

2. Implement Lower Bound and Upper Bound

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Lower Bound: Smallest Index such that $\text{arr}[\text{ind}] \geq x$.

$\text{arr}[] = \{3, 5, 8, 15, 19\}$ $n=5$
 0 1 2 3 4

E.g.

if $x = 8$, then lower bound is now = 2 (index)

$$\text{arr}[\text{ind}] \geq x$$

$$\text{arr}[2] \geq 8$$

if $x = 9$, then lower bound is now = 3 (index)

$$\text{arr}[3] \geq 9$$

$$15 \geq 9$$

if $x = 16$, $\text{arr}[4] \geq 16$

$$19 \geq 16$$

$x = 20$, last hypothetically index = 5

now suppose,

arr [] = { 3 5 8 15 19 19 19 } n = 7
0 1 2 3 4 5 6

$$x = 1^q, \quad \text{then} \quad b = 4 \quad \text{arr}[4] \geq 1^q$$

$$19 \geq -19$$

Even then is no element we will have hypothetical index which is last index.

$$\text{ans} = \emptyset, \quad \text{mid} = \frac{0+9}{2} = 4 \quad \text{arr}[4] \geq x$$

$\# z = 1$

now eliminating the right space

high = mid - 1

(you need to find
the smallest)
(probably can be ans
so update it)

$$\begin{array}{ccccccccccccc} & & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ \text{d}_1 & \text{TT} & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 & 13 \end{array}$$

$\text{arr}[\cdot] = \{1, 2, 3, 3, 7, 8, 9, 9, 9, 11\}$ $x=1$
 low mid high

$\text{ans} = 4/1$, $\text{mid} = 1$

$\text{arr}[0] \geq x$

$2 \geq 1$ (can be ans so
update pt)

$\text{arr}[\cdot] = \{1, 2, 3, 3, 7, 8, 9, 9, 9, 11\}$ $x=1$
 low, high,
 mid

$\text{ans} = 1/0$, $\text{mid} = 0$

$\text{arr}[0] \geq 1$

$1 \geq 1$ (might be
ans so update it)

now high will cross low, so you stop and

return the ans, and the lower bound = 0

Lower Bound Implementation:

$f(\text{arr}, \text{target}, n)$

{ $\text{low}=0, \text{high}=n-1;$

```

ans = n;

while (low <= high)

{
    mid = (low + high) / 2;

    if (arr[mid] >= x)
        ans = mid; // may be an answer

        high = mid - 1;

    else
        low = mid + 1;

}

return ans;

```

```

1 int lowerBound(vector<int> arr, int n, int x) {
2     int low = 0, high = n - 1;
3     int ans = n;
4
5     while (low <= high) {
6         int mid = (low + high) / 2;
7
8         // Check if the element at the middle index is greater than or equal to the target
9         if (arr[mid] >= x) {
10             ans = mid; // Update the potential lower bound.
11             high = mid - 1; // Adjust the search range to the left of the middle index.
12         } else {
13             low =
14                 mid + 1; // Adjust the search range to the right of the middle index.
15         }
16     }
17
18     return ans;
19 }

```

Lower Bound C++ STL:

vector \downarrow

$\rightarrow lb = \text{lower_bound}(\text{arr.begin}(), \text{arr.end}(), x) - \text{arr.begin}();$

if it's arr
(arr, arr+n, 2)

now if you asked to find out lb for particular index for e.g from 2 to 7.

[2, 7]

lb = lower_bound (arr+2, arr+7, x)

T.C $\Rightarrow \Theta(\log_2 n)$

Upper Bound: smallest index such that arr [ind] $\geq x$

arr [] = [2 3 6 7 8 8 11 11 11 12]

x = 6

x = 11

x = 12

x = 13

x = 0

up = ind = 3 \Rightarrow 7 \geq 6 ind = 9 ind = 10 ind = 10 ind = 0

if (a[mid] $\geq x$)
this is only change

ans = mid

high = mid - 1

else

low = mid + 1

Upper Bound C++ STL:

Vector ↴

→ $ub = \text{upper_bound}(\text{arr.begin}(), \text{arr.end}(), x) - \text{arr.begin}();$

if it's arr
(arr, arr+n, x)

```
1 int upperBound(vector<int> &arr, int x, int n){      Press Esc to exit full screen
2     int low = 0, high = n - 1;
3     int ans = n;
4
5     while (low <= high) {
6         int mid = (low + high) / 2;
7
8         // Check if the element at the middle index is greater than or equal to the target.
9         if (arr[mid] > x) {
10             ans = mid; // Update the potential lower bound.
11             high = mid - 1; // Adjust the search range to the left of the middle index.
12         } else {
13             low =
14                 mid + 1; // Adjust the search range to the right of the middle index.
15         }
16     }
17
18     return ans;
19 }
```

Q) Search Insert Position

Problem Statement Suggest Edit

You are given a sorted array 'arr' of distinct values and a target value 'm'. You need to search for the index of the target value in the array.

If the value is present in the array, then return its index.
If the value is not present, determine the index where it would be inserted in the array while maintaining the sorted order.

Example:

Input: arr = [1, 2, 4, 7], m = 6

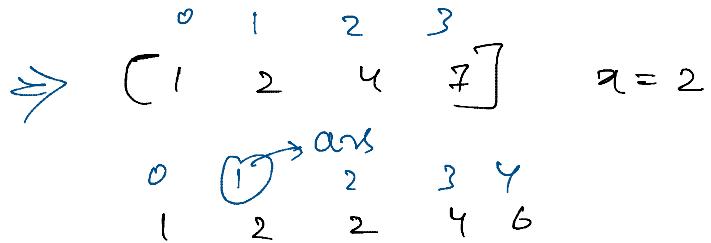
Output: 3

Explanation: If the given array 'arr' is: [1, 2, 4, 7] and m = 6. We insert m = 6 in the array and get 'arr' as: [1, 2, 4, 6, 7]. The position of 6 is 3 (according to 0-based indexing)

arr[] = [0 1 2 3]
 $x = 6$

0 1 2 3
arr[] = [1 2 4 ?]
 $x = 6$

0 1 2 3 4
1 2 4 6 7



It's similar to the lower bound. ($\text{arr}[\text{ind}] \geq x$)

```

1 int searchInsert(vector<int>& arr, int x)
2 {
3     int n = arr.size();
4     int low = 0, high = n - 1;
5     int ans = n;
6
7     while (low <= high) {
8         int mid = (low + high) / 2;
9
10        if (arr[mid] >= x) {
11            ans = mid;
12            high = mid - 1;
13        } else {
14            low =
15            | mid + 1;
16        }
17    }
18
19    return ans;
20 }
```

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Q) Floor and Ceil in Sorted Array

Problem Statement [Suggest Edit](#)

You're given an unsorted array A of N integers and an integer X. Find the floor and ceiling of X in Arr[0..N-1].

Floor of X is the largest element in the array which is smaller than or equal to X.

Ceiling of X is the smallest element in the array greater than or equal to X.

floor and ceil



largest no in array $[<= x]$



smallest no in array $[>= x] \Rightarrow \text{lower_bound}$

largest no in array $\boxed{L=x}$

smallest no in array $\boxed{x \leq L} \Rightarrow \text{lower_bound}$

$$\text{arr}[] = [10 \quad 20 \quad 30 \quad 40 \quad 50] \quad x = 25$$

$$\text{floor} = 20 \quad (20 \leq 25)$$

$$\text{ceil} = 30 \quad (30 > 25)$$

$$\text{arr}[] = [10 \quad 20 \quad 25 \quad 30 \quad 40] \quad x = 25$$

$$\text{floor} = 25 \quad (25 \leq 25)$$

$$\text{ceil} = 25 \quad (25 > 25)$$

Pseudo code for floor:

floor (arr, x)

{

ans = -1;

low = 0, high = n-1;

while (low <= high)

{

mid = (low + high) / 2;

if (arr[mid] <= mid)

ans = arr[mid]

low = mid + 1

else

high = mid - 1;

else

high = mid - 1;

}

How is BS making sure that we get the smallest index?

→ BS eliminates left / right to find the smallest index

→ By default it does that.