

Design and Analysis of Algorithms

Tutorial - 7

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Section - CE

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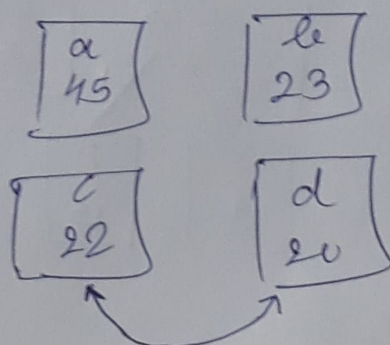
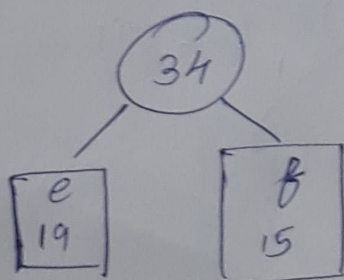
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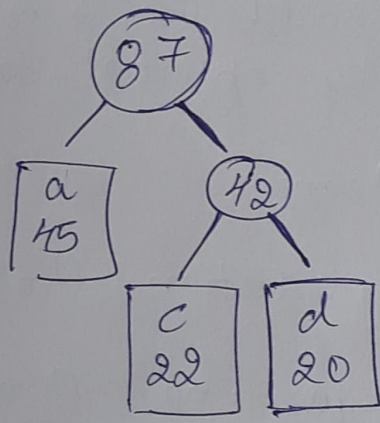
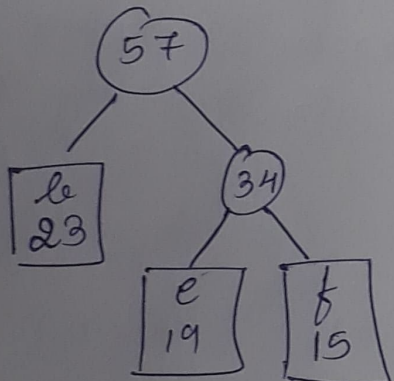
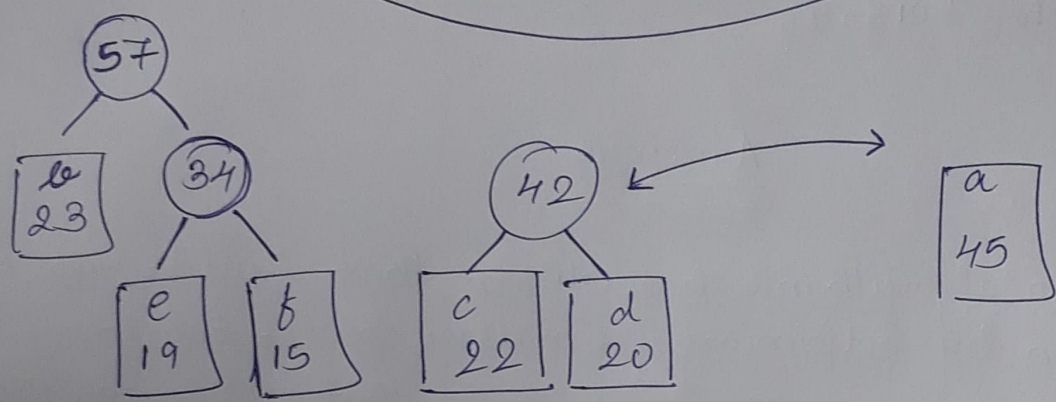
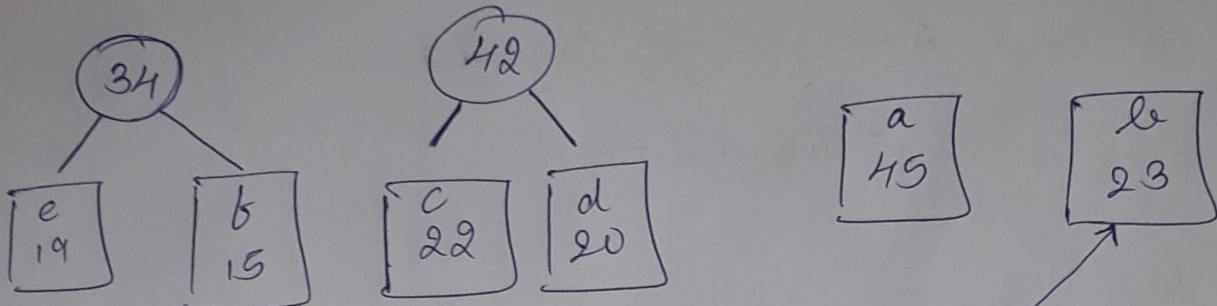
Answers

Ans 1 It is an algorithmic paradigm that builds up a solution by adjoining smaller pieces together, always choosing the next piece that offers the most obvious and immediate benefit.
We should use greedy approach whenever a locally optimal solution is also globally optimal.

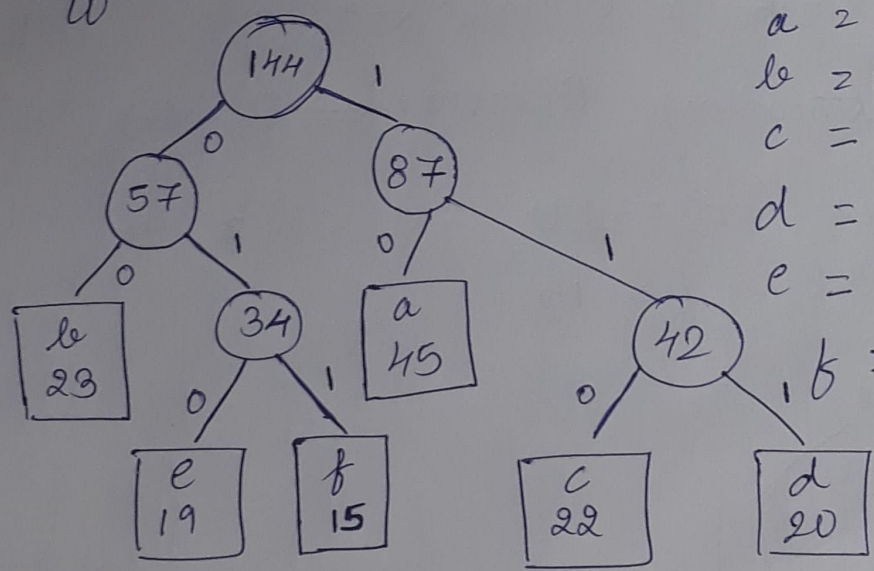
<u>Ans 2</u>	Name	TC	SC
Activity Selection		$O(m \log m) \leftrightarrow O(m)$	$O(m)$
Job Sequencing		$O(m^2) \leftrightarrow O(m \log m)$	$O(m)$
Fractional Knapsack		$O(m \log m) \leftrightarrow O(m)$	$O(m)$
Huffman Encoding		$O(m \log m) \leftrightarrow O(\log m)$	$O(m)$

Ans 3
 $a = 45$ $d = 20$ $le = 23$ $e = 19$
 $c = 22$ $f = 15$





Final Huffman Tree



- a = 10
- b = 00
- c = 110
- d = 111
- e = 010
- f = 111

$$\text{Total bits used} = (45 \times 2) + (23 \times 2) + (22 \times 3) + (20 \times 3)^3 + (19 \times 3) + (15 \times 3) = \underline{\underline{364 \text{ bits}}}$$

Ans 4 A 2-tree is used to implement Huffman encoding algorithm. It is a binary tree where every node has either 2-child or no child.

Applications of Huffman Encoding

- Data compression in long files without any loss
- To implement traffic routes with traffic magnitude.

Ans 5

V	10	5	15	7	6	18	3
W	2	3	5	7	1	4	1
V/W	5	1.67	3	1	6	4.5	3

$$k = 15 - 1 - 2 - 4 - 5 - 1 - 2 = 0$$

$$\text{Profit} = 30 + 10 + 18 + 15 + 3 + 30.34 = 79.34$$

V	6	10	18	15	3	5
W	1	2	4	5	1	3
V/W	6	5	4.5	3	3	1.67

Ans 6. Fractional Knapsack: It is using a greedy approach as we have divided our profits to the smallest unit possible & then builds upon it. 31

Huffman Encoding - It is using the greedy approach as we have divided our profits to the smallest unit possible & then builds upon it.

Huffman Encoding - It is using the greedy approach as it always places the node with the lower frequency further from the parent node.

Ans 7.

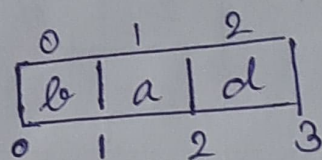
Start	1	2	0	6	9	10
End	3	5	7	8	11	12
Index	0	1	2	3	4	5

Tales to do = [0], [3], [4] or [5]

i.e. \Rightarrow Max = 4

Ans 8

	Profit	Deadline
a	20	2
b	15	2
c	10	1
d	5	3
e	1	3

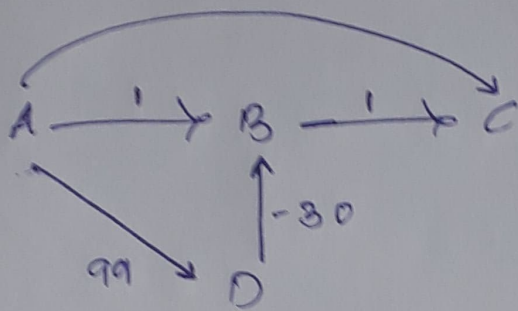


$$\text{Profit} = 20 + 15 + 5$$

$$= \underline{\underline{40}}$$

- Ans 9 Times when not to use greedy approach
- When the approach involves a lot of assumptions, such as "pick always the ..."
 - We should avoid greedy approach on complex implementation.
 - When we are making performance-critical applications.

Eg. Dijkstra's algorithm is very unoptimised for graphs with negative edges.



We can't find the distance of the pair $[A, C]$ it gives 0, though it is -200

Ans 10 Normally, the time complexity of Job sequencing is $O(n^2)$ but we can improve it using a Priority Queue, made of Max Heap.

Algorithm

- ① Sort the jobs based on deadlines.
- ② Iterate the end & calculate the available slots between two consecutive deadlines include all data in Max-Heap.
- ③ If there are slots available & there are jobs in the Max Heap, include the job ID with max profit & deadlines in the result.

④ Sort the array based on deadlines.

Time complexity: $O(m \log(m))$

Space complexity: $O(m)$