

232. Implement Queue using Stacks

Implement a first in first out (FIFO) queue using only two stacks. The implemented queue should support all the functions of a normal queue (`push`, `peek`, `pop`, and `empty`).

Implement the `MyQueue` class:

- `void push(int x)` Pushes element `x` to the back of the queue.
- `int pop()` Removes the element from the front of the queue and returns it.
- `int peek()` Returns the element at the front of the queue.
- `boolean empty()` Returns `true` if the queue is empty, `false` otherwise.

Notes:

- You must use **only** standard operations of a stack, which means only `push to top`, `peek/pop from top`, `size`, and `is empty` operations are valid.
- Depending on your language, the stack may not be supported natively. You may simulate a stack using a list or deque (double-ended queue) as long as you use only a stack's standard operations.

Example 1:

Input

```
["MyQueue", "push", "push", "peek", "pop", "empty"]  
[[], [1], [2], [], [], []]
```

Output

```
[null, null, null, 1, 1, false]
```

Explanation

```
MyQueue myQueue = new MyQueue();  
myQueue.push(1); // queue is: [1]  
myQueue.push(2); // queue is: [1, 2] (leftmost is front of the queue)  
myQueue.peek(); // return 1  
myQueue.pop(); // return 1, queue is [2]  
myQueue.empty(); // return false
```

Queue \rightarrow FIFO

Stack \rightarrow LIFO

Queue

$p(3)$

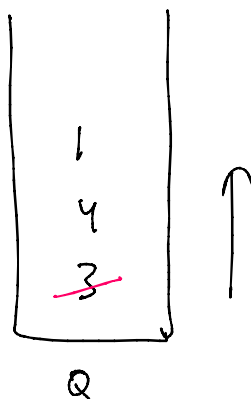
$p(4)$

$p(1)$

$top() \Rightarrow 3$

$pop()$ // remove 3

$top() \Rightarrow 4$



Stack

1

$p(3)$

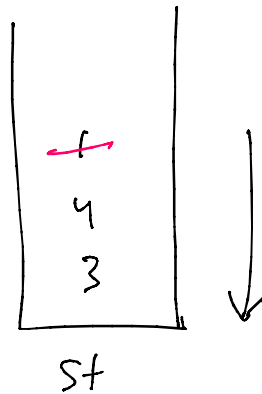
$p(4)$

$p(1)$

$top() \Rightarrow 1$

$pop() // \text{remove } 1$

$top() \Rightarrow 4$



Implement Queue using Stack

1st Approach

$push(4)$

$push(3)$

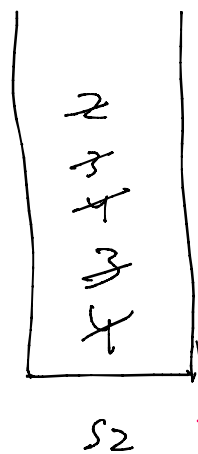
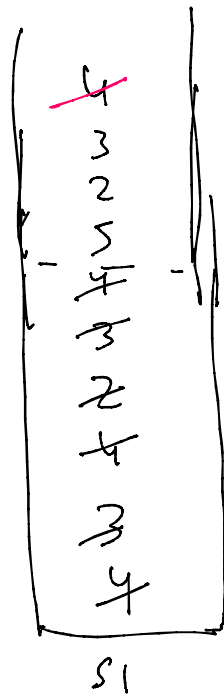
$push(2)$

$push(5)$

$top() \rightarrow 4$

$pop() // 4$

$top() \rightarrow 3$



push(x)

(.) $s1 \rightarrow s2$

(.) $x \rightarrow s1$

(.) $s2 \rightarrow s1$

pop()

$s1 = pop()$

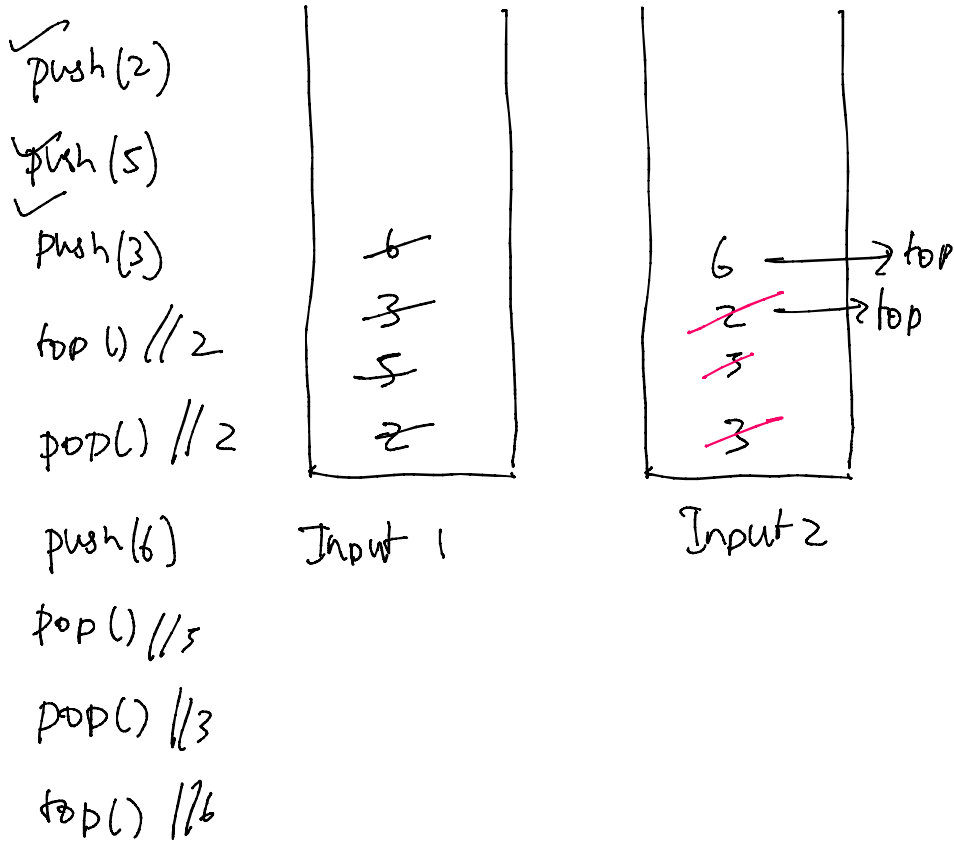
\rightarrow follow this

T.C $\Rightarrow O(N)$

S.C $\Rightarrow O(2N)$

2nd Approach

2nd Approach

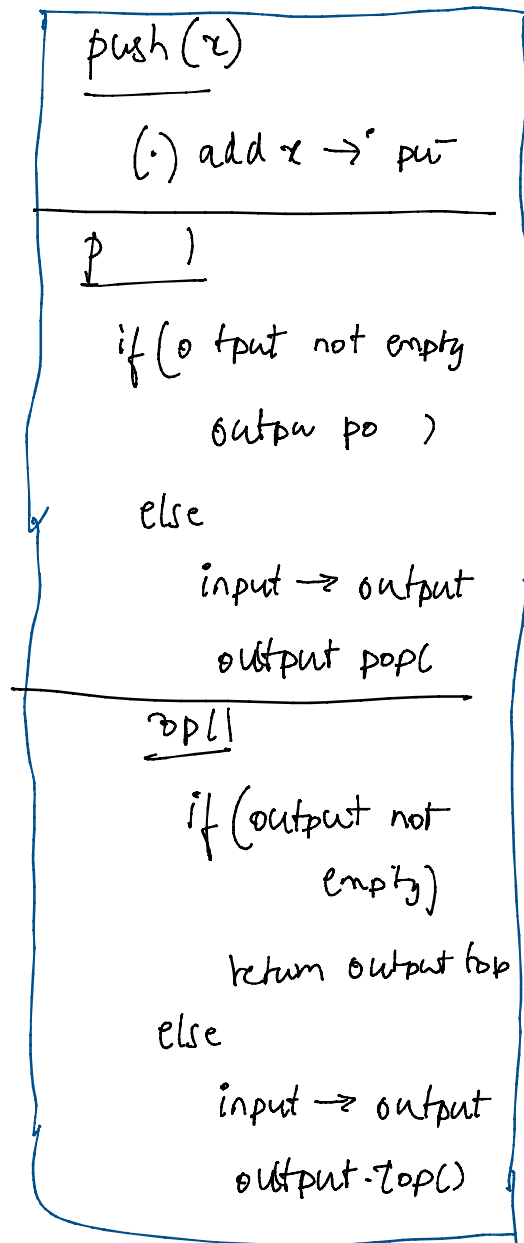


$T.C \Rightarrow \text{push} \Rightarrow O(1)$

$\text{pop/top} \Rightarrow O(1)$ Sometimes $O(N)$

So $O(1)$ amortised.

S.C $\Rightarrow O(2N)$



```

class MyQueue {
    stack<int> in;
    stack<int> out;
public:
    /** Initialize your data structure here. */
    MyQueue() {

        in = stack<int>();
        out = stack<int>();

    }

    /** Push element x to the back of queue. */
    void push(int x) {

        in.push(x);

    }

    /** Removes the element from in front of queue and returns that
    element. */
    int pop() {

        if(out.empty()) {

            //taxing
            while(!in.empty()) {

                //loading the whole stack in out
                out.push(in.top());
                in.pop();

            }

        }

        int val = out.top();
        out.pop();
        return val;

    }

    /** Get the front element. */
    int peek() {

        if(out.empty()) {

            //taxing
            while(!in.empty()) {

                //loading the whole stack in out
                out.push(in.top());
                in.pop();

            }

        }

        return out.top();

    }

    /** Returns whether the queue is empty. */
    bool empty() {

        return out.empty() && in.empty();

    }

};

```