

## 34. 3 Sum | Brute - Better - Optimal with Codes

15. 3Sum

Medium 25136 2267 Add to List Share

Given an integer array `nums`, return all the triplets  $[nums[i], nums[j], nums[k]]$  such that  $i \neq j$ ,  $i \neq k$ , and  $j \neq k$ , and  $nums[i] + nums[j] + nums[k] = 0$ .

Notice that the solution set must not contain duplicate triplets.

**Example 1:**

**Input:** `nums = [-1, 0, 1, 2, -1, -4]`

**Output:**  $\{[-1, 0, 1], [-1, 0, 1]\}$

**Explanation:**

$nums[0] + nums[1] + nums[2] = (-1) + 0 + 1 = 0$ .

$nums[1] + nums[2] + nums[4] = 0 + 1 + (-1) = 0$ .

$nums[0] + nums[3] + nums[4] = (-1) + 2 + (-1) = 0$ .

The distinct triplets are  $[-1, 0, 1]$  and  $[-1, 0, 1]$ .

Notice that the order of the output and the order of the triplets does not matter.

$$arr[i] + arr[j] + arr[k] = 0$$

$$(i \neq j \neq k)$$

You cannot take same element at once.

$$arr[J] = \{-1, 0, 1, 2, -1, -4\}$$

$\left. \begin{matrix} \{-1, 2, -1\} \\ \{0, 1, -1\} \end{matrix} \right\}$  You have to return unique triplet

i) Brute Force:

→ Try out all triplet

$$\begin{array}{c} arr[J] = \{-1, 0, 1, 2, -1, -4\} \\ \uparrow \quad \uparrow \quad \uparrow \quad \uparrow \quad \uparrow \quad \uparrow \quad \text{so on...} \\ \text{pick } \quad \text{2} \quad \text{2} \quad \text{2} \quad \text{2} \quad \text{2} \quad \text{2} \end{array}$$

$$\begin{bmatrix} -, -, - \\ 1 \quad 2 \quad 3 \end{bmatrix}$$

for ( $i = 0$ ;  $i < n$ ;  $i + 1$ )

{

    for ( $j = i + 1 \rightarrow n$ )

        for ( $k = j + 1 \rightarrow n$ )

$\dots$

if ( $\text{arr}[i] + \text{arr}[j] + \text{arr}[k] = 0$ )

L  
 $\{-, -, -\}$  ((triplet)

}

}

// To make sure not to store duplicate

Sort this duplicate.

as we need unique use set duplicate.

```
1 #include<bits/stdc++.h>
2 vector<vector<int>> triplet(int n, vector<int> &num)
3 {
4     set<vector<int>> st; // store all the triplets in form of set ds in vector
5     for(int i = 0; i < n; i++){
6         for(int j = i+1; j < n; j++){
7             for(int k = j+1; k < n; k++){
8                 if(num[i] + num[j] + num[k] == 0){
9                     vector<int> temp = {num[i], num[j], num[k]};
10                    sort(temp.begin(), temp.end());
11                    st.insert(temp);
12                }
13            }
14        }
15    }
16    vector<vector<int>> ans(st.begin(), st.end());
17    return ans;
18 }
19 }
```

T.C  $\Rightarrow O(n^3 \times \log(\text{no of unique}))$

S.C  $\Rightarrow O(\text{no of triplets})$

2) Better Solution:

$$\text{arr}[J] = \{-1, 0, 1, 2, -1, -4\}$$

To do in  $O(n^2)$ , you need to get rid of 3rd loop,

$$\text{arr}[i] + \text{arr}[j] + \text{arr}[k] = 0$$

$$\text{arr}[k] = -(\text{arr}[i] + \text{arr}[j])$$

$$\begin{matrix} i & & j \\ \text{arr}[J] = \{ -1, 0, 1, 2, -1, -4 \} \end{matrix}$$

$$\text{arr}[k] = -(-1 + -1)$$

$$= -(-2)$$

$$\text{arr}[k] = 2 \quad (\text{is then } +2 \text{ in the array})$$

In order to look for  $\text{arr}[k]$  without iterating is hashing.

Dry Run for Hashing,

- You have to keep in mind you cannot pick one element twice.

E.g.,

$$\{-1, 0, 1, 2, -1, -4\}$$



$$\text{arr}[k] = - (2 - 4)$$

$$\text{arr}[k] = - (2 - 4)$$

$$= - (-2)$$

$$= 2$$

so be mindful that if you are looking for 2 in the array then we have to be sure that it's not the current '2', '-4' that you have taken for  $\text{arr}[i]$  and  $\text{arr}[j]$  something different from those two.

You can do it by

$$[-1, [0, 1, 2], -1, -4]$$

Keep this element in hashmap.

$$\text{arr}[] = [-1, 0, 1, 2, -1, -4]$$

i    j

$$\text{arr}[k] = - (\text{arr}[i] + \text{arr}[j])$$

$$= - (-1 + 0)$$

1 2 3 4 5 6 7 8 9 10

0

$i = -1$   
 $= 1$  (is it in set NO)   
 then move  $j$  and keep the prev  $j$  in  
 the hash map.

$= -(-1 + 1)$

$= -0$

$= 0$  (it's true in hashset)

So one of the ans after sorting  $(-1, 0, 1)$

Keep doing that for every element.

T.C  $\Rightarrow O(n^2 \log n)$  S.C  $\Rightarrow O(n) + O(\text{no of unique triple}) \times 2$

```
i C++ Autocomplete
1 class Solution {
2 public:
3     vector<vector<int>> threeSum(vector<int>& num) {
4         int n = num.size();
5         set<vector<int>> st;
6         for(int i = 0; i < n; i++) {
7             set<int> hashset;
8             for(int j = i+1; j < n; j++) {
9                 int third = -(num[i] + num[j]);
10                if(hashset.find(third) != hashset.end()) {
11                    vector<int> temp = {num[i], num[j], third};
12                    sort(temp.begin(), temp.end());
13                    st.insert(temp);
14                }
15                hashset.insert(num[j]);
16            }
17        }
18        vector<vector<int>> ans(st.begin(), st.end());
19        return ans;
20    }
21 }
```

```
1 #include<bits/stdc++.h>
2 vector<vector<int>> triplet(int n, vector<int> &num)
3 {
4     set<vector<int>> st;
5     for(int i = 0; i < n; i++) {
6         set<int> hashset;
7         for(int j = i+1; j < n; j++) {
8             int third = -(num[i] + num[j]);
9             if(hashset.find(third) != hashset.end()) {
10                 vector<int> temp = {num[i], num[j], third};
11                 sort(temp.begin(), temp.end());
12                 st.insert(temp);
13             }
14             hashset.insert(num[j]);
15         }
16     }
17     vector<vector<int>> ans(st.begin(), st.end());
18     return ans;
19 }
20 }
```

Optimal Approach:

$\text{arr}[] = [-1, 2, 2, -2, -1, 0, -1, 0, 0, -2, 2, -2, 2]$

Now you need to get rid of set d.s.

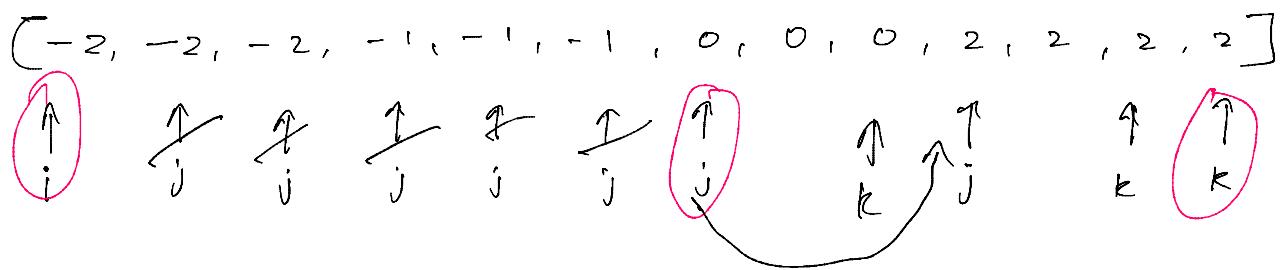
Now you need to get rid of set d.s.

sort it

$$[-2, -2, -2, -1, -1, -1, 0, 0, 0, 2, 2, 2, 2]$$

Algorithm used here is two pointer approach:

$$\text{arr}[i] + \text{arr}[j] + \text{arr}[k]$$



$i \rightarrow$  will be always constant (at first)

$j \rightarrow (i+1)$

$k \rightarrow$  at last

add up all three element

$$-2 -2 + 2 = \boxed{-2 < 0}$$

We need  $= 0$ , but here  $< 0$ , so you need to increase

so we will do  $j++$ ,

$$-2 - 1 + 2 = \boxed{-1 < 0}$$

again same

$$-2 + 0 + 2 = \boxed{0 = \leq 0}$$

so when it's  $\leq 0$  that one of the triplet.

so store it  $\{-2, 0, 2\}$

once  $\boxed{\text{it's } \leq 0 \text{ move } j++ \text{ and } k--}$

$\hookrightarrow$  until the value is not '0', because

the triplet will be same.

once  $k$  cross  $j$ , you stop as it's no more stored.

Now again you will move 'i', but the element which is not ' $-2$ '.

$\left[ -2, -2, -2, -1, -1, -1, 0, 0, 0, 2, 2, 2, 2 \right]$

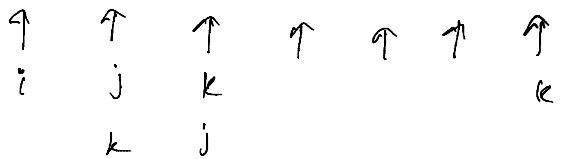
$\begin{matrix} \uparrow & \uparrow & & \uparrow & \uparrow & \uparrow \\ i & j & j & k & j & k \end{matrix}$

$$-1 + (-1) + 2 = = 0$$

so store it  $\{-1, -1, 2\}$

$$-1 + 0 + 0 = -1 < 0$$

$[-2, -2, -2, -1, -1, -1, 0, 0, 0, 2, 2, 2, 2]$



$$0 + 0 + 2 = 2 \geq 0$$

So  $k--$

You'll get another triplet  $\{0, 0, 0\}$

```

1 #include <bits/stdc++.h>
2 vector<vector<int>> triplet(int n, vector<int> &num) {
3     vector<vector<int>> ans;
4     sort(num.begin(), num.end());
5     for (int i = 0; i < n; i++) {
6         if (i > 0 && num[i] == num[i - 1])
7             continue;
8         int j = i + 1;
9         int k = n - 1;
10        while (j < k) {
11            int sum = num[i] + num[j] + num[k];
12            if (sum < 0) {
13                j++;
14            } else if (sum > 0) {
15                k--;
16            } else {
17                // one of the triplet
18                vector<int> temp = {num[i], num[j], num[k]};
19                ans.push_back(temp);
20                j++;
21                k--;
22                while (j < k && num[j] == num[j - 1])
23                    j++;
24                while (j < k && num[k] == num[k + 1])
25                    k--;
26            }
27        }
28    }
29    return ans;
30 }
31 
```

T.C  $\Rightarrow (n \log n) + O(n \times n)$

S.C  $\Rightarrow O(\text{no of unique})$