

Shenandoah GC 2.0

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Overview

- Shenandoah introduction
- · Old vs new barrier scheme
- New platforms
- Elimination of forwarding pointer word
- · Self-fixing barriers



Shenandoah Introduction

Shenandoah is:

- A garbage collector in OpenJDK
- · Concurrent GC
- Universal
- · Aimed to solve the GC pause problem
 - · Regardless of heap size
 - Slow/tiny hardware, cloud, constrained environments, etc



Shenandoah Introduction

You may have heard:

- That Shenandoah requires more memory
- · Has complicated barriers
- · Lacks some optimizations
- → No longer true



Shenandoah Introduction

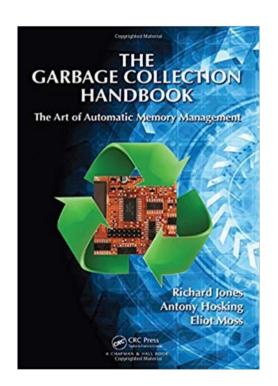
Old Young Mark Copy Compact Ser/Par Conc-Mark Conc-Sweep Copy **CMS** Conc-Mark Copy Compact G1 Shenandoah Conc-Mark Conc-Compact **ZGC**



Concurrent Marking

Not covered here, no time, sorry

If you're interested, it's using "Snapshot at the beginning" algorithm, similar to G1





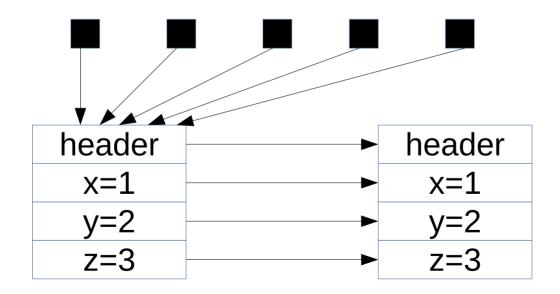
Compaction:

- 1. Copy all reachable objects from the collection set to empty regions
- 2. Adjust all references to old copies to point to new copies
 - → This is easy while the program is stopped.

Problem: how to do this while program is running?

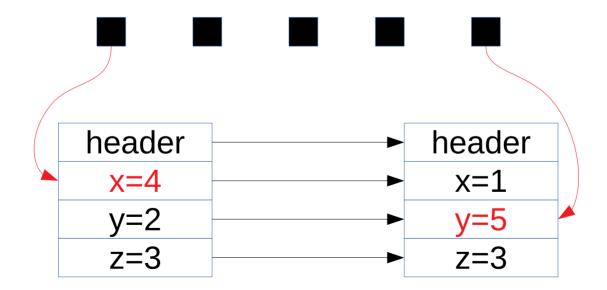






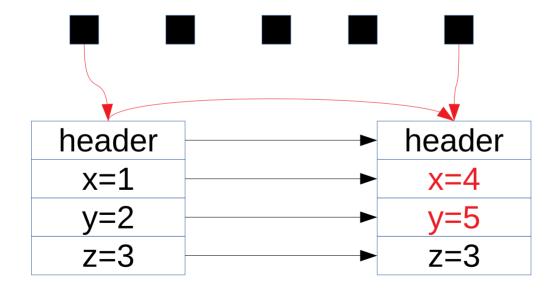


Problem:



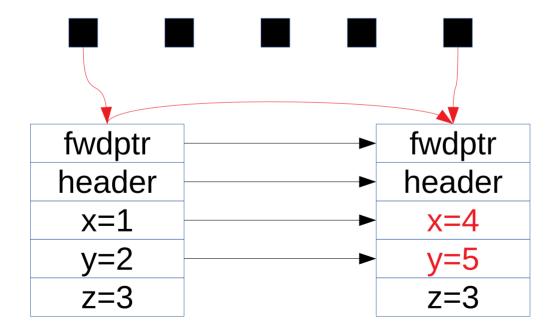


Solution:





Solution (Shenandoah 1.0):





Requires:

- · Read barriers: Intercept reads, resolve by reading forwarding pointer.
- Write barriers: Intercept writes, establish canonical copy by CASing forwarding pointer
- · Object-Equals barriers: to ensure correct result when comparing two copies of same object
- · CAS barriers: combined problem of above
- Arraycopy barriers: Special case for bulk array copy





Example:

```
void doSomething() {
  Foo foo = bar.foo;
  while (...) { // hot loop
    a = foo.x;
    foo.y = b;
  }
}
```



Example:

```
void doSomething() {
  Foo foo = bar.foo;
  while (...) { // hot loop
    foo'= rb(foo); // Required for primitive fields too
    a = foo'.x;
    foo'' = wb(foo); // Required for primitive fields too
    foo''.y = b;
  }
}
```



We can optimize this with some JIT heroics:

```
void doSomething() {
  Foo foo = bar.foo;
  foo'= rb(foo);
  foo'' = wb(foo);
  while (...) { // hot loop
    a = foo'.x;
    foo''.y = b;
  }
}
```



We can optimize this with some JIT heroics:

```
void doSomething() {
  Foo foo = bar.foo;
  foo'= wb(foo);
  while (...) { // hot loop
    a = foo'.x;
    foo'.y = b;
  }
}
```



Idea: Establish canonical copy right when object is loaded

```
void doSomething() {
  Foo foo' = bar.foo;
  foo= lrb(foo');
  while (...) { // hot loop
    a = foo.x;
    foo.y = b;
  }
}
```



Load-reference-barriers:

- · Read-barriers: Intercept reads, resolve by reading forwarding pointer.
- Write-barriers: Intercept writesreference-loads, establish canonical copy by CASing forwarding pointer
- · Above two also required for primitive fields
- Object-equals-barriers: to ensure correct result when comparing two copies of same object
- CAS-object-barriers: combined problem of above





Load-reference-barriers:

- LRB: Intercept loads of objects, establish canonical copy by CASing forwarding pointer
- · CAS-object-barriers: still necessary but simpler





Load-reference-barriers:

Much simpler barrier scheme



- Strong invariant
- Much less frequent (only object-loads!)
- Much simpler optimization story



Load-reference-barriers:

- Much simpler barrier scheme
 - → Makes porting much easier
- Strong invariant
- Much less frequent (only object-loads)
 - → Enables elimination of forwarding pointer
- Much simpler optimization story





Shenandoah Architectures

x86_64	•	Primary target; continuously tested.
x86_32	•	Secondary target; continuously tested.
AArch64	•	Primary target; continuously tested.
ARM32	*	In development; help welcome.
PPC64	8	Not supported; contributions welcome.
S390X	8	Not supported; contributions welcome.
SPARC	8	Not supported. No hardware to test on.
Others	8	Please contact Shenandoah devs for guidance if you are willing to port Shenandoah to another platform.

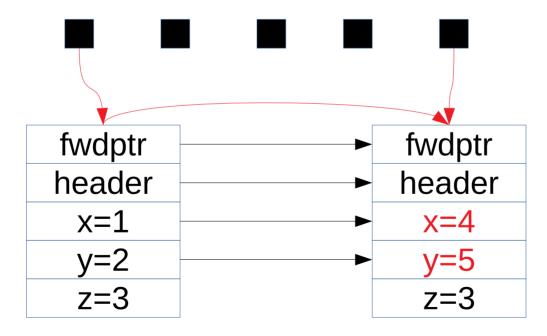


Shenandoah OSes

Linux	•	Primary target; continuously tested.
Windows	•	Secondary target; continuously tested.
macOS	•	Additional target; tested by community.
Solaris	•	Additional target; tested by community.
Others	?	The porting should be trivial, please try and contact Shenandoah devs with your success and failure reports.



Concurrent Compaction (Shenandoah 1.0):





Shenandoah 1.0 – weak invariant:

- All reads may read from old copy (if no canoncical copy has been established yet)
- · All writes **must** write to new copy

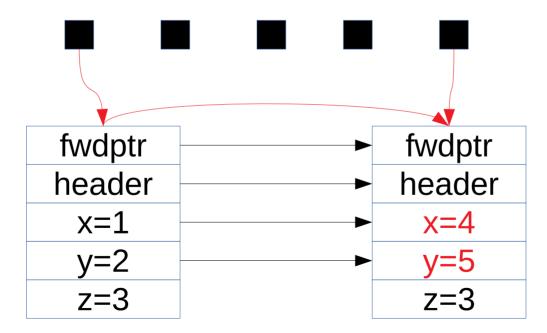


Shenandoah 2.0 – strong invariant:

- · All reads **must** read from new copy
- · All writes **must** write to new copy
 - → old copy is 100% unused
 - → can use it to keep forwarding information

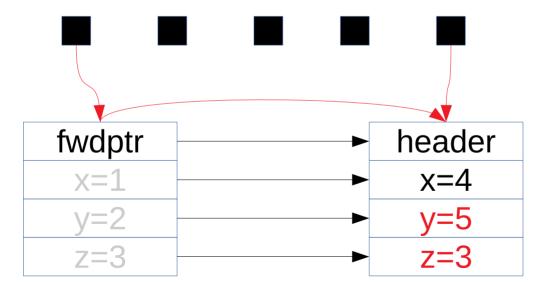


Concurrent Compaction (Shenandoah 1.0):





Concurrent Compaction (Shenandoah 2.0):





```
T load_ref_barrier(T* addr) {
  T obj = *addr;
  if (in_cset(obj)) {
    T fwd = decode_fwdptr(obj);
    if (oop == fwd) {
      fwd = lrb_slowpath(oop);
    return fwd;
  return obj;
```



```
// Shenandoah 1.0 impl
T decode_fwdptr(T obj) {
  return ((T*)obj)[fwdptr offset];
// Shenandoah 2.0 impl
T decode_fwdptr(T obj) {
  intptr_t hdr = ((intptr_t*)obj)[header_offset];
  if ((hdr \& 3) == 3) \{ // Forwarded \}
    return (T)(hdr \& ~3);
  } else {
    Return obj;
```



- · Reduces memory footprint per-object
- Down to 0.66x (best case), realistically 0.75x-0.95x
 - → reduces allocation pressure
 - → fewer GC cycles for same amount of allocs



```
T load_ref_barrier(T* addr) {
  T obj = *addr;
  if (in_cset(obj)) {
    T fwd = decode_fwdptr(obj);
    if (oop == fwd) {
      fwd = lrb_slowpath(oop);
    return fwd;
  return obj;
```



```
T load_ref_barrier(T* addr) {
  T obj = *addr;
  if (in_cset(obj)) {
    T fwd = decode_fwdptr(obj);
    if (oop == fwd) {
      fwd = lrb_slowpath(oop);
    return fwd;
                             Infrequent case: Object not
  return obj;
                             forwarded yet
```



```
T load_ref_barrier(T* addr) {
  T obj = *addr;
  if (in_cset(obj)) {
    T fwd = decode_fwdptr(obj);
    if (oop == fwd) {
      fwd = lrb_slowpath(oop);
    return fwd;
                            Frequent case: Object forwarded,
  return obj;
                            but field not updated, yet
```



While we're there, couldn't we update the field?



```
T load_ref_barrier(T* addr) {
  T obj = *addr;
  if (in_cset(obj)) {
   T fwd = decode_fwdptr(obj);
    if (oop == fwd) {
      fwd = lrb_slowpath(oop);
    } else {
      cas(addr, fwd, obj);
                              Frequent case: Object forwarded,
    return fwd;
                              field updated already
  return obj;
```



Availability

- · Shenandoah 1.0 in JDK12
- Shenandoah 2.0 in JDK13
- · Backports to shenandoah/jdk11 (Fedora, RHEL, ...)
- · Backports to shenandoah/jdk8 (Fedora, RHEL, ...)
- · Upstream to jdk11u proper in progress
- · Upstream to jdk8u proper soon



Contacts

- · Wiki:
 - https://wiki.openjdk.java.net/display/shenandoah/Main
- · Mailing list:
 - https://wiki.openjdk.java.net/display/shenandoah/Main
- · Twitter: @rkennke @shipilev



Questions?

