

i) Aim : To implement banker's Algorithm.

	Max			Allocation			Available		
	A	B	C	A	B	C	A	B	C
P ₀	0	0	1	0	0	1			
P ₁	1	7	5	1	0	0			
P ₂	2	3	5	1	3	5			
P ₃	0	6	5	0	6	3			
Total				2	9	9	1	5	2

$$\text{Need} = \text{Max} - \text{Allocation}$$

	Need		
	A	B	C
P ₀	0	0	0
P ₁	0	7	5
P ₂	1	0	0
P ₃	0	0	2

Safety Algorithm.

$$\text{Work} = \text{Available}$$

$$\text{Finish}[i] = \text{false} \quad i = 0, 1, 2, 3$$

Finish :

F	F	F	F
0	1	2	3

$$\text{Need}_1 \leq \text{work}$$

$$\text{Need}_0 = 0, 0, 0$$

$$\text{Need}_0 \leq \text{work} \quad \checkmark$$
$$0, 0, 0 \quad 1, 5, 2$$

$$\text{finish}[0] = \text{True}$$

P0

$$\text{Work} = (1, 5, 2) + (0, 0, 1)$$

$$\text{Work} = 1 \quad 5 \quad 3$$

$$\text{Need}_1 = 0, 7, 5$$

$$\text{P1 } \text{Need}_1 > \text{Work}$$

$$(0, 7, 5) > (1, 5, 3) \quad \times$$

$$\text{Need}_2 = 1, 0, 0$$

P2

$$\text{Need}_2 \leq \text{Work}$$

$$(1, 0, 0) \leq (1, 5, 3)$$

$$\text{finish}[2] = \text{True} \quad \checkmark$$

$$\text{Work} = (1, 5, 3) + (1, 3, 5)$$

$$= (2, 8, 8)$$

P3

$$\text{Need}_3 = 0, 0, 2$$

$$\text{Need}_3 \leq \text{Work}$$

$$(0, 0, 2) \leq (2, 8, 8) \quad \checkmark$$

$$\text{finish}[3] = \text{true}$$

$$\text{Work} = \cancel{(0, 0, 2)} + \cancel{(2, 8, 8)} + (0, 6, 3)$$

$$\text{Work} = (2, 14, 11)$$

P2

$$\text{Need}_2 = 0, 7, 5$$

$$\text{Need}_2 \leq \text{Work}$$

$$(0, 7, 5) \leq (2, 14, 11)$$

$$\text{Work} \leftarrow \text{finish}[2] = \text{True}$$

$$\text{Work} = (2, 14, 11) + (1, 0, 0)$$

$$\text{Work} = (3, 14, 11)$$

$$\text{finish}[i] = \text{true} \quad \text{for } 0 \leq i \leq 3$$

hence system in safe state

$$\text{Finish} = \begin{bmatrix} T & T & T & T \end{bmatrix}$$

Safe sequence : P_0, P_2, P_3, P_1

P1 request additional resource (0, 5, 2)

$$req_1 \leq Need$$

$$(0, 5, 2) \leq (0, 7, 5) \quad \checkmark$$

$$req_1 \leq Available$$

$$(0, 5, 2) \leq (1, 5, 2) \quad \checkmark$$

$$\begin{aligned} Available &= Available - request_i \\ &= (1, 5, 2) - (0, 5, 2) \\ &= (1, 0, 0) \end{aligned}$$

$$\begin{aligned} Allocation &= Allocation + request_i \\ &= (1, 0, 0) + (0, 5, 2) \\ &= (1, 5, 2) \end{aligned}$$

$$\begin{aligned} Need &= Need - request \\ &= (0, 7, 5) - (0, 5, 2) \\ &= (0, 2, 3) \end{aligned}$$

Check for safe sequence

	Max			Allocation			Available		
	A	B	C	A	B	C	A	B	C
P0	0	0	1	0	0	1			
P1	1	7	5	1	5	2			
P2	2	3	3	1	3	5			
P3	0	6	5	0	6	3			
Total				2	14	11	1	0	0

	Need		
	A	B	C
P0	0	0	0
P1	0	2	3
P2	1	0	0
P3	0	0	2

safety Algorithm

Work = Available

finish =

F	F	F	F
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Step -1

Need₀ ≤ Work

P0 (0, 0, 0) (1, 0, 0) ✓

Work = (1, 0, 0) + (0, 0, 1)
= (1, 0, 1) finish[0] = True

Step 2

Need₁ > Work

X

P1

$$(0, 2, 3) > (1, 0, 1)$$

Finish[1] = false

Step 3

P2

Need₂ ≤ Work

$$(1, 0, 0) \leq (1, 0, 1)$$

✓

Finish[2] = True

$$\text{Work} = (1, 0, 1) + (1, 3, 5)$$

$$= (2, 3, 6)$$

Step 4

Need₃ ≤ Work

$$P3 \quad (0, 0, 2) \leq (2, 3, 6)$$

✓

Finish[3] = True

$$\text{Work} = (2, 3, 6) + (0, 6, 3)$$

$$= (2, 9, 9)$$

Step 5

Need₁ ≤ Work

$$(0, 2, 3) \leq (2, 9, 9)$$

✓

P1

Finish[1] = True

$$\text{Work} = (2, 9, 9) + (1, 5, 2) \\ = (3, 14, 11)$$

$$\text{finish}[i] = \text{true} \quad 0 \leq i \leq 3$$

Hence system is in safe state

The safe sequence is $\langle P_0, P_2, P_3, P_1 \rangle$

Yes we can grant request immediately

$$(i) \quad \begin{matrix} (1, 5, 2) \\ A \quad B \quad C \end{matrix}$$

$$(ii) \quad \begin{array}{c|c|c|c} & A & B & C \\ \hline P_0 & 0 & 0 & 0 \\ \hline P_1 & 0 & 7 & 5 \\ \hline P_2 & 0 & 7 & 5 \\ \hline P_3 & 0 & 0 & 2 \end{array}$$

(iii) Yes, It satisfied safety algorithm

(iv) $\langle P_0, P_2, P_3, P_1 \rangle$

(v) Yes, ~~It~~ Banker's Algorithm grant request immediately

(vi) New Table + safe sequence.