

37. Merge Overlapping Intervals

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56. Merge Intervals

Medium 18670 634 Add to List Share

Given an array of `intervals` where `intervals[i] = [starti, endi]`, merge all overlapping intervals, and return an array of the non-overlapping intervals that cover all the intervals in the input.

Example 1:

```
Input: intervals = [[1,3],[2,6],[8,10],[15,18]]  
Output: [[1,6],[8,10],[15,18]]  
Explanation: Since intervals [1,3] and [2,6] overlap, merge them into [1,6].
```

Example 2:

```
Input: intervals = [[1,4],[4,5]]  
Output: [[1,5]]  
Explanation: Intervals [1,4] and [4,5] are considered overlapping.
```

Merge Overlapping Subintervals ?

$\left[(1,3) \ (2,6) \ (8,9) \ (9,11) \ (8,10) \ (1,4) \ (15,18) \ (16,18) \right]$

Subinterval $\Rightarrow (1, 3)$



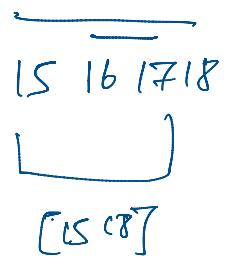
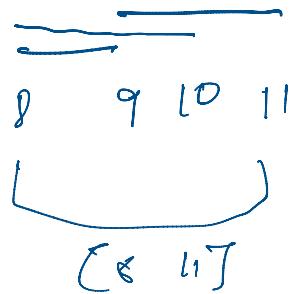
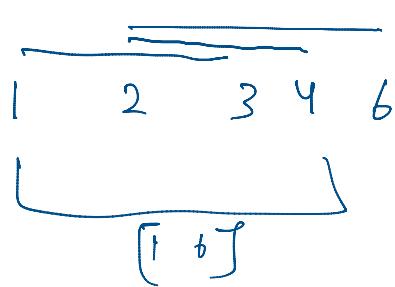
Start with

, and end with 3

start with

1 and end with 3

$$\{1, 6\} \quad \{8, 11\} \quad \{15, 18\}$$



1) Bruk Force:

→ Sort them

$$(1, 3) (2, 4) (2, 6) (8, 9) (8, 10) (9, 11) (15, 18) (16, 17)$$

→ Start iterating over

$$(1, 3) (2, 4) (2, 6) (8, 9) (8, 10) (9, 11) (15, 18) (16, 17)$$

↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑

→ (1, 3) now in (2, 4) if 2 comes b/w (1, 3) then it's overlapping

- so now it expands to $(1, 4)$
- same for $(2, 6) \Rightarrow (1, 6)$
- $(8, 9)$ is not overlapping as $8 \geq 6$, so stop one ans is $(1, 6)$
- again take one interval $(8, 9)$
- $(8, 10)$ is part of $(8, 9) \Rightarrow (8, 10)$
- $(9, 11)$ is part of $(8, 10) \Rightarrow \boxed{(8, 11)}$
↓ one more interval
- $(15, 17)$ another interval start,
 $(16, 18)$ can be part of $(15, 18)$
so $(15, 18)$ is one more interval.

```

1 class Solution {
2 public:
3     vector<vector<int>> merge(vector<vector<int>>& arr) {
4         int n = arr.size();
5         sort(arr.begin(), arr.end());
6
7         vector<vector<int>> ans;
8         for (int i = 0; i < n; i++) {
9             int start = arr[i][0];
10            int end = arr[i][1];
11            if (!ans.empty() && end <= ans.back()[1]) {
12                continue;
13            }
14            for (int j = i + 1; j < n; j++) {
15                if (arr[j][0] <= end) {
16                    end = max(end, arr[j][1]);
17                } else {
18                    break;
19                }
20            }
21            ans.push_back({start, end});
22        }
23        return ans;
24    }
25}

```

$$TC \Rightarrow O(n \log n) + O(2n)$$

$$SC \Rightarrow O(n)$$

2) Optimal Solution:

(1,3) (2,4) (2,6) (8,9) (8,10) (9,11) (15,18) (16,17)

↑ ↑ ↑ ↑

Try to do in single iteration.

(1,3) \Rightarrow (1,4)

[2,4]

2 \leftarrow 3 (so attach)

(2,6]

2 \leftarrow 4

$\cup^{(1,2)} \rightarrow \cup^{(1,2)}$
 $\cup^{(1,2)} \rightarrow \cup^{(1,3)}$
 $\cup^{(1,3)} \rightarrow \cup^{(1,4)}$
 $\Rightarrow (1,6)$

$(1,9) \Rightarrow (8,10) \Rightarrow (8,11) (8,10)$
 $(9,11) \Rightarrow (8,9) (8,10)$
 $8 < 9 (\text{so attach}) \quad 9 < 10$

$(15,18) \Rightarrow (16,17)$
 $\text{as it inside keep it.}$

$i \leftarrow i+1 \Rightarrow O(n \log n)$

```

1 class Solution {
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4         int n = arr.size();
5         sort(arr.begin(), arr.end());
6
7         vector<vector<int>> ans;
8         for (int i = 0; i < n; i++) {
9             if(ans.empty() || arr[i][0] > ans.back()[1]){
10                 ans.push_back(arr[i]);
11             } else {
12                 ans.back()[1] = max(ans.back()[1], arr[i][1]);
13             }
14         }
15         return ans;
16     }
17 };

```