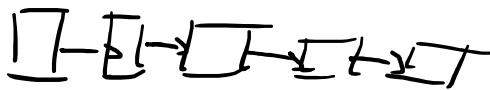
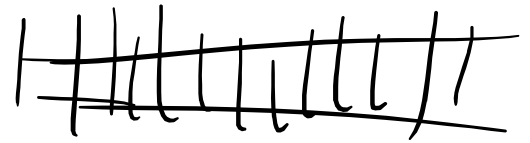
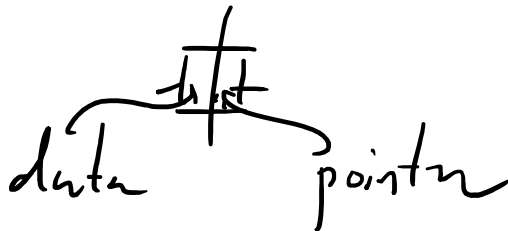


Linked Lists:

List \Leftrightarrow Array



Object oriented Space

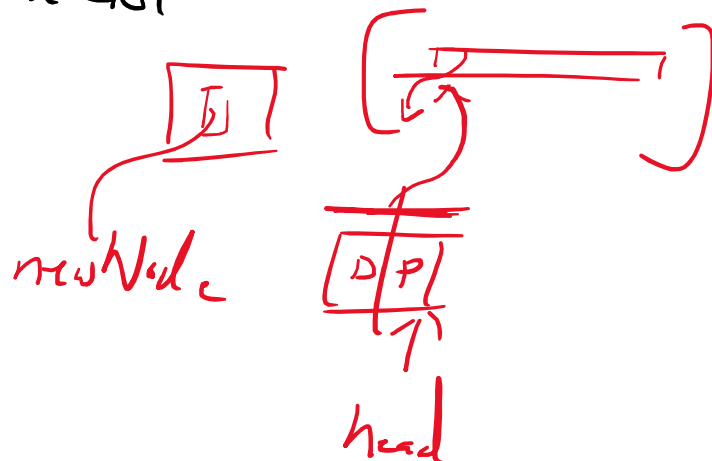


Function orientation

Null $\approx \phi$

Node \Rightarrow Linked List

```
def insertion(self, data):
    newNode = Node(data)
    newNode.next = self.head
    self.head = newNode
```



```
def insertion(self, data):
    newNode = Node(data)
    if not self.head:
        self.head = newNode
    else:
        current = self.head
        while current.next:
```



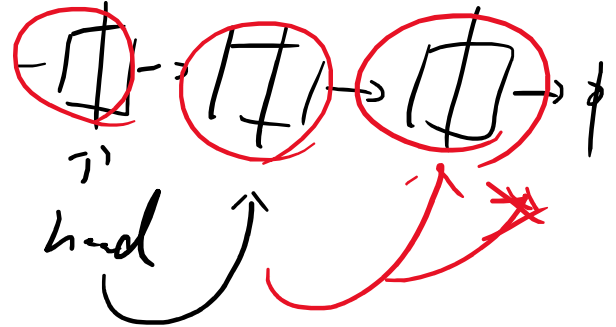
```

current = self.head
while current.next:
    current = current.next
current.next = newNode

```

$head \Rightarrow \underline{current} \rightarrow next$

$\rightarrow fn(head)$
 $\rightarrow fn(head)$
 $\rightarrow fn(head) \times$



```

def deletion (self, element):

```

```

    if not self.head:
        return

```

```

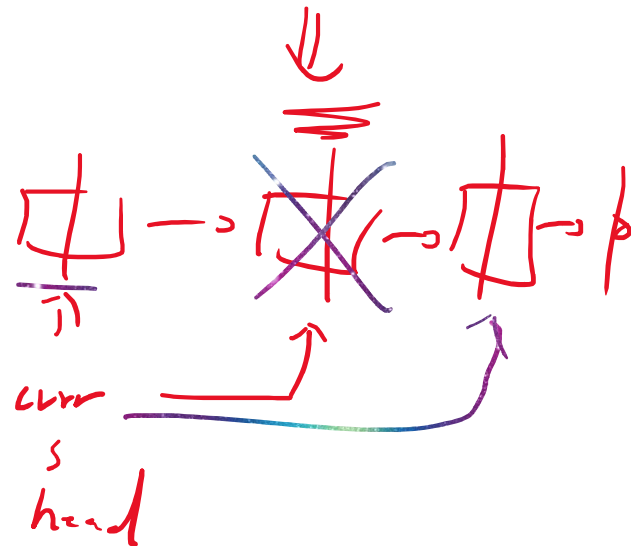
    if self.head == element:
        self.head = self.head.next
        return

```

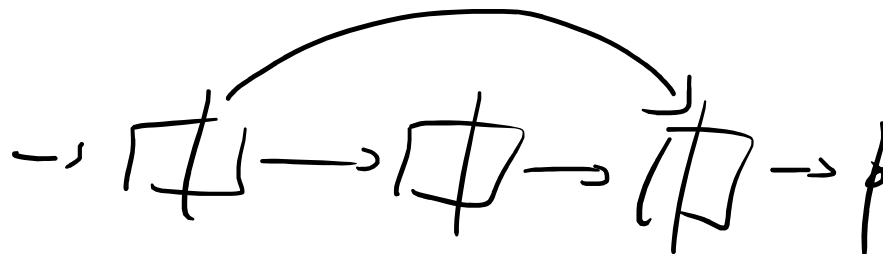
```

    current = self.head
    while current.next:
        if current.next.data == element:
            current.next = current.next.next
            return
        current = current.next

```



\Rightarrow Detect a cycle in LL :



Hashmap based

Floyd Cycle
(Hare & Tortoise Alg)

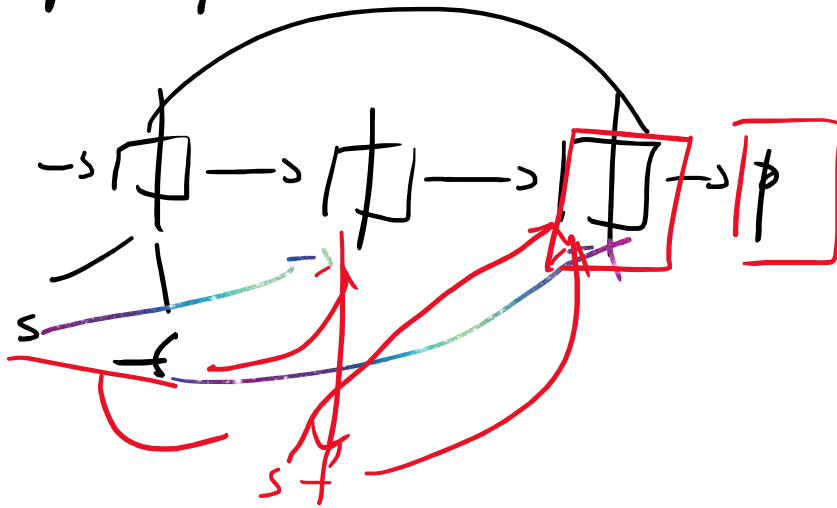
=> Hashmap :

- > Traverse the list individually & keep putting node addresses in a hashtable.
- > At any point if NULL is reached return False.
- > If the next of current points to any prev, nodes in hashtable return true.
- > Time Complexity $O(n)$
- > Auxiliary Space $O(n)$

=> Floyd Cycle \approx Hare & Tortoise Algorithm.

- > Traverse LL using 2 pointers.
- > Move one pointer slow & another by 2.

- Move one pointer slow 1 and then by 2.
- If 2 pointers meet there is a cycle.



⇒ Reverse Linked List :

↳ Start through 3 pointers : prev, current & next

↳ Start through LL :

↳ Before changing next of curr, store it

→ next = curr → next

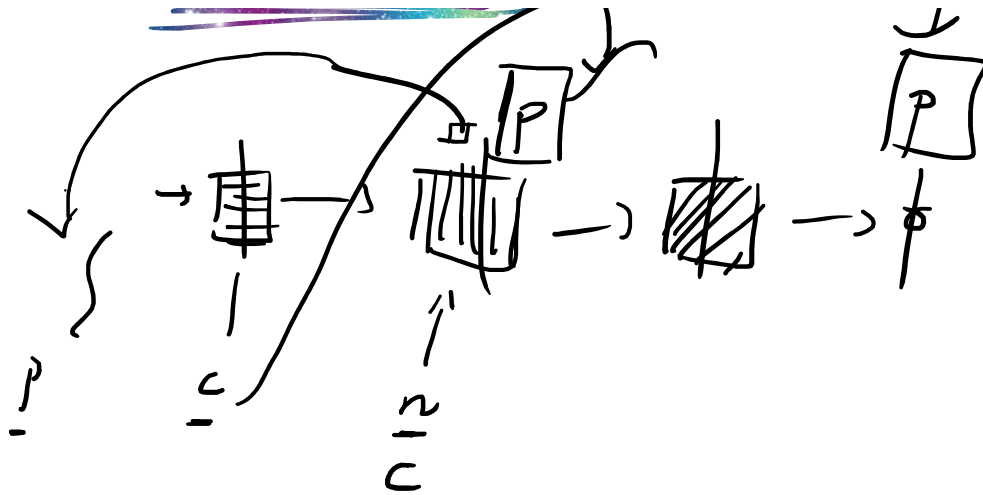
↳ Now update the next pointer of curr to prev

→ curr → next = prev

↳ Update prev as curr & curr as next

→ prev = curr
→ curr = next





Time Complexity $O(N)$
 Auxiliary Space $O(1)$