1. Fill in the Blank (5000pts)

Note: Items 1-4 deal with Red-Black properties.

- 1. Every node is either red or black
- 2. Color-wise, the root and leaves are black
- 3. If a node is red, then its parent is black
- 4. All simple paths from any node x to a descendant leaf have the same number of black nodes
- 5. The pop operation is used to remove an element from a stack
- 6. The element removed from a queue is the first-most element added to the queue
- 7. Dijkstra's algorithm produces shortest paths between nodes.
- 8. If the next empty location is found in a squared number sequence, then you are using quadratic probing.

2.	Matching	(30,000)	points)	

- a. Comments (D) Accessed by a parameter to main.
- b. Block (C) In Java, integer, floating-point, Boolean, and character.
- c. Primitive Types (B) A sequence of statements within braces.
- d. Command-line argument (A) Make code easier for humans to read but have no semantic meaning.
- e. Null reference (E) The value of an object reference that does not refer to any object.
- f. Mergesort (G) A recursive instance that can be solved without recursion.
- g. Base Case (F) A divide-and-conquer algorithm that obtains an O(N log N) sort.
- h. Bellman-Ford algorithm (I) An alg that is used to solve the positive-weighted, shortest-path proble
- i. Dijkstra's algorithm (H) An alg that is used to solve the negative-weighted, shortest-path problem

3. (A Single Point) Fill in the runtimes of each sorting algorithm.

Sorting Algo	Big Oh	Theta
Selection Sort		
	$\mathrm{O}(n^2)$	Yes
Bubble Sort		
	$\mathrm{O}(n^2)$	Yes
Merge Sort		
	$\mathrm{O}(n^2)$	Yes
Heap Sort		
	O(nlogn)	No
Quick Sort		
	O(nlogn)	No

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4. (1600 pts) Give the expected Big-O runtime for the following operations on using the given data structures.

Red-Black Tree: search, min, max, successor, and predecessor O(logn)

Bubble Sort: Best, worst, and average case

 $n, n^2, and n^2$

Selection Sort: Best, worst, and average case

All n^2

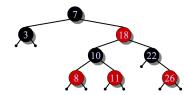
Heapsort: Best, worst, and average case n (if all keys are distinct, nlogn), nlogn, nlogn

5. (12 pts) **Red-Black** Below is the adjacency matrix for a Red-Black tree. The outer edges indicate the value at a given node (the first value is the root, the next two are the node's children (left to right), and so on in an inorder traversal).

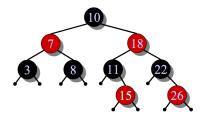
Draw the red-black tree below, indicating black nodes as boxes and red nodes as circles.

Γ	7	3	18	10	22	8	11	26]
7	0	1	1	0	0	0	0	0
3	1	0	0	0	0	0	0	0
18	1	0	0	1	1	0	0	0
10	0	0	1	0	0	1	1	0
22	0		1		0	0	0	1
8	0	0	0	1	0	0	0	0
11	0	0	0	1	0	0	0	0
L 26	0	0	0	0	1	0	0	0]

sorry, this one was a little bit of a disaster



Now, insert a node x=15 into the tree. *Idea:* Recolor, moving violation up the tree. Right-Rotate(18), Left-Rotate(7), and recolor.



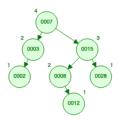
6

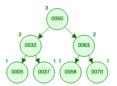
6. **Trees** (5! points)Perform a preorder traversal on this tree Remember: Preorder is root, left, right



 $_{\rm a,c,x,y,m,g,p,w,q}$

7. AVL (4! points) Perform the AVL instructions listed below. Perform correct rotations.





8. **AVL** (6? points) Do that again. Perform correct rotations. Here, make 2 the parent of 1 and 14 and the left child of 32.

9. (10 or 18 points) **Heaps**

(a) Describe the structure of a min-heap, how do each nodes compare to their children? The children in a min-heap will be larger than their parent.

(b) Fill the runtimes of min-heap operations

GetMin() O(1): returns the root element

RemoveMin() O(logn): remove min element (remove the root AND maybe heapify)
Insert() O(logn): add a new key, may have to traverse up to fix any errors

(c) Fill in the array index locations of each node

The Root A[0]

Node i's Left Child A[(2*i) +1]

Node i's Right Child A[(2*i)+2]

Node i's Parent A[(i-1)/2]

10. **Hash Slinging Slasher** (π points) Insert the following into a hash table using linear probing with a hash function $H(x) = k \mod 7$. [76, 93, 40, 47, 10, 55]

```
You should end up with this array: [47, 55, 93, 10, -40, 76]
```

What is quadratic probing? An open addressing method to avoid collisions. Uses a quadratic function to find the next available space for an item.

11. Code (2^e points) Describe the runtime of the following psedocode functions.

This runs in $O(\log n)$ time. The value of n is halved on each iteration of the loop. If $n = 2^x$, then $\log n = x$.

```
def IReallyLoveOwls(n)
    result = 0
    for i in range (n):
        for j in range (i, n)
            for k in range (n):
                 result += 1
    return result
```

 $O(n^3)$ since we run through 3 for loops, each of size n.