

36. Number of Subarrays with xor K

Tuesday, April 25, 2023 4:30 PM

Problem Statement

[Suggest Edit](#)

Given an array 'A' consisting of 'N' integers and an integer 'B', find the number of subarrays of array 'A' whose bitwise XOR(\oplus) of all elements is equal to 'B'.

Example:

Input: 'N' = 4 'B' = 2
'A' = [1, 2, 3, 2]

Output: 3

Explanation: Subarrays have bitwise xor equal to '2' are: [1, 2, 3, 2], [2], [2].

A subarray of an array is obtained by removing some(zero or more) elements from the front and back of the array.

contiguous part of array
Count Subarrays with XOR as k.
arr [] = { 4 2 2 6 4 } k = 6

$$\{4, 2\} = 4 \wedge 2 = 6.$$

$$\{6\} = 6$$

$$\{2, 2, 6\} \Rightarrow \{2 \wedge 2\} = 0 \wedge 6 = 6$$

$$\{4, 2, 2, 6, 4\} = \{6\} = 6$$

ans = 6

$$\begin{array}{r} 8 \ 4 \ 2 \ 1 \\ \hline 0 \ 1 \ 0 \ 0 = 4 \\ 0 \ 0 \ 1 \ 0 = 2 \\ \hline 0 \ 1 \ 1 \ 0 = 6 \end{array}$$

i) Brute Force:

arr [] = [4 2 2 6 4] k = 6

* Generate Subarrays

for (i=0; i < n; i++)

{

 for (int j=i; j < n; j++)

```
for (int j = i; j < n; j++)
```

```
{ XOR = 0
```

```
pr(k = i → j)
```

```
XOR = XOR ^ arr[j];
```

```
if (XOR == k) cnt++;
```

```
}
```

```
}
```

```
return cnt;
```

$T.C \Rightarrow O(n^2)$ $S.C \Rightarrow O(1)$

2) Better Solution :

→ We don't need k -loop.

```
for (i = 0; i < n; i++)
```

```
{ XOR = 0;
```

```
for (j = i; j < n; j++)
```

```
{ XOR = XOR ^ arr[j];
```

```
if (XOR == k) cnt++;
```

```
}
```

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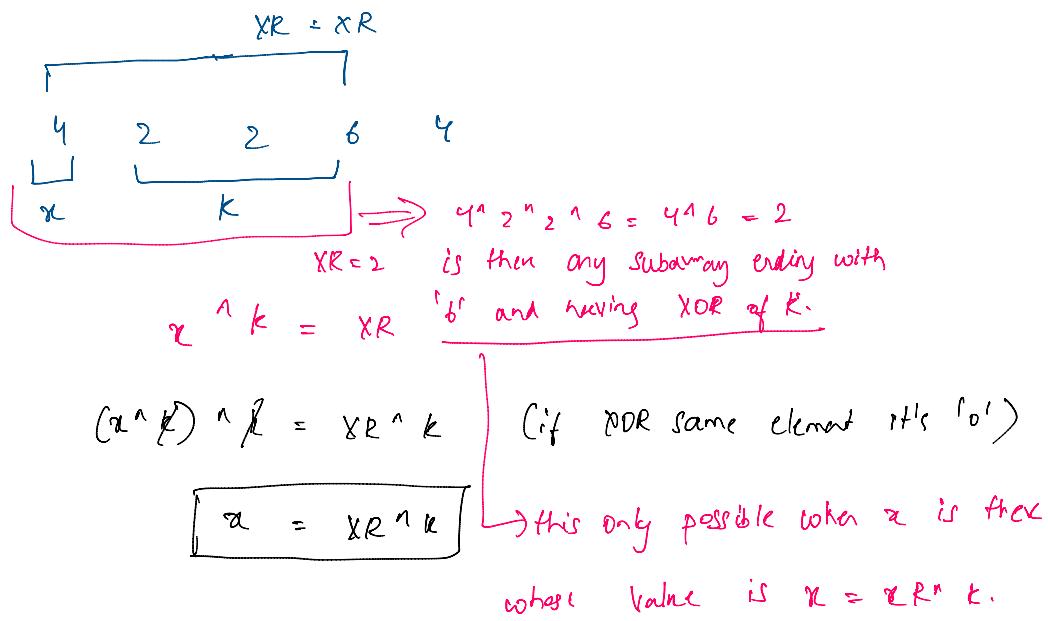
return cut;

$$TC \Rightarrow O(n^2) \quad SC \Rightarrow O(1)$$

3) Optimal Solution

$$\text{avr} \{ \cdot \} = [4 \quad 2 \quad 2 \quad 6 \quad 4] \quad k = 6$$

- Every subarray has starting point and ending point
we will focus on ending point.

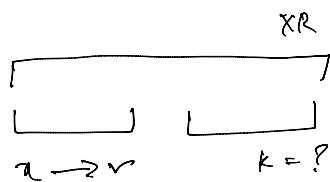


So we have someone from start which is giving '4'

Yes we have so now we can say that there

sniff - some element which will give me the 'K'.

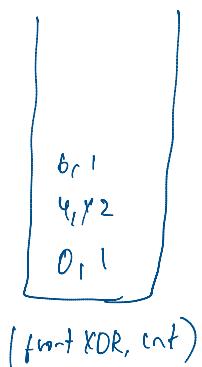
exists some element which will give me the 'k'.
 we will try to move, and keep a track of
 pre XOR.



To keep a track we need hashmap dict

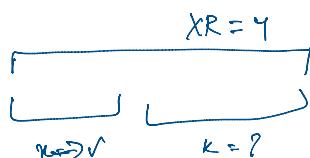
$$\text{arr}[] = [4 \ 2 \ 2 \ 6 \ 4] \quad k = 6$$

$\uparrow \ \uparrow \ \uparrow \ \uparrow$



$$XR = \emptyset \oplus 4, 2(4 \oplus 2) = 6, 2(6 \oplus 2) = 4, -6(6 \oplus 4) = 2, 4(2 \oplus 4) = 6.$$

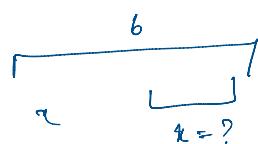
$\text{cnt} = \emptyset \cup \{3\}$,



$$x = XR \oplus k$$

$$= 4 \wedge 6$$

$$= 2 \text{ (is there '2' previously NO)}$$

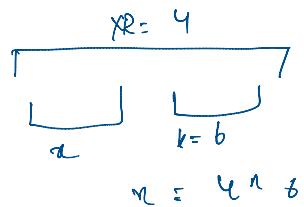
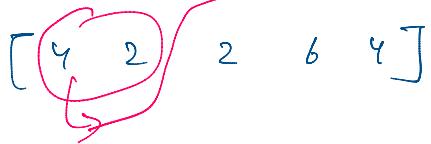


$$x = 6 \wedge 6$$

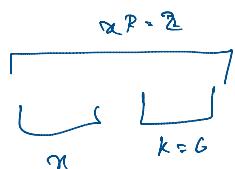
$$x = 0 \text{ (yes value stored)}$$

so entire thing is subarray

so entire thing is subarray

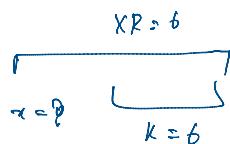
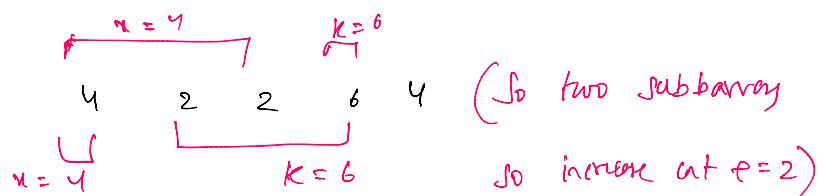


$$n = 2 \quad (\text{No front as } 2 \text{ is last})$$



$$n = 2 \wedge 6$$

$= 4$ (4 as we have two element)



$$n = 6 \wedge 6 = 0 \quad (\text{we have '0'})$$

So entire array. $\text{cnt} = 1$

$\text{cnt} = 4$

$$\text{T.C} \Rightarrow O(n \log n) \quad \text{P.C} \Rightarrow O(n)$$

```
1 #include<bits/stdc++.h>
2 int subarraysWithSumK(vector < int > a, int k) {
3     int xr = 0;
4     map<int, int> mpp;
5     mpp[xr]++;
6     int cnt = 0;
7     for(int i = 0; i < a.size(); i++){
8         xr = xr ^ a[i];
9         // looking for k
10        int x = xr ^ k;
11        cnt += mpp[x];
12        mpp[xr]++;
13    }
14    return cnt;
15 }
```