#### Announcements

- Reading
  - Today: Chapter 3 (3.3-3.4)
    - Skip details of code
  - Thursday: Chapter 5 (5.1-5.2)
- Program #1 Due at 10 PM not 10AM
- TA Office Hours
  - Th 1-3
  - F 4-6
  - phone x5-2776

## Error Codes (cont.)

#### Error Recovery

- Given m bits of data and r bits of error code
- Want to correct any one bit error
- There are n words one bit from each valid message
  - so need n+1 words for each valid message
  - thus  $(n + 1) 2^m <= 2^n$
  - but n = m + r so  $(m + r + 1) <= 2^r$

#### Hamming Code

- recovers from any one bit error
- number bits from left (starting at 1)
  - power of two bits are parity
  - rest contain data
- bit is checked by all parity bits in its sum of power expansion
  - bit 11 is used to compute parity bits 1, 2, and 8

## Hamming Code Example

Char	ASCII	Hamming
Н	1001 000	0011 0010 000
а	1100 001	1011 1001 001
m	1101 101	1110 1010 101
	1101 001	0110 1011 001

#### Burst Errors

- can send hamming codes by column rather than row
- if use k rows, then can detect any burst error up to k bits
  - uses kr bits to check a block km bits long

# Computing a Hamming Code

Bit #s	1	2	3	4	5	6	7	8	9	10	11
Parity/Data	Р	Р	D	Р	D	D	D	Р	D	D	D
Data To Snd			1		0	0	1		0	0	0
Parity Bit 1	0		1		0		1		0		0
Parity Bit 2		0	1			0	1			0	0
Parity Bit 4				1	0	0	1				
Parity Bit 8								0	0	0	0
Message	0	0	1	1	0	0	1	0	0	0	0

## Checking & Correcting a Hamming Code

Bit #s	1	2	3	4	5	6	7	8	9	10	11
Parity/Data	Р	Р	D	Р	D	D	D	Р	D	D	D
Data Sent	0	0	1	1	0	0	1	0	0	0	0
Data Recv	0	0	0	1	0	0	1	0	0	0	0
Parity Bit 1	1		0		0		1		0		0
Parity Bit 2		1	0			0	1			0	0
Parity Bit 4				1	0	0	1				
Parity Bit 8								0	0	0	0
XOR Paritys	1	1		0				0			
Corrected Msg	0	0	1	1	0	0	1	0	0	0	0

Binary # when XOR the parity is the bit position with the error (e.g. 0011 = bit 3 is wrong)

#### **Error Detection**

- Less bits are required
  - if errors are infrequent, then then this works better
  - assumes that re-transmission is possible
- Cyclic Redundancy Codes (CRC)
  - Use a generator function G(x) of degree r
    - r+1 bits long
    - x<sup>5</sup> + x<sup>2</sup> + 1 is degree 5 and represented as 100101
  - let M' be the message with r 0's on the end of it
  - divide M' into G(x) and compute remainder
    - use this as the r bit CRC code
  - a code with r bits will detect all burst errors less than r bits

### CRC's

#### several G's are standardized

$$- CRC-12 = x^{12} + x^{11} + x^3 + x^2 + x + 1$$

$$- CRC-16 = x^{16} + x^{15} + x^2 + 1$$

$$- CRC-CCITT = x^{16} + x^{12} + x^5 + 1$$

#### 16 bit CRC will catch

- all single and double bit errors
- all errors with an odd number of bits
- all burst errors of length less than 16

## **CRC** Example

Frame : 1101011011

Generator: 10011

Message after appending 4 zero bits: 1 1 0 1 0 1 1 0 0 0 0

Division is done using XOR

1 0 1 0 0 1 0 0 1 1 0 1 1 1 0 0 0 0 0 0 1 1 1 0

Remainder

Transmitted frame: 1101011111110

#### PPP Protocol

- Link Protocol for Serial Lines
  - Supports multiple network protocols: IP, IPX, CLNP, ...
  - designed for dialup or leased lines
- Link Establishment (via LCP Link Control Protocol)
  - Negotiate Options
    - configure-request: list of proposed options and values
    - configure-{ack/nack}: will (won't) use the requested option
  - Allows for authentication

1	. 1	. 1	1 or 2	variable	2 or 4	. 1
flag	Address	control	protocol	payload	checksum	flag
01111110	11111111	00001110				01111110

From: Computer Networks, 3rd Ed. by Andrew S. Tanenbaum, (c)1996 Prentice Hall.

### PPP Cont.

- NCP protocol
  - per network level protocol
  - used to establish network attributes (e.g. addresses)
  - high bit of protocol # is a one
- Notes on Link Format
  - character stuff flag byte in data
    - Escape Character is 0x7d (0111 1101)
    - Escape Character and Frame Marker sent at
      - <Esc-Char><data XOR 0x20>
  - option to skip address and control fields (since constant)
- IP
  - Protocol byte (0x21) or 0x8021 for IP NCP

### **ATM Datalink Protocol**

- Header
  - use CRC over the 32 bits of the header
- How to find cell boundary?
  - use shift register to check for valid checksum
    - 1/256 chance of a random match
  - use HUNT mode to increase chances
    - after a good cell, skip to the next cell boundary
    - must receive δ cells with checksum matches
- Detecting loss of synchronization
  - one bad cell is probably an error
  - many bad cells is likely a slip (loss of sync)
  - if  $\alpha$  bad cells are seen in a row, switch to hunt mode