Distributionally-Aware Exploration for CVaR Bandits

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Mainly based on:

Tamkin, A., Keramati, R., Dann, C., & Brunskill, E. (2019). Distributionally-aware exploration for cvar bandits. In NeurIPS 2019 Workshop on Safety and Robustness on Decision Making.

Introduction

- Traditional multi-armed bandits aims at finding the optimal arm with maximal mean reward.
- ▶ However, risk sensitive objectives are often desirable in some high-stakes settings.
 - e.g. health-care, finance and machine control
- ▶ A popular risk-sensitive measure is the Conditional Value at Risk (CVaR).
- Consider MAB with CVaR as objective called CVaR bandit.

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Notations

- lacktriangle Consider a stochastic K -armed bandit setting with rewards contained in [0, U].
- $ightharpoonup T_i(n)$ the number of times arm i has been pulled up to round n
- $ightharpoonup A_t$ the action taken during round t; $[m]:=\{1,2,...,m\}$
- $ightharpoonup P_i$ the PDF of the distribution of rewards from the *i*-th arm
- $lackbox{$\setminus$} (X_{i,t})_{i\in[K],t\in[n]}$ denote a collection of independent random variables, with the pdf of X_{it} equal to P_i
- $ightharpoonup X_t = X_{A_t,T_{A_t}(t)}$ is the reward in round t
- ▶ The empirical distribution function of $X_{i,t}$ is $\hat{F}_{i,t}(x) = \frac{1}{t} \sum_{s=1}^{t} \mathbb{I}\left\{X_{i,s} \leq x\right\}$

Background

- \blacktriangleright Let X be a bounded random variable with CDF $F(x)=\mathbb{P}[X\leq x]$
- lacktriangle The CVaR at level lpha of a random variable X is then defined as

$$\operatorname{CVaR}_{\alpha}(X) := \sup_{\nu} \left\{ \nu - \frac{1}{\alpha} \mathbb{E} \left[(\nu - X)^{+} \right] \right\}.$$

- ▶ Define the inverse CDF as $F^{-1}(u) = \inf\{x : F(x) \ge u\}$.
- ▶ When X has a continuous distribution, $\mathrm{CVaR}_{\alpha}(X) = \mathbb{E}_{X \sim F} \left[X \mid X \leq F^{-1}(\alpha) \right]$
- ▶ Write CVaR as a function of the CDFF, $CVaR_{\alpha}(F)$.

Background

 \triangleright For continuous random variable X,

$$\operatorname{CVaR}_{\alpha}(X) = \sup_{\nu} \left\{ \nu - \frac{1}{\alpha} \mathbb{E} \left[(\nu - X)^{+} \right] \right\}$$

$$= \sup_{\nu} \left\{ \nu - \frac{1}{\alpha} \int_{-\infty}^{\nu} (\nu - x) f(x) dx \right\}$$

$$= F^{-1}(\alpha) - \frac{1}{\alpha} \int_{-\infty}^{F^{-1}(\alpha)} (F^{-1}(\alpha) - x) f(x) dx$$

$$= F^{-1}(\alpha) - \frac{F^{-1}(\alpha)}{\alpha} F(x) \Big|_{-\infty}^{F^{-1}(\alpha)} + \frac{\int_{-\infty}^{F^{-1}(\alpha)} x f(x) dx}{\alpha}$$

$$= \frac{\int_{-\infty}^{F^{-1}(\alpha)} x f(x) dx}{\int_{-\infty}^{F^{-1}(\alpha)} f(x) dx} = \mathbb{E}_{X \sim F} \left[X \mid X \leq F^{-1}(\alpha) \right]$$

CVaR-regret

ightharpoonup Define the CVaR-regret at time n as

$$R_n^{\alpha} := \mathbb{E}\left[\sum_{t=1}^n \max_i \left(\text{CVaR}_{\alpha}\left(F_i\right)\right) - \text{CVaR}_{\alpha}\left(F_{A_t}\right)\right]$$
$$= \mathbb{E}\left[\sum_{t=1}^n \Delta_{A_t}^{\alpha}\right]$$
$$= \sum_{i=1}^K \mathbb{E}[T_i(n)]\Delta_i^{\alpha},$$

where the third line mimics the regret decomposition in risk-neutral MAB.

Motivation of algorithm

- ► CVaR-UCB computes an optimistic estimate of the CVaR of each arm, and then chooses the arm with the highest UCB in each turn.
- This optimistic estimate is based on the concentration of the empirical CDF via the DKW inequality: With probability at least 1δ ,

$$||\widehat{F}_{i,t}(\cdot) - F_i(\cdot)||_{\infty} \le \sqrt{\frac{1}{2t} \ln(\frac{2}{\delta})}$$

- Specifically, the UCB of CVaR is constructed as follows
 - computes an optimistic estimate of the CDF via DKW inequality
 - the UCB of the CVaR value is set to be the CVaR of that optimistic CDF

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Algorithm

Algorithm 1: CVaR-UCB **Input:** Risk level α , reward range U, horizon n 1 Choose each arm once: 2 Set \hat{F}_a as the CDFs of each arm a on [0, U] for all $a \in [K];$ 3 Set $T_a \leftarrow 1$: 4 for t = 1, ..., n do for $a = 1, \ldots, K$ do $\epsilon_a \leftarrow \sqrt{\frac{\ln(2n^2)}{2T}};$ $\tilde{F}_a(x) \leftarrow \left(\hat{F}_a(x) - \epsilon_i \mathbf{1}\{x \in [0, U)\}\right)^+;$ $UCB_a^{DKW}(t) \leftarrow CVaR_{\alpha}(\tilde{F}_a);$ Play action $A_t = \operatorname{argmax}_i \operatorname{UCB}_i^{\operatorname{DKW}}(t)$; $T_{A_{+}} \leftarrow T_{A_{+}} + 1;$ Update empirical CDF \hat{F}_{A_t} of arm A_t ;

Comparison with Direct Bonuses on the CVaR

- ▶ View the CDF as a set of samples.
- ightharpoonup The optimistic CDF can be found very simply by shifting the lowest-reward samples to the maximum reward U

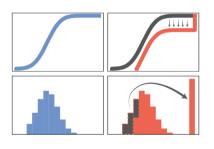




Figure: illustration of the method (left) and comparison with direct bonuses on the sample CVaR (right).

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Comparison with Direct Bonuses on the CVaR

- A natural alternative to the proposal (Cassel et al.'18) directly compute the empirical CDF, extract the empirical CVaR and then add a bonus based on the number of samples.
- ▶ Procedurally this is equivalent to right-shifting each observed sample.
- In contrast, the DKW-based algorithm compute a lower bound on the empirical CDF, effectively shifting probability mass from the lower-reward tail to the max reward.
- ► The latter approach depends on the shape of the CDF itself while the former one is agnostic of the CDF structure, and relies only on the number of samples observed.

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CVaR regret upper bound

Theorem 1.

Consider CVaR-UCB on a stochastic K-armed bandit problem with rewards bounded in [0, U]. For any given horizon n the expected CVaR-regret after this horizon is bounded as

$$R_n^{\alpha} \leq \sum_{i \in [K]: \Delta_i^{\alpha} > 0} \frac{4U^2 \ln(\sqrt{2}n)}{\alpha^2 \Delta_i^{\alpha}} + 3\sum_{i=1}^K \Delta_i^{\alpha}; \quad R_n^{\alpha} \leq \frac{4U}{\alpha} \sqrt{nK \ln(\sqrt{2}n)} + 3KU$$

- \blacktriangleright The bounds differ on their dependence on the number of samples n and risk level α :
 - the problem-dependent bound is $O\left(U^2 \log n/\alpha^2\right)$
 - the problem-free bound grows as $O(U\sqrt{n}/\alpha)$
- For $\alpha = 1$, recover (in dominant terms) the well known UCB regret results

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Lemma 2 (An alternative representation of CVaR).

Let F be a CDF of a bounded non-negative random variable and $\nu \in \mathbb{R}$ be arbitrary. Then $\mathbb{E}_F\left[(\nu-X)^+\right] = \int_0^\nu F(y) dy$. Hence, one can write the CVaR of $X \sim F$ with F(0) = 0 as

$$\text{CVaR}_{\alpha}(F) = \sup_{\nu} \left\{ \frac{1}{\alpha} \int_{0}^{\nu} (\alpha - F(y)) dy \right\}$$

Proof.

First

$$\mathbb{E}_{F}\left[(\nu - X)^{+}\right] = \int_{0}^{\nu} (\nu - x)dF(x) = \nu \int_{0}^{\nu} dF(x) - \int_{0}^{\nu} xdF(x)$$
$$= \nu F(x)|_{0}^{\nu} - (xF(x)|_{0}^{\nu} - \int_{0}^{\nu} F(x)dx) = \int_{0}^{\nu} F(x)dx$$

Proof.

Plugging this identity into

$$\nu - \frac{1}{\alpha} \mathbb{E}_F \left[(\nu - X)^+ \right] = \frac{1}{\alpha} \left(\nu \alpha - \int_0^\nu F(y) dy \right) = \frac{1}{\alpha} \int_0^\nu (\alpha - F(y)) dy$$

Lemma 3 (Bounding difference of CVaR via distance between CDFs).

Let F and G be the CDFs of two non-negative random variables and let ν_F, ν_G be a maximizing value of ν in the definition of $\mathrm{CVaR}_{\alpha}(F)$ and $\mathrm{CVaR}_{\alpha}(G)$ respectively. Then:

$$\begin{split} |\mathrm{CVaR}_{\alpha}(F) - \mathrm{CVaR}_{\alpha}(G)| & \leq \frac{1}{\alpha} \int_{0}^{\max\left\{F^{-1}(\alpha), G^{-1}(\alpha)\right\}} |G(y) - F(y)| dy \\ & \leq \frac{\max\left\{F^{-1}(\alpha), G^{-1}(\alpha)\right\}}{\alpha} \sup_{x} |F(x) - G(x)| \leq \frac{U}{\alpha} ||F(x) - G(x)||_{\mathcal{D}} ||_{\mathcal{D}} ||_{$$

Proof.

Assume w.l.o.g. that $\text{CVaR}_{\alpha}(F) - \text{CVaR}_{\alpha}(G) \geq 0$. A possible value of ν_F is $F^{-1}(\alpha)$.

$$\operatorname{CVaR}_{\alpha}(F) - \operatorname{CVaR}_{\alpha}(G) \leq \nu_{F} - \alpha^{-1} \mathbb{E}_{F} \left[(\nu_{F} - X)^{+} \right] - \left(\nu_{F} - \alpha^{-1} \mathbb{E}_{G} \left[(\nu_{F} - X)^{+} \right] \right) \\
= \frac{1}{\alpha} \left(\mathbb{E}_{G} \left[(\nu_{F} - X)^{+} \right] - \mathbb{E}_{F} \left[(\nu_{F} - X)^{+} \right] \right) \\
= \frac{1}{\alpha} \left(\int_{0}^{\nu_{F}} G(y) dy - \int_{0}^{\nu_{F}} F(y) dy \right) \\
\leq \frac{1}{\alpha} \int_{0}^{\nu_{F}} |G(y) - F(y)| dy \leq \frac{\nu_{F}}{\alpha} \sup_{y} |F(y) - G(y)|$$

We can in full analogy upper-bound $\mathrm{CVaR}_{\alpha}(G) - \mathrm{CVaR}_{\alpha}(F)$ and arrive at the statement.

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Lemma 4 (Optimistic CDF results in optimistic estimate of CVaR).

Let G and F be CDFs of non-negative random variables so that $\forall x \geq 0 : F(x) \geq G(x)$. Then for any $\alpha \in [0,1]$, we have $\mathrm{CVaR}_{\alpha}(F) \leq \mathrm{CVaR}_{\alpha}(G)$.

Lemma 5 (Difference in CVaR).

Let \hat{F} be the empirical CDF obtained by n, i.i.d samples drawn from F. Let $\epsilon>0$ and $\mathcal{G}=\left\{\sup_x|F(x)-\hat{F}(x)|\leq\epsilon\right\}$ be the event that the empirical CDF is uniformly ϵ -close to F. Define $\tilde{F}(x)=[\hat{F}(x)-\epsilon 1\{x\in[0,U]\}]^+$. Then in event \mathcal{G} the following inequality holds

$$\left| \mathrm{CVaR}_{\alpha}(F) - \mathrm{CVaR}_{\alpha}(\tilde{F}) \right| \leq \frac{2\tilde{F}^{-1}(\alpha)\epsilon}{\alpha} \leq \frac{2U\epsilon}{\alpha}$$

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Proof.

By Lemma 3, the triangle-inequality and the definition of ${\cal G}$ and $\tilde F$

$$\begin{aligned} \left| \text{CVaR}_{\alpha}(F) - \text{CVaR}_{\alpha}(\tilde{F}) \right| &\leq \frac{\tilde{F}^{-1}(\alpha)}{\alpha} \sup_{x} |F(x) - \tilde{F}(x)| \\ &\leq \frac{\tilde{F}^{-1}(\alpha)}{\alpha} \sup_{x} |F(x) - \hat{F}(x)| + \frac{\tilde{F}^{-1}(\alpha)}{\alpha} \sup_{x} |\hat{F}(x) - \tilde{F}(x)| \\ &\leq \frac{2\tilde{F}^{-1}(\alpha)\epsilon}{\alpha}. \end{aligned}$$

Lemma 6 (Down-shift is optimistic for CVaR).

In event G the following inequality holds

$$\mathrm{CVaR}_{\alpha}(F) \leq \mathrm{CVaR}_{\alpha}(\tilde{F})$$

- ▶ The proof closely follows the proof of UCB from [Lattimore'20]
- Let c_i^{α} denote the CVaR of arm i and $\hat{F}_{i,t}$ denote the empirical CDF of the i th arm before timestep t
- ▶ Define $\tilde{c}_i^{\alpha}(t)$ as $\tilde{c}_i^{\alpha}(t) = \text{CVaR}_{\alpha}\left(\tilde{F}_{i,t}\right)$ where $\tilde{F}_{i,t}$ is defined as follows,

$$\tilde{F}_{i,t}(x) = \left(\hat{F}_{i,t} - \sqrt{\frac{\ln(2/\delta)}{2T_i(t)}} 1\{x \in [0, U]\}\right)^+$$

$$\epsilon_i(t) = \frac{U}{\alpha} \sqrt{\frac{2\ln(2/\delta)}{T_i(t)}}$$

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- ▶ CVaR regret decomposes as $R_n^{\alpha} = \sum_{i=1}^K \Delta_i^{\alpha} \mathbb{E}[T_i(n)].$
- ▶ Bound $\mathbb{E}[T_i(n)]$ for each suboptimal arm i.
- ► Assume arm 1 is the optimal arm
- ightharpoonup Define the good event G_i as:

$$G_{i} = \left\{ c_{1}^{\alpha} \leq \min_{t \in [n]} \tilde{c}_{1}^{\alpha}(t) \right\} \cap \left\{ \tilde{c}_{i}^{\alpha}\left(u_{i}\right) \leq c_{1}^{\alpha} \right\},$$

where $u_i \in [n]$ will be chosen later.

- ▶ Show by contradiction that if G_i then $T_i(n) \leq u_i$

- ▶ Suppose $T_i(n) > u_i$ on event G_i , then arm i was played more than u_i times over n rounds
- ▶ There must be a round $t \in [n]$ where $T_i(t-1) = u_i$ and $A_t = i$.

$$\tilde{c}_{i}^{\alpha}(t-1) = \text{CVaR}_{\alpha} \left(\hat{F}_{i,t-1} - \sqrt{\frac{\ln(2/\delta)}{2T_{i}(t-1)}} \right)
= \text{CVaR}_{\alpha} \left(\hat{F}_{i,u_{i}} - \sqrt{\frac{\ln(2/\delta)}{2u_{i}}} \right) = \tilde{c}_{i}^{\alpha}(u_{i}) < \tilde{c}_{1}^{\alpha} < \tilde{c}_{1}^{\alpha}(t-1)$$

- ▶ Hence $A_t = \arg \max_j \tilde{c}_j^{\alpha}(t-1) \neq i$, which is a contradiction.
- ▶ It is left to show $\mathbb{P}(G_i^c)$ is low.

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- $G_i^c = \{c_1^{\alpha} > \min_{t \in [n]} \tilde{c}_1^{\alpha}(t)\} \cup \{\tilde{c}_i^{\alpha}(u_i) > c_1^{\alpha}\}$
- Bound the probability of the first event

$$\mathbb{P}\left(c_1^{\alpha} > \min_{t \in [n]} \tilde{c}_1^{\alpha}(t)\right) = \mathbb{P}\left(\exists t \in [n] : c_1^{\alpha} > \tilde{c}_1^{\alpha}(t)\right)$$

$$\leq \mathbb{P}\left(\exists t \in [n] : \sup_{x} \left|\hat{F}_{1,t}(x) - F_1(x)\right| > \sqrt{\frac{\ln(2/\delta)}{2T_1(t)}}\right)$$

$$\leq n\delta$$

lacktriangle Choose u_i such that $\Delta_i^lpha \geq \epsilon_i\left(u_i
ight)$, t_i the round at which arm i was chosen the u_i -th time

$$\mathbb{P}\left(\tilde{c}_{i}^{\alpha}\left(u_{i}\right) > c_{1}^{\alpha}\right) = \mathbb{P}\left(\tilde{c}_{i}^{\alpha}\left(u_{i}\right) - c_{i}^{\alpha} > \Delta_{i}^{\alpha}\right) \leq \mathbb{P}\left(\tilde{c}_{i}^{\alpha}\left(u_{i}\right) - c_{i}^{\alpha} > \epsilon_{i}\left(u_{i}\right)\right)$$

$$\leq \mathbb{P}\left(\sup_{x} \left|\hat{F}_{i,t_{i}}(x) - F_{i}(x)\right| > \sqrt{\frac{\ln(2/\delta)}{2u_{i}}}\right) \leq \delta$$

Substituting the two bound into

$$\mathbb{E}\left[T_i(n)\right] \le u_i + n(n+1)\delta$$

▶ Set $u_i = \left\lceil \frac{2\ln(2/\delta)U^2}{\alpha^2\Delta_i^{\alpha^2}} \right\rceil$ so that $\Delta_i^{\alpha} \geq \epsilon_i\left(u_i\right)$ and choose $\delta = \frac{1}{n^2}$

$$\mathbb{E}\left[T_i(n)\right] \le \left\lceil \frac{2\log\left(2n^2\right)\right)U^2}{\alpha^2 \Delta_i^{\alpha^2}} \right\rceil + 2 \le 3 + \frac{4\ln(\sqrt{2}n)U^2}{\alpha^2 \Delta_i^{\alpha^2}}$$

Substituting this into CVaR-regret decomposition

$$R_n^{\alpha} = \sum_{i=1}^k \Delta_i^{\alpha} \mathbb{E}\left[T_i(n)\right] \le \sum_{i=1}^K \frac{4\ln(\sqrt{2}n)U^2}{\alpha^2 \Delta_i^{\alpha}} + 3\sum_{i=1}^K \Delta_i^{\alpha}$$

$$R_n^{\alpha} = \sum_{i=1}^k \Delta_i^{\alpha} \mathbb{E}\left[T_i(n)\right] = \sum_{i:\Delta_i^{\alpha} < \Delta} \Delta_i^{\alpha} \mathbb{E}\left[T_i(n)\right] + \sum_{i:\Delta_i^{\alpha} \ge \Delta} \Delta_i^{\alpha} \mathbb{E}\left[T_i(n)\right]$$

$$\leq n\Delta + \sum_{i:\Delta_i^{\alpha} \ge \Delta} \Delta_i^{\alpha} \mathbb{E}\left[T_i(n)\right]$$

$$\leq n\Delta + \sum_{i:\Delta_i^{\alpha} \ge \Delta} \left(3\Delta_i^{\alpha} + \frac{4\ln(\sqrt{2}n)U^2}{\alpha^2 \Delta_i^{\alpha}}\right)$$

$$\leq n\Delta + \frac{4K\ln(\sqrt{2}n)U^2}{\alpha^2 \Delta} + \sum_{i=1}^K 3\Delta_i^{\alpha}$$

$$\leq 4\sqrt{nK\ln(\sqrt{2}n)\frac{U}{\alpha}} + 3\sum_{i=1}^K \Delta_i^{\alpha}$$

$$\leq 4\frac{U}{\sqrt{nK\ln(\sqrt{2}n)} + 3KU}$$

Truncated Normal Environments

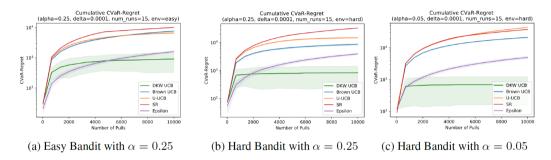


Figure: Compare CVaR-UCB with four others: 1) an ϵ -greedy algorithm with $\epsilon=0$:1; 2) the CVaR best-arm identification algorithm from [Kolla'19]; 3) the U-UCB algorithm from [Cassel'18].; and 4) a variant of U-UCB called Brown-UCB. Means and 95% confidence intervals shown for fifteen runs, with $\delta=10^{-4}$. Y-axis has log scale.

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Comparison against a Tuned *ϵ*-Greedy Baseline

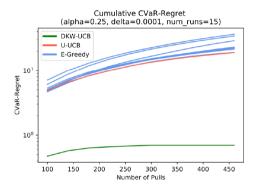


Figure: The ϵ -greedy algorithm was run with a wide range of starting epsilons and decay constants. It is important to verify that finding a successful decay schedule for ϵ -greedy is not easy. In the risk-neutral case, knowledge of the optimality gaps can be leveraged to create an decaying ϵ -greedy algorithm with logarithmic regret growth.

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Dependence on Number of Arms

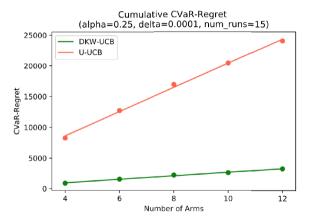


Figure: Cumulative CVaR-regret of our algorithm on the One Good Arm environment for different numbers of arms. Values were collected after 3500 pulls and averaged over 15 runs.

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Proxy regret

Cassel et al. introduced the notion of proxy regret as:

$$\bar{R}_{\pi}(n) = \text{CVaR}_{\alpha}\left(F_{p^{\star}}\right) - \mathbb{E}\left[\text{CVaR}_{\alpha}\left(F_{n}^{\pi}\right)\right]$$

where $p^* = \operatorname{argmax}_{p \in \Delta_{K-1}} \operatorname{CVaR}_{\alpha}(F_p)$ where Δ_{K-1} is the K-1 dimensional simplex

► Here

$$F_p = \sum_{i=1}^{K} p_i F_i$$
$$F_n^{\pi} = \frac{1}{n} \sum_{t=1}^{n} F_{\pi_t}$$

Proxy regret bounds for CvaR-UCB and U-UCB

Proposition 1.

Consider a stochastic K -armed bandit problem with rewards bounded in [0,U]. For any given horizon n and risk level α , both $\mathrm{CVaR} - UCB$ and U - UCB incur proxy regret with $O\left(\frac{\log n}{n}\right)$ and $O\left(1/\alpha^2\right)$ dependency on the horizon and risk level, respectively.

It rules out the possibility that the algorithm's superior performance is due to the use of a different objective.

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