WEEK - 9

GREEDY

- When a problem does not exhibit the optimal substructure property, it means that a straightforward application of dynamic programming may not be appropriate or may fail to provide the optimal solution. Here are a few scenarios where the optimal substructure property fails:
- Scenario 1: Overlapping Subproblems
- Dynamic programming relies on storing solutions to subproblems to avoid redundant computations.
- If there are overlapping subproblems, meaning the same subproblem is solved multiple times, DP can be inefficient.
- In such cases, memoization (top-down DP) or tabulation (bottom-up DP) can still be used, but the efficiency gains may not be significant.

- Scenario 2: Greedy Choice is Not Always Optimal
- Some problems require making greedy choices at each step, but these local optimal choices do not lead to a global optimal solution.
- Greedy algorithms, which make the best choice at each step, might not work if the problem doesn't have the property that a global optimum can be reached by selecting local optima.
- In such cases, a different algorithmic approach may be needed.

- Scenario 3: Non-Monotonicity
- Problems where the optimal solution does not necessarily include the optimal solution to subproblems.
- An example is the "Longest Increasing Subsequence" problem, where simply choosing the longest subsequence at each step does not guarantee the overall longest increasing subsequence.
- This violates the optimal substructure because the optimal solution of the problem cannot be directly constructed from the optimal solutions of its subproblems.

- Scenario 4: Dependency on Future Choices
- Problems where the optimal solution at a particular step depends on future choices, which dynamic programming cannot anticipate.
- An example is the "0-1 Knapsack Problem" with a constraint on the weight of items.
- Choosing an item in the current step may affect the choices available for future steps, making it challenging to use DP directly.
- Scenario 5: NP-Hard or NP-Complete Problems
- Problems that are inherently complex, such as "Traveling Salesman Problem" or "Subset Sum Problem".
- These problems do not have efficient solutions in general and dynamic programming might not be a suitable approach due to the exponential time complexity.

Dynamic Programming: Longest Common Subsequence

- Problem Statement:
- Given two sequences of characters, find the length of the longest subsequence present in both sequences.
- Example:
- Sequence 1: "ABCDGH"
- Sequence 2: "AEDFHR"
- Longest Common Subsequence (LCS): "ADH" (length = 3)

- Approach:
- Dynamic Programming is used to solve this problem efficiently.
- We create a 2D table to store the lengths of LCS for different prefixes of the two sequences.
- Dynamic Programming Table:
- Let's use the sequences "ABCBDAB" and "BDCAB" for illustration.

- B D C A B
- 000000
- A 0 0 0 0 1 1
- B 0 0 0 0 1 2
- C 0 0 1 1 1 2
- B 0 0 1 1 1 2
- D 0 1 1 1 1 2
- A 0 1 1 1 2 2
- B 0 1 1 1 2 3

- Algorithm:
- Initialization: Fill the first row and column of the table with zeros.
- Build the Table:
 - If characters match:
 - Increment the value in the diagonal by 1.
 - If characters don't match:
 - Take the maximum of the value above or left.
- Result: The bottom-right cell contains the length of the LCS.

- Traceback:
- To find the actual LCS:
 - Start from the bottom-right cell.
 - Move diagonally up whenever the characters match.
 - Follow the path until reaching the top-left cell.
- Complexity:
- Time Complexity: O(mn) where m and n are the lengths of the two sequences.
- Space Complexity: O(mn) for the 2D table.

Introduction to Greedy Algorithm

- What is a Greedy Algorithm?
- A Greedy Algorithm is an algorithmic paradigm that makes a series of locally optimal choices at each step with the hope of finding a global optimum.
- It makes the best choice at each step without regard for the global structure, with the belief that this will lead to the best possible solution.
- Greedy algorithms are quite efficient and straightforward to implement, making them useful in a variety of problems.

Greedy Interval Scheduling

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- Greedy Interval Scheduling
- Problem Statement:
- Given a set of tasks with start and finish times, the goal is to find the maximum number of tasks that can be completed without conflicts.
- Greedy Solution:

- Sort: First, sort the tasks by their finish times in ascending order.
- Select: Iterate through the sorted tasks. At each step, choose the task that finishes earliest and does not conflict with previously selected tasks.
- Example:
- Tasks:
 - Task 1: Start 1, Finish 4
 - Task 2: Start 3, Finish 5
 - Task 3: Start 0, Finish 6
 - Task 4: Start 5, Finish 7
 - Task 5: Start 3, Finish 8
 - Task 6: Start 5, Finish 9

Prefix Codes (Huffman Code)

- What are Prefix Codes?
- Prefix codes are codes with the property that no code word is a prefix of another.
- Huffman coding is a method to construct such prefix codes efficiently.
- Huffman Coding Steps:
- Frequency Table: Create a frequency table of characters based on their occurrence in the given data.

Build Huffman Tree:

- Start with all characters as individual nodes with their frequencies.
- Combine two nodes with the lowest frequencies into a new node. Repeat until there's only one node left.

Assign Codes:

- Traverse the Huffman tree:
 - Left edge -> '0'
 - Right edge -> '1'
- Assign codes by the path from the root to each character.

- Sorted by Finish Time:
 - Task 3, Task 1, Task 2, Task 5, Task 4, Task 6
- Greedy Selection:
 - Select Task 3, Task 5, Task 6 for a total of 3 tasks.
- Explanation:
- Greedy Interval Scheduling works because selecting the task that finishes earliest allows room for more tasks to fit into the schedule without overlap.
- It doesn't always provide the optimal solution, but in this case, it does.

- Example:
- Data: "ABBCCCDDDDEEEEE"
- Frequency Table:
 - A: 1, B: 2, C: 3, D: 4, E: 5

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/\
/\
E:5 /\
A:1 B:2
/\
C:3 D:4
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- Codes:
 - A: 00, B: 01, C: 10, D: 11, E: 1
- Explanation:
- Huffman Coding optimally assigns shorter codes to more frequent characters.
- Prefix property ensures no ambiguity in decoding, as no code is a prefix of another.

Summary:

- Greedy algorithms make locally optimal choices with the hope of finding a global optimum.
- Greedy Interval Scheduling selects tasks based on earliest finish times without conflicts.
- Huffman Coding creates prefix codes efficiently for data compression.
- Feel free to use these points and examples to create a PowerPoint presentation on Greedy Algorithms, Greedy Interval Scheduling, and Huffman Coding!

- What is the main concept behind Greedy Algorithms?
- A) Making decisions based on current information
- B) Randomly selecting options
- C) Considering all possible solutions
- D) Using trial and error

- Which of the following is an example of a Greedy Interval Scheduling problem?
- A) Graph traversal
- B) Job Scheduling
- C) Sorting algorithms
- D) None

- Which algorithm is commonly used to construct Huffman Codes?
- A) Kruskal's algorithm
- B) Prim's algorithm
- C) Dijkstra's algorithm
- D) Huffman's algorithm

- What is the advantage of using Huffman Codes for data compression?
- Reduced file size
- Faster data transmission
- Improved data security
- Efficient coding

- Which scheduling problem can be solved using the Interval Scheduling algorithm?
- A) Job Scheduling
- B) Task Allocation
- C) Network Routing
- D) Resource Allocation

- What is the time complexity of the Interval Scheduling algorithm?
- ▶ A) O(n^2)
- ▶ B) O(log n)
- ► C) O(2^n)
- ▶ D) **O(n)**

- Which data structure is commonly used to implement Huffman Codes?
- A) Heap
- B) Linked List
- C) Hash Table
- D) Binary Search Tree

- In Huffman coding, data in a tree always occur?
 a) roots
 - b) leaves
 - c) left sub trees
 - d) right sub trees