

Sorting: arrangement of data in some order on the basis of some parameters

2 3 9 12 17 19 Sorted in asc order (magnitude)

19 6 5 2 -1 -19 sorted in desc order (magnitude)

	1	13	9	6	12	sorted in asc
count	1	2	3	4	6	order (count of
of factors						factors)

(Advanced)
X How to sort?

Why sorting? searching becomes easier

library / inbuilt function \rightarrow sort()

\downarrow
TC : $O(n \log_2 n)$

• $O(n^2) / O(n^3)$

\downarrow
 $O(n \log n) / n^2 \log n \rightarrow$ Always try to think about sorting

• $O(n)$ X Sorting

\downarrow
 $O(\log n)$

1. Given an array of N integers, we've to delete all elements of the array. Before deleting an element, pay cost = sum of elements in the array (at that point). Find min cost. Distinct elements

Ex $[~~2~~, ~~1~~, ~~4~~]$

delete 1	$2 + 1 + 4 = 7$
delete 2	$2 + 4 = 6$
delete 4	$4 = 4$
	<u>17</u>

Total cost = 17

delete 4	$2 + 1 + 4 = 7$
delete 2	$2 + 1 = 3$
delete 1	$1 = 1$
	<u>11</u>

min cost \nearrow

Ex $[~~4~~, ~~6~~, 1]$

delete 6	$4 + 6 + 1 = 11$
delete 4	$4 + 1 = 5$
delete 1	$1 = 1$
	<u>17</u>

Ex $[~~3~~, ~~5~~, ~~1~~, -3]$

delete 5	$3 + 5 + 1 + (-3) = 6$
delete 3	$3 + 1 + (-3) = 1$
delete 1	$1 + (-3) = -2$
delete -3	$-3 = -3$
	<u>2</u>

~~[a b c d]~~

delete a
delete b
delete c
delete d

$a+b+c+d$
 $b+c+d$
 $c+d$
 d

$a+2b+3c+4d$

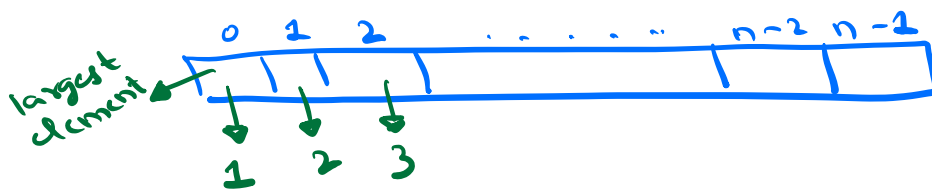
largest element

min. cost

$a > b > c > d$

Start removing from largest end

• sort(arr) $\rightarrow T_c: O(n \log n)$



$i \rightarrow (i+1)$ times

- 1) sort data in desc order
- 2) cost = 0

$\rightarrow n \log n$

```
for (i = 0 ; i < n ; i++) {
    cost += a[i] * (i+1);
}
```

} n

$T_c: O(n \log_2 n)$
 $SC: O(n)$

sort() $\rightarrow SC: O(n)$
 \downarrow
depend on sorting algo

2. In an array of N element, find count of noble integers.

$A[i]$ is noble if

count of elements $< A[i] = A[i]$

Distinct nos

	0	1	2	3	4	5	
Ex	1	-5	3	5	-10	4	ans = 3
count	2	1	3	5	0	4	

	0	1	2	3	
Ex	-3	0	2	5	ans = 1
	0	1	2	3	

• A -ve element can't be noble.

count = $a[i]$
 \downarrow
 0 or +ve \neq -ve

Idea 1: For every element, count smaller elements

ans = 0

for ($i=0$; $i < n$; $i++$) {

count = 0

for ($j=0$; $j < n$; $j++$) {

if ($a[j] < a[i]$)

count++

}

if (count == $a[i]$)

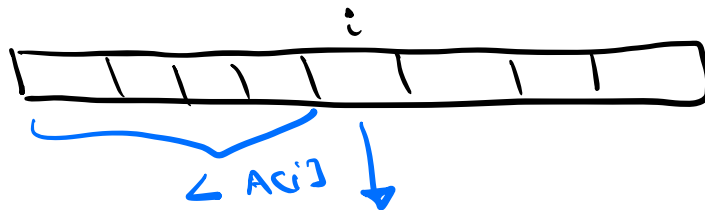
ans++

}

TC: $O(n^2)$

SC: $O(1)$

Idea 2 sort data (asc)

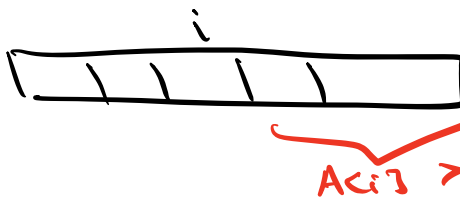


[0 i-1]
 $i-1-0+1$

```
// sort data
ans = 0
for (i=0 ; i < n ; i++) {
    if (a[i] == i)
        ans++
}
```

count of smaller elements

desc



$N-i-1$
 $[i+1 \quad N-1]$
 $(N-1) - (i+1) + 1$
 $= N-1-i-X+X$

what if there are duplicates?

	0	1	2	3	4	
Ex	-10	1	1	3	100	
count	0	1	1	3	4	

ans = 3

	0	1	2	3	4	5	6	7	8	
Ex	-10	1	1	2	4	4	4	8	10	
count	0	1	1	3	4	4	4	7	8	

ans = 5

	0	1	2	3	4	5	6	7	8	9	10	11	12	13
Ex	-3	0	2	2	5	5	5	5	8	8	10	10	10	14
count	0	1	2	2	4	4	4	4	8	8	10	10	10	13

ans = 7

if cur - elem != prev - ele
count = index

1. sort the data $\rightarrow n \log n$

2. ans = 0

cnt = 0

if (a[0] == 0)

ans ++

cnt of smaller elements

TC: $O(n \log n)$

for (i = 1; i < n; i++)

if (a[i] != a[i-1])

cnt = i;

if (a[i] == cnt)

ans ++

$\rightarrow n$

7

7	8	9
5	8	8
4	8	8

	0	1	2	3	4	5	6	7	8
-3	0	2	2	5	5	7	8	8	
cnt = 0	1	2	2	4	4	6	7	7	
ans = 0	0	1	2	2	2	2	2	2	(2)

3. Sort data in asc order by **count of factors**

if **count is equal**, sort in asc order based on **magnitude**

	9	3	10	6	4	⇒	3	4	9	6	10
cnt	3	2	4	4	3		2	3	3	4	4

sort () → by default, **asc / desc order**
by **magnitude**

sort (— , — , —)

↓ ↓

start index end index

↑ greater <int>

comparator fn
↓
rules of your sorting

★ Custom sorting

x, y

$x < y$ (asc)

↓
 x should come first

$x, y \rightarrow 2$ nos. are given

if (count-f(x) < count-f(y))
x should come first

count-f(x) > count-f(y)
y should come first

count-f(x) == count-f(y)
 $x \leq y$

x should come 1st

$x > y$
y should come 1st

x	y
5	5
cnt 2	2

int /

bool

comp(int x, int y) \rightarrow data type
of element
that you
sort

if first
argument (x)
should come
first in
sorted data
 \downarrow
return true
else
return false

int cntx = count-factors(x);
int cnty = count-factors(y);

if (cntx < cnty)
return true;

else if (cnty < cntx)
return false;

else <

if (x <= y)
return true

else
return false;

for every comparison
 \downarrow
TC \rightarrow comparator
fn

>

\sqrt{m}
 \downarrow
max
ele

Total time = no. of comparisons \times time in comparator

- sort in desc order of magnitude

```
bool comp (int x, int y) <
{
    if (x < y)
        return false;
    else
        return true;
}
```

x y
2 < 5
5 comes first
return false

x y
5 > 2
return true

- array of strings
sort acc. to length of strings in asc order

```
bool comp (string x, string y) <
{
    if (x.length < y.length)
        return true;
    else
        return false;
}
```

x y
cat camel
3 5
x to come 1st
return true

x y
monkey dog
6 3
return false

dog cat
lexicographical
↓
char vs char

Doubts

5 ← [2 5 10 20 30]

↓

n
↓
size of
array

count - f(30)
↓
for (i = 1 ; i ≤ √30 ; i++)

count - f(m) → [1 √m]
↓
O(√m) → max element