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# Assignment 2 - Peak Detector

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# 1 Introduction

## 1.1 Group Composition:

- 1) Command Processor
  - *Will Myers*
  - *Oscar Fuller*
  - *Oscar Munday*
- 2) Data Processor
  - *Reuben Bartley*
  - *Frank Gaze*

## 1.2 Individual contributions:

### **Oscar Fuller -**

In this project I was responsible for building the shifter required to store the three numbers that follow the initial a/A and send them to the data processor which correspond the number of bytes needed to be processed. Secondly, I was responsible for receiving the eight-bit data word sent from the data processor in hexadecimal and transmitting it in ascii to the PuTTY Console.

### **Will Myers -**

My contribution to this project was to find an efficient way to format the PuTTY console and detect the end of the byte sequence from the data processor. This required considerable work on the top-level state machine, a look up table to output the different format sequences (multiplexer when synthesized) and a counter for the 'seqDone' signal.

### **Oscar Munday -**

### **Reuben Bartley -**

### **Frank Gaze -**

In this peak detection project, I was responsible for part of the development of the Data Processor. I was responsible for the Data Processors communications with the Data Generator. This involved integrating the required Two-Phase protocol and working with the 'numwords' and the 'SeqDone' signals.

# 2 Project Plan

## 2.1 Group Meetings

We plan to meet every Monday, Tuesday, and Wednesday during our designated lab times. This will provide an opportunity to discuss our project as a group, while also having the opportunity to discuss our design decisions with staff and to receive valuable feedback. Furthermore, starting on week 18 the separate groups will have to organise a weekend session to continue work on their sections. This will aid in meeting our milestones and objectives. Also, it will mean that any problems encountered can be passed on to staff on Monday which will aid with our project's progression. These weekend sessions will become whole group meetings from week 21 to aid with integration of the processors and to allow for extra assistance on the command processor if required.

## 2.2 COVID-19 Compromise

Due to the ongoing Covid-19 outbreak, we had to find remote methods to continue to work effectively as a team. This had to allow effective group collaboration, while also avoiding hindering the projects development. This was achieved through weekly meetings held on Zoom (A free video calling platform). Every Sunday we met as a group, which provided an opportunity to discuss our progress on meeting our deadlines and to receive feedback from other members. At the same time, we also organized an additional weekly meeting between our separate groups. This proved to be very effective, as it permitted members to discuss problems in great depths of detail. The productivity resulting from these meetings resulted greatly due to our groups use of GitHub as a safe central storage for our code. This allowed members to review specific sections of code while others discussed the problems they were encountering. This greatly aided in the progression and development of our project.

## 2.3 Milestones & Objectives

TASKS	DURATION (WEEKS)	MEMBERS	Week Timetable No.									
			16	17	18	19	20	21	22	23		
<b>Preparation Phase</b>												
Understanding Vivado	1	All										
Understand System Simulation & Synthesis	2	All										
<b>Data Processor</b>												
Make ASM Chart	1	Frank & Reuben										
Create Data Retrieval	3	Frank, G										
Create Peak Detection	3	Reuben, B										
Combine sections	1	Frank & Reuben										
<b>Command Processor</b>												
Make ASM Chart	1	A, D										
Data Echoing (ANNN)	2	Oscar, M										
Communication with Data Processor	3	Oscar, F										
Formatting Putty Terminal	4	Will, M										
Outputting Byte	3	Oscar, F										
Checking SeqDone	3	Will, M & Oscar, M										
<b>Report</b>												
Data Processor	3	Frank & Reuben										
Command Processor	3	Will Meyers, Oscar, F and Oscar, M										
Integration	2	All										

Figure 1: Gantt Chart

We decided to designate the first week to allowing ourselves to get comfortable with Vivado. At the same time, we also decided we would all start to investigate the full system simulation and synthesis. This is because it would help us later in the integration phase when the two processors are combined. After the second week, we will split into our separate groups and begin work on our ASM charts. Next, the data processor team will have until the end of week 20 to finish their code. This will then allow them to begin work on the report while also aiding with the design of the command processor if required. The objective is to finish the control processor by the end of week 21. However, if this milestone is not completed, then we are prepared to continue working on it in week 22. This will then leave us with a final week to finish the group report and complete the integration of the two processors to complete the full system. These objectives and milestones can be observed in more detail in the Gantt chart below.

## 3 Command Processor

### 3.1 Echoing ANNN into PuTTY Console - Oscar Munday

#### Top Level State Transitions: INIT – CORRECT\_WORD

This section dictates the initialisation of the command processor and how it waits for a correct sequence of starting inputs while echoing all inputs back to the PuTTY terminal.

In the first state INIT, all value-holders are set to default starting values before moving to the next state, RXNOW\_WAIT.

The command processor waits for the rxNow line to go high for one clock cycle – signaling an input byte is ready to be received from the Rx module – before moving to the next state.

In COUNT\_CHECK the input byte is registered. A counter allows the processor to track how much of a valid input sequence has been received so far. If the counter is on 0, the processor is checking for the first byte of the input sequence and so the next state is A\_CHECK. If the counter is greater than 0, the processor must be checking for a number and so the next state is NUM\_CHECK.

In the state A\_CHECK, the byte is checked to be the start command: ASCII-code for ‘a’ or ‘A’. If it is not, then the next state will be ERROR; the processor moves back to INIT where it is reset, and the above process repeated. If it is, then the processor moves to CORRECT\_WORD. Additionally, the txnow line is set high for 1 clock cycle signaling to the Tx module that a byte has been registered and is ready to be transmitted. In turn, the module sends it to the computer, thus the input data is echoed in the PuTTY terminal.

In NUM\_CHECK, the input byte is checked to be an ASCII-coded number from 0-9. Similarly to when in A\_CHECK, if it is not then the next state will be ERROR and the processor is reset, and if it is then the next state is CORRECT\_WORD. As before, txnow is set high to echo the registered byte back to the PuTTY terminal.

In CORRECT\_WORD the counter is incremented by 1 which indicates another valid byte of the sequence has been received. The rxdone line is set high for a clock cycle signaling that the byte has been successfully read from the Rx module and the processor is ready for the next byte. The counter is checked and if it is 0, the processor returns to RXNOW\_WAIT. If not, the current byte is a number and so the processor moves to the state SHIFTER. After which it returns to RXNOW\_WAIT unless the sequence is finished. Functionality of the SHIFTER state is explained in the next section.

#### **Components required: Output Register**

The output register stores a byte ready to be sent to the Tx module, and therefore to the PuTTY terminal. D is assigned to a value holder and follows its value for the duration of a clock cycle. At the same time, load is asserted high to signal Q’s reevaluation. In the next clock cycle D holds a constant value and Q is assigned to replicate D’s value. In the clock cycle after, Q holds the assigned value and maintains it indefinitely which allows for stable data storage in Q. The byte stored in Q is therefore registered and ready for transmission to the Tx module.

If regreset is set high, Q is assigned to its default value ‘11111111’. This is asserted during initialisation to give Q a defined initial value.

#### **Components required: Counter**

This counter is used to track how far through a valid input sequence ‘aNNN’/ ‘ANNN’ the command processor has received.

### **3.2 NNN TO Numwords - Oscar Fuller**

#### **Top Level State Transitions: CORRECT\_WORD – WAIT\_SHIFT**

In this section a shifter is used to store the numbers coming from the rxData that will become numWords\_BCD.

When in the state CORRECT\_WORD, if the counter is on 1, 2 or 3, the data coming in is a number required in numWords\_BCD so the next state will be SHIFTER.

In the SHIFTER state the shifter is enabled and the number will be loaded into the shifter. This state also checks the counter again to see if it has a value of greater than or equal to 4, this signifies the shifter having all three numbers required. If this is true it will proceed to the WAIT\_SHIFT state, if not, the command processor will go back to the RXNOW\_WAIT state.

When in the WAIT\_SHIFT state load1 goes high to output the numbers stored in the shifter in a BCD array in the signal numWords\_BCD to the data processor. This state also waits for txdone to go high to signify the transmitter has finished outputting the final N from NNN and the command processor can proceed to the FORMAT1 state. Also, the counter is reset.

#### **Components required: Shifter**

In order to store the numbers sent by Rx so that they could be converted to BCD and sent through numwords, we decided to use a shift register. This shifter will take in three ascii numbers from rxData (after the initial A/a) and then transmit all three numbers in binary coded decimal to the data processor.

The shifter will be made up of three separate registers each corresponding to the significance of the number. Once the shifter is enabled it will store the four least significant bits of the eight-bit ascii character provided into the first register, this is because for all the numbers in ascii only the four least significant bits show the value of the number.

When the shifter is enabled a second time the four bits will be shifted up into the next more significant register and a new four bits will be saved. This carries on until all three numbers are saved.

When load\_shift goes high the shifter will output the three numbers in an array of three binary coded decimal numbers.

### **3.3 Formatting PuTTY Console - Will Myers**

After echoing ANNN from Rx back to the PuTTY console, we decided that we should display the retrieved bytes on a new line and have spaces between each byte transmitted. Also, because there is a maximum of 999 bytes that could be outputted, we included functionality to start a new line in PuTTY every 20 bytes to aid in the readability of the displayed data.

#### **Top Level State Transitions: FORMAT1– SEQ\_CHECK**

The sequence of symbols outputted after the ANNN is:

“SPACE, =, =, SPACE, \, r, \, n, SPACE”. As previously mentioned, the WAIT\_SHIFT state resets the counter so that we can use the counter to cycle through the look up table.

State FORMAT1 acts as a buffer state to allow the clock driven format look up table time to load the correct value of “000000” at its input and output the correct symbol in FORMAT2. Enable is also set high in FORMAT1 to allow the FSM to iterate through the table.

The next state is FORMAT2 where the output of the format multiplexer is loaded into the output register. This is then sent out of the transmitter modules in FORMAT3.

After this, the command processor waits for the transmitter to finish sending the symbol by waiting until txdone is set high in WAIT\_FORMAT.

The next state is FORMAT\_CHECK where it loops back to FORMAT1 to output a new symbol, unless cnt0Out is greater than or equal to 8, at which point the final SPACE has been outputted and the controller moves on to start outputting bytes. Please note the remaining decision blocks in the ASM chart during this state will be explained in the last section.

As explained in the following section, the FSM loops back to FORMAT1 after the second digit from byte has been outputted in TRANSMIT2. This causes it to go through the same process, increasing the count and outputting a space between the byte. This is because when the format table receives a number greater than 8 at the address, it outputs a space.

In the next state SEQ\_CHECK, the counter is checked again and if it equals 29, then 20 bytes would’ve been outputted and the processor moves to NEW\_LINE (when first byte is outputted the count is 9, and it is incremented with each space between bytes).

In NEW\_LINE, reset\_to\_4 is set high which sets the counter to 4. This results in the sequence “\, r, \, n” being outputted in PuTTY, starting a new line to output the next 20 bits. If the count doesn’t equal 29 then the next state equals START1 and a new byte from the data processor will be received. Note the remaining decision block in the ASM chart during this state will be explained in the last section of the command processor.

#### **Components required: FomatMUX**

This look up table is used to output various sequences of symbols to format the PuTTY terminal by connecting the address of the table to the counters output. Then, for different counts, the table outputs a different

symbol in ascii code.

In the combinational process of the multiplexer, the output is an internal signal q1 instead of formatout. This is then assigned to formatout in the sequential process to ensure that the component is clock driven.

### 3.4 Outputting Byte - Oscar Fuller

#### Top Level State Transitions: START1 – TRANSMIT2

In this section the command processor communicates with the data processor and displays the bytes retrieved from the data generator. Once the command processor is in the START1 state it sets the start signal high for a clock cycle, causing the data processor to start retrieving data. Once start is enabled the next state is DATA\_WAIT, this is a wait state that waits for the signal dataready from the data processor to go high, indicating that it has valid data to be read on the byte line.

The next state is WAIT\_BYTE. In this state the four most significant bits of the byte signal from the data processor are loaded into the ASCII look up table in the byte multiplexer. This is because the first four bits of byte correspond to the first hexadecimal digit which is then converted into ascii using the table. In the next state LOAD\_BYTE, the output of the byte multiplexer is connected to the input of the output register. The registers enable is simultaneously set to high, loading the number across to the output of the register and thus to txData.

In TRANSMIT1 txnow goes high for one clock cycle enabling the transmitter to transmit the ascii coded, hexadecimal digit to the computer.

In the next state called WAIT\_TX, the state machine waits for txDone to go high. Also, the remaining four least significant bits of the byte signal are loaded into the byte multiplexer to be converted into ascii code. Once txDone goes high the command processor goes into the LOAD\_BYTE2 state, which has the same function as LOAD\_BYTE and then to the final wait state WAIT\_TX2, which is the same as WAIT\_TX, before looping back to FORMAT1.

#### Components required: ByteMUX

The byteMux is a multiplexer which converts a four-bit hexadecimal number into ascii code. It functions with each four-bit hexadecimal input having a corresponding eight-bit ascii character. When a four-bit number is inputted for example 1010 which corresponds to an A in hexadecimal the multiplexer will output 01000001, an A in ascii.

Similarly to the functionality of the format multiplexer, in the combinational process of the byte multiplexer, the output is an internal signal q1 is used instead of asciiout. This is then assigned to asciiout in the sequential process.

### 3.5 Checking for SeqDone– Will Myers

As the data processor sends its last byte to the command processor, it sends the seqDone signal high for just one clock cycle, right after sending dataReady. This therefore requires checking for seqDone and ‘remembering’ that it has gone high after the last byte has been sent. This ‘remembering’ will be achieved by using another counter.

#### Top Level States

While in the LOAD\_BYTE the processor also checks if seqDone has gone high in which case the next state will be LAST\_BYTE. In LAST\_BYTE, enSeqDone is set high which causes the seqDone counter to count once.

The processor then outputs the final byte and then returns to FORMAT1 before outputting the final SPACE. It then moves to state SEQ\_CHECK, where SEQcntOut is checked and if it equals “01”, then the next state is FINAL\_FORMAT, where the counter is reset to 4 and the seqDone counter counts again to “10”. After FINAL\_FORMAT, the processor moves to FORMAT1, which begins the final formatting loop to output “\,

r, \, n”, starting a new line where new commands may be inputted to the console. The FSM finally loops back to the INIT state when the seqDone counter is at “10” and when the main counter is at 8 to signify that the formatting sequence has been completed.

### Components required: Seqdone Counter

The architecture of the seqDone counter is the same as that of the main counter except that is only made of two registers as it is only enabled twice.

The byteMux functions when the signal loadDATA goes high, this signifies the data being loaded into the multiplexer and then being outputted on the next clock cycle on the q signal.

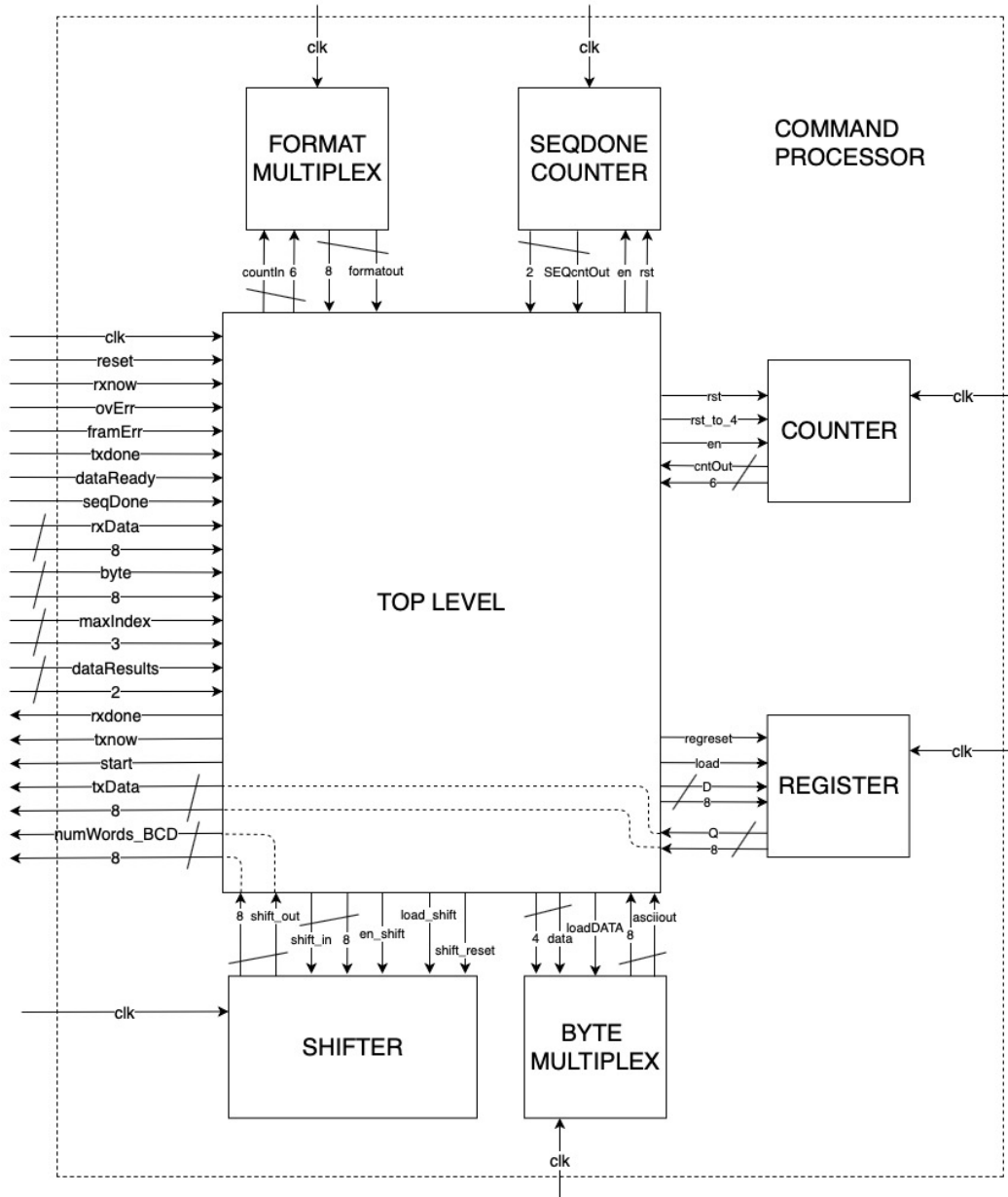


Figure 2: Block level diagram of command processor



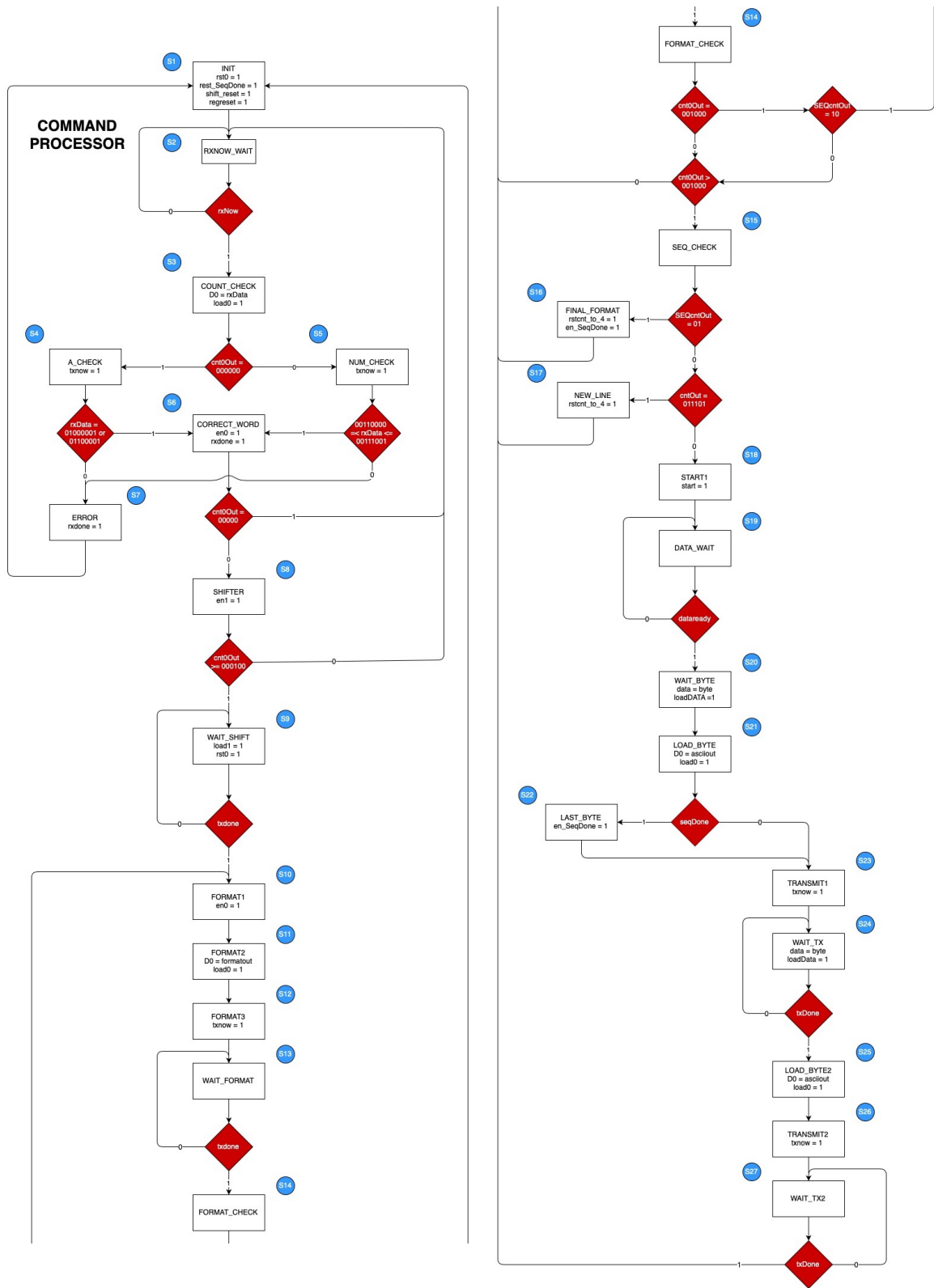


Figure 3: ASM for command processor

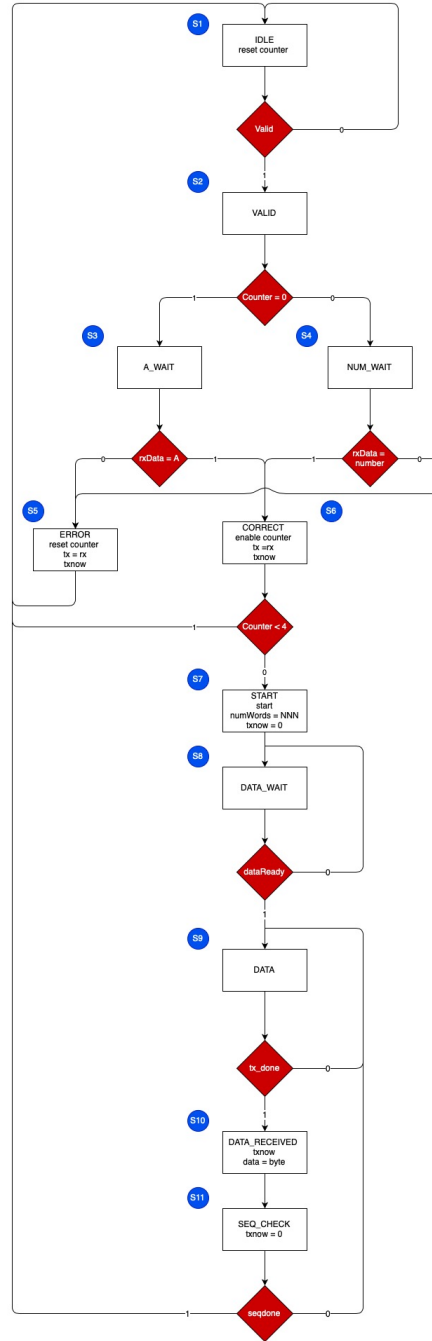


Figure 4: shortened ASM of command processor

## 4 Data Processor

### DIGITAL DESIGN, GROUP PROJECT

ASSIGNMENT 2, DATA PROCESSOR  
GROUP 19

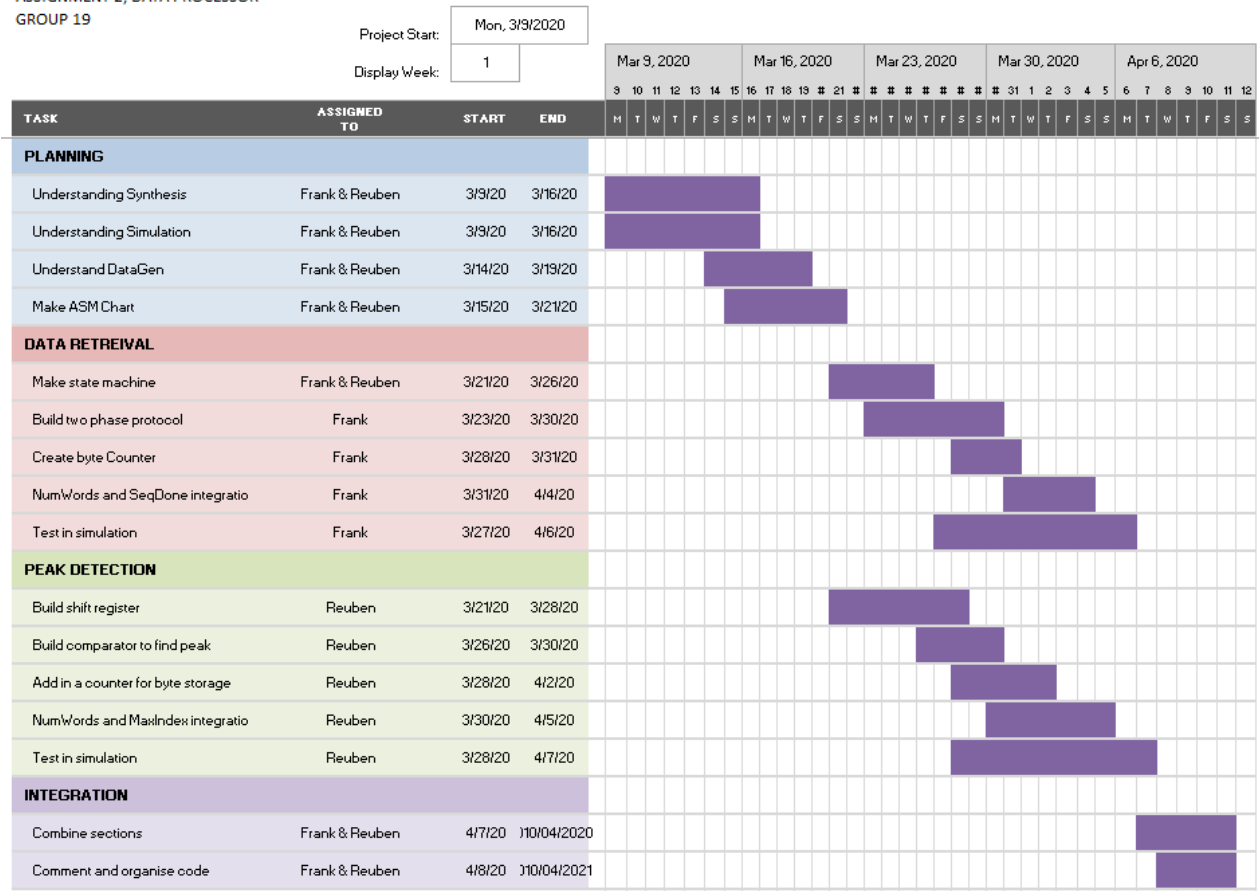


Figure 5: Gantt chart for data processor

### 4.1 Data Retrieval – Frank Gaze

Prior to starting our code, we decided to split the Data Processor into two submodules: Data Retrieval and Peak Detection. Two ASM charts were constructed to aid with the coding process and provided a more visual and practical representation of what we wished to implement. The ASM chart for the Data Retrieval submodule is shown below.

#### Data Gen

To gain information on successfully implementing the two-phase protocol, I devoted time to fully understanding the Data Generator and its single-phase protocol. This provided crucial insight into how to approach changing CtrlOut and also detecting a change in CtrlIn. Especially with the implementation of NOT and XOR gates used in flip flops.

[h!]

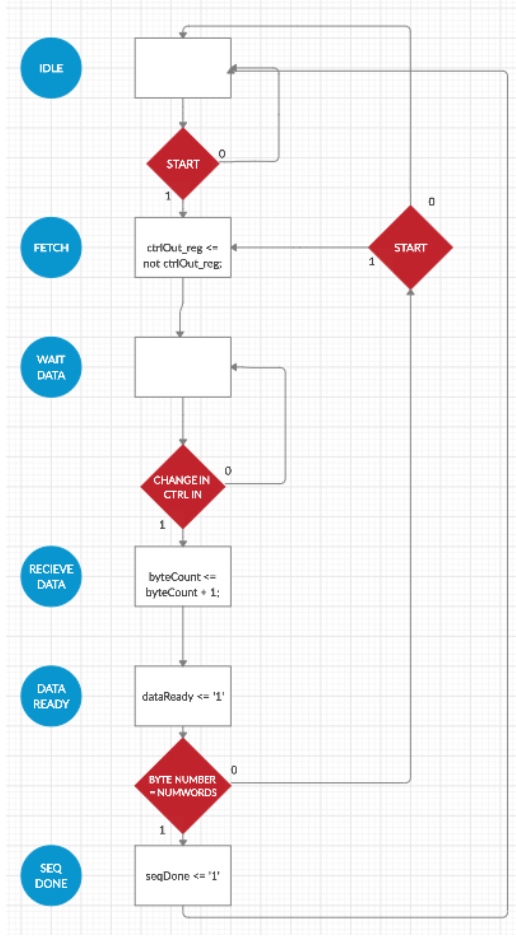


Figure 6: ASM for data retrieval

### Byte Counter

A byte counter was important to the implementation of the Data Processor. This is because it is essential in determining when the Peak Detector had reached the end of its given sequence. The byte counter was implemented in a process which was edge triggered and its output varied depending on which state it was currently in. For example, the byte counter would only increase if it entered the process while in the 'RECEIVE DATA' state, as that meant it had just passed through the two phase protocol and had received valid data from the Data Generator. Else, it would either be reset or remain unchanged. Once the machine entered the 'DATA READY' state, the 'dataReady' signal was set high and the data was passed onto the byte lines for the command processor. After this, the system used the byte counters value and a comparator to check if the sequence was complete or not.

### Sequence Complete Comparator

A comparator was implemented to detect whether the Peak Detector had finished the given sequence. This was accomplished by comparing the value of the byte counter to that supplied by the 'numwords' signal. This was implemented in an asynchronous process which was triggered by a change in the byte counter. However, to successfully compare the two signals, 'numwords' had to be converted from BCD (Binary Coded Decimal) to an integer. This was performed in a separate asynchronous process which was executed whenever the 'numwords' signal was updated. If the byte counters value was equal to that of which 'numwords'

### Two-Phase Protocol

To request a byte from the data generator, a change in CtrlOut was required. This was implemented by using a flip-flop with a NOT gate which affected the CtrlOut signal whenever the system entered the 'FETCH' state. This would only occur after the 'start' signal from the command processor was high on the rising edge of a clock cycle. Else, it would remain in the 'IDLE' state.

Once CtrlOut was changed the system would enter the 'WAIT DATA' state, where it would wait for a change in the CtrlIn signal from the Data Generator. To detect a change in CtrlIn a process was made that assigned the variable of CtrlIn to a new signal on the rising edge of every clock cycle. This new signal and CtrlIn was then used in an XOR gate. If this returned as high, then CtrlIn had changed and meant the data on the data lines was valid and the state machine could enter the 'RECEIVE DATA' state, where the byte counter would increase.

specified, then the machine would enter the 'SEQ DONE' state where the 'SeqDone' signal is set high and 'dataResults' and 'maxIndex' would contain the correct data. However, if the sequence was not complete then the machine would either enter the 'FETCH' or 'IDLE' state depending on the value of the 'start' signal.

## 4.2 Peak Detector – Reuben Bartley

### Peak Detection

A comparator is used to detect whether the most recent byte that will be input into the shift register is greater than the Byte that is currently in middle position of the larger Data Register. If this is the case then the 'peakDetect' signal will go high. A shift register stores the value of the most recent byte and will shift this to the right each time a new byte of data is retrieved for the data generator. Therefore, when 'peakDetect' is high this triggers a process that loads all 4 bytes into the first 4 positions of the dataResults\_Reg a register, that is a 7 CHAR Array designed to store the value of the Peak byte in the 3rd CHAR with the 3 bytes that precede and proceed the Peak stored in the first 3 CHAR's and the last 3 CHAR respectively.

### Storing the Peak Values Index

The numWords value is given to the data Processor in BCD and therefore needs to be converted into an integer so that it can be used by the comparator. This is discussed in the previous section. However, the peak value index is required by the command processor in BCD format therefore, an integer to BCD converter was needed. In order to do this each of the different powers needs to be separated from each other and then converted into binary separately. This was done by making use of the modulus and division operators available in the numeric package. Once these were isolated they could be converted to unsigned binary and stored in the corresponding position of the BCD array of Max Index\_reg, a register that stores the most recent peak Value index before outputting it back to the Command Processor at the end of the sequence.

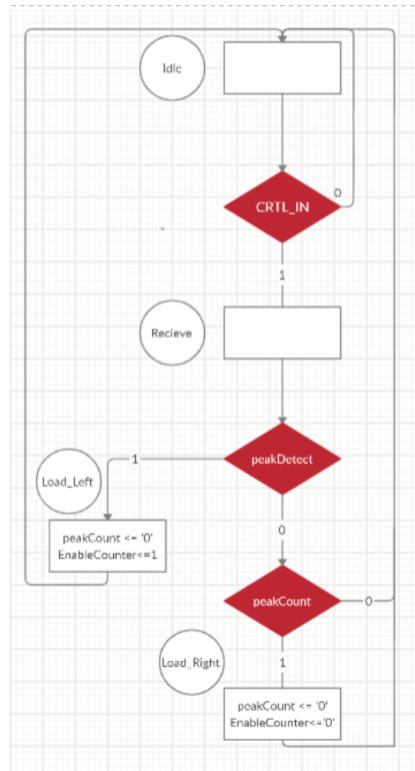


Figure 7: ASM for peak detector