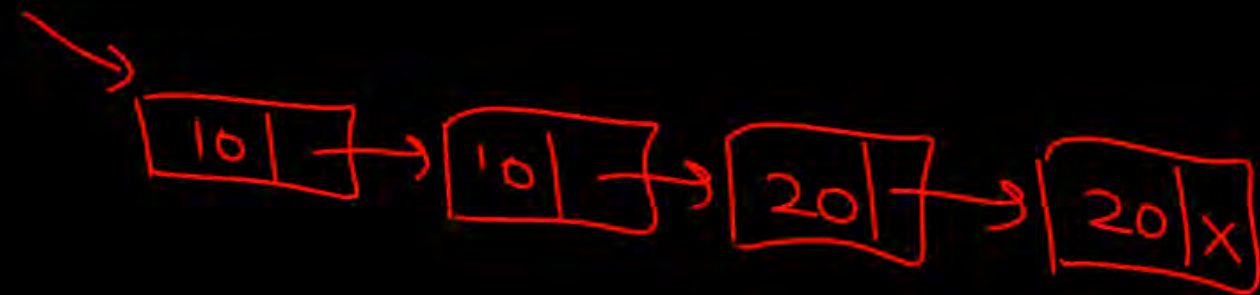
```
    return ( (P == NULL) || (P->next == NULL) ) ||
```

```
    ( (P->data <= P->next->data) &&
      f(P->next) );
```



return 0 || 0 || (10 <= 20)
0 || 0 || 1 &&

return ① && f(2000)



(distinct ele)

Asc./Increasing

10, 10, 20, 20

→ बढ़ती

(Duplicate ele)

Non-decreasing

1, 2, 5, 7, 8

1, 1, 1, 3, 5, 5

(2) In worst case, the no. of comparisons needed to search a singly linked list of length n for a given element is :-

A) $\log n$

B) $n/2$

C) $\log n - 1$

~~D) n~~



worst case:

3.

what is the worst case time complexity to reverse a singly linked list in $O(1)$ space?

\downarrow
Constant memory $\leftarrow \begin{matrix} 1 \\ 2 \\ 3 \end{matrix}$

limited no. of var.

$O(n)$

4



$$1 + 2 + 3 + \dots + (n-1)$$

$$= \frac{(n-1)(n)}{2}$$

$$= \frac{n^2 - n}{2}$$

1st → 0
2nd → 1
3rd ele

⇒ 2 com

4th

⇒ 3 comp.

5th

⇒ 4 comp

⋮

nth

⇒ (n-1) comp



```

void join(node *m, node *n) {
    node *p = n;
    while(p->next != NULL)
    {
        p = p->next;
    }
    p->next = m;
}

```

Assuming that m and n points to
valid NULL terminated LL.
 Invocation of join will.

- A) append list m to the end of list n for all i/p's
- ☒ B) Either cause a null pointer dereference or append list m to the end of list n
- C) Cause a null pointer dereference.
- D) Append list m to the end of list n.

```
void join( node *m, node *n ) {
```

```
    node *p = n;
```

```
    while ( p->next != NULL )
```

```
    {
```

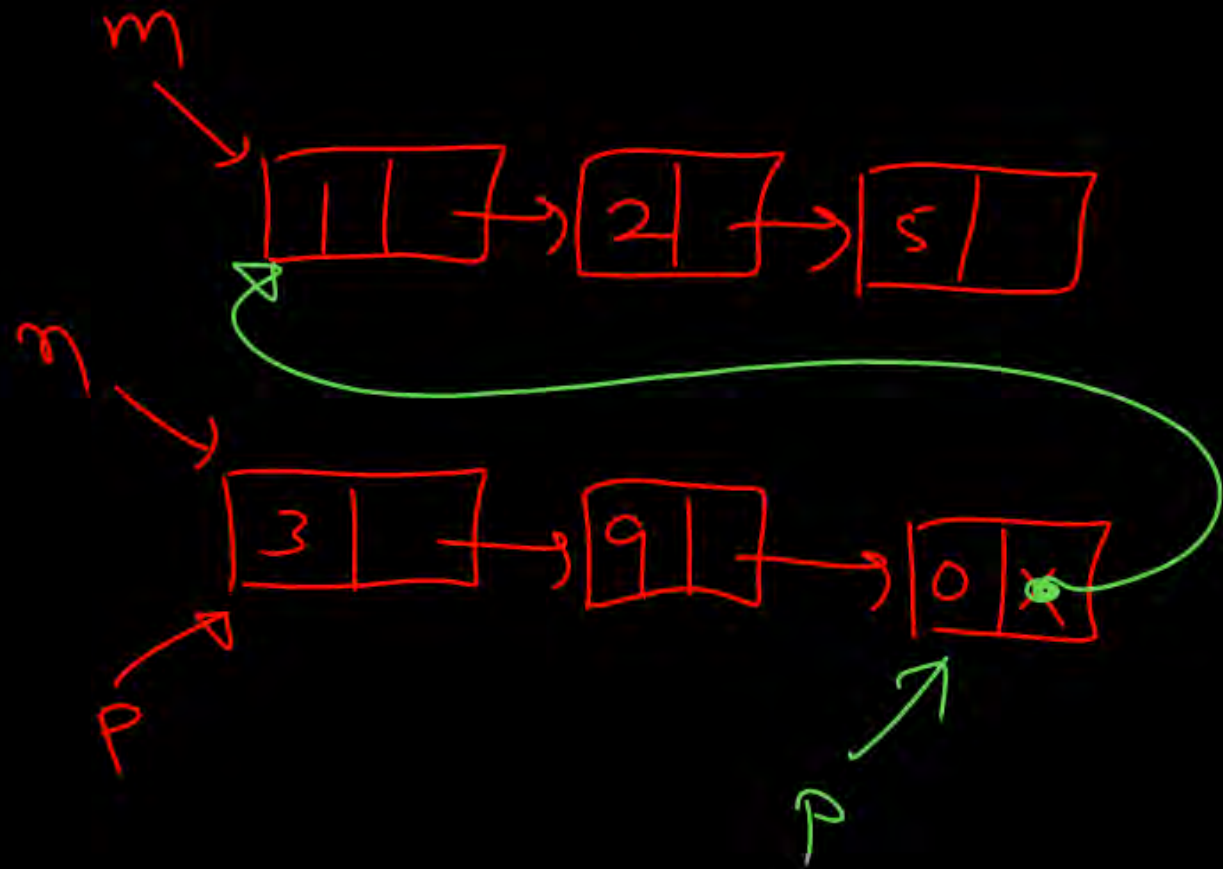
```
        p = p->next;
```

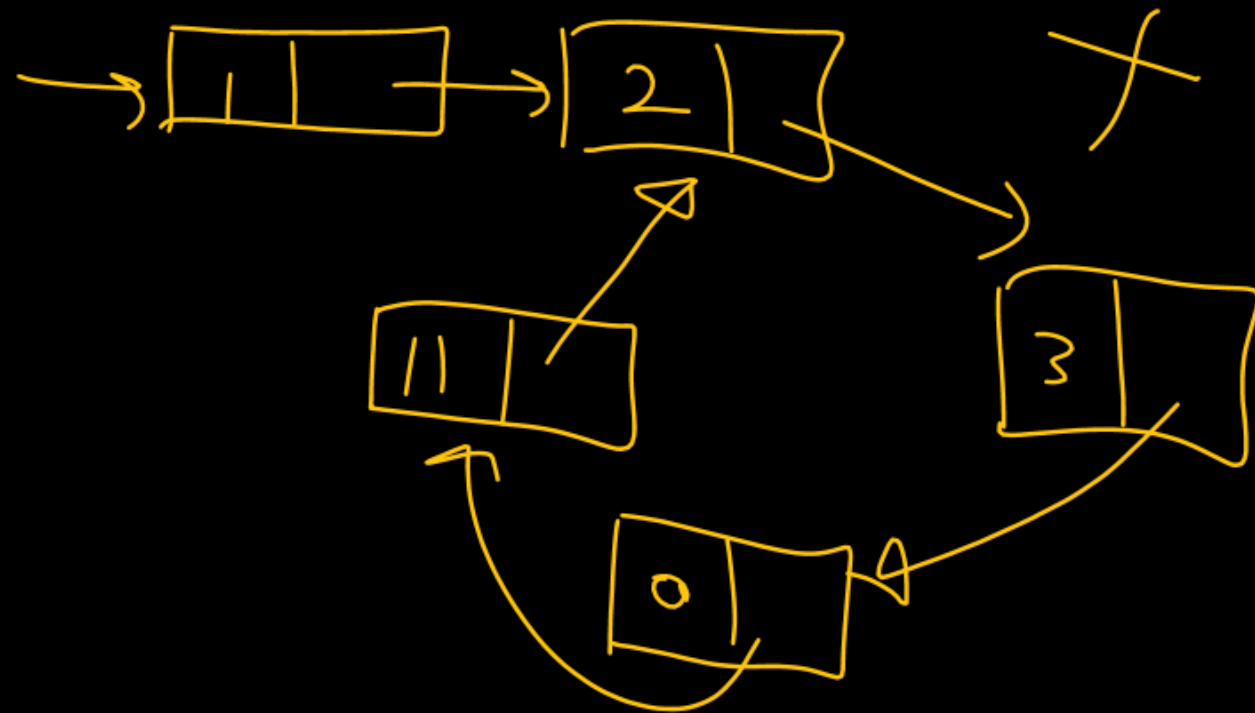
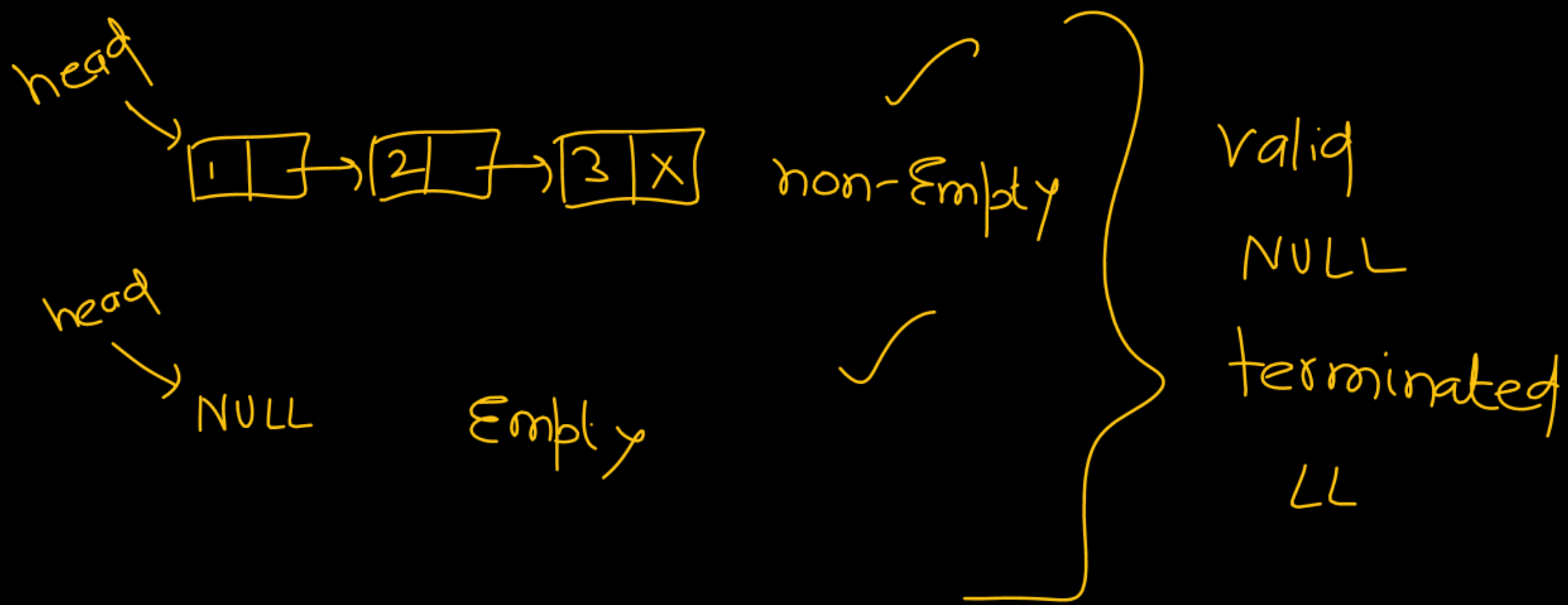
```
    }
```

```
    p->next = m;
```

```
}
```

Assuming that m and n points to
valid NULL terminated LL.
Invocation of join will.





7)

rearrange the ele.

1 → 2 → 3 → 4 → 5 → 6 → 7

```
void rearrange(struct node *list)
{
```

```
    struct node *p, *q;
    int temp;
```

```
    if (!list || !list->next) return;
```

```
    p = list; q = list->next;
```

```
    while(q) {
```

```
        temp = p->value;
```

```
        p->value = q->value;
```

```
        q->value = temp;
```

```
        p = q->next;
```

```
        q = p ? p->next : 0;
```

```
    }
```

```
}
```

list == NULL

list->next == NULL

5 min

7)

rearrange the ele.

2 1 4 3 6 5
1 → 2 → 3 → 4 → 5 → 6 → 7

2 1 4 3 6 5 7

↑
p q

NULL

5 min

P = list; q = list → next;

while(q) {

temp = P → value;

P → value = q → value;

q → value = temp;

P = q → next;

q = P ? P → next : 0;

}

swap

True/non-zero

q = P ? P → next : 0;

LL-7 08:30 PM

[Types of LL
PYQs

