

18. Longest Subarray With Sum K (CodeStudio)

Monday, March 13, 2023 3:36 PM

$$\text{arr}[] = \{1, 2, 3, 1, 1, 1, 4, 2, 3\}$$

$$k = 3$$

Subarray \Rightarrow contiguous part of array. Eg $\Rightarrow \{1, 2, 3\} \checkmark \{4, 2, 3\} \checkmark$

This not subarray $\Rightarrow \{3, 4, 2\} \times$

but this is subsequence.

$$\{1, 2, 3, \boxed{1, 1, 1}, 4, 2, 3\} \quad k = 3$$

3 len = 3 (longest)

1) Brute force

\rightarrow Generate all subarrays

$$\begin{array}{c} j \\ \downarrow \\ \{1, 2, 3, 1, 1, 1, 4, 2, 3\} \\ \uparrow \\ i \end{array} \quad (i - j)$$

for first $i \leftrightarrow$ with whole j

for 2nd $i \rightarrow$ with whole j

do so until i goes out the end.

```

len = 0
for (i = 0 → i + t)
{
    for (j = i → j + t)
        {
            S = 0
            for (k = i → j)
                S + = a[k]
            if (S == k) len = max(len, j - i + 1)
        }
}
print(k)

```

T.C $\Rightarrow O(n^3)$
S.C $\Rightarrow O(1)$

2) Optimization

{ 1, 2, 3, 1, 1, 1, 4, 2, 3 }

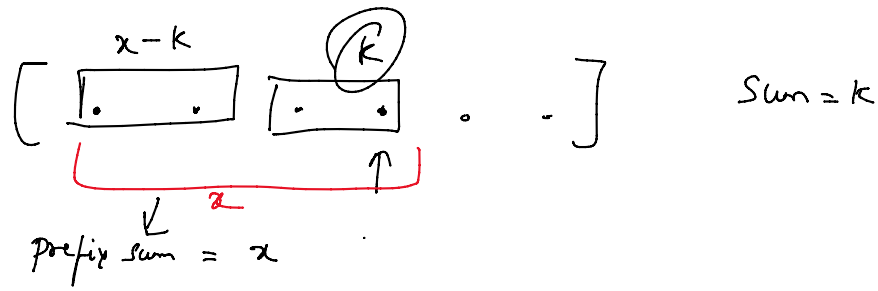
```

len = 0
for (i = 0 → i + t)
{
    S = 0
    for (j = i → j + t)
        {
            S + = a[j];
            if (S == k) len = max(len, j - i + 1)
        }
}
print(k)

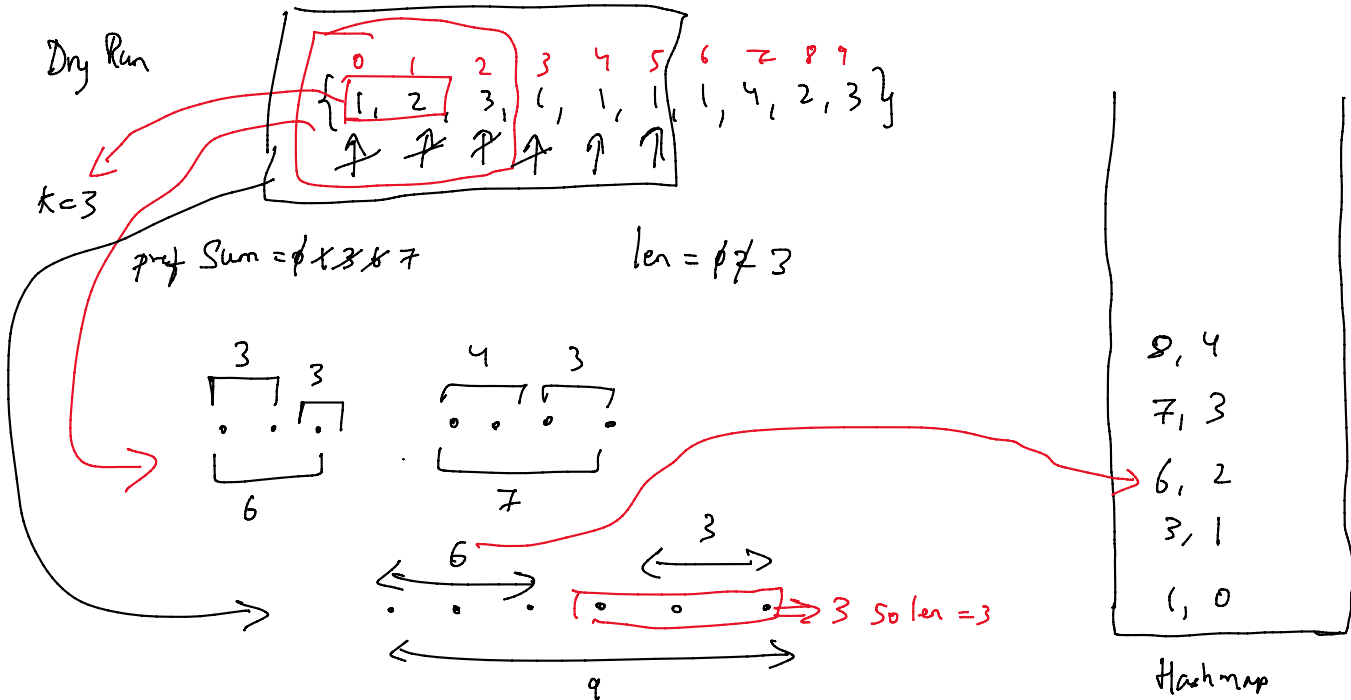
```

T.C $\Rightarrow O(n^2)$

Hashing:



if there exists a subarray with sum k as (i) as the last element



Longest Subarray With Sum K

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Easy 36/40 Avg time to solve 20 mins Success Rate 75% Share 3 upvotes

Problem Statement [Suggest Edit](#)

You are given an array 'A' of size 'N' and an integer 'K'. You need to print the length of the longest subarray of array 'A' whose sum = 'K'.

Example:

Input: 'N' = 7 'K' = 3
'A' = [1, 2, 3, 1, 1, 1, 1]

Output: 3

Explanation: Subarrays whose sum = '3' are:
[1, 2], [3], [1, 1, 1], [1, 1, 1]
Here, the length of the longest subarray is 3, which is our final answer.

```
1 #include <bits/stdc++.h>
2 int longestSubarrayWithSumK(vector<int> a, long long k) {
3     map<long long, int> preSumMap;
4     long long sum = 0;
5     int maxLen = 0;
6     for(int i = 0; i < a.size(); i++){
7         sum += a[i];
8         if(sum == k){
9             maxLen = max(maxLen, i+1);
10        }
11        long long remaining = sum - k;
12        if(preSumMap.find(remaining) != preSumMap.end()){
13            int len = i - preSumMap[remaining];
14            maxLen = max(maxLen, len);
15        }
16        preSumMap[sum] = i;
17    }
18    return maxLen;
19 }
```

This code will not work for some test case

eg.

./. 1 1.

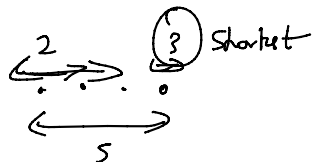
this case will not work for some cases

eg.

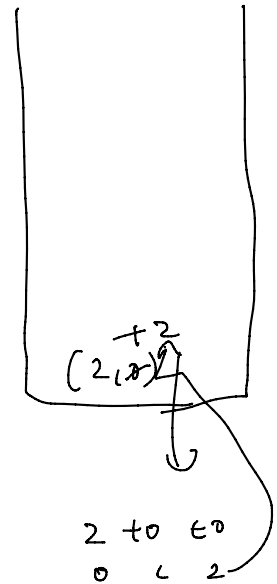
$$\text{arr}[] = \{ 2, 0, 0, 3 \} \quad k=3$$

$$\text{presum} = 0 \leq 5$$

$$\text{len} = 0$$



we need longest subarray not shortest



$$T.C \Rightarrow O(N \log N) \quad (\text{ordered map})$$

$$S.C \Rightarrow O(N)$$

Optimal

$$\text{arr}[] = \{ 1, 2, 3, 1, 1, 1, 1, 3, 3 \}$$

Two pointer

$$k=6$$

$$\text{len} = 8 - 4 = 4$$

sum = 13 [6], 7 So we will reduce it, so we will remove 1

$$6 \neq 6$$

$$5 \neq 6$$

$$4 \neq 6$$

$$3 \neq 6$$

$$2 \neq 6$$

$$1 \neq 6$$

```
1 int longestSubarrayWithSumK(vector<int> a, long long k) {
2     int left = 0, right = 0;
3     long long sum = a[0];
4     int maxLen = 0;
5     int n = a.size();
6
7     while(right < n){
8         while(left <= right && sum > k){
9             sum -= a[left];
10            left++;
11        }
12        if(sum == k){
13            maxLen = max(maxLen, right - left + 1);
14        }
15        right++;
16        if(right < n){
17            sum += a[right];
18        }
19    }
20    return maxLen;
21 }
```