# Monte Carlo Matrix Inversion

Complexity proof and implementation

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Reinforcement Learning class

# Policy evaluation revisited

• Given a policy  $\pi$ , the value function  $V_{\pi}$  is the fixed point of Bellman equation: if d is the number of state, we have  $V_{\pi} \in \mathbb{R}^d$ 

$$V_{\pi} = R_{\pi} + \gamma P_{\pi} V_{\pi}$$

where

- $\cdot R_{\pi} \in \mathbb{R}^d, R_{\pi}(s) = \mathbb{E}[R_t|S_t = s, A_{t:\infty} \sim \pi].$
- $P_{\pi} \in \mathbb{R}^{d\times d}$  transition matrix:  $(P_{\pi})_{i,j} = \mathbb{P}[s_i|s_i]$
- The solution of Bellman equation is (complexity is  $d^3$ ):

$$V_{\pi}^* = (I - \gamma P_{\pi})^{-1} R_{\pi}$$

· Policy iteration: For each k:

$$V_{\pi}^{k+1} = R_{\pi} + \gamma P_{\pi} V_{\pi}^k$$

· By recursion

$$||V_{\pi}^* - V_{\pi}^k|| \le \frac{1}{1 - \gamma} ||V_{\pi}^* - V_{\pi}^0||$$

# Policy evaluation revisited

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$$||V_{\pi}^* - V_{\pi}^k|| \le \frac{1}{1 - \gamma} ||V_{\pi}^* - V_{\pi}^0||$$

• If  $\epsilon$  is the amount of reduction desired i.e

$$||V_{\pi}^* - V_{\pi}^k|| \le \epsilon ||V_{\pi}^* - V_{\pi}^0||$$

then the number of multiplication required is:

$$(1 + \frac{\log(\epsilon)}{\log(\gamma)})d^2$$

 It is better than exact method, but could we more improve the complexity?

### Monte Carlo methods

#### Idea!!

A simple sum  $\sum_k a_k$  could be interpreted as the expected value of random variable

$$\sum_{k} a_{k} = \sum_{k} \frac{a_{k}}{p_{k}} p_{k} = \mathbb{E}[Z]$$

where Z is random variable defined by  $\mathbb{P}(Z = \frac{a_k}{p_k}) = p_k$  and  $\{p_k\}$  a probability mass.

#### Neuman expansion of inverses

If  $\rho(A) < 1$  than  $(I - A)^{-1}_{M}$  exists and satisfies:

$$(I - A)^{-1} = \lim_{N \to \infty} \sum_{n=0}^{N} A^n$$

• 
$$V = (I - \gamma P)R = R + \gamma PR + \gamma^2 P^2 R + \dots$$

· ith component:

$$V_i = R_i + \gamma \sum_{i_1} p_{ii_1} R_{i_1} + ... + \gamma^k \sum_{i_1...i_k} p_{ii_1} ... p_{i_{k-1}i_k} R_{i_k} + ...$$

$$V_i = R_i + \sum_{k} \sum_{i_1...i_k} \gamma^k \prod_{i=1}^k p_{i_{j-1}i_j} R_{i_k}$$

### Ulman and von-Neumann technique 1950

- Let's define a Markov chain with transition matrix  $\tilde{P}$  and state set  $\{1, 2, ..., d\}$ .
- The chain starts in state i and is allowed o make k transitions.
- the chain's length k is a geometric distributed random variable with parameter  $p_{step}$ :

$$\mathbb{P}(k \text{ state transitions}) = p_{step}^k (1 - p_{step})$$

- Each trajectory starting in state i,  $x_0 = i \rightarrow x_1 = i_1 \rightarrow ... x_k = i_k$  corresponds to a unique term in the sum defining  $V_i$
- For our case the RV Z (defined by  $\mathbb{P}(Z = \frac{a_k}{p_k}) = p_k$ ) is:

$$\mathbb{P}(Z = \frac{\gamma^k \prod_{j=1}^k p_{i_{j-1}i_j} R_{i_k}}{p_{\text{step}}^k (1 - p_{\text{step}}) \prod_{j=1}^k \tilde{p}_{i_{j-1}i_j}}) = p_{\text{step}}^k (1 - p_{\text{step}}) \prod_{j=1}^k \tilde{p}_{i_{j-1}i_j}$$

### Monte Carlo Method

• If we take  $\tilde{P} = P$  and  $p_{step} = \gamma$ , we obtain:

$$\mathbb{P}(Z = \frac{R_{i_k}}{1 - \gamma}) = \gamma^k (1 - \gamma) \prod_{i=1}^k p_{i_{j-1}i_j}$$

· Our MC estimate is for a state i:

$$V_{MC}^{n} = R_i + \frac{1}{n} \sum_{k}^{n} Z_k$$

## Complexity analysis

- work = number of multiplications required to achieve a given amount of reduction error  $\epsilon$
- $\epsilon$  is defined by:

$$|V^*(i) - V_{MC}^n| \le \epsilon ||V^* - V_{MC}^0||$$

$$|V^*(i) - V_{MC}^n| \le \epsilon ||(I - \gamma P)^{-1}R|| \le \epsilon \frac{||R||}{1 - \gamma}$$

- Z is bounded,  $|Z| \leq \frac{||R||}{1-\gamma}$
- ·  $Var(Z) = \mathbb{E}[Z^2] \mathbb{E}[Z]^2 \le (\frac{||R||}{1-\gamma})^2$
- If the distribution of Z is somewhat bell-shaped, the variance may be much less. Assume so that:

$$Var(Z) \le \frac{1}{2} \left(\frac{||R||}{1-\gamma}\right)^2$$

### **Central Limit Theorem**

 As the expectation and the variance of Z are finite, we can apply the CLT theorem:

$$\frac{1}{\sqrt{\operatorname{Var}(\frac{\sum_{k}^{n} Z_{k}}{n})}} \left[\frac{\sum_{k}^{n} Z_{k}}{n} - \mathbb{E}[Z]\right] \Rightarrow N(0,1)$$

• 
$$\operatorname{Var}(\frac{\sum_{k}^{n} Z_{k}}{n}) = \frac{\operatorname{Var}(Z)}{n}$$

•

$$\mathbb{P}[|V^* - V_{MC}^n| \le \epsilon \frac{||R||}{1 - \gamma}] = \mathbb{P}\left[\frac{\left|\frac{\sum_{k}^{n} Z_{k}}{n} - \mathbb{E}[Z]\right|}{\sqrt{\text{Var}\left(\frac{\sum_{k}^{n} Z_{k}}{n}\right)}} \le \epsilon \frac{\frac{||R||}{1 - \gamma}}{\sqrt{\text{Var}\left(\frac{\sum_{k}^{n} Z_{k}}{n}\right)}}\right] \\
\le \mathbb{P}\left[\frac{\left|\frac{\sum_{k}^{n} Z_{k}}{n} - \mathbb{E}[Z]\right|}{\sqrt{\text{Var}\left(\frac{\sum_{k}^{n} Z_{k}}{n}\right)}} \le \frac{\epsilon \sqrt{n}}{\sqrt{2}}\right] \\
(=?) \to \int_{-\frac{\epsilon \sqrt{n}}{\sqrt{2}}}^{\frac{\epsilon \sqrt{n}}{\sqrt{2}}} \frac{1}{\sqrt{2\pi}} e^{\frac{-t^{2}}{2}} dt = 1 - 2 \int_{\frac{\epsilon \sqrt{n}}{\sqrt{2}}}^{\infty} \frac{1}{\sqrt{2\pi}} e^{\frac{t^{2}}{2}} dt$$

- In order to obtain 95 % confidence level,  $\frac{\epsilon\sqrt{n}}{\sqrt{2}}$  should be greater than 2. i.e  $n \ge 1 + \frac{2}{\epsilon}$ .
- Actually, in our derivation, n represents the number of Markov chain realizations (or trajectories). Each trajectory of length k requires k elementary operations. As the length is geometrically distributed with parameter  $\gamma$ ,  $\mathbb{E}[k] = \frac{1}{1-\gamma}$ .
- At the end, the number of operations required to achieve  $\epsilon$  error reduction is:

$$work = \frac{1}{1 - \gamma} (1 + \frac{2}{\epsilon})$$

which proves the formula given in the article.

## Concentration inequalities

- The previous proof is based on the CLT which is an asymptotic result so the provided formula of the work is valid only when n is very large.
- We could give a better estimate of the work using the Hoeffding's inequality.

### Hoeffding's inequality

Let  $X_1, ..., X_n$  be independent random variables bounded by the intervals  $[a_i, b_i]$ :  $a_i \le X_i \le b_i$ . We define the empirical mean of these variables by  $\overline{X} = \frac{1}{n}(X_1 + \cdots + X_n)$ . Then,

$$\mathbb{P}\left(|\overline{X} - \mathbb{E}\left[\overline{X}\right]| \ge t\right) \le 2 \exp\left(-\frac{2n^2t^2}{\sum_{i=1}^{n}(b_i - a_i)^2}\right)$$

and if  $b_i - a_i \le C$ , then:

$$\mathbb{P}\left(\left|\overline{X} - \mathbb{E}\left[\overline{X}\right]\right| \ge t\right) \le 2\exp\left(-\frac{2nt^2}{C^2}\right)$$

### Better estimate of the work

• Let's apply Hoeffding's inequality: In our case  $C = \frac{||R||}{1-\gamma}$ .

$$\mathbb{P}[|V^* - V_{MC}^n| \ge \epsilon \frac{||R||}{1 - \gamma}] \le 2 \exp\left(-\frac{2n(\epsilon \frac{||R||}{1 - \gamma})^2}{C^2}\right)$$
$$= 2 \exp\left(-2n\epsilon^2\right)$$

- In order to obtain 95 % confidence level,  $n \geq 1 \log(\frac{1 0.95}{2})\frac{1}{\epsilon^2}$
- · Finally:

$$work = \frac{1}{1 - \gamma} \left( 1 + \frac{3.68}{\epsilon^2} \right)$$

# Let's check experimentally our beautiful formulas

$$\text{work}_{\text{Monte-Carlo}} = \frac{1}{1-\gamma} (1 + \frac{3.68}{\epsilon^2}) \longleftrightarrow \text{work}_{\text{iterative}} = (1 + \frac{\log(\epsilon)}{\log(\gamma)}) d^2$$



