Contents

1	1 Basic	1
	1.1 vimrc	1
	1.2 Pragma	1
_	2. Pata Standard	
2	2 Data Structure	1
	2.1 Black Magic	1
	2.2 Lazy Segment Tree	
	2.3 Lichao Tree	
	2.4 Linear Basis	
	2.5 Heavy Light Decomposition	
	2.6 Link Cut Tree	2
3	3 Graph	3
,	3.1 Bridge CC	
	3.2 Vertex BCC	3
	3.3 Strongly Connected Component	
	3.4 Two SAT	
	3.5 Virtual Tree	
	3.6 Dominator Tree	
	3.7 Dinic	
	3.8 Min Cost Max Flow	5
	3.9 Stoer Wagner Algorithm	
	3.10General Matching	
	3.11Hopcroft Karp Algorithm	
	3.12Directed MST	
	3.13Edge Coloring	7
1	4 Geometry	7
4		-
	4.1 Basic	7
	4.2 2D Convex Hull	/
5	5 String	8
	•	
	5.1 KMP	8
	5.1 KMP	
	5.1 KMP	
	5.1 KMP. 5.2 Z Value	8 8
	5.1 KMP. 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem	
	5.1 KMP. 5.2 Z Value	
	5.1 KMP. 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers	
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number	8
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes	8
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd	8
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem	
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6 Math 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve	
	5.1 KMP. 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve 6.6 Fast Walsh Transform	
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve 6.6 Fast Walsh Transform 6.7 Floor Sum	
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve 6.6 Fast Walsh Transform 6.7 Floor Sum 6.8 Linear Programming	
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve 6.6 Fast Walsh Transform 6.7 Floor Sum 6.8 Linear Programming 6.9 Miller Rabin	
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve 6.6 Fast Walsh Transform 6.7 Floor Sum 6.8 Linear Programming 6.9 Miller Rabin 6.10 Pollard's Rho	
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve 6.6 Fast Walsh Transform 6.7 Floor Sum 6.8 Linear Programming 6.9 Miller Rabin 6.10 Pollard's Rho 6.11 Gauss Elimination	
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve 6.6 Fast Walsh Transform 6.7 Floor Sum 6.8 Linear Programming 6.9 Miller Rabin 6.10Pollard's Rho 6.11Gauss Elimination 6.12Fast Fourier Transform	
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve 6.6 Fast Walsh Transform 6.7 Floor Sum 6.8 Linear Programming 6.9 Miller Rabin 6.10Pollard's Rho 6.11Gauss Elimination 6.12Fast Fourier Transform 6.133 Primes NTT	
	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve 6.6 Fast Walsh Transform 6.7 Floor Sum 6.8 Linear Programming 6.9 Miller Rabin 6.10Pollard's Rho 6.11Gauss Elimination 6.12Fast Fourier Transform	
6	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve 6.6 Fast Walsh Transform 6.7 Floor Sum 6.8 Linear Programming 6.9 Miller Rabin 6.10Pollard's Rho 6.11Gauss Elimination 6.12Fast Fourier Transform 6.133 Primes NTT 6.14Number Theory Transform	8
6	5.1 KMP 5.2 Z Value 5.3 Suffix Array 5.4 Booth Algorithm 5.5 Manacher Algorithm 6.1 Lemma And Theory 6.1.1 Pick's Theorem 6.1.2 Euler's Planar Graph Theorem 6.1.3 Modular inversion recurrence 6.2 Numbers 6.2.1 Catalan number 6.2.2 Primes 6.3 Extgcd 6.4 Chinese Remainder Theorem 6.5 Linear Sieve 6.6 Fast Walsh Transform 6.7 Floor Sum 6.8 Linear Programming 6.9 Miller Rabin 6.10Pollard's Rho 6.11Gauss Elimination 6.12Fast Fourier Transform 6.133 Primes NTT	8

1 Basic

1.1 vimrc

1.2 Pragma

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,tune=native")
```

2 Data Structure

```
2.1 Black Magic
template<typename T>
using pbds_tree = tree<T, null_type, less<T>,
   rb_tree_tag, tree_order_statistics_node_update>;
// find_by_order: like array accessing, order_of_key
2.2 Lazy Segment Tree
// 0-based, [l, r)
// Remember to call init
struct tag {
  // Construct identity element
  tag() { }
  // apply tag
  tag& operator+=(const tag &b) {
    return *this;
struct node {
  // Construct identity element
  node() { }
  // Merae two nodes
  node operator+(const node &b) const {
    node res = node();
    return res;
  // Apply tag to this node
  void operator()(const tag &t) {
};
template<typename N, typename T>
struct lazy_segtree {
  N arr[maxn << 1];
  T tag[maxn];
  int n;
  void init(const vector<N> &a) {
    n = a.size();
    for (int i = 0; i < n; i++)</pre>
      arr[i + n] = a[i], tag[i] = T();
    for (int i = n - 1; i; i--)
      arr[i] = arr[i << 1] + arr[i << 1 | 1];</pre>
  void upd(int p, T v) {
    if (p < n)
      tag[p] += v;
    arr[p](v);
  void pull(int p) {
    for (p >>= 1; p; p >>= 1) {
      arr[p] = arr[p << 1] + arr[p << 1 | 1];
      arr[p](tag[p]);
    }
  void push(int p) {
    for (int h = __lg(p); h; h--) {
      int i = p >> h;
      upd(i << 1, tag[i]);
upd(i << 1 | 1, tag[i]);
      tag[i] = T();
  void edt(int 1, int r, T v) {
    int tl = 1 + n, tr = r + n - 1;
    push(t1); push(tr);
    for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
      if (1 & 1)
        upd(l++, v);
      if (r & 1)
        upd(--r, v);
    pull(tl); pull(tr);
  N que(int 1, int r) {
    N resl = N(), resr = N();
    int tl = 1 + n, tr = r + n - 1;
    push(t1); push(tr);
    for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
      if (1 & 1)
        resl = resl + arr[l++];
      if (r & 1)
```

resr = arr[--r] + resr;

```
return resl + resr;
  }
};
2.3 Lichao Tree
struct lichao { // maxn: range
  struct line {
    ll a, b;
    line(): a(0), b(0) { } // or LINF
    line(ll a, ll b): a(a), b(b) { }
    11 operator()(11 x) { return a * x + b; }
  } arr[maxn << 2];</pre>
  void insert(int 1, int r, int id, line x) {
    int m = (1 + r) >> 1;
    if(arr[id](m) < x(m))
      swap(arr[id], x);
    if(1 == r - 1)
      return;
    if(arr[id].a < x.a)</pre>
      insert(m, r, id << 1 | 1, x);
    else
      insert(1, m, id << 1, x);
  } // change to > if query min
  void insert(ll a, ll b) { insert(0, maxn, 1, line(a,
    b)); }
  11 que(int 1, int r, int id, int p) {
  if(1 == r - 1)
      return arr[id](p);
    int m = (1 + r) >> 1;
    if(p < m)
      return max(arr[id](p), que(l, m, id << 1, p));</pre>
```

return max(arr[id](p), que(m, r, id << 1 | 1, p));</pre>

11 que(int p) { return que(0, maxn, 1, p); }

2.4 Linear Basis

} tree;

} // chnage to min if query min

```
template<int BITS>
struct linear_basis {
  array<uint64_t, BITS> basis;
  linear_basis() { basis.fill(0); }
  void insert(uint64_t x) {
    for (int i = BITS - 1; i >= 0; i--) if ((x >> i) &
    1) {
      if (basis[i] == 0) {
        basis[i] = x;
        return;
      x ^= basis[i];
    }
  }
  bool valid(uint64_t x) {
    for (int i = BITS - 1; i >= 0; i--)
      if ((x >> i) & 1) x ^= basis[i];
    return x == 0;
  }
  uint64_t operator[](int i) { return basis[i]; }
}; // max xor sum: greedy from high bit
// min xor sum: zero(if possible) or min_element
```

2.5 Heavy Light Decomposition

```
/* Requirements:
* N := the count of nodes
* edge[N] := the edges of the graph
* Can be modified:
* tree := Segment Tree or other data structure
struct heavy_light_decomposition {
 int dep[N], pa[N], hea[N], hev[N], pos[N], t;
 int dfs(int u) {
    int mx = 0, sz = 1;
    hev[u] = -1;
    for(int v : edge[u]) {
     if(v == pa[u])
        continue;
     pa[v] = u;
     dep[v] = dep[u] + 1;
     int c = dfs(v);
      if(c > mx)
       mx = c, hev[u] = v;
```

```
sz += c;
    return sz;
  void find_head(int u, int h) {
    hea[u] = h;
    pos[u] = t++; // 0-indexed !!!
    if(~hev[u])
      find_head(hev[u], h);
    for(int v : edge[u])
      if(v != pa[u] && v != hev[u])
        find_head(v, v);
  void init(int rt) {
    dfs(rt, rt);
    find_head(rt, rt);
  /* It is necessary to edit below for every use */
  void edt(int a, int b, int v) {
  int query(int a, int b) { // query path sum
    int res = 0;
    for(; hea[a] != hea[b]; a = pa[hea[a]]) {
      if(dep[hea[a]] < dep[hea[b]])</pre>
        swap(a, b):
      res += tree.que(pos[hea[a]], pos[a] + 1);
    if(dep[a] > dep[b])
      swap(a, b);
    return res + tree.que(pos[a], pos[b] + 1);
} hld;
2.6 Link Cut Tree
namespace LCT {
  const int N = 1e5 + 25;
  int pa[N], ch[N][2];
  11 dis[N], prv[N], tag[N];
  vector<pair<int, int>> edge[N];
  vector<pair<11, 11>> eve;
```

```
inline bool dir(int x) { return ch[pa[x]][1] == x; }
inline bool is_root(int x) { return ch[pa[x]][0] != x
   && ch[pa[x]][1] != x; }
inline void rotate(int x) {
  int y = pa[x], z = pa[y], d = dir(x);
  if(!is root(y))
    ch[z][dir(y)] = x;
  pa[x] = z;
  ch[y][d] = ch[x][!d];
  if(ch[x][!d])
   pa[ch[x][!d]] = y;
  ch[x][!d] = y;
 pa[y] = x;
inline void push_tag(int x) {
  if(!tag[x])
   return;
  prv[x] = tag[x];
  if(ch[x][0])
    tag[ch[x][0]] = tag[x];
  if(ch[x][1])
    tag[ch[x][1]] = tag[x];
 tag[x] = 0;
void push(int x) {
  if(!is_root(x))
    push(pa[x]);
 push_tag(x);
inline void splay(int x) {
 push(x);
  while(!is_root(x)) {
    if(int y = pa[x]; !is_root(y))
      rotate(dir(y) == dir(x) ? y : x);
    rotate(x);
 }
inline void access(ll t, int x) {
 int lst = 0, tx = x;
  while(x) {
    splay(x);
```

```
if(lst) {
        ch[x][1] = lst;
        eve.push_back({prv[x] + dis[x], t + dis[x]});
      lst = x;
      x = pa[x];
    splay(tx);
    if(ch[tx][0])
      tag[ch[tx][0]] = t;
  void dfs(int u) {
    prv[u] = -LINF;
    for(const auto &[v, c] : edge[u]) {
      if(v == pa[u])
        continue;
      pa[v] = u;
      ch[u][1] = v;
      dis[v] = dis[u] + c;
      dfs(v);
    }
 }
};
```

3 Graph

3.1 Bridge CC

```
namespace bridge_cc {
 vector<int> tim, low;
  stack<int, vector<int>> st;
  int t, bcc_id;
  void dfs(int u, int p, const vector<vector<pair<int,</pre>
    int>>> &edge, vector<int> &pa) {
    tim[u] = low[u] = t++;
    st.push(u);
    for(const auto &[v, id] : edge[u]) {
      if(id == p)
        continue;
      if(tim[v])
        low[u] = min(low[u], tim[v]);
        dfs(v, id, edge, pa);
        if(low[v] > tim[u]) {
          int x;
          do {
            pa[x = st.top()] = bcc_id;
            st.pop();
          } while(x != v);
          bcc_id++;
        }
        else
          low[u] = min(low[u], low[v]);
      }
   }
 }
  vector<int> solve(const vector<vector<pair<int, int</pre>
    >>> &edge) { // (to, id)
    int n = edge.size();
    tim.resize(n);
   low.resize(n);
   t = bcc_id = 1;
    vector<int> pa(n);
    for(int i = 0; i < n; i++) {</pre>
      if(!tim[i]) {
        dfs(i, -1, edge, pa);
        while(!st.empty()) {
          pa[st.top()] = bcc_id;
          st.pop();
        bcc_id++;
     }
    return pa;
 } // return bcc id(start from 1)
```

3.2 Vertex BCC

```
class bicon_cc {
  private:
    int n, ecnt;
```

```
vector<vector<pair<int, int>>> G;
    vector<int> bcc, dfn, low, st;
    vector<bool> ap, ins;
    void dfs(int u, int f)
      dfn[u] = low[u] = dfn[f] + 1;
      int ch = 0;
      for (auto [v, t]: G[u]) if (v != f) {
        if (!ins[t]) {
          st.push_back(t);
          ins[t] = true;
        if (dfn[v]) {
          low[u] = min(low[u], dfn[v]);
          continue;
        }
        ++ch;
        dfs(v, u);
        low[u] = min(low[u], low[v]);
        if (low[v] >= dfn[u]) {
          ap[u] = true;
          while (true) {
            int eid = st.back();
            st.pop_back();
            bcc[eid] = ecnt;
            if (eid == t) break;
          ecnt++;
        }
      if (ch == 1 && u == f) ap[u] = false;
    }
  public:
    void init(int n_) {
      G.clear(); G.resize(n = n_);
      ecnt = 0; ap.assign(n, false);
      low.assign(n, 0); dfn.assign(n, 0);
    void add_edge(int u, int v) {
      G[u].emplace_back(v, ecnt);
      G[v].emplace_back(u, ecnt++);
    void solve() {
      ins.assign(ecnt, false);
      bcc.resize(ecnt); ecnt = 0;
      for (int i = 0; i < n; ++i)</pre>
        if (!dfn[i]) dfs(i, i);
    // The id of bcc of the x-th edge (0-indexed)
    int get_id(int x) { return bcc[x]; }
    // Number of bcc
    int count() { return ecnt; }
    bool is_ap(int x) { return ap[x]; }
}; // 0-indexed
```

3.3 Strongly Connected Component

```
namespace scc {
  vector<int> edge[maxn], redge[maxn];
  stack<int, vector<int>> st;
  bool vis[maxn];
  void dfs(int u) {
    vis[u] = true;
    for(int v : edge[u])
      if(!vis[v])
        dfs(v);
    st.push(u);
  void dfs2(int u, vector<int> &pa) {
    for(int v : redge[u])
      if(!pa[v])
        pa[v] = pa[u], dfs2(v, pa);
  void add_edge(int u, int v) {
    edge[u].push_back(v);
    redge[v].push_back(u);
  // pa[i]: scc id of all nodes in topo order
  vector<int> solve(int n) {
    vector<int> pa(n + 1);
    for(int i = 1; i <= n; i++)</pre>
      if(!vis[i])
        dfs(i);
```

int id = 1; // start from 1

while(!st.empty()) {

st.pop();

int u = st.top();

```
if(!pa[u])
        pa[u] = id++, dfs2(u, pa);
    return pa;
 } // 1-based
3.4 Two SAT
// maxn >= 2 * n (n: number of variables)
// clauses: (x, y) = x V y, -x if neg, var are 1-based
// return empty is no solution
vector<bool> solve(int n, const vector<pair<int, int>>
    &clauses) {
  auto id = [\&](int x) { return abs(x) + n * (x < 0);
    };
  for(const auto &[a, b] : clauses) {
   scc::add_edge(id(-a), id(b));
    scc::add_edge(id(-b), id(a));
 auto pa = scc::solve(n * 2);
 vector<bool> ans(n + 1);
 for(int i = 1; i <= n; i++) {</pre>
   if(pa[i] == pa[i + n])
      return vector<bool>();
   ans[i] = pa[i] > pa[i + n];
  return ans;
}
3.5 Virtual Tree
```

```
// dfn: the dfs order, vs: important points, r: root
vector<pair<int, int>> build(vector<int> vs, int r) {
  vector<pair<int, int>> res;
  sort(vs.begin(), vs.end(), [](int i, int j) {
      return dfn[i] < dfn[j]; });</pre>
  vector<int> s = {r};
  for (int v : vs) if (v != r) {
    if (int o = lca(v, s.back()); o != s.back()) {
      while (s.size() >= 2) {
        if (dfn[s[s.size() - 2]] < dfn[o]) break;</pre>
        res.emplace_back(s[s.size() - 2], s.back());
        s.pop_back();
      if (s.back() != o) {
        res.emplace_back(o, s.back());
        s.back() = o;
      }
    s.push back(v);
  for (size_t i = 1; i < s.size(); ++i)</pre>
    res.emplace_back(s[i - 1], s[i]);
  return res; // (x, y): x->y
} // The returned virtual tree contains r (root).
```

3.6 Dominator Tree

```
'* Find dominator tree with root s in O(n)
* Return the father of each node, **-2 for unreachable
     ** */
struct dominator_tree { // 0-based
 int tk;
 vector<vector<int>> g, r, rdom;
 vector<int> dfn, rev, fa, sdom, dom, val, rp;
  dominator\_tree(int n): tk(0), g(n), r(n), rdom(n),
 dfn(n, -1), rev(n, -1), fa(n, -1), sdom(n, -1),
  dom(n, -1), val(n, -1), rp(n, -1) {}
  void add_edge(int x, int y) { g[x].push_back(y); }
 void dfs(int x) {
    rev[dfn[x] = tk] = x;
    fa[tk] = sdom[tk] = val[tk] = tk;
    for (int u : g[x]) {
     if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
     r[dfn[u]].push_back(dfn[x]);
   }
  void merge(int x, int y) { fa[x] = y; }
 int find(int x, int c = 0) {
```

```
if (fa[x] == x) return c ? -1 : x;
     if (int p = find(fa[x], 1); p != -1) {
       if (sdom[val[x]] > sdom[val[fa[x]]])
         val[x] = val[fa[x]];
       fa[x] = p;
       return c ? p : val[x];
     } else {
       return c ? fa[x] : val[x];
  }
  vector<int> build(int s, int n) {
     dfs(s);
     for (int i = tk - 1; i >= 0; --i) {
       for (int u : r[i])
        sdom[i] = min(sdom[i], sdom[find(u)]);
       if (i) rdom[sdom[i]].push_back(i);
       for (int u : rdom[i]) {
         int p = find(u);
         dom[u] = (sdom[p] == i ? i : p);
       if (i) merge(i, rp[i]);
     vector<int> p(n, -2);
     p[s] = -1;
     for (int i = 1; i < tk; ++i)</pre>
      if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
     for (int i = 1; i < tk; ++i)</pre>
      p[rev[i]] = rev[dom[i]];
    return p;
};
```

3.7 Dinic

```
// Return max flor from s to t. INF, LINF and maxn
    required
template < typename T> // maxn: edge/node counts
struct dinic { // T: int or ll, up to range of flow
  const T IN_INF = (is_same_v<T, int>) ? INF : LINF;
  struct E {
    int v; T c; int r;
    E(int v, T c, int r):
      v(v), c(c), r(r){}
  vector<E> adj[maxn];
  pair<int, int> is[maxn]; // counts of edges
  void add_edge(int u, int v, T c, int i = 0) {
    is[i] = {u, adj[u].size()};
    adj[u].push_back(E(v, c, (int) adj[v].size()));
adj[v].push_back(E(u, 0, (int) adj[u].size() - 1));
  int n, s, t;
  void init(int nn, int ss, int tt) {
    n = nn, s = ss, t = tt;
    for (int i = 0; i <= n; ++i)
      adj[i].clear();
  int le[maxn], it[maxn];
  int bfs() {
    fill(le, le + maxn, -1); le[s] = 0;
    queue<int> q; q.push(s);
    while (!q.empty()) {
      int u = q.front(); q.pop();
      for (auto [v, c, r] : adj[u]) {
         if (c > 0 \&\& le[v] == -1)
           le[v] = le[u] + 1, q.push(v);
      }
    return ~le[t];
  T dfs(int u, T f) {
    if (u == t) return f;
    for (int &i = it[u]; i < (int) adj[u].size(); ++i)</pre>
      auto &[v, c, r] = adj[u][i];
if (c > 0 && le[v] == le[u] + 1) {
         T d = dfs(v, min(c, f));
        if (d > 0) {
           c -= d:
           adj[v][r].c += d;
           return d;
```

```
return 0;
  T flow() {
    T ans = 0, d;
    while (bfs()) {
      fill(it, it + maxn, 0);
      while ((d = dfs(s, IN_INF)) > 0) ans += d;
    return ans:
  T rest(int i) {
    return adj[is[i].first][is[i].second].c;
|};
```

Min Cost Max Flow 3.8

```
struct cost_flow { // maxn: node count
  static const int64_t INF = 102938475610293847LL;
  struct Edge {
    int v, r;
    int64_t f, c;
    Edge(int a,int b,int _c,int d):v(a),r(b),f(_c),c(d)
  int n, s, t, prv[maxn], prvL[maxn], inq[maxn];
  int64_t dis[maxn], fl, cost;
  vector<Edge> E[maxn];
  void init(int _n, int _s, int _t) {
    n = _n; s = _s; t = _t;
for (int i = 0; i < n; i++) E[i].clear();</pre>
    fl = cost = 0;
  void add_edge(int u, int v, int64_t f, int64_t c) {
    E[u].push_back(Edge(v, E[v].size() , f, c));
    E[v].push_back(Edge(u, E[u].size()-1, 0, -c));
  pair<int64_t, int64_t> flow() {
    while (true) {
      for (int i = 0; i < n; i++) {
        dis[i] = INF;
        inq[i] = 0;
      dis[s] = 0;
      queue<int> que;
      que.push(s);
      while (!que.empty()) {
        int u = que.front(); que.pop();
        inq[u] = 0;
        for (int i = 0; i < E[u].size(); i++) {</pre>
          int v = E[u][i].v;
          int64_t w = E[u][i].c;
          if (E[u][i].f > 0 && dis[v] > dis[u] + w) {
            prv[v] = u; prvL[v] = i;
            dis[v] = dis[u] + w;
            if (!inq[v]) {
              inq[v] = 1;
               que.push(v);
            }
          }
        }
      if (dis[t] == INF) break;
      int64_t tf = INF;
      for (int v = t, u, 1; v != s; v = u) {
        u = prv[v]; l = prvL[v];
        tf = min(tf, E[u][1].f);
      for (int v = t, u, 1; v != s; v = u) {
        u = prv[v]; l = prvL[v];
        E[u][1].f -= tf;
        E[v][E[u][1].r].f += tf;
      cost += tf * dis[t];
      fl += tf;
    return {fl, cost};
  }
};
```

auto lca = [&](int x, int y) { static int tk = 0; tk++; x = fnd(x);y = fnd(y);for(;; swap(x, y)) **if**(x != n) { if(mark[x] == tk)return x: mark[x] = tk;x = fnd(pre[match[x]]); } auto blossom = [&](int x, int y, int 1) { while(fnd(x) != 1) { pre[x] = y;y = match[x]; if(s[y] == 1)que.push_back(y), s[y] = 0; if(pa[x] == x) pa[x] = 1;**if**(pa[y] == y) pa[y] = 1; x = pre[y];} };

Stoer Wagner Algorithm

```
// return global min cut in O(n^3)
struct SW { // 1-based
  int edge[maxn][maxn], wei[maxn], n;
  bool vis[maxn], del[maxn];
  void init(int _n) {
    n = _n; MEM(edge, 0); MEM(del, 0);
  void add_edge(int u, int v, int w) {
    edge[u][v] += w; edge[v][u] += w;
  void search(int &s, int &t) {
    MEM(wei, 0); MEM(vis, 0);
    s = t = -1;
    while(true) {
      int mx = -1;
      for(int i = 1; i <= n; i++) {
   if(del[i] || vis[i]) continue</pre>
         if(mx == -1 || wei[mx] < wei[i])</pre>
           mx = i;
      if(mx == -1) break;
      vis[mx] = true;
      s = t; t = mx;
      for(int i = 1; i <= n; i++)</pre>
         if(!vis[i] && !del[i])
           wei[i] += edge[mx][i];
    }
  int solve() {
    int ret = INF;
    for(int i = 1; i < n; i++) {</pre>
      int x, y;
      search(x, y);
      ret = min(ret, wei[y]);
      del[y] = true;
      for(int j = 1; j <= n; j++) {</pre>
        edge[x][j] += edge[y][j];
         edge[j][x] += edge[y][j];
      }
    }
    return ret;
  }
} sw;
```

3.10 General Matching

```
// Find max matching on general graph in O(|V|^3)
vector<int> max_matching(vector<vector<int>> g) {
  int n = g.size();
  vector < int > match(n + 1, n), pre(n + 1, n), que;
  vector < int > s(n + 1), mark(n + 1), pa(n + 1);
  function<int(int)> fnd = [&](int x) {
    if(x == pa[x]) return x;
    return pa[x] = fnd(pa[x]);
  auto bfs = [&](int r) {
    fill(s.begin(), s.end(), -1);
    iota(pa.begin(), pa.end(), 0);
    que = \{r\}; s[r] = 0;
```

```
for(int it = 0; it < que.size(); it++) {</pre>
      int x = que[it];
      for(int u : g[x]) {
        if(s[u] == -1) {
          pre[u] = x;
          s[u] = 1;
          if(match[u] == n) {
            for(int a = u, b = x, lst;
                b != n; a = lst, b = pre[a]) {
              lst = match[b];
              match[b] = a;
              match[a] = b;
            }
            return;
          }
          que.push_back(match[u]);
          s[match[u]] = 0;
        else if(s[u] == 0 \&\& fnd(u) != fnd(x)) {
          int 1 = 1ca(u, x);
          blossom(x, u, 1);
          blossom(u, x, 1);
        }
      }
   }
  }:
  for(int i = 0; i < n; i++)</pre>
   if(match[i] == n) bfs(i);
  match.resize(n);
  for(int i = 0; i < n; i++)</pre>
    if(match[i] == n) match[i] = -1;
  return match:
} // 0-based
3.11 Hopcroft Karp Algorithm
```

```
// Find maximum bipartite matching in O(Esqrt(V))
// g: edges for all nodes at left side
vector<int> hopcroft_karp(vector<vector<int>> g, int 1,
     int r) {
  vector<int> match_l(l, -1), match_r(r, -1);
 vector<int> dis(1);
 vector<bool> vis(1);
 while(true) {
    queue<int> que;
    for(int i = 0; i < 1; i++) {</pre>
      if(match_l[i] == -1)
        dis[i] = 0, que.push(i);
      else
        dis[i] = -1;
      vis[i] = false;
    while(!que.empty()) {
      int x = que.front();
      que.pop();
      for(int y : g[x])
        if(match_r[y] != -1 && dis[match_r[y]] == -1) {
          dis[match_r[y]] = dis[x] + 1;
          que.push(match_r[y]);
    }
    auto dfs = [&](auto dfs, int x) {
      vis[x] = true;
      for(int y : g[x]) {
        if(match_r[y] == -1) {
          match_1[x] = y;
          match_r[y] = x;
          return true;
        else if(dis[match_r[y]] == dis[x] + 1
            && !vis[match_r[y]]
            && dfs(dfs, match_r[y])) {
          match_l[x] = y;
          match_r[y] = x;
          return true;
      }
      return false;
    bool ok = true;
    for(int i = 0; i < 1; i++)</pre>
      if(match_l[i] == -1 && dfs(dfs, i))
        ok = false;
```

```
if(ok)
      break:
  return match_1;
} // 0-based
```

3.12 Directed MST

```
// Find minimum directed minimum spanning tree in O(
// DSU rollback is reugired
// Return parent of all nodes, -1 for unreachable ones
    and root
struct dmst_edge { int a, b; ll w; };
struct dmst_node { // Lazy skew heap node
  dmst_edge key;
  dmst_node *1, *r;
  ll delta;
  void prop() {
    key.w += delta;
    if (1) 1->delta += delta;
    if (r) r->delta += delta;
    delta = 0;
  dmst_edge top() { prop(); return key; }
}:
dmst_node *dmst_merge(dmst_node *a, dmst_node *b) {
 if (!a || !b) return a ?: b;
  a->prop();
  b->prop();
  if (a->key.w > b->key.w) swap(a, b);
  swap(a->1, (a->r = dmst_merge(b, a->r)));
void dmst_pop(dmst_node*& a) {
  a->prop();
  a = dmst_merge(a->1, a->r);
pair<11, vector<int>> dmst(int n, int r, const vector<</pre>
    dmst_edge>& g) {
  dsu_undo uf(n);
  vector<dmst_node*> heap(n);
  vector<dmst_node*> tmp;
  for (dmst_edge e : g) {
    tmp.push_back(new dmst_node {e});
    heap[e.b] = dmst_merge(heap[e.b], tmp.back());
  11 \text{ res} = 0;
  vector<int> seen(n, -1), path(n), par(n);
  seen[r] = r;
  vector<dmst_edge> Q(n), in(n, \{-1, -1\}), comp;
  deque<tuple<int, int, vector<dmst_edge>>> cycs;
  for (int s = 0; s < n; s++) {</pre>
    int u = s, qi = 0, w;
    while (seen[u] < 0) {</pre>
      if (!heap[u]) return {-1, {}};
      dmst_edge e = heap[u]->top();
      heap[u]->delta -= e.w;
      dmst_pop(heap[u]);
      Q[qi] = e;
      path[qi++] = u;
      seen[u] = s;
      res += e.w;
      u = uf.find(e.a);
      if (seen[u] == s) { // found cycle, contract
        dmst_node* cyc = 0;
        int end = qi, time = uf.time();
        do {
          cyc = dmst_merge(cyc, heap[w = path[--qi]]);
        } while (uf.join(u, w));
        u = uf.find(u);
        heap[u] = cyc;
        seen[u] = -1;
        cycs.push\_front({u, time, {&Q[qi], &Q[end]}});
      }
    for (int i = 0; i < qi; i++)</pre>
      in[uf.find(Q[i].b)] = Q[i];
  for (auto& [u, t, comp] : cycs) { // restore sol (
    optional)
    uf.rollback(t);
```

```
dmst_edge indmst_edge = in[u];
    for (auto& e : comp) in[uf.find(e.b)] = e;
    in[uf.find(indmst_edge.b)] = indmst_edge;
 for (int i = 0; i < n; i++)</pre>
   par[i] = in[i].a;
  for (auto &a : tmp)
   delete a;
  return {res, par};
3.13 Edge Coloring
/* Find a edge coloring using at most d+1 colors, where
     d is the max deg, in O(V^3)
 * mat[i][j] is the color between i, j in 1-based (0
     for no edge)
  use recolor() to add edge. Calculation is done in
     every recolor */
struct edge_coloring { // 0-based
 int n:
  int mat[maxn][maxn];
 bool vis[maxn], col[maxn];
  void init(int _n) { n = _n; } // remember to init
 int check_conflict(int x, int loc) {
    for (int i = 0; i < n; i++)</pre>
      if (mat[x][i] == loc)
       return i;
    return n:
 int get_block(int x) {
    memset(col, 0, sizeof col);
    for (int i = 0; i < n; i++) col[mat[x][i]] = 1;</pre>
    for (int i = 1; i < n; i++) if (!col[i]) return i;</pre>
    return n;
  void recolor(int x, int y) {
    int pre_mat = get_block(y);
    int conflict = check_conflict(x, pre_mat);
    memset(vis, 0, sizeof vis);
    vis[y] = 1;
    vector<pair<int, int>> mat_line;
    mat_line.push_back({y, pre_mat});
    while (conflict != n && !vis[conflict]) {
      vis[conflict] = 1;
      y = conflict;
      pre_mat = get_block(y);
      mat_line.push_back({y, pre_mat});
      conflict = check_conflict(x, pre_mat);
    if (conflict == n) {
      for (auto t : mat_line) {
       mat[x][t.first] = t.second;
        mat[t.first][x] = t.second;
     }
    }
    else {
      int pre_mat_x = get_block(x);
      int conflict_x = check_conflict(conflict,
    pre_mat_x);
      mat[x][conflict] = pre_mat_x;
      mat[conflict][x] = pre_mat_x;
      while (conflict_x != n) {
        int tmp = check_conflict(conflict_x, pre_mat);
        mat[conflict][conflict_x] = pre_mat;
        mat[conflict_x][conflict] = pre_mat;
```

4 Geometry

4.1 Basic

}

} } mg;

```
struct point {
 ld x, y;
 point() { }
 point(ld a, ld b): x(a), y(b) { }
```

conflict = conflict_x;

swap(pre_mat_x, pre_mat);

recolor(x, mat_line[0].first);

conflict_x = tmp;

```
point operator-(const point &b) const {
   return point(x - b.x, y - b.y);
  point operator+(const point &b) const {
    return point(x + b.x, y + b.y);
  point operator*(ld r) const {
   return point(x * r, y * r);
  point operator/(ld r) const {
    return point(x / r, y / r);
  bool operator<(const point &b) const {</pre>
  return x == b.x ? y < b.y : x < b.x; }
ld dis2() { return x * x + y * y; }</pre>
  ld dis() { return sqrt(dis2()); }
  point perp() { return point(-y, x); }
  point norm() {
    ld d = dis();
    return point(x / d, y / d);
ld cross(const point &a, const point &b, const point &c
  auto x = b - a, y = c - a;
  return x.x * y.y - y.x * x.y;
ld dot(const point &a, const point &b, const point &c)
  auto x = b - a, y = c - a;
  return x.x * y.x + x.y * y.y;
ld area(const point &a, const point &b, const point &c)
  return ld(cross(a, b, c)) / 2;
}
static inline bool eq(ld a, ld b) { return abs(a - b) <</pre>
     EPS; }
int sgn(ld v) {
  return v > 0 ? 1 : (v < 0 ? -1 : 0);
int ori(point a, point b, point c) {
  return sgn(cross(a, b, c));
bool collinearity(point a, point b, point c) {
  return ori(a, b, c) == 0;
bool btw(point p, point a, point b) {
  return collinearity(p, a, b) && sgn(dot(p, a, b)) <=</pre>
    0:
point projection(point p1, point p2, point p3) {
  return (p2 - p1) * dot(p1, p2, p3) / (p2 - p1).dis2()
using Line = pair<point, point>;
bool seg_intersect(Line a, Line b) {
  point p1, p2, p3, p4;
  tie(p1, p2) = a;
  tie(p3, p4) = b;
  if (btw(p1, p3, p4) || btw(p2, p3, p4) || btw(p3, p1,
     p2) || btw(p4, p1, p2))
    return true:
  return ori(p1, p2, p3) * ori(p1, p2, p4) < 0 &&</pre>
    ori(p3, p4, p1) * ori(p3, p4, p2) < 0;
point intersect(Line a, Line b) {
  point p1, p2, p3, p4;
  tie(p1, p2) = a;
  tie(p3, p4) = b;
  1d a123 = cross(p1, p2, p3);
  ld a124 = cross(p1, p2, p4);
return (p4 * a123 - p3 * a124) / (a123 - a124);
```

4.2 2D Convex Hull

```
// returns a convex hull in counterclockwise order
// for a non-strict one, change cross >= to >
vector<point> convex hull(vector<point> p) {
  sort(p.begin(), p.end());
  if (p[0] == p.back()) return {p[0]};
  int n = p.size(), t = 0;
```

5 String

5.1 KMP

```
vector<int> kmp(const string &s) {
   int n = s.size();
   vector<int> dp(n);
   for(int i = 1, j = 0; i < n; i++) {
     while(j && s[i] != s[j])
        j = dp[j - 1];
     if(s[i] == s[j])
        j++;
     dp[i] = j;
   }
   return dp;
}</pre>
```

5.2 Z Value

```
// Return Z value of string s in O(|s|)
// Note that z[0] = |s|
vector<int> Zalgo(const string &s) {
  vector<int> z(s.size(), (int) s.size());
  for (int i = 1, l = 0, r = 0; i < z[0]; ++i) {
    int j = clamp(r - i, 0, z[i - l]);
    while (i + j < z[0] && s[i + j] == s[j])
        j++;
    if (i + (z[i] = j) > r)
        r = i + z[l = i];
  }
  return z;
}
```

5.3 Suffix Array

else {

```
int sa[maxn], tmp[2][maxn], c[maxn];
void get_sa(const string &s) { // m: char set
 int *x = tmp[0], *y = tmp[1], m = 256, n = s.size();
  for (int i = 0; i < m; i++) c[i] = 0;</pre>
  for (int i = 0; i < n; i++) c[x[i] = s[i]]++;</pre>
  for (int i = 1; i < m; i++) c[i] += c[i - 1];
  for (int i = n - 1; i >= 0; --i) sa[--c[x[i]]] = i;
  for (int k = 1; k < n; k <<= 1) {</pre>
   for (int i = 0; i < m; i++) c[i] = 0;</pre>
    for (int i = 0; i < n; i++) c[x[i]]++;</pre>
    for (int i = 1; i < m; i++) c[i] += c[i - 1];</pre>
    int p = 0;
    for (int i = n - k; i < n; i++) y[p++] = i;
    for (int i = 0; i < n; i++)</pre>
     if (sa[i] >= k) y[p++] = sa[i] - k;
    for (int i = n - 1; i >= 0; --i) sa[--c[x[y[i]]]] =
     y[i];
    y[sa[0]] = p = 0;
    for (int i = 1; i < n; i++) {
      int a = sa[i], b = sa[i - 1];
      if (x[a] == x[b] \&\& a + k < n \&\& b + k < n \&\& x[a]
     + k] == x[b + k]) { }
      else p++;
      y[sa[i]] = p;
    if (n == p + 1)
     break;
    swap(x, y);
    m = p + 1;
} // sa[i]: index which ranks i
int rk[maxn], lcp[maxn]; // lcp[i] : lcp with i-1
void get_lcp(const string &s) {
  int n = s.size(), val = 0;
  for (int i = 0; i < n; i++) rk[sa[i]] = i;</pre>
  for (int i = 0; i < n; i++) {</pre>
    if (rk[i] == 0) lcp[rk[i]] = 0;
```

```
if (val) val--;
   int p = sa[rk[i] - 1];
   while (val + i < n && val + p < n && s[val + i]
   == s[val + p])
      val++;
   lcp[rk[i]] = val;
  }
}</pre>
```

5.4 Booth Algorithm

```
// return start index of minimum rotation in O(|s|)
int min_rotation(string s) {
    s += s;
    int k = 0;
    vector<int> f(s.size(), -1);
    for(int j = 1; j < s.size(); j++) {
        int i = f[j - k - 1];
        for(i = f[j - k - 1];
            i != -1 && s[j] != s[i + k + 1]; i = f[i])
            if(s[k + i + 1] > s[j])
            k = j - i - 1;
        if(i == -1 && s[j] != s[k + i + 1]) {
            if(s[j] < s[k + i + 1])
            k = j;
            f[j - k] = -1;
        }
        else
            f[j - k] = i + 1;
    }
    return k;
}</pre>
```

5.5 Manacher Algorithm

```
vector<int> manacher_algorithm(string s) {
  int n = 2 * s.size() + 1;
  string t(n, 0);
  vector<int> len(n);//len[i]: max length when mid at i
  for(int i = 0; i < n; i++) {</pre>
    if(i & 1)
      t[i] = s[i / 2];
  for(int i = 0, l = 0, r = -1; i < n; i++) {
    len[i] = (i <= r ? min(len[2 * 1 - i], r - i) : 0);</pre>
    while(i - len[i] >= 0 && i + len[i] < n && t[i -</pre>
        len[i]] == t[i + len[i]])
      len[i]++;
    len[i]--;
    if(i + len[i] > r)
      l = i, r = i + len[i];
  return len:
```

6 Math

6.1 Lemma And Theory

6.1.1 Pick's Theorem

For a simple polygon, its area A can be written as $A=i+\frac{b}{2}-1$ in which i is the number of points that are strictly interior to the polygon and b is the number of points that are on the polygon's boundary.

6.1.2 Euler's Planar Graph Theorem

 $F\colon$ number of regions bounded by edges. $V-E+F=C+1, E\leq 3V-6$

6.1.3 Modular inversion recurrence

For some prime p,

$$inv_i = \begin{cases} 1 & i = 1 \\ p - \lfloor \frac{p}{i} \rfloor \times inv_{(p \mod i)} & 1 < i < p \end{cases}$$

6.2 Numbers

6.2.1 Catalan number

```
Start from n=0:1,1,2,5,14,42,132,429,1430,4862,16796,58786,\dots C_n=\frac{1}{n+1}\binom{2n}{n}=\frac{(2n)!}{(n+1)!n!}=\prod_{k=2}^n\frac{n+k}{k} C_n=\binom{2n}{n}-\binom{2n}{n+1} Recurrence C_0=1 C_{n+1}=\sum_{i=0}^nC_iC_{n-i} C_{n+1}=\frac{2(2n+1)}{n+2}C_n
```

6.2.2 Primes

```
12721, 13331, 14341, 755779999997771, 999991231, 1000000007, 1000000009, 100069696910^{12} + 39, 10^{15} + 37
```

6.3 Extgcd

```
// return (d, x, y) s.t. ax+by=d=gcd(a,b)
template<typename T>
tuple<T, T, T> extgcd(T a, T b) {
   if(!b) return make_tuple(a, 1, 0);
   auto [d, x, y] = extgcd(b, a % b);
   return make_tuple(d, y, x - (a / b) * y);
}
```

6.4 Chinese Remainder Theorem

```
// x % m1 = x1, x % m2 = x2
ll chre(ll x1, ll m1, ll x2, ll m2){
    ll g = __gcd(m1, m2);
    if ((x2 - x1) % g) return -1; // no solution
    m1 /= g; m2 /= g;
    ll p = get<1>(extgcd(m1, m2));
    ll lcm = m1 * m2 * g;
    ll res = p * (x2 - x1) * m1 + x1;
    // might overflow for above two lines, be cautious
    return (res % lcm + lcm) % lcm;
}
```

6.5 Linear Sieve

```
int least_prime_divisor[maxn];
vectorxint> pr;
void linear_sieve() {
  for(int i = 2; i < maxn; i++) {
    if(!least_prime_divisor[i]) {
      pr.push_back(i);
      least_prime_divisor[i] = i;
    }
  for(int p : pr) {
    if(1LL * i * p >= maxn) break;
    least_prime_divisor[i * p] = p;
    if(i % p == 0) break;
    }
}
```

6.6 Fast Walsh Transform

```
/* do not move ta,tb, default for xor
  * remove last 2 lines for non-xor
  * or convolution:
  * x[i]=ta,x[j]=ta+tb; x[i]=ta,x[j]=tb-ta for inv
  * and convolution:
  * x[i]=ta+tb,x[j]=tb; x[i]=ta-tb,x[j]=tb for inv */
void fwt(int x[], int N, bool inv = false) {
  for(int d = 1; d < N; d <<= 1) {
    for(int s = 0, d2 = d * 2; s < N; s += d2)
    for(int i = s, j = s + d; i < s + d; i++, j++) {
      int ta = x[i], tb = x[j];
      x[i] = modadd(ta, tb);
      x[j] = modsub(ta, tb);
  }</pre>
```

```
if(inv) for(int i = 0, invn = modinv(N); i < N; i++)</pre>
    x[i] = modmul(x[i], invn);
} // N: array len
6.7 Floor Sum
// @param n `n < 2^32`
// @param m `1 <= m < 2^32`
// @return sum_{i=0}^{n-1} floor((ai + b)/m) mod 2^64
ull floor_sum_unsigned(ull n, ull m, ull a, ull b) {
 ull ans = 0:
 while (true) {
  if (a >= m) {
   ans += n * (n - 1) / 2 * (a / m); a %= m;
  if (b >= m) {
   ans += n * (b / m); b %= m;
  ull y_max = a * n + b;
  if (y_max < m) break;</pre>
  // y_{max} < m * (n + 1)
  // floor(y_max / m) <= n
  n = (ull)(y_max / m), b = (ull)(y_max % m);
  swap(m, a);
 return ans:
11 floor_sum(ll n, ll m, ll a, ll b) {
 ull ans = 0;
 if (a < 0) {
 ull a2 = (a % m + m) % m;
  ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a) / m);
  a = a2:
 if (b < 0) {
  ull b2 = (b \% m + m) \% m;
  ans -= 1ULL * n * ((b2 - b) / m);
  b = b2;
 }
 return ans + floor_sum_unsigned(n, m, a, b);
6.8 Linear Programming
/* M constraints, i-th constraint is:
  \sum_{j=0}^{n-1} A[i][j] * x_j <= B[i]
Let v = \sum_{j=0}^{n-1} C[j] * x_j
  maximize v satisfying constraints
  sol[i] = x_i
  remind the precision error */
struct Simplex { // 0-based
  using T = long double;
  static const int N = 410, M = 30010;
  const T eps = 1e-7;
  int n, m;
  int Left[M], Down[N];
  T a[M][N], b[M], c[\bar{N}], v, sol[N];
  bool eq (T a, T b) { return fabs(a - b) < eps; }</pre>
  bool ls (T a, T b) { return a < b && !eq(a, b); }</pre>
  void init(int _n, int _m) {
    n = n, m = m, v = 0;
    for (int i = 0; i < m; ++i) for (int j = 0; j < n;
    ++j) {
      a[i][j] = 0;
    for (int i = 0; i < m; ++i) b[i] = 0;</pre>
    for (int i = 0; i < n; ++i) c[i] = sol[i] = 0;</pre>
  void pivot (int x, int y) {
    swap(Left[x], Down[y]);
    T k = a[x][y]; a[x][y] = 1;
    vector <int> nz;
    for (int i = 0; i < n; ++i) {</pre>
      a[x][i] /= k;
      if(!eq(a[x][i], 0)) nz.push_back(i);
    b[x] /= k;
    for (int i = 0; i < m; ++i) {
  if (i == x || eq(a[i][y], 0)) continue;</pre>
      k = a[i][y], a[i][y] = 0;
b[i] -= k * b[x];
```

for (int j : nz) a[i][j] -= k * a[x][j];

```
if (eq(c[y], 0)) return;
    k = c[y], c[y] = 0, v += k * b[x];
    for (int i : nz) c[i] -= k * a[x][i];
  // 0: found solution, 1: no feasible solution, 2:
    unbounded
  int solve() {
    for (int i = 0; i < n; ++i) Down[i] = i;</pre>
    for (int i = 0; i < m; ++i) Left[i] = n + i;</pre>
    while (1) {
      int x = -1, y = -1;
      for (int i = 0; i < m; ++i) if (ls(b[i], 0) && (x
     == -1 || b[i] < b[x])) x = i;
      if (x == -1) break;
      for (int i = 0; i < n; ++i) if (ls(a[x][i], 0) &&</pre>
     (y == -1 \mid | a[x][i] < a[x][y])) y = i;
      if (y == -1) return 1;
      pivot(x, y);
    while (1) {
      int x = -1, y = -1;
      for (int i = 0; i < n; ++i) if (ls(0, c[i]) && (y
     == -1 \mid \mid c[i] > c[y])) y = i;
      if (y == -1) break;
      for (int i = 0; i < m; ++i) if (ls(0, a[i][y]) &&</pre>
     (x == -1 \mid | b[i] / a[i][y] < b[x] / a[x][y])) x =
    i;
      if (x == -1) return 2;
      pivot(x, y);
    for (int i = 0; i < m; ++i) if(Left[i] < n) sol[</pre>
    Left[i]] = b[i];
    return 0;
} LP;
```

6.9 Miller Rabin

```
bool is_prime(ull x) { // need modular pow(mpow)
  static auto witn = [](ull a, ull u, ull n, int t) {
    if(!a) return false;
    while(t--) {
      ull a2 =
                 _uint128_t(a) * a % n;
      if(a2 == 1 && a != 1 && a != n - 1) return true;
      a = a2;
    }
    return a != 1;
  if(x < 2) return false;</pre>
  if(!(x & 1)) return x == 2;
  int t = __builtin_ctzll(x - 1);
  ull odd = (x - 1) \gg t;
  for(ull m:
       {2, 325, 9375, 28178, 450775, 9780504,
     1795265022})
    if(witn(mpow(m % x, odd, x), odd, x, t))
      return false;
  return true:
| }
```

6.10 Pollard's Rho

```
ull f(ull x, ull k, ull m) {
  return (__uint128_t(x) * x + k) % m;
// does not work when n is prime
// return any non-trivial factor
ull pollard_rho(ull n) {
  if(!(n & 1)) return 2;
  mt19937 rnd(120821011);
  while(true) {
    ull y = 2, yy = y, x = rnd() % n, t = 1;
    for(ull sz = 2; t == 1; sz <<= 1, y = yy) {</pre>
      for(ull i = 0; t == 1 && i < sz; ++i) {</pre>
        yy = f(yy, x, n);
        t = \_gcd(yy > y ? yy - y : y - yy, n);
    if(t != 1 && t != n) return t;
  }
}
```

6.11 Gauss Elimination

```
void gauss_elimination(vector<vector<double>>> &d) {
   int n = d.size(), m = d[0].size();
   for (int i = 0; i < m; ++i) {
      int p = -1;
      for (int j = i; j < n; ++j) {
        if (fabs(d[j][i]) < eps) continue;
        if (p == -1 || fabs(d[j][i])>fabs(d[p][i])) p=j;
      }
      if (p == -1) continue;
      for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
      for (int j = 0; j < n; ++j) {
        if (i == j) continue;
        double z = d[j][i] / d[i][i];
      for (int k = 0; k < m; ++k) d[j][k] -= z*d[i][k];
      }
    }
}
// Not tested</pre>
```

6.12 Fast Fourier Transform

```
using cplx = complex<double>;
const double pi = acos(-1);
cplx omega[maxn * 4];
void prefft(int n) {
 for(int i = 0; i <= n; i++)</pre>
  omega[i] = cplx(cos(2 * pi * i / n),
     sin(2 * pi * i / n));
void fft(vector<cplx> &v, int n) {
  int z = 
           __builtin_ctz(n) - 1;
  for(int i = 0; i < n; i++) {
    int x = 0, j = 0;
    for(; (1 << j) < n; j++) x ^= (i >> j & 1) << (z -</pre>
    j);
    if(x > i) swap(v[x], v[i]);
  for(int s = 2; s <= n; s <<= 1) {
    int z = s \gg 1;
    for(int i = 0; i < n; i += s) {</pre>
      for(int k = 0; k < z; k++) {</pre>
        cplx x = v[i + z + k] * omega[n / s * k];
        v[i + z + k] = v[i + k] - x;
        v[i + k] = v[i + k] + x;
      }
    }
  }
}
void ifft(vector<cplx> &v, int n) {
  fft(v, n); reverse(v.begin() + 1, v.end());
  for(int i = 0; i < n; i++) v[i] = v[i] * cplx(1.0 / n</pre>
    , 0);
vl convolution(const vl &a, const vl &b) {
  // Should be able to handle N <= 10^5, C <= 10^4
  int sz = 1, tot = a.size() + b.size() - 1;
  while(sz < tot) sz <<= 1;</pre>
  prefft(sz);
  vector<cplx> v(sz);
  for(int i = 0; i < sz; i++) {</pre>
    double re = i < a.size() ? a[i] : 0;</pre>
    double im = i < b.size() ? b[i] : 0;</pre>
    v[i] = cplx(re, im);
  fft(v, sz);
  for(int i = 0; i <= sz / 2; i++) {</pre>
    int j = (sz - i) & (sz - 1);
    cplx x = (v[i] + conj(v[j])) * (v[i] - conj(v[j]))
     cplx(0, -0.25);
    if(j != i) v[j] = (v[j] + conj(v[i])) * (v[j] -
    conj(v[i])) * cplx(0, -0.25);
    v[i] = x;
  ifft(v, sz);
  vl c(sz);
  for(int i = 0; i < sz; i++)c[i] = round(v[i].real());</pre>
  c.resize(tot);
  return c;
```

6.13 3 Primes NTT

// MOD: arbitrary prime

```
const int M1 = 998244353;
                                                            return k;
const int M2 = 1004535809;
                                                          } // both not tested
const int M3 = 2013265921;
int super_big_crt(int64_t A, int64_t B, int64_t C) {
 static_assert(M1 <= M2 && M2 <= M3);
  11 r12 = mpow(M1, M2 - 2, M2);
  11 r13 = mpow(M1, M3 - 2, M3);
 11 r23 = mpow(M2, M3 - 2, M3);
 11 M1M2 = 1LL * M1 * M2 % MOD;
 B = (B - A + M2) * r12 % M2;
 C = (C - A + M3) * r13 % M3;
 C = (C - B + M3) * r23 % M3;
 return (A + B * M1 + C * M1M2) % MOD;
} // return ans % MOD
```

6.14 Number Theory Transform

```
/* mod | g | maxn possible values:
998244353 | 3 | 8388608
1004535809 | 3 } 2097152
2013265921 | 31 | 134217728 */
template <int mod, int G, int maxn>
struct NTT {
  ll mpow(ll a, ll b) {
    11 \text{ res} = 1;
    for(; b; b >>= 1, a = a * a % mod)
      if(b & 1)
        res = res * a % mod;
    return res;
  static_assert(maxn == (maxn & -maxn));
  int roots[maxn];
  NTT() {
    ll r = mpow(G, (mod - 1) / maxn);
    for(int i = maxn >> 1; i; i >>= 1) {
       roots[i] = 1;
      for(int j = 1; j < i; j++)</pre>
        roots[i + j] = roots[i + j - 1] * r % mod;
      r = r * r % mod;
    }
  }
  // n must be 2^k, and 0 <= f[i] < mod
  // n >= the size after convolution
  void operator()(vector<ll> &f, int n, bool inv =
    false) {
    for(int i = 0, j = 0; i < n; i++) {</pre>
      if(i < j) swap(f[i], f[j]);</pre>
      for(int k = n >> 1; (j ^= k) < k; k >>= 1);
    for(int s = 1; s < n; s *= 2) {
      for(int i = 0; i < n; i += s * 2) {
         for(int j = 0; j < s; j++) {</pre>
           ll a = f[i + j];
           ll b = f[i + j + s] * roots[s + j] % mod;
           f[i + j] = (a + b) \% mod;
           f[i + j + s] = (a - b + mod) \% mod;
        }
      }
    if(inv) {
      int invn = mpow(n, mod - 2);
      for(int i = 0; i < n; i++)
  f[i] = f[i] * invn % mod;</pre>
      reverse(f.begin() + 1, f.end());
 }
};
```

7 Misc

7.1 Josephus Problem

```
// n people kill m for each turn
int f(int n, int m) {
  int s = 0;
  for (int i = 2; i <= n; i++)
    s = (s + m) % i;
  return s;
}
// died at kth
int kth(int n, int m, int k){
  if (m == 1) return n-1;
  for (k = k*m+m-1; k >= n; k = k-n+(k-n)/(m-1));
```