

CS & IT ENGINEERING

Discrete maths
Graph Theory



Lecture No.

06

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TOPICS TO BE COVERED

01 Definition In Connectivity

02 Connected vs Disconnected

03 Range of Edges

04 Concepts of tree

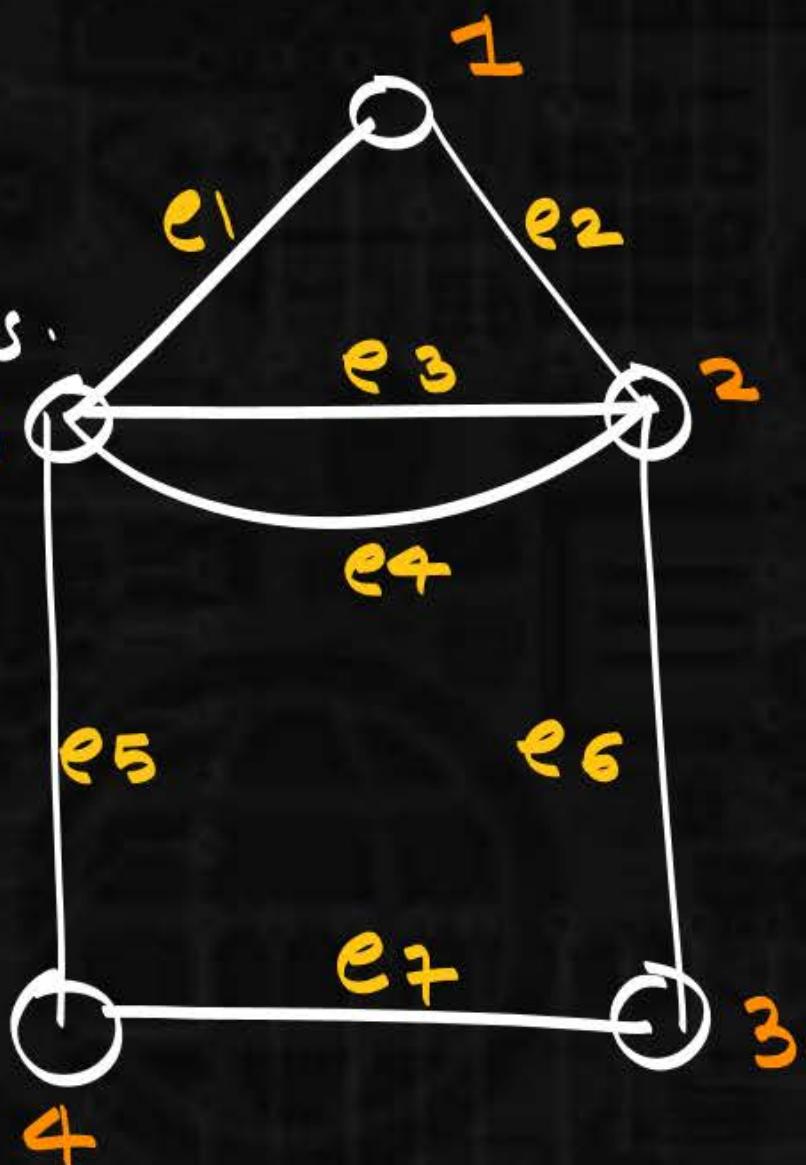
05 Connectivity theorem

Connectivity in Graphs

Walk: alternating sequences of vertices & edges.

R.V | R.E

1 e₂ 2 e₃ 5 e₃ 2 e₆ 3.



Trial:

R.V | R.E

1 e₂ 2 e₃ 5 e₄ 2 e₆ 3.

Path:

R.V | R.E

1 e₂ 2 e₆ 3.

Connectivity in Graphs



When path is not available,
at least 1. pair of vertices.

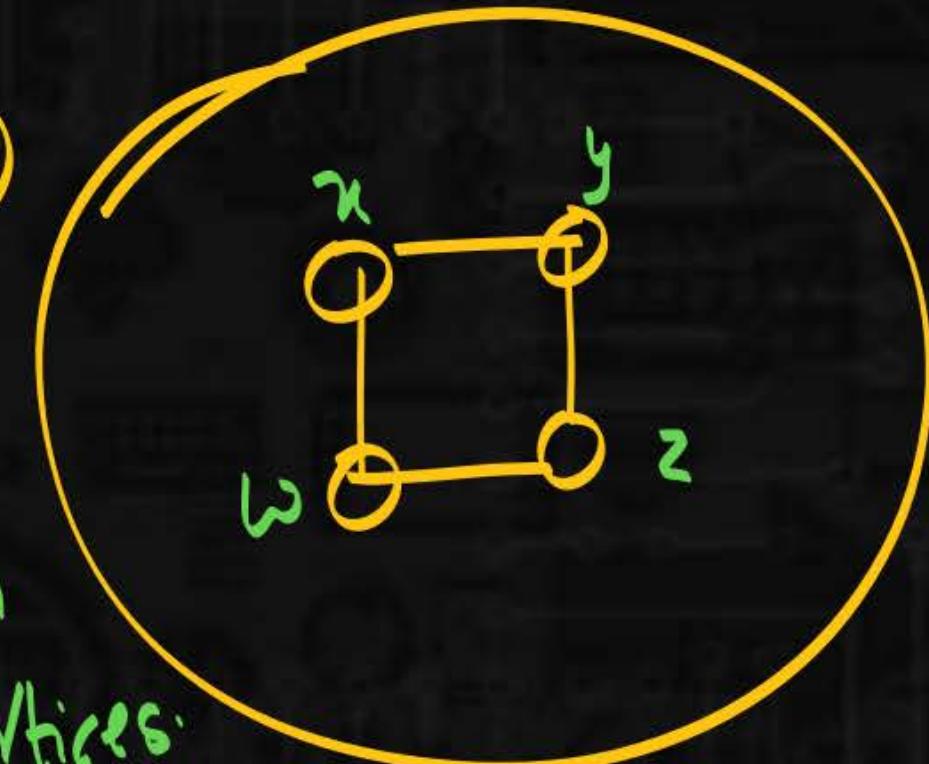
Graph:

Connected.

$$G = (V, E)$$

$$|V| = 4$$

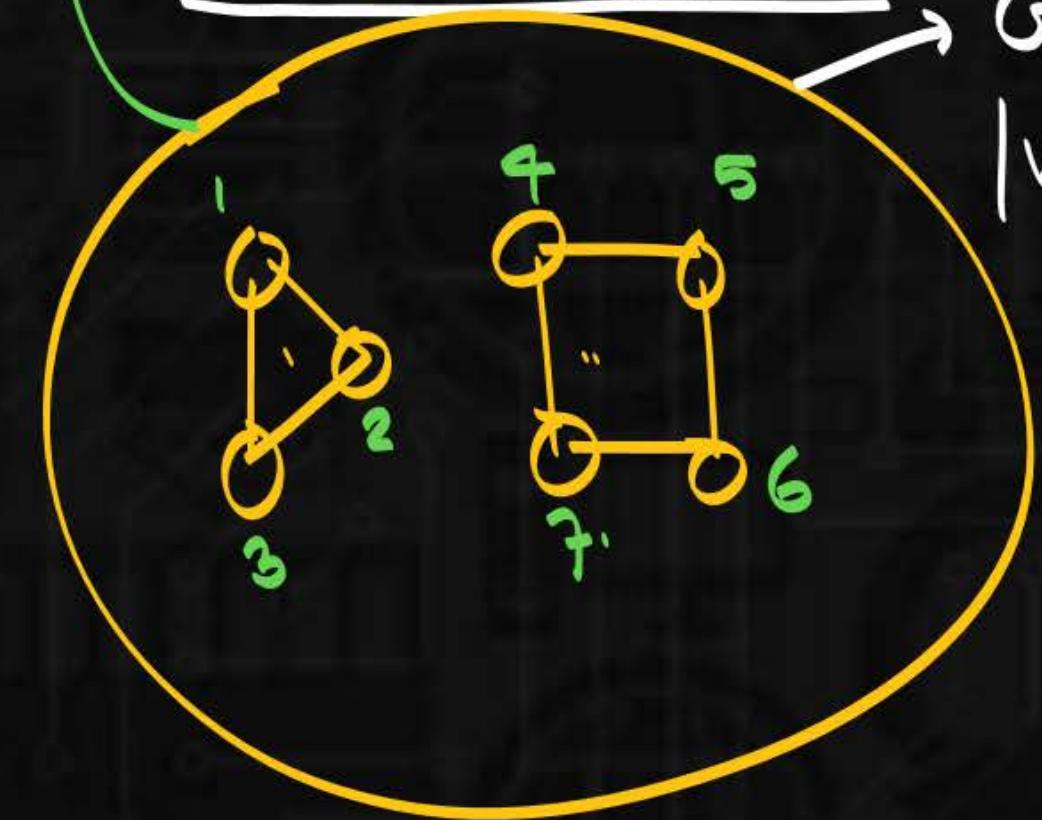
When path is
available betn
all pair of vertices.



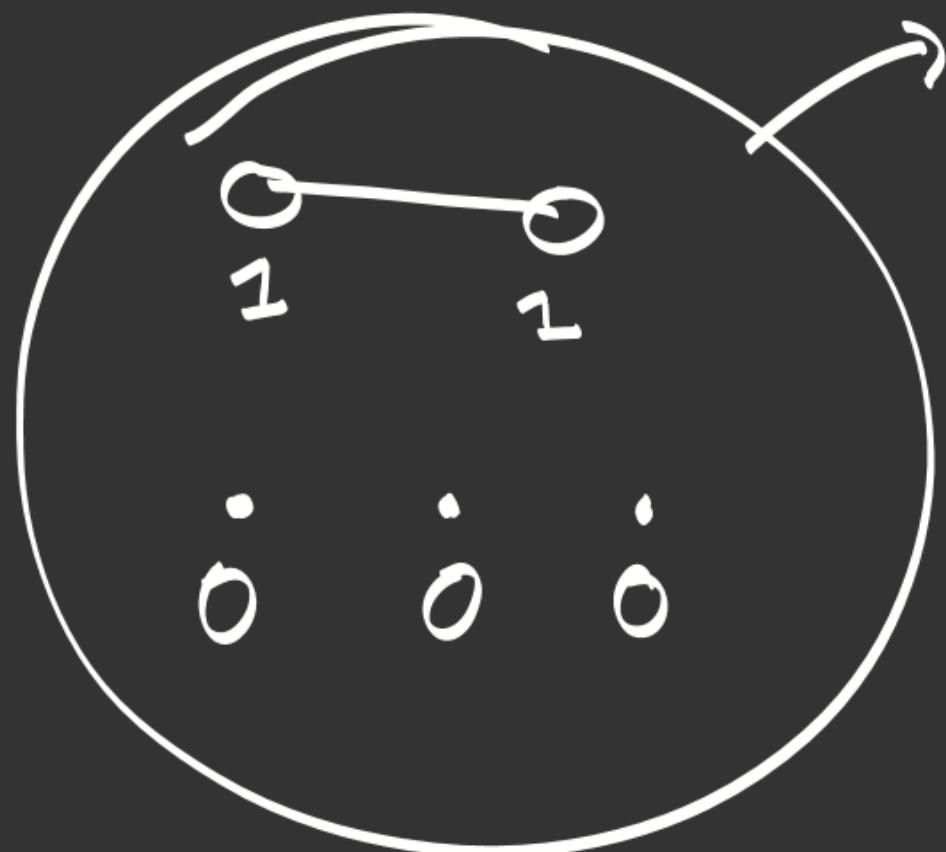
Disconnected.

$$G = (V, E)$$

$$|V| = 7$$



1, 1, 0, 0, 0



disconnected

Connectivity in Graphs

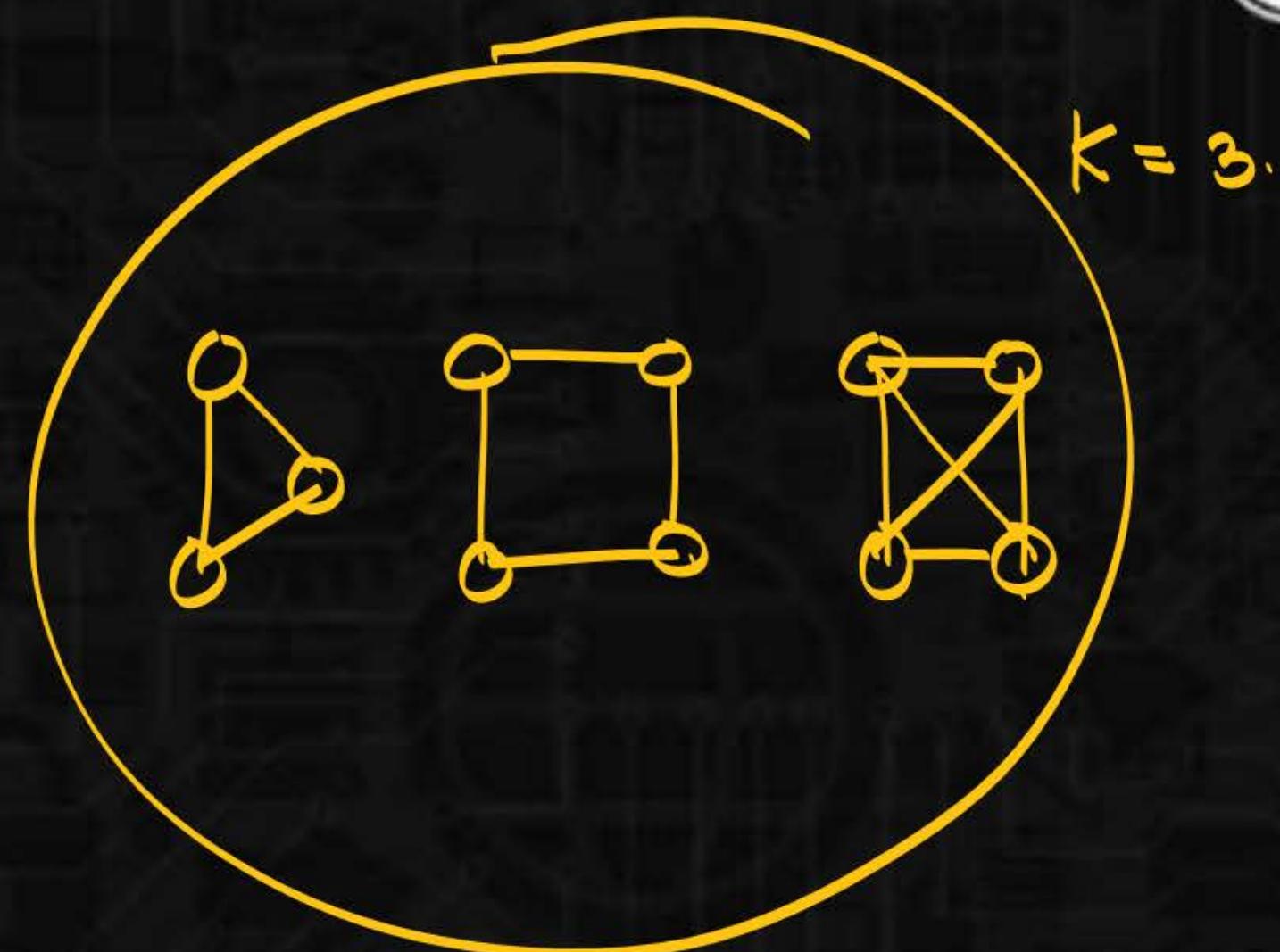
Disconnected graph contains

Connected subparts.

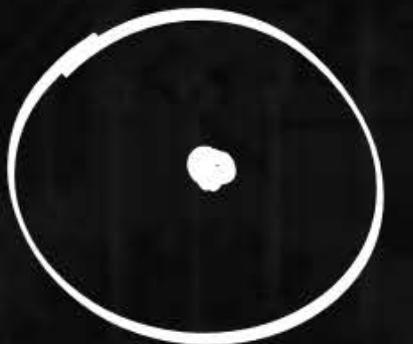
Component (K)

Connected. ($K=1$)

Disconnected ($K \geq 2$)



Connectivity in Graphs



$k = 1 \rightarrow$ connected
Graph.



$k = 2$.
D.C.



$k = 4$.

Connectivity in Graphs

Consider a Graph where vertices are represented as {1 ... 100}

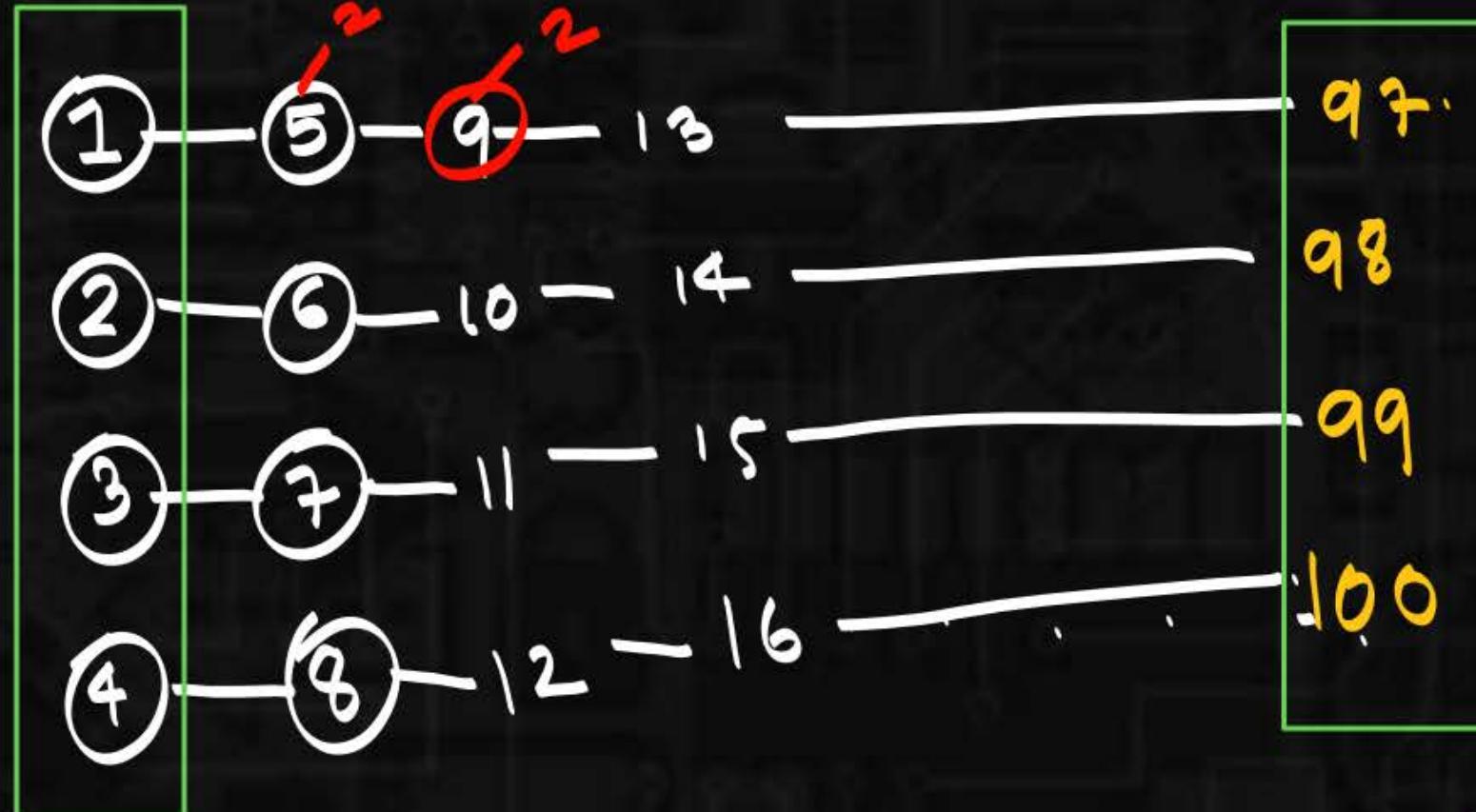
two vertices are connected with each other.

when $|x-y|=4$, what will be no. of components?

Total edges?

$$4 \times 1 + 4 \times 1 + 92 \times 2 = 2e$$

$$e = 96$$



Connectivity in Graphs

A is non empty set. $|A|=n$.

(GATE)

Graph vertices are represented as elements of (powerset of A)

two vertices are connected if there intersection will give exactly 2 elements, what will be no. of

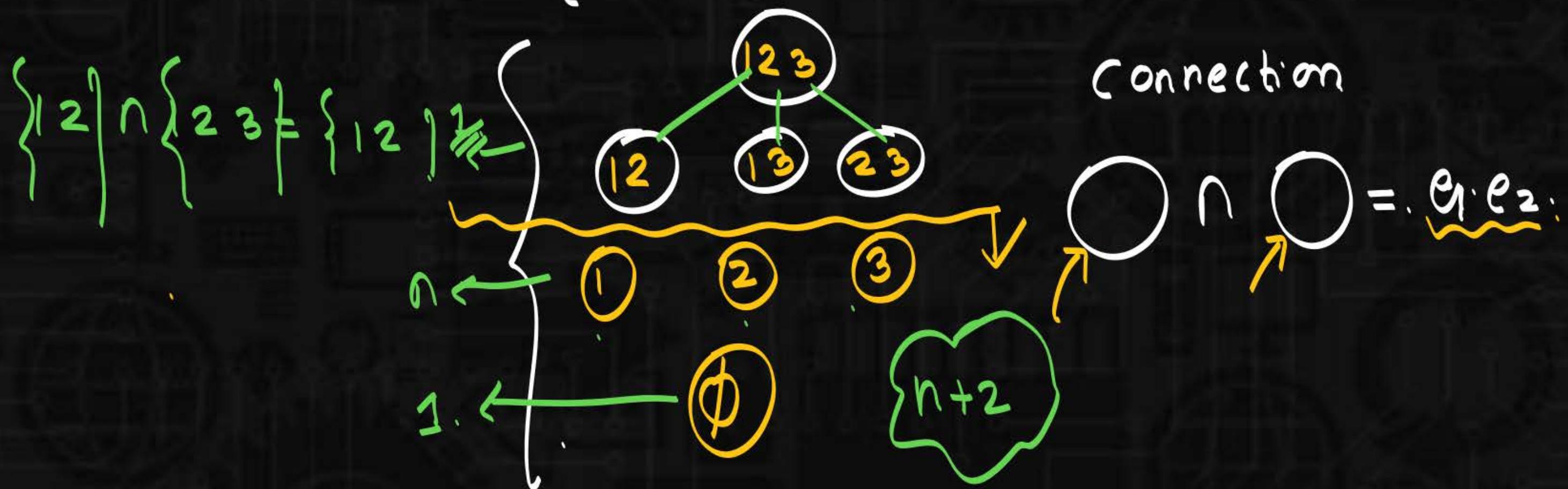
components? ($n \geq 3$)

Connectivity in Graphs

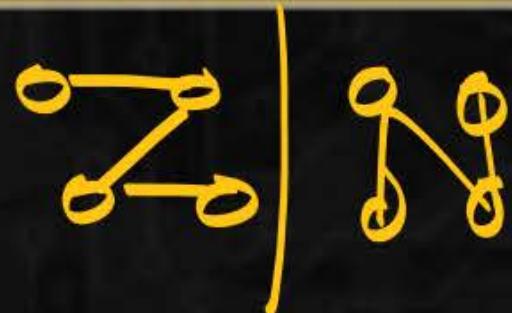
$$A = \{1, 2, 3\}, |A| = n$$

Total no. of vertices = 2^n
Total elements in 2^A = 2^n

$$\text{powerset}(P(A)/_{2^A}) : \left\{ \emptyset, \{1\}, \{2\}, \{3\}, \dots, \{1, 2, 3\} \right\}$$



Connectivity in Graphs



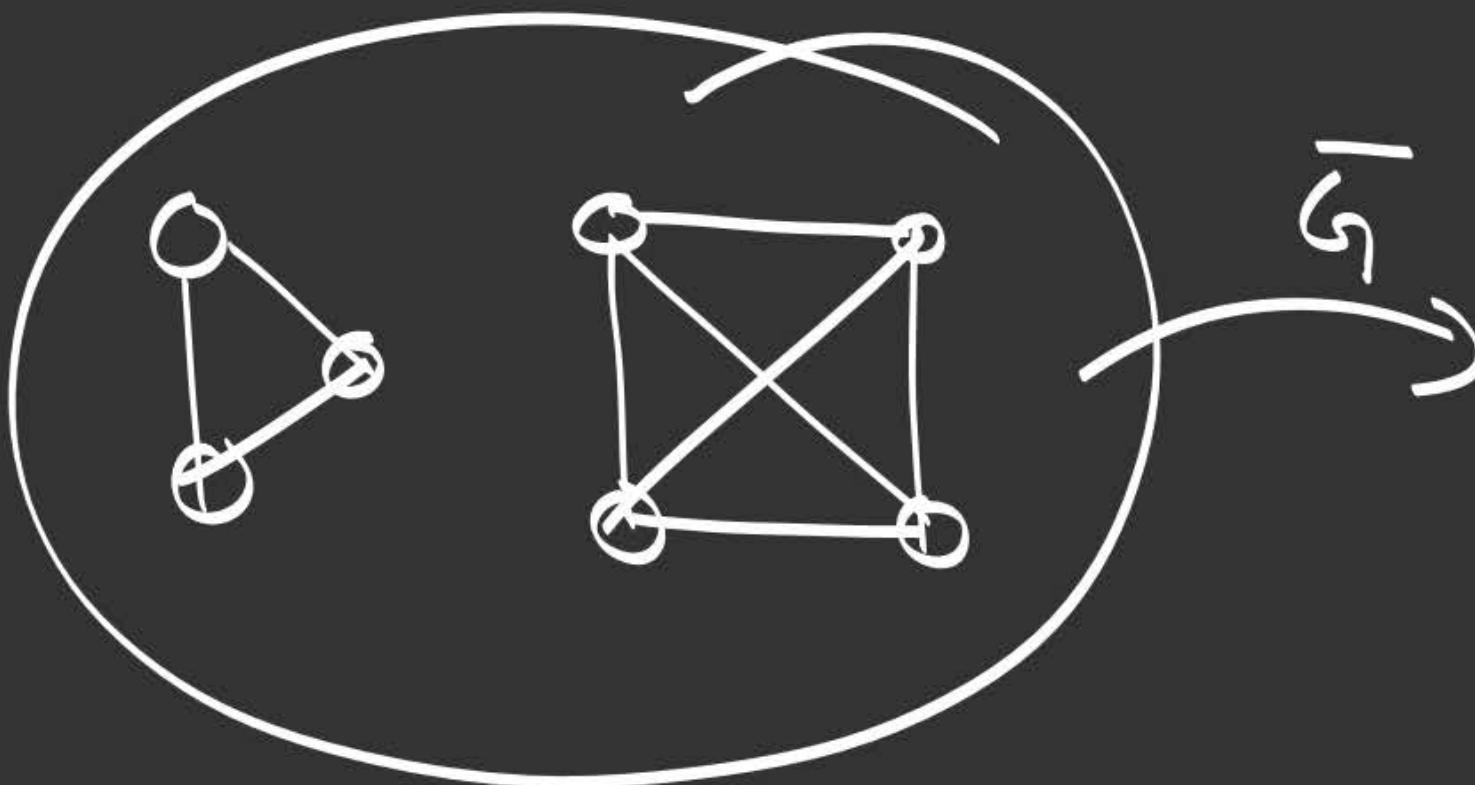
$K = 2$.

S₁ : if G is connected then \bar{G} will also be connected (false)

S₂ : if G is disconnected then \bar{G} will be connected.

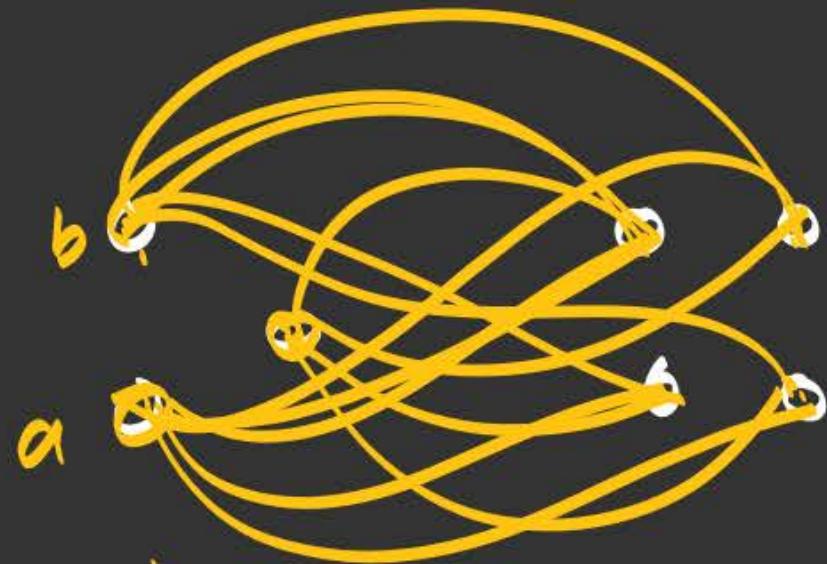
Assumption





$$G = (V, E)$$

$$|V| = 7$$



Path is available.
all pair.
Connected.

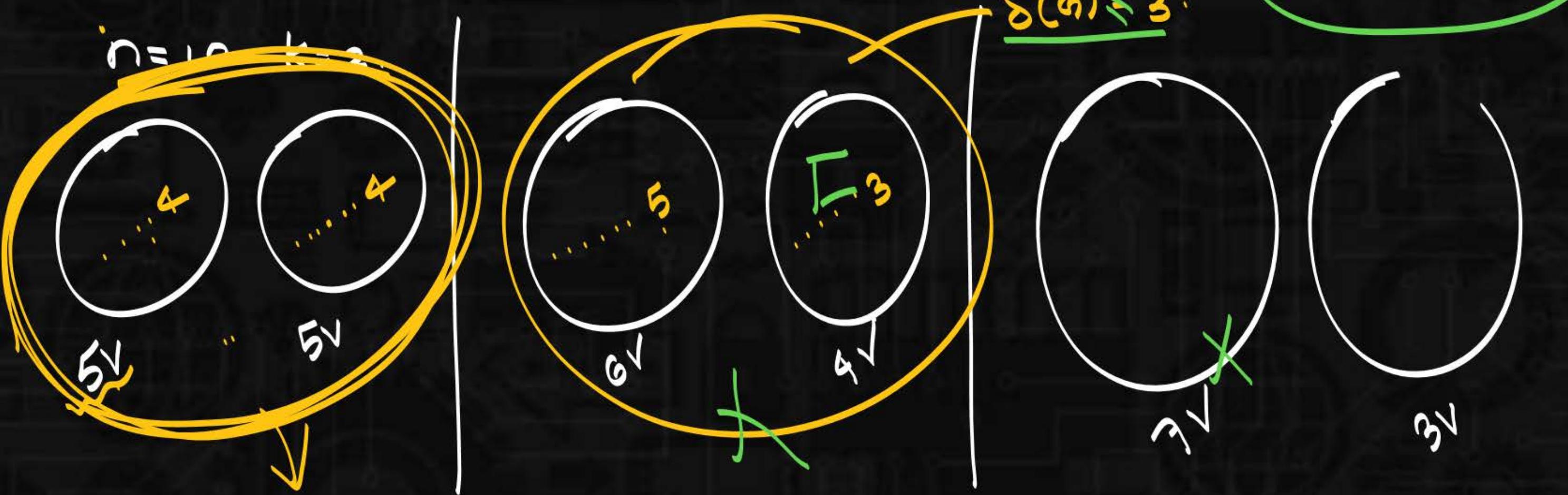
Connectivity in Graphs

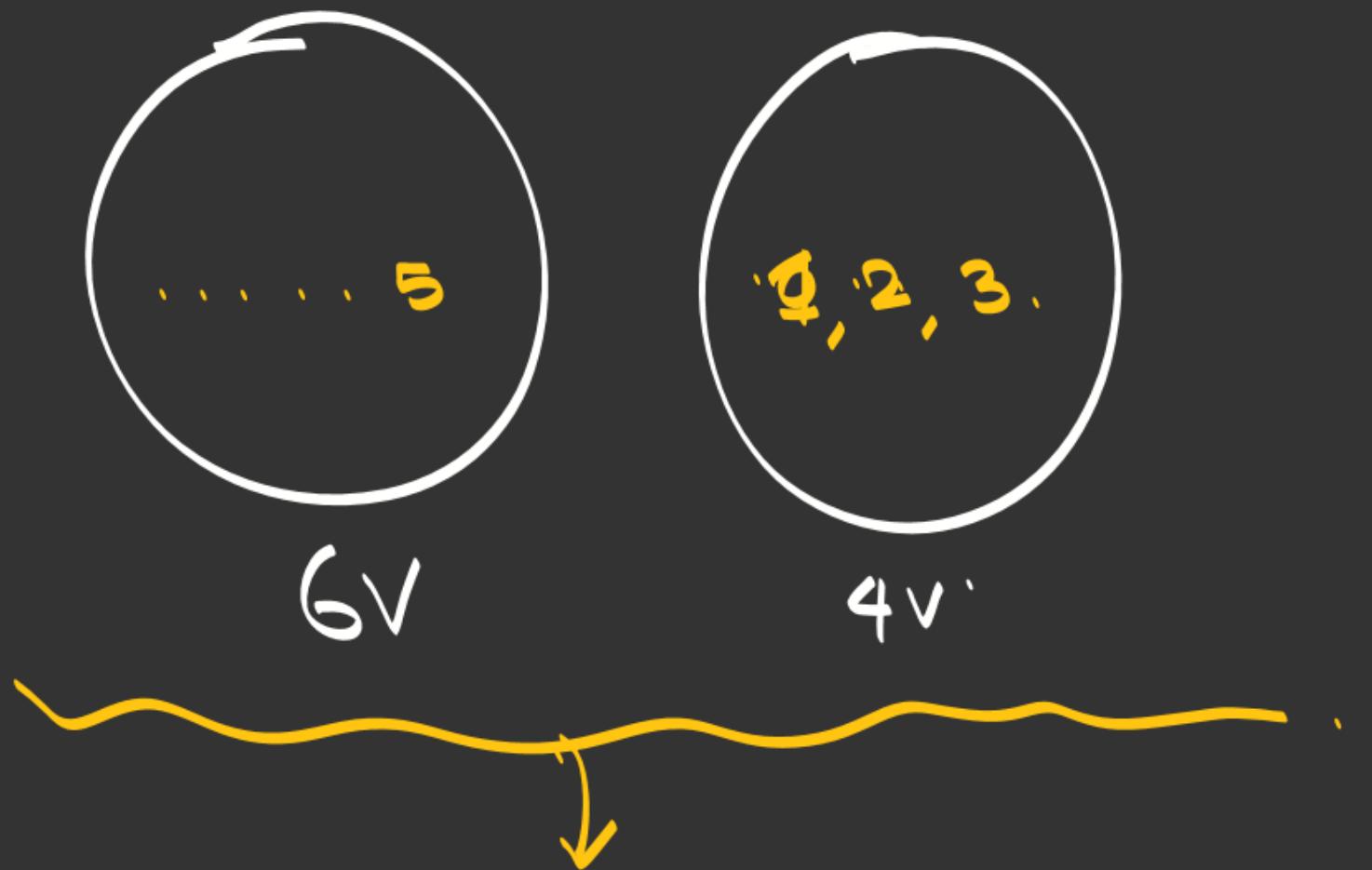
If $\delta(G) \geq 5$ with $n=10$ then G is connected. (True)

Assumption: this is Disconnected with $K=2$.

$$-\delta(a) \leq 3$$

no division is
possible.





$$\delta(G) \geq 5$$

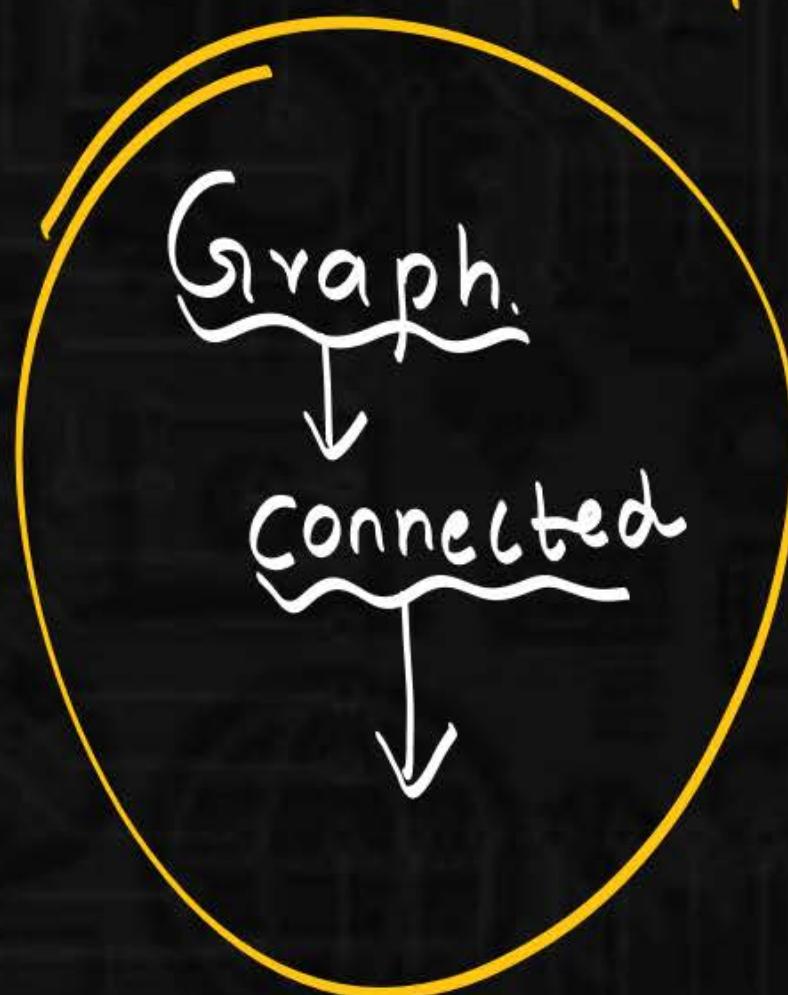
minimum degree ≤ 3

Connectivity in Graphs

Thm: if $\delta(G) \geq \frac{n-1}{2}$ then it is connected. (vice versa is not True).

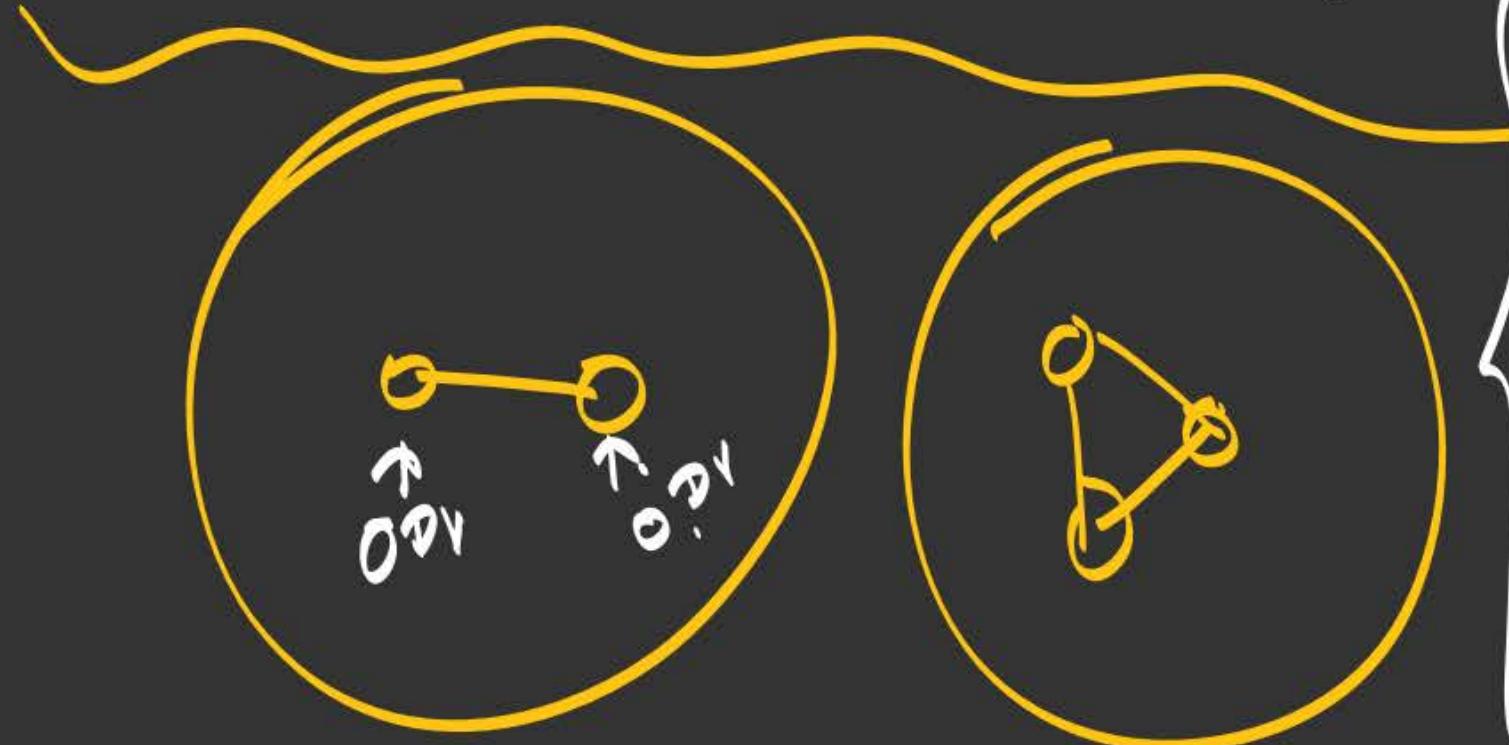
Connectivity in Graphs

Consider a Graph is having exactly 2 odd degree vertices
Check path is available betn those vertices ?.



Graph \rightarrow Disconnected.





In Disconnected Graph,
2 exactly Odd degree vertex.
will be in same component (Thm2)
hence path is available b/w them.

Thm2:

no. of odd degree
vertices should be even.

Connectivity in Graphs

Range of edges ($K = 1$)

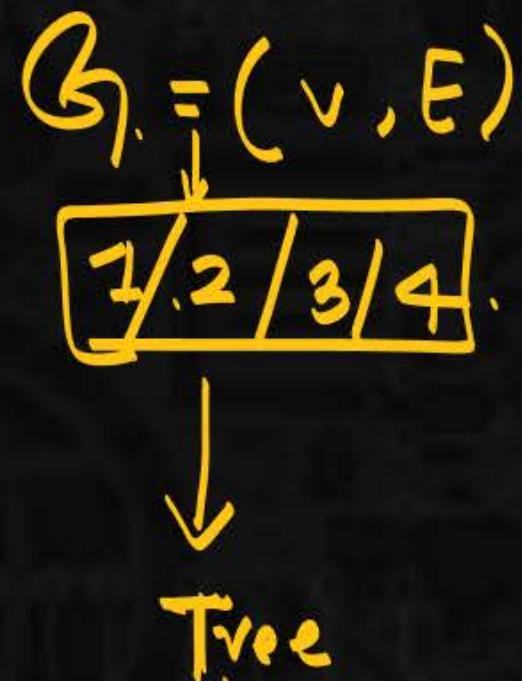
$$1 \leq 2 \leq 3 \leq 4$$

$$\left\{ \begin{array}{l} 1. \quad \left\{ \begin{array}{l} n-1 \leq e \leq \frac{n(n-1)}{2} \\ + \end{array} \right. \rightarrow (K_n) \\ \text{connected} \end{array} \right.$$

Tree $\left\{ \begin{array}{l} 2. \text{ connected acyclic graph.} \end{array} \right.$

3. unique path b/w all pair of vertices

4. minimally connected



$n = 4 \ e = 0$



$n = 4 \ e = 1$



$n = 4 \ e = 2$



$n = 4$



$e = 3$

$e = 3$ (min.)



$e = 3$

no cy

$n = 4 \ e = 3$ connected
acycl. Graph.



Connectivity in Graphs

Disconnected.

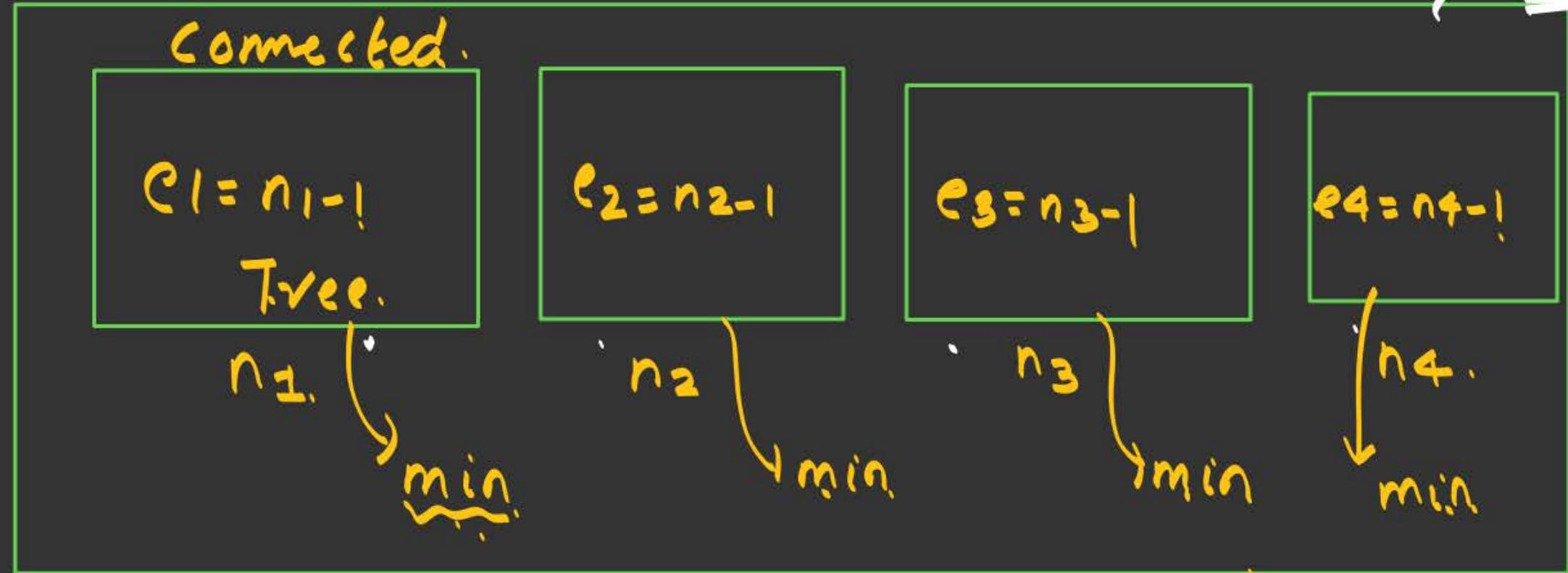
Range of edges ($K \geq 2$)

Forest


$$n - K \leq e \leq \frac{(n - K)(n - K + 1)}{2}$$

max no. of edges

\rightarrow forest:



$$\begin{aligned} \text{Total edges} &= n_1 - 1 + n_2 - 1 + n_3 - 1 + n_4 - 1 \\ &= (\underbrace{n_1 + n_2 + n_3 + n_4}) - 4. \\ n - 4 &= n - K. \end{aligned}$$

Total vertices = n .

1st component $\rightarrow n_1$.

$2n$ " n_2 .

Component = ($K = 4$)

forest = collection of Trees.

no. of edges
in Forest = $n - K$.

Connectivity in Graphs

$G = (V, E)$ with 10 vertices & 3 components.
min & max no. of edges?

Ans: 7, 28.

