

19. Two Sum

1. Two Sum

Hint

Easy



44.2K

1.4K



Companies

Given an array of integers `nums` and an integer `target`, return indices of the two numbers such that they add up to `target`.

You may assume that each input would have **exactly one solution**, and you may not use the same element twice.

You can return the answer in any order.

Example 1:

Input: `nums = [2,7,11,15]`, `target = 9`

Output: `[0,1]`

Explanation: Because `nums[0] + nums[1] == 9`, we return `[0,1]`.

Old Video

1) Brute force

0	1	2	3	4
2	6	5	8	11
i	j			

T.C $\Rightarrow O(n^2)$

S.C $\Rightarrow O(1)$

target = 14

check if this exist in the array or not.
 $14 - 2 = 12$

$14 - 6 = 8$ → this exist in array

2) Using hashtable

0	1	2	3	4
2	6	5	8	11
i	i	i	i	

$(1, 3) \Rightarrow$ o/p
 → from the pointer i .
 → from the hashtable

target = 14

Check if it's there in hashtable or not
 (if not there then push it in hashtable the array)

$14 - 2 = 12$

$14 - 6 = 8$

$14 - 5 = 9$

$14 - 8 = 6$ → this exist in index '1'.

(5, 2)
(6, 1)
(2, 0)

hashtable.

T.C $\Rightarrow O(n)$

S. C $\Rightarrow O(N)$ (Storing value in hashtable)

So in C++, we use unordered map

```
1 class Solution {
2 public:
3     vector<int> twoSum(vector<int>& nums, int target) {
4         vector<int> ans;
5         unordered_map<int, int> mpp; // hashtable
6         for(int i = 0; i < nums.size(); i++){
7             if(mpp.find(target - nums[i]) != mpp.end()){
8                 ans.push_back(mpp[target - nums[i]]);
9                 ans.push_back(i);
10                return ans;
11            }
12            mpp[nums[i]] = i;
13        }
14        return ans;
15    }
16 };
```

New video

arr[] = { 2, 6, 5, 8, 11 }

target = 14

1st part \Rightarrow Yes/No

2nd part \Rightarrow return the index of two element.

Brute force

for (i = 0 \rightarrow n)

{ for (j = i + 1 \rightarrow n)

{

if (i == j) continue; // Because we cannot pick up same elements

if (arr[i] + arr[j] == target)

{

print(Yes) // 1st part

print{i, j} // 2nd part

}

}

}
}

T.C $\Rightarrow O(n^2)$ S.C $\Rightarrow O(1)$

2nd Approach

Hashing

arr[] = { 2, 6, 5, 8, 11 }

target = 14

Check if it's there in hash table or not
(if not there then push it in hash table the array)

[3, 1]
↳ from hash map.
↳ from i pointer

14 - 2 = 12

14 - 6 = 8

14 - 5 = 9

14 - 8 = 6

↳ there in hash map

(5, 2)
(6, 1)
(2, 0)

Hash map
(ele, index)
↓ ↓
key Value

T.C $\Rightarrow O(n \log n)$ S.C $\Rightarrow O(n)$

```
1 #include <bits/stdc++.h>
2 string read(int n, vector<int> book, int target)
3 {
4     map<int, int> mpp;
5     for(int i = 0; i < n; i++){
6         int a = book[i];
7         int more = target - a;
8         if(mpp.find(more) != mpp.end()){
9             return "YES"; // if index {mpp[more], i}
10        }
11        mpp[a] = i;
12    }
13    return "NO";
14 }
```

3rd Approach

Without using map.

Using two pointer approach

arr[] = { 2, 6, 5, 8, 11 }

target = 14

∴ LL

$arr[7] = \{ 2, 6, 5, 8, 11 \}$ target = 14

Sort
 $\{ 2, 5, \boxed{6, 8}, 11 \}$
 ↓ ↓ ↓ ↓
 left left left right

$2 + 11 = 13 < 14$ (So move left)

$5 + 11 = 16 \geq 14$ (move right)

$5 + 8 = 13 < 14$ (move left)

$6 + 8 = 14$ ✓

It solves part-1, but for part-2 (index) you need to put $arr[7]$ in another d.s. $\{(2,0) (6,1) (5,2) (8,3) (11,4)\}$ along with index and sort it. for part-2 this is not best.

```

1 string read(int n, vector<int> book, int target)
2 {
3     int left = 0, right = n-1;
4     sort(book.begin(), book.end());
5     while(left < right) {
6         int sum = book[left] + book[right];
7         if(sum == target) {
8             return "YES";
9         }
10        else if(sum < target) left++;
11        else right--;
12    }
13    return "NO";
14 }
15

```

T.C $\Rightarrow O(n) + O(n \log n)$

S.C $\Rightarrow O(1)$