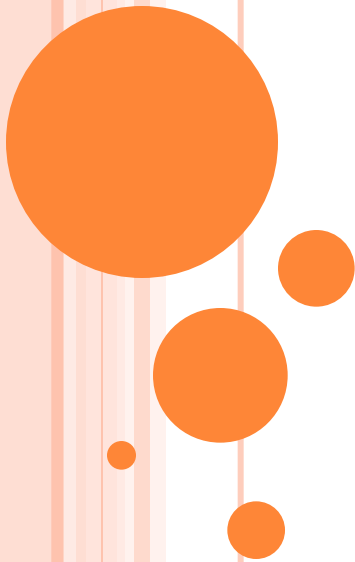


# **CSE 4102**

## **LEXICAL ANALYSIS**

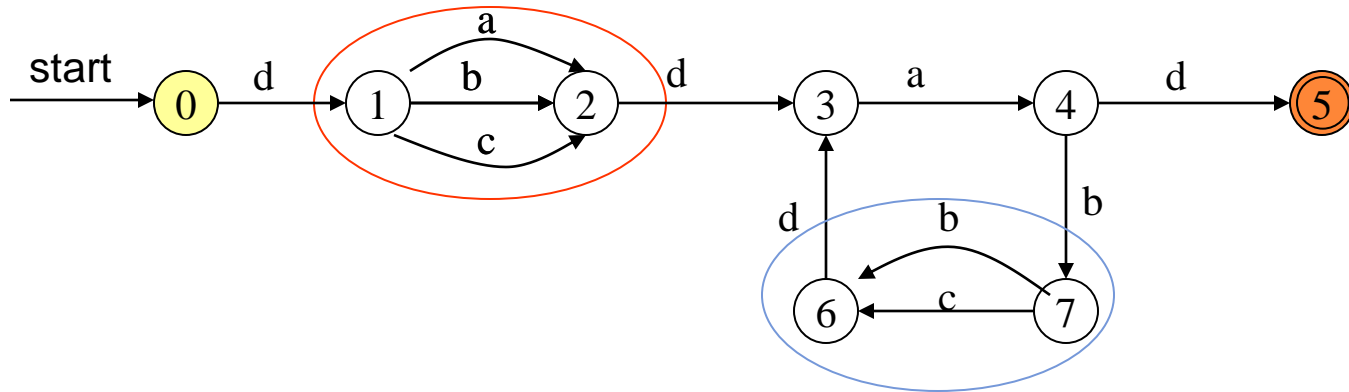
### **Lecture 03**



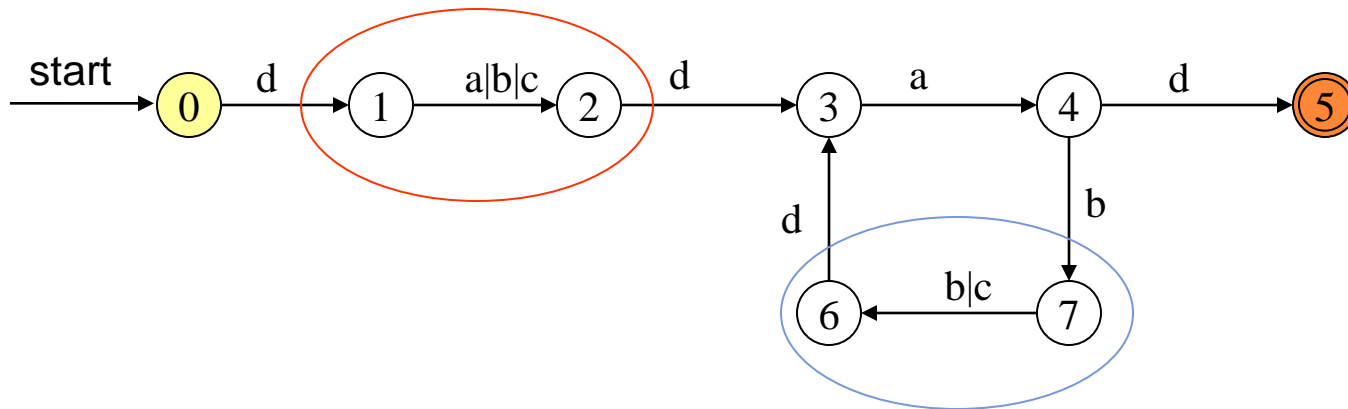
# CONVERTING DFAS TO REs

1. Combine serial links by concatenation
2. Combine parallel links by alternation
3. Remove self-loops by Kleene closure
4. Select a node (other than initial or final) for removal.  
Replace it with a set of equivalent links whose path expressions correspond to the in and out links
5. Repeat steps 1-4 until the graph consists of a single link between the entry and exit nodes.

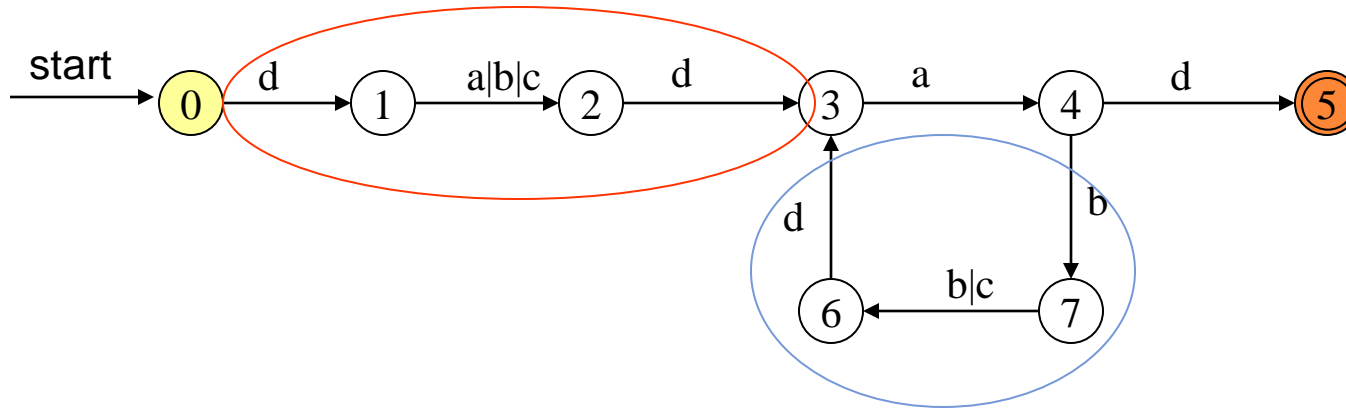
# EXAMPLE



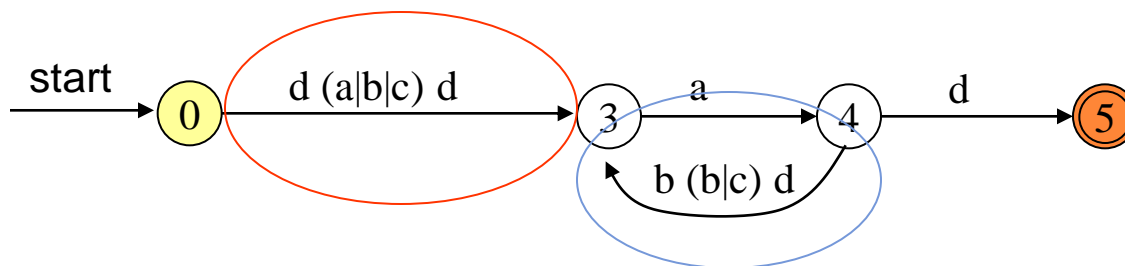
parallel edges become alternation



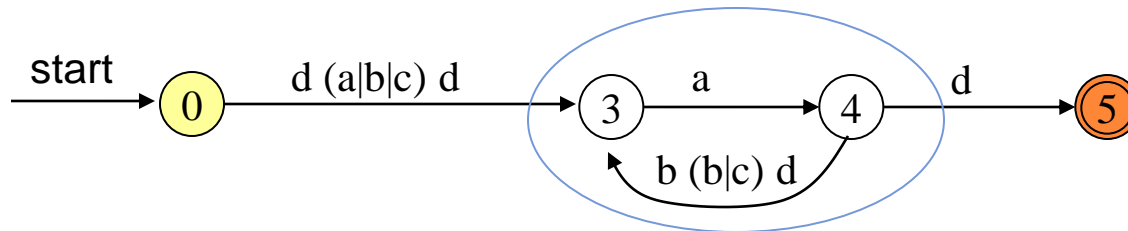
# EXAMPLE



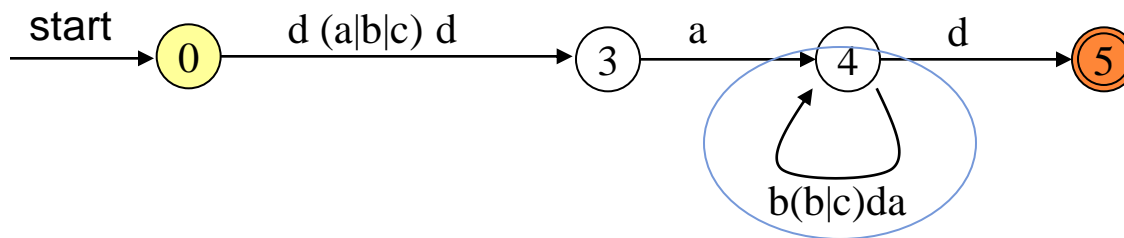
serial edges become concatenation



# EXAMPLE



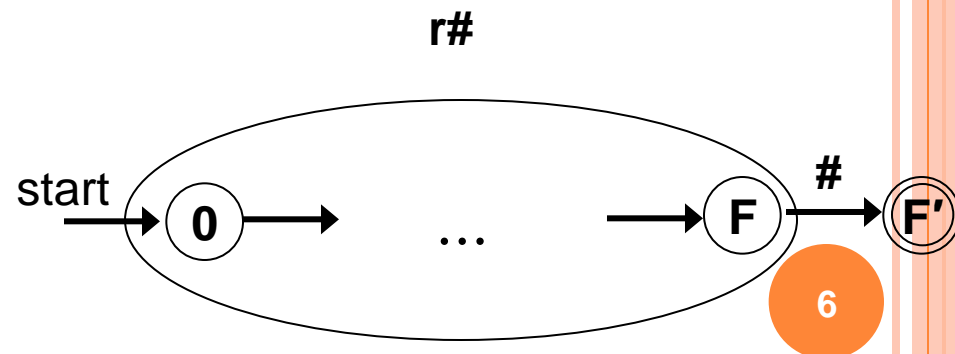
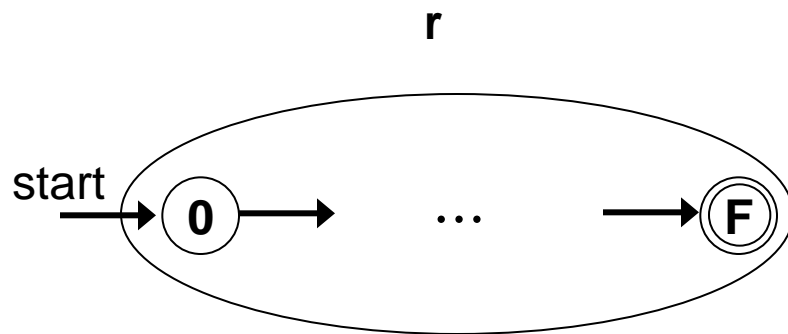
Find paths that can be “shortened”



# REGULAR EXPRESSION TO DFA

## Important States of NFA

- If it has a non- $\epsilon$  out-transition
- $move(s, a)$  is non-empty if  $s$  is important
- Accepting states are not important states
  - Adding a unique marker  $\#$  after the RE  $r$  (i.e.  $r\#$ ) we can make the accepting states important
  - Now a state with a transition on  $\#$  will be accepting state



# SYNTAX TREE

- Augmented RE ( $r\#$ ) can be represented by a syntax tree

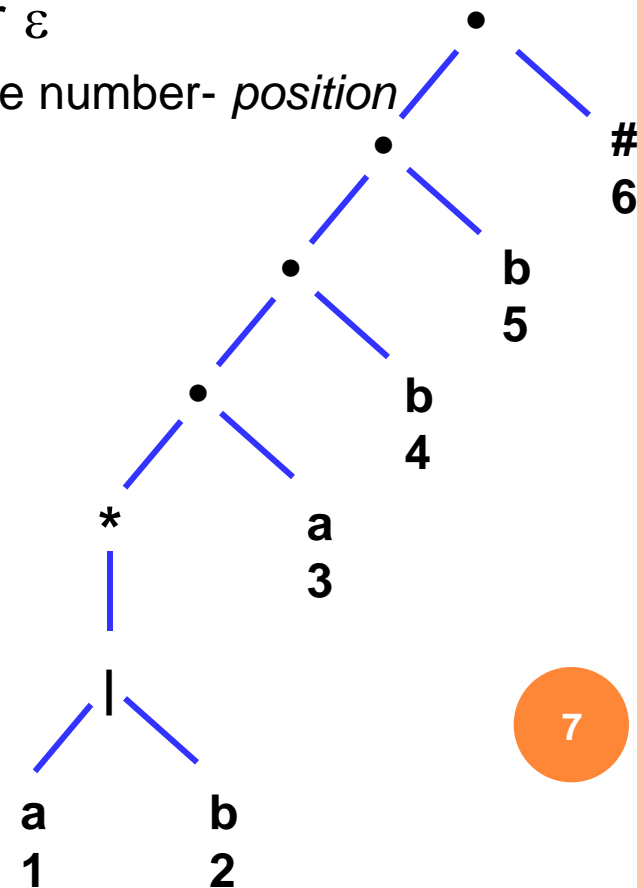
- Leaves contain: Alphabet symbols or  $\varepsilon$

- Each non- $\varepsilon$  leaf is associated with a unique number- *position* of the leaf and *position* of the symbol

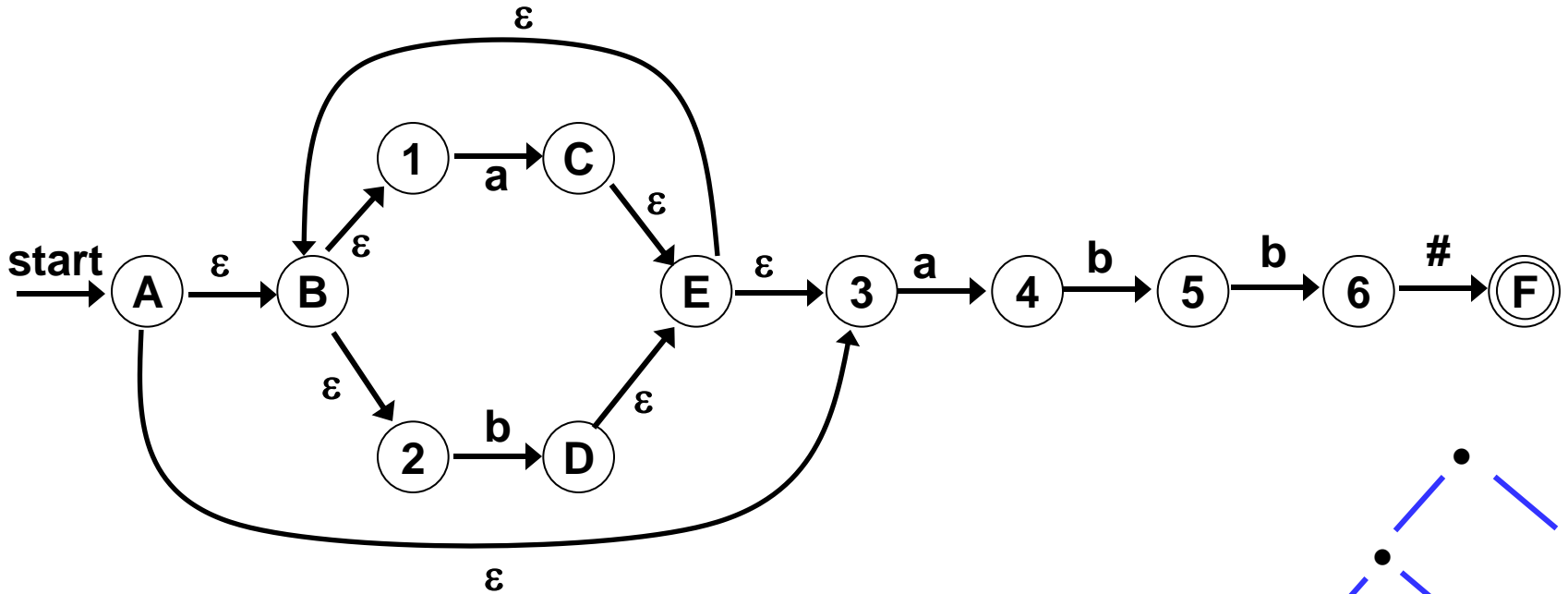
- Internal nodes contain: Operators

- cat-node*, *or-node* or *star-node*

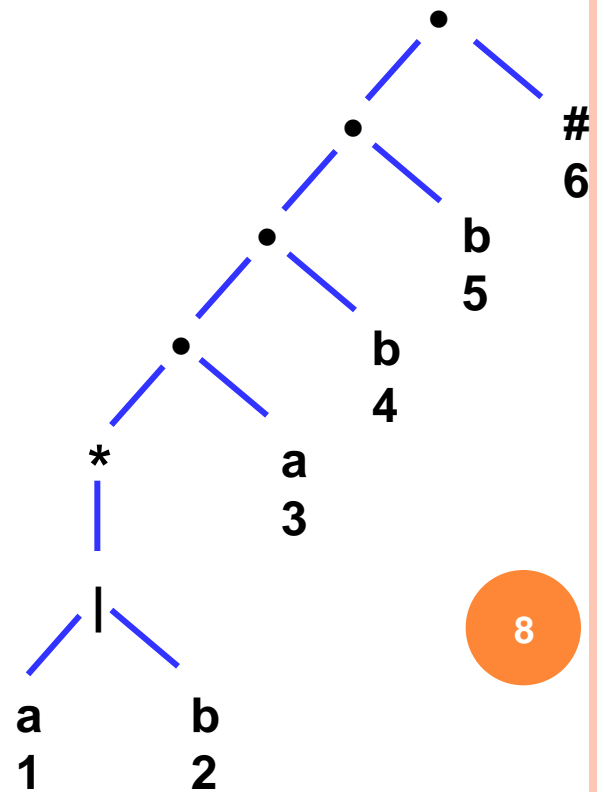
- Syntax tree for  $r\# = (a|b)^*abb\#$



# NFA FOR $(A|B)^*ABB\#$

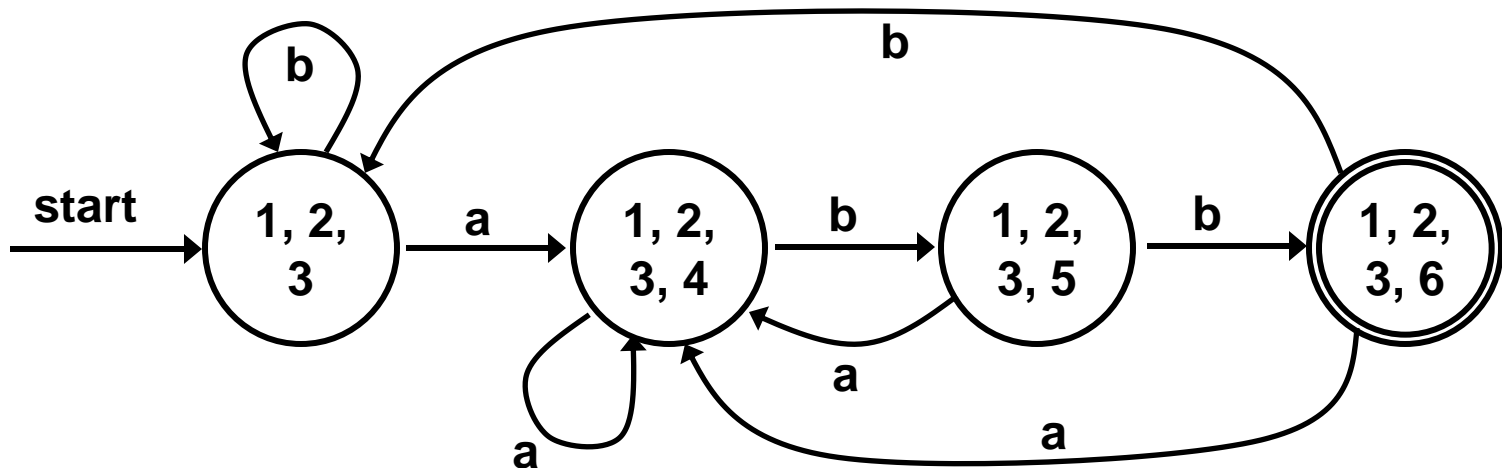
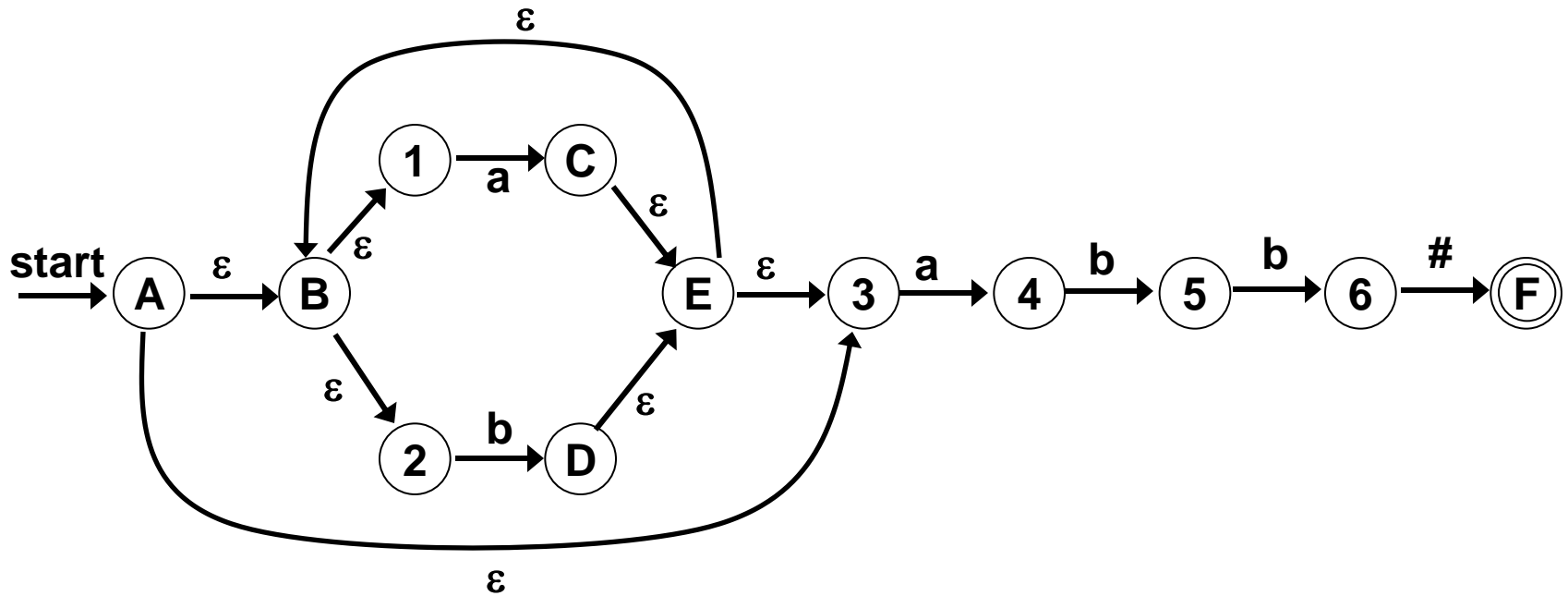


- Lettered states are non-important states
- Number states are important states
  - Numbers correspond to the number in syntax tree



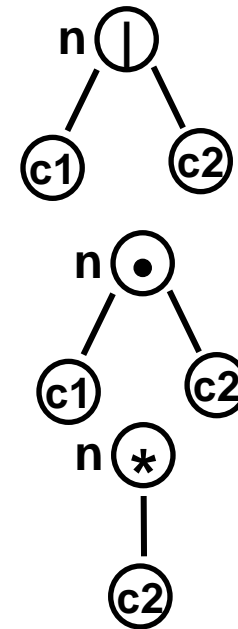


# DFA FOR $(A|B)^*ABB\#$



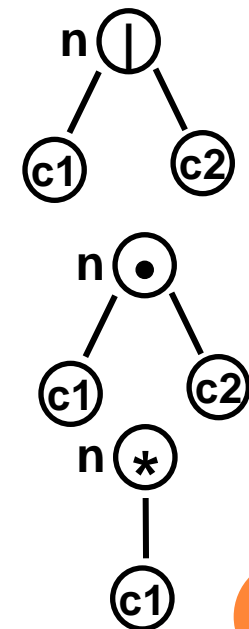
# TERMINOLOGY

- Nullable:
  - Nodes that are the root of some sub-expression that generate empty string
- If  $n$  is a leaf labeled by  $\varepsilon$  then
  - **nullable ( $n$ ) = true**
- If  $n$  is a leaf labeled with position  $i$ 
  - **nullable ( $n$ ) = false**
- If  $n$  is an or-node ( $|$ ) with children  $c1$  and  $c2$ 
  - **nullable ( $n$ ) = nullable( $c1$ ) or nullable ( $c2$ )**
- If  $n$  is an cat-node ( $\bullet$ ) with children  $c1$  and  $c2$ 
  - **nullable ( $n$ ) = nullable( $c1$ ) and nullable ( $c2$ )**
- If  $n$  is an star-node ( $*$ ) with children  $c1$ 
  - **nullable ( $n$ ) = true**



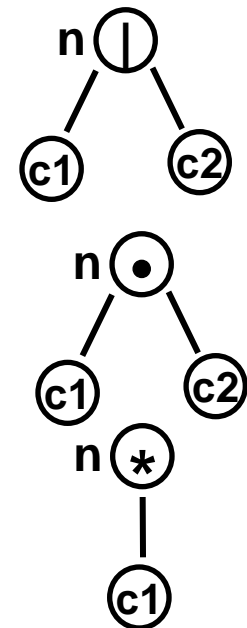
# TERMINOLOGY

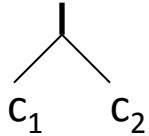
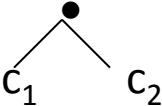

- **Firstpos(n):**
  - Set of positions that can match the first symbol of a string generated by the sub-expression rooted at n
- If n is a leaf labeled by  $\varepsilon$  then
  - **firstpos (n) =  $\emptyset$**
- If n is a leaf labeled with position  $i$ 
  - **firstpos (n) = {i}**
- If n is an or-node (|) with children c1 and c2
  - **firstpos (n) = firstpos(c1)  $\cup$  firstpos (c2)**
- If n is a cat-node (•) with children c1 and c2
  - **firstpos(n) = If nullable (c1) then firstpos(c1)  $\cup$  firstpos (c2) else firstpos(c1)**
- If n is an star-node (\*) with children c1
  - **firstpos (n) = firstpos(c1)**



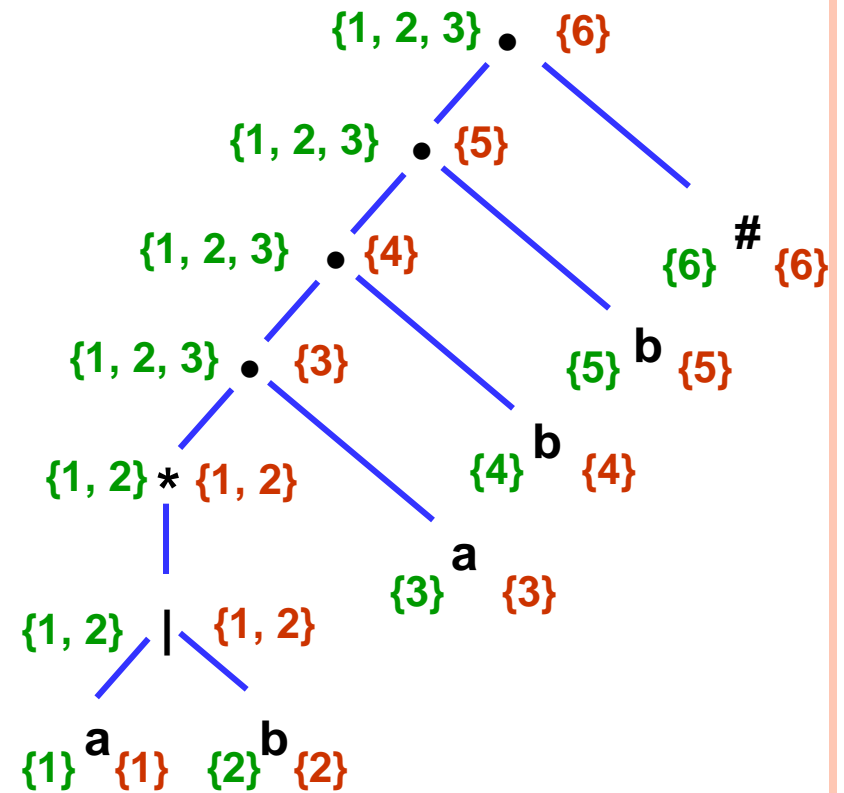
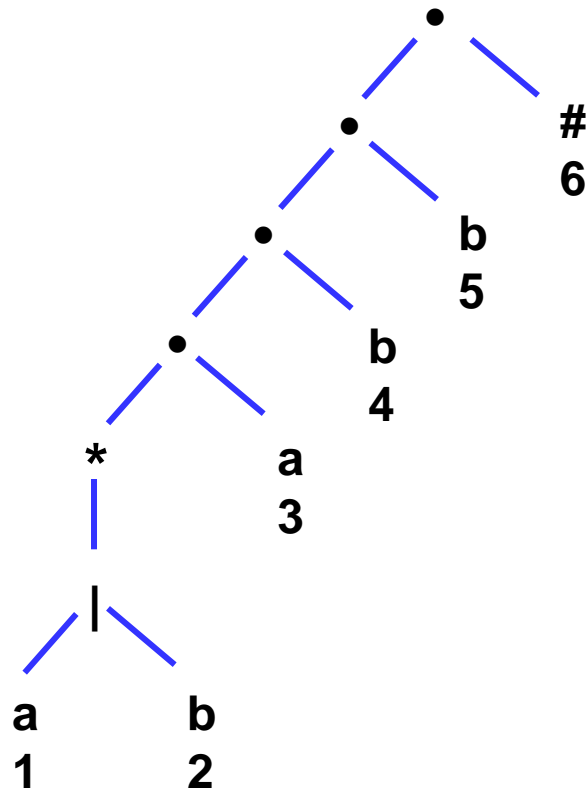
# TERMINOLOGY

- Lastpos(n):
  - Set of positions that can match the last symbol of a string generated by the sub-expression rooted at n
- If n is a leaf labeled by  $\varepsilon$  then
  - **lastpos (n) =  $\emptyset$**
- If n is a leaf labeled with position  $i$ 
  - **lastpos (n) = {i}**
- If n is an or-node (|) with children c1 and c2
  - **lastpos (n) = lastpos(c1)  $\cup$  lastpos (c2)**
- If n is an cat-node ( $\bullet$ ) with children c1 and c2
  - **lastpos(n) = If nullable (c2) then lastpos(c1)  $\cup$  lastpos (c2)**
  - **else lastpos(c2)**
- If n is an star-node ( $*$ ) with children c1
  - **lastpos (n) = lastpos(c1)**



n	nullable(n)	firstpos(n)	lastpos(n)
leaf labeled $\varepsilon$	true	$\Phi$	$\Phi$
leaf labeled with position i	false	$\{i\}$	$\{i\}$
	$\text{nullable}(c_1) \text{ or } \text{nullable}(c_2)$	$\text{firstpos}(c_1) \cup \text{firstpos}(c_2)$	$\text{lastpos}(c_1) \cup \text{lastpos}(c_2)$
	$\text{nullable}(c_1) \text{ and } \text{nullable}(c_2)$	if ( $\text{nullable}(c_1)$ ) $\text{firstpos}(c_1) \cup \text{firstpos}(c_2)$ else $\text{firstpos}(c_1)$	if ( $\text{nullable}(c_2)$ ) $\text{lastpos}(c_1) \cup \text{lastpos}(c_2)$ else $\text{lastpos}(c_2)$
	<b>true</b>	<b>firstpos(<math>c_1</math>)</b>	<b>lastpos(<math>c_1</math>)</b>

## FIRSTPOS AND LASTPOS EXAMPLE



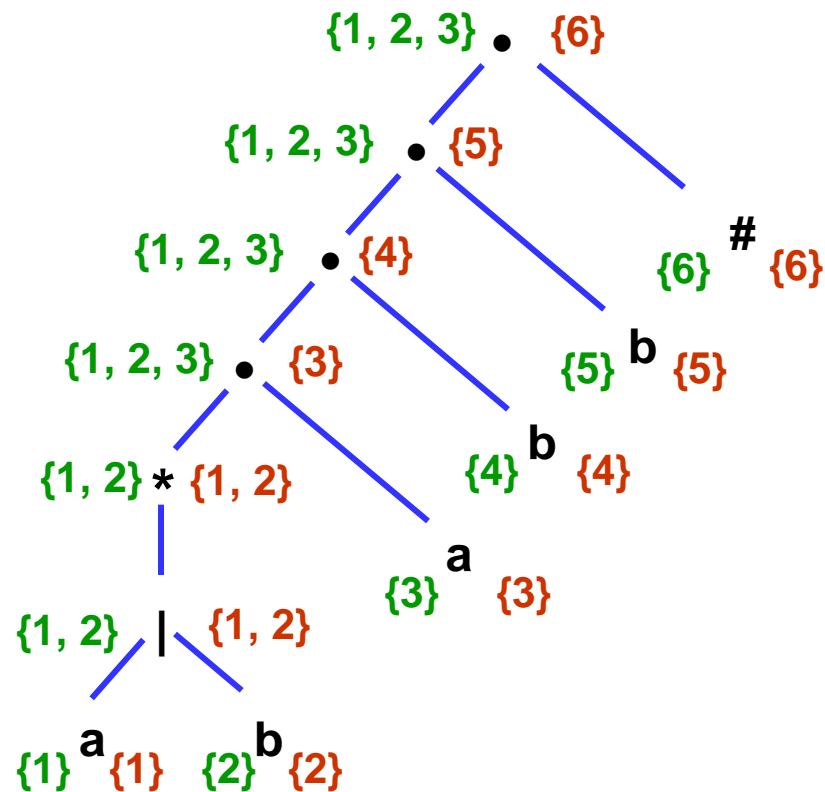
# TERMINOLOGY

- Followpos( $i$ ):
  - Tells what positions can follow position  $i$  in the syntax tree
- **Rule 1:**

If  $n$  is a cat-node with left child  $c1$  and right child  $c2$  and  $i$  is a position in lastpos( $c1$ ), then all positions in firstpos( $c2$ ) are in followpos( $i$ )
- **Rule 2:**

If  $n$  is a star node, and  $i$  is a position in lastpos( $n$ ), then all positions in firstpos( $n$ ) are in followpos( $i$ )

# FOLLOWPOS EXAMPLE



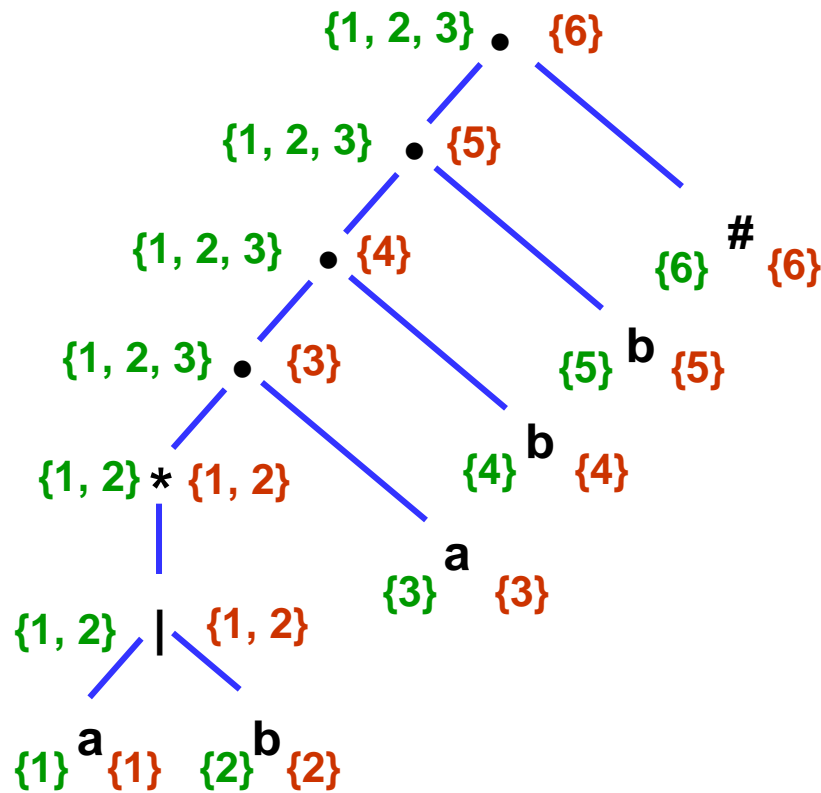
Node	followpos
1	{1,2,
2	{1,2,
3	{
4	{
5	{
6	{

## ○ At star-node:

- $lastpos(*) = \{1,2\}$  and  $firstpos(*) = \{1,2\}$
- According to Rule 2:
  - $followpos\{1\} = \{1,2\}$
  - $followpos\{2\} = \{1,2\}$



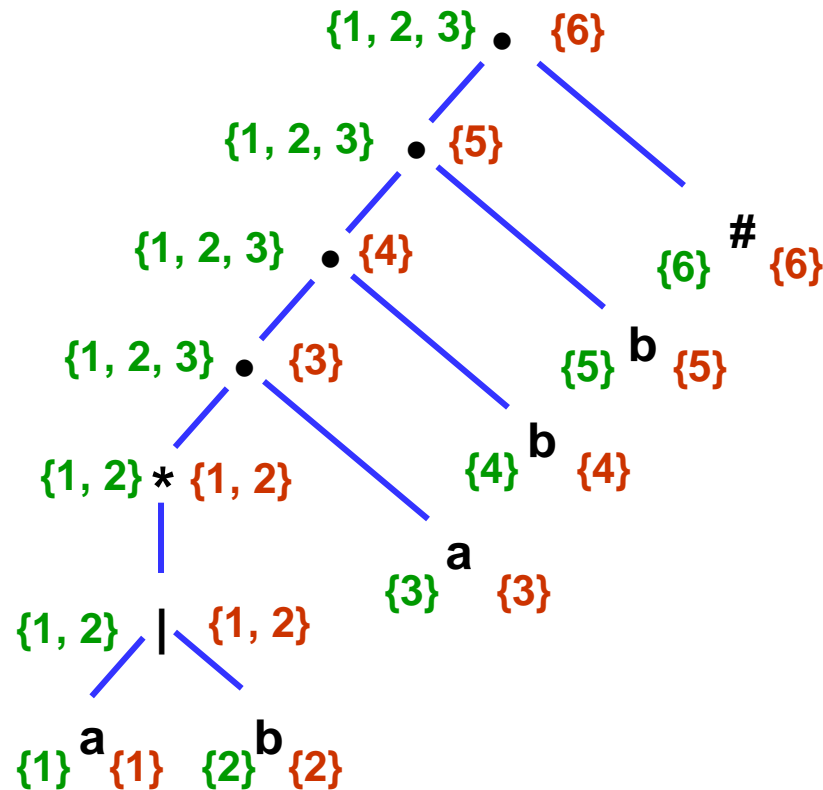
# FOLLOWPOS EXAMPLE



Node	Followpos
1	{1,2,3
2	{1,2,3
3	{
4	{
5	{
6	{

- At cat-node above the star-node, '\*' is left child and 'a' is right child
  - $lastpos(*) = \{1,2\}$  and  $firstpos(a) = \{3\}$
  - According to Rule 1:
    - $followpos\{1\} = \{3\}$
    - $followpos\{2\} = \{3\}$

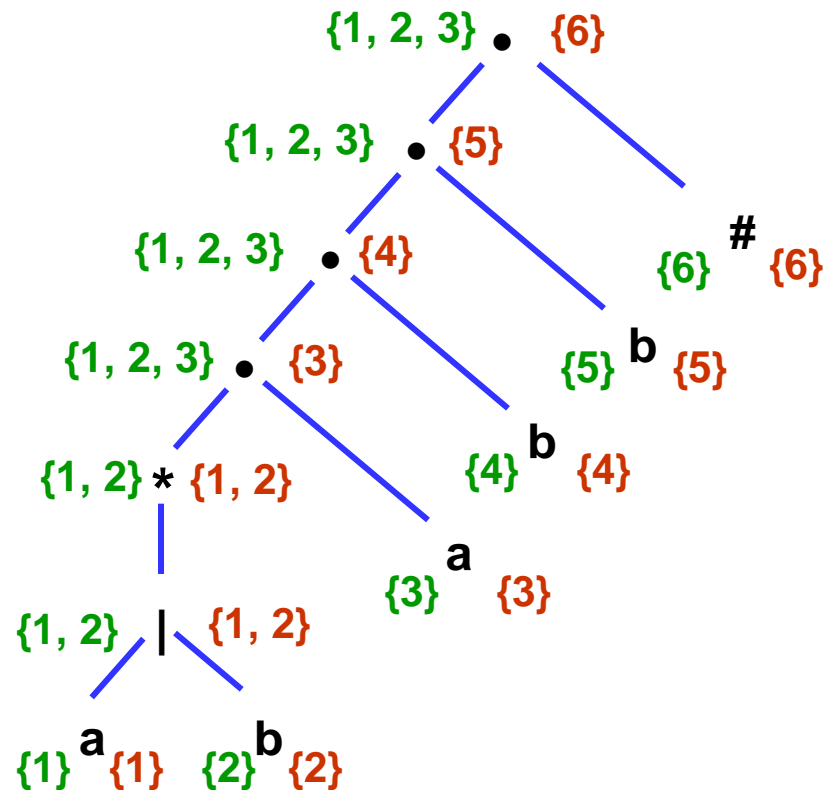
# FOLLOWPOS EXAMPLE



Node	Followpos
1	{1,2,3
2	{1,2,3
3	{4
4	{5
5	{6
6	{

- At next cat-node '•' is left child and 'b' is right child
  - $lastpos(\bullet) = \{3\}$  and  $firstpos(b) = \{4\}$
  - According to Rule 1:
    - $followpos\{3\} = \{4\}$
- Similarly,  $followpos\{4\} = \{5\}$  and  $followpos\{5\} = \{6\}$

## FOLLOWPOS EXAMPLE

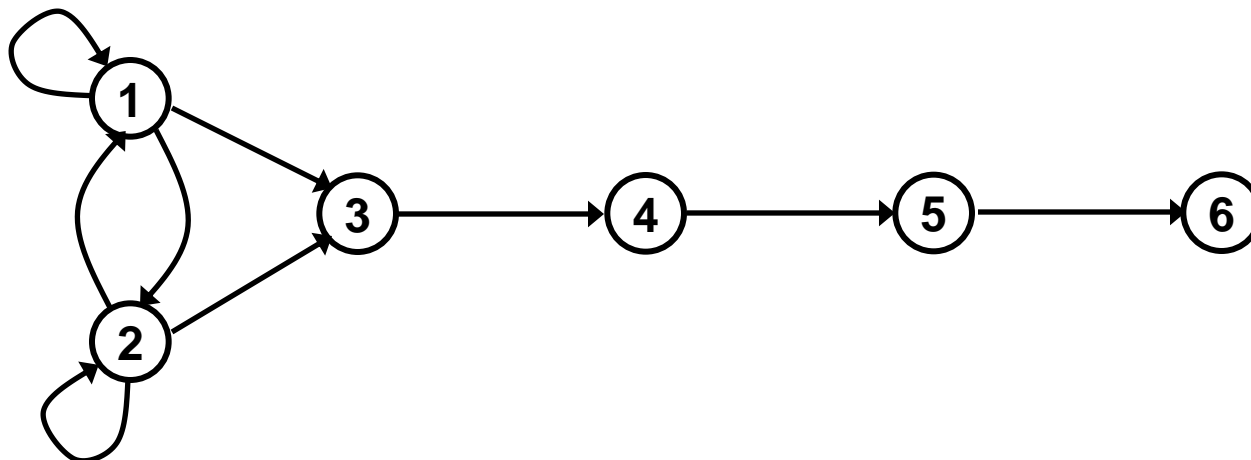


Node	followpos
1	{1,2,3}
2	{1,2,3}
3	{4}
4	{5}
5	{6}
6	-

# FOLLOWPOS GRAPH

- A node for each position
- Edge from node  $i$  to node  $j$  if  $j \in \text{followpos}\{i\}$
- *followpos* graph becomes equivalent NFA without  $\varepsilon$ -transition if
  - All positions in *firstpos* of root become start state
  - Label edge  $\{i,j\}$  by the symbol at position  $j$
  - Position associated with  $\#$  only accepting state

Node	followpos
1	{1,2,3}
2	{1,2,3}
3	{4}
4	{5}
5	{6}
6	-



# CONSTRUCTION OF DFA FROM RE

- Input: A regular expression  $r$
- Output: A DFA  $D$  that recognizes  $L(r)$
- Method:
  1. Construct syntax tree  $ST$  for augmented RE  $r\#$
  2. Construct the functions nullable, firstpos, lastpos and followpos for  $ST$
  3. Construct  $Dstates$ : set of states of  $D$   
 $Dtrans$ : transition table for  $D$

# CONSTRUCTION OF DFA FROM RE

## ○ Algorithm

Initially, the only unmarked state in **Dstates** is *firstpos(root)*  
while there is an unmarked state **T** in **Dstates** do begin

    Mark **T**;

    For each input symbol **a** do begin

        Let **U** be the set of positions that are in followpos(**p**) for  
        some position **p** in **T** such that the symbol at position  
        **p** is **a**

        If **U** is not empty and is not in **Dstates** then

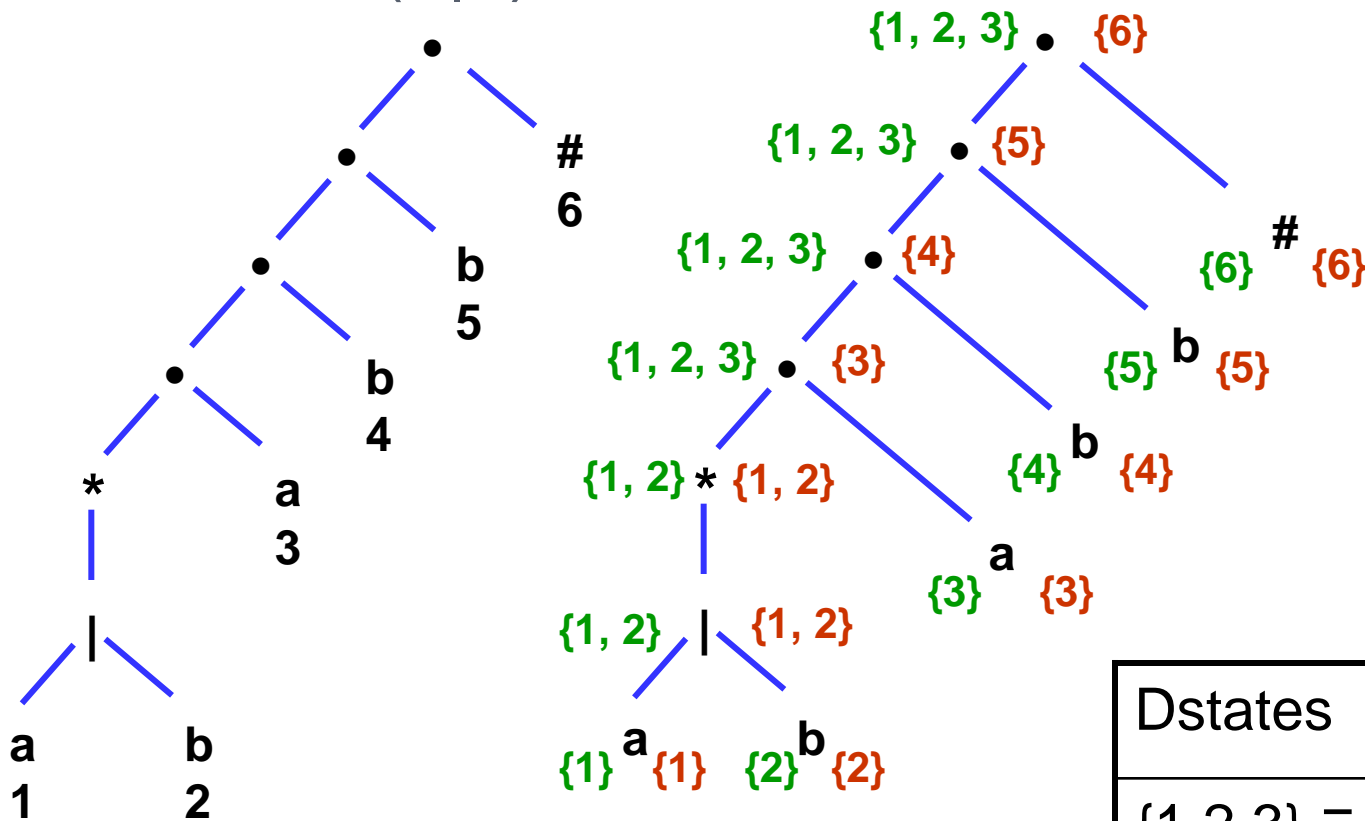
            Add **U** as an unmarked states to **Dstates**

**Dtrans[T,a]=U**

    End

end

# DFA FOR $(A|B)^*ABB\#$



firstpos{root} = {1,2,3}  $\equiv$  A (unmarked)

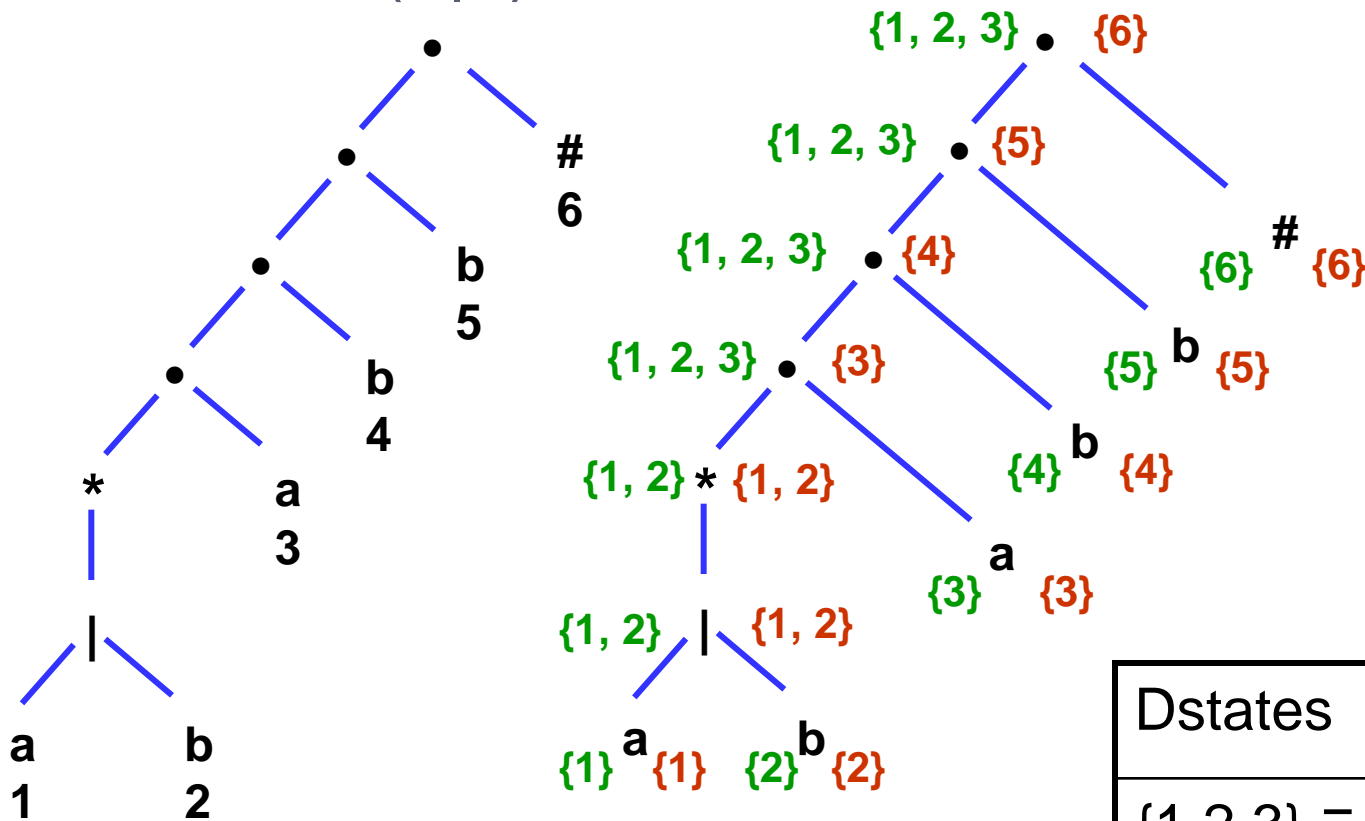
For the input symbol **a**, positions are 1, 3  
 $\therefore$  followpos(1)  $\cup$  followpos{3}  
 $= \{1,2,3,4\} \equiv B$

For the input symbol **b**, positions are 2  
 $\therefore$  followpos(2) = {1,2,3,}  $\equiv A$

Node	followpos
1	{1,2,3}
2	{1,2,3}
3	{4}
4	{5}
5	{6}
6	-

Dstates	a	b
{1,2,3} $\equiv$ A	B	A
{1,2,3,4} $\equiv$ B		

# DFA FOR $(A|B)^*ABB\#$




$\{1,2,3,4\} \equiv B$  (unmarked)

For the input symbol **a**, positions are 1, 3  
 $\therefore \text{followpos}(1) \cup \text{followpos}\{3\}$   
 $= \{1,2,3,4\} \equiv B$

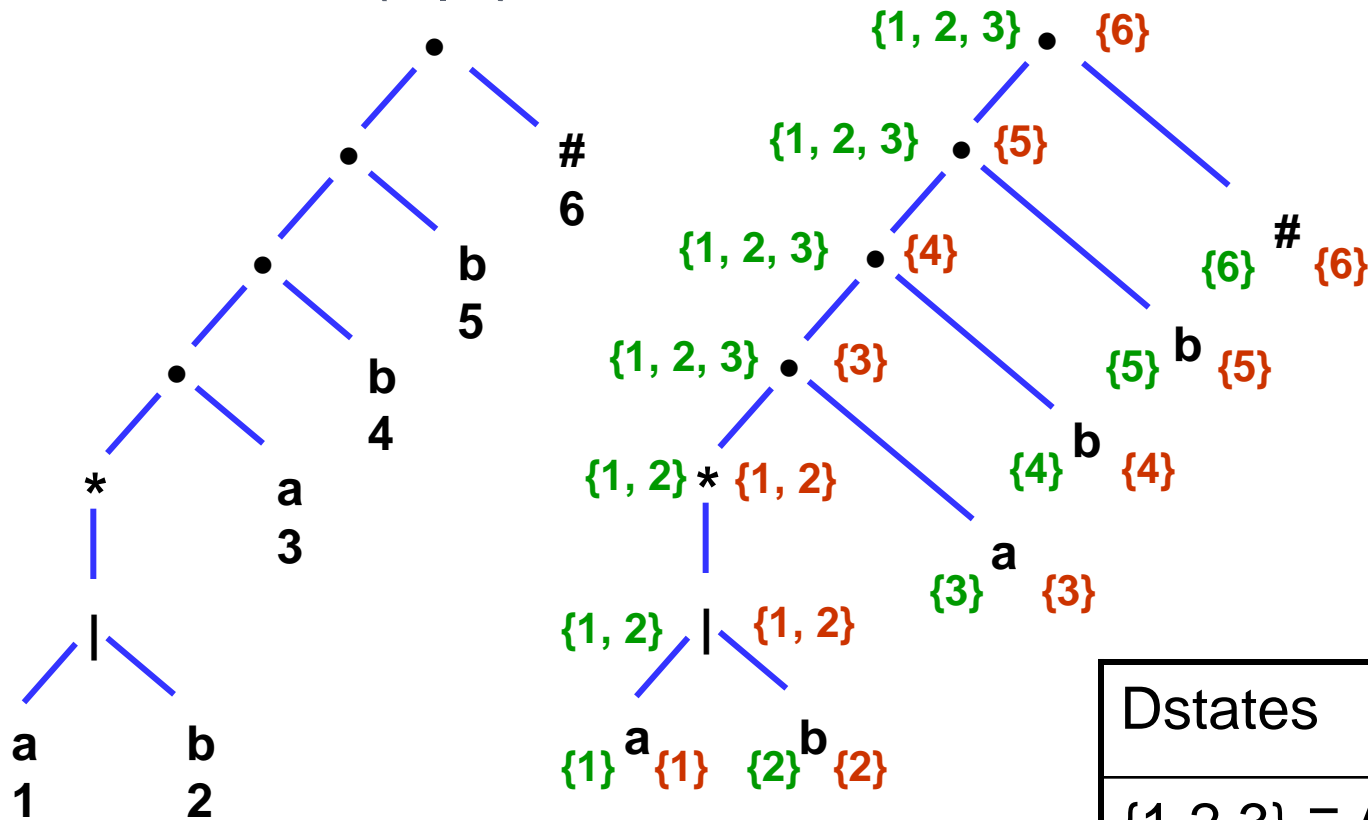
For the input symbol **b**, positions are 2, 4  
 $\therefore \text{followpos}(2) \cup \text{followpos}\{4\}$   
 $= \{1,2,3,5\} \equiv C$

Node	followpos
1	{1,2,3}
2	{1,2,3}
3	{4}
4	{5}
5	{6}
6	-

Dstates	a	b
$\{1,2,3\} \equiv A$	B	A
$\{1,2,3,4\} \equiv B$	B	C
$\{1,2,3,5\} \equiv C$		



# DFA FOR $(A|B)^*ABB\#$



Node	followpos
1	{1,2,3}
2	{1,2,3}
3	{4}
4	{5}
5	{6}
6	-

$\{1,2,3,5\} \equiv C$  (unmarked)

For the input symbol **a**, positions are 1, 3

$\therefore \text{followpos}(1) \cup \text{followpos}\{3\}$

$= \{1,2,3,4\} \equiv B$

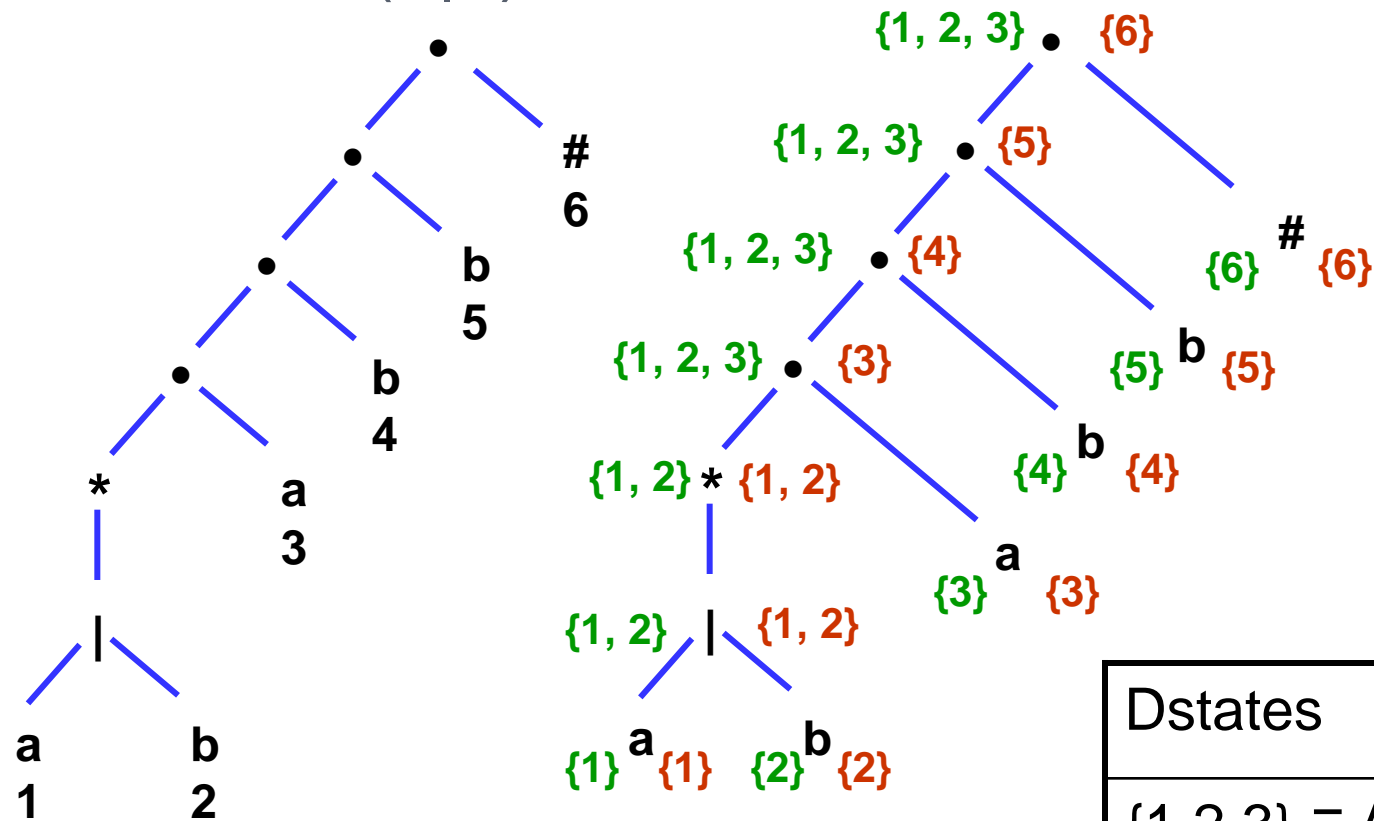
For the input symbol **b**, positions are 2, 5

$\therefore \text{followpos}(2) \cup \text{followpos}\{5\}$

$= \{1,2,3,6\} \equiv D$

Dstates	a	b
$\{1,2,3\} \equiv A$	B	A
$\{1,2,3,4\} \equiv B$	B	C
$\{1,2,3,5\} \equiv C$	B	D
$\{1,2,3,6\} \equiv D$		

# DFA FOR $(A|B)^*ABB\#$



Node	followpos
1	{1,2,3}
2	{1,2,3}
3	{4}
4	{5}
5	{6}
6	-

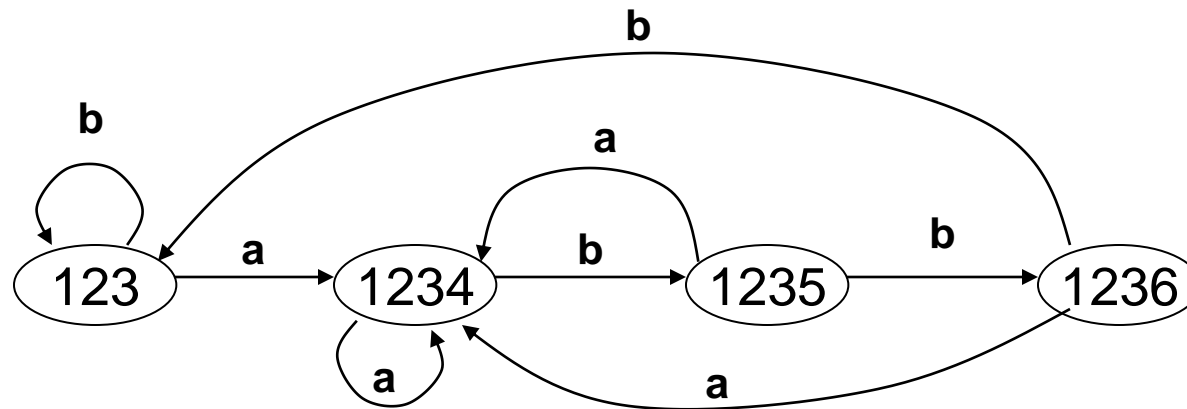
$\{1,2,3,6\} \equiv D$  (unmarked)

For the input symbol **a**, positions are 1, 3  
 $\therefore \text{followpos}(1) \cup \text{followpos}\{3\}$   
 $= \{1,2,3,4\} \equiv B$

For the input symbol **b**, positions are 2  
 $\therefore \text{followpos}(2)$   
 $= \{1,2,3\} \equiv A$

Dstates	a	b
$\{1,2,3\} \equiv A$	B	A
$\{1,2,3,4\} \equiv B$	B	C
$\{1,2,3,5\} \equiv C$	B	D
$\{1,2,3,6\} \equiv D$	B	A

# DFA FOR $(A|B)^*ABB\#$



# Thank You