

CS & IT ENGINEERING

Discrete maths
Mathematical logic



Lecture No.

06



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TOPICS TO BE COVERED

01 Inference Rule

02 Type 3 Questions in logic

03 Type 3 with Type 1

04 GATE QUESTIONS on 3 &1

05 Practice

All mothers are female.

Case 1:

$$\begin{array}{c} \text{All of } n \quad n \text{ is mother} \quad n \text{ is female.} \\ \hline \forall n. \quad \text{MT}(n) \quad F(n) \end{array}$$

Case 2:

$$\forall n [\text{MT}(n) \rightarrow F(n)]$$

BABA
female.
mother/female.

All mothers are female

$$\underline{\forall x [MT(x) \wedge F(x)]}$$

$$\cancel{\forall x [MT(x) \rightarrow F(x)]} \quad \text{True.}$$

All (\rightarrow)

$$F. \wedge F$$

$$F \rightarrow$$

True.

$$F \wedge F$$

$$F \rightarrow$$

True

$$T \wedge T$$

$$T \rightarrow T$$

T.

False

Some cats are black.

$$\text{Some of } x, \underline{x \text{ is cat}} \quad \underline{x \text{ is black.}} \quad \text{True}$$

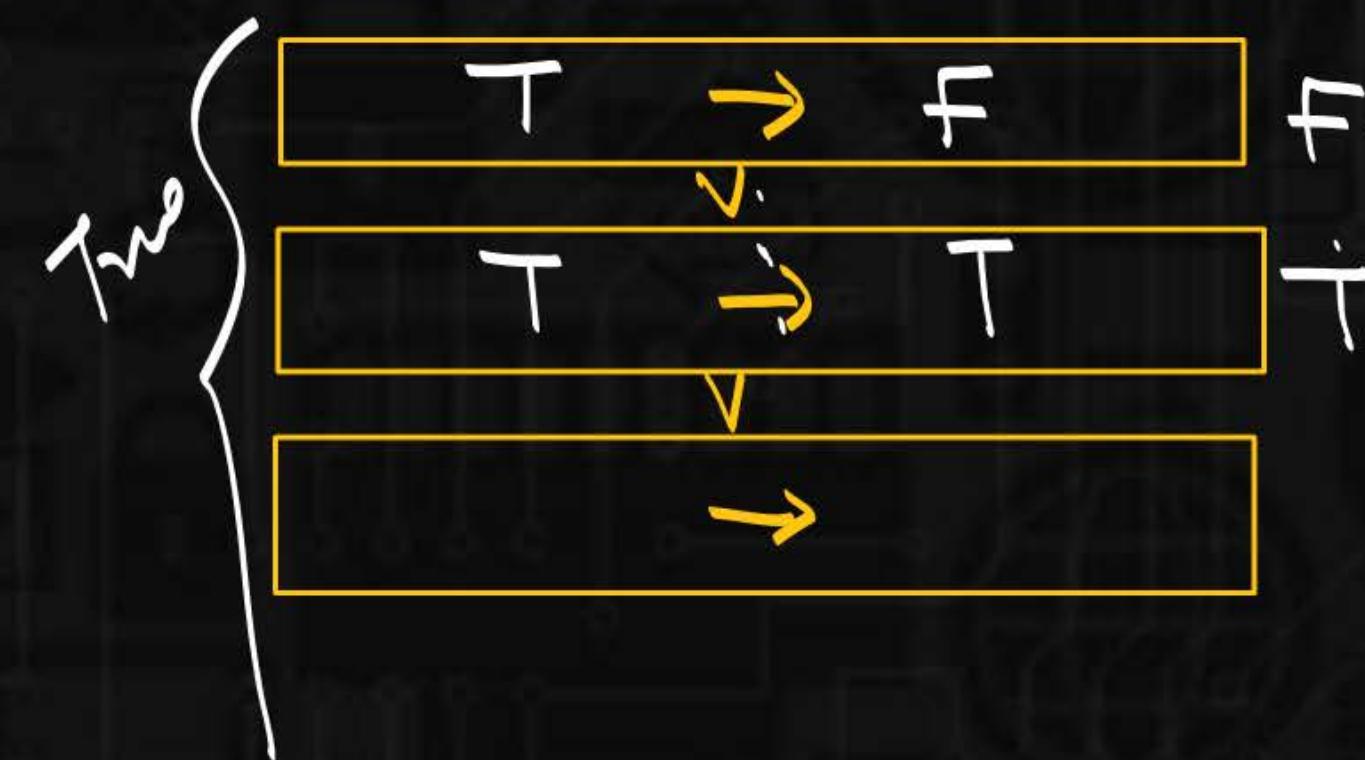
$$\exists x \quad CT(x) \quad BL(x)$$

- maggie/white.
- Alia/bl
- cow

$$\exists x [CT(x) \wedge bl(x)]$$



$$\exists x [CT(x) \rightarrow bl(x)]$$



Some cats are black.

False.

$$\text{Some of } x, \underline{x \text{ is cat}} \quad \underline{x \text{ is black.}}$$

$$\exists x \quad CT(x) \quad BL(x)$$

- Ali / w
- cow
- kat / w

$$\exists x [CT(x) \wedge bl(x)]$$

$$\exists x [CT(x) \rightarrow bl(x)]$$

X.
Some →

T T

F T

Some ∧

T T

F F

False

T \wedge F	F
F \wedge	F
T \wedge F	F

True

T \rightarrow F	F
F \rightarrow	

Same behaviour :

	\wedge	\vee	\rightarrow	\leftrightarrow
Some(T)	T	T	T	T.
Some(F)	F	F	F	F.

Some real no. are rationals (GATE)

WARNING:

English \rightarrow logical

we just convert
we don't find

the truth
 \therefore value.

A)
~~B)~~

$\exists x [\text{Real}(x) \wedge \text{Rational}(x)]$

C)

D)

All graphs are connected.

$$\forall n [\text{Graph}(n) \rightarrow \text{Connected}(n)]$$

1. All graphs are connected ✓
 $\forall n [G(n) \rightarrow C(n)]$

3. not (all graphs are connected)
 $\neg (\forall n [G(n) \rightarrow C(n)])$

2. All graphs are not connected ≡ 4. no graphs are connected
 $\forall n [G(n) \rightarrow \neg C(n)]$ $\forall n [G(n) \rightarrow \neg C(n)]$
not all $\neg \forall n [\rightarrow]$ {no/none} $\forall n [\rightarrow \neg]$

all of my friends are perfect

not all of my friends are
perfect

$$\neg \forall x [\rightarrow]$$

all of my friends are not perfect \equiv none of my friends are
perfect.

$$\forall x [\rightarrow \neg]$$

$P(n)$: $n > 0$

$Q(n)$: n is even

$\gamma(n)$: n is perfect square.

$S(n)$: n is \div by 4

$t(n)$: n is \div by 5

$\forall n [Q(n) \wedge \gamma(n) \rightarrow S(n)]$

$\exists n [P(n) \wedge Q(n)]$

1) There exist a positive int that is even.

2) if n is even then n is not divisible by 5
 $\forall n [Q(n) \rightarrow \neg t(n)]$

3) no even int is \div by 5

4) There exist even int \div by 5

5) If n is even and n is perfect square then n is \div by 4.

not all

$$\neg \forall x [\quad \rightarrow \quad]$$

no / none.

$$\forall x [\quad \rightarrow \neg \quad]$$

no graphs are
connected.

$$\forall u [G(u) \rightarrow \neg C(u)]$$

OR

$$\forall u [\neg G(u) \vee \neg C(u)]$$

$$\forall u [\neg [G(u) \wedge C(u)]]$$

$\exists u [G(u) \wedge C(u)]$

!

Gold and Silver ornaments are precious.

(GATE)

A) $\forall n [\underline{G(n)} \vee \underline{S(n)} \rightarrow P(n)]$

B) $\forall n [\boxed{\underline{G(n)} \wedge \underline{S(n)}} \rightarrow P(n)]$

{English:
mathematics}

$$\forall u \forall y [(u > y) \rightarrow (u - y > 0)]$$

$$\forall u \forall y [(u < y) \rightarrow \exists z (u < z < y)]$$

$$\forall u \forall y [|u| = |y| \rightarrow (y = \underline{\pm} u)]$$

P and Q are two propositions. Which of the following logical expressions are equivalent?

- I. $P \vee \sim Q$
- II. $\sim(\sim P \wedge Q)$
- III. $(P \wedge Q) \vee (P \wedge \sim Q) \vee (\sim P \wedge \sim Q)$
- IV. $(P \wedge Q) \vee (P \wedge \sim Q) \vee (\sim P \wedge Q)$

(GATE-08)

- (a) Only I and II
- (b) Only I, II and III
- (c) Only I, II and IV
- (d) All of I, II, III and IV

Establish the validity of the following arguments.

a) $[p \wedge (p \rightarrow q) \wedge (\neg q \vee r)] \rightarrow r$

b) $p \rightarrow q$

$$\neg q$$

$$\neg r$$

$$\therefore \neg(p \vee r)$$

c) $p \rightarrow q$

$$r \rightarrow \neg q$$

$$r$$

$$\therefore \neg p$$

Let p, q, r and s be four primitive statements. Consider the following arguments:

$$P: [(\sim p \vee q) \wedge (r \rightarrow s) \wedge (p \vee r)] \rightarrow (\sim s \rightarrow q)$$

$$Q: [(\sim p \wedge q) \wedge [q \rightarrow (p \rightarrow r)]] \rightarrow \sim r$$

$$R: [(q \wedge r) \rightarrow p] \wedge (\sim q \vee p) \rightarrow r$$

$$S: [p \wedge (p \rightarrow r) \wedge (q \vee \sim r)] \rightarrow q$$

Which of the above arguments are valid?(GATE - 04)

- (a) P and Q only
- (b) P and R only
- (c) P and S only
- (d) P, Q, R and S

d) $p \wedge q$
 $p \rightarrow (r \wedge q)$
 $r \rightarrow (s \vee t)$

$$\frac{\neg s}{\therefore t}$$

e) $p \rightarrow (q \rightarrow r)$
 $p \vee s$
 $t \rightarrow q$

$$\frac{\neg s}{\therefore \neg r \rightarrow \neg t}$$

Which one of the following is NOT equivalent to $p \leftrightarrow q$?

(GATE-15-Set1)

- (a) $(\sim p \vee q) \wedge (p \vee \sim q)$
- (b) $(\sim p \vee q) \wedge (q \rightarrow p)$
- (c) $(\sim p \wedge q) \vee (p \wedge \sim q)$
- (d) $(\sim p \wedge \sim q) \vee (p \wedge q)$

The Simplest form of $(p \wedge (\sim r \vee q \vee \sim q)) \vee ((r \vee t \vee \sim r) \wedge \sim q)$ is

- (a) $p \wedge \sim q$
- (b) $p \vee \sim q$
- (c) t
- (d) $(p \rightarrow \sim q)$

Let p, q, r and s be four primitive statements. Consider the following arguments:

$$P: [(\sim p \vee q) \wedge (r \rightarrow s) \wedge (p \vee r)] \rightarrow (\sim s \rightarrow q)$$

$$Q: [(\sim p \wedge q) \wedge [q \rightarrow (p \rightarrow r)]] \rightarrow \sim r$$

$$R: [(q \wedge r) \rightarrow p] \wedge (\sim q \vee p) \rightarrow r$$

$$S: [p \wedge (p \rightarrow r) \wedge (q \vee \sim r)] \rightarrow q$$

Which of the above arguments are valid?

(GATE - 04)

- (a) P and Q only
- (b) P and R only
- (c) P and S only
- (d) P,Q,R and S

The Simplest form of

$$(p \vee (p \wedge q) \vee (p \wedge q \wedge \sim r)) \wedge ((p \wedge r \wedge t) \vee t)$$

(a) $p \wedge t$

(b) $q \wedge t$

(c) $p \wedge r$

(d) $p \wedge q$

P
W

$$S_1: \{(\sim p \rightarrow (q \rightarrow \sim W)) \wedge (\sim S \rightarrow q) \wedge \sim t \wedge (\sim p \vee t)\} \rightarrow (w \rightarrow s)$$

$$S_2: \{(q \rightarrow t) \wedge (s \rightarrow r) \wedge (\sim q \rightarrow s)\} \rightarrow (\sim t \rightarrow r)$$

The statement formula $\{(a \rightarrow c) \wedge (b \rightarrow d) \wedge (c \rightarrow \neg d)\} \rightarrow (\neg a \vee \neg b)$ is

- (a) satisfiable but not-valid
- (b) valid
- (c) not satisfiable
- (d) none of these

The statement formula

$$\{((\sim p \vee q) \rightarrow r) \wedge (r \rightarrow (s \vee t)) \wedge (\sim s \wedge \sim u) \wedge (\sim u \rightarrow \sim t)\} \rightarrow p$$
 is

- (a) valid
- (b) not satisfiable
- (c) satisfiable but not valid
- (d) none of these

