

Cascading Rollback | Cascadeless Schedule

👤 Akshay Singhal 📁 Database Management System

Recoverability-

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Before you go through this article, make sure that you have gone through the previous article on [Recoverability in DBMS](#).

We have discussed-

- Non-serial schedules which are not serializable are called as non-serializable schedules.
- Non-serializable schedules may be recoverable or irrecoverable.

Recoverable Schedules-

If in a schedule,

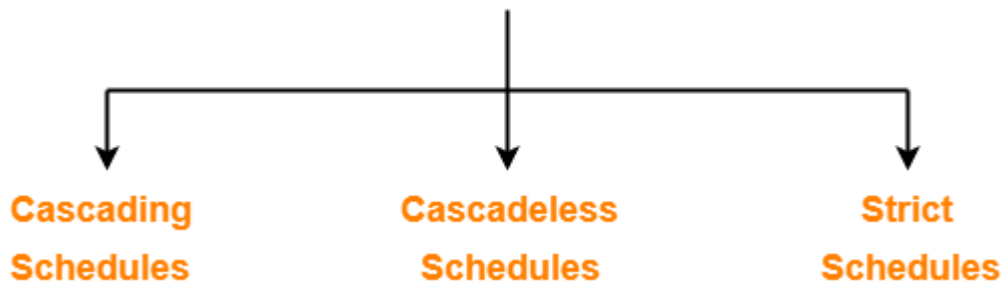
- A transaction performs a dirty read operation from an uncommitted transaction
- And its commit operation is delayed till the uncommitted transaction either commits or roll backs

then such a schedule is called as a **Recoverable Schedule**.

Types of Recoverable Schedules-

A recoverable schedule may be any one of these kinds-

Recoverable Schedules



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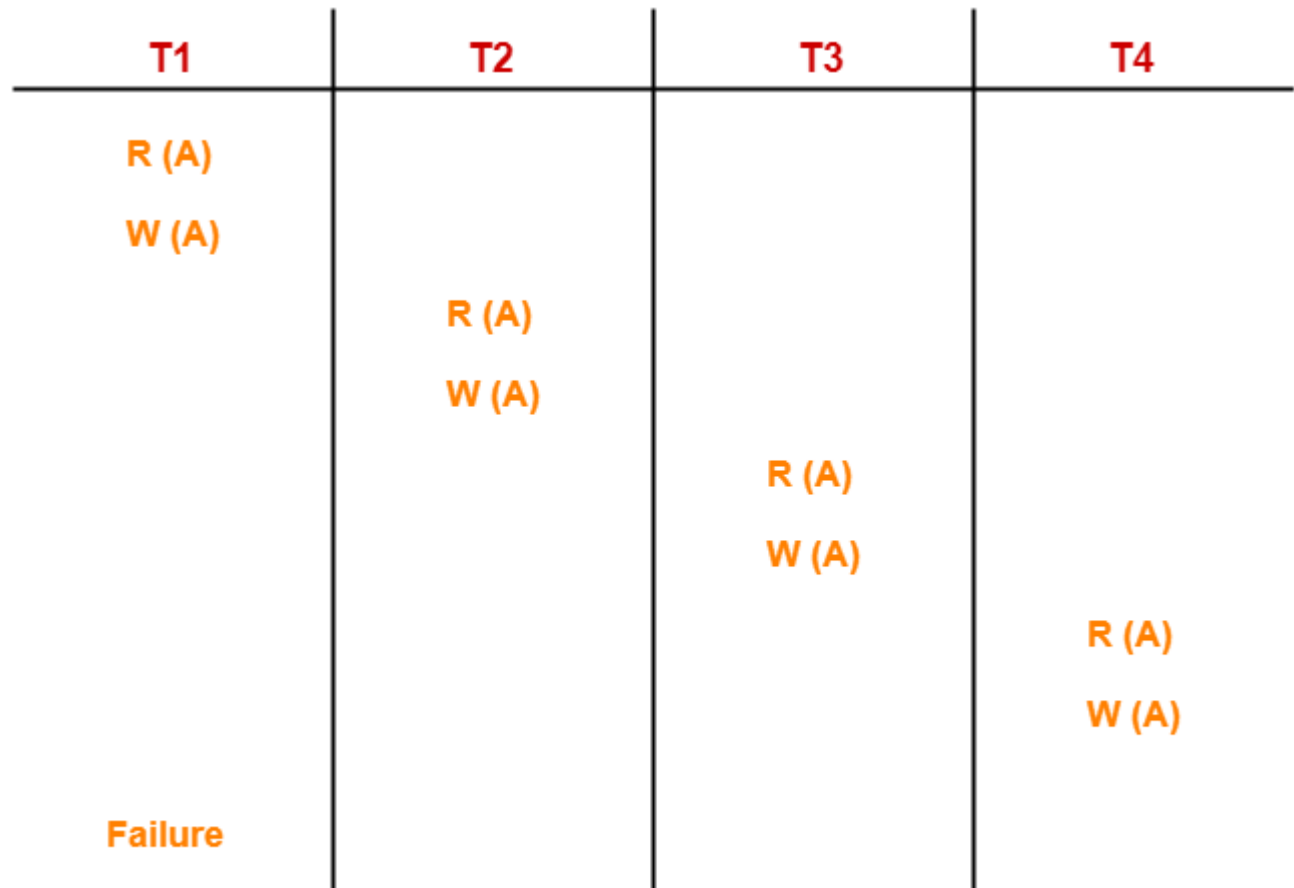
time management

1. Cascading Schedule
2. Cascadeless Schedule
3. Strict Schedule

Cascading Schedule-

- If in a schedule, failure of one transaction causes several other dependent transactions to rollback or abort, then such a schedule is called as a **Cascading Schedule** or **Cascading Rollback** or **Cascading Abort**.
- It simply leads to the wastage of CPU time.

Example-



Cascading Recoverable Schedule

Here,

- Transaction T2 depends on transaction T1.
- Transaction T3 depends on transaction T2.
- Transaction T4 depends on transaction T3.

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In this schedule,

- The failure of transaction T1 causes the transaction T2 to rollback.
- The rollback of transaction T2 causes the transaction T3 to rollback.
- The rollback of transaction T3 causes the transaction T4 to rollback.

Such a rollback is called as a **Cascading Rollback**.

NOTE-

If the transactions T2, T3 and T4 would have committed before the failure of transaction T1, then the schedule would have been irrecoverable.

Cascadeless Schedule-

If in a schedule, a transaction is not allowed to read a data item until the last transaction that has written it is committed or aborted, then such a schedule is called as a **Cascadeless Schedule**.

In other words,

- Cascadeless schedule allows only committed read operations.
- Therefore, it avoids cascading roll back and thus saves CPU time.

Example-

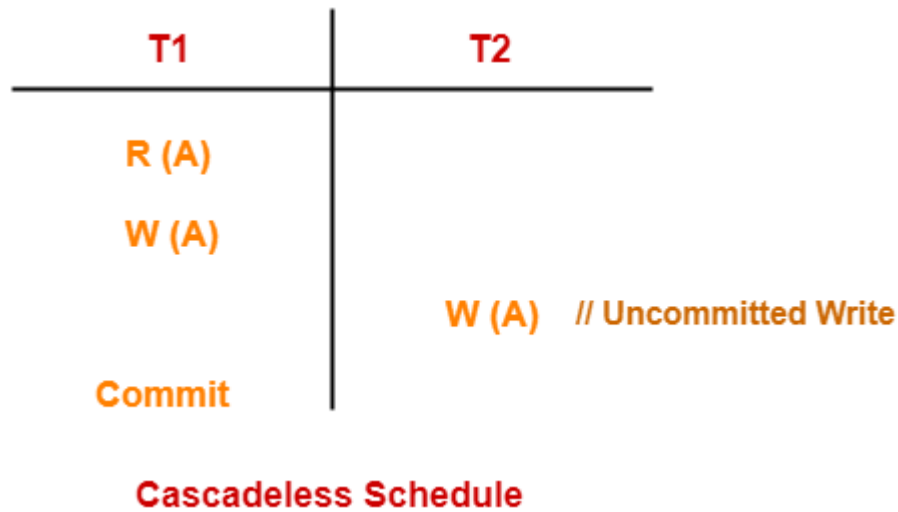
T1	T2	T3
R (A)		
W (A)		
Commit		
	R (A)	
	W (A)	
	Commit	
		R (A)
		W (A)
		Commit

Cascadeless Schedule

NOTE-

- Cascadeless schedule allows only committed read operations.
- However, it allows uncommitted write operations.

Example-



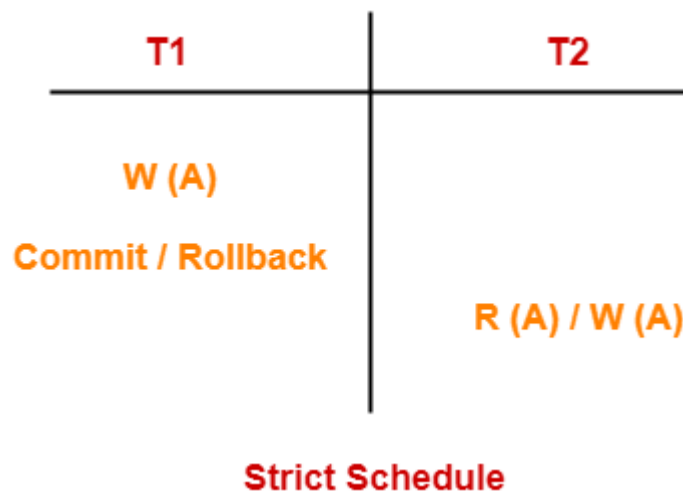
Strict Schedule-

If in a schedule, a transaction is neither allowed to read nor write a data item until the last transaction that has written it is committed or aborted, then such a schedule is called as a **Strict Schedule**.

In other words,

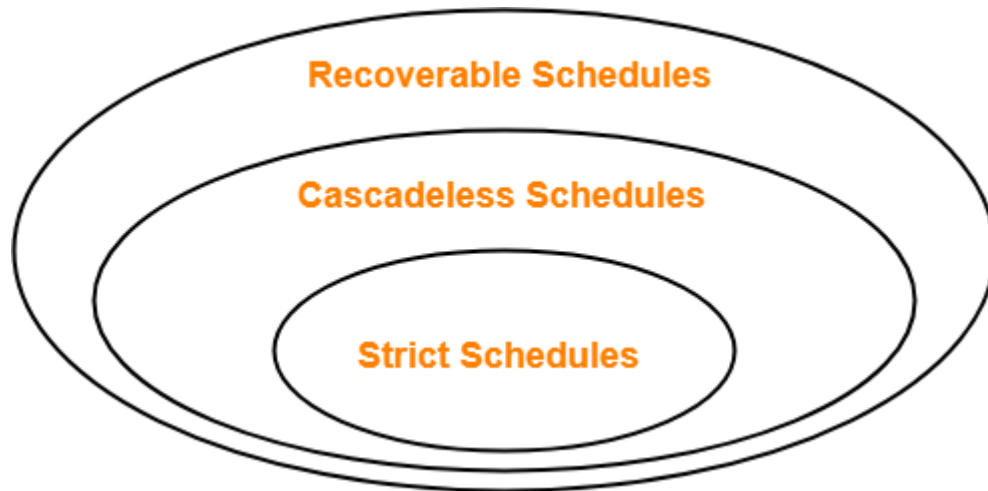
- Strict schedule allows only committed read and write operations.
- Clearly, strict schedule implements more restrictions than cascadeless schedule.

Example-



Remember-

- Strict schedules are more strict than cascadeless schedules.
- All strict schedules are cascadeless schedules.
- All cascadeless schedules are not strict schedules.



Next Article- [Equivalence of Schedules](#)

Get more notes and other study material of [Database Management System \(DBMS\)](#).

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T1	T2	T3	T4
R (A)			
	R (A)		
		R (A)	
			R (A)
W (B)			
	W (B)		
		W (B)	
			W (B)

T1	T2	T3	T4
	R(X)		
W(X)		W(X)	
Commit		Commit	
	W(Y)		
	R(Z)		
	Commit		
		R(X)	
		R(Y)	
		Commit	



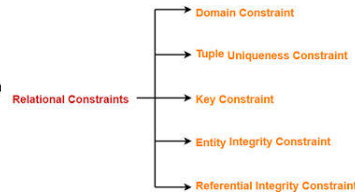
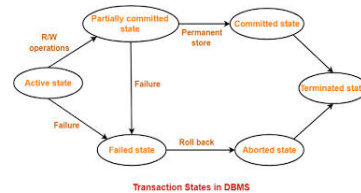
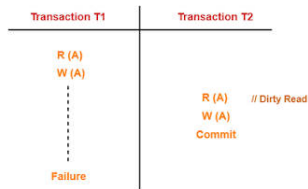
Schedule S1		Schedule S2		Schedule S3	
T1	T2	T1	T2	T1	T2
R(X)	R(X)	R(X)	R(X)	R(X)	R(X)
X = X + 5	X = X + 5	X = X + 5	X = X + 5	X = X + 5	X = X + 5
W(X)	W(X)	W(X)	W(X)	W(X)	W(X)
R(Y)	R(Y)	R(X)	R(X)	R(X)	R(X)
Y = Y + 5	Y = Y + 5	X = X + 5	X = X + 5	X = X + 5	X = X + 5
W(Y)	W(Y)	W(X)	W(X)	W(X)	W(X)
		R(Y)	R(Y)	R(Y)	R(Y)
		X = X + 5	Y = Y + 5	Y = Y + 5	Y = Y + 5
		W(X)	W(X)	W(Y)	W(Y)

View Serializability in DBMS | Practice Problems

Conflict Serializability | Practice Problems

View Serializability in DBMS

Equivalence of Schedules | Equivalent Schedules in DBMS

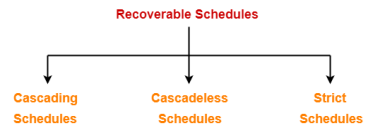


Concurrency Problems | DBMS

Transaction States in DBMS

Constraints in DBMS | Database Types of Constraints in System DBMS

Summary



Article Name Cascading Rollback | Cascadeless Schedule

Description Recoverable Schedules in DBMS are of three types- Cascading Rollback Schedule or Cascading Abort, Cascadeless Schedule and Strict Schedule. Cascadeless schedules are less strict than strict schedules.

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