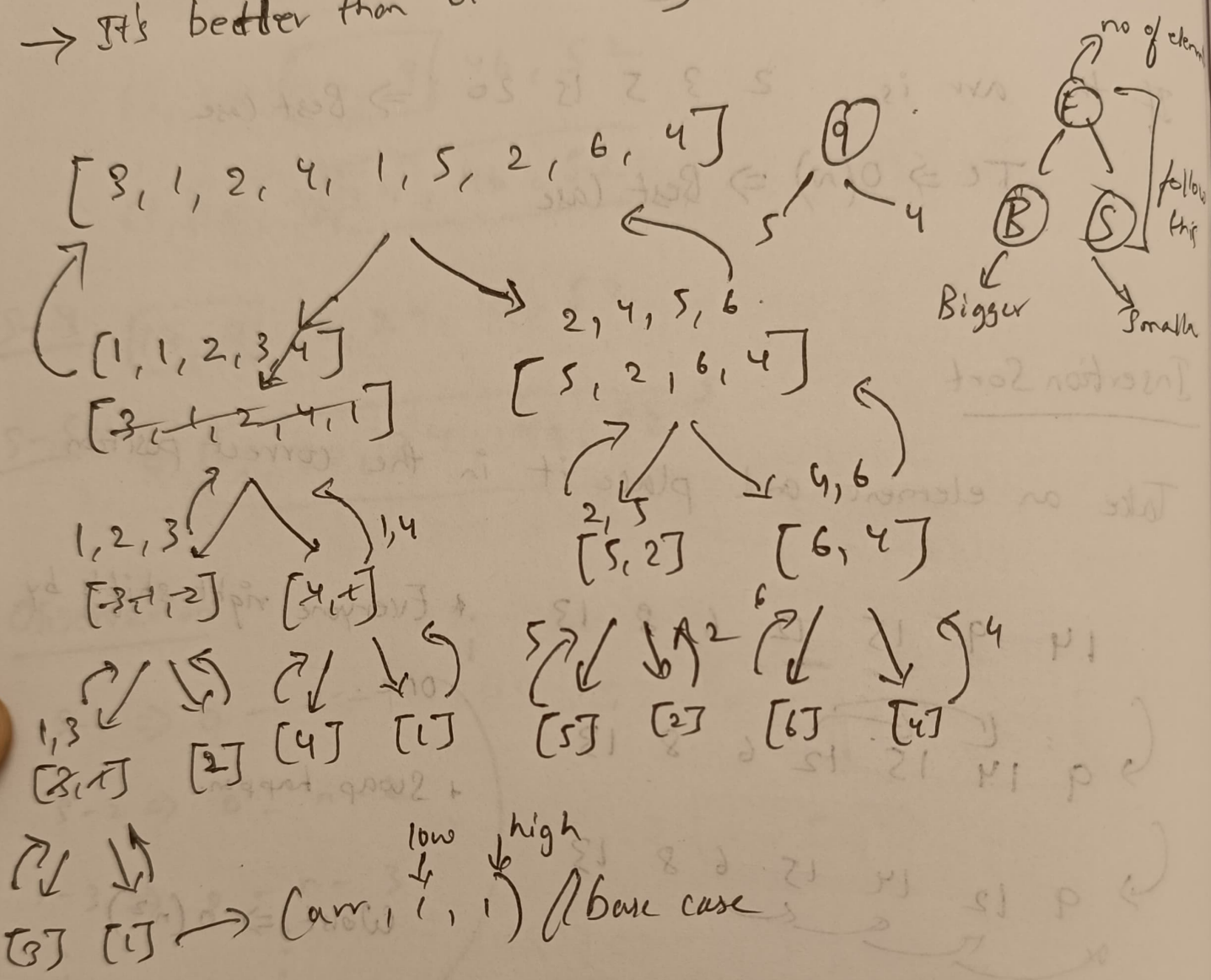


# Merge Sort → Divide & Merge.

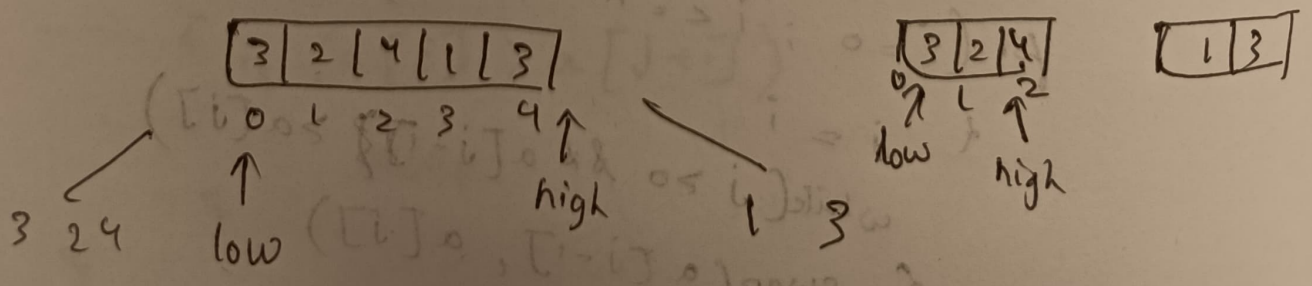
→ It's better than other sorting algorithm.



low high  
 (arr, low, high) // base case

Sorted  $\Rightarrow [1, 1, 2, 2, 3, 4, 4, 5, 6]$

PseudoCode (Play around with index, instead of breaking it)



$\rightarrow 0 \rightarrow 4$   
 $\text{merge\_sort}(\text{arr}, \text{low}, \text{high})$   
 { // base case if ( $\text{low} \geq \text{high}$ ) return;  
 $\text{mid} = (\text{low} + \text{high}) / 2$   $\rightarrow$  index  
 $\text{merge\_sort}(\text{arr}, \text{low}, \text{mid})$   
 $\text{merge\_sort}(\text{arr}, \text{mid} + 1, \text{high})$   
 $\text{merge}(\text{arr}, \text{low}, \text{mid}, \text{high});$

Dry Run

0, 4

$\text{ms}(\text{arr}, l, h)$

if ( $l \geq h$ ) return;

$m = (l + h) / 2 \rightarrow 2$

$\text{ms}(\text{arr}, \text{low}, \text{mid}) \rightarrow$

$\text{ms}(\text{arr}, m + 1, h)$

$\text{m}(\text{arr}, l, m, h)$

$\text{ms}(0, 2)$

if () x

$m = 1$

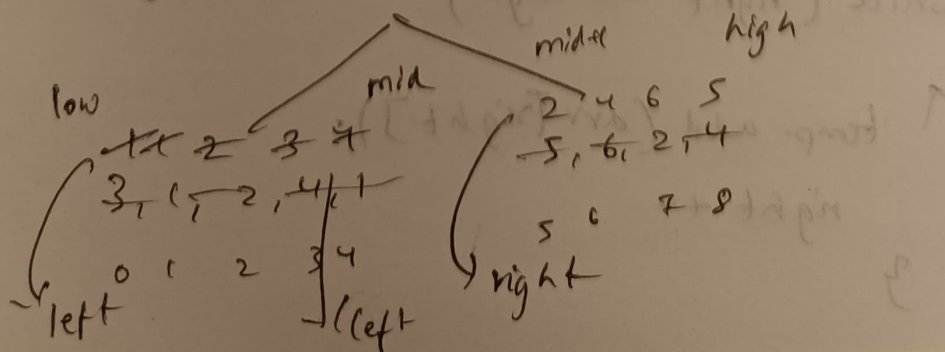
ms

}

Approach of merge

3, 1, 2, 4, 1, 5, 6, 2, 4

0 1 2 3 4 5 6 7 8



merge(arr, low, mid, high)

{ temp  $\rightarrow$  [ ]

left = low

right = mid + 1;

while (left  $\leq$  mid & right  $\leq$  high)

{ if (arr[left]  $\leq$  arr[right])

temp.add(arr[left])

left ++;

else

temp.add(arr[right])

right ++;

}

while (left  $\leq$  mid)

{ temp.add(arr[left]);

left ++;

}

while (right  $\leq$  high)

{ temp.add(arr[right]);

right ++

}

for (i = low  $\rightarrow$  high)

{ arr[i] = temp[i - low];

}

T.C

[n]

$\frac{n}{2}$

$\frac{n}{4}$

$\frac{n}{8}$

$\frac{n}{16}$

1

T.C  $\Rightarrow O(\log_2 n)$

S.C  $\Rightarrow O(n)$

if it  $\frac{n}{2}$  the log.