

ADS

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Roll no: 127

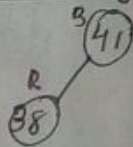
13. If we want to delete the Red black tree element in 8, 12, 19, 31, 38, 41. the insertion process must be take place in 41, 38, 31, 19, 12, 8.
So the insertion steps are given below.

Insertion

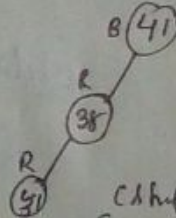
Inserting 41



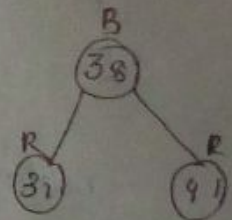
Inserting 38



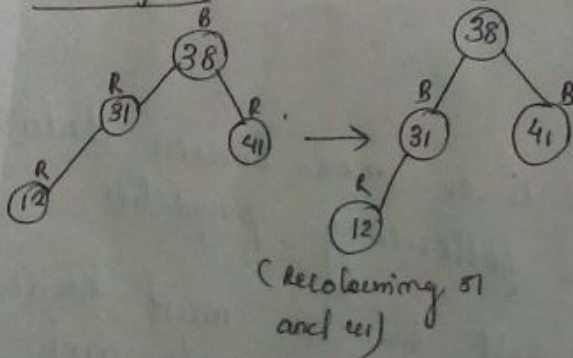
Inserting 31



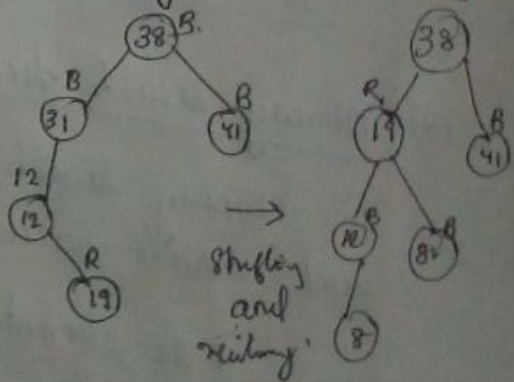
(shifting)
(recoloring)



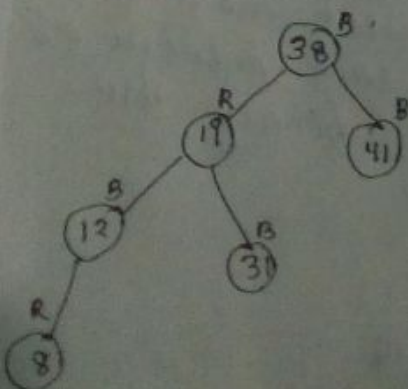
Inserting 12



Inserting 19

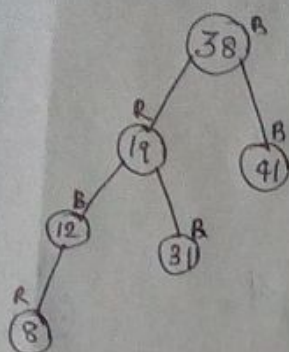


So the final tree is

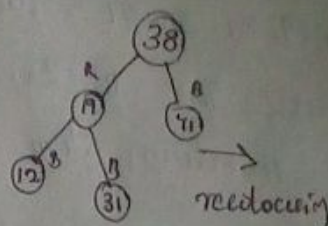
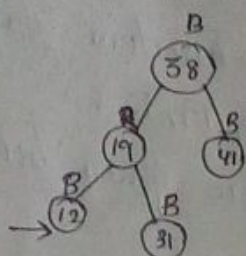


Now we are deleting the Red black tree elements in 8, 12, 19, 31, 38, 41.

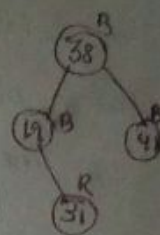
deleting 8



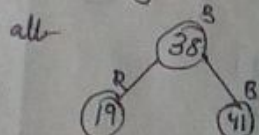
deleting 12



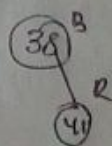
red colouring



deleting 19



deleting 31



deleting 38



deleting 41

Null tree.

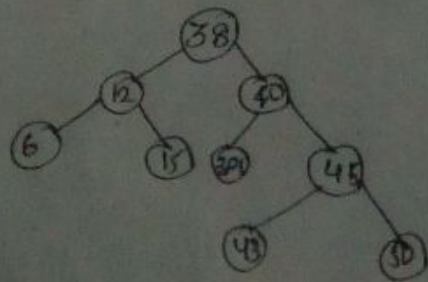
38 deleted.

12) Binary search tree.

Binary search tree is a non-linear data structure and which have the following properties

* The elements in left sub tree must have the value less than that of the root node, and the right elements must have the greater value than the root node. the parent node has atleast two children. and also, each sub tree of binary search tree is also a binary tree.

for example:



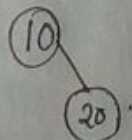
for example, if we want to insert the following elements in a binary search tree, we must follow the following steps

elements: 10 20 15 40 5

Step 1: (10)

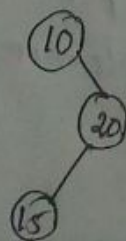
Step 2: Inserting 20.

20 is greater than 10 so it is in the right sub tree of 10



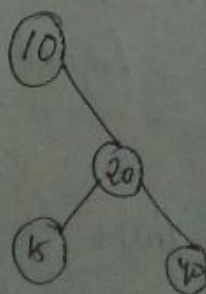
Step 3: Inserting 15.

15 is also greater than 10, and it is also in right sub tree of 10, and less than 20 - right left sub tree of 20



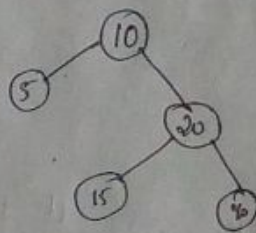
Step 4:

40 is the largest element in the list. so it is in the right sub tree of 20



Step 5:

to inserting 5, it is less than that of the root node. So it is the left sub tree.



the final binary search tree is given above.

Insertion steps:

Step 1: start.

Step 2: Insert the root element

Step 3: Compare the next element with root. If it is less than the root, it must be placed left side of the root, else place the element in right subtree.

Step 4: stop.

11. a) Insertion beginning and end mid - Singly linked list and doubly linked list

Singly linked list - beginning insertion

Step 1: start

Step 2: check if there is already a node exist

Step 2: if there is no node, then create a new node and set it as ptr

step 3: set ~~the~~ ~~new~~ ~~node~~ ~~to~~ ~~ptr~~ ~~in~~

$ptr = ptr \rightarrow next;$ (to store the address of the 1st node).

step 4: set new node,
insert the value to that node.

$new_node \rightarrow data = value$

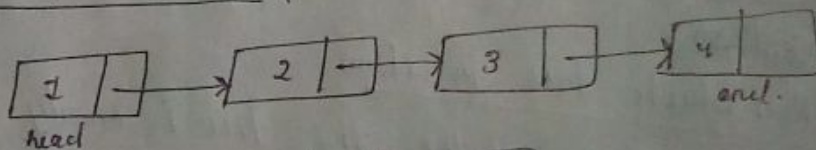
steps: if there is no element there then set,

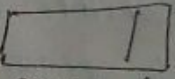
$new_node \rightarrow next = head$

$head = new_node$.

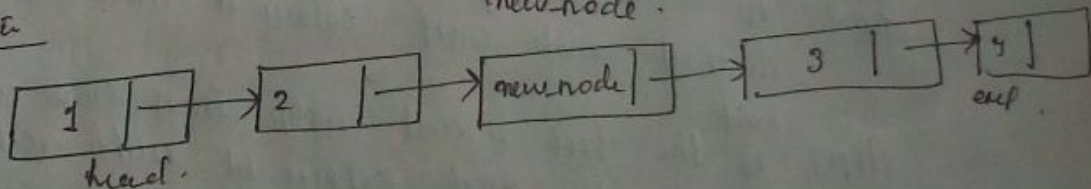
step 5: exit.

Inserting at middle



add  new-node.

after



Algorithm

step 1: start

step 2: check the ~~element~~ ~~is~~ list is empty.
If it is empty, then create a new node
and follow the above mentioned steps

step 3: If it is not empty, then create a new
node with attributes data and links.

step 4: If the list ~~empty~~ is not empty and the
list has only one element, then calculate
the size of the list and divide it by

2 To get mid-point of the list
step 5: define a current node and it will point
to the current new node.

step 6: define another node temp which points to
node of next to current.

step 7: the new node will be inserted after
current and before temp.

step 8: stop

Insertion in a doubly linked list

Insertion

10. Red black tree Insertion

step 1: check whether the tree is empty or not

step 2: if it is empty then create a new node
as root node and colour it with black.

step 3: ^{exit} if the tree is not empty, then create a
new node and colour it with red.
if the parent's node is black then exit.

step 4: if the parent node is ~~at~~ Red then
check the colour of parent's sibling or
new node.

step 5: the parent's sibling is ~~not~~ black
or null, then rotate and recolor it.

step 6: if parent's sibling is red then recolor.

step 7: exit

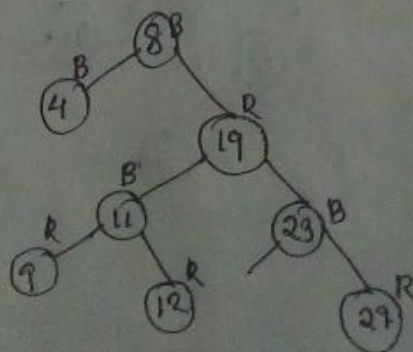
7. Red black tree

Red black tree is a type of binary search tree with some additional properties that are listed below.

The root node of the red black tree is always black coloured. Every path from the root node to leaf node must have the same number of black nodes. No two red nodes can be adjacent, which means that, the red node cannot be a parent of the another red node.

. And a Red black tree is a kind of self balancing tree, where each node has an extra bit. Each node has some colour is red or black. These colours are used to ensure that the tree remains balanced during insertion and deletion.

example of red black tree.

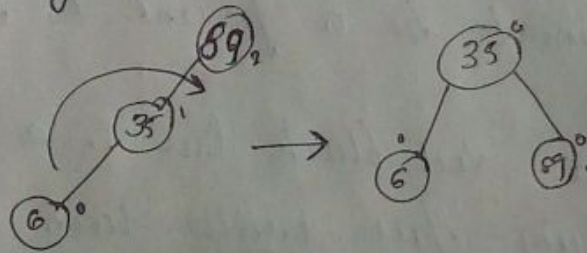


6. Balancing can be performed, ~~at~~ ~~on~~ using the following rotations.

- 1) Right rotation
- 2) Left rotation
- 3) Right left rotation
- 4) Left right rotation

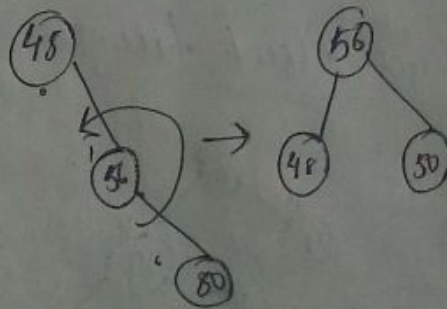
1) Right rotation

In the following binary search tree, there is a single rotation is required.

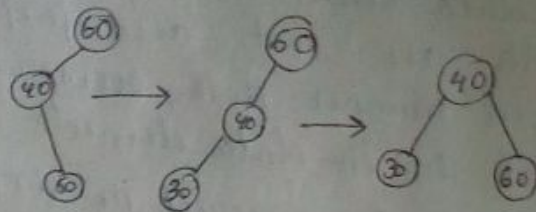


2) Left rotation

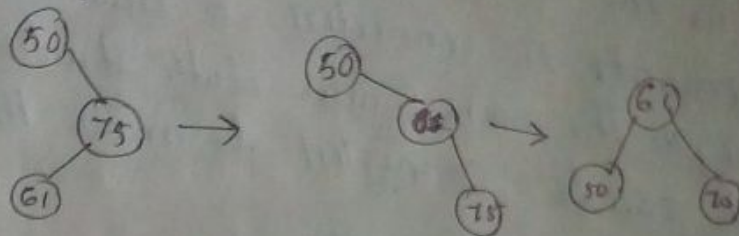
In the following binary search tree, there is single left-rotation is required



3 Right-left rotation



4. Left-right rotation



1) Linear data structure.

A linear data structure is a structure in which the elements are stored sequentially, the elements are connected to the previous and the next element. The elements are stored sequentially, so they can be traversed or accessed in a single run. The implementation process are very easy compared to the non-linear data structure. The data elements are ~~the~~ traversed one after another and can access only one element at a time.

examples:

- * Array
- * Queue
- * Stack.
- * linked list

Non linear data structure

A non-linear data structure is also another type of data structure in which the data elements are not arranged in a contiguous manner. The arrangement is non-sequential, so the data elements cannot be traversed or accessed in single run. In the case of linear data structures, the elements are connected to two elements, whereas, in the non-linear data structure, an element can be connected to more than two elements.

Eg: tree and graph.

3) Binary search tree

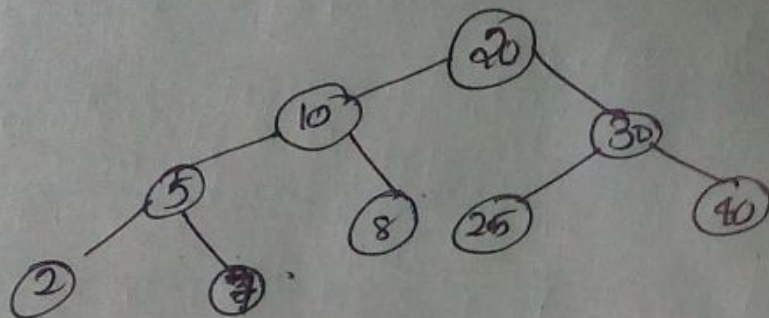
A binary search tree is a type of tree in data structure which has the following properties

→ The value of the left subtree is less than the value of its parent's node's key.

→ The value of the key of the right subtree is greater than or equal to the node.

The binary search tree is a collection of nodes arranged in a way where they maintain BST properties.

example of binary search tree



Basic operations

- Insertion
- deletion
- searching.
- Traversal
 - Inorder
 - preorder
 - post order

2) A stack is a linear data structure can be implemented using one-dimensional array. Stack implemented using array stores only a fixed number of data values. The implementation of it is very simple. Just define a one dimension array of a size and insert and delete the values into that array by using LIFO principle with the help of a specific variable. Initially, the top is set to -1. To add a value from the stack, then delete the top value and decrement the top by value 1.

Algorithm

Step 1: start

Step 2: declare a 1 dimensional array with a fixed size.

Step 3: define a integer variable called top and initialize $top = -1$.

Step 4: And perform the stack operations

1) push

2) pop

Step 5: stop