HBL551 (Numerical Analysis):

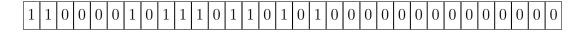
Homework 1

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Problem

1. Show that the IEEE Standard 754 single precision floating-point representation



is equal to $(-118.625)_{10}^*$

2. Convert $(-118.625)_{10}$ into IEEE Standard 754 single precision floating-point representation.

Solution

Problem 1:

a. This will be the notation that I will be using throughout my solution:

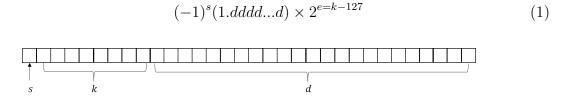


Figure 1: IEEE 754 single precision floating-point representation

where $s \equiv \text{sign}$, $e \equiv \text{true}$ exponent, $k \equiv biased$ exponent, and $d \equiv \text{significand}$.

b. From the IEEE Standard 754 single precision floating-point representation

 $s = (1), k = (10000101)_2, \text{ and } d = (11011010100000000000000)_2^{**}.$

^{*}I used the subscript 10 to denote the decimal system.

^{**}I used the subscript 2 to denote the binary system

c. The general formula that I will be using is

$$(-1)^{s} [1 + (d)_{10}] \times 2^{e=(k)_{10}-127}.$$
 (2)

Note: Since there are 8 bins allocated for the exponent, the range of values of the exponent is $[0, 2^8 - 1] = [0, 255]$. To extend the exponent's range to negative numbers, the first bin in the exponent is allocated for the sign. Note, however, that it will be redundant to denote 10000000 as -0, so it is assigned instead to 128. Hence, the exponent's range becomes $[-(2^7 - 1), (2^7 - 1) + 1] = [-127, 128]$. But the exponent, k, that is stored in the computer is the sum of the true exponent and the number $(2^7 - 1) = 127$. Therefore, if k is the value of the biased exponent that is stored in the computer, the number 127 must be deducted from k to get the true exponent, e; i.e., e = k - 127.

d. Converting k and d into decimal system gives:

$$(k)_{10} = 2^7 + 2^2 + 2^0 = 133,$$

 $(d)_{10} = 2^{-1} + 2^{-2} + 2^{-4} + 2^{-5} + 2^{-7} + 2^{-9} = 0.853515625.$

Therefore, the IEEE Standard 754 single precision floating-point representation

is equal to $(-1)^1(1+0.853515625) \times 2^{133-127} = (-118.625)_{10}$

Problem 2:

- a. In converting $(-118.625)_{10}$ into IEEE Standard 754 single precision floating-point representation, I will break my solution into two parts. The first part involves the conversion of the whole number $(118)_{10}$ into binary system and the second part involves the conversion of the decimal number $(0.625)_{10}$ into binary system.
- b. First part:

A positive base-10 integer I can be expressed as an expansion of powers of 2 as:

$$I = (b_n \times 2^n) + (b_{n-1} \times 2^{n-1}) + \dots + (b_1 \times 2^1) + (b_0 \times 2^0).$$
 (3)

Therefore, if I = 118, then the process of converting I into binary system involves:

$$\begin{array}{llll} 118 = 59 \times 2 + 0 & \to & b_0 = 0, \\ 59 = 29 \times 2 + 1 & \to & b_1 = 1, \\ 29 = 14 \times 2 + 1 & \to & b_2 = 1, \\ 14 = 7 \times 2 + 0 & \to & b_3 = 0, \\ 7 = 3 \times 2 + 1 & \to & b_4 = 1, \\ 3 = 1 \times 2 + 1 & \to & b_5 = 1, \\ 1 = 0 \times 2 + 1 & \to & b_6 = 1. \end{array}$$

[†]The code that I made for converting IEEE standard 754 *single precision* floating-point representation to decimal can be found in page 4..

c. Second part

A fractional base-10 number R can be expressed as an expansion of powers of 2^{-1} as:

$$R = (b_1 \times 2^{-1}) + (b_2 \times 2^{-2}) + \dots + (b_n \times 2^{-n}) + \dots$$
 (5)

such that

$$R = 0.b_1 b_2 ... b_n ... (6)$$

Therefore, if R = 0.625, then the process of converting R into binary system involves:

$$0.625 - 2^{-1} = 0.125 \ge 0 \quad \to \quad b_1 = 1,$$

$$0.125 - 2^{-2} = -0.125 < 0 \quad \to \quad b_2 = 0,$$

$$0.125 - 2^{-3} = 0 \ge 0 \quad \to \quad b_3 = 1.$$

$$(7)$$

d. To convert $(-118.625)_{10}$ into IEEE 754 single precision floating-point representation, the sum of eqs. (4) and (7) is multiplied by $(-1)^1$:

$$(-1)^{1}(118.625)_{10} = (-1)^{1}((118)_{10} + (0.625)_{10}),$$

$$= (-1)^{1}((11101101)_{2} + (101)_{2}),$$

$$= (-1)^{1}((1110110.101)_{2}),$$

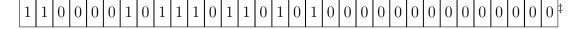
$$= (-1)^{1}(1.110110101 \times 2^{6}),$$

$$= (-1)^{1}(1 + 110110101) \times 2^{6}.$$
(8)

Comparing the result in eq. (8) with eq. (2), it is easy to see that s = 1, $d = (110110101)_2$, $e = (6)_{10}$, and $k = e + 127 = (133)_{10}$. The final step is to convert k into binary system, which follows the same steps done in part b:

$$\therefore k = (133)_{10} = (10000101)_2 \tag{9}$$

e. Therefore, using the notation presented in Figure (1), it is easy to see that the IEEE 754 single precision floating-point representation of $(118.625)_{10}$ is given by



[‡]The code that I made for converting decimal to IEEE 754 *single precision* floating-point representation can be found in page 6.

Codes

Problem 1

```
function dec = SingPrec2Dec(bin)
format long
bin = bin(~isspace(bin)); % remove spaces
bits = bin(:) - '0';
len = size(bits);
% This is done because zeros (1,32) won't give an accurate ans
zer = zero(:) - `0';
if len(1) == 32
   if bits = zer \% this guarantees [`000...0']_{-2} = (0)_{-10}
      dec = 0:
   else
      s = bits(1);
      k = bits(2:9); \% 8 bins are allocated for exp
      d = bits(10:32); \% 23 bins for the significand
     for i = 1:8
         if (s = 1 | s = 0) & (k(i) = 1 | k(i) = 0)
            k_{-i}(i) = k(i) * 2 ^ (8-i);
         else
            error ('a bit can only be 0 or 1')
         end
      end
     for i = 1:23
         if d(i) = 1 | d(i) = 0
            d_i(i) = d(i) * 2 ^ (-i);
         else
            error ('a bit can only be 0 or 1')
         end
      end
     k_{-}10 = sum(k_{-}i);
      d_{-}10 = sum(d_{-}i);
```

Command Prompt:

```
>> SingPrec2Dec(['110000101110110101000000000000000'])
ans =
      -118.625
ans =
 0
ans =
 -1
ans =
 1
ans =
 10
```

Problem 2

```
function bin = Dec2SingPrec(dec)
if dec = 0
   bin = num2str(zeros(1, 32));
else
  if dec < 0
      s = [1]; \% s = 1
   else
      s = [0]; \% s = 0
   end
   dec = abs(dec); % the sign is ignored
  whole_number = floor(dec);
   w_n = whole_number;
   i = 1;
   while (1)
      w_{-n_{-}}2(i) = mod(w_{-n}, 2);
      w_n = floor(w_n / 2);
      i = i + 1;
      if w_n = 0, \dots
            break, end
   end
   w_n_2 = wrev(w_n_2); \% whole number in binary
   l = size(w_n_2); \% size of the whole num vect
  frac_number = dec - whole_number;
   f_n = frac_number;
   for i = 1:100
      if (f_n - 2^-i) >= 0
         f_{n_2}(i) = 1;
         f_n = f_n - 2(-i);
      else
         f_n_2(i) = 0;
         f_n = f_n;
      end
   end
```

```
if whole_number = 0
      e = l(2) - 1; % (true exp) going to the left: pos
      k = e + 127; % biased exponent
      index = [2: 1(2)];
      d = [w_n_2(index) f_n_2];
      %%%%considers dec num whose whole num is zero%%%%%%
   else
      %go through the values of f_n_2 to look for the 1st 1
      m = 1;
      while (1)
         if f_n_2(m) = 1, break, end
         m = m + 1;
      end
      %true exp = index of 1st 1
      e = -m; % (true exp) going to the right: neg
      k = e + 127; %biased expoenent
      len_f_n = size(f_n_2); % len of frac number in bin
      index_l_f_n = [m + 1: len_f_n(2)]; \% from num after 1
      f_{n-2} = f_{n-2} (index_{l-f_n});
      d = f<sub>n-2</sub>; %significand contains bits from frac num
   end
   k_{loop} = k;
   k_2 = zeros(1, 8); \% k_2 is set to contain 8 elements
   j = 1;
   while (1)
      k_2(j) = mod(k_{loop}, 2);
      k_{loop} = floor(k_{loop} / 2);
      j = j + 1;
      if k_{loop} = 0, break, end
   end
   k_{-2} = wrev(k_{-2}); % biased exp in binary
   index_sig = [1 : 23];
   bin = num2str([s k_2 d(index_sig)]);
end
end
```

Relative Percent Error Function

```
function rel_per_accu()
true_value = input('True value = ');
num_value = input('Numerical value = ');

eps = (abs(true_value - num_value)/true_value)*100;
fprintf('\n')
fprintf('Relative percent error is %.20f%%\n', eps)
fprintf('\n')
end
```

Command Prompt

```
>> c = Dec2SingPrec(-1)
c =
 0 0 0 0 0 0 0 0 0 0 0 0 0 0,
>> SingPrec2Dec(c)
ans =
 -1
\gg d = Dec2SingPrec(1)
d =
 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
>> SingPrec2Dec(d)
ans =
  1
>> e = Dec2SingPrec(-10)
e =
  0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
>> SingPrec2Dec(e)
ans =
-10
\gg f = Dec2SingPrec(10)
f =
  0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
```

```
>> SingPrec2Dec(f)
ans =
  10
\gg g = Dec2SingPrec(-0.5)
g =
 >> SingPrec2Dec(g)
ans =
           -0.5
\gg h = Dec2SingPrec (0.5)
h =
 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
>> SingPrec2Dec(h)
ans =
           0.5
\gg i = Dec2SingPrec(-0.2)
i =
 >> SingPrec2Dec(i)
ans =
    -0.199999988079071
```

```
>> rel_per_accu()
True value = -0.2
Numerical value = -0.199999988079071
Relative percent error is 0.00000596046451084575%
\gg j = Dec2SingPrec(0.2)
j =
   >> SingPrec2Dec(j)
ans =
       0.199999988079071
>> rel_per_accu()
True value = 0.2
Numerical value = 0.199999988079071
Relative percent error is 0.00000596046451084575%
\gg k = Dec2SingPrec(-0.0001)
k =
   1 1 0 1 1 1 0 0 0 1 0 1 1 1'
>> SingPrec2Dec(k)
ans =
    -9.99999974737875e-05
>> rel_per_accu()
True value = -0.0001
Numerical value = -9.99999974737875e-05
Relative percent error is 0.00000252621250198919%
```

Reference

1. M. R. King and N. A. Mody, Numerical and Statistical Methods for Bioengineering: Applications in MATLAB, Cambridge University Press, Cambridge, United Kingdom (2011).